

Hashing

Perfect Hash Functions

- All the hash functions we have considered up to now allow for multiple keys to hash to the same index
 - not only did we have to deal with this problem
 - but it cost us performance as well
- The reason for collisions is because we assumed we knew little about the values of keys

Perfect Hash Functions

- If we know something about the keys, it's possible to write a hash function that will never have collisions
 - this is called a *perfect hash function*
- For example,
 - if you had a class with exactly 100 students and each was given a 2 digit ID, your hash function could simply be to use the students ID number to index into the table

Perfect Hash Function

- The previous example was very simple
 - and not going to be very common
- A real world example
 - compilers need to check for reserved words
 - there are a limited number of reserved words
 - C has about 32; Java about 50
 - it is possible to examine each word and assign it a unique value
 - the performance penalty for this is small because n is small
 - if you were doing this for the entire dictionary, it would be much more time consuming

Cichelli's Algorithm

- *Cichelli's Algorithm* is a commonly used solution to the compiler problem
 - before Cichelli, a binary search was used
- Basic idea
 - assign a value to each letter appearing at the beginning and at the end of each key word
 - this is called a g-value
 - then use the following hash function
 - $h(\text{word}) = (\text{length}(\text{word}) + g(\text{firstletter}) + g(\text{lastletter})) \% \text{size}$

Cichelli's Algorithm

- The real trick is to assign the g-values
 - guess the value of the first and last letter of the first word
 - compute the first word's hash value and reserve it
 - guess the value of the first and last letter of the second word
 - if either letter has already been assigned a g-value, do not assign it a new one – use the assigned value
 - compute the second word's hash value and reserve it
 - if it collides with the first's hash value, make two new guesses
 - repeat this process until all words have a unique hash value

Pseudo-Code

```
// count the frequency that each letter appears as a first or last letter
// order the words by their frequency values – highest value first
    // frequency value = freq(first) + freq(last)
// pick a maxValue – usually the number of words divided by 2
boolean cichelli(Stack wordStack) {
    while(!wordStack.isEmpty()) {
        // pop the first word from wordStack
        if( // both first and last letter have been assigned g-values ) {
            if( // hash value for word is valid ) {
                // assign hash value to word
                if( // recursive call to cichelli() returns true ) { return true; }
                else { detach the hash value for word }
            }
            // push word back on top of wordStack and return false
        }
    }
}
```

Psuedo-Code (continued)

```
else if( // neither letter assigned g-value AND first != last letter ) {  
    // for every value of m and n from 0 to maxVal  
        // assign first letter the g-value of m and second letter gets n  
        if( // hash value for word is valid ) {  
            // assign hash value to word  
            if( // recursive call to cichelli() returns true ) { return true; }  
            else { detach the hash value for word }  
        }  
    }  
    // reset g-value for letters so they are unassigned  
    // push word back on top of wordStack and return false  
}
```


Pseudo-Code

```
else { // only one letter assigned g-value OR first = last letter
    // for every value of m from 0 to maxValve {
        // give unassigned letter the g-value of m
        if( // hash value for word is valid ) {
            // assign hash value to word
            if( // recursive call to cichelli() returns true ) { return true; }
            else { detach the hash value for word }
        }
    }
    // reset g-value for letter so it is unassigned
    // push word back on top of wordStack and return false
}
} // end of while(!wordStack.isEmpty())
return true; // empty stack means we have a solution
}
```

Example

- Consider the following list of states
 - Alabama, Maine, Montana, Nevada, Idaho
- Step one, find frequencies (case insensitive)
 - a: 4; m: 2; n: 1; e: 1; i: 1; o: 1
- Step two, order words based on frequency
 - Alabama-8, Montana-6, Maine-3, Nevada-3, Idaho-2
- Step three, pick a max value
 - $\text{maxValue} = 4 / 2 = 2$

Example

- Step 4, call cichelli()

Alabama: $a = 0$, $h = 2$ hash values $\rightarrow \{ 2 \}$

Montana: $m = 0$, $h = 2$ hash values $\rightarrow \{ 2 \}$

Montana: $m = 1$, $h = 3$ hash values $\rightarrow \{ 2, 3 \}$

Nevada: $n = 0$, $h = 1$ hash values $\rightarrow \{ 1, 2, 3 \}$

Maine: $e = 0$, $h = 1$ hash values $\rightarrow \{ 1, 2, 3 \}$

Maine: $e = 1$, $h = 2$ hash values $\rightarrow \{ 1, 2, 3 \}$

Maine: $e = 2$, $h = 3$ hash values $\rightarrow \{ 1, 2, 3 \}$

Nevada: $n = 1$, $h = 2$ hash values $\rightarrow \{ 2, 3 \}$

Nevada: $n = 2$, $h = 3$ hash values $\rightarrow \{ 2, 3 \}$

Montana: $m = 2$, $h = 4$ hash values $\rightarrow \{ 2, 4 \}$

Nevada: $n = 0$, $h = 1$ hash values $\rightarrow \{ 1, 2, 4 \}$

Maine: $e = 0$, $h = 2$ hash values $\rightarrow \{ 1, 2, 4 \}$

Maine: $e = 1$, $h = 3$ hash values $\rightarrow \{ 1, 2, 3, 4 \}$

Idaho: $i = 0$, $o = 0$, $h = 0$ hash values $\rightarrow \{ 0, 1, 2, 3, 4 \}$

Concerns

- Picking a maxValue is not always easy
 - what if the previous example had used 1?
 - no solution would have been found
 - if this happens, just pick a larger maxValue and try again
- Even with a large maxValue, may not always find a solution
 - if two words are the same length and have the same first and last letter, no solution
 - consider “brick” and “block”
 - no matter what, they will always hash to the same value