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1)

(a) To identify a test case that does not execute the fault, we can choose a test case where the dimensions of matrices **a** and **b** are compatible. For example

(b) To identify a test case that executes the fault but does not cause an error, we can choose a test case where the dimensions of matrices a and b are incompatible, but the error is not raised. For example

$$a = [[1, 2, 3], [3, 4, 5]]$$

 $b = [[7, 8], [6, 7]]$

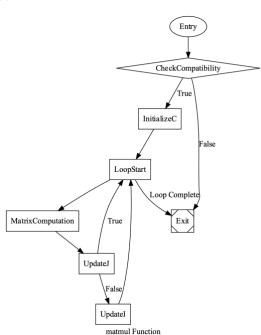
(c) To identify a test case that results in an error but not in a failure, we need a test case where the error is caught and handled gracefully. This means that the code raises an exception but does not crash the program.

$$a = [[5, 7], [8, 21]]$$

$$b = [[8], [4]]$$

(d) The first error occurs when the program is at line 9 where p! =p1. Here p is 2 and p2 is 1. The matrices are compatible and can be multiplied with each other. The expected output is [68, 148]. The output from the program is Value Error ("Incompatible dimensions").

(e)



Entry: Line 7 and 8

Check Compatibility: Line 9
False: Skip to Exit Node

Initialize C: Line 11

Loop Start: Line 12

Loop Start: Line 13

Matrix Computation: Line 14

Update J: Update loop variable (j)

Update I: Update loop. Variable (i)

Exit: Line 15 and terminates the code if there are any

errors

(b) The semanties of the repeat-until loop is that in each iteration -> & & executed -> b is evaluated -> If the current value of 6 is false, the loop continues to the next oferation -> If the current value of b is true, the loop termanates Cand statements following the loop are executed). (8,9) W<b, q'> U false < repeat Smith byg'> q'1 < repeat 8 until 6,9> 49" Consder When bis palse

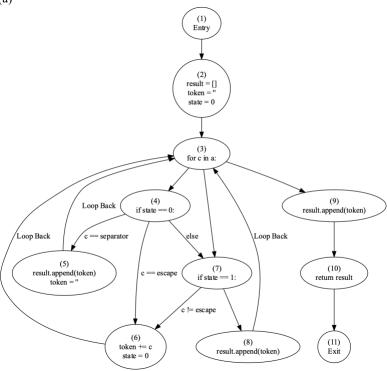
(c)
(c) Show that the following judgement is valid $\langle n:=2; \text{ repeat } n:=n-1 \text{ orden } n \leq DC] \times U(n:=0]$ Prove $\langle n:=2; \text{ repeat } n:=n-1 \text{ Until } n \leq 0, C] \times U(n:=0]$ $\langle n:=2; \text{ repeat } n:=n-1 \text{ Until } n \leq 0, C] \times U(n:=0]$ $\langle n:=2; \text{ repeat } n:=n-1 \text{ Until } n \leq 0, C] \times U(n:=0] \times U(n:=0)$ $\langle n:=2; \text{ repeat } n:=n-1 \text{ (n:=)} \times U(n:=0)$ $\langle n:=2; \text{ repeat } n:=n-1 \text{ Until } x \leq 0, Cn:=0 \text{ filse } \langle \text{ repeat } n:n-1 \text{ until } n \leq 0, C] \times U(n:=0)$ $\langle n:=2; \text{ repeat } n:=n-1 \text{ until } n \leq 0, C] \times U(n:=0)$

(d) Prove that the statement d Ithus topper or egranutrically equivalent to 8; of 6 thun stop else (repeat Suntil 6) The semantics of repeat until loop When the condition is true: 297 49" <b, 9"> 1 true -0 When the wordston is falses LS, a>V a" < ba"> false < repeat S with by9"> V9' Lrepeat S until byay Wa! Consider this as the first statement It has a descriptive tree, denoted as T. Depending on the condition b is true as fake it am have two frams. So, for the false condition we can rewrite D as follows Zrepeat S until by a> Day $\langle S_1 a \rangle \psi a' \longrightarrow T_1$ $\langle repeat S until b_1 a' \rangle \psi a'' \longrightarrow T_2$

Sewantics & of 6 then skop else (repeat Smoth 6) when the condition is true 299749" <6,9"> Whene <000.9">9" Les of 6 teum chap else Crepeat Suntil 6) a > Wall In 3 we conceder as two parts which is denoted as g and &. In the state or, first 8 get executed and get state a)", The code goes to if condition semanter for it when the condition is true esqueen below LS,, 97 01 Lif to teren q else (2/4> V q' The skap statement is being executed. The semantic from exap < skap, 9>49 when the undayon is fake 52,9> 19 9" < by9"> 1 fake < repeat Smith 69/2/201 2019 6 teren scap else Crepeat S with 6), N/ 1/2 In semanties & we consider as two pants which is G & executed in crate of and then gets state of. senantes for of cond" when fake LS2, 07 40° if b teun of else on a) Da, repeat I with 6 & executed in state of and then to state at when wing above rule from S2 Hence provedy

3)





(b)

- Node Coverage (TR_{NC}):

Node coverage aims to ensure that each node in the CFG is visited at least once during testing.

TR_{NC}: {1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11}

All nodes in the CFG are reachable, so there are no infeasible test requirements for node coverage.

- Edge Coverage (TREC):

Edge coverage focuses on traversing every edge in the CFG at least once.

Infeasible test requirements for edge coverage:

(5, 3): This edge is infeasible because it represents a loop back to node 3 from node 5, but this loop is controlled by the loop structure in the code and cannot be directly covered in a single test case. Achieving this edge coverage would require multiple iterations of the loop, which is not a typical way to measure edge coverage.

- Edge-Pair Coverage (TREPC):

Edge-pair coverage involves testing pairs of edges in the CFG to ensure that specific paths and transitions are covered.

TR_{EPC}: $\{(1,2,3),(2,3,9),(2,3,4),(2,3,7),(4,5,3),(4,6,3),(4,7,6),(4,7,8),(5,3,9),(6,3,9),(7,6,3),(7,8,3),(8,3,9),(9,10,11)\}$