

CS557: Cryptography

Elementary Number Theory-VI

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Finite Fields

- finite fields play a key role in cryptography
- The number of elements in a finite field **must** be a power of a prime P (p^n)
 - known as Galois field
 - denoted $GF(p^n)$
- in particular often use the fields:
 - $GF(p)$
 - $GF(2^n)$

Summary- Fields

Def (field): A set F with two binary operations $+$ (addition) and \cdot (multiplication) is called a *field* if

$$1 \quad \forall a, b \in F, a+b \in F$$

$$2 \quad \forall a, b, c \in F, (a+b)+c=a+(b+c)$$

$$3 \quad \forall a, b \in F, a+b=b+a$$

$$4 \quad \exists 0 \in F, \forall a \in F, a+0=a$$

$$5 \quad \forall a \in F, \exists -a \in F, a+(-a)=0$$

$$6 \quad \forall a, b \in F, a \cdot b \in F$$

$$7 \quad \forall a, b, c \in F, (a \cdot b) \cdot c=a \cdot (b \cdot c)$$

$$8 \quad \forall a, b \in F, a \cdot b=b \cdot a$$

$$9 \quad \exists 1 \in F, \forall a \in F, a \cdot 1=a$$

$$10. \quad \forall a \neq 0 \in F, \exists a^{-1} \in F, a \cdot a^{-1}=1$$

$$11 \quad \forall a, b, c \in F, a \cdot (b+c)=a \cdot b+a \cdot c$$

- Equivalently, $(F, +)$ is a commutative (additive) group and $(F \setminus \{0\}, \cdot)$ is a commutative (multiplicative) group.
- A field is a commutative ring with identity where each non-zero element has a multiplicative inverse.

Polynomials over Fields

Polynomials:

Finite field): A field $(F, +, \cdot)$ is called a **finite** field if the set F is **finite**.

Example: \mathbb{Z}_p denotes $\{0, 1, \dots, p-1\}$. We define $+$ and \cdot as addition and multiplication modulo p , respectively.

One can prove that $(\mathbb{Z}_p, +, \cdot)$ is a **field** iff p is prime.

Theorem: There is a unique polynomial $r(x)$ of degree $< m$ over F such that

$$f(x) = h(x) \cdot g(x) + r(x).$$

where, $r(x)$ is called the **remainder** of $f(x)$ modulo $g(x)$.

Galois Fields $GF(p^k)$

Theorem: For every prime power p^k ($k=1,2,\dots$) there is a **unique** finite field containing p^k elements. These fields are denoted by $GF(p^k)$.

There are **no finite fields** with **other** cardinalities.



Remarks:

1. For $F=GF(p^k)$, **$\text{char}(F)=p$** .

2. $GF(p^k)$ and \mathbb{Z}_{p^k} are **not** the same!

Évariste Galois (1811-1832) **NB>: Operations (+, x) require additional steps in Galois field**

Galois Fields GF(p)

- $GF(p)$ is the set of integers $\{0, 1, \dots, p-1\}$ with arithmetic operations modulo prime p
- these form a finite field
 - since have multiplicative inverses

Multiplication in GF(7)

×	0	1	2	3	4	5	6
0	0	0	0	0	0	0	0
1	0	1	2	3	4	5	6
2	0	2	4	6	1	3	5
3	0	3	6	2	5	1	4
4	0	4	1	5	2	6	3
5	0	5	3	1	6	4	2
6	0	6	5	4	3	2	1

Polynomial Arithmetic

- can compute using polynomials

$$f(x) = a_n x^n + a_{n-1} x^{n-1} + \dots + a_1 x + a_0 = \sum a_i x^i$$

- not interested in any specific value of x

- Ordinary Polynomial Arithmetic

- add or subtract corresponding coefficients
- multiply all terms by each other

- Eg: $f(x) = x^3 + x^2 + 2$ and $g(x) = x^2 - x + 1$

$$f(x) + g(x) = x^3 + 2x^2 - x + 3$$

$$f(x) - g(x) = x^3 + x + 1$$

$$f(x) \times g(x) = x^5 + 3x^2 - 2x + 2$$

Polynomial Arithmetic with Modulo Coefficients

- when computing value of each coefficient
 - do calculation modulo some value
 - forms a polynomial ring
 - could be modulo any prime
- but we are most interested in mod 2
 - i.e all coefficients are 0 or 1
 - eg. let $f(x) = x^3 + x^2$ and $g(x) = x^2 + x + 1$

$$f(x) + g(x) = x^3 + x + 1$$

$$f(x) \times g(x) = x^5 + x^2$$

Polynomial Division

- can write any polynomial in the form:
 - $f(x) = q(x)g(x) + r(x)$
 - can interpret $r(x)$ as being a remainder
 - $r(x) = f(x) \bmod g(x)$
- if no remainder, we say $g(x)$ divides $f(x)$
- if $f(x)$ has no divisors other than itself & 1 we say it is **irreducible** polynomial
- Arithmetic modulo an irreducible polynomial forms a field

Polynomial GCD

- can find greatest common divisor for polys
 - $c(x) = \text{GCD}(a(x), b(x))$ if $c(x)$ is the poly of greatest degree which divides both $a(x), b(x)$
- can adapt Euclid's Algorithm to find it:
EUCLID[$a(x), b(x)$]
 1. $A(x) = a(x); B(x) = b(x)$
 2. if $B(x) = 0$ **return** $A(x) = \text{gcd}[a(x), b(x)]$
 3. $R(x) = A(x) \bmod B(x)$
 4. $A(x) \leftarrow B(x)$
 5. $B(x) \leftarrow R(x)$
 6. **goto** 2

Computational Consideration

- since coefficients are 0 or 1, can represent any such polynomial as a bit string
- addition becomes XOR of these bit strings
- multiplication is shift & XOR
- modulo reduction done by repeatedly substituting highest power with remainder of irreducible poly (also shift & XOR)

Computational Example

- in $GF(2^3)$ have (x^2+1) is 101_2 & (x^2+x+1) is 111_2
- so addition is
 - $(x^2+1) + (x^2+x+1) = x$
 - $101 \text{ XOR } 111 = 010_2$
- and multiplication is
 - $(x+1).(x^2+1) = x.(x^2+1) + 1.(x^2+1)$
 $= x^3+x+x^2+1 = x^3+x^2+x+1$
 - $011.101 = (101) \ll 1 \text{ XOR } (101) \ll 0 =$
 $1010 \text{ XOR } 101 = 1111_2$
- polynomial modulo reduction is
 - $(x^3+x^2+x+1) \bmod (x^3+x+1) = 1.(x^3+x+1) + (x^2) = x^2$
 - $1111 \bmod 1011 = 1111 \text{ XOR } 1011 = 0100_2$

Galois Field (polynomial)

- The field defined over the set of residues $F[x]/p(x)$ with the addition and multiplication modulo $p(x)$, where $p(x)$ is irreducible, is called Galois field GF .
- If the field F is Z_N (N is prime) then the corresponding Galois field $Z_N[x]/p(x)$ is denoted by $GF(N^n)$ ($n = \text{degree}(p(x))$)
- $GF(N)$ is the set of integers $\{0, 1, \dots, N-1\}$ with arithmetic operations modulo prime N
 - these form a finite field
 - since have multiplicative inverses
 - **Most General use GFs are $GF(2^n)$ and $GF(N)$ | N is prime**

Example GF(7): GF(N: N is prime)

+	0	1	2	3	4	5	6
0	0	1	2	3	4	5	6
1	1	2	3	4	5	6	0
2	2	3	4	5	6	0	1
3	3	4	5	6	0	1	2
4	4	5	6	0	1	2	3
5	5	6	0	1	2	3	4
6	6	0	1	2	3	4	5

(a) Addition modulo 7

×	0	1	2	3	4	5	6
0	0	0	0	0	0	0	0
1	0	1	2	3	4	5	6
2	0	2	4	6	1	3	5
3	0	3	6	2	5	1	4
4	0	4	1	5	2	6	3
5	0	5	3	1	6	4	2
6	0	6	5	4	3	2	1

(b) Multiplication modulo 7

w	$-w$	w^{-1}
0	0	—
1	6	1
2	5	4
3	4	5
4	3	2
5	2	3
6	1	6

(c) Additive and multiplicative inverses modulo 7

Example GF(2³): p(x) = x³ + x + 1

		000	001	010	011	100	101	110	111
	+	0	1	2	3	4	5	6	7
000	0	0	1	2	3	4	5	6	7
001	1	1	0	3	2	5	4	7	6
010	2	2	3	0	1	6	7	4	5
011	3	3	2	1	0	7	6	5	4
100	4	4	5	6	7	0	1	2	3
101	5	5	4	7	6	1	0	3	2
110	6	6	7	4	5	2	3	0	1
111	7	7	6	5	4	3	2	1	0

(a) Addition

	×	0	1	2	3	4	5	6	7
0	0	0	0	0	0	0	0	0	0
1	0	1	2	3	4	5	6	7	
2	0	2	4	6	3	1	7	5	
3	0	3	6	5	7	4	1	2	
4	0	4	3	7	6	2	5	1	
5	0	5	1	4	2	7	3	6	
6	0	6	7	1	5	3	2	4	
7	0	7	5	2	1	6	4	3	

(b) Multiplication

	w	-w	w ⁻¹
0	0	0	—
1	1	1	1
2	2	2	5
3	3	3	6
4	4	4	7
5	5	5	2
6	6	6	3
7	7	7	4

(c) Additive and multiplicative inverses

Thanks