

CS557: Cryptography

Elementary Number Theory-V

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- Elementary Number Theory
 - GCD
 - Euclidean Algorithm.
 - Group Theory
 - Modular arithmetic

Previous Class

• F

$$8^{-1} \text{ mod } 11$$

$$\text{GCD}(8, 11) = 1$$

• E

$$11 = 8(1) + 3$$

$$3 = 3(2) + 2$$

$$2 = 2(1) + 1$$

$$1 = 1(2) + 0$$

$$\text{GCD} = 1 \text{ remainder} = 0$$

$$a = 6a + 8$$

$$3 = 11 - 8(1)$$

$$2 = 8 - 3(2)$$

$$1 = 3 - 2(1)$$

that $r.a + s.b$

$$\text{GCD} = 1 = 3 - 2(1)$$

$$= 3 - [8 - 3(2)](1)$$

$$= -8(1) + 3(3)$$

$$= -8(-1) + [11 - 8(1)](3)$$

$$= 11(3) + 8(-4)$$

Fast Exponential: Square and multiply

Compute $30^{37} \pmod{77}$

Find $x^c \pmod{n}$

C as bit string: $b_{k-1} \dots b_0$

$z = 1$

For $i = k-1$ downto 0 do

$z = z^2 \pmod{n}$

 if $b_i = 1$ then

$z = z * x \pmod{n}$

Endfor

i	b	z	
5	1	30	$=1*1*30 \pmod{77}$
4	0	53	$=30*30 \pmod{77}$
3	0	37	$=53*53 \pmod{77}$
2	1	29	$=37*37*30 \pmod{77}$
1	0	71	$=29*29 \pmod{77}$
0	1	2	$=71*71*30 \pmod{77}$

Chinese Remainder Theorem

- used to speed up modulo computations
 - working modulo a product of numbers
 - eg. mod $M = m_1 m_2 \dots m_k$
 - Chinese Remainder theorem lets us work in each moduli m_i separately
 - since computational cost is proportional to size, this is faster than working in the full modulus M
-
- can implement CRT in several ways
 - to compute $(A \bmod M)$ can firstly compute all $(a_i \bmod m_i)$ separately and then combine results to get answer using:

Chinese Remainder Theorem (Contd.)

- Given $x \equiv b_i \pmod{m_i}$ to solve for mod N , $N = \prod m_i$
- Algo:
 - 1. Solve for each (N/m_i) . $y_i \equiv 1 \pmod{m_i}$ i.e.
 $y = (N/m_i)$ inverse modulo m_i
 - 2. $A \equiv \sum_i N/m_i \cdot b_i \pmod{N}$
- Example: Given $x \equiv 1 \pmod{5}$ and $x \equiv 10 \pmod{11}$
 - Compute $x \equiv ? \pmod{55}$

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 - Compute $x \equiv ? \pmod{55}$
 - Ans: $21 \pmod{55}$

- Find $x \equiv (2,3,2) \pmod{(3,5,7)}$ respectively.
- Sol:

Example

i	m_i	b_i	M_i	$y_i = M_i^{-1} \pmod{m_i}$	$b_i M_i y_i$
1	3	2	$m/3=35$	$35 y_1 \equiv 1 \pmod{3}$ $\Rightarrow -1$	$2 \times 35 \times -1$
2	5	3	$m/5=21$	$21 y_2 \equiv 1 \pmod{5}$ $\Rightarrow 1$	$3 \times 21 \times 1$
3	7	2	$m/7=15$	$15 y_3 \equiv 1 \pmod{7}$ $\Rightarrow 1$	$2 \times 15 \times 1$
	$M =$ 105				$x = -70 + 63 +$ $30 = 23.$

Modular e'th roots

We know how to solve modular linear equations:

$$a \cdot x + b = 0 \quad \text{in } Z_N$$

$$\text{Solution: } x = -b \cdot a^{-1} \quad \text{in } Z_N$$

What about higher degree polynomials?

Example: Let p be a prime and $c \in Z_p$.

Can we solve:

$$x^2 - c = 0 \quad , \quad y^3 - c = 0 \quad , \quad z^{37} - c = 0 \quad \text{in } Z_p$$

Modular e'th roots

Let p be a prime and $c \in \mathbb{Z}_p$.

Def: $x \in \mathbb{Z}_p$ s.t. $x^e = c$ in \mathbb{Z}_p is called an e'th root of c .

$$7^{1/3} = 6 \text{ in } \mathbb{Z}_{11} \quad 6^3 = 216 = 7 \text{ in } \mathbb{Z}_{11}$$

Examples:

$$3^{1/2} = 5 \text{ in } \mathbb{Z}_{11}$$
$$1^{1/3} = 1 \text{ in } \mathbb{Z}_{11}$$

$2^{1/2}$ does not exist in \mathbb{Z}_{11}

How to Compute e^{th} root?

When does $c^{1/e}$ in Z_p exist?

If $\gcd(e, p-1) = 1$, then

$c^{1/e}$ exists in Z_p for all c in $(Z_p)^*$.

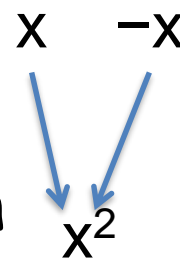
Proof: let $d = e^{-1}$ in Z_{p-1} . i.e, $d \cdot e = 1$ in $Z_{p-1} \Rightarrow$

$$c^d = c^{1/e} \text{ in } Z_p$$

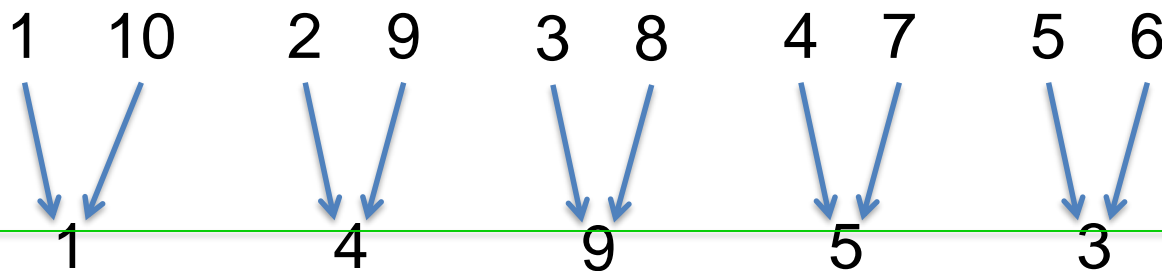
The case $e=2$: square roots

If p is an odd prime then $\gcd(2, p-1) \neq 1$

Fact: in \mathbb{Z}_p^* , $x \rightarrow x^2$ is a 2-to-1 function



Example: in \mathbb{Z}_{11}^* :



Def: x in \mathbb{Z}_p is a **quadratic residue (Q.R.)** if it has a square root in \mathbb{Z}_p

p odd prime \Rightarrow the # of Q.R. in \mathbb{Z}_p is $(p-1)/2 + 1$

Euler's theorem

Thm: x in $(\mathbb{Z}_p)^*$ is a Q.R. $\iff x^{(p-1)/2} = 1$ in \mathbb{Z}_p (p odd prime)

Example:

in \mathbb{Z}_{11}	:	1^5	2^5	3^5	4^5	5^5	6^5	7^5	8^5	9^5	10^5
=		1	-1	1	1	1	-1	-1	-1	1	-1

Note: $x \neq 0 \Rightarrow x^{(p-1)/2} = (x^{p-1})^{1/2} = 1^{1/2} \in \{1, -1\}$ in \mathbb{Z}_p

Def: $x^{(p-1)/2}$ is called the Legendre Symbol of x over p

Computing square roots mod p

Suppose $p \equiv 3 \pmod{4}$

Lemma: if $c \in (\mathbb{Z}_p)^*$ is Q.R. then $\sqrt{c} = c^{(p+1)/4}$ in \mathbb{Z}_p

Proof: $\left[c^{\frac{p+1}{4}} \right]^2 = c^{\frac{p+1}{2}} = \underbrace{c^{\frac{p-1}{2}}}_{=1} \cdot c = c \quad \text{in } \mathbb{Z}_p$

When $p \equiv 1 \pmod{4}$, can also be a bit harder and may be solved using randomized algorithm.

Thanks