



CHANNEL CODING CODE DESIGN

EE539: Principles of Information Theory and
Coding

Abstract

A report on the fulfillment of Exercise 5, focusing on Channel coding code design which takes as input a binary channel and designs a high-rate code.

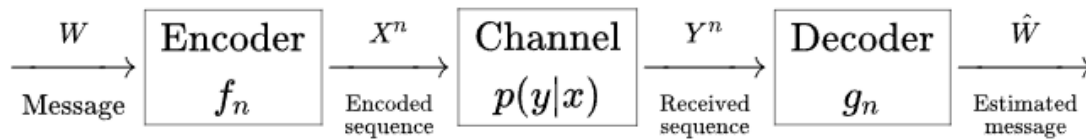
Jayant Som

MS Student in Quantum Science and Nanotechnology

Error Correction (Channel coding)

Error detection and correction are techniques that enable reliable delivery of digital data over unreliable communication channels. Many communication channels are subject to channel noise, and thus errors may be introduced during transmission from the source to a receiver. Error detection techniques allow detecting such errors, while error correction enables detection of errors and reconstruction of the original, error-free data.

The basic mathematical model for a communication system:



A message W is transmitted through a noisy channel by using encoding and decoding functions. An encoder maps W into a pre-defined sequence of channel symbols of length n . In its most basic model, the channel distorts each of these symbols independently of the others. The output of the channel is fed into a decoder which maps the sequence into an estimate of the message. In this setting, the probability of error is defined as:

$$P_e = \Pr \left\{ \hat{W} \neq W \right\}$$

Shannon's noisy-channel coding theorem

This theorem was stated by Claude Shannon in 1948. It establishes that for any given degree of noise contamination of a communication channel, it is possible to communicate discrete data nearly error-free up to a computable maximum rate through the channel. Shannon's theorem has wide-ranging applications in both communications and data storage.

The Shannon theorem states that given a noisy channel with channel capacity C and information transmitted at a rate R , then if $R < C$, there exist codes that allow the probability of error at the receiver to be made arbitrarily small. This means that, theoretically, it is possible to transmit information nearly without error at any rate below a limiting rate, C .

The converse: If $R > C$, an arbitrarily small probability of error is not achievable. All codes will have a probability of error greater than a certain positive minimal level, and this level increases as the rate increases. So, information cannot be guaranteed to be transmitted reliably across a channel at rates beyond the channel capacity.

For every discrete memoryless channel, the channel capacity, defined in terms of the mutual information $I(X;Y)$ as $C = \max I(X;Y)$. Formally, for any $\epsilon > 0$ and $R < C$, for large enough N , there exists a code of length N and rate $\geq R$ and a decoding algorithm, such that the maximal probability of block error is $\leq \epsilon$.

Channel coding Code Design

I have designed the channel coding and segregated my program into 4 parts:

0. **Orchestrator:** This program reads the transition probabilities: $p(0|0)$, $p(1|0)$, $p(0|1)$, $p(1|1)$ and coordinates the entire workflow, running all components systematically, while managing data flow between them.
1. **Channel Capacity Calculator:** This program numerically calculates the maximum mutual information or channel capacity C , corrected to 4 places of decimal. For the channel given as input, it prints the value of C and the optimizing q^* .
2. **Codebook generator:** This creates M random codewords of length n using the optimal distribution q^* , saving them to a file.
3. **Simulator:** This program randomly selects messages, simulates their transmission through the channel to generate outputs, performs maximum likelihood decoding, and calculates the success rate across 1000 trials while displaying sample transmissions.

1. Channel Capacity Calculator

The sequence of operations executed within the code and my explanations are as follows:

a. Input:

- Accepted channel transition probabilities as comma-separated values.
- Validated that each conditional probability distribution sums to 1.

b. Structuring the Data:

- Organized the probabilities into a nested dictionary $p_{Y|X}$ for efficient access during calculations.

c. Calculating Mutual Information:

- Computed the output distribution $\Pr(Y)$ by marginalizing over inputs.
- Calculated the output entropy $H(Y)$.
- Determined the conditional entropy $H(Y|X)$ by averaging over input symbols.

d. Optimizing Capacity:

- Used numerical optimization (`minimize_scalar`) to find the input distribution q^* that maximizes mutual information.
- Ensured the optimization stays within valid probability bounds (0 to 1).

I have used `minimize_scalar` to numerically find the input distribution q^* that maximizes the channel capacity C . Since the mutual information $I(X;Y)$ depends on a single variable $q=P(X=0)$, this method efficiently searches for the optimal q within the range $[0, 1]$. It is the simplest and most efficient approach for finding the capacity-achieving input distribution q^* for a binary channel.

e. Output:

- Stored the results ($p_{Y|X}$, C , q^*) in the IPython namespace for downstream use.
- Printed the channel capacity C (to 4 decimal places) and the optimal input distribution q^* .

```

In [ ]: import numpy as np
        from scipy.optimize import minimize_scalar

        # I am taking channel probabilities as input
        input_str = input("Enter p(0|0), p(1|0), p(0|1), p(1|1) (comma-separated): ")
        p0_0, p1_0, p0_1, p1_1 = map(float, input_str.split(','))

        # I am validating probability distributions
        # Checking if the first pair of probabilities sums to approximately 1
        assert np.isclose(p0_0 + p1_0, 1), "First probability pair does not sum to 1"
        # Checking if the second pair of probabilities sums to approximately 1
        assert np.isclose(p0_1 + p1_1, 1), "Second probability pair does not sum to 1"

        # I am structuring channel probabilities for easy access
        pYgX = {
            0: {0: p0_0, 1: p1_0},
            1: {0: p0_1, 1: p1_1}
        }

        def mutual_info(q):
            # I am calculating output distribution Pr(Y)
            pY0 = q * p0_0 + (1-q) * p0_1
            pY1 = q * p1_0 + (1-q) * p1_1

            # I am computing output entropy H(Y)
            HY = 0.0
            if pY0 > 0:
                HY -= pY0 * np.log2(pY0)
            if pY1 > 0:
                HY -= pY1 * np.log2(pY1)

            # I am computing conditional entropy H(Y|X)
            HYgX = 0.0
            if p0_0 > 0:
                HYgX += q * (-p0_0 * np.log2(p0_0))
            if p1_0 > 0:
                HYgX += q * (-p1_0 * np.log2(p1_0))
            if p0_1 > 0:
                HYgX += (1-q) * (-p0_1 * np.log2(p0_1))
            if p1_1 > 0:
                HYgX += (1-q) * (-p1_1 * np.log2(p1_1))

            return HY - HYgX

        # I am finding capacity by optimizing mutual information
        result = minimize_scalar(lambda q: -mutual_info(q), bounds=(0,1), method='bracket')
        C = mutual_info(result.x)
        q_opt = result.x

        # I am storing results for later use
        get_ipython().push({'pYgX': pYgX, 'C': C, 'q_opt': q_opt})
        print(f"\n• Computed Capacity: C = {C:.4f} bits\n• Optimal q* = {q_opt:.4f}")

```

2. Codebook generator

The sequence of operations executed within the code and my explanations are as follows:

a. Initializing parameters:

- Set the block-length n to 24 symbols.
- Calculated the code rate R as 75% of the channel capacity C .
- Calculated the number of codewords M using $M = 2^{(n \cdot R)}$

b. Generating Codebook:

- Created each codeword by randomly generating n bits using the optimal input distribution q^* .
- Stored all codewords in a list structure.

c. Writing codebook:

- Wrote the complete codebook to Codebook.txt, with one codeword per line.

d. Completion and Storage:

- Displayed a success message confirming codebook generation.
- Then finally, I pushed key variables (M , n , R) to the IPython namespace for downstream use.

```

In [ ]: import numpy as np

# I am setting up code parameters using channel capacity results
n = 24 # Blocklength
R = 0.75 * C # Rate (75% of capacity)
M = int(2 ** (n * R)) # Number of codewords

# I am displaying key parameters for verification
print(f"Parameters:")
print(f"- Blocklength (n): {n}")
print(f"- Rate (R): {R:.4f} (75% of capacity C = {C:.4f})")
print(f"- Number of codewords (M): {M}")
print(f"- Input distribution q*: [0 = {q_opt:.4f}, 1 = {1-q_opt:.4f}]")

# I am generating random codewords using the optimal distribution
codebook = []
for i in range(M):
    codeword = np.random.choice(['0', '1'], size = n, p = [q_opt, 1-q_opt])
    codebook.append(''.join(codeword))

# I am saving the codebook to a text file
with open("Codebook.txt", "w") as f:
    f.write("\n".join(codebook))

# I am confirming successful generation and storing variables
print(f"\033[1;32m\nSuccessfully saved {M} codewords to Codebook.txt !\033[0m")
get_ipython().push({'M': M, 'n': n, 'R': R})

```

3. Simulator

The sequence of operations executed within the code and my explanations are as follows:

a. Initialization:

- Loaded the pre-generated codebook from Codebook.txt file and extracted its parameters (number of codewords M and blocklength n).

b. Defining simulation and decoding functions:

- Created the channel simulation function that corrupts transmitted bits according to given probabilities.
- Created the maximum likelihood decoder function that compares received sequences against all possible codewords.

c. Simulation Preparation:

- Set the total number of trials to 1000 and selecting 10 random trials to display.
- Initialized counters for success tracking and timing the execution.

d. Execution:

- For each trial, I randomly selected a message and its corresponding codeword.
- Then I simulated the channel transmission to generate a received sequence.
- Then I applied ML decoding to recover the most likely transmitted message.
- Finally checked whether the decoding was successful and displaying sample results.

e. Results:

- At last, I calculated the total runtime and success rate percentage.


```

In [ ]: import numpy as np
import time
import random

# I am loading the codebook
with open("Codebook.txt") as f:
    codebook = [line.strip() for line in f]
M, n = len(codebook), len(codebook[0])

# I am using channel probabilities from capacity calculations
# pYgX = {0: {0: p0_0, 1: p1_0}, 1: {0: p0_1, 1: p1_1}}

# I am simulating the channel transmission for each bit
def simulate_channel(x_str):
    return ''.join(
        str(np.random.choice([0, 1], p = [pYgX[int(xi)][0], pYgX[int(xi)][1]
        for xi in x_str
    )

# I am implementing maximum likelihood decoding
def ml_decode(y_n):
    return max(range(M),
                key=lambda c: sum(np.log(pYgX[int(xi)][int(yi)])
                                for xi, yi in zip(codebook[c], y_n)
                                if pYgX[int(xi)][int(yi)] > 0))

# I am initializing simulation parameters
total_trials = 1000
sample_trials = 10
success_count = 0
start_time = time.time()

# I am setting up the results display header
print("="*60)
print(f"RUNNING {total_trials} TRANSMISSION TRIALS\n")
print(f"Displaying {sample_trials} random samples below:\n")

print(f"{'Trial':<6} | {'w':<4} | {'x^n(w)':<{n}} | {'y^n':<{n}} | {'w^':<4}"
print("-" * (6 + 4 + n + n + 4 + 8 + 5*3))

# I am selecting random trials to display
sample_indices = set(random.sample(range(total_trials), sample_trials))

# I am running the main simulation loop
for trial in range(total_trials):

    # I am randomly selecting a message to transmit
    w = np.random.randint(0, M)
    x_w = codebook[w]

    # I am simulating channel transmission
    y_n = simulate_channel(x_w)

    # I am performing ML decoding

```

```

w_ = ml_decode(y_n)

# I am tracking successful decodings
success = (w_ == w)
success_count += success

# I am displaying sample trial results
if trial in sample_indices:
    result = "SUCCESS" if success else "ERROR"
    print(f"{trial+1:<6} | {w:<4} | {x_w:<{n}} | {y_n:<{n}} | {w_:<4} |

# I am calculating and displaying final results
runtime = time.time() - start_time
success_rate = 100 * success_count / total_trials

print("\n" + "="*60)
print(f"\033[1;32mSimulation completed in {runtime:.2f} seconds!\033[0m")
print(f"\033[1;31m\n==== FINAL RESULTS ==== \n\033[0m")
print(f"- Total trials: {total_trials}")
print(f"\033[1;32m- Success rate: {success_rate:.2f}%\033[0m")
print("="*60)

```

Orchestration and Results

1. Simulating a Binary Symmetric Channel (BSC) with 10% Error

$p(0|0), p(1|0), p(0|1), p(1|1)$: **0.9, 0.1, 0.1, 0.9**

- Flips bits with 10% probability
- Capacity ≈ 0.531 bits

2. Simulating a Binary Symmetric Channel (BSC) with crossover probability 0.2

$p(0|0), p(1|0), p(0|1), p(1|1)$: **0.8, 0.2, 0.2, 0.8**

- Flips bits with 20% probability
- Capacity ≈ 0.278 bits

```
In [1]: # I am configuring the notebook files
%config InteractiveShell.ast_node_interactivity = 'all'
%reload_ext autoreload
%autoreload 2

# Running the components in the following sequence :
print("\033[1;31m\n\n==== 1. CHANNEL CAPACITY CALCULATION ==== \n\033[0m")
%run -i 1_Channel_Capacity_Calculator.ipynb

print("\033[1;34m\n\n==== 2. CODEBOOK GENERATION ==== \n\033[0m")
%run -i 2_Codebook_Generator.ipynb

print("\033[1;35m\n\n==== 3. TRANSMISSION SIMULATION ==== \n\033[0m")
%run -i 3_Simulator.ipynb
```

==== 1. CHANNEL CAPACITY CALCULATION ====

- Computed Capacity: $C = 0.5310$ bits
- Optimal $q^* = 0.5000$

==== 2. CODEBOOK GENERATION ====

Parameters:

- Blocklength (n): 24
- Rate (R): 0.3983 (75% of capacity $C = 0.5310$)
- Number of codewords (M): 753
- Input distribution q^* : [$0 = 0.5000$, $1 = 0.5000$]

Successfully saved 753 codewords to Codebook.txt !

==== 3. TRANSMISSION SIMULATION ====

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RUNNING 1000 TRANSMISSION TRIALS

Displaying 10 random samples below:

Trial	w	$x^n(w)$	y^n	w^*
Result				

21	25	110000110101101111111101	110000110101101111110101	25
SUCCESS				
205	4	100011010010110111000101	00001101001011011100101	4
SUCCESS				
306	468	11110111111011110111101	11110111111010110011101	468
SUCCESS				
331	629	001100110001010111010100	001100110001010111010000	629
SUCCESS				
391	388	001000001101110001010110	001000001101110001011110	388
SUCCESS				
469	549	101010110111010111100011	101010110111111111100010	549
SUCCESS				
699	684	101110111111010001000001	10111011111110111001001	79
ERROR				
759	194	100100000100000000101101	101100000100000000101101	194
SUCCESS				
777	160	100101001110011001111000	100000001110011001111000	160
SUCCESS				
786	591	010110101000111011100110	010110111000111001100110	591
SUCCESS				

=====

Simulation completed in 26.68 seconds!

==== FINAL RESULTS ====

- Total trials: 1000
 - **Success rate: 89.00%**
- =====

```
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print("\033[1;31m\n\n==== 1. CHANNEL CAPACITY CALCULATION ==== \n\033[0m")
%run -i 1_Channel_Capacity_Calculator.ipynb

print("\033[1;34m\n\n==== 2. CODEBOOK GENERATION ==== \n\033[0m")
%run -i 2_Codebook_Generator.ipynb

print("\033[1;35m\n\n==== 3. TRANSMISSION SIMULATION ==== \n\033[0m")
%run -i 3_Simulator.ipynb
```

==== 1. CHANNEL CAPACITY CALCULATION ====

- Computed Capacity: $C = 0.2781$ bits
- Optimal $q^* = 0.5000$

==== 2. CODEBOOK GENERATION ====

Parameters:

- Blocklength (n): 24
- Rate (R): 0.2086 (75% of capacity $C = 0.2781$)
- Number of codewords (M): 32
- Input distribution q^* : [$0 = 0.5000$, $1 = 0.5000$]

Successfully saved 32 codewords to Codebook.txt !

==== 3. TRANSMISSION SIMULATION ====

=====

RUNNING 1000 TRANSMISSION TRIALS

Displaying 10 random samples below:

Trial	w	$x^n(w)$	y^n	w^*	
Result					

118	16	101010111101111011110010	101110111000110011010010	16	
SUCCESS					
412	26	000000000110101111001011	010000000101101111001000	26	
SUCCESS					
427	20	000001111011101101010100	011011111010100000010100	20	
SUCCESS					
492	28	010111101001100010001010	010111000001100110001110	28	
SUCCESS					
542	14	001100100110111001111011	01110011011011111110101	2	
ERROR					
560	12	110000000001001010110010	000100000011001011111111	12	
SUCCESS					
691	30	111001101001101110010000	111001001101101010010000	30	
SUCCESS					
696	24	011010010110010010111001	011000010110000110111001	24	
SUCCESS					
725	29	001000111001001110101000	001111101011001100001000	17	
ERROR					
869	5	101001101100010001101000	101000101000010111101011	5	
SUCCESS					

=====

Simulation completed in 1.59 seconds!

==== FINAL RESULTS ====

- Total trials: 1000
- **Success rate: 81.00%**

=====



Codebook.txt - Notepad

File Edit Format View Help

```
010101110111101010101111
000111011110000011101011
000100101010110000011001
111001001000111101101100
101000111001010000111001
010011110001110111110000
111100100110101010110000
110000100101111011110101
101011001110000001101001
100111000011110011010011
010010000001100011001011
110110010100100000010001
110010011000011100111100
110110101011011100111101
011111000100101111011001
001010011110111110010000
101100111010100100011101
001000000111001000000011
101001101000111110011100
110100101110101101010000
010010001010111011001111
100110011000100000101000
001000100100000010111101
111000100101000011000010
001111110001100110111011
110000010010110100001111
100100110100110011000010
111110000001110000101001
011011100010110101101101
110100001110011110100111
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