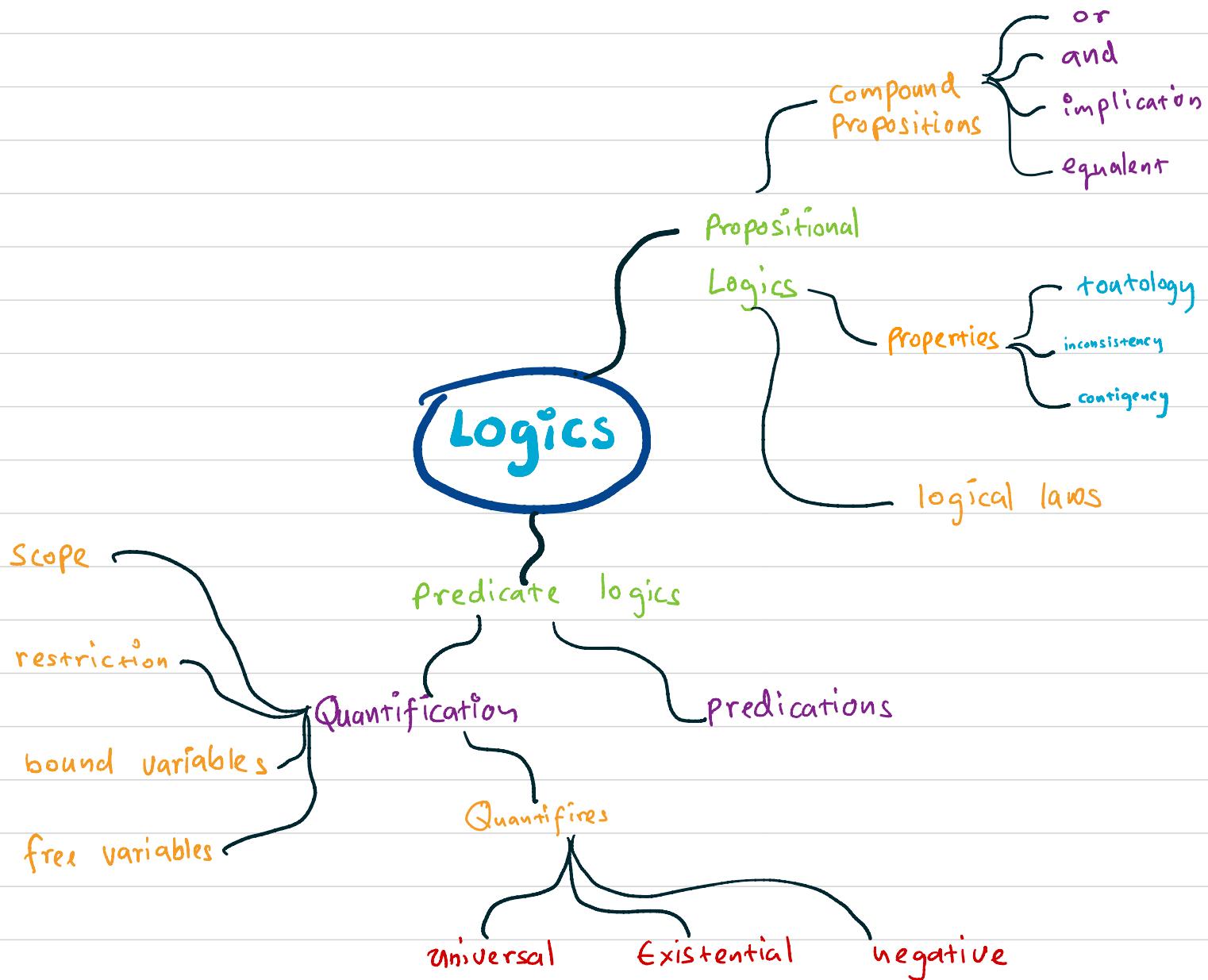


Logics



Jayasanka Weerasinghe
- 2020 -

PROPOSITIONAL LOGICS

66 A proposition is any sentence written down in mathematics or English which can only be TRUE or FALSE

Example :-

01. Today is Wednesday
02. It is raining

These propositions have only TRUE or FALSE answers; No other values. We call it "LAW OF EXCLUDED MIDDLE"

We use "propositional symbols" to represent a proposition

$$P \equiv \text{It is raining} \quad \text{:-}$$

The opposite of a proposition represented by the logical not (" \neg ") symbol (negation)

$$\neg P \equiv \text{It is } \underline{\text{not}} \text{ raining} \quad \text{:-}$$

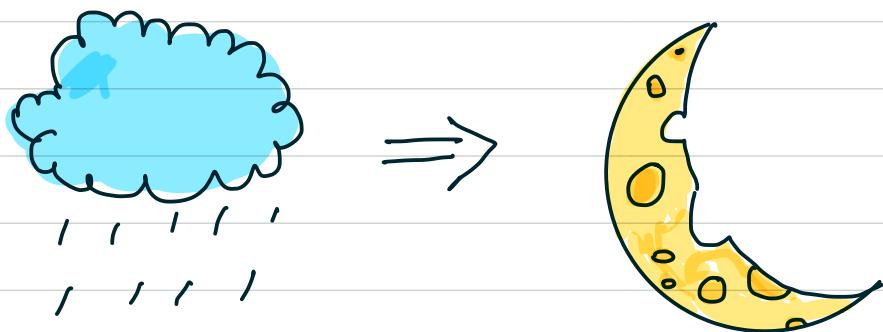


COMPOUND PROPOSITIONS

66 Compound propositions made up simple propositions combined using logical connectives.

- ✓ or
- ∧ and
- ⇒ if... then / implies
- ↔ is equivalent to

LOGICAL IMPLICATION (\Rightarrow)



If it's raining then the moon is made of cheese

or

It is raining implies the moon is made of cheese



This doesn't mean if it is not raining, the moon should not be made of cheese.

But, if it is raining, then the moon is made of cheese for sure!

It's raining and the moon is made of cheese **TRUE**

It's not raining and the moon is made of cheese **TRUE**

It's not raining and the moon is not made of cheese **TRUE**

It's raining and the moon is not made of cheese **FALSE**

This can happen, who knows!

The truth table :-

P	Q	$P \Rightarrow Q$
true	true	true
true	false	false
false	true	true
false	false	true

$P \Rightarrow Q$ is true only if (P is false) OR (Q is true)

Another example :-

If Jimmy is a Dog , then Jimmy has 4 legs

P = Jimmy is a Dog

Q = Jimmy has 4 legs

P	Q	$P \Rightarrow Q$
true	true	true
true	false	false
false	true	true
false	false	true

Jimmy can be a Cat

Jimmy can be a bird

LOGICAL EQUIVALENCE (\Leftrightarrow)

Equivalence states whether two logical statements have the same truth value.



This is Jayantha.
Jayantha is an animal.

We need to know whether he is a mammal or not!

If Jayantha drinks milk then he is a mammal.

If Jayantha is a mammal, he is drinking milk.

A = Jayantha drinks milk

B = Jayantha is a mammal

A	B	$A \Leftrightarrow B$
true	true	true
true	false	false
false	true	false
false	false	true

$A \Leftrightarrow B$ is true if and only if $A \Rightarrow B$ AND $B \Rightarrow A$ are true

Bracketed Propositions

To ensure that a complex proposition has the intended meaning, we must use brackets.

ex :- $P \wedge (Q \Rightarrow R)$

Contingencies, Inconsistencies & Tautologies

Tautology \rightarrow a logical statement which is true in every interpretation

ex :- $\neg P \vee P$

Inconsistency \rightarrow a logical statement which is false in every interpretation

ex :- $\neg P \wedge P$

Contingency \rightarrow a logical statement which is

- true in at least one interpretation
- false in at least one interpretation

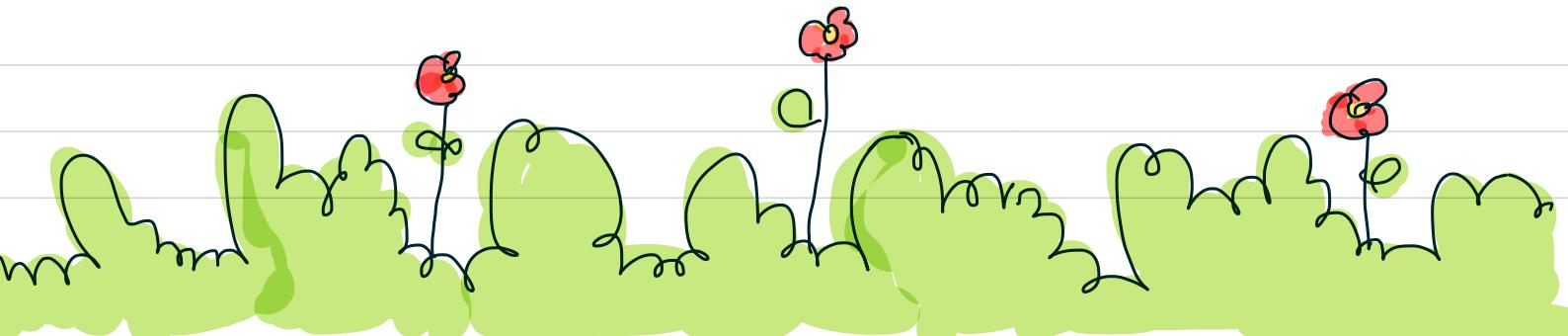
ex :- $P \Rightarrow Q$

Logical Laws

Here's a list of standard logical laws :-

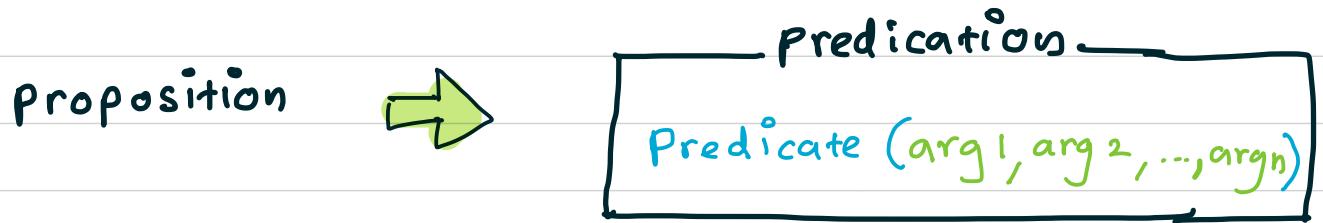
$P \wedge P = P$	[\wedge Idempotence]
$P \vee P = P$	[\vee Idempotence]
$P \wedge \neg P = \text{false}$	[Contradiction]
$P \vee \neg P = \text{true}$	[Excluded Middle]
$(P \wedge Q) \wedge R = P \wedge (Q \wedge R)$	[\wedge Associativity]
$(P \vee Q) \vee R = P \vee (Q \vee R)$	[\vee Associativity]
$P \wedge Q = Q \wedge P$	[\wedge Commutativity]
$P \vee Q = Q \vee P$	[\vee Commutativity]
$P \wedge \text{true} = P$	[\wedge Identity]
$P \vee \text{false} = P$	[\vee Identity]
$P \wedge \text{false} = \text{false}$	[\wedge Zero]
$P \vee \text{true} = \text{true}$	[\vee Zero]

$P \wedge (Q \vee R) = (P \wedge Q) \vee (P \wedge R)$	[Distribution over \vee]
$P \vee (Q \wedge R) = (P \vee Q) \wedge (P \vee R)$	[\vee Distribution over \wedge]
$P \wedge (P \vee Q) = P$	[Absorption I]
$P \vee (P \wedge Q) = P$	[Absorption II]
$\neg(P \wedge Q) = \neg P \vee \neg Q$	[de Morgan I]
$\neg(P \vee Q) = \neg P \wedge \neg Q$	[de Morgan II]
$\neg(\neg P) = P$	[Double Negation]
$P \Rightarrow Q = \neg P \vee Q$	[Implication]
$P \Rightarrow Q = \neg P \Rightarrow \neg Q$	[Contrapositive]
$P \vee Q = (\neg P) \Rightarrow Q$	[\vee Definition]
$P \Leftrightarrow Q = (P \Rightarrow Q) \wedge (Q \Rightarrow P)$	[\Leftrightarrow Definition]



PREDICATE Logic

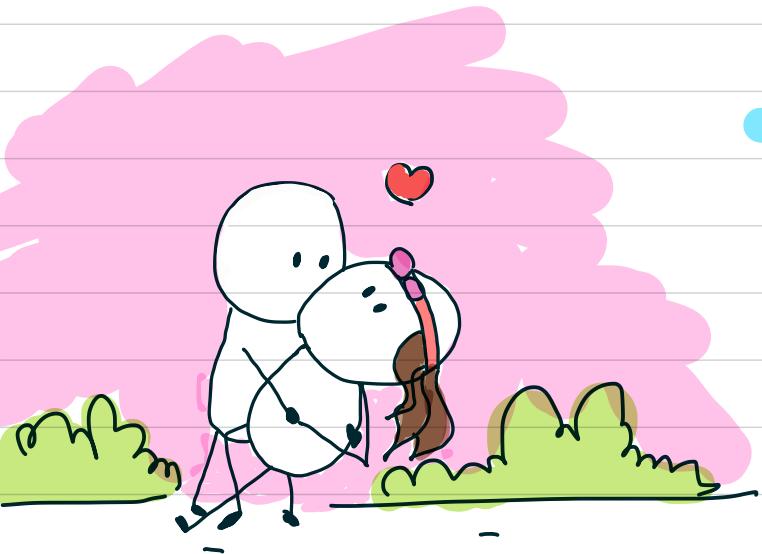
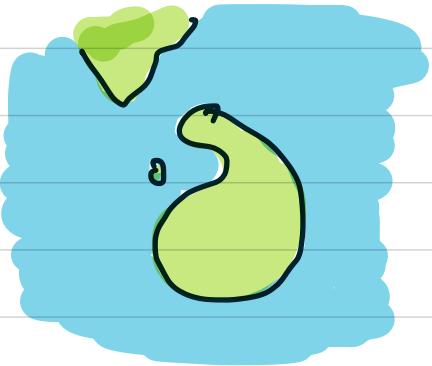
any proposition can be defined as
a predication



$P(x)$ can be read as x is a P

Ex :-

Sri Lanka is a country
Country (sri_lanka)



Upali loves Sumana

Love (Upali, Sumana)
↑ ↓
constants

or
 $\alpha = 'Upali'$, $\beta = 'Sumana'$

Love (α , β)

↑ ↓
variables

Rathne sends Jayanthi a letter

Send (rathne, Jayanthi, letter)

A predicate with $1 \leq n$ arguments is called
n place predicate

Ex :- Love (upali, samantha) is 2 place predicate

ARGUMENT STRUCTURE MATTERS!



Thissa shot Muwa
Shoot (thissa, muwa)

Shoot (muwa, thissa)
Muwa shot Thissa



Same predicate, different meanings!

COMPLEX PREDICATES

The teacher saw that the children were reading a book.

proposition 1 : The teacher saw something

proposition 2 : The children were reading a book

See (teacher, Read (children, book))

Read more : youtu.be/lhodKMPwShc

QUANTIFICATION

Consider the following examples:



p: All girls love Tharindu
Love (girls, tharindu)



q: Some girls love Tharindu
Love (girls, tharindu)



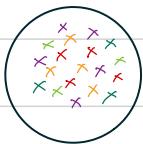
r: Girls don't love Tharindu
Love (girls, tharindu)

This predicates doesn't reflect the different meanings. We need an addition mechanism to formalize the relationships between sets...

That's where the Quantifiers comes to the picture!

\forall_x

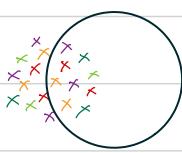
universal quantifier



For all x it holds that...

\exists_x

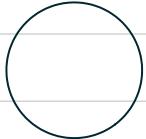
existential quantifier



There exists at least one x such that ...

$\neg_x \sim_x$

negative quantifier

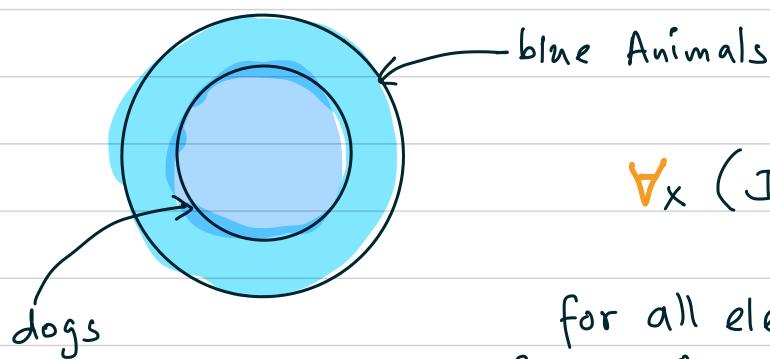


For no x it holds that...



UNIVERSAL QUANTIFIER

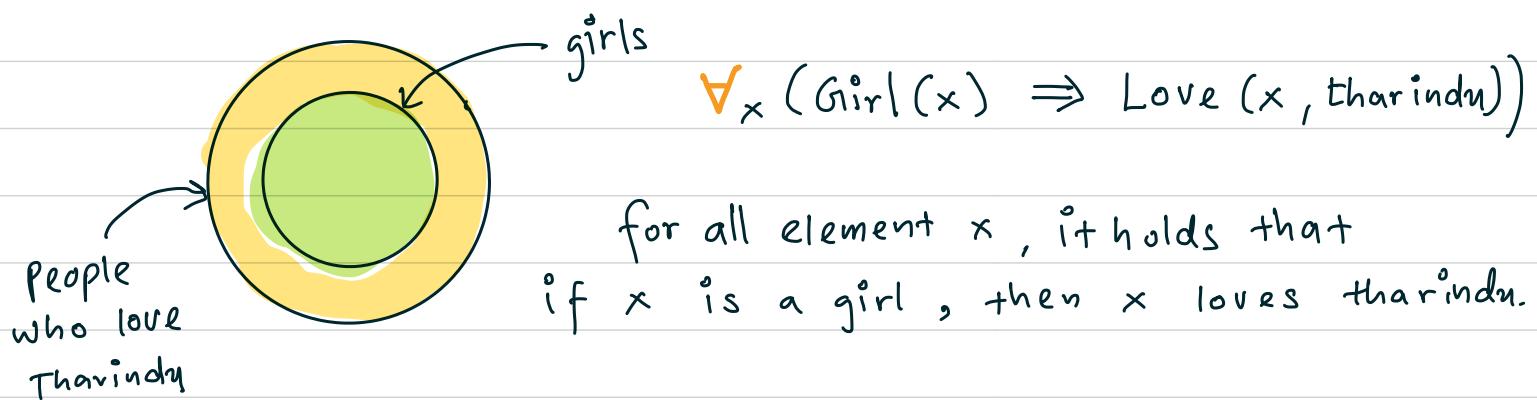
P: All dogs are blue



$$\forall x (\text{Dog}(x) \Rightarrow \text{Blue}(x))$$

for all element x , it holds that
if x is a dog, then x is blue.

q: All girls love Tharindu

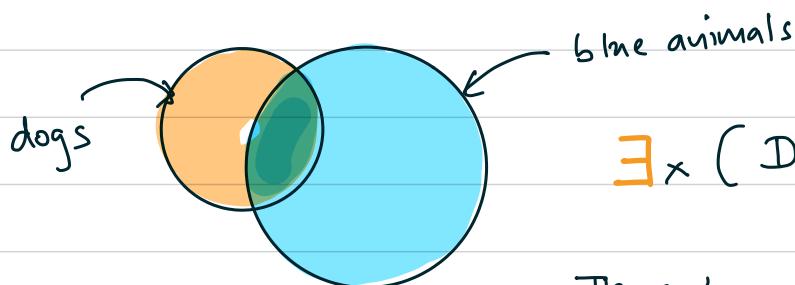


$$\forall x (\text{Girl}(x) \Rightarrow \text{Love}(x, \text{Tharindu}))$$

for all element x , it holds that
if x is a girl, then x loves Tharindu.

EXISTENTIAL QUANTIFIER

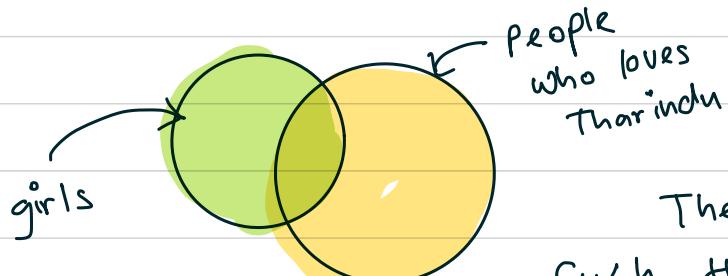
P: Some dogs are blue



$$\exists x (\text{Dog}(x) \wedge \text{Blue}(x))$$

There's at least one element such that
 x is a dog and x is blue.

q: Some girls love Tharindu

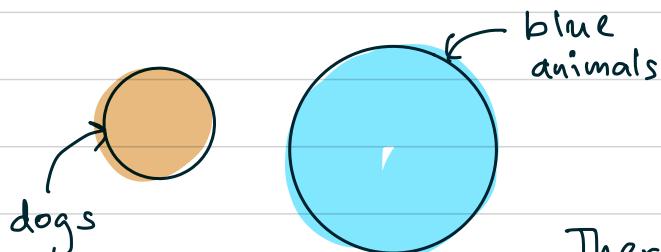


$$\exists x (\text{Girl}(x) \wedge \text{Love}(x, \text{Tharindu}))$$

There's at least one element x , such that x is a girl and also x loves Tharindu.

NEGATIVE QUANTIFIER (\neg , \sim)

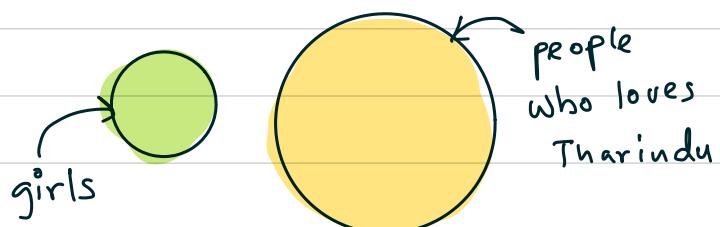
p: No dog is blue



$$\neg \exists x (\text{Dog}(x) \Rightarrow \text{Blue}(x))$$

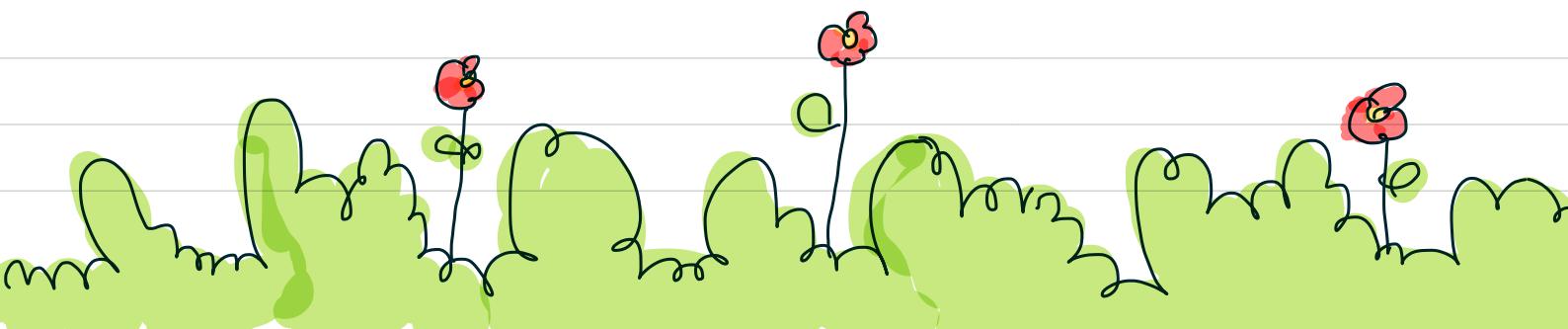
There's no element x , such that if x is a dog; then x is also blue.

q: No girl loves Tharindu



$$\neg \exists x (\text{Girl}(x) \Rightarrow \text{Love}(x, \text{Tharindu}))$$

There's no element x , such that if x is a girl; then x is also loves Tharindu.



66

Quantifiers cannot refer to anything by themselves.

It's always necessary to specify what is being quantified. The **RESTRICTION** establishes the set of entities that are effected (restricted).



Ex :-

$$\forall x (\boxed{\text{Dog}(x)} \Rightarrow \boxed{\text{Blue}(x)})$$

restriction scope

- Bhathiya & Santhosh

The **Scope** of the quantifier is a set of entities denoted by the main predication of the proposition in which the quantifier appears.

further restrictions can be defined by adding adjectives.

Ex :-

All smoking dogs are blue.

$$\forall x (\underbrace{\text{Dog}(x) \wedge \text{Smoke}(x)}_{\text{restriction}} \Rightarrow \underbrace{\text{Blue}(x)}_{\text{scope}})$$

01. Everybody is happy

$$\forall x (\text{Person}(x) \Rightarrow \text{Happy}(x))$$

02. Kumara likes some girls

$$\exists x (\text{Girl}(x) \wedge \text{Like}(\text{kumara}, x))$$

03. Nobody shoot Thissa

$$\neg \exists x (\text{Person}(x) \Rightarrow \text{Shot}(x, \text{thissa}))$$

$$\neg \exists x (\text{Person}(x) \wedge \text{Shot}(x, \text{thissa}))$$

Read as:

"It is not the case that

there's at least one element x

such that x is a person and also
x shoot Thissa"

BOUND VARIABLES AND FREE VARIABLES

66 A variable in a formula is defined as "bound" if it is named by the quantifier & within the scope. 

- Kanka Herath

ex:-

$$\forall x (Dog(x) \Rightarrow Drunk(x))$$

x is a bound variable!

66 A variable that occurs in a formula is said to be "free" if it is not a bound variable 

- Shashika Nisansala

ex:-

$$\forall x (Person(x) \Rightarrow Love(x, y))$$

y is a free variable!

x is a bound variable...

note :- A single variable may occur both bound and free in the same formula.

ex:-

$$(\forall x (A(x) \Rightarrow B(x))) \wedge Q(x)$$

 = bound

 = free

$$\forall_x (\text{Dog}(x) \Rightarrow \text{Drunk}(x))$$

can be also written as:-

$$\forall(x) \cdot (x \in \text{DOG} \Rightarrow \text{Drunk}(x))$$
$$\exists_{(n)} (\text{Natural-number}(n) \wedge \text{Odd}(n))$$
$$\exists(n) \cdot (n \in \mathbb{N} \wedge \text{Odd}(n))$$
