

Course title : **CSE2001**  
Course title : **Data Structures and Algorithms**  
Module : **6**  
Topic : **1**

# **Graph Data Structure**

# Objectives

This session will give the knowledge about

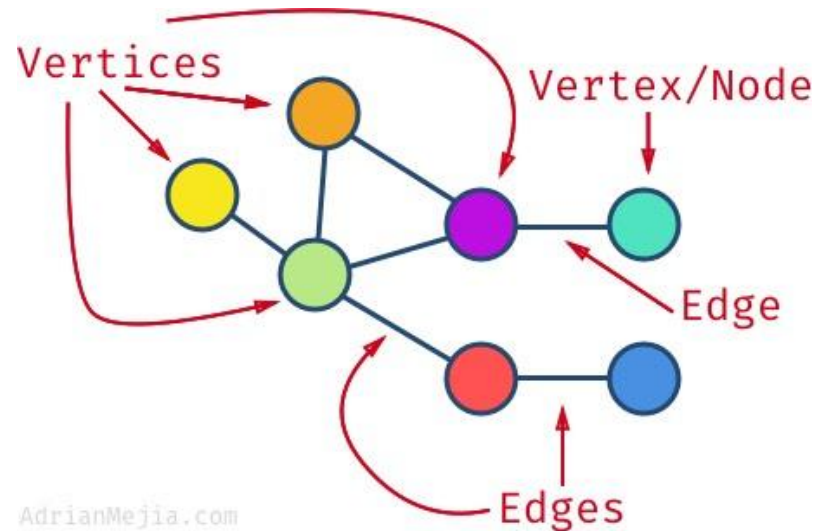
- Introduction to Graphs
- Breadth First Search - BFS
- Depth First Search - DFS

# Introduction to Graph

A graph is a collection of **nodes (or vertices) and edges (or arcs)**

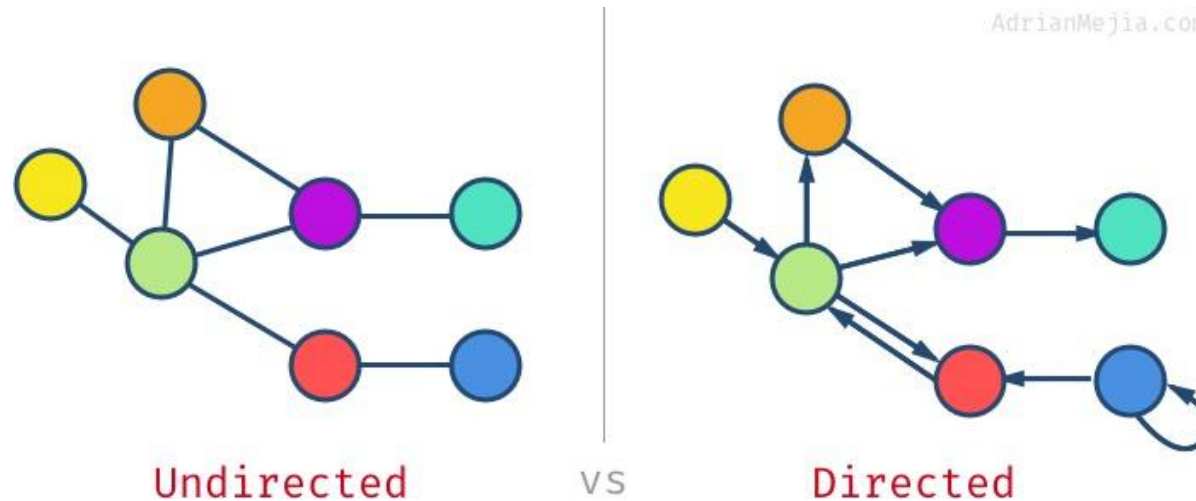
Each node contains an element

Each edge connects two nodes together (or possibly the same node to itself) and may contain an edge attribute



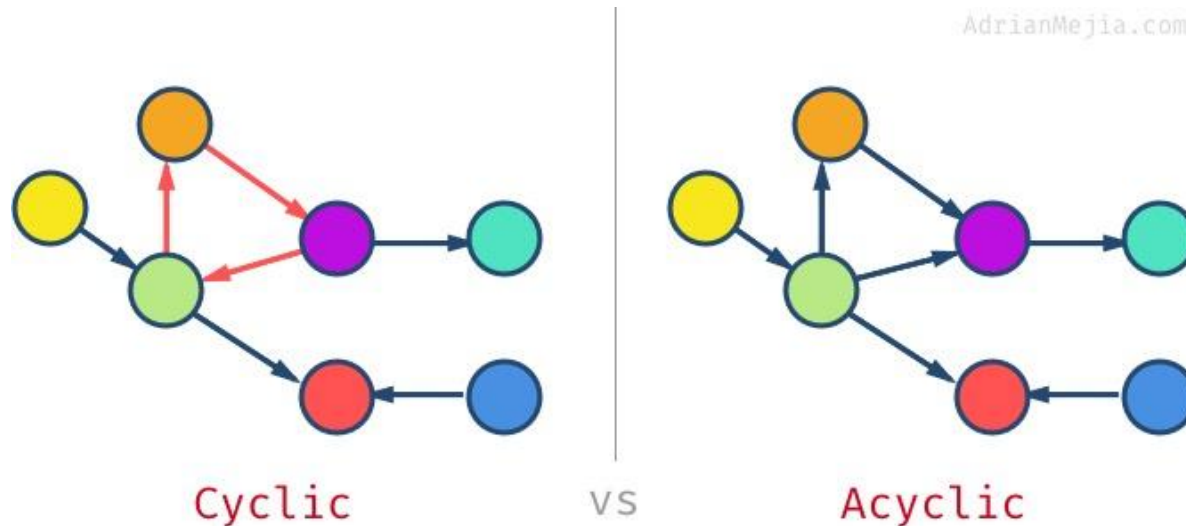
# Introduction to Graph

- The **degree** is the number of edges connected to a vertex. E.g., the purple vertex has a degree of 3 while the blue one has a degree of 1.
- If the edges are bi-directional, then we have an **undirected graph**.
- If the edges have a direction, then we have a **directed graph** or di-graph for short.
- Vertex can have edges that go to itself (e.g., blue node), this is called **self-loop**.



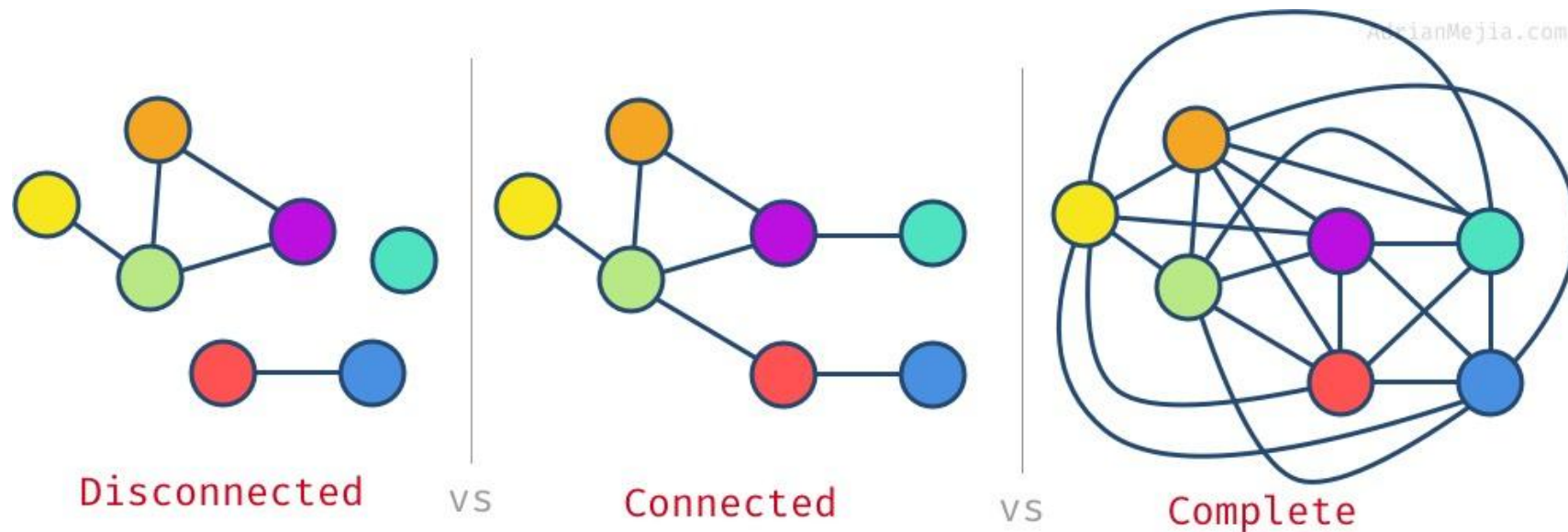
# Introduction to Graph

- A graph can have cycles which means that if you traverse through the node, you could get the same node more than once. The **graph without cycles is called acyclic graph**.
- Also, acyclic undirected graphs are called **tree**. We are going to cover trees in depth in the next post.



# Introduction to Graph

- If all nodes have at least one edge, then we have a **connected graph**.
- When all nodes are connected to all other nodes, then we have a **complete graph**.



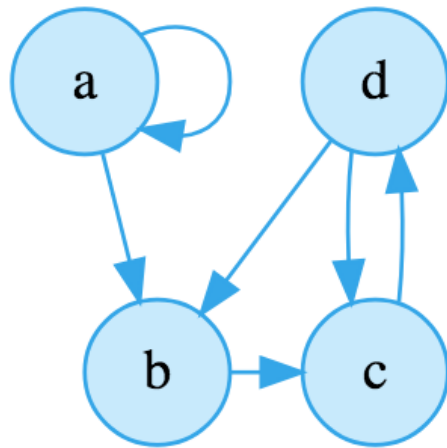
# Graph Terminologies

- The **size** of a graph is the number of *nodes* in it
- The **empty graph** has size zero (no nodes)
- If two nodes are connected by an edge, they are **neighbors** (and the nodes are **adjacent** to each other)
- The **degree of a node** is the number of edges it has
- For directed graphs,
  - If a directed edge goes from node S to node D, we call S the **source** and D the **destination** of the edge
    - The edge is an **out-edge** of S and an **in-edge** of D
    - S is a **predecessor** of D, and D is a **successor** of S
  - The **in-degree** of a node is the number of in-edges it has
  - The **out-degree** of a node is the number of out-edges it has

# Representing graphs

There are two primary ways of representing graph:

- Adjacency list
- Adjacency Matrix



Adjacency Matrix

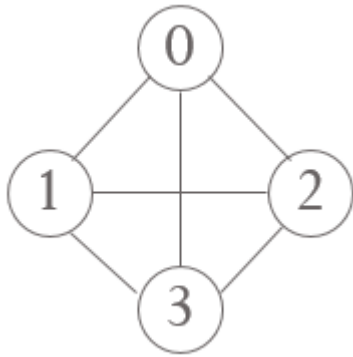
1	a	b	c	d	e
2	a	1	1	-	-
3	b	-	-	1	-
4	c	-	-	-	1
5	d	-	1	1	-

Adjacency List

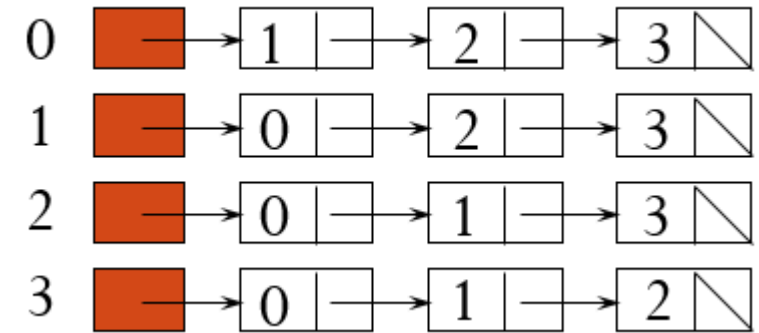
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1  a -> { a b }
2  b -> { c }
3  c -> { d }
4  d -> { b c }
```



# Representing graphs



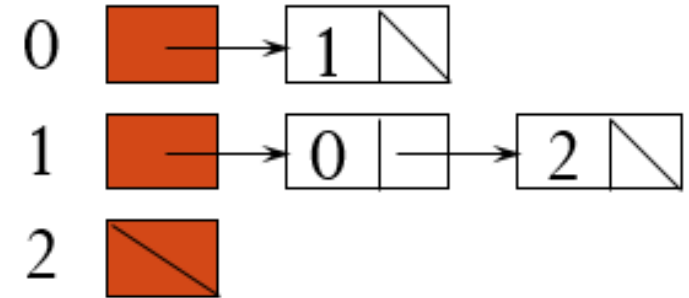
$$\begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix}$$



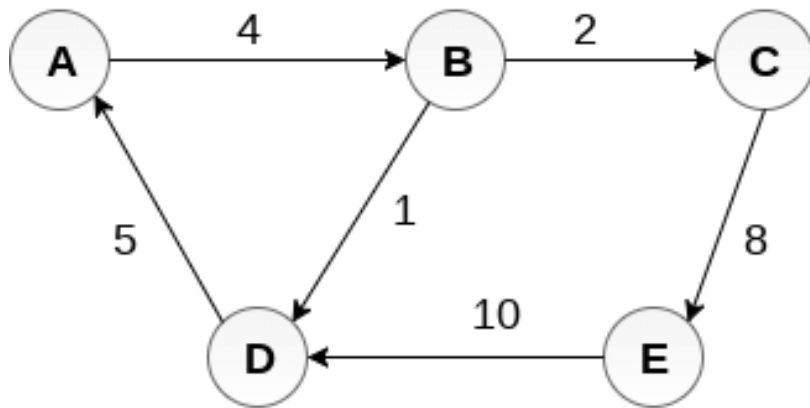
# Representing graphs



$$\begin{bmatrix} 0 & 1 & 0 \\ 1 & 0 & 1 \\ 0 & 0 & 0 \end{bmatrix}$$



# Representing graphs

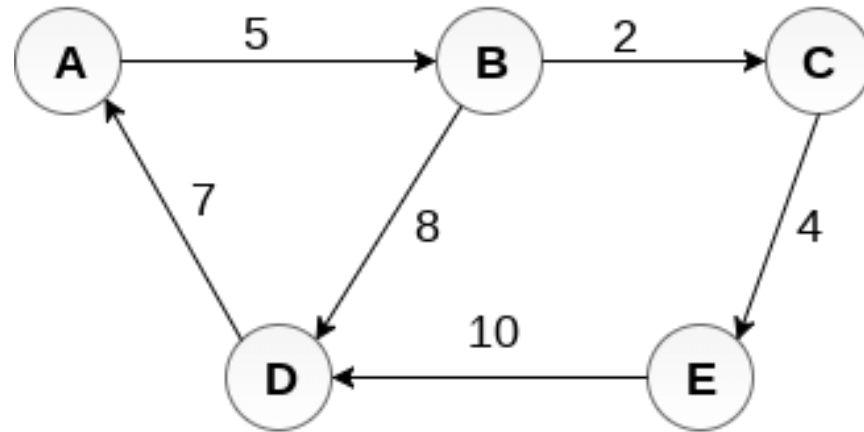


**Weighted Directed Graph**

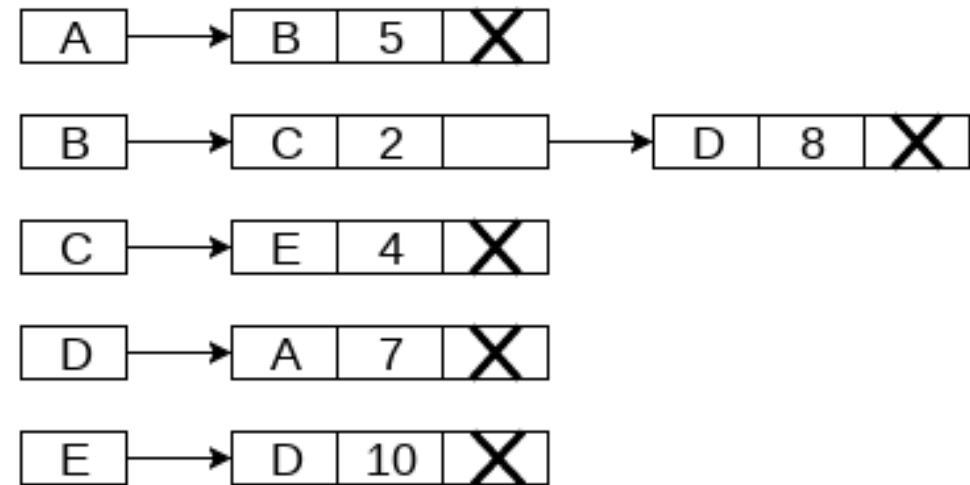
	A	B	C	D	E
A	0	4	0	0	0
B	0	0	2	1	0
C	0	0	0	0	8
D	5	0	0	0	0
E	0	0	0	10	0

**Adjacency Matrix**

# Representing graphs



**Weighted Directed Graph**



**Adjacency List**

# Graph Traversal

- BFS (Breadth First Search)
  - Start from a vertex, visit all the reachable vertices in a breadth first manner
  - Uses Queue for non-recursive implementation
- DFS (Depth First Search)
  - Start from a vertex, visit all the reachable vertices in a depth first manner
  - Uses Stack for non-recursive implementation

# Depth First Search (DFS) Algorithm

## Algorithm

Step 1: SET STATUS = 1 (ready state) for each node in G

Step 2: Push the starting node A on the stack and set its STATUS = 2 (waiting state)

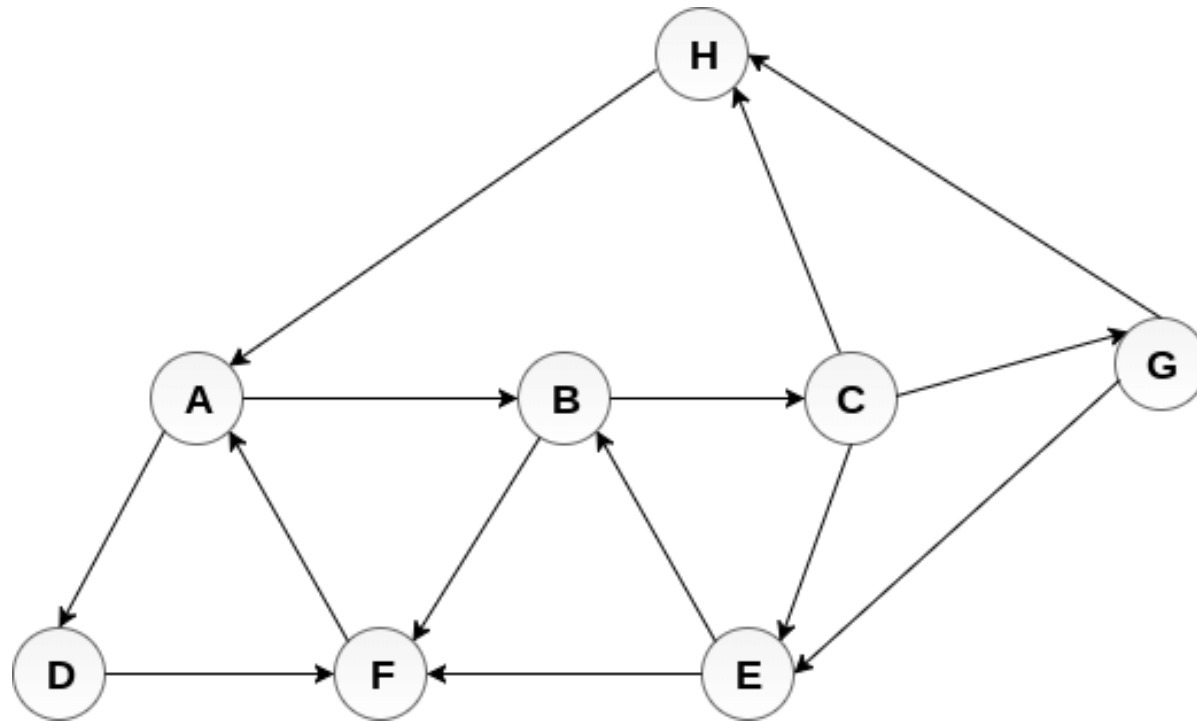
Step 3: Repeat Steps 4 and 5 until STACK is empty

Step 4: Pop the top node N. Process it and set its STATUS = 3 (processed state)

Step 5: Push on the stack all the neighbours of N that are in the ready state (whose STATUS = 1) and set their STATUS = 2 (waiting state)  
[END OF LOOP]

Step 6: EXIT

# DFS - Example



## Adjacency Lists

A : B, D

B : C, F

C : E, G, H

G : E, H

E : B, F

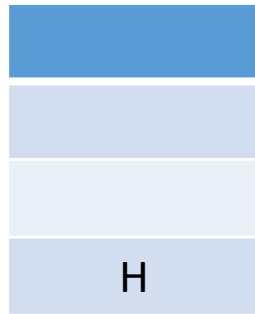
F : A

D : F

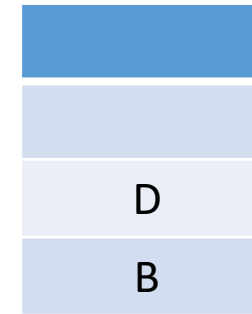
H : A

# DFS - Example

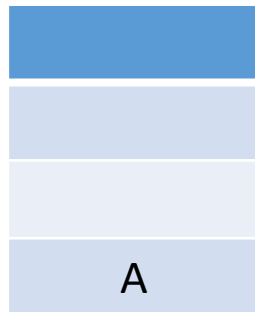
Step1: Push H onto the stack



Step3: **Pop A** and Push all its adjacent



Step2: **Pop H** and Push all its adjacent



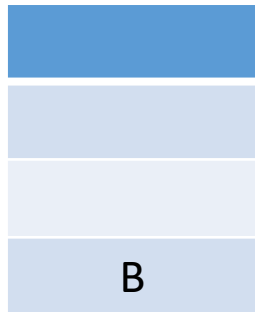
Step4: **Pop D** and Push all its adjacent



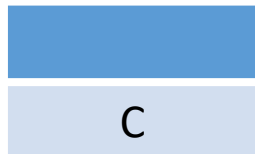


# DFS - Example

Step5: **Pop F** and Push all its adjacent. But A is already visited. So no Push



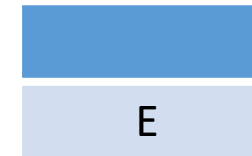
Step6: **Pop B** and Push all its adjacent. F is already visited



Step7: **Pop C** and Push all its adjacent. H is already visited



Step8: **Pop G** and Push all its adjacent. But B, F already visited. So no push



## DFS - Example

Step9: **Pop E** and Push all its adjacent. But B, E is already visited. So no Push



Now the stack is empty. Stop the process. List all the popped elements

**DFS = H → A → D → F → B → C → G → E**

# Breadth First Search (BFS) Algorithm

## Algorithm

Step 1: SET STATUS = 1 (ready state) for each node in G

Step 2: Enqueue the starting node A and set its STATUS = 2 (waiting state)

Step 3: Repeat Steps 4 and 5 until QUEUE is empty

Step 4: Dequeue a node N. Process it and set its STATUS = 3 (processed state).

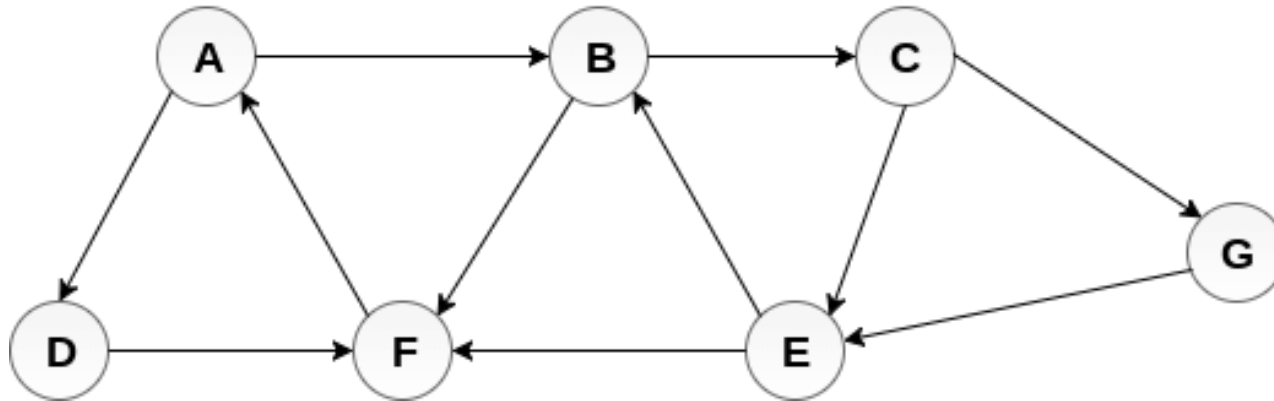
Step 5: Enqueue all the neighbours of N that are in the ready state

(whose STATUS = 1) and set  
their STATUS = 2 (waiting state)

[END OF LOOP]

Step 6: EXIT

# BFS - Example



## Adjacency Lists

A : B, D

B : C, F

C : E, G

G : E

E : B, F

F : A

D : F

# BFS - Example

Step1: Enque A onto the Queue Q1

Q1: 

A				
---	--	--	--	--

Step2: Deque A and Enque all its adjacent

Q1: 

B	D			
---	---	--	--	--

Step3: Deque B and Enque all its adjacent

Q1: 

D	C	F		
---	---	---	--	--

## BFS - Example

Step4: **Deque D** and Enque all its adjacent. But F is already inserted. So no Enque

Q1: 

C	F			
---	---	--	--	--

Step5: **Deque C** and Enque all its adjacent

Q1: 

F	E	G		
---	---	---	--	--

Step6: **Deque F** and Enque all its adjacent. But A already visited so no Enque.

Q1: 

E	G			
---	---	--	--	--

## BFS - Example

Step7: **Deque E** and Enque all its adjacent. But B, F is already inserted. So no Enque

Q1: 

Step8: **Deque G** and Enque all its adjacent. But E already visited

Q1: 

Now queue is empty, Display all deque element is order

**BFS: A, B, D, C, F, E, G**

# Summary

At the end of this session we learned about

- Graphs
- Breadth First Search - BFS
- Depth First Search - DFS