



# **Design & Analysis of Algorithms**

## **Lecture 7**

# **Mathematical Analysis of Non-recursive Algorithms**

# Mathematical Analysis of Non-recursive Algorithms

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- We systematically apply the general framework (discussed earlier) to analyzing the time efficiency of **non-recursive** algorithms.

# Mathematical Analysis of Non-recursive Algorithms

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## General Plan for Analyzing the Time Efficiency

- **1.** Decide on a parameter (or parameters) indicating an input's size.
- **2.** Identify the algorithm's **basic operation**. (As a rule, it is located in the **innermost** loop.)
- **3.** Check whether the **number of times** the **basic operation** is executed depends only on the **size** of an **input**. If it also depends on some additional property, the **worst-case**, **average-case**, and, if necessary, **best-case** efficiencies have to be investigated separately.

# Mathematical Analysis of Non-recursive Algorithms

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## General Plan for Analyzing the Time Efficiency

- **4.** Set up a **sum** expressing the **number of times** the algorithm's **basic operation is executed**.
- **5.** Using standard formulas and rules of sum manipulation, either find a closed-form formula for the count or, at the very least, establish its **order of growth**.

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## Example 1

- Consider the problem of finding the value of the largest element in a list of  $n$  numbers.
- For simplicity, we assume that the list is implemented as an **array**.

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**ALGORITHM** *MaxElement*( $A[0..n - 1]$ )

//Determines the value of the largest element in a given array

//Input: An array  $A[0..n - 1]$  of real numbers

//Output: The value of the largest element in  $A$

*maxval*  $\leftarrow A[0]$

**for**  $i \leftarrow 1$  **to**  $n - 1$  **do**

**if**  $A[i] > \textit{maxval}$

*maxval*  $\leftarrow A[i]$

**return** *maxval*

- input's size here is the number of elements in the array, i.e.,  $n$ .

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- Algorithm's for loop-
  - There are two operations in the loop's body: the comparison  $A[i] > \text{maxval}$  and
  - the assignment  $\text{maxval} \leftarrow A[i]$ .
- Which of these two operations should we consider **basic**?



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- Since the comparison is executed on each repetition of the loop and the assignment is not, we should consider the **comparison** to be the algorithm's basic operation.
- Note that the number of comparisons will be the same for all arrays of size  $n$ ; therefore, in terms of this metric, there is **no need** to distinguish among the **worst**, **average**, and **best** cases here.

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- Let us denote  $C(n)$  the number of times this comparison is executed.
- The algorithm makes one comparison on each execution of the loop, which is repeated for each value of the loop's variable  $i$  within the bounds 1 and  $n - 1$ , inclusive.
- Therefore, we get the following sum for  $C(n)$ :

$$C(n) = \sum_{i=1}^{n-1} 1$$

$$C(n) = \sum_{i=1}^{n-1} 1 = n - 1 \in \Theta(n)$$

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## EXAMPLE 2

- Consider the element uniqueness problem:  
check whether all the elements in a given  
array of  $n$  elements are distinct.
- This problem can be solved by the following  
straightforward algorithm.

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**ALGORITHM** *UniqueElements*( $A[0..n - 1]$ )

//Determines whether all the elements in a given array are distinct

//Input: An array  $A[0..n - 1]$

//Output: Returns “true” if all the elements in  $A$  are distinct

//           and “false” otherwise

**for**  $i \leftarrow 0$  **to**  $n - 2$  **do**

**for**  $j \leftarrow i + 1$  **to**  $n - 1$  **do**

**if**  $A[i] = A[j]$  **return** false

**return** true

- input's size here is again  $n$ , the number of elements in the array.

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- Since the innermost loop contains a single operation (the **comparison** of two elements), we should consider it as the algorithm's **basic** operation.
- **Note**, however, that the number of comparisons depends not only on **n** but also on whether there are equal elements in the array and, if there are, which **array positions** they occupy.
- Let us investigate the **worst case** scenario.

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$$\begin{aligned}C_{worst}(n) &= \sum_{i=0}^{n-2} \sum_{j=i+1}^{n-1} 1 = \sum_{i=0}^{n-2} [(n-1) - (i+1) + 1] = \sum_{i=0}^{n-2} (n-1-i) \\&= \sum_{i=0}^{n-2} (n-1) - \sum_{i=0}^{n-2} i = (n-1) \sum_{i=0}^{n-2} 1 - \frac{(n-2)(n-1)}{2} \\&= (n-1)^2 - \frac{(n-2)(n-1)}{2} = \frac{(n-1)n}{2} \approx \frac{1}{2}n^2 \in \Theta(n^2).\end{aligned}$$

# Mathematical Analysis of Non-recursive Algorithms

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## Some formula to consider

$$\sum_{i=l}^u ca_i = c \sum_{i=l}^u a_i,$$

$$\sum_{i=l}^u (a_i \pm b_i) = \sum_{i=l}^u a_i \pm \sum_{i=l}^u b_i,$$

$$\sum_{i=l}^u 1 = u - l + 1 \quad \text{where } l \leq u \text{ are some lower and upper integer limits,}$$

$$\sum_{i=0}^n i = \sum_{i=1}^n i = 1 + 2 + \cdots + n = \frac{n(n+1)}{2} \approx \frac{1}{2}n^2 \in \Theta(n^2).$$

# References

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- **Chapter 2:** Anany Levitin, “Introduction to the Design and Analysis of Algorithms”, Pearson Education, Third Edition, 2017
- **Chapter 2:** Thomas H. Cormen, Charles E. Leiserson, Ronald L. Rivest and Clifford Stein, “Introduction to Algorithms”, MIT Press/PHI Learning Private Limited, Third Edition, 2012.



# Homework

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## **Analyze the following algorithms:**

- Matrix multiplication of dimensions  $n \times n$ .
- Finding the number of binary digits in the binary representation of a positive decimal integer.