

# Code Spectrum

## 1. Monotonic Function

```
#include <bits/stdc++.h>
using namespace std;
int main() {
    int x;
    cin>>x;
    int f = pow(x,3)+2*pow(x,2)+3*x+5;
    cout<<f<<endl;
    return 0;
}
```

## 2. Seh Lenge Thoda

```
#include <bits/stdc++.h>
using namespace std;

int main()
{
    int n;
    cin>>n;
    vector<pair<int,int>> arr(n);
    for(int i=0;i<n;i++) cin>>arr[i].first;
```

```

    for(int i=0;i<n;i++) cin>>arr[i].second;

    int sum=0;
    for(int i=0;i<n;i++)
    {
        sum+=(arr[i].first*arr[i].second);
    }
    cout<<sum<<endl;
}

```

## 3. Basant - The festival of roses

```

#include <bits/stdc++.h>
using namespace std;
int main() {
    int n,x;
    cin>>n>>x;
    vector<pair<int,string>> freshers(n);
    for(int i=0;i<n;i++){
        cin>>freshers[i].second>>freshers[i].first;
    }
    sort(freshers.begin(),freshers.end());
    reverse(freshers.begin(),freshers.end());
}

```

```
    cout<<freshers[x].second<<endl;
    return 0;
}
```

### Alternate Approach:

```
#include <bits/stdc++.h>
using namespace std;
int main() {
    int n,x;
    cin>>n>>x;
    vector<pair<int,string>> freshers(n);
    for(int i=0;i<n;i++){
        cin>>freshers[i].second>>freshers[i].first;
    }
    sort(freshers.begin(),freshers.end());
    cout<<freshers[n-x].second<<endl;
    return 0;
}
```

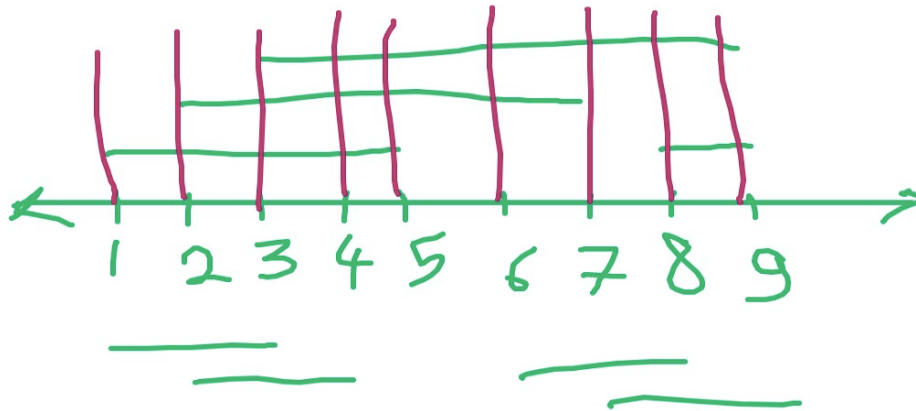
## 4. Despo Raju

```
#include <bits/stdc++.h>
using namespace std;
int main() {
    int q;
    cin>>q;
    while(q--){
        int n;
```

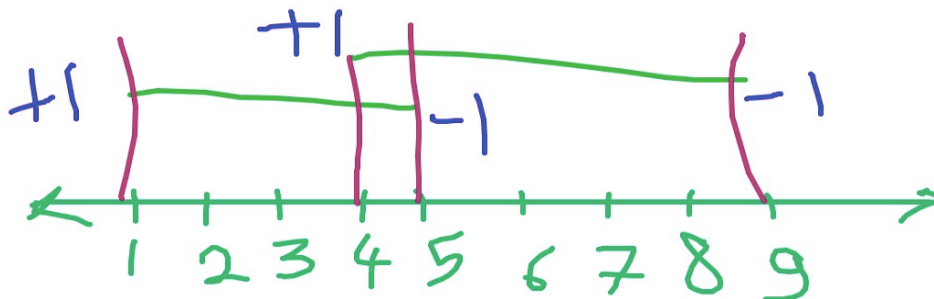
```
    cin>>n;
    string s;
    cin>>s;
    long long int swaps=0;
    int countOne=0;
    for(int i=0;i<n;i++){
        if(s[i]=='1'){
            countOne++;
        }
        else{
            swaps+=countOne;
        }
    }
    cout<<swaps<<endl;
}
return 0;
}
```

## 5. Lalit in Trouble

We need to find maximum number of overlapping classes at any instant.



We can consider the overlapping classes at the start and end-times. No need to consider all the points on number line.



First, take a variable for counting overlapping segments,  $cnt=0$

When u find a start time of any class, just increase  $cnt$  by 1.

When u find a end time of any class, just decrease  $cnt$  by 1.

And when calculating  $cnt$ , also store the maximum number of overlapping segments as the answer.

```
#include <bits/stdc++.h>
using namespace std;
int main() {
    int n;
    cin>>n;
    int s,f;
    vector<pair<int,int> > vec;
    for(int i=0; i<n; i++)
    {
        cin>>s>>f;
        vec.push_back({s,+1});
        vec.push_back({f,-1});
    }
    sort(vec.begin(), vec.end());
    int cnt=0; // number of overlapping classes/intervals
    int ans=0; // max. no. of overlapping classes/intervals
    int sz=vec.size();

    for(int i=0; i<sz; i++)
    {
        cnt=cnt+vec[i].second;
        ans=max(ans,cnt);
    }
    cout << ans-1 <<'\n';
    return 0;
}
```

# 6. Raju Bhai and Group Photo

## Naive Idea (Brute Force Approach):

Go through all possible subarrays and if sum is divisible by x, update ans as maximum size of subarray.

```
#include <bits/stdc++.h>
using namespace std;
int main() {
    int n,x;
    cin>>x>>n;

    vector<int> vec(n);
    for(int i=0; i<n; i++)
    {
        cin>>vec[i];
    }

    int ans=0;

    for(int l=0; l<n; l++)
    {
        for(int r=l; r<n; r++)
        {
```

```

        long long sum=0;
        for(int i=1; i<=r; i++)
        {
            sum+=vec[i];
        }
        if(sum % x == 0)
        {
            ans=max(ans, r-l+1);
        }
    }

    cout<<ans;
    return 0;
}

```

Time Complexity:  $O(N^3)$

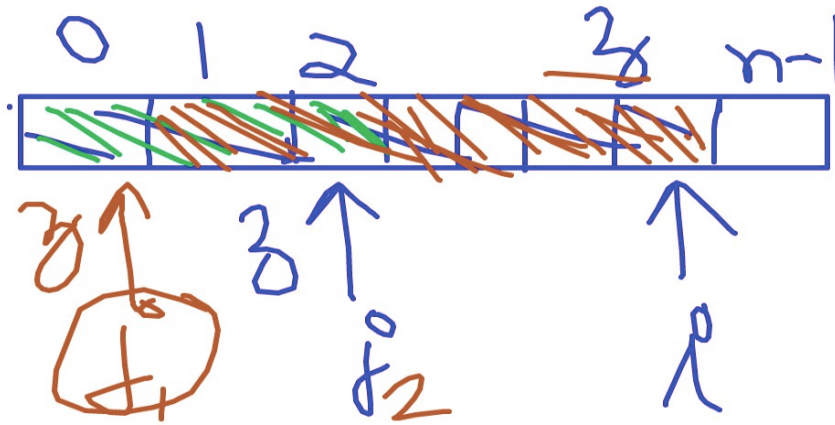
In general, you can perform  $10^7 - 10^8$  operations in 1 second.  
 Here, in worst case,  $N = 10^5 \Rightarrow N^3 = 10^{15}$ . **So, it gives TLE.**

### **A fast or optimised approach:**

For every prefix, store the prefix sum mod x.

This prefix sum will always lie between  $[0, x-1]$ .





1. If for any prefix ending at index  $i$ , if there exists any prefix ending at  $j$  ( $j < i$ ) and  $\text{pref}[i] == \text{pref}[j]$ , then the subarray from  $[j+1, i]$  should be divisible by  $x$ .
2. Suppose there are 2 or more than 2 such possible values of  $j$ , satisfying the above criteria, then I need to consider only the smallest value of  $j$ .

```
#include <bits/stdc++.h>

using namespace std;

int main() {
    int n,x;
    cin>>x>>n;

    vector<int> vec(n);
    for(int i=0; i<n; i++)
```

```

{
    cin>>vec[i];
}

vector<int> low(x);
// low[y] = The smallest value of j,
// such that sum of prefix ending at j mod x = y

for(int i=0; i<x; i++)
{
    low[i]=-1; // -1 indicates no such prefix is available
}

int sum=0;
int ans=0;
for(int i=0; i<n; i++)
{
    //      (a+b)%x = ((a%x) + (b%x))%x
    sum=((sum%x)+(vec[i]%x))%x;
    if(low[sum]!=-1)
    {
        int j=low[sum];
        ans=max(ans, i-j);
    }
    else
    {
        low[sum]=i;
    }
    if(sum%x==0)

```

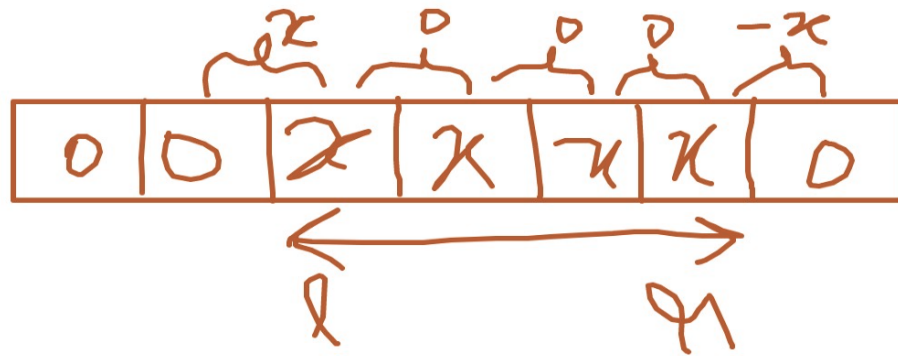
```
    {  
        ans=max(ans, i+1);  
    }  
}  
  
cout<<ans;  
  
return 0;  
}
```

**Time Complexity:**  $O(n) + O(x) = O(\max(n, x))$

**Follow Up:** What if  $x \leq 10^9$ ?

The problem is you can't create such a large array.  
(Use map)

## 7. The Endgame



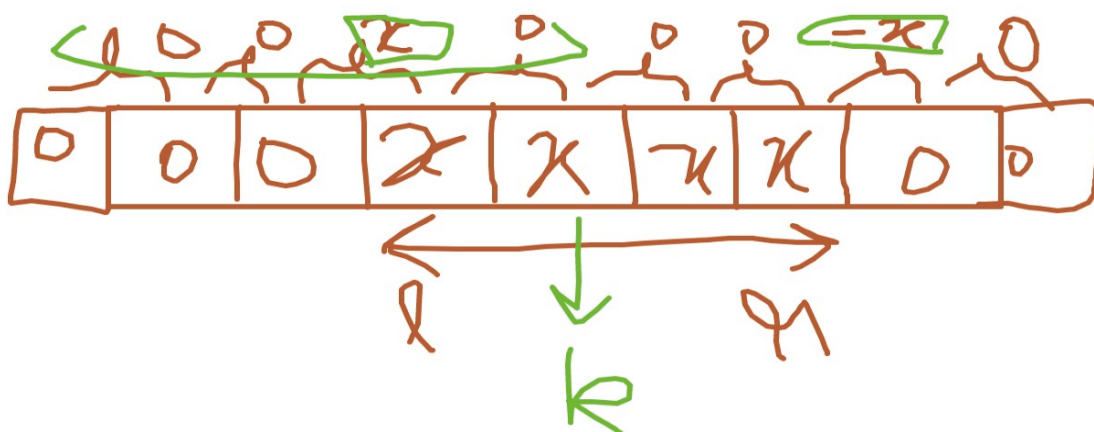
If we build a **difference array** diff, (an array containing difference of consecutive elements),

If we need to increase all elements of the given array in range  $[L, R]$ , then in difference array, we need to make only 2 changes:

1. Increase  $\text{diff}[L]$  by  $x$
2. Decrease  $\text{diff}[R+1]$  by  $x$

Now, each of the  $p$  operations can be performed in  $O(1)$ .

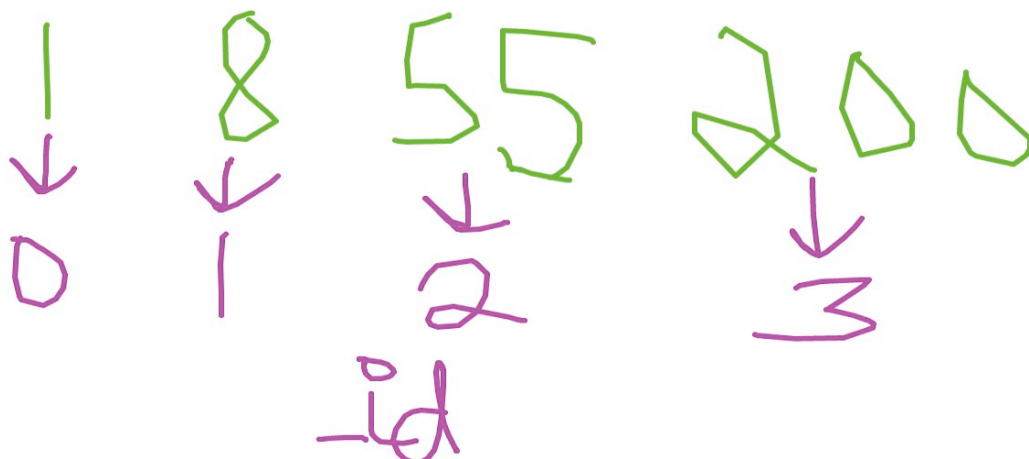
For all  $p$  operations, time complexity is:  $O(p)$



To find  $\text{arr}[k]$ , just find the prefix sum of the difference array upto  $k^{\text{th}}$  index.

Now, the problem is, we can't create a difference array of size  $10^9$  (Because, you can create integer arrays of size  $10^6$ - $10^7$ )

- Use map and give an id to all the unique elements  
(Since, there are  $\leq 10^3$  distinct elements)



- See Setter's code in Editorial

### Follow Up:

Suppose, there was no such constraint on distinct values of  $l$  and  $r$ . Then, how would you approach ?

- You use a map and use `lower_bound()`

See from here:

1. `lower_bound()` in vector:

[https://www.geeksforgeeks.org/upper\\_bound-and-lower\\_bound-for-vector-in-cpp-stl/](https://www.geeksforgeeks.org/upper_bound-and-lower_bound-for-vector-in-cpp-stl/)

2. `lower_bound()` in map:

[https://www.geeksforgeeks.org/map-lower\\_bound-function-in-c-stl/](https://www.geeksforgeeks.org/map-lower_bound-function-in-c-stl/)

- See tester's code in Editorial