

GRAPHS

Day 4

(DSU and MST)

Youtube link :

<https://youtu.be/QHCGQ371eXs>

Contents:

1. DSU and related operations
2. Applications of DSU
3. Kruskal's Algorithm for MST

DSU (Disjoint Set Union)

- A data structure used for performing the union of disjoint sets very efficiently
- **Disjoint sets:** Any 2 sets are disjoint when their intersection is an empty set.

Eg. $A = \{ 1, 2, 6 \}$

$B = \{ 4, 5 \}$

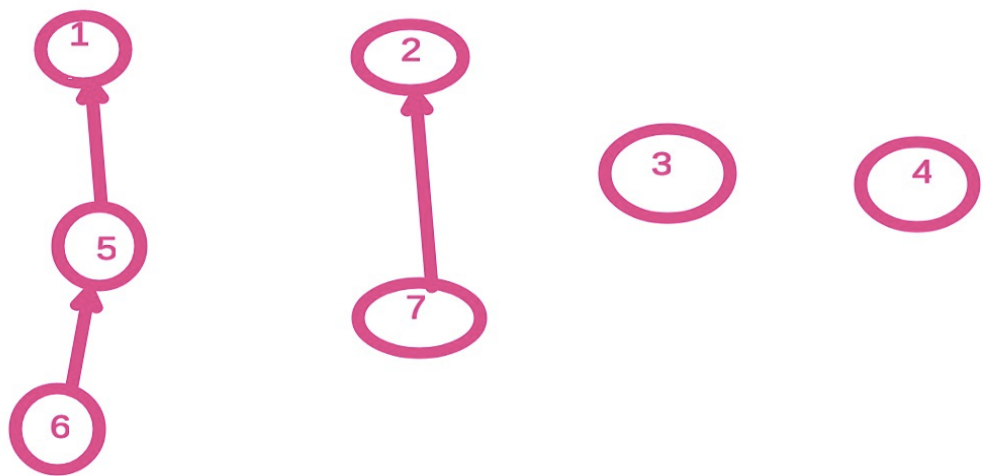
$A \cap B = \phi$ (empty set)

So, A and B are disjoint sets

3 operations:

1. `make_set()` ; // To create DSU initially
2. `find_set()`; // Find the set to which an element belongs (To find representative of a set)
3. `union_set()`; // Take Union of 2 sets or merge 2 sets

**Representative of a set is the topmost element :
(Parent of representative is that element itself)**



Naive implementation

```
const int MAX=100005;
int par[MAX];

void make_set(int n)
{
    for(int i=0; i<n; i++)
    {
        par[i]=i;
    }
}
```

```
int find_set(int x)
{
    if( par[x] == x)
    {
        return x;
    }
    return find_set(par[x]);
}
```

**// You can also write this in a while loop like:
(Both will work the same way, use anyone which you like)**

```
int find_set(int x)
{
    while( par[x] != x)
    {
        x=par[x];
    }
    return x;
}
```

```
void union_set(int a, int b)
{
    int p1 = find_set(a);
    // representative of set of a
```

```

    int p2 = find_set(b);
    // representative of set of b
    if ( p1 != p2)
    {
        par[p1]=p2;
    }
}

```

Time complexity: $O(d)$

Where d = depth (distance from the representative)

Faster Implementation

1. Path compression in find_set():

(Store the representative in parent directly)

```

int find_set(int x)
{
    if( par[x] == x)
    {
        return x;
    }
    par[x]=find_set(par[x]);
    return par[x];
}

```

It reduces time complexity to $O(\log N)$ approximately

2. Union by Rank in union_set() :

- Create another array for size of set. (rnk[])
- Make smaller set as a child of larger set.

```
int rnk[ 100005];
void make_set(int n)
{
    for(int i=0; i<n; i++)
    {
        par[i] = i;
        rnk[i] = 1;
    }
}

void union_set(int a, int b)
{
    int p1 = find_set(a);
    int p2 = find_set(b);
    if ( p1 == p2)
        return;
    if (rnk[p1] > rnk[p2])
    {
        par[p2]=p1;
        rnk[p1] = rnk[p1] + rnk[p2];
    }
    else
    {
```

```
    par[p1] = p2;  
    rnk[p2] = rnk[p1] + rnk[p2];  
}  
}
```

Time Complexity of find_set() and union_set(), when you use both path compression and union by rank:

$O(\alpha(N))$

where $\alpha(N)$ = Inverse Ackermann Function

(Approximately , it is equal to $O(1)$)

Try this problem:

<https://www.hackerrank.com/challenges/merging-communities/problem>

Solution:

```
const int MAX=100005;  
int par[MAX];  
int rnk[MAX];  
void make_set(int n)  
{  
    for(int i=1; i<=n; i++)  
    {  
        par[i] = i;  
        rnk[i] = 1;  
    }  
}
```

```
int find_set(int x)
{
    if( par[x] == x)
    {
        return x;
    }
    par[x]=find_set(par[x]);
    return par[x];
}

void union_set(int a, int b)
{
    int p1 = find_set(a);
    int p2 = find_set(b);
    if ( p1 == p2)
        return;
    if (rnk[p1] > rnk[p2])
    {
        par[p2]=p1;
        rnk[p1] = rnk[p1] + rnk[p2];
    } else
    {
        par[p1] = p2;
        rnk[p2] = rnk[p1] + rnk[p2];
    }
}
```

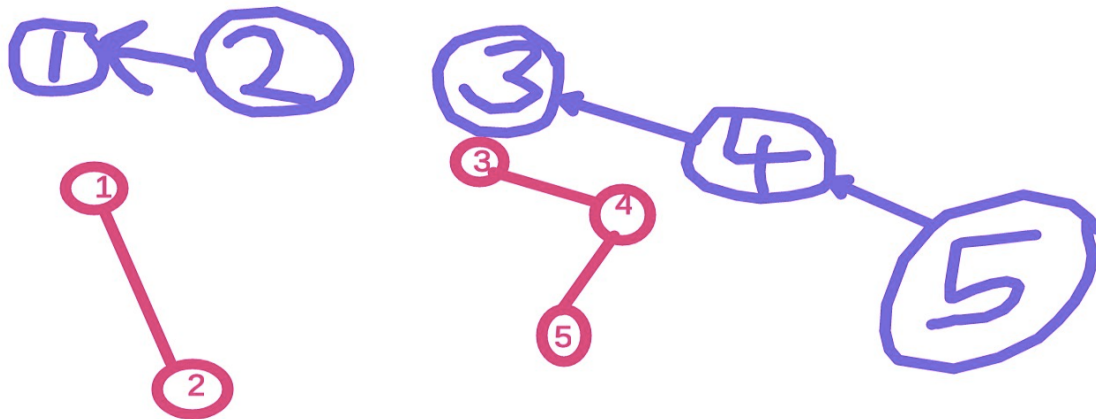
```

int main()
{
    int n,q;
    cin>> n >> q;
    make_set(n);
    while(q--)
    {
        char ch;
        cin>>ch;
        if(ch=='M')
        {
            int a,b;
            cin>>a>>b;
            union_set(a,b);
        }
    else {
        int a;
        cin>>a;
        int representative = find_set(a);
        cout<<rnk[representative]<<'\n';
    }
}
return 0;
}

```

Applications of DSU

1. Finding the number of components in a graph



Try:

<https://cses.fi/problemset/task/1676>

Solution:

```
int mx=1;

void union_set(int a, int b)
{
    int p1 = find_set(a);
    int p2 = find_set(b);
    if ( p1 == p2)
        return;
    if (rnk[p1] > rnk[p2])
    {
        par[p2]=p1;
        rnk[p1] = rnk[p1] + rnk[p2];
        mx=max(mx,rnk[p1]);
    } else
```

```

{
    par[p1] = p2;
    rnk[p2] = rnk[p1] + rnk[p2];
    mx=max(mx,rnk[p2]);
}
}

int components = n;
cin>> n >> m;
make_set(n);
int u,v;
while(m--)
{
    cin>>u>>v;
    u--;
    v--;
    if( find_set(u) == find_set(v) )
    {
        cout<<components<< " "<<mx<<'\n';
        continue;
    }
    union_set(u,v);
    components--;
    cout<<components<< " "<<mx<<'\n';
}

```

2. Finding a cycle in undirected graph

Code:

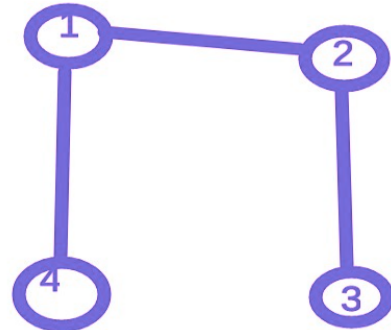
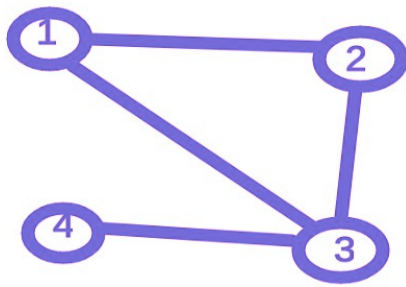
```
make_set(n);
for( auto e: edges)
{
    int u=e.first;
    int v=e.second;
    if ( find_set(u) == find_set(v) )
    {
        cout<<"Cycle found";
        break;
    }
    union_set(u,v);
}
```

MST (Minimum Spanning Tree)

Spanning Tree of a graph :

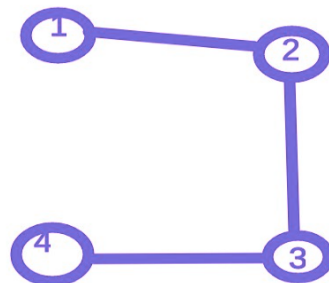
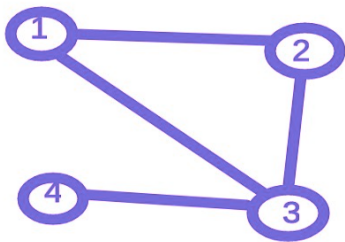
A spanning tree is a tree containing all nodes of the given graph and is a subgraph of the given graph.

Example of not a spanning tree:



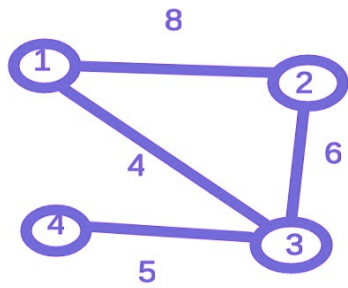
(Since edge (1,4) is not present in the original graph)

Example of a Spanning Tree:

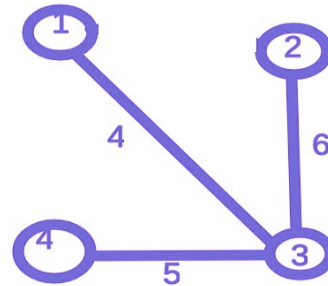


Minimum Spanning Tree:

We need to find a spanning tree whose sum of edges is minimum.



Cost of MST
 $= 4 + 6 + 5$
 $= 15$



Kruskal's Algorithm (A Greedy Algorithm)

1. Sort edges in increasing order of weights
- 2.

```
make_set(n);
int sum=0;
int cnt=0;
for (auto e: edges)
{
    int u=e.first;
    int v=e.second.first;
    int w=e.second.second;
    if ( find_set(u) == find_set(v) )
    {
        continue;
    }
    union_set(u,v);
    sum = sum + w;
}
```

```

        cnt++;
    }

    if( cnt == n-1)
        cout<<sum<<'\n';
    else
    {
        cout<<"NO MST Possible";
    }
}

```

Try to solve:

<https://cses.fi/problemset/task/1675>

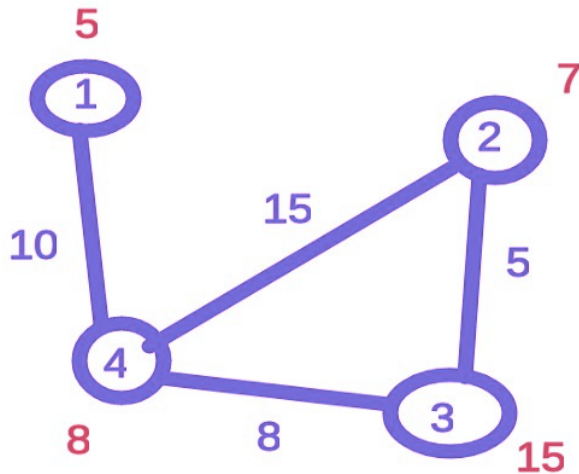
Q. Find a spanning tree of a given graph with minimum products of weights?

Solution:

- Find MST with Kruskal's algorithm and print the products of its weights.

Q. Given n cities and m pipe connections between those cities . For digging a well in a city i, it costs C[i]. Find minimum cost to provide water supply to all.

Example:

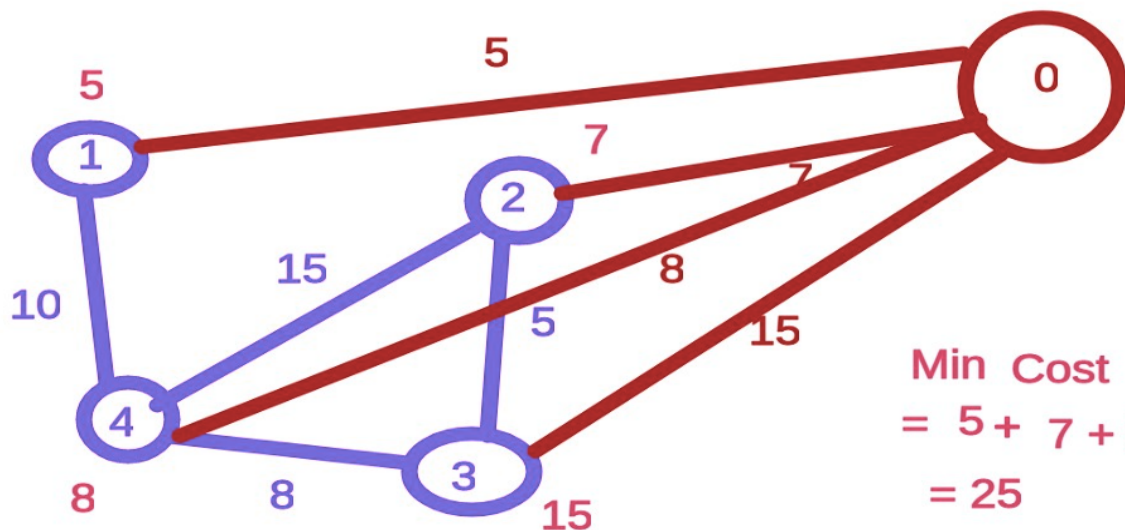


$$\begin{aligned} \text{Min Cost} &= 5 + 7 + 5 + 8 \\ &= 25 \end{aligned}$$

Solution Approach:

Consider a imaginary node 0 and connect it to all remaining nodes with a edge cost of $C[i]$.

Now, find the MST of this new graph.



$$\begin{aligned} \text{Min Cost} &= 5 + 7 + 5 + 8 \\ &= 25 \end{aligned}$$