

Exercise 11

Greedy Algorithm

11. Write a C++ program for implementation of following algorithms
a) Knapsack algorithm b) Job sequence with dead line

Objective:

The objective of this exercise is to enable you to perform knapsack algorithm and job sequence with dead line algorithm.

Procedure and Description

a) Knapsack algorithm

The knapsack problem is a problem in combinatorial optimization (finding an optimal object from a finite set of objects). Given a set of items, each with a weight and a value, determine the number of each item to include in a collection so that the total weight is less than or equal to a given limit and the total value is as large as possible. It derives its name from the problem faced by someone who is constrained by a fixed-size knapsack and must fill it with the most valuable items.

Algorithm:

n: No. of objects

m&u: capacity

i&j: used as index of objects

P[i]: profit index i

W[j]: Weight index j

X: benefit

The steps for knapsack problem are given below

Step 1: Greedy knapsack (m, n)

Step 2: p [1: n] and w [1: n] contain the profits and weights respectively

Of the n objects ordered such that $p[i]/w[i] \geq p[i+1]/w[i+1]$

Step 3: m is the knapsack size and x [1: n] is the solution vector

Step 4: for i: =1 to n do x[i]:=0.0; //Initialize x

Step 5: u: =m;

Step 6: for i: =1 to n do

Step 7: if (w[i]>u) then break;

Step 8: x[i]:=1.0; U: =U-w[i];

If (i ≤ n) then x[i]:=u/w[i];

Step 9: end

Expected output:

After executing program enter set of items, each with ratio, weight and a profit, then by performing algorithm steps it determines the benefit.

Input: capacity: 7

ratio:2.5	Wgt: 2	profit: 5.0
ratio: 2.0	wgt: 2	profit: 4.0
ratio: 1.67	wgt: 3	profit: 5.0
ratio: 1.5	wgt: 4	profit: 6.0

Output: Cumulative benefit: 14.0

b) Job sequence with dead line

In this job sequence with dead line, given a set of n jobs, each having a “deadline” (an integer) $d[i]$ and a profit $p[i]$ associated with it. For a job i the profit is earned if and only if the job is completed within its deadline. Each job takes one unit of time on a machine (processor) and only one machine is available. We want to maximize the profit. The jobs are arranged in decreasing order of profit in an array.

Algorithm:

n : set of jobs

i & j : used as index of jobs

$d[i]$: an integer deadline

$P[i]$: job i profit

$J[i]$: i th job with optimal solution

r & k : variables used to store result values.

The steps for Job sequence with dead line are given below

Step 1: Algorithm JS (d, j, n)

Step 2: $d[i] \geq 1, 1 \leq i \leq n$ are the deadlines, $n \geq 1$. The jobs are ordered such that $p[1] \geq p[2] \geq \dots \geq P[n]$.

Step 3: $J[i]$ is the i th job in the optimal solution, $1 \leq i \leq k$.

Step 4: Also, at termination $d[J[i]] \leq d[J[i+1]]$, $1 \leq i \leq k$.

Step 5: $d[0] := J[0] := 0$; //initialize

Step 6: $J[1] := 1$; include job 1

$K := 1$;

Step 7: for $i := 2$ to n do

Step 8: Consider jobs in non-increasing order of $p[i]$
Find position for i and check feasibility for insertion

Step 9: $r := k$;

Step 10: While $((d[J[r]] > d[i]) \text{ and } (d[J[r]] \neq r))$ do $r := r - 1$;

Step 11: if $((d[J[r]] \leq d[i]) \text{ and } (d[i] > r))$ then

Step 12: Insert i into $J[]$.

Step 13: for $q := k$ to $(r+1)$ step -1 do $J[q+1] := J[q]$;
 $J[r+1] := i$; $k := k+1$;

Step 14: end

Step 15: return k ;

Expected output:**Input:**

After executing program enter the Time slot, number of jobs, profit deadline and profit/Time.

Example:

Time slots 1, 2, and 3. (Slot 0 is sentinel)

Number of jobs=5

Job (i)	Profit	Deadline	Profit/Time
A	100	2	100
B	19	1	19
C	27	2	27
D	25	1	25
E	15	33	15

Output:

Jobs done: Job c at time 1

Job a at time 2

Job e at time 3

Number of jobs done: 3, total profit: 142