

#include <senior week 11>

DP Extra

Today, we'll go over some extra DP topics.

Game Theory

In game theory, given some current state of the game, we want to figure out who wins. In a two-player game, there are two positions: a *winning* position or a *losing* position.

A winning position means a position where the player who goes first wins.

A losing position means a position where the player who goes first loses, i.e. the person who goes later wins.

After handling bases cases, when we want to transition there are 2 possibilities:

1. If I can make a move and go to a losing position, I am currently in a winning position
2. If no matter what move I make, the game transitions to a winning position, I am currently in a losing position.

Because one must consider all possible next positions, generally we loop based on position rather than move. So it should look something like this:

```
for (position):  
    for (move):  
        consider move
```

Try this: [Educational DP Contest AtCoder K - Stones - DMOJ: Modern Online Judge](#)

```
#include <bits/stdc++.h>  
using namespace std;  
  
int main() {  
    int n, k; cin >> n >> k;  
    bool dp[k+1];  
  
    memset(dp, false, sizeof(dp));  
  
    int stones[n];  
    for (int i = 0; i < n; i++) cin >> stones[i];  
  
    for (int i = 1; i <= k; i++) {  
        for (int j = 0; j < n; j++){  
            int s = stones[j];  
            if (i-s >= 0 && dp[i-s] == false) {
```

```

        dp[i] = true;
    }
}

if (dp[k] == true) {
    cout << "First";
} else {
    cout << "Second";
}
return 0;
}

```

DP w/ states

In DP, we don't know what our last move was, which is a problem when encountering situations where we have restrictions on possible moves. For example, not being able to use the same move twice in a row.

In this case, we want to assign specific states to a DP position. In general, we'll be filling out a DP[position][state].

The idea of states is to somehow update them to make sure we never reach an illegal position.

Question: [Educational DP Contest AtCoder C - Vacation - DMOJ: Modern Online Judge](#)

How can we assign states to make sure we never violate the rule of 'don't do the same activity twice in a row?'

ANS: have states 1,2,3 which represent activity A,B,C. Then we don't want the same activity twice, so transition looks like this:

$DP[d][1] = \max(DP[d-1][2], DP[d-1][3]) + \text{activity A today}$

$DP[d][2] = \max(DP[d-1][1], DP[d-1][3]) + \text{activity B today}$

$DP[d][3] = \max(DP[d-1][1], DP[d-1][2]) + \text{activity C today}$

```

#include <bits/stdc++.h>

using namespace std;

```

```

int n,a,b,c;
int dp[100000][3];

int main() {

    cin >> n >> a >> b >> c;
    dp[0][0] = a;
    dp[0][1] = b;
    dp[0][2] = c;
    for (int i = 1; i < n; i++){
        cin >> a >> b >> c;
        dp[i][0] = a + max(dp[i-1][1],dp[i-1][2]);
        dp[i][1] = b + max(dp[i-1][0],dp[i-1][2]);
        dp[i][2] = c + max(dp[i-1][0],dp[i-1][1]);
    }
    cout << max(dp[n-1][0],max(dp[n-1][1],dp[n-1][2]));

    return 0;
}

```

Interval DP

This is another popular type of DP. $DP[\text{left}][\text{right}] = \text{ans}$.

It's common in this type of problem to loop by range, so it looks like this:

for (range):

 for ([i to i+range]):

 consider range

When there are 2 lists, a common way to solve it is to use $DP[a][b] = \text{answer for list1}[0\dots a] + \text{list2}[0\dots b]$.

[Educational DP Contest AtCoder L - Deque - DMOJ: Modern Online Judge](#)

We can set $DP[L][R] = (\text{First player} - \text{Second player})$, or (Player about to move - other player).

Then for each transition, we either take the left or the right most values.

So $[L\dots R] \leftarrow [L+1\dots R]$ or $[L\dots R-1]$.

If the first player (say A) takes from the left, we are considering $[L+1\dots R]$. Now it is the second player (say B)'s turn to move. Now it is B's turn for the array $[L+1\dots R]$. From our DP, we know the answer to this subproblem. So $B-A = DP[L+1][R]$. We want to solve for A-B though, so $-DP[L+1][R] = A-B$. Adding on the value taken, we have:

$DP[L][R] = arr[L] - DP[L+1][R].$

Similarly, the other case is $DP[L][R] = arr[R] - DP[L][R-1].$

So our transition is:

$DP[L][R] = \max(arr[L]-DP[L+1][R], arr[R]-DP[L][R-1]).$

```
#include <bits/stdc++.h>

using namespace std;

int n, arr[3000];
long long dp[3000][3000];

int main() {

    cin >> n;
    for (int i = 0; i < n; i++) {
        cin >> arr[i];
        dp[i][i] = arr[i];
    }

    for (int len = 1; len < n; len++){
        for (int i = 0; i < n-len; i++){
            dp[i][i+len] =
max(arr[i]-dp[i+1][i+len], arr[i+len]-dp[i][i+len-1]);
        }
    }

    cout << dp[0][n-1];

    return 0;
}
```

HW:

[DMOJC '18 Contest 3 P2 - Bob and French Class - DMOJ: Modern Online Judge](#)

[A Game - DMOJ: Modern Online Judge](#)

[IOI '96 P1 - A Game - DMOJ: Modern Online Judge](#) (Very similar to above, for the point farmers)

[COCI '21 Contest 1 #2 Kamenčiči - DMOJ: Modern Online Judge](#) (may want to use the data struct pair)

Bonus: [Longest Common Subsequence - DMOJ: Modern Online Judge](#)