

# Assignment 2

## Q1 Perform Integer Arithmetic :

$$a] 15 + (-30)$$

$$(75)_{10} = (01001011)_2$$

$$(30)_{10} = (00011110)_2$$

ones complement = (11100001)<sub>2</sub>

$$\begin{aligned} \text{twos complement} &= (11100001 + 1)_2 \\ &= (11100010)_2 \end{aligned}$$

$$\therefore (-30)_{10} = (11100010)_2$$

$$\begin{array}{r} \therefore 75 + (-30) = 01001011 + 11100010 \\ = 01001011 \\ + 11100010 \\ \hline 00101101 \end{array}$$

$$\therefore (00101101)_2 = (45)_{10}$$

Q2 Perform the Unsigned Multiplication:

Multiplicand = 14

Multiplicand = 12

$$(14)_{10} = (1110)_2$$

$$(12)_{10} = (1100)_2$$

$$14 \times 12 = (168)_2$$

Multiplicand =  $(1110)_2$

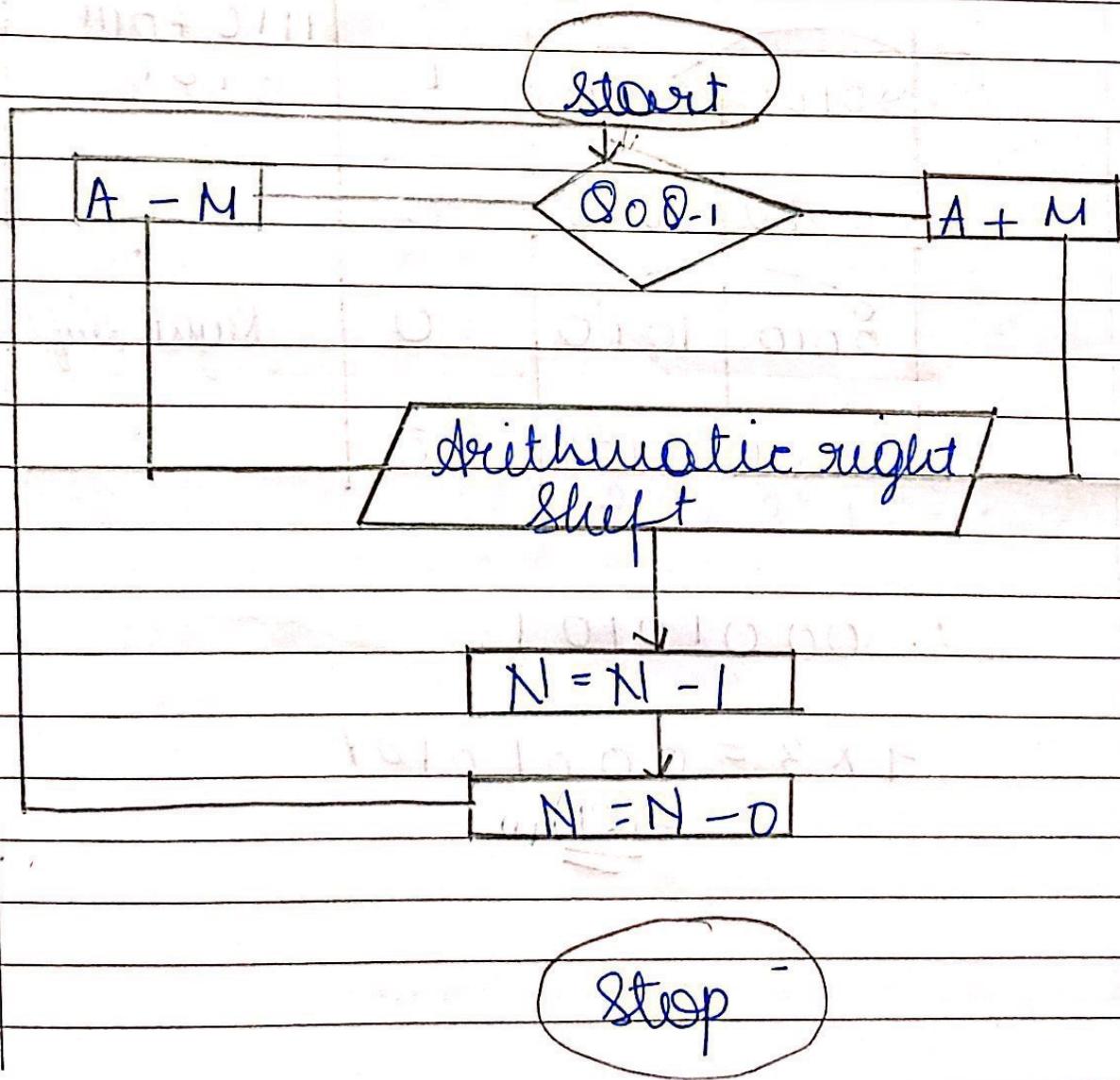
Multiplier =  $(1100)_2$

$$\begin{array}{r}
 1110 \\
 1100 \\
 \hline
 0000000 \\
 0000000 - \\
 01110 - - \\
 1110 - - - \\
 \hline
 10101000
 \end{array}$$

Multiplicand  
 = Multiplier

$$\therefore (1110)_2 \times (1100)_2 = (1010101000)_2$$

$$\therefore (10101000)_2 = (168)_2$$



Q3 Perform the signed Multiplication using  
butter's algorithm of  
multiplicand = 14  
multiplier = -12

$$M = (14)_{10} = (01110)_2$$

$$Q_0 = (-12)_{10} = (10100)_2$$

$$\begin{aligned} -M &= (-14)_{10} = 1\text{'s complement} = (10001)_2 \\ &\quad 2\text{'s} \\ &\quad = (10010)_2 \end{aligned}$$

$Q_0 Q_{-1}$	Action
00	$A \cdot S \cdot R$
11	$A \cdot S R$
01	$AC + M$
10	$AC + (-M)$

n	Ac	$Q_0$	$Q_{-1}$	M	Action
(5)	00000	10100	0	01110	i] $Q_0 Q_{-1} = 00$ $\therefore A \cdot S \cdot R$
	00000	01010	0	01110	
(4)	00000	00101	0	01110	i] $Q_0 Q_{-1} = 00$ $\therefore A \cdot S \cdot R$
(3)	10010	00101	0	01110	i] $Q_0 Q_{-1} = 10$ $Ac = Ac + (-M)$
					$\begin{array}{r} 00000 \\ + 10010 \\ \hline 10010 \end{array}$
	11001	00010	1	01110	ii] $A \cdot S R$

n	Ac	$Q_0$	$Q_{-1}$	M	Action
(2)	00111	00010	1	01110	$Q_0 \cdot Q_{-1} = 01$ $Ac = Ac + M$ 11001 $+ 01110$ <u>100111</u>
	00011	10001	0	01110	ii) J.A.S.R
(1)	10101	10001	0	01110	$Q_0 \cdot Q_{-1} = 10$ $Ac = Ac + (-M)$ 00011 $+ 10010$ <u>10101</u> ii) J.A.S.R

$$Ac Q_0 = (11010110000)_2$$

As M.S.B is 1 hence product is (-ve)  
so we take 2's complement

$$1's \text{ complement} = 0010100111$$

$$2's \text{ complement} = 0010100111$$

$$\begin{array}{r} + \\ \hline 0010101000 \end{array}$$

$$\therefore (-14)_{10} \times (-12)_{10} = (-168)_{10}$$

Q4 Perform division of the following no. using Restoring algorithm

$$\text{Dividend} = (10001)_2 = 8$$

$$\text{Divisor} = (0011)_2 = M$$

To handle borrow

$$M \leftarrow (00011)_2$$

$$-M = 1\text{'s complement} = 11100$$

$$2\text{'s complement} = 11100$$

$$\begin{array}{r}
 11100 \\
 + 1 \\
 \hline
 11101
 \end{array}$$

n	C	Ac	Q	M	Action
(1)	0	0000010000	00011	00011	i] Initialization
	0	000010000	00011	00011	ii] Left shift
	1	111000000	00011	00011	iii] C, Ac + (-M) 00001 + 11101 11110
					iv] Restoring C, AC + M 11110 + 00011 00001
	0	000100000	00011	00011	
(2)	0	001000000	00011	00011	i] Left shift
	1	111100000	00011	00011	ii] C, AC + (-M) 00010 + 11101 11111
	0	001000000	00011	00011	iii] Restoring 11111 + 00011 000010

n	C	Ac	Q	M	Operations
(2)	0	0100	0001	00011	i] Shift left
	0	0001	0001	00011	ii] C, AC+M
				00100	
				+11101	
				00001	
(1)	0	0010	001	00011	i] Shift Left
	1	1111	0010	00011	ii] C, AC+C-M
				00010	
				+11101	
				11111	
	0	0010	0010	00011	iii] Restoring
				11111	
				+00011	
				00010	

$Q \leftarrow \text{Quotient} = (0010)_2 = (2)_{10}$

$R \leftarrow \text{Remainder} = (0010)_2 = (2)_{10}$

start

$n \leftarrow$  no. of bits  
 $M \leftarrow$  Divisor  
 $A \leftarrow 0$   
 $Q \leftarrow$  dividend

shift-left  $A \&$

$$A = A - M$$

$= 0$        $A[n]?$        $= 1$

$Q[0] \leftarrow I$

$Q[0] \leftarrow 0$

Restore A

$$n = n - 1$$

NO

$n=0$

yes

Quotient is  $Q$  & remainder is  $A$

stop

Q5 Perform division of the following no. using non restoring algorithm.

$$\text{Dividend} = 1010 \rightarrow Q$$

$$\text{Divisor} = 00011 \rightarrow M$$

$$-M = 11101$$

$$R = \underline{\underline{0001}} = (1)_{10}, Q = (0011)_2 = (3)_{10}$$

n	C	Ac	Q	M	Action
(1)	0	0000	1010	00011	i] Initialization
	0	0001	010□	00011	ii] Shift left iii] C, Ac + C - M
	1	1110	0100	00011	$\begin{array}{r} 00001 \\ +11101 \\ \hline 11110 \end{array}$

(2)	1	1100	100□	00011	i] Shift left
	1	1111	1000	00011	ii] C, Ac + M $\begin{array}{r} 011100 \\ +00011 \\ \hline 11111 \end{array}$

(3)	1	10111	000□	00011	i] Shift left
	0	0010	0001	00011	ii] C, Ac + M $\begin{array}{r} 11111 \\ +00011 \\ \hline 000100 \end{array}$

(4)	0	0100	001 □	00011	i] Shift left
	0	0001	0011	00011	ii] C, Ac + M - M $\begin{array}{r} 00100 \\ +11101 \\ \hline 00001 \end{array}$

## Flowchart :-

Start

$n \leftarrow$  no. of bits

$M \leftarrow$  Divisor

$A \leftarrow 0$

$Q \leftarrow$  Dividend

= 0

= 1

sign  
bit of  
A?

Shift left A Q

$$A = A - M$$

Shift left A Q

$$A = A + M$$

sign  
of bit of  
A

Q<sub>0</sub> ← 1

Q<sub>0</sub> ← 0

NO

$$n = n - 1$$

yes

sign  
bit of A

= 0

= 1

$$A = A + M$$

Quotient in Q &  
Remainder in A

Stop