## Phylogenetic Inference: Building Trees

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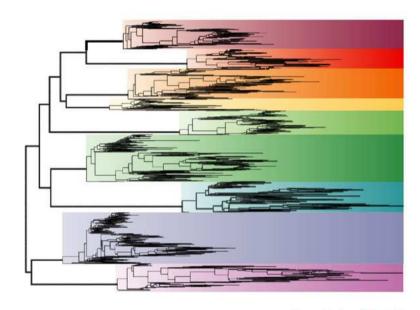
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### **Intra-Host Viral Evolution**



Nature Reviews | Genetics

1195 *env* sequences from 9 HIV+ patients [taken from Rambaut et al. (2004)]

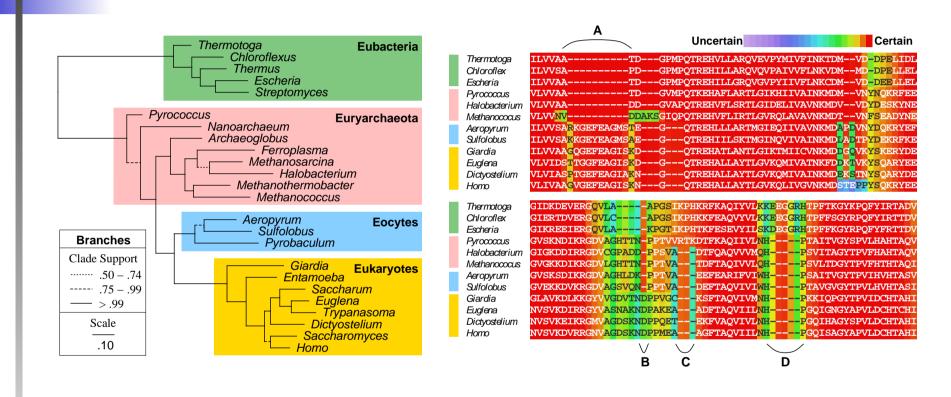
Retroviruses (and HBV) exist as a quasi-species within infected patients:

 Shared substitutions may be insufficient to resolve intra-host phylogenies

Improve resolution using joint model:

- Indel rates ≥ substitution rates
- Opportunity to detect intra-host recombination

# Reconstructing the Tree of Life: Are Humans Just Big Slime Molds?

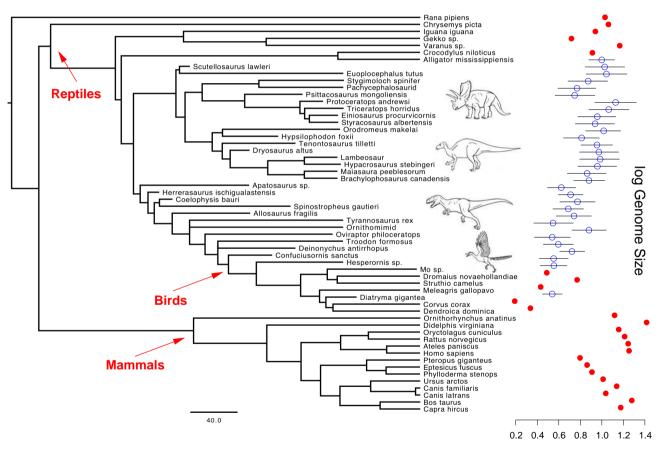


Contentious issue among paleobiologists: Do Archaea
 (Euryarchaeota/Eocytes) form one or two domains? Weekly World
 News calls humans slime molds.

**MAP Tree** 

**Partial Au Plot** 

## The Chicken or the (Small) Genome: Which Came First?



**Evolutionary History and Genome Sizes of Reptiles, Dinosaurs, Birds and Mammals** 

Issue: Bird genomes are markedly smaller than those from other vertebrates.

### Question: Did small

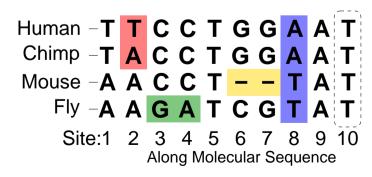
genomes
precede
flight or
co-evolve?

### **Maximum Parsimony (MP)**

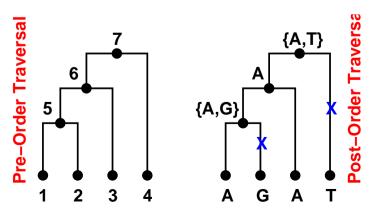
Most often used  $\neq$  "best", not even statistically consistent, but fast, fast, fast ... if you know the tree

**Key**: Find tree with minimal # of "suspected" substitutions (internal states are not observed, 0/1 model process)

- Counting minimum # of substitutions is easy
- Enumerating (searching through) all possible trees is hard



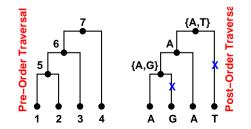
Sites are independent



### **Maximum Parsimony (MP)**

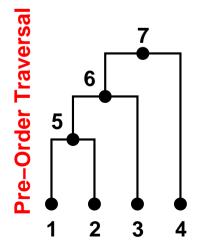
#### A little history:

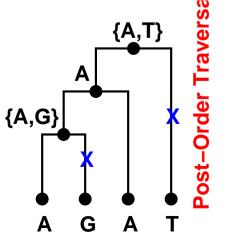
- Anthony Edwards/Luca Cavalli-Sforza (1963,1964)
  - Both students of R.A Fisher
  - Introduced both parsimony and likelihood methods (for continuous quantities, e.g. gene frequencies) in one paper
- Camin and Sokal (1965) provide first program for molecular sequences
- Fitch and Margoliash (1967) provide efficient algorithm



### **Maximum Parsimony Algorithm**

```
procedure Fitch and Margoliash (1967) Algorithm
  cost C \leftarrow 0 {Initialization}
  pointer k \leftarrow 2N - 1 {at the root node}
  To obtain the set R_k of possible states at node k {Recursion}
  if k is leaf then
     R_k \leftarrow observed character for taxon k
  else
     Compute R_i, R_j for daughters i, j of k
     if R_i \cap R_j \neq 0 then
       R_k \leftarrow R_i \cap R_i
     else
       R_k \leftarrow R_i \cup R_j
       C \leftarrow C + 1
     end if
  end if
  minimum cost is C {Termination}
```

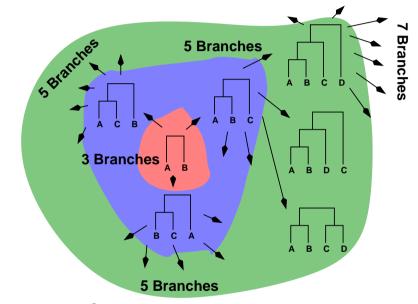




### **Searching for the MP Tree**

#### Complexity

- Find MP score is
   NP-complete
- Find MP tree is NP-hard



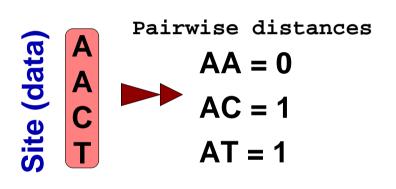
Recall that # of N-taxon rooted trees is  $3 \times 5 \times \cdots \times 2N-3$ 

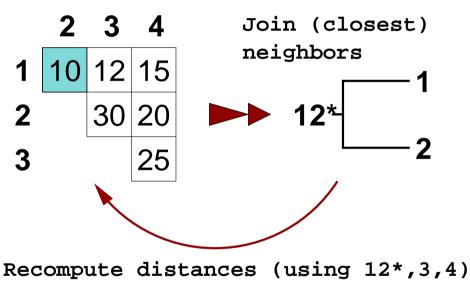
Attack exponential-order space Branch-and-Bound:

- Monotonic order: min  $PS_2 \leq min PS_3 \leq ...$
- Bound if min  $PS_k > best n$ -taxon PS found so far.

## Neighbor-Joining (Saitou and Nei, 1987)

Computational algorithm: alignment → single tree



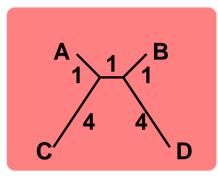


- Advantages: very fast, great for 1000s of sequences
- Disadvantages: no site-to-site rate variation, no natural ways to compare trees/measure data support

### **Neighbor-Joining**

Caveat: Pairs i, j with min  $d_{ij}$  are not necessarily nearest neighbors.

E.g., 
$$d_{\rm AB} = 3 < d_{\rm AC} = 5$$



Solution: Subtract off the average distances to all other leaves via

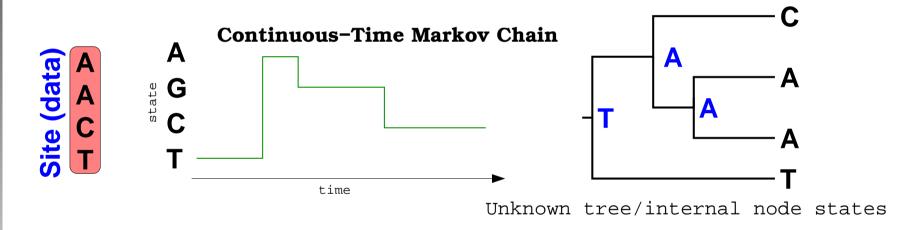
$$D_{ij} = d_{ij} - (r_i + r_j), \quad r_i = \frac{1}{|L| - 2} \sum_{k \in L} d_{ik},$$

where L is the current set of leaves. Proof in Studier and Keppler (1988).

Computational:  $O(N^3)$ 

## Likelihood-based Methods (Felsenstein, 1973)

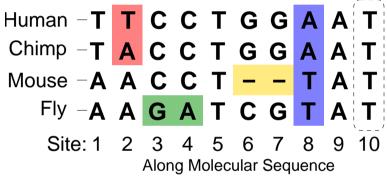
Statistical technique: assumes an unknown tree and a stochastic model for character change along the tree

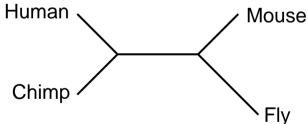


- Advantages: site-to-site rate/tree variation is easy, can formulate probability statements
- Disadvantages: must "search" tree-space → slow

## Traditional Phylogenetic Reconstruction

#### **Reconstruction Example**





- Substitution: single residue replaces another
- Insertion/deletion: residues are inserted or deleted

#### **Statistical Model**

Assume: Homologous sites are iid and site patterns (e.g. dotted box)

$$XY \dots Z \sim \mathsf{Multinomial}(p_{XY \dots Z})$$

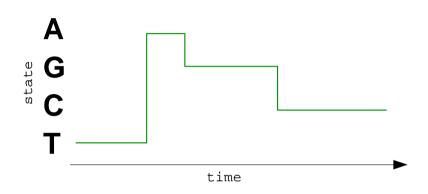
where  $p_{XY...Z}$  is determined by an unknown tree  $\tau$ , branch lengths  $t \in \mathbf{T}$  and continuous-time Markov chain model (for residue substitution) given by infinitesimal rate matrix  $\mathbf{Q}$ 

$$P(X \to Y \text{ in time } t) = e^{t\mathbf{Q}}$$

## $CTMC(Q) = \epsilon \sim Normal(\mu, \sigma^2)$ of Phylogenetics

Continuous in elapsed time t, discrete in starting/ending state!

Memory-less process in which the probability that state breplaces state a during  $(t, t + s) = sq_{ab} + o(s)$ 



• Infinitesimal generator matrix Q has off-diagonal entries  $q_{ab}$  and row sums = 0

Think: Exponential waiting time with rate  $R_a = \sum_b q_{ab}$  until chain leaves a. Then the new state b is independently chosen with probabilities  $q_{ab}/R_a$ 

## From Infinitesimal to Finite Time

Let  $p_{ab}(t)$  = the finite-time probability of the chain moving from state a at time 0 to state b at time t, then matrix  $\mathbf{P}(t) = \{p_{ab}(t)\}$  satisfies

$$\frac{d}{dt} P(t) = P(t)Q$$
 where  $P(0) = I$ 

with solution

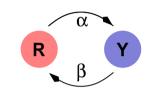
$$P(t) = e^{tQ} = I + tQ + \frac{1}{2}(tQ)^2 + \dots = \sum_{k=0}^{\infty} \frac{1}{k!}(tQ)^k$$

as

$$\frac{d}{dt}e^{tQ} = Qe^{tQ} = e^{tQ}Q$$
 for  $t$  real

### **Example: Two-State Model**

Consider purines (R) → pyrimidines (Y). Kolmogorov forward equation:



$$p_{\mathrm{RY}}(t+s) = p_{\mathrm{RR}}(t)\alpha s + p_{\mathrm{RY}}(t)(1-\beta s) + o(s)$$

yielding

$$\frac{\mathrm{d}}{\mathrm{d}t}p_{\mathrm{RY}}(t) = \alpha p_{\mathrm{RR}}(t) - \beta p_{\mathrm{RY}}(t)$$

$$\mathbf{Q} = \left( \begin{array}{cc} -\alpha & \alpha \\ \beta & -\beta \end{array} \right)$$

with eigenvalues 0 and  $-(\alpha + \beta)$ 

Solutions of  $P(t) = e^{tQ}$  have the form

$$c + de^{-(\alpha+\beta)t}$$

## Standard CTMCs for Phylogenetics

• Jukes and Cantor (JC69),  $\pi_a = \frac{1}{4}, \kappa_1 = \kappa_2 = 1$ 

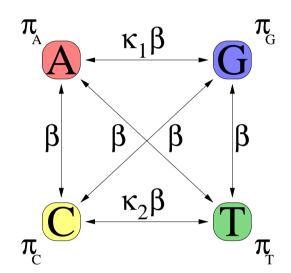
• Kimura (K80), 
$$\pi_a = \frac{1}{4}, \kappa_1 = \kappa_2$$

- Hasegawa, Kishino and Yano (HKY85),  $\kappa_1 = \kappa_2$  (most common)
- Tamura and Nei (TN93), right
- General Time Reversible (GTR)

Note identifiability concern in  $e^{tQ}$ . Common solution is to fix 1 d.f. such that

$$\sum_{a} q_{aa} \pi_a = -1$$

Scaling:  $t = 1 \Rightarrow 1$  expected substitution per site

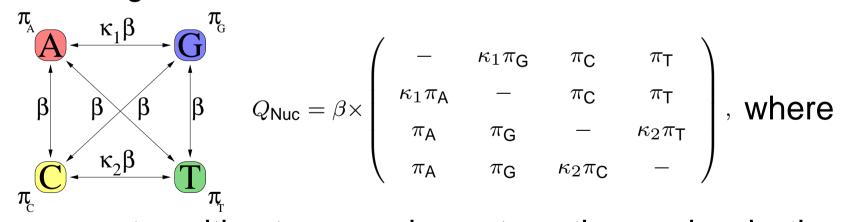


## **Explicit Parameterization of TN93**

Nucleotides mutate according to a Markovian process

$$\Pr(X \to Y \text{ in time } t) = e^{tQ_{\mathsf{Nuc}}}$$

where  $Q_{\text{Nuc}}$  is a 4x4 infinitesimal rate matrix and t is a branch length.

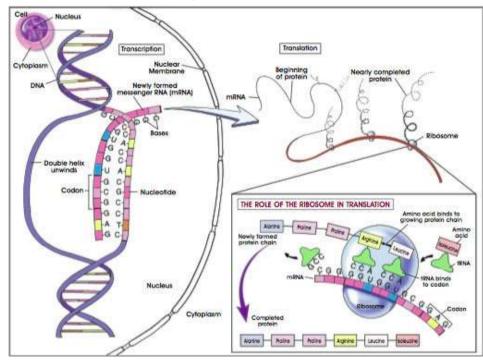


 $\kappa_1, \kappa_2$  are transition:transversion rate ratios and  $\pi$  is the stationary distribution of {A,G,C,T}.  $\beta$  controls the overall rate and can vary from site-to-site.

### Site-to-Site Rate Variation

Variation occurs quite naturally and is also an important inference

- short range: codon phase (slow-slow-fast)
- long range:
   enzymatic active
   sites, protein folding,
   immunological
   pressures/selection



Assume: infinitesimal rates for site k are  $r_k \times t \times q_{ab}$ . Various priors on  $r_k$  with  $\mathsf{E}(r_k) = 1$ . Implicitly Bayesian

• Yang (1994) - discretized Gamma distribution

### **General Time Reversible CTMC**

Let

$$oldsymbol{Q} = oldsymbol{R} oldsymbol{D}_{\pi}$$

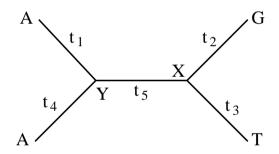
where R is symmetric and  $D_{\pi}$  is a diagonal matrix composed of the stationary distribution  $\pi$ .

- Detailed balance  $\Leftrightarrow \pi_a q_{ab} = \pi_b q_{ba}$ . Balance + irreducibility  $\Leftrightarrow$  reversible
- ullet Note  $oldsymbol{Q}$  is similar to  $oldsymbol{R}$ , as  $oldsymbol{D}^{1/2}oldsymbol{Q}oldsymbol{D}^{-1/2}=oldsymbol{R}$
- Hence, Q must have real eigenvalues and real eigenvectors

The properties speed up computation of the finite-time transition matrix  $\mathbf{P}(t) = e^{t\mathbf{Q}}$ 

## Calculating the Probability of a Single Site Pattern $Y_i$

Given the tree and unobserved internal node states, the probability is the product of the finite time mutation probabilities over all branches:



$$L(Y_i) \propto p_{\mathsf{AAGT}} = \sum_X \sum_Y \Pr(Y \to \mathsf{A}, t_1) \Pr(X \to \mathsf{G}, t_2) *$$

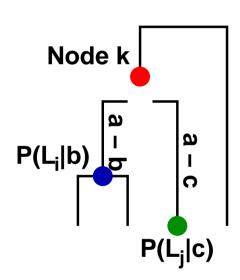
$$\Pr(X \to \mathsf{T}, t_3) \Pr(Y \to \mathsf{A}, t_4) \Pr(X \to Y, t_5) \pi_X \tag{1}$$

• Number of sumants grow rapidly in  $N \to \text{sum-product/peeling}$  algorithm to distribute sums across the product

## **Pruning Algorithm Felsenstein (1981)**

Let  $P(L_k|a)$  = likelihood of leaves below node k given k is in state a. Then, recursively compute  $P(L_k|a)$  given  $P(L_i|b)$  and  $P(L_j|c)$  for daughters i,j of k:

```
Set pointer k \leftarrow 2N-1 {the root, initialization} Compute P(L_k|a) \ \forall \ a as follows: {recursion} if k is a leaf node then if a is observed then P(L_k|a) = 1 else P(L_k|a) = 0 end if
```



else

Compute  $P(L_i|a)$  and  $P(L_j|a) \forall a$  for daughters i, j of k {post-order traversal}

$$P(L_k|a) = \sum_b \sum_c \Pr(a \to b, t_i) P(L_i|b) \times \Pr(a \to c, t_j) P(L_j|c)$$

end if

$$L(Y_i) \leftarrow \sum_a P(L_{2n-1}|a)\pi_a$$
 {termination}

### **ML Tree or MAP Tree?**

#### Reporting uncertainty on tree estimates:

- The Bootstrap
  - Most common
  - Assumes evolutionary events are reproducible. "If I went back out to the field and recollected exchangeable data . . ."
- Bayesian inference
  - Returns the probability of a tree given the observed data and model
  - Requires MCMC (e.g., MrBayes or BEAST)
  - Advantages
    - Does not rely on asymptotics (hypothesis testing)
    - \* Naturally incorporates uncertainty in all parameters (including discrete quantities: trees, site-classifications, etc.)
    - \* Arguably faster algorithms
  - Disadvantages
    - \* Must specify (justifiable) prior distributions