

Chapter 4: Intermediate SQL

Database System Concepts, 6th Ed.

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Chapter 4: Intermediate SQL

- Join Expressions
- Views
- Transactions
- Integrity Constraints
- SQL Data Types and Schemas
- Authorization



Joined Relations

- Join operations take two relations and return as a result another relation.
- A join operation is a Cartesian product which requires that tuples in the two relations match (under some condition).
 It also specifies the attributes that are present in the result of the join
- The join operations are typically used as subquery expressions in the from clause



Join operations – Example

Relation course

course_id	title	dept_name	credits
BIO-301	Genetics	Biology	4
CS-190	Game Design	Comp. Sci.	4
CS-315	Robotics	Comp. Sci.	3

Relation prereq

course_id	prereg_id
BIO-301	BIO-101
CS-190	CS-101
CS-347	CS-101

Observe that

prereq information is missing for CS-315 and course information is missing for CS-437



Outer Join

- An extension of the join operation that avoids loss of information.
- Computes the join and then adds tuples form one relation that does not match tuples in the other relation to the result of the join.
- Uses null values.



Left Outer Join

• course natural left outer join prereq

course_id	title	dept_name	credits	prereq_id
BIO-301	Genetics	Biology	4	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-101
CS-315	Robotics	Comp. Sci.	3	null



Right Outer Join

• course natural right outer join prereq

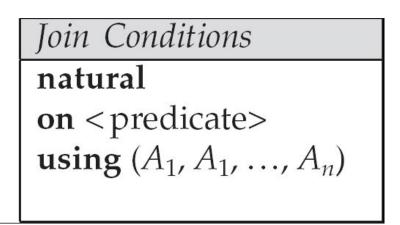
course_id	title	dept_name	credits	prereg_id
BIO-301	Genetics	Biology	4	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-101
CS-347	null	null	null	CS-101



Joined Relations

- Join operations take two relations and return as a result another relation.
- These additional operations are typically used as subquery expressions in the **from** clause
- Join condition defines which tuples in the two relations match, and what attributes are present in the result of the join.
- Join type defines how tuples in each relation that do not match any tuple in the other relation (based on the join condition) are treated.

Join types
inner join
left outer join
right outer join
full outer join





Full Outer Join

• course natural full outer join prereq

course_id	title	dept_name	credits	prereq_id
BIO-301	Genetics	Biology	4	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-101
CS-315	Robotics	Comp. Sci.	3	null
CS-347	null	null	null	CS-101



Joined Relations - Examples

 course inner join prereq on course.course_id = prereq.course_id

course_id	title	dept_name	credits	prereq_id	course_id
BIO-301	Genetics	Biology	4	BIO-101	BIO-301
CS-190	Game Design	Comp. Sci.	4	CS-101	CS-190

- What is the difference between the above, and a natural join?
- course left outer join prereq on course.course_id = prereq.course_id

course_id	title	dept_name	credits	prereq_id	course_id
BIO-301		Biology	4	BIO-101	BIO-301
CS-190	Game Design	Comp. Sci.	4	CS-101	CS-190
CS-315	Robotics	Comp. Sci.	3	null	null



Joined Relations – Examples

course natural right outer join prereq

course_id	title	dept_name	credits	prereg_id
BIO-301	Genetics	Biology	4	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-101
CS-347	null	null	null	CS-101

course full outer join prereq using (course_id)

course_id	title	dept_name	credits	prereg_id
BIO-301	Genetics	Biology	4	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-101
CS-315	Robotics	Comp. Sci.	3	null
CS-347	null	null	null	CS-101



Views

- In some cases, it is not desirable for all users to see the entire logical model (that is, all the actual relations stored in the database.)
- Consider a person who needs to know an instructors name and department, but not the salary. This person should see a relation described, in SQL, by

select *ID*, *name*, *dept_name* **from** *instructor*

- A view provides a mechanism to hide certain data from the view of certain users.
- Any relation that is not of the conceptual model but is made visible to a user as a "virtual relation" is called a view.



View Definition

 A view is defined using the create view statement which has the form

create view *v* **as** < query expression >

where <query expression> is any legal SQL expression. The view name is represented by *v*.

- Once a view is defined, the view name can be used to refer to the virtual relation that the view generates.
- View definition is not the same as creating a new relation by evaluating the query expression
 - Rather, a view definition causes the saving of an expression;
 the expression is substituted into queries using the view.



Example Views

- A view of instructors without their salary create view faculty as select ID, name, dept_name from instructor
- Find all instructors in the Biology department select name
 from faculty
 where dept_name = 'Biology'
- Create a view of department salary totals
 create view departments_total_salary(dept_name, total_salary) as
 select dept_name, sum (salary)
 from instructor
 group by dept_name;



Views Defined Using Other Views

- create view physics_fall_2009 as
 select course.course_id, sec_id, building, room_number
 from course, section
 where course.course_id = section.course_id
 and course.dept_name = 'Physics'
 and section.semester = 'Fall'
 and section.year = '2009';
- create view physics_fall_2009_watson as select course_id, room_number from physics_fall_2009 where building= 'Watson';



View Expansion

Expand use of a view in a query/another view

```
create view physics_fall_2009_watson as
(select course_id, room_number
from (select course.course_id, building, room_number
    from course, section
    where course.course_id = section.course_id
        and course.dept_name = 'Physics'
        and section.semester = 'Fall'
        and section.year = '2009')
where building= 'Watson';
```



Views Defined Using Other Views

- One view may be used in the expression defining another view
- A view relation v_1 is said to depend directly on a view relation v_2 if v_2 is used in the expression defining v_1
- A view relation v_1 is said to depend on view relation v_2 if either v_1 depends directly to v_2 or there is a path of dependencies from v_1 to v_2
- A view relation v is said to be recursive if it depends on itself.



View Expansion

- A way to define the meaning of views defined in terms of other views.
- Let view v₁ be defined by an expression e₁ that may itself contain uses of view relations.
- View expansion of an expression repeats the following replacement step:

repeat

Find any view relation v_i in e_1 Replace the view relation v_i by the expression defining v_i until no more view relations are present in e_1

 As long as the view definitions are not recursive, this loop will terminate



Update of a View

 Add a new tuple to faculty view which we defined earlier insert into faculty values ('30765', 'Green', 'Music');

This insertion must be represented by the insertion of the tuple

('30765', 'Green', 'Music', null)

into the *instructor* relation



Some Updates cannot be Translated Uniquely

- create view instructor_info as
 select ID, name, building
 from instructor, department
 where instructor.dept_name= department.dept_name;
- insert into instructor_info values ('69987', 'White', 'Taylor');
 - 4 which department, if multiple departments in Taylor?
 - 4 what if no department is in Taylor?
- Most SQL implementations allow updates only on simple views
 - The from clause has only one database relation.
 - The select clause contains only attribute names of the relation, and does not have any expressions, aggregates, or distinct specification.
 - Any attribute not listed in the select clause can be set to null
 - The query does not have a group by or having clause.



And Some Not at All

- create view history_instructors as select * from instructor where dept_name= 'History';
- What happens if we insert ('25566', 'Brown', 'Biology', 100000) into history_instructors?



Materialized Views

- Materializing a view: create a physical table containing all the tuples in the result of the query defining the view
- If relations used in the query are updated, the materialized view result becomes out of date
 - Need to maintain the view, by updating the view whenever the underlying relations are updated.



Transactions

- Unit of work
- Atomic transaction
 - either fully executed or rolled back as if it never occurred
- Isolation from concurrent transactions
- Transactions begin implicitly
 - Ended by commit work or rollback work
- But default on most databases: each SQL statement commits automatically
 - Can turn off auto commit for a session (e.g. using API)
 - In SQL:1999, can use: begin atomic end
 - 4 Not supported on most databases