

## 2 Nash equilibrium (NE)

### 2.1 Examples

Prisoner's dilemma.

		P2	
		Share	Steal
P1	Share	5, 5	0, 10
	Steal	10, 0	0.1, 0.1

Let  $s^* = (\text{steal}, \text{steal})$

P1:  $u_1(s^*) = 0.1$

$u_1(\text{share}, s_{-1}^*) = 0 < u_1(s^*)$ .

↑ P1 switches to share  
↑ P2 remains steal

P2:  $u_2(s^*) = 0.1$   $u_2(\text{share}, s_{-2}^*) = 0 < u_2(s^*)$

So  $s^*$  is a Nash equilibrium.

Bach or Stravinsky.

		P2	
		Bach	Stravinsky
P1	Bach	2, 1	0, 0
	Stravinsky	0, 0	1, 2

Both (B,B) and (S,S) are NEs.

Rock paper scissors.

		P2		
		R	P	S
P1	R	0, 0	-1, 1	1, -1
	P	1, -1	0, 0	-1, 1
	S	-1, 1	1, -1	0, 0

For any strategy profile, one player with utility 0 or -1 can get higher utility by switching to a winning gesture.

No NE ... (sort of a lie, to be resolved later).

**2/3 average game.** 3 players. Each player simultaneously pick from  $\{1, \dots, 10\}$ . A \$1 prize is split among all players closest to  $2/3$  of the average of the 3 numbers, other players get \$0.

$S = (6, 4, 1)$  average  $\frac{11}{3}$ .  $\frac{2}{3}$  average  $\frac{2}{3} \cdot \frac{11}{3} = 2.4...$

$u_3(s) = 1$ ,  $u_1(s) = u_2(s) = 0$ .

P1:  $u_1(1, s_{-1}) = 0.5 > u_1(s)$ . Not a NE.

$(1, 1, 1)$  is the only NE (needs proof)

## 2.2 Best response function

Definition: best response function.

Player  $i$ 's best response function (BRF) for  $s_{-i} \in S_{-i}$  is

$$B_i(s_{-i}) = \{s_i' \in S_i : \underbrace{u_i(s_i', s_{-i})}_{\text{utility of a best response}} \geq \underbrace{u_i(s_i, s_{-i})}_{\text{utility of all possible responses}} \forall s_i \in S_i\}.$$

Examples. Prisoner's dilemma.

$$B_1(\text{share}) = \{\text{steal}\} \quad B_1(\text{steal}) = \{\text{steal}\}$$

$\uparrow$   $\uparrow$   
 $P_2$ 's strat  $u_1(\text{steal}, \text{share}) > u_1(\text{share}, \text{share})$

2/3 average game.

$$S_{-1} = (5, 5) \quad (P_2, P_3 \text{ play } 5). \quad u_1(x, 5, 5) = \begin{cases} 0 & x \geq 6 \\ 1/3 & x = 5 \\ 1 & x \leq 4 \end{cases}$$

$$B_1(5, 5) = \{1, 2, 3, 4\}.$$

Lemma 1.

A strategy profile  $s^* \in S$  is a Nash equilibrium if and only if for each player  $i$ ,  $s_i^* \in B_i(s_{-i}^*)$ .

[No incentive for any player to switch, since each player played their best response already.]

Examples.

		P2			
		Share	Steal		
P1	Share	5, 5	0, 10	○ best response for P1 given P2's choices	□ " " P2 " P1's "
	Steal	10, 0	0.1, 0.1		

$\uparrow$  BR given P2 share  
 $\rightarrow$  best responses to each other, NE.

		P2				
		D	E	F		
P1	A	1, 2	2, 1	1, 0		
	B	2, 1	0, 1	0, 0		
	C	0, 1	0, 0	1, 2		

NE: (B, D), (C, F).

best responses to each other.

## 2.3 Cournot's oligopoly model

Firms  $N = \{1, \dots, n\}$  are producing a single type of goods sold on the common market.

- Each firm  $i$  decides the number of units of goods  $q_i$  to produce.
- Production cost for firm  $i$  is  $C_i(q_i)$  where  $C_i$  is a given increasing function.
- Given a strategy profile  $q = (q_1, \dots, q_n)$ , a unit of the goods sell for the price of  $P(q)$ , where  $P$  is a given non-increasing function on  $\sum_i q_i$ .
- The utility of firm  $i$  is  $u_i(q) = q_i P(q) - C_i(q_i)$ .

SzidaroVszky and Yakowitz proved that a NE exists under some continuity and differentiability assumptions on  $P, C_i$ .

Special case: linear costs and prices

Assume  $C_i(q_i) = cq_i \quad \forall i \in N$ , fixed  $c$ .

$P(q) = \max\{0, \alpha - \sum_j q_j\}$  fixed  $\alpha$ .  $0 < c < \alpha$ .

Utility is  $u_i(q) = q_i P(q) - C_i(q_i) = \begin{cases} q_i(\alpha - \sum_j q_j - c) & \alpha - \sum_j q_j \geq 0 \\ -cq_i & \alpha - \sum_j q_j < 0 \end{cases}$

Find BR for firm  $i$ : If  $\alpha - \sum_{j \neq i} q_j \leq 0$ , then BR is 0 (price is 0, no need to produce).

When  $\alpha - \sum_j q_j - c > 0$ , we make a profit. Take  $q_i$  out of the sum.

$\alpha - c - q_i - \sum_{j \neq i} q_j > 0$ . The utility is  $q_i(\alpha - c - q_i - \sum_{j \neq i} q_j)$ .

Given fixed  $q_j, j \neq i$ , we want  $q_i$  that maximizes utility (BR).

This is maximized at  $q_i = (\alpha - c - \sum_{j \neq i} q_j)/2$ . (Using calculus?)

$$B_i(q_{-i}) = \begin{cases} (\alpha - c - \sum_{j \neq i} q_j)/2 & \alpha - c - \sum_{j \neq i} q_j > 0 \\ \{0\} & \text{otherwise.} \end{cases}$$