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(1)  $xs.map(id) = xs$

LHS =  $[] .map(id) = []$ . (by def)  
RHS =  $[] =$  LHS

Induction case:  $xs = (x:xs')$

Assume:  $xs'.map(id) = xs'$

Prove:  $(x:xs').map(id) = xs'$

LHS =  $(x:xs').map(id)$   
=  $(x:xs'.map(id))$   
=  $(x:xs')$

RHS =  $(x:xs') =$  LHS

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(2)  $(xs.append(ys)).map(f) = (xs.map(f)).append(ys.map(f))$

Base Case:  $xs = []$   
LHS =  $([] .append(ys)).map(f)$   
=  $ys.map(f)$   
RHS =  $([] .map(f)).append(ys.map(f))$   
=  $[] .append(ys.map(f))$   
=  $ys.map(f) =$  LHS

Induction case:  
 $xs = (x:xs')$

Assume:  $(xs'.append(ys)).map(f) =$   
 $(xs'.map(f)).append(ys.map(f))$

Prove:  $(x:xs'.append(ys)).map(f) =$   
 $(x:xs'.map(f)).append(ys.map(f))$

LHS =  $((x:xs').append(ys)).map(f)$   
=  $(x:xs'.append(ys)).map(f)$   
=  $x:(xs'.append(ys)).map(f)$

```

RHS = ((x:xs').map(f)).append(ys.map(f))
      = (x:xs'.map(f)).append(ys.map(f))
      = (x:xs').append(ys.map(f))
      = x:(xs'.append(ys)).map(f)
      = LHS

```

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(3)  $(xs.append(ys)).fold(g, a) = xs.fold(g, ys.fold(g, a))$

Base Case:  $xs = []$

LHS =  $([].append(ys)).fold(g, a) = ys.fold(g, a)$

RHS =  $[].fold(g, ys.fold(g, a)) = ys.fold(g, a)$

Induction case:  $xs = (x:xs')$

Assume:  $(xs'.append(ys)).fold(g, a) = xs'.fold(g, ys.fold(g, a))$

Prove:  $((x:xs').append(ys)).fold(g, a) = (x:xs').fold(g, ys.fold(g, a))$

```

LHS = ((x:xs').append(ys)).fold(g, a)
      = x:xs'.append(ys).fold(g, a)
      = x:xs'.fold(g, ys.fold(g, ys.fold(g, a)))

```

```

RHS = (x:xs').fold(g, ys.fold(g, a))
      = g.apply(x, xs'.fold(g, ys.fold(g, a))) = LHS

```

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(4)  $(xs.append(ys)).length() = xs.length() + ys.length()$

Base Case:  $xs = []$

LHS =  $[].append(ys).length() = ys.length()$

RHS =  $[].length() + ys.length() = ys.length() = LHS$

Induction step:  $xs = (x:xs')$

Assume:  $(xs'.append(ys)).length() = xs'.length() + ys.length()$

Prove:  $((x:xs').append(ys)).length() = (x:xs').length() + ys.length()$

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(5)

Prove:  $(xs.reverseH\ ys).length() = xs.length() + ys.length()$

${}.reverseH(ys) = ys$

$(x:xs).reverseH(ys) = xs.reverseH(x:ys)$

Base case:  $xs = []$

LHS =  ${}.reverseH(ys).length() = ys.length()$

RHS =  ${}.length() + ys.length() = ys.length() = \text{LHS}$

Inductive case:  $xs = (x:xs')$

Assume this is true IH:

forall  $ys$ .  $xs'.reverseH(ys).length() = xs'.length() + ys.length()$

To prove:  $(x:xs').reverseH(ys).length() = (x:xs').length() + ys.length()$

LHS =  $(x:xs').reverseH(ys).length()$   
       $xs'.reverseH(x:ys).length()$   
       $xs'.length() + (x:ys).length()$   
       $xs'.length() + 1 + ys.length()$

RHS =  $(x:xs').length() + ys.length()$   
       $= 1 + xs'.length() + ys.length()$   
       $= \text{LHS}$

---

(6)

Prove:  $xs.length() = xs.reverse2().length()$

List reverse2 () { return reverseH(new EmptyL()); }

LHS =  $xs.length()$

RHS =  $xs.reverseH([]).length()$

(5) forall  $ys$ .  $xs.reverseH(ys).length() = xs.length() + ys.length()$

(5')  $xs.reverseH([]).length() = xs.length() +$

`[]`.length() = `xs.length()`  
Done.

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(7)  $(xs.append(ys)).reverse() =$   
 $ys.reverse().append(xs.reverse())$

Base case: `xs = []`

LHS = `[]`.append(`ys`).reverse() = `ys.reverse()`

RHS = `ys.reverse().append([].reverse())` = `ys.reverse()` =  
LHS

Induction case: `xs = (x:xs')`

Assume: `xs'.append(ys).reverse() =`  
`ys.reverse().append(xs'.reverse())`

Prove:  $(x:xs').append(ys).reverse() =$   
 $ys.reverse().append((x:xs').reverse())$

LHS = `(x:xs')`.append(`ys`).reverse()  
=

RHS = `ys.reverse().append((x:xs').reverse())`  
=

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(8)  $(xs.reverse()).reverse() = xs$

Base Case: `xs=[]`

LHS = (`[]`.reverse()).reverse()  
= `[]`.reverse()  
= `[]`

RHS = `[]` = LHS

Induction step: `xs=(x:xs')`

Assume: `(xs'.reverse()).reverse() = xs'`

Prove:  $((x:xs').reverse()).reverse() = (x:xs')$

LHS = `(x:xs')`.reverse()).reverse()

```

    = xs'.reverse().append(x:[])).reverse
    = (x:
[]).reverse().append((xs'.reverse()).reverse()).reverse()
    = [].reverse().append(x:
[])).append(xs'.reverse().reverse())
    = (x:[]).append(xs'.reverse().reverse())
    = (x:[]).append(xs') (IH)
    = (x:xs')

```

RHS = (x:xs') = LHS

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(9) xs.reverseH (ys) = (xs.reverse()).append(ys)

Base Case: xs = []

LHS = [].reverseH (ys)

= ys

RHS = ([].reverse()).append(ys)

= [].append(ys)

= ys = LHS

Induction case: xs = (x:xs')

Assume: xs'.reverseH (ys) = (xs'.reverse()).append(ys)

Prove: (x:xs').reverseH (ys) =

((x:xs').reverse()).append(ys)

LHS = (x:xs').reverseH(ys)

= xs'.reverseH(x:ys)

= (xs'.reverse()).append(x:ys) (IH)

= (xs'.reverse()).append(x:ys).reverse())

= (xs'.reverse ()).append(x:ys).reverse().reverse())

=

((x:ys).reverse()).append((xs'.reverse()).reverse()).reverse()

= (((x:ys).reverse()).append(xs').reverse())

= (ys.reverse().append(x:[]).append(xs').reverse())

= (ys.reverse()).append(x:

[]).append((xs').append(xs').reverse())

= (ys.reverse()).append(x:xs').reverse()

= ((x:xs').reverse()).append(ys.reverse().reverse())

= (x:xs').reverse()).append(ys)

$$\text{RHS} = ((x:xs').\text{reverse}()).\text{append}(ys) = \text{LHS}$$


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$$(10) \quad \overline{(xs.\text{reverseH}(ys)).\text{reverseH}(zs) = ys.\text{reverseH}(xs.\text{append}(zs))}$$

Base Case:  $xs = []$

$$\begin{aligned} \text{LHS} &= ([].\text{reverseH}(ys)).\text{reverseH}(zs) = ys.\text{reverseH}(zs) \\ \text{RHS} &= ys.\text{reverseH}([].\text{append}(zs)) = ys.\text{reverseH}(zs) = \text{LHS} \end{aligned}$$

Induction step:  $xs = (x:xs')$

Assume:  $(xs'.\text{reverseH}(ys)).\text{reverseH}(zs) = ys.\text{reverseH}(xs'.\text{append}(zs))$

Prove:  $((x:xs').\text{reverseH}(ys)).\text{reverseH}(zs) = ys.\text{reverseH}((x:xs').\text{append}(zs))$

$$\begin{aligned} \text{LHS} &= ((x:xs').\text{reverseH}(ys)).\text{reverseH}(zs) \\ &= (xs'.\text{reverseH}(x:ys)).\text{reverseH}(zs) \\ &= (x:ys).\text{reverseH}(xs'.\text{append}(zs)) \quad (\text{IH}) \\ &= ys.\text{reverseH}(xs'.\text{append}(zs)) \\ \text{RHS} &= ys.\text{reverseH}(x:xs').\text{append}(zs) \\ &= ys.\text{reverseH}(xs'.\text{append}(zs)) = \text{LHS} \end{aligned}$$


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$$(11) \quad \overline{xs.\text{reverseH}(ys.\text{append}(zs)) = (xs.\text{reverseH}(ys)).\text{append}(zs)}$$

Invalid argument:  $xs.\text{reverseH}(ys.\text{append}(zs)) \neq (xs.\text{reverseH}(ys)).\text{append}(zs)$

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$$(12) \quad \overline{(xs.\text{append}(ys)).\text{reverseH}(zs) = ys.\text{reverseH}(xs.\text{reverseH}(zs))}$$

Base Case:  $xs = []$

$$\begin{aligned} \text{LHS} &= [].\text{append}(ys).\text{reverseH}(zs) = ys.\text{reverseH}(zs) \\ \text{RHS} &= ys.\text{reverseH}([].\text{reverseH}(zs)) = ys.\text{reverseH}(zs) \end{aligned}$$

Induction step:  $xs = (x:xs')$

Assume:  $(xs'.append(ys)).reverseH(zs) =$   
 $ys.reverseH(xs'.reverseH(zs))$

Prove:  $((x:xs').append(ys)).reverseH(zs) =$   
 $ys.reverseH((x:xs').reverseH(zs))$

LHS =  $(xs'.append(ys)).reverseH(zs)$   
=  $(xs'.append(ys)).reverse().append(zs)$   
=  $ys.reverse.append(xs', reverse()).append(zs)$   
=  $ys.reverse().append((xs'.reverse()).append(zs))$   
=  $ys.reverseH(xs.reverseH(zs))$

RHS =  $ys.reverseH(xs'.reverseH(zs))$

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(13)  $(xs.append(ys)).reverse2() =$   
 $ys.reverse2().append(xs.reverse2())$

Base Case:  $xs = []$

LHS =  $[].append(ys).reverse2() = ys.reverse2()$   
RHS =  $ys.reverse2().append([].reverse2()) = ys.reverse2() =$   
LHS

Induction step:  $xs = (x:xs')$

LHS =  $xs'.append(ys).reverse2()$   
=  $xs'.append(ys).reverseH([])$   
=  $ys.reverseH(xs'.reverseH([]))$   
=  $ys.reverseH(xs'.reverse2())$   
=  $ys.reverseH({}.append(xs', reverse2()))$   
=  $ys.reverse2().append(xs', reverse2())$   
RHS =  $ys.reverse().append(xs', reverse2()) =$  LHS

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(14)  $(xs.reverse2()).reverse2() = xs$

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(15)  $t.flattenH(xs) = t.flatten().append(xs)$

By induction on  $t$ :

Base case:  $t = d$

$$\text{LHS} = d.\text{flattenH}(xs) = d:xs$$

$$\begin{aligned}\text{RHS} &= d.\text{flatten}().\text{append}(xs) = (d:[]).\text{append}(xs) = d: \\ &[].\text{append}(xs) = \text{LHS}\end{aligned}$$

Inductive case:  $t = N(d, t_1, t_2)$

Lemma:

$$xs.\text{append}(ys.\text{append}(zs)) = (xs.\text{append}(ys)).\text{append}(zs)$$

$$\begin{aligned}\text{To prove: } &N(d, t_1, t_2).\text{flattenH}(xs) = \\ &N(d, t_1, t_2).\text{flatten}().\text{append}(xs)\end{aligned}$$

$$\text{IH1: } t_1.\text{flattenH}(xs) = t_1.\text{flatten}().\text{append}(xs)$$

$$\text{IH2: } t_2.\text{flattenH}(xs) = t_2.\text{flatten}().\text{append}(xs)$$

$$\begin{aligned}\text{LHS} &= N(d, t_1, t_2).\text{flattenH}(xs) \\ &= t_1.\text{flattenH}(d:t_2.\text{flattenH}(xs)) \\ &= t_1.\text{flatten}().\text{append}(d:t_2.\text{flattenH}(xs)) \\ &= t_1.\text{flatten}().\text{append}(d:t_2.\text{flatten}().\text{append}(xs))\end{aligned}$$

$$\begin{aligned}\text{RHS} &= N(d, t_1, t_2).\text{flatten}().\text{append}(xs) \\ &= (t_1.\text{flatten}().\text{append}(d:t_2.\text{flatten}().\text{append}(xs))).\text{append}(xs) \\ &= t_1.\text{flatten}().\text{append}(d:t_2.\text{flatten}().\text{append}(xs)) \\ &= \text{LHS}\end{aligned}$$

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$$(16) \quad t.\text{flatten2}() = t.\text{flatten}()$$

Base Case:  $t = d$

$$\begin{aligned}\text{LHS} &= d.\text{flatten2}() \\ &= d.\text{flattenH}([]) \\ &= d:[]\end{aligned}$$

$$\begin{aligned}\text{RHS} &= d.\text{flatten}() \\ &= d:[]\end{aligned}$$

Induction step:  $t = N(d, t_1, t_2)$

$$\text{Prove: } n(d, t_1, t_2).\text{flatten2}() = N(d, t_1, t_2).\text{flatten}()$$



```
IH1:t1.flatten2() = t1.flatten()
IH2:t2.flatten2()= t2.flatten()
```

```
LHS =
    = N(d,t1,t2).flatten2()
    = N(d,t1,t2).flattenH([])
    = N(d,t1,t2).flatten().append([])
    = N(d,t1,t2).flatten()
RHS = N(d,t1,t2).flatten.append(XS)
```

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```
(17) t.map(f1).sum() = t.nodes()
```

Base case:  $t = d$

Invalid argument:  $t.map(f1).sum() \neq t.nodes()$

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```
(18) t.nodes() = t.longestPath().length() + 1
```

Invalid argument:  $t.nodes() \neq t.longestPath().length() + 1$

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```
(19)
```

Prove: For non-empty trees  $t$ , it is the case that  
 $t.internalNodes() + 1 = t.leaves()$ .

Base case:  $t = d$

LHS =  $d.internalNodes() + 1 = 1$

RHS =  $d.leaves() = 1 = \text{LHS}$

Induction case:  $t = N(d,t1,t2)$

Where  $t1 = d, t2 = d$

To prove:  $N(d,t1,t2).internalNodes() + 1 =$   
 $N(d,t1,t2).leaves()$

LHS =  $N(d,t1,t2).internalNodes() + 1$   
      =  $d.internalNodes() + t1.internalNodes() +$   
       $t2.internalNodes() + 1$   
      =  $1 + 0 + 0 + 1 = 2$

$$\begin{aligned}
\text{RHS} &= N(d, t_1, t_2). \text{leaves}() \\
&= t_1. \text{leaves}() + t_2. \text{leaves}() \\
&= 2 = \text{LHS}
\end{aligned}$$


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(20)

Prove: A full  $m$ -ary with  $n$  nodes has  $(n - 1)/m$  internal nodes and  $((m - 1)n + 1)/m$  leaves.

$$\begin{aligned}
\text{Leaf} &= d \\
\text{Leaves} &= 1 \\
\text{check: } 1 &= ((1 - 1)1 + 1)/1 = 1 \text{ YES} \\
\text{internal nodes} &= 0 \\
\text{check: } 0 &= (1-1)/1 \text{ YES}
\end{aligned}$$

$$\begin{aligned}
&\text{Node}(d, t_1, t_2, \dots, t_m) \\
&\text{leaves} = l_1 + l_2 + \dots + l_m \\
&\text{internal nodes} = i_1 + i_2 + \dots + i_m + 1 \\
&\text{For each subtree } k, \text{ we have } i_k = (n_k - 1)/m \\
&\text{For each subtree } k, \text{ we have } n_k = ((m-1)n_k + 1)/m
\end{aligned}$$

check:

$$\begin{aligned}
\text{LHS} &= i_1 + i_2 + \dots + i_m + 1 \\
&= (n_1 - 1)/m + (n_2 - 1)/m + \dots + (n_m - 1)/m + 1 \\
&= (n_1 - 1 + n_2 - 1 + \dots + n_m - 1 + m) / m \\
&= (n_1 + n_2 + \dots + n_m) / m
\end{aligned}$$

$$\begin{aligned}
\text{RHS} &= l_1 + l_2 + \dots + l_m \\
&= ((m-1)n_1 + 1) + ((m-1)n_2 + 1) + \dots + ((m-1)n_m + 1) \\
&= (m*n_1 - n_1 + 1) + (m*n_2 - n_2 + 1) + \dots + (m*n_m - n_m + 1) \\
&= (m*n_1 - n_1 + m*n_2 - n_2 + \dots + m*n_m - n_m) / m \\
&= (n_1 + n_2 + \dots + n_m) / m \\
&= \text{LHS}
\end{aligned}$$


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(21) A full  $m$ -ary with  $i$  internal nodes has  $mi + 1$  nodes and  $(m - 1)i + 1$  leaves.

$$\text{Leaf} = d$$

Internal nodes = 0  
 nodes = 1  
 check:  $1 = m_i + 1 = 1$  YES  
 leaves = 1  
 check:  $1 = (m - 1)i + 1 = 1$  YES

Node( $d, t_1, t_2, \dots, t_m$ )  
 nodes =  $n_1 + n_2 + \dots + n_m + 1$   
 leaves =  $l_1 + l_2 + \dots + l_m$   
 For each subtree  $k$ , we have  $n_k = m_i k + 1$   
 For each subtree  $k$ , we have  $l_k = (m - 1)i k + 1$

check:  
 LHS =  $n_1 + n_2 + \dots + n_m + 1$   
       =  $(m i_1 + 1) + (m i_2 + 1) + \dots + (m i_m + 1)$   
       =  $(n_1 + n_2 + \dots + n_m) / m$   
  
 RHS =  $l_1 + l_2 + \dots + l_m$   
       =  $((m-1)i_1 + 1) + ((m-1)i_2 + 1) + \dots + ((m-1)i_m + 1)$   
       =  $(m i_1 - i_1 + m i_2 - i_2 + \dots + m i_m - i_m) + m$   
       = LHS

(22) A full  $m$ -ary with  $l$  leaves has  $(m l - 1) / (m - 1)$  nodes  
 and  $(l - 1) / (m - 1)$  internal nodes.

Leaf =  $d$   
 Internal nodes = 0  
 nodes = 1  
 check:  $1 = m_i + 1 = 1$  YES  
 leaves = 1  
 check:  $1 = (m - 1)i + 1 = 1$  YES

Node( $d, t_1, t_2, \dots, t_m$ )  
 nodes =  $n_1 + n_2 + \dots + n_m + 1$   
 leaves =  $l_1 + l_2 + \dots + l_m$   
 For each subtree  $k$ , we have  $n_k = m_i k + 1$   
 For each subtree  $k$ , we have  $l_k = (m - 1)i k + 1$

check:

$$\begin{aligned}
\text{LHS} &= n_1 + n_2 + \dots + n_m + 1 \\
&= (m_1 + 1) + (m_2 + 1) + \dots + (m_m + 1) \\
&= \\
&= (n_1 + n_2 + \dots + n_m) / m \\
\\
\text{RHS} &= l_1 + l_2 + \dots + l_m \\
&= ((m-1)i_1 + 1) + ((m-1)i_2 + 1) + \dots + ((m-1)i_m + 1) \\
&= (m_1 - i_1 + m_2 - i_2 + \dots + m_m - i_m) + m \\
&= \\
&= \text{LHS}
\end{aligned}$$

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(23)

$$l = 100$$

$$n = ?$$

$$i = ?$$

$$n = ((m-1)n + 1)/m$$

$$100 = ((4-1)n + 1)/4$$

$$100 = (3n + 1)/4$$

$$400 = 3n + 1$$

$$399 = 3n$$

$$n = 133$$

133 people have read the letter

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(24)

$$l = 100$$

$$n = 133$$

$$i = ?$$

$$i = (l-1)/(m-1)$$

$$i = (100-1)/(4-1)$$

$$i = 99/3$$

$$i = 33$$

33 people sent out the letter