

Let X_{ij} = # of truckloads to move from high point i to low point j

$$i = \{1, 2, 3\} \quad j = \{A, B\}$$

Objective

$$\min 10X_{1A} + 29X_{1B} + 15X_{2A} + 4X_{2B} + 24X_{3A} + 5X_{3B}$$

s.t.

$$X_{1A} + X_{1B}$$

$$= 3$$

$$X_{2A} + X_{2B}$$

$$= 3$$

$$X_{3A} + X_{3B} = 4$$

$$X_{1A}$$

$$+ X_{2A}$$

$$+ X_{3A}$$

$$= 7$$

$$X_{1B}$$

$$+ X_{2B}$$

$$+ X_{3B} = 3$$

$$X_{ij} \geq 0$$

The Model

- Want to distribute stuff from any set of sources to any set of sinks/destinations

- Each source has a supply

Assume: - fixed

- entire supply must be distributed

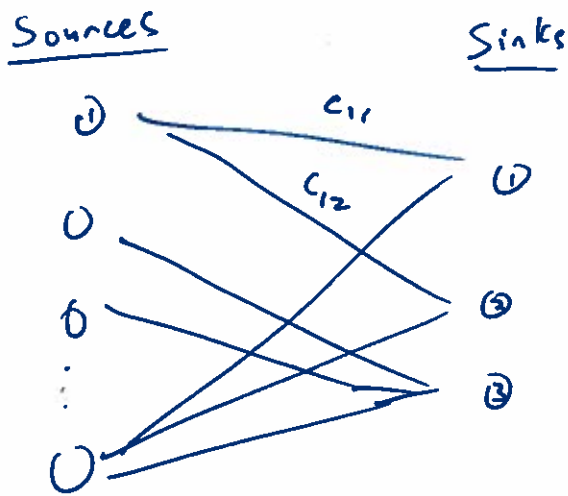
- Each sink has a demand

Assume: - fixed

- entire demand must be met

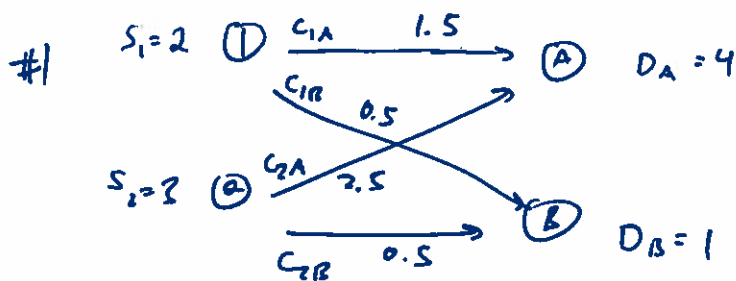
- Cost c_{ij} per unit shipped from i to j

- directly proportional to # of units shipped

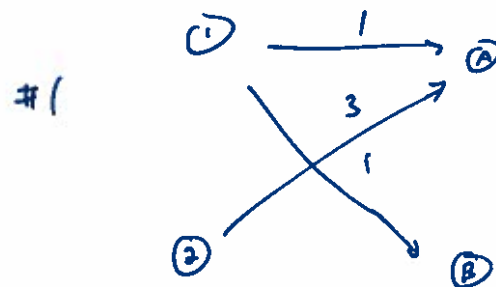
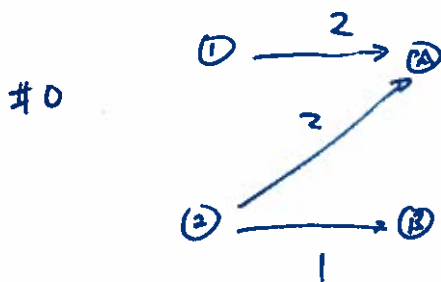


- Special property of transportation problems
- ordinary LP's may have non-integer sol'n's
 - transportation problems guarantee and integer opt. sol'n when supplies + demands are integer

Proof Sketch: Consider the following fractional flow



There are two feasible integer flows



$$\left. \begin{aligned} \text{Cost}_0 &= 2C_{1A} + 0C_{1B} + 2C_{2A} + 1C_{2B} \\ \text{Cost}_1 &= 1.5C_{1A} + 0.5C_{1B} + 2.5C_{2A} + 0.5C_{2B} \\ \text{Cost}_2 &= 1C_{1A} + 1C_{1B} + 3C_{2A} + 0C_{2B} \end{aligned} \right\}$$

Note that

$$\text{Cost}_1 = \frac{\text{Cost}_0 + \text{Cost}_2}{2}$$

So $\text{Cost}_1 \leq \max\{\text{Cost}_0, \text{Cost}_2\}$

* Either fractional flow isn't optimal (b/c #0/#2 optimal)

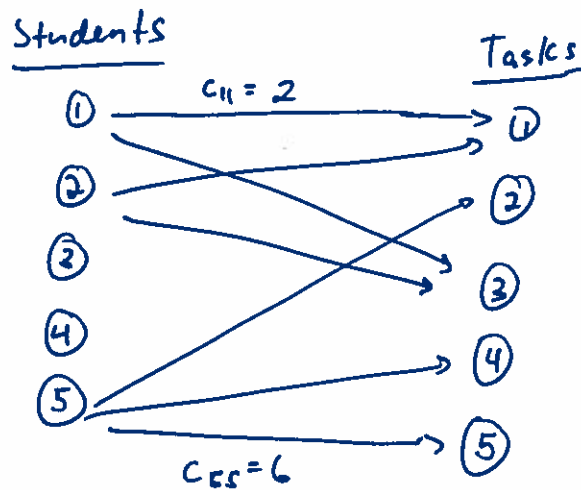
* or it is optimal, in which case $\text{Cost}_0 = \text{Cost}_1 = \text{Cost}_2$

What if supply > demand

$\sum \text{supply} = S_i \Rightarrow \leq S_i$ or create dummy demand
node w/ $D = \sum S_i - \sum D_i$
w/ zero cost to ship to

II. Assignment Problems

* This is a transportation problem where all S_i 's + D_i 's are 1
(students are "shipping" their labor to the tasks)



This example actually
doesn't have a feasible
assignment!

III. Min Cost Flow Problems

Formulate as an LP

$$\begin{aligned} \min \quad & 2x_{13} + x_{14} + 3x_{21} + 6x_{23} + 4x_{34} + x_{35} \\ \text{s.t.} \quad & 200 + x_{21} = x_{13} + x_{14} && \left(\begin{array}{l} \text{Flow balance at} \\ \text{node 1} \end{array} \right) \\ & 500 = x_{21} + x_{23} && (\text{node 2}) \\ & x_{13} + x_{23} = x_{34} + x_{35} && (\text{node 3}) \\ & x_{14} + x_{34} = 300 && (\text{node 4}) \\ & x_{35} = 400 && (\text{node 5}) \\ & x_{ij} \geq 0 && x_{14} \leq 100 \quad x_{13} \leq 300 \\ & && x_{21} \leq 100 \quad x_{35} \leq 400 \end{aligned}$$

-Integrality still holds! (for integer S + D + UB)

IV. Shortest Path Problem

* MCF = min cost flow

Source: (A) - 1 unit of supply

Sink: (F) - 1 unit of demand

Do we have to put a capacity on the arcs?