

13.1.1

(a) m, n, i, p

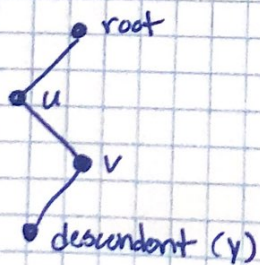
(b) f, g, e, d, a, b, c

(c) i, g, n, e, a, b, c, g

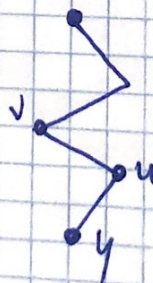
(d) if p is level 0, d is level 3

13.1.3

(a)



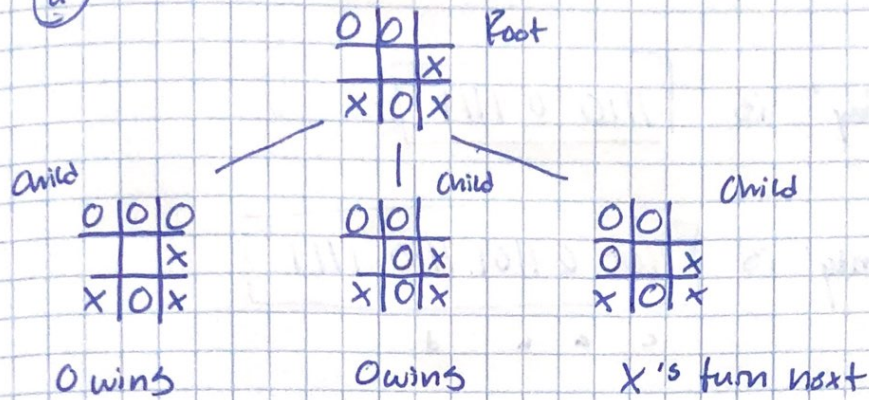
or



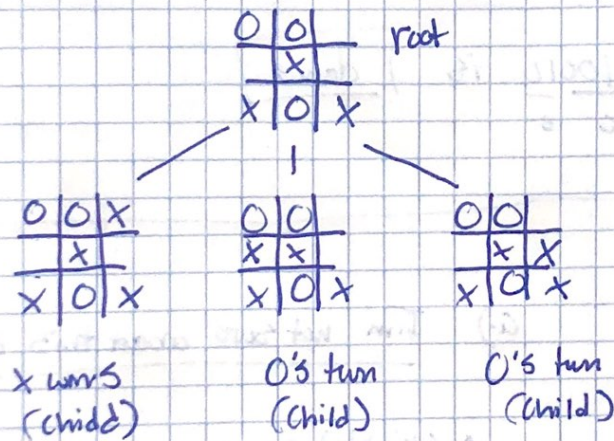
y has a unique path then u and v must be on that path. thus, u is a descendant of v or vice versa.

13.2.1

(a.)



(b.)



13.2.2

a. 'day' is 1110 0 1111

b. 'Candy' is 1100 0 1101 1110 1111
c a n d y

c. 1110101101 is den

d. 11100110110010 is dance
d a n c e

13.2.3

	bits
a :	2
b :	2
c :	3
d :	3
e :	3
f :	3

a.) I'm not sure about this question *

• weighted avg. =

$$\frac{2(.05) + 2(.05) + 3(.1) + 3(.15) + 3(.25) + 3(.4)}{1}$$

= 2.9

• normal average

$$(2 + 2 + 3 + 3 + 3 + 3) / 6 = \span style="border: 1px solid black; padding: 2px;">2.66$$

13.2.3

(b.)

No tree would result in fewer bits

13.3.1

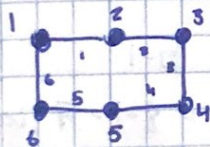
(a.)

There is no tree w/ 7 vertices and total degree 14

each vertex must have ~~a~~ degree of 2 but

7 is odd. or $\sum \deg(v) = 2 \cdot |E| = 2 \cdot 6 = 12 \neq 14$

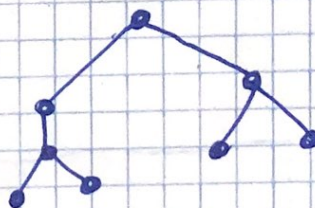
(b.)



(c.)



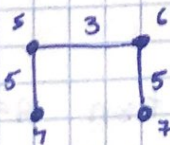
or



not counting root

counting root

(d.)



13.3.2

(a.)

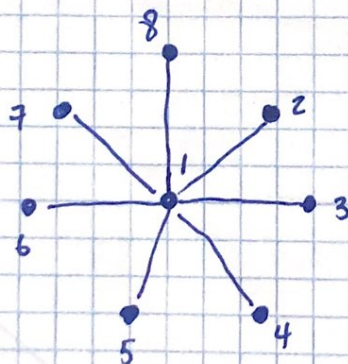
$n=1$ 1 leaf

$n=2$ 2 leaves

$n=3$ 2 leaves

For $n \geq 3$ most leaves $= n-1$

(b.)



$$n-1 \\ = 8-1 = 7 \text{ leaves}$$

13.3.3

(c.)

n vertices

m edges

$$m = n-1$$

} graph is a tree

Converse: m edges then $n = m+1 \Rightarrow m = n-1$

so yes

13.3.5

We know that for a connected graph

~~$m = n - 1$~~ where $m = \text{edges}$ and $n = \text{vertices}$
 $m \geq n - 1$

m can be equal to or greater than $n - 1$



$$\begin{aligned} n &= 3 \\ m &= n - 1 = 2 \\ m &= n - 1 \end{aligned}$$



$$\begin{aligned} n &= 4 \\ m &= 4 \\ m &\geq n - 1 \end{aligned}$$

\therefore if $m < n - 1$ you will have a disconnect

somewhere $\rightarrow \bullet \leftarrow$

try to create a connected graph w/ the two examples above using $m < n - 1 \Rightarrow m = n - 2$

13.4.1

(a.)

Postorder: f, i, h, e, b, g, c, a, d

(b.)

Preorder: d, f, b, i, h, e, a, c, g