

# NAIL062 P&P LOGIC: WORKSHEET 5 – THE RESOLUTION METHOD

**Teaching goals:** After completing, the student

- knows the necessary concepts of the resolution method (resolution rule, resolvent, resolution proof/refutation, resolution tree), can formally define them and provide examples
- can work with formulas in CNF and their models in set representation
- can construct a resolution refutation of a given (even infinite) CNF formula (if it exists), and also draw the corresponding resolution tree
- knows the notion of a substitution tree, can formally define it and construct it for a concrete CNF formula
- can apply the resolution method to solve a given problem (word problem, etc.)

## IN-CLASS PROBLEMS

**Problem 1.** Let  $\varphi$  denote the formula  $\neg(p \vee q) \rightarrow (\neg p \wedge \neg q)$ . Show that  $\varphi$  is a tautology:

- (a) Convert  $\neg\varphi$  to CNF and write the resulting formula as  $S$  in set representation.
- (b) Find a resolution refutation of  $S$ .

**Problem 2.** Prove by resolution that in  $T = \{\neg p \rightarrow \neg q, \neg q \rightarrow \neg r, (r \rightarrow p) \rightarrow s\}$  the statement  $s$  holds.

**Problem 3.** Let the propositional variables  $r, s, t$  represent (respectively) that “*Radka / Sára / Tom is at school*” and write  $\mathbb{P} = \{r, s, t\}$ . We know that:

- *If Tom is not at school, then Sára is not there either.*
  - *Radka does not go to school without Sára.*
  - *If Radka is not at school, then Tom is there.*
- (a) Formalize our knowledge as a theory  $T$  in the language  $\mathbb{P}$ .
  - (b) Using the resolution method, prove that  $T$  entails that *Tom is at school*: Write a formula  $S$  in set representation that is unsatisfiable exactly when this holds, and find a resolution refutation of  $S$ . Draw the resolution tree.
  - (c) Determine the set of models of the theory  $T$ .

**Problem 4.** Construct a *substitution tree* for the following formula. Based on this tree, construct a resolution refutation according to the procedure from the proof of the Completeness Theorem for resolution.

$$S = \{\{p, r\}, \{q, \neg r\}, \{\neg q\}, \{\neg p, t\}, \{\neg s\}, \{s, \neg t\}\}$$

## EXTRA PRACTICE

**Problem 5.** Find a resolution refutation of the following formulas:

- (a)  $\neg(((p \rightarrow q) \rightarrow \neg q) \rightarrow \neg q)$
- (b)  $(p \leftrightarrow (q \rightarrow r)) \wedge ((p \leftrightarrow q) \wedge (p \leftrightarrow \neg r))$

**Problem 6.** Tonia and Fabio describe to us their latest recipe for the best pizza in the world.

- Tonia said: “The recipe includes anchovies or basil or garlic.”
- Tonia also said: “If cooked ham is not included, then basil is not included either.”
- Fabio said: “Cooked ham is included in the recipe.”
- Fabio further said: “Neither anchovies nor basil are included, but garlic is included.”

We know that Tonia always tells the truth, while Fabio always lies.

- (a) Express our knowledge as a propositional theory  $T$  in the language  $\mathbb{P} = \{a, b, c, d\}$ , where the propositional variables mean (in order) “anchovies / basil / garlic / cooked ham are included in the recipe”.
- (b) Using the resolution method, prove that  $T$  entails that “anchovies are included in the recipe”. Draw the resolution tree.

**Problem 7.** The integers are afflicted by a mysterious disease spreading (in discrete steps) according to the following rules (valid for all integers at all time steps).

- (i) *A healthy number becomes ill exactly when precisely one neighboring number was ill (in the previous time).*
- (ii) *An ill number recovers exactly when the previous number was ill (in the previous time).*
- (iii) *At time 0 the number 0 was ill, all other numbers were healthy.*
- (a) Write theories  $T_1, T_2, T_3$  expressing (respectively) statements (i), (ii), (iii) over the set of propositional atoms  $\mathbb{P} = \{p_i^t \mid i \in \mathbb{Z}, t \in \mathbb{N}_0\}$ , where the atom  $p_i^t$  expresses that “number  $i$  is ill at time  $t$ .”
- (b) Convert the axioms from  $T_1, T_2, T_3$  to CNF and write a theory  $S$  in set representation that is unsatisfiable exactly when  $T_1 \cup T_2 \cup T_3 \models \neg p_1^2$ , i.e. when “Number 1 is healthy at time 2.” (It is enough to convert only the specific axioms from  $T_1, T_2, T_3$  that entail  $\neg p_1^2$ , and include only the corresponding clauses in  $S$ .)
- (c) Prove by resolution that  $S$  is unsatisfiable. Show the refutation as a resolution tree.

#### FOR FURTHER THOUGHT

**Problem 8.** Prove in detail that if  $S = \{C_1, C_2\}$  is satisfiable and  $C$  is the resolvent of  $C_1$  and  $C_2$ , then  $C$  is satisfiable as well.