NAIL062 P&P Logic: Worksheet 5 – The resolution method

Teaching goals: After completing, the student

- knows the necessary concepts of the resolution method (resolution rule, resolvent, resolution proof/refutation, resolution tree), can formally define, provide examples
- can work with propositions in CNF and their models in set representation
- can construct a resolution refutation of a given (even infinite) CNF formula (if it exists), and also draw the corresponding resolution tree
- knows the notion of a tree of reductions, can formally define it and construct it for a concrete CNF formula
- can apply the resolution method to solve a given problem (word problem, etc.)

In-class problems

Problem 1. Let φ denote the proposition $\neg(p \vee q) \to (\neg p \wedge \neg q)$. Show that φ is a tautology:

(a) Convert $\neg\varphi$ to CNF and write the resulting formula as S in set representation.
(b) Find a resolution refutation of S.

Problem 2. Prove by resolution that in $T = \{\neg p \to \neg q, \neg q \to \neg r, (r \to p) \to s\}$ the proposition s is valid.

Problem 3. Let the propositional variables r, s, t represent (respectively) that *"Robin / Sam / Tom is at school"* and write $\mathbb{P} = \{r, s, t\}$. We know that:

- *If Tom is not at school, then Sam is not there either.*
- *Robin does not go to school without Sam.*
- *If Robin is not at school, then Tom is there.*

(a) Formalize our knowledge as a theory T in the language $\mathbb{P}$.
(b) Using the resolution method, prove that it is a consequence of T that *Tom is at school*: Write a formula S in set representation that is unsatisfiable exactly when this holds, and find a resolution refutation of S. Draw the resolution tree.
(c) Determine the set of models of the theory T.

Problem 4. Construct a *tree of reductions* for the following formula. Based on this tree, construct a resolution refutation according to the procedure from the proof of the theorem on Completeness of Resolution.

$$S = \{\{p, r\}, \{q, \neg r\}, \{\neg q\}, \{\neg p, t\}, \{\neg s\}, \{s, \neg t\}\}$$

Extra practice

Problem 5. Find a resolution refutation of the following formulas:

(a) $\neg(((p \to q) \to \neg q) \to \neg q)$
(b) $(p \leftrightarrow (q \to r)) \wedge ((p \leftrightarrow q) \wedge (p \leftrightarrow \neg r))$

Problem 6. Tonia and Fabio describe to us their latest recipe for the best pizza in the world.

- Tonia said: "The recipe includes anchovies or basil or capers."
- Tonia also said: "If duck is not included, then basil is not included either."
- Fabio said: "Duck is included in the recipe."
- Fabio further said: "Neither anchovies nor basil are included, but capers are included."

We know that Tonia always tells the truth, while Fabio always lies.

(a) Express our knowledge as a propositional theory T in the language $\mathbb{P} = \{a, b, c, d\}$, where the propositional variables mean (in order) "anchovies / basil / capers / duck are included in the recipe".

(b) Using the resolution method, prove that T implies that "anchovies are included in the recipe". Draw the resolution tree.

Problem 7. The integers are afflicted by a mysterious disease spreading (in discrete steps) according to the following rules (valid for all integers at all time steps).

(i) A healthy number becomes ill exactly when precisely one neighboring number was ill (in the previous time step).

(ii) An ill number recovers exactly when the previous number was ill (in the previous time step).

(iii) At time 0 *the number* 0 *was ill, all other numbers were healthy.*

(a) Write theories T_1, T_2, T_3 expressing (respectively) statements (i), (ii), (iii) over the set of propositional variables $\mathbb{P} = \{p_i^t \mid i \in \mathbb{Z}, t \in \mathbb{N}_0\}$, where p_i^t expresses that "*number i is ill at time t.*"

(b) Convert the axioms from T_1, T_2, T_3 to CNF and write a theory S in set representation that is unsatisfiable exactly when $T_1 \cup T_2 \cup T_3 \models \neg p_1^2$, i.e. when "*Number* 1 *is healthy at time* 2." (It is enough to convert only the specific axioms from T_1, T_2, T_3 that imply $\neg p_1^2$, and include only the corresponding clauses in S.)

(c) Prove by resolution that S is unsatisfiable. Show the refutation as a resolution tree.

For further thought

Problem 8. Prove in detail that if $S = \{C_1, C_2\}$ is satisfiable and C is the resolvent of C_1 and C_2, then C is satisfiable as well.