

# NOPT042 Constraint programming: Tutorial 8 - Global constraints, Routing

## What was in Lecture 5

### Path consistency

- **arc consistency**: never destroys solutions, sometimes can find a solution (without backtracking) -- iff all domains reduced to 1 element
- **path consistency**: for any path in variables, if assignment of start point and end point satisfy all binary constraints between them, then there is a consistent path in the constraint network
- enough to enforce for paths of length two ("every edge [in the constraint network] extends to any triangle")
- PC is stronger, removes pairs of inconsistent values
- but it is more expensive
- algorithms: PC-1, PC-2 [, PC-3, PC-4, PC-5]
- directional path consistency
- drawbacks of PC: memory consumption, bad strength/efficiency ration, modifies the constraint network (adds redundant constraints, changes connectivity, ruins graph-structure-based heuristics), still not complete
- restricted path consistency (AC, only check PC for pairs which are the only support for one of the values)

```
In [1]: %load_ext ipicat
```

Picat version 3.9

## Exercise: Magic sequence

A magic sequence of length  $n$  is a sequence of integers  $x_0, \dots, x_{n-1}$  between 0 and  $n - 1$ , such that for all  $i \in \{0, \dots, n - 1\}$ , the number  $i$  occurs exactly  $x_i$ -times in the sequence. For instance, 6,2,1,0,0,0,1,0,0,0 is a magic sequence since 0 occurs 6 times in it, 1 occurs twice, etc.

(Problem from [the book](#).)

We want to maximize the sum of the numbers in the sequence. Write a constraint model that does that (for a given length  $n$ ).

The constraint `global_cardinality`

```
global_cardinality(List, Pairs)
```

Let `List` be a list of integer-domain variables `[X1, . . . , Xd]`, and `Pairs` be a list of pairs `[K1-V1, . . . , Kn-Vn]`, where each key `Ki` is a unique integer, and each `Vi` is an integer-domain variable. The constraint is true if every element of `List` is equal to some key, and, for each pair `Ki-Vi`, exactly `Vi` elements of `List` are equal to `Ki`. This constraint can be defined as follows:

```
global_cardinality(List,Pairs) =>
  foreach($Key-V in Pairs)
    sum([B : E in List, B#<=>(E#=Key)]) #= V
  end.
```

---from [the guide](#)

```
In [2]: !time picat magic-sequence/.solution/magic-sequence.pi 4
```

```
CPU time 0.0 seconds. Backtracks: 2
```

```
[1,2,1,0]
```

```
real    0m0.020s
user    0m0.010s
sys     0m0.005s
```

```
In [3]: !cat magic-sequence/.solution/magic-sequence.pi
```

```
% Adapted from Constraint Solving and Planning with Picat, Springer
% by Neng-Fa Zhou, Hakan Kjellerstrand, and Jonathan Fruhman
import cp.
```

```
main([N]) =>
  N := N.to_int,
  magic_sequence(N,Sequence),
  println(Sequence).
```

```
magic_sequence(N, Sequence) =>
  Sequence = new_list(N),
  Sequence :: 0..N-1,
```

```
% create list: [0-Sequence[1], 1-Sequence[2], ...]
Pairs = [$I-Sequence[I+1] : I in 0..N-1],
global_cardinality(Sequence,Pairs),
```

```
time2(solve(Sequence)).
```

```
In [4]: !time picat magic-sequence/.solution/magic-sequence.pi 64
!time picat magic-sequence/.solution/magic-sequence2.pi 64
!time picat magic-sequence/.solution/magic-sequence2.pi 400
```

[illegible][illegible][illegible]

```
!cat magic-sequence/.solution/magic-sequence2.pi
```

```

% Adapted from Constraint Solving and Planning with Picat, Springer
% by Neng-Fa Zhou, Hakan Kjellerstrand, and Jonathan Fruhman
import cp.

main([N]) =>
    N := N.to_int,
    magic_sequence(N,Sequence),
    println(Sequence).

magic_sequence(N, Sequence) =>
    Sequence = new_list(N),
    Sequence :: 0..N-1,

    % create list: [0-Sequence[1], 1-Sequence[2], ...]
    Pairs = [$I-Sequence[I+1] : I in 0..N-1],
    global_cardinality(Sequence,Pairs),

    % extra/redudant (implicit) constraints to speed up the model
    N #= sum(Sequence),
    Integers = [I : I in 0..N-1],
    scalar_product(Integers, Sequence, N),

    time2(solve([ff], Sequence)).

```

```

In [6]: !time picat magic-sequence/.solution/magic-sequence2.pi 1024
!time picat magic-sequence/.solution/magic-sequence3.pi 1024

```

CPU time 2.308 seconds. Backtracks: 7

[illegible]

```
real    3m33.519s
user    3m19.980s
sys     0m13.465s
```

CPU time 0.483 seconds. Backtracks: 7

[illegible]

[illegible]

```
real    0m3.987s
user    0m3.641s
sys     0m0.340s
```

```
In [7]: !cat magic-sequence/.solution/magic-sequence3.pi
```

```
% Adapted from Constraint Solving and Planning with Picat, Springer
% by Neng-Fa Zhou, Hakan Kjellerstrand, and Jonathan Fruhman
import cp.
```

```
main([N]) =>
  N := N.to_int,
  magic_sequence(N,Sequence),
  println(Sequence).

magic_sequence(N, Sequence) =>
  Sequence = new_list(N),
  Sequence :: 0..N-1,

  % extra/redudant (implicit) constraints to speed up the model
  N #= sum(Sequence),
  Integers = [I : I in 0..N-1],
  scalar_product(Integers, Sequence, N),

  %% create list: [0-Sequence[1], 1-Sequence[2], ...]
  Pairs = [$I-Sequence[I+1] : I in 0..N-1],
  global_cardinality(Sequence,Pairs),

  time2(solve([ff], Sequence)).
```

## The order of constraints

The order might matter to the solver, the above model is an example. When the `cp` solver parses a constraint, it tries to reduce the domains. The implicit (redundant) constraint using `scalar_product` can reduce the model a lot, which is much better to do before parsing `global_cardinality`. (Note: e.g. MiniZinc doesn't preserve the order of constraints during compilation, the behaviour is a bit unpredictable.)

*Heuristic:* easy constraints and constraints that are strong [in reducing the search space] should go first (?)

# Routing

Routing, aka the [Arc Routing Problem](#), refers to a variety of constraint satisfaction problems related to the well-known [Vehicle Routing Problem](#). Such

problems appear in a variety of applications such as public transport (bus routing), garbage collection (not that one), snow ploughing, and many others.

## Exercise: Knight tour

Given an integer  $N$ , plan a tour of the knight on an  $N \times N$  chessboard such that the knight visits every field exactly once and then returns to the starting field. You can assume that  $N$  is even.

(Hint: For a matrix `M` is a matrix, you can use `M.vars()` to extract its elements into a list.)

## The `circuit` constraint

The constraint `circuit(L)` requires that the list  $L$  represents a permutation of  $1, \dots, n$  consisting of a single cycle, i.e., the graph with edges  $i \rightarrow L[i]$  is a cycle. A similar constraint is `subcircuit(L)` which requires that elements for which  $L[i] \neq i$  form a cycle.

```
In [8]: !picat knight-tour/.solution/knight-tour.pi 8
```

```
x = {{11,12,13,10,15,21,22,14},{3,20,26,18,23,4,30,6},{2,1,29,5,27,32,8,7},{19,9,37,34,35,40,16,38},{43,17,50,53,52,28,24,46},{51,25,49,61,60,36,62,31},{59,33,57,58,63,64,45,39},{42,41,44,54,55,56,48,47}}
```

```
X:
```

```
11 12 13 10 15 21 22 14
 3 20 26 18 23  4 30  6
 2  1 29  5 27 32  8  7
19  9 37 34 35 40 16 38
43 17 50 53 52 28 24 46
51 25 49 61 60 36 62 31
59 33 57 58 63 64 45 39
42 41 44 54 55 56 48 47
```

```
Tour:
```

```
1 62  5 10 13 24 55  8
 4 11  2 63  6  9 14 23
61 64 35 12 25 56  7 54
34  3 26 59 36 15 22 57
39 60 37 18 27 58 53 16
30 33 40 43 46 17 50 21
41 38 31 28 19 48 45 52
32 29 42 47 44 51 20 49
```

```
In [9]: !cat knight-tour/.solution/knight-tour.pi
```

```

/*****
Adapted from
knight_tour.pi
from Constraint Solving and Planning with Picat, Springer
by Neng-Fa Zhou, Hakan Kjellerstrand, and Jonathan Fruhman
*****/
import cp.

main([N]) =>
    N := N.to_int,
    knight(N,X),
    println(x=X),
    println("X:"),
    print_matrix(X),
    extract_tour(X,Tour),
    println("Tour:"),
    print_matrix(Tour).

% Knight's tour for even N*N.
knight(N, X) =>
    X = new_array(N,N),
    X :: 1..N*N,
    XVars = X.vars(),
    % restrict the domains of each square
    foreach (I in 1..N, J in 1..N)
        D = [-1,-2,1,2],
        Dom = [(I+A-1)*N + J+B : A in D, B in D,
                abs(A) + abs(B) == 3,
                member(I+A,1..N), member(J+B,1..N)],
        Dom.length > 0,
        X[I,J] :: Dom
    end,
    circuit(XVars),
    solve([ff,split],XVars).

extract_tour(X,Tour) =>
    N = X.length,
    Tour = new_array(N,N),
    K = 1,
    Tour[1,1] := K,
    Next = X[1,1],
    while (K < N*N)
        K := K + 1,
        I = 1+((Next-1) div N),
        J = 1+((Next-1) mod N),
        Tour[I,J] := K,
        Next := X[I,J]
    end.

print_matrix(M) =>
    N = M.length,
    V = (N*N).to_string().length,
    Format = "% " ++ (V+1).to_string() ++ "d",
    foreach(I in 1..N)
        foreach(J in 1..N)
            printf(Format,M[I,J])

```



```
end,  
n1  
end,  
n1.
```

## Exercise: Chinese Postman

Given a weighted graph, find the shortest closed path that traverses every edge at least once. (Assume that the graph is undirected, loopless, and the distances are positive integers and satisfy the triangle inequality.) See [Wikipedia](#) for a sample solved instance. Unlike the [Eulerian Circuit Problem](#) (where each edge has to be traversed exactly once), there is a polynomial-time algorithm for the Chinese Postman. But our goal is to write a constraint model.

```
In [10]: !cat chinese-postman/instance.pi
```

```
instance(NumVertices, AdjacencyMatrixWithWeights) =>  
    NumVertices = 5,  
    AdjacencyMatrixWithWeights = [  
        [0, 7, 0, 5, 1],  
        [7, 0, 6, 0, 3],  
        [0, 6, 0, 8, 4],  
        [5, 0, 8, 0, 2],  
        [1, 3, 4, 2, 0]  
    ].  
% optimal length: 45, shortest closed path: [1,2,3,4,5,3,2,5,1,5,4,1]
```