## NTIN071 A&G: Tutorial 10 – Pushdown automata

Solve 1a-h, 2ab, 3a first (the rest is for practice).

**Problem 1** (Constructing PDA). Design pushdown automata for the following languages. (Acceptance can be either by final state or by empty stack, in some cases construct both.)

(a) 
$$L = \{ w \mid w \in \{0, 1\}^*, |w|_1 \ge 3 \}$$

(b) 
$$L = \{ww^R \mid w \in \{0, 1\}^*\}$$

(c) 
$$L = \{w \in \{(,)\}^* \mid w \text{ is correctly parenthesized}\}$$

(d) 
$$L = \{w \in \{0, 1\}^* \mid w = w^R\}$$

(e) 
$$L = \{a^i b^j c^k \mid i = j \text{ or } j = k\}$$

(f) 
$$L = \{a^i b^j c^k \mid i + j = k\}$$

(g) 
$$L = \{a^{2n}b^{3n} \mid n \ge 0\}$$

(h) 
$$L = \{w \in \{0,1\}^* \mid |w|_0 = |w|_1\}$$

(i) 
$$L = \{u2v \mid u, v \in \{0, 1\}^* \text{ and } |u| \neq |v|\}$$

(j) 
$$L = \{w \in \{(,),[,]\}^* \mid w \text{ is correctly parenthesized}\}$$

**Problem 2** (Acceptance by final state vs. empty stack). Convert selected PDA constructed in the previous problem from acceptance by final state to acceptance by empty stack, and vice versa. (Try both constructions.)

**Problem 3** (CFG to PDA). For a given context-free grammar G construct pushdown automata  $P_1, P_2$  such that  $L(G) = N(P_1) = L(P_2)$ .

(a) 
$$G = (\{S, T, X\}, \{a, b\}, \mathcal{P}, S)$$

$$\mathcal{P} = \{ S \to aTXb, \\ T \to XTS \mid \epsilon, \\ X \to a \mid b \}$$

(b) 
$$G=(\{S,T,X\},\{(,),*,+,,1\},\mathcal{P},S)$$
 
$$\mathcal{P}=\{S\to S+T\mid T,$$
 
$$T\to T*X\mid X,$$
 
$$X\to 1\mid (S)\}$$

For a reasonably long word  $w \in L(G)$  find its leftmost derivation from G and simulate the automaton  $P_1$  on the input w.

**Problem 4** (PDA to CFG). Convert selected (small) pushdown automata constructed in Problem 1 to context-free grammars. For a reasonably long word w accepted by the automaton find a leftmost derivation of w from the grammar.

BONUS: CONTEXT GRAMMARS

**Problem 5** (A context grammar). Consider the grammar  $G = (\{S, A, B, C\}, \{a, b, c\}, S, P)$ , where

$$P = \{S \rightarrow aSBC \mid aBC, B \rightarrow BBC, C \rightarrow CC, CB \rightarrow BC, \\ aB \rightarrow ab, bB \rightarrow bb, bC \rightarrow bc, cC \rightarrow cc\}$$

What language does it generate? Is G a context grammar? If not, find an equivalent context grammar.