NTIN071 A&G: Tutorial 4 – Closure properties of regular languages

Teaching goals: The student is able to

- formally describe a construction of an automaton based on other automata
- decide whether regular languages are closed under various set and string operations, including more complex ones, and prove or disprove it

IN-CLASS PROBLEMS

Problem 1 (Closure under set and string operations). Given DFAs A, B, construct an automaton C recognizing the given language. (Give a formal description of the automaton.)

(a) $L(A) - L(B)$	A	a	b	В	a	b
(b) $L(A).L(B)$	$\rightarrow 0$	1	2	$\rightarrow 0$	0	5
	* 1	3	0	* 1	1	3
(c) $L(A)^{+}$	2	4	5	2	2	5
(d) $L(A)^*$	3	0	2	3	3	2
	4	2	5	* 4	6	1
(e) $L(A)^R$	5	0	3	5	5	1
		'		* 6	4	2

Problem 2 (Delete). Let L be a regular language over the alphabet $\Sigma = \{a, b\}$. Desribe the following languages in set notation. Decide if they are (necessarily) also regular, prove or disprove. The language of all words obtained from words of the language L by...

- (a) ... deleting all occurrences of the letter a.
- (b) ... deleting the initial letter and writing this letter at the end of the word.
- (c) ... deleting the longest contiguous sequence of a's from the beginning of the word.

EXTRA PRACTICE AND THINKING

Problem 3 (Shift). Given a regular language L over an alphabet Σ , define the language L' as follows. Is the language L' necessarily regular?

$$L' = \{uv \mid u, v \in \Sigma^*, vu \in L\}$$

Problem 4 (Cut). Consider two regular languages L and M over an alphabet Σ , and define the language K as follows. Is the language K necessarily regular?

$$K = \{uw \mid u, w \in \Sigma^*, (\exists v \in M) \ uvw \in L\}$$

Problem 5 (Switch final and nonfinal states). If we switch accepting and nonaccepting states in a given NFA, will the language recognized by the resulting automaton be the complement of the language recognized by the original NFA? Justify your answer.