NTIN071 A&G: Tutorial 6 – Formal grammars, regular and context-free GRAMMARS

Teaching goals: The student is able to

- explain the formal definition of a grammar, and the language it generates,
- give definitions and examples grammars of all types in the Chomsky hierarchy,
- describe a language generated by a given context-free grammar,
- construct a grammar for a language given in set notation,
- convert a finite automaton to a right-linear grammar,
- convert a right-linear grammar to a finite automaton,
- design algorithms to test basic properties of context-free grammars.

In-class problems

Problem 1 (Constructing grammars). Design grammars (of the highest possible type) which generate the following languages ($\Sigma = \{a, b\}$ unless specified otherwise):

(a)
$$L = \{ w \in \Sigma^* \mid |w|_b \text{ is even} \}$$

(d)
$$L = \{a^i b^j \mid 0 \le i \le j \le 2i\}$$

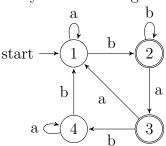
(b)
$$L = \{ww^R \mid w \in \Sigma^*\}$$

(e)
$$L = \{uabbav \mid u, v \in \Sigma^* \text{ and } |u| = |v|\}$$

(c)
$$L = \{a^i b^j c^k \mid i = j \text{ or } j = k\}$$

$$\begin{array}{lll} \text{(a)} \ L = \{w \in \Sigma^* \mid |w|_b \text{ is even}\} & \text{(d)} \ L = \{a^i b^j \mid 0 \leq i \leq j \leq 2i\} \\ \text{(b)} \ L = \{w w^R \mid w \in \Sigma^*\} & \text{(e)} \ L = \{uabbav \mid u, v \in \Sigma^* \text{ and } |u| = |v|\} \\ \text{(c)} \ L = \{a^i b^j c^k \mid i = j \text{ or } j = k\} & \text{(f)} \ L = \{w \in \Sigma^* \mid -1 \leq |w|_a - |w|_b \leq 1\} \\ \end{array}$$

Problem 2 (FA to grammar). For the following automaton, find an equivalent grammar. Which class of the Chomsky hierarchy does it belong to?



Problem 3 (Regular grammar to FA). Convert the following right-linear grammar to a finite automaton: $G = (\{S, A, B, C\}, \{a, b\}, \mathcal{P}, S)$ where \mathcal{P} consists of the following:

$$S \rightarrow abS \mid babA \mid \epsilon$$

$$A \rightarrow abA \mid aB \mid bC$$

$$B \rightarrow abS \mid B \mid bC \mid \epsilon$$

$$C \rightarrow aab \mid A \mid aA \mid \epsilon$$

Problem 4 (Testing properties of context-free languages). Design an (efficient) algorithm which decides whether a given context-free grammar satisfies the given property:

(a)
$$L(G) \neq \emptyset$$

(b)
$$\epsilon \in L(G)$$

(c) L(G) is a finite language

EXTRA PRACTICE AND THINKING

Problem 5 (Constructing grammars). Design grammars (of the highest possible type) which generate the following languages ($\Sigma = \{a, b\}$ unless specified otherwise):

- (a) $L = \Sigma^*$
- (a) $L = \{a^{2i}b^{j} \mid i \leq j\}$ (b) $L = \{a^{2i}b^{j} \mid i \leq j\}$ (c) $L = \{w \in \Sigma^{*} \mid |w|_{a} = 2|w|_{b}\}$
- (d) $L = \{uabbav \mid u, v \in \Sigma^* \text{ and } |u| \neq |v|\}$ (e) $L = \{w \# s^R \mid w, s \in \Sigma^* \text{ and } s \text{ is a subword of } w\}$

Problem 6 (Small grammars generating large (finite) languages). Find a sequence of context-free grammars G_1, G_2, G_3, \ldots (over a given alphabet Σ) such that G_n generates exactly all words of length $\leq 2^n$ (and no other words), and the size of G_n (for simplicity, say the number of symbols in bodies of production rules) is in O(n).