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# Mathematical Optimization — Solution 12

https://moodle-app2.let.ethz.ch/course/view.php?id=2180

#### Exercise 1: Submodular Functions

(a) " $\Rightarrow$ ": Let  $S := A \cup \{j\}, T := A \cup \{k\}$ . Then,

$$f(S \cup T) - f(T) \le f(S) - f(S \cap T) \quad \Leftrightarrow \quad f(A \cup \{j,k\}) - f(A \cup \{k\}) \le f(A \cup \{j\}) - f(A).$$

"
$$\Leftarrow$$
": Let  $A, B \subseteq N$ ,  $S := A \cap B$ ,  $A \setminus B := \{j_1, \ldots, j_r\}$ ,  $B \setminus A := \{k_1, \ldots, k_s\}$ . Then,

$$f(B) - f(A \cap B) = f(S \cup \{k_1, \dots, k_s\}) - f(S)$$

$$= \sum_{i=1}^{s} f(S \cup \{k_1, \dots, k_i\}) - f(S \cup \{k_1, \dots, k_{i-1}\})$$

$$\geq \sum_{i=1}^{s} f(S \cup \{k_1, \dots, k_i\} \cup \{j_1\}) - f(S \cup \{k_1, \dots, k_{i-1}\} \cup \{j_1\})$$

$$\geq \dots \geq \sum_{i=1}^{s} f(S \cup \{k_1, \dots, k_i\} \cup \{j_1, \dots, j_r\}) - f(S \cup \{k_1, \dots, k_{i-1}\} \cup \{j_1, \dots, j_r\})$$

$$= \sum_{i=1}^{s} f(A \cup \{k_1, \dots, k_i\}) - f(A \cup \{k_1, \dots, k_{i-1}\}) = f(A \cup B) - f(A).$$

(b) " $\Rightarrow$ ": Let  $S, T \subseteq N, T \setminus S := \{j_1, \dots, j_r\}$ . Then,

$$f(T) \stackrel{\text{non-decreasing}}{\leq} f(S \cup T) = f(S) + (f(S \cup T) - f(S))$$

$$= f(S) + \sum_{i=1}^{r} (f(S \cup \{j_1, \dots, j_i\}) - f(S \cup \{j_1, \dots, j_{i-1}\}))$$

$$\stackrel{\text{(a)}}{\leq} f(S) + \sum_{i=1}^{r} (f(S \cup \{j_i\}) - f(S)).$$

" $\Leftarrow$ ": Setting  $T := S \cup \{j, k\}$  gives

$$f(S \cup \{j, k\}) \le f(S) + f(S \cup \{j\}) - f(S) + f(S \cup \{k\}) - f(S)$$
  
  $\Leftrightarrow f(S \cup \{j\}) - f(S) \ge f(S \cup \{j, k\}) - f(S \cup \{k\}),$ 

which is (a) and therefore implies submodularity.

For monotonicity, let  $T \subseteq S$ . Now,

$$f(T) \le f(S) + \sum_{j \in T \setminus S} f(S \cup \{j\}) - f(S) = f(S).$$

### Exercise 2: Matroid-union

Let  $I, J \in \mathcal{F}$  with |I| < |J|. Without loss of generality, assume that  $|I \cap E_1| < |J \cap E_1|$  (otherwise, exchange the roles of  $M_1$  and  $M_2$ ). Using the hint, there exists  $e \in (J \cap E_1) \setminus (I \cap E_1)$  such that  $(I \cap E_1) \cup e \in \mathcal{F}_1$ . Therefore,  $I \cup \{e\} \in \mathcal{F}$ , as  $(I \cap E_1) \cup e = I \cup e \cap E_1 \in \mathcal{F}_1$  and  $(I \cup e) \cap E_2 = I \cap E_2 \in \mathcal{F}_2$ , where we used that  $e \in E_1 \Rightarrow e \notin E_2$ .

#### Exercise 3: Hamiltonian Paths and Matroid Intersection

We consider the following three matroids for intersection:

 $\mathcal{M}_1$ : Graphic matroid  $\mathcal{M}_1 = (\mathcal{F}_1, E)$ . Note: We ignore the direction of the edges.

 $\mathcal{M}_2$ : Partition matroid  $\mathcal{M}_2 = (\mathcal{F}_2, E)$  where each node  $v \in V$  has at most one outgoing edge:  $\mathcal{F}_2 = \{A \subseteq E \mid \forall v \in V : |\{(v, w) \mid (v, w) \in A\}| \leq 1\}$ 

 $\mathcal{M}_3$ : Partition matroid  $\mathcal{M}_3 = (\mathcal{F}_3, E)$  where each node  $v \in V$  has at most one incoming edge:  $\mathcal{F}_3 = \{A \subseteq E \mid \forall v \in V : |\{(w, v) \mid (w, v) \in A\}| \leq 1\}$ 

Any independent set in  $\mathcal{F}_1 \cap \mathcal{F}_2 \cap \mathcal{F}_3$  represents a disjoint union of directed, simple paths. Therefore, G admits a directed Hamiltonian path if and only if a maximally independent set of cardinality |V|-1 exists in  $\mathcal{F}_1 \cap \mathcal{F}_2 \cap \mathcal{F}_3$ . Note that in the case where there doesn't exist a directed path between any two vertices of G, an independent set of cardinality |V|-1 cannot exist. Determining whether a Hamiltonian path (or hamiltonian cycle) exists in a given graph is  $\mathcal{NP}$ -complete. Hence finding the maximal independent set in the intersection of three matroids is  $\mathcal{NP}$ -hard.