

Consider the following algorithm.

Algorithm 1 function

```
Mystery(n)
r ← 0
for i ← 1 to n - 1 do
  for j ← i + 1 to n do
    for k ← 1 to j do
      r ← r + 1
    end for
  end for
end for
return r
```

- (a) What is the value returned by the function Mystery? Express your answer as a function of n and give the closed form.
- (b) Using $O()$ notation, give the worst-case running time of the function Mystery.
-

This is a very interesting problem, as the name says Mystery.

If you write the same problem in a simple programming language, then it looks something like

```
#include <iostream>
```

```
using namespace std;
```

```
int main() {
```

```
    int n;
```

```
    cin >> n;
```

```
    int r = 0;
```

```
    for (int i = 1; i <= n-1; i++){  
        for (int j = i+1; j <= n; j++){  
            for (int k = 1; k <= j; k++){  
                r++;  
            }  
        }  
    }  
    cout << r << endl;  
    return 0;  
}
```

So it basically has 3 for loops (in nested fashion)

Let's take an example to feel the problem

Let $n = 5$

The outermost 'for' loop goes from $i = 1$ to $i = 4$

And for each for these 'i', the inner 'for' loops work,

Say for $i = 1$

j goes from $i+1$ to n , that is, 2 to 5

and for each value of j , innermost loop goes from $k = 1$ to $k = j$ and increments the value of 'r' in each iteration.

So for $j = 2$

k goes from 1 to 2.

Finally, value of r is returned.

So, r gets incremented two times for j = 2, three times for j = 3, ..., five times for j = 5

In similar manner for other values of i

So final value of $r = (2 + 3 + 4 + 5)$ {corresponding to $i = 1$ } + $(3 + 4 + 5)$ {corresponding to $i = 2$ } + $(4 + 5)$ {corresponding to $i = 3$ } + (5) {corresponding to $i = 4$ }

$$= 1 * 2 + 2 * 3 + 3 * 4 + 4 * 5$$

$$= 40$$

GENERALIZATION {EXPRESSION FOR VALUE of r}:

IF we write in terms of n;

$$\text{Value of } r = 1 * 2 + 2 * 3 + 3 * 4 + \dots + (n-1) * n$$

$$\text{Or write } r = \sum_{i=1}^{n-1} (i) * (i+1) \quad [\text{That is summation from } i = 1 \text{ to } i = n-1]$$

TIME COMPLEXITY:

Since there are 3 for loops, so in worst case,

Time complexity can go as bad as $O(n^3)$.

BONUS: [This is just for extra knowledge, not asked in question]

You can return the same value of r by doing an optimization

Say

Initialize $r = 0$

```
For (int i = 1; i <= n-1; i++){
```

```
     $r = r + (i) * (i+1)$ 
```

```
}
```

Return r

This will also return same value of r, but with complexity just $O(n)$, linear time.

I hope it helps 😊