Seed Matrices

SIAM CSE 2021







Seed Matrices: Theory

Recall:

Forward mode computes

JS

at cost proportional to number of columns in S.

Reverse mode computes

$$W^TJ$$

at cost proportional to number of columns in W.

Seed Matrices: Practice

In reality, it is rarely that simple.

- ► data structure pragmatics
- multiple output variables
- ► multiple input variables
- ► Jacobian compression

Seed Matrices: data structure pragmatics

Depending on

- programming language (row-major vs. column-major arrays)
- choice of AD tool

the J, S, W matrices might be stored in a transposed layout or as part of a struct and the derivative dimension might be flattened.

Consequently, what is logically X_{ij} might be stored as dX(j,i) or x[i].grad[j].

Seed Matrices: multiple scalar outputs, scalar input

- Consider function foo(in:x,out:f,out:g)
- Then, forward mode AD computes

$$\dot{f} = \frac{\partial f}{\partial x}\dot{x}$$
 and $\dot{g} = \frac{\partial g}{\partial x}\dot{x}$

ullet Initializing $\dot{x}=1$ yields the full (logical) Jacobian, split over two data structures.

$$J = \left[egin{array}{c} \dot{f} \ \dot{g} \end{array}
ight]$$

Seed Matrices: multiple scalar outputs, vector input

- Consider function foo(in:x(N),out:f,out:g)
- That is, x is now a vector of length N.
- Then, vector forward mode AD computes

$$\dot{f} = \frac{\partial f}{\partial x}\dot{x}$$
 and $\dot{g} = \frac{\partial g}{\partial x}\dot{x}$

- Initializing \dot{x} to the identity matrix yields the full Jacobian, at a cost proportional to N.
- Alternatively, one could loop over $\dot{x} = e_i$, computing one column of the Jacobian at a time.

Seed Matrices: multiple scalar outputs, vector input

- Consider function foo(in:x(N),out:f,out:g)
- Then, reverse mode AD computes

$$\overline{x} = \overline{f} \frac{\partial f}{\partial x} + \overline{g} \frac{\partial g}{\partial x}$$

- Initializing $\overline{f}=1$ and $\overline{g}=0$ yields $\overline{x}=\frac{\partial f}{\partial x}$ Initializing $\overline{f}=0$ and $\overline{g}=1$ yields $\overline{x}=\frac{\partial g}{\partial x}$

Seed Matrices: multiple scalar outputs, vector input

- Initializing $\overline{f}=1$ and $\overline{g}=0$ yields $\overline{x}=\frac{\partial f}{\partial x}$
- ullet Initializing $\overline{f}=0$ and $\overline{g}=1$ yields $\overline{x}=rac{\partial \widehat{g}}{\partial x}$
- Thus, we are computing W^TJ for each of

$$W = \left[egin{array}{c} 1 \ 0 \end{array}
ight] \ \ ext{and} \ \ W = \left[egin{array}{c} 0 \ 1 \end{array}
ight]$$

• In vector reverse mode, we can let $\overline{f} = [1 \ 0]$ and $\overline{g} = [0 \ 1]$ and compute the full Jacobian all at once, using W = I.

Seed Matrices: multiple scalar inputs, scalar output

- Now, consider function foo(in:x,in:y,out:f)
- Then, reverse mode AD computes

$$\overline{x} = \overline{f} \frac{\partial f}{\partial x}$$
 and $\overline{y} = \overline{f} \frac{\partial f}{\partial x}$

ullet Initializing $\overline{f}=1$ yields the full Jacobian, or gradient, split over two data structures.

$$J = [\overline{x} \ \overline{y}]$$

Seed Matrices: multiple scalar inputs, scalar output

- Consider function foo(in:x,in:y,out:f)
- Then, forward mode AD computes

$$\dot{f} = \frac{\partial f}{\partial x}\dot{x} + \frac{\partial f}{\partial y}\dot{y}$$

- Initializing $\dot{x}=1$ and $\dot{y}=0$ yields $\dot{f}=\frac{\partial f}{\partial x}$ Initializing $\dot{x}=0$ and $\dot{y}=1$ yields $\dot{f}=\frac{\partial f}{\partial v}$

Seed Matrices: multiple scalar inputs, scalar output

- Initializing $\dot{x}=1$ and $\dot{y}=0$ yields $\dot{f}=\frac{\partial f}{\partial x}$ • Initializing $\dot{x}=0$ and $\dot{y}=1$ yields $\dot{f}=\frac{\partial f}{\partial y}$
- Thus, we are computing JS for each of

$$S = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$$
 and $S = \begin{bmatrix} 0 \\ 1 \end{bmatrix}$.

• In vector forward mode, we can let $\dot{x} = [1 \ 0]$ and $\dot{y} = [0 \ 1]$ and compute the full Jacobian all at once, using S = I.

Seed Matrices: multiple vector inputs

- Consider foo(in:x(3),in:y(3),out:f)
- Then, to compute the full Jacobian using vector forward mode AD, we must initialize

$$\dot{x} = \left[egin{array}{ccccc} 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 \end{array}
ight]$$

and

$$\dot{y} = \left[egin{array}{ccccc} 0 & 0 & 0 & 1 & 0 & 0 \ 0 & 0 & 0 & 0 & 1 & 0 \ 0 & 0 & 0 & 0 & 0 & 1 \end{array}
ight]$$

Seed Matrices: Jv and J^Tv products

- Numerical algorithms may only need Jv and/or J^Tv products
- ► Can initialize *S* or *W* to *v* and compute this product at a small multiple of function cost
- ➤ AD tool may not support vector forward or vector reverse mode
- ► Can iterate through columsn of *S* or *W*

Suppose the Jacobian is tridiagonal. That is,

$$J = \begin{bmatrix} a_{11} & a_{12} & 0 & 0 & 0 & 0 \\ a_{21} & a_{22} & a_{23} & 0 & 0 & 0 \\ 0 & a_{32} & a_{33} & a_{34} & 0 & 0 \\ 0 & 0 & a_{43} & a_{44} & a_{45} & 0 \\ 0 & 0 & 0 & a_{54} & a_{55} & a_{56} \\ 0 & 0 & 0 & 0 & a_{65} & a_{66} \end{bmatrix}$$

Suppose the Jacobian is tridiagonal. Then, initializing the seed matrix S to

$$S = \left[egin{array}{cccc} 1 & 0 & 0 \ 0 & 1 & 0 \ 0 & 0 & 1 \ 1 & 0 & 0 \ 0 & 1 & 0 \ 0 & 0 & 1 \end{array}
ight]$$

yields

$$JS = \begin{bmatrix} a_{11} & a_{12} & 0 & 0 & 0 & 0 \\ a_{21} & a_{22} & a_{23} & 0 & 0 & 0 \\ 0 & a_{32} & a_{33} & a_{34} & 0 & 0 \\ 0 & 0 & a_{43} & a_{44} & a_{45} & 0 \\ 0 & 0 & 0 & a_{54} & a_{55} & a_{56} \\ 0 & 0 & 0 & 0 & a_{65} & a_{66} \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

yields

$$\begin{bmatrix} a_{11} & a_{12} & 0 & 0 & 0 & 0 \\ a_{21} & a_{22} & a_{23} & 0 & 0 & 0 \\ 0 & a_{32} & a_{33} & a_{34} & 0 & 0 \\ 0 & 0 & a_{43} & a_{44} & a_{45} & 0 \\ 0 & 0 & 0 & a_{54} & a_{55} & a_{56} \\ 0 & 0 & 0 & 0 & a_{65} & a_{66} \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} = \begin{bmatrix} a_{11} & a_{12} & 0 \\ a_{21} & a_{22} & a_{23} \\ a_{34} & a_{32} & a_{33} \\ a_{44} & a_{45} & a_{43} \\ a_{54} & a_{55} & a_{56} \\ 0 & a_{65} & a_{66} \end{bmatrix}$$

- ▶ We can compute a tridiagonal Jacobian using a seed matrix with 3 columns, independent of N
- ► For general sparse *J*, we can perform a distance-2 coloring of the bipartite graph representing *J* to identify the *structurally orthogonal* columns of *J*
- Cost of computing the compressed Jacobian is proportional to the number of colors
- ► Special colorings for sparsity arising from stencil-based discretizations

- We can compute a tridiagonal Jacobian using a seed matrix with 3 columns, independent of N
- Cost of computing the compressed Jacobian is proportional to the number of colors

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\begin{bmatrix} a_{11} & a_{12} & 0 & 0 & 0 & 0 \\ a_{21} & a_{22} & a_{23} & 0 & 0 & 0 \\ 0 & a_{32} & a_{33} & a_{34} & 0 & 0 \\ 0 & 0 & a_{43} & a_{44} & a_{45} & 0 \\ 0 & 0 & 0 & a_{54} & a_{55} & a_{56} \\ 0 & 0 & 0 & 0 & a_{65} & a_{66} \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix} = \begin{bmatrix} a_{11} & a_{12} & 0 \\ a_{21} & a_{22} & a_{23} \\ a_{34} & a_{32} & a_{33} \\ a_{44} & a_{45} & a_{43} \\ a_{54} & a_{55} & a_{56} \\ 0 & a_{65} & a_{66} \end{bmatrix}
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