

2.72

- 1) Because word is defined as an unsigned integer, all integers in this line will be casted to unsigned making it impossible to return a signed integer. In addition, because we are doing a logical right shift and then & 0xFF, this implementation is unable to account for negative numbers.

2) `int xbyte (packed_t word, int bytenum)`  
`{`  
`return (int)word << ((3-bytenum) << 3) >> 24;`  
`}`

2.81

- a.  $(x < y) == (-x > -y)$   
 Can yield 0 because if x is -2147483648 and y is -2, -x would still be -2147483648 and -y would be 2. In this case,  $(x < y)$  is true but  $(-x > -y)$  is not true.
- b.  $((x + y) << 4) + y - x == 17*y + 15*x$   
 Always yields 1 because the left side of this expression,  $(x+y) << 4$  is equivalent to  $16(x+y)$ . If we distribute, add y, and subtract x, we will be left with  $15*x + 17*y$  which is equal to the right side of the equation,  $17*y+15*x$ .
- c.  $\sim x + \sim y + 1 == \sim(x + y)$ :  
 Always yields 1 because when we take the  $\sim$  of a two's complement number, we are actually switching it's sign and subtracting one ( $\sim x = -x - 1$ ). In the context of this expression,  $\sim x + \sim y + 1 = -x - 1 + -y - 1 + 1 = -x + -y - 1$ . On the other side of the expression  $\sim(x+y) = -x + -y - 1$ . The left and right side of this expression are equal, so this test evaluates to true.
- d.  $(ux - uy) == -(\text{unsigned})(y - x)$   
 Always yields 1 because taking the negative of any unsigned integer equals  $2^{32}$  minus itself ( $-ux = 2^{32} - ux$ ). Knowing this, the left side of this expression can be rewritten as  $(ux + 2^{32} - uy)$ . Likewise, the right side can be written as  $2^{32} - (uy - ux) \rightarrow (2^{32} - uy + ux)$ . Comparing the left and right sides, we should get  $(2^{32} + ux - uy) == (2^{32} + ux - uy)$  which evaluates to true.
- e.  $((x >> 2) << 2) <= x$ :  
 Always yields 1 because when we shift x to the right, we perform an arithmetic shift which preserves the sign of the integer ( $01111. . . >> 2 \rightarrow 00011. . .$  or  $11010. . . >> 2 \rightarrow 11110. . .$ ). However, when we shift back to the left, we fill the lowest significant bits with 0s making the new x ultimately smaller than the original x.  
 Ex.  $x = 10101010 = -86$   
 $x >> 2 = 11101010 = -22$   
 $\rightarrow << 2 = 10101000 = -88$