LIFE OF A HELLO PROGRAM

% CSAPP-CHAPTER 1

TRANSLATION

SOURCE FILE

- interpretation of source file
  - sequence of bits (as file)
  - text characters (as text file)
  - program of C statements, syntax (as C source file)
- encoding: context decides representation

TRANSLATION

```
gcc hello.c
```

1. from: C statement @source file
2. to: machine-language instructions @executable object file

COMPILATION SYSTEMS

Four phases

1. Preprocessor (cpp): directives (#), `.c`,`.h` -> `.i` (text file)
2. compiler (cc1): `.i`->`.s` (text file)
   - `.s` is assembly-lang. program, each statement is 1-to-1 mapped to a machine-lang. instruction
3. assembler (as): `.s`->`.o`
   - relocatable object program, binary file: encode instructions not characters
4. linker (ld): `.o`, `.o`->`.o`.
   - merge multiple relocatable objects to a single executable object

EXECUTION

```
./a.out
```

- shell: command-line interpreter
- shell loads and execute object `a.out`

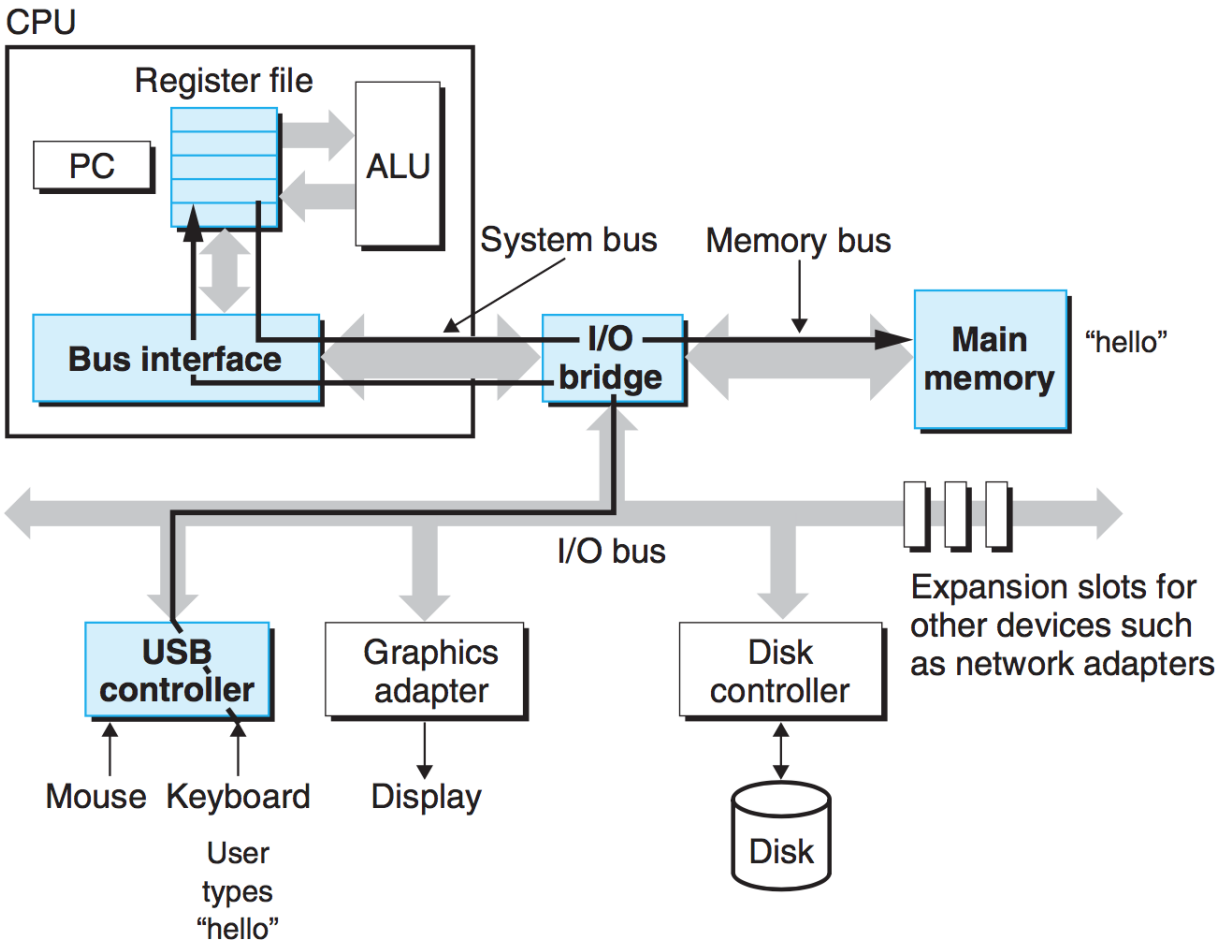
HW OVERVIEW

- Bus: words between HW components
- IO devices: keyboard, mouse, display, disk, network,
  - controller: transfer info. btwn IO bus and devices
- Main memory (not virtual memory!):
  - DRAM chips (physically)
  - linear array of bytes (logically)
  - stores: 1. machine instructions, 2. C program variables
- Processor:
  - register: word-size storage
  - register PC: where to load instruction

- Processor (continued)
  - instruction execution model
    - appear to execute in sequence (actually pipelined, out-of-order)
    - instruction read by PC
    - Turing machine
  - internal:
    - ALU: arithmetic/logic unit
    - register file: a bunch of named registers
  - ISA
    1. Load/Store: a word from main memory to a register
    2. Operate: read register content to ALU, arithmetic op on two words, store result to a register
    3. Jump: overwriting PC

- Processor (continued 2)
  - Cache
    - process-memory gap: register >100 faster than memory
    - L1 size: 10KB
    - L2 size: 1MB
    - SRAM chip
    - Locality: program to access code/data in localized regions

Figure 1.5
Reading the `hello` command from the keyboard.

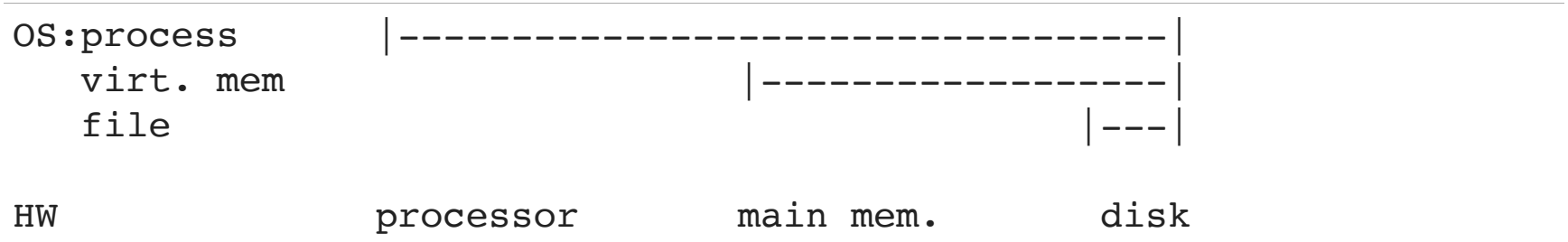


EXECUTION FROM HW PERSPECTIVE

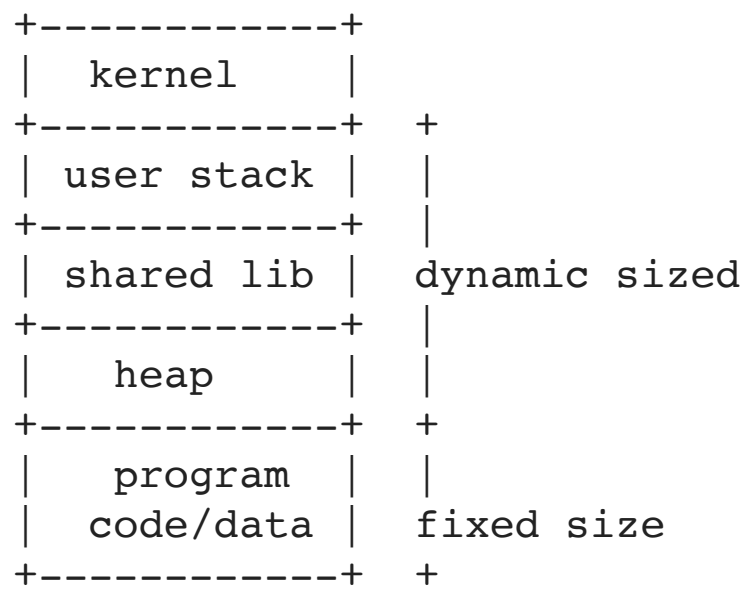
- the shell read through keyboard IO characters (`./a.out`) into register, and then store them to main memory
- the shell read character `ENTER`, starts to load `a.out`
  - load.o: load executable by copy program code/data from file to main memory.
  - DMA: disk device to memory without CPU

SYSTEMS/OS OVERVIEW

- OS goal:
  1. protecting HW from misuse by runaway app
  2. providing app easy way to manipulate HW
- fundamental abstractions: **process**, **virtual memory**, **file**



1. Process
    - look like exclusive use of HW (no interrupt, only obj in memory)
    - actually, run concurrently
        - instructions of different process interleaved
        - context switch: OS as mediator when switching from one process to another
2. Virtual memory
    - virtual address space
        1. top-most region: kernel
        2. lower region: user
3. file: sequence of bytes
    - read/write file through syscall (Unix IO)


kernel
user stack
shared lib
heap
program
code/data
dynamic sized
fixed size

EXECUTION FROM SYSTEMS PERSPECTIVE

- Shell and `./hello` are two processes, run concurrently