

- Maximal Contiguous Subsequence sum.

A. Big-Oh $O(n^3)$

$$\sum_{i=0}^{n-1} \left(\sum_{j=i}^{n-1} \left(\sum_{k=i}^j (1) + 2 \right) \right)$$

$$\begin{aligned}\sum_{k=i}^j (1) &= j - i + 1 \\ &= j - (i - 1)\end{aligned}$$

$$\sum_{i=0}^{n-1} \left(\sum_{j=i}^{n-1} (j - (i - 1) + 2) \right)$$

$$\sum_{i=0}^{n-1} \left(\sum_{j=i}^{n-1} (j - (i - 3)) \right)$$

$$\begin{aligned}\sum_{j=i}^{n-1} (j - (i - 3)) &= \sum_{j=i}^{n-1} (j) - \sum_{j=i}^{n-1} (i - 3) \\ &= (\sum_{j=0}^{n-1} (j) - \sum_{j=0}^{i-1} (j)) - (i - 3) \sum_{j=i}^{n-1} (1) \\ &= \left(\frac{n(n+1)}{2} - \frac{i(i-1)}{2} \right) + (i-3)(i-n) \\ &= \left(\frac{i^2 - in - 3i}{2} - \frac{i^2 - i}{2} \right) + 3n + \frac{n(n-1)}{2} \\ &= \frac{i(i+1)}{2} - i(n+3) + 3n + \frac{n(n-1)}{2} \\ &= \frac{i^2}{2} + \frac{i}{2} - i(n+3) + 3n + \frac{n(n-1)}{2} \\ &= \frac{i^2}{2} - i(n + \frac{5}{2}) + \frac{n(n+5)}{2}.\end{aligned}$$

$$\sum_{i=0}^{n-1} \left(\frac{i^2}{2} - i(n + \frac{5}{2}) + \frac{n(n+5)}{2} \right)$$

$$= \left(\frac{1}{2}\right) \sum_{i=0}^{n-1} (i^2) - (n + \frac{5}{2}) \sum_{i=0}^{n-1} (i) + \left(\frac{1}{2}\right) (n(n+5)) \sum_{i=0}^{n-1} (1)$$

$$\begin{aligned}&= \frac{1}{2} \left(\frac{n(n-1)(2n-1)}{6} \right) - \frac{(n + \frac{5}{2})(n)(n-1)}{2} + \frac{n^2(n+5)}{2} \\ &= \frac{n(n-1)}{2} \left(\frac{2n-1}{6} - \frac{2n+5}{2} \right) + \frac{n^2(n+5)}{2}\end{aligned}$$

$$C(n) = \frac{n(n^2 + 9n + 8)}{6} \approx \frac{n^3}{6} \in O(n^3)$$

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Because best case, average case, worst case are irrelevant b/c entire array must be traversed all runs

$$\Theta(n^3) = O(n^3)$$

B. Big-Oh $O(n^2)$

$$\begin{aligned} \cdot \sum_{i=0}^{n-1} \left(\sum_{j=i}^{n-1} (2) + 1 \right) &= \\ &= \sum_{i=0}^{n-1} (2(n-1-i+1) + 1) \\ &= \sum_{i=0}^{n-1} (2(n-i) + 1) \\ &= \sum_{i=0}^{n-1} (2n+1) - \sum_{i=0}^{n-1} (2i) \\ &= (2n+1) \sum_{i=0}^{n-1} (1) - 2 \sum_{i=0}^{n-1} (i) \\ &= n(2n+1) - 2 \frac{n(n-1)}{2} \\ C(n) &= n(n+2) \end{aligned}$$

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$$\begin{aligned} C(n) &= n(n+2) \\ &= \frac{n^2 + 2n}{n^2} \in O(n^2) \end{aligned}$$

C. Big-Oh $O(n)$

$$\begin{aligned} \cdot \sum_{i=0}^{n-1} (2) &= \\ &= \sum_{i=0}^{n-1} (2) \\ &= 2 \sum_{i=0}^{n-1} (1) \end{aligned}$$

$$C(n) = 2n$$

$$C(n) = 2n \underset{n \in O(n)}{\approx}$$

| Cubed | | |
|-------|-----------|-------------|
| N | T(n) | T(N)/N^3 |
| 125 | 881300 | 0.4512256 |
| 250 | 2051700 | 0.1313088 |
| 500 | 14755200 | 0.1180416 |
| 1000 | 115616700 | 0.1156167 |
| 2000 | 790822500 | 0.098852813 |

- Because the function is of $O(n^3)$, $T(n) = c(n^3)$.
- Based on the data acquired, we use the empirical data point $T(n) = 790822500$ ns when $n = 2000$ to find that the constant $c = T(n)/n^3 = 0.0988$.

| Square | | |
|--------|-------------|-------------|
| N | T(n) | T(N)/N^2 |
| 25000 | 822245800 | 1.31559328 |
| 50000 | 3392653600 | 1.35706144 |
| 100000 | 13090112700 | 1.30901127 |
| 200000 | 20547614100 | 0.513690353 |
| 400000 | 35042271400 | 0.219014196 |
| 800000 | 1.20169E+11 | 0.187764 |

- The value of C is converging to a positive and tiny value, meaning our function is more like $O(n)$ than $O(n^2)$

| Linear | | |
|---------|----------|-------------|
| N | T(n) | T(N)/N |
| 100000 | 619300 | 6.193 |
| 200000 | 801700 | 4.0085 |
| 400000 | 968100 | 2.42025 |
| 800000 | 1418200 | 1.77275 |
| 1600000 | 5658000 | 3.53625 |
| 3200000 | 10044700 | 3.13896875 |
| 6400000 | 24527900 | 3.832484375 |