

W202 - Homework 1

Problem 1a : What does Bob send Alice?

$e = 3$ (Bob's public key)

$N = 1081$

$m = 13$ (Alice's message)

Solution : Use RSA encryption algorithm to encrypt Alice's message using Bob's public key.

Step 1 : Calculate " C ", given values " e ", " N ", " m ".

- $\hookrightarrow C = m^e \bmod N$ (RSA encryption algorithm)
- $\hookrightarrow C = 13^3 \bmod 1081$ (substitute m, e, N variables)
- $\hookrightarrow C = 2197 \bmod 1081$ (calculate 13^3)
- $\hookrightarrow \underline{\underline{C = 35}}$ (calculate $2197 \bmod 1081$)
using WolframAlpha

Answer : Alice sends ciphertext " 35 " to Bob.

Problem 1b : Show Bob decrypting message.

$$d = 675 \quad (\text{Bob's decryption key})$$

$$N = 1081$$

$$C = 35 \quad (\text{Ciphertext from problem 1a})$$

Solution : Use RSA decryption algorithm to decrypt Alice's message with Bob's private key.

Step 1 : Calculate "m", given values "c", "d", "n".

$$\hookrightarrow M = C^d \bmod n \quad (\text{RSA decryption algorithm})$$

$$\hookrightarrow M = 35^{675} \bmod 1081 \quad (\text{substitute } c, d, n \text{ variables})$$

$$\hookrightarrow \underline{\underline{m = 13}} \quad (\text{used WolframAlpha to calculate } m)$$

Answer: Yes. When Bob decrypts ciphertext "35" from Alice, he gets "13" which is the original message from Alice.

Problem 2a : Encode "hi" using scheme.

Following the text-to-number encoding scheme...

01	02	03	04	05	06	07	08	09
↓	↓	↓	↓	↓	↓	↓	↓	↓
a	b	c	d	e	f	g	h	i

10	11	12	13	14	15	16	17	18
↓	↓	↓	↓	↓	↓	↓	↓	↓
j	k	l	m	n	o	p	q	r

19	20	21	22	23	24	25	26
↓	↓	↓	↓	↓	↓	↓	↓
s	t	u	v	w	x	y	z

Solution : Find encoding for "h" and "i".

↳ h → 08 (encoding for "h" and "i")
i → 09

↳ 08 09 (combined encoding of "hi")

↳ 809 (encoding after removing spaces)

Answer : "hi" encoded with scheme is :

↳ 0809 (combined without spaces removed)

↳ 809 (combined with spaces removed)

Problem 2b : What does Alice send Bob?

$$e = 7 \quad (\text{Bob's public key})$$

$$p = 1003001$$

$$q = 1000033$$

$$m = 809 \quad (\text{from problem 2a})$$

Solution : Use RSA encryption algorithm to encrypt Alice's message "hi" before sending to Bob.

$$C = m^e \bmod n$$

(RSA encryption algorithm)

\downarrow
plaintext

\downarrow
message

\downarrow
P · q

Step 1 : Solve for "n", given values for "p" & "q".

$$\hookrightarrow n = p \cdot q$$

$$\hookrightarrow n = (1003001)(1000033) \quad \left(\begin{array}{l} \text{calculate value } n \\ \text{given } p \text{ and } q \end{array} \right)$$

$$\hookrightarrow n = 1003034099033$$

Step 2 : Calculate "c", given values "m", "e", "n".

$$\hookrightarrow C = 809^7 \bmod 1003034099033$$

$$\hookrightarrow C = 253100088695 \quad \left(\begin{array}{l} \text{calculate value } C \\ \text{using Wolfram Alpha} \end{array} \right)$$

Answer : Alice sends " 253100088695 " as ciphertext to Bob.

Problem 2 c :

Question : What algorithm would Bob use to calculate his decryption key "d" ?

Answer : Extended Euclidean Algorithm

More specifically, referring to the RSA equation " $e \cdot d \bmod \phi(n) = 1$ ", where "d" is the decryption key, we can use the Extended Euclidean Algorithm to solve for "d". The equation can be written as :

$$\phi(n)(x) + e(y) = 1,$$

where the resulting value of (y) is the the corresponding value of decryption key "d".

An alternative equation can be derived from re-writing " $e \cdot d \bmod \phi(n) = 1$ " to " $d = e^{-1} \bmod \phi(n)$ ", which will also calculate the decryption value "d".

Problem 2d :

Part 1 : show $d = 716451497143$.

Solution : Solve for "d" using RSA equation.

$e = 7$ (Bob's public key)

$P = 1003001$

$q = 1000033$

RSA equations for "d" : $d = e^{-1} \bmod \phi(n)$
(from problem 2c)

Step 1 : Find $\phi(n)$ given values "p" and "q".

$$\hookrightarrow \phi(n) = (p-1)(q-1)$$

$$\hookrightarrow \phi(n) = (1003001-1)(1000033-1)$$

$$\hookrightarrow \phi(n) = (1003000)(1000032)$$

$$\hookrightarrow \phi(n) = 1003032096000$$

Step 2 : Use RSA equation to solve for "d".

$$\hookrightarrow e \cdot d \bmod \phi(n) = 1 \quad (\text{RSA equation})$$

$$\hookrightarrow d = e^{-1} \bmod \phi(n) \quad (\text{RSA rewritten to find "d"})$$

$$\hookrightarrow d = 7^{-1} \bmod 1003032096000 \quad (\text{calculated using})$$

$$\hookrightarrow d = 716451497143 \quad (\text{WolframAlpha})$$

Answer (part 1) : Yes. The decryption key is valid.

Problem 2d :

part 2 : Decrypt Alice's message.

Solution : Use RSA decryption algorithm.

$$d = 71645149743 \quad (\text{Bob's decryption key})$$

$$n = 1003034099033 \quad (\text{from problem 2b})$$

$$C = 253100088695 \quad (\text{from problem 2b})$$

$$m = C^d \bmod n \quad (\text{RSA decryption formula})$$

\downarrow message \downarrow ciphertext \downarrow decryption key \downarrow p.g (cert)

Step 1 : Solve for "m", given values "C", "d", "n".

$$\hookrightarrow m = 253100088695^{71645149743} \bmod 1003034099033$$

$$\hookrightarrow \underline{\underline{m = 809}} \quad (\text{Calculated using WolframAlpha})$$

Answer : Yes. Bob was able to decrypt Alice's
(part 2) message "809" using his decryption
key "d".

Problem 3 :

Explain why $e \cdot d \equiv 1 \pmod{(p-1)(q-1)}$
and not $e \cdot d \equiv 1 \pmod{pq}$

Answer :

The RSA algorithm utilizes properties of inverses in order to encrypt and decrypt a message. More specifically, the RSA equation " $m^e \pmod n$ " shows that the encryption exponent " e " and decryption exponent " d " are modular inverses of each other, derived from a special group of integers. Using the RSA equation " $m^e \pmod n$ " and the Fermat-Euler Theorem " $a^{\phi(n)} \equiv 1 \pmod n$ ", we can mathematically prove the equations :

$e \cdot d \equiv 1 \pmod{(p-1)(q-1)}$ is true and,
 $e \cdot d \equiv 1 \pmod{pq}$ is not true.

(see next page)

Proof 1 : If $m^{ed} \bmod n$, and
 $e \cdot d \equiv 1 \bmod (p-1)(q-1)$, then :

- $\hookrightarrow m^{ed} \bmod n$ (RSA equation)
- $\hookrightarrow m^{k(p-1)(q-1)+1} \bmod n$ ($1 \bmod \phi(n) \approx 1 + k\phi(n)$ property)
- $\hookrightarrow m^{k\phi(n)+1} \bmod n$ ($(p-1)(q-1) \approx \phi(n)$)
- $\hookrightarrow (m') (m^{\phi(n)})^k \bmod n$ (Distributive Property)
- $\hookrightarrow (m') (1)$ (Reduce with Euler theorem)

\approx
 m Answer: Its the original message
can be recovered, the equation
 $e \cdot d \equiv 1 \bmod (p-1)(q-1)$ is true.

Proof 2 : If $m^{ed} \bmod n$, and
 $e \cdot d \equiv 1 \bmod (pq)$, then :

- $\hookrightarrow m^{ed} \bmod n$ (RSA equation)
- $\hookrightarrow m^{k(pq)+1} \bmod n$ (Cannot factor further)

* As we cannot factor equation further, we
are unable to prove that equation
" $e \cdot d \equiv 1 \bmod (pq)$ " is true. *

Problem 4 :

Prove " m " is prime if and only if $\phi(m) = m-1$.

Proof # 1 : If m is prime, then $\phi(m) = m-1$.

↳ Euler's Totient Function $\phi(m)$ counts the number of positive integers up to " m ", that are co-prime to " m ".

↳ If " m " is a prime, then :

↳ Factors of " m " are 1 and " m ".

↳ All other positive integers are co-prime to " m ".

↳ By definition of a prime, we can express Euler's Totient Function as the equation :

$\phi(m) = m-1$, if " m " is a prime. Thus, " m " is prime, if and only if $\phi(m) = m-1$.

Proof # 2 : If $\phi(m) = m-1$, then " m " is prime.

↳ Fermat's Little Theorem states that if " m " is prime (and " a " is not divisible by " m "), then the equation : $a^{m-1} \equiv 1 \pmod{m}$ is true.

↳ By definition of Fermat's Little Theorem, " m " is prime if and only if $a^{m-1} \equiv 1 \pmod{m}$ evaluates as true.

Problem 5 : Compute values of ϕ ;

a) $\phi(6)$

Solution :

Step 1 : List all prime factors of 6.

↳ 6 has prime factors $2 \cdot 3$

Step 2 : Use Euler's Phi Function for calculation

Equation: $\phi(n) = n \prod_{p|n} \left(1 - \frac{1}{p}\right)$

↳ $\phi(6) = 6 \left(1 - \frac{1}{2}\right) \left(1 - \frac{1}{3}\right)$

↳ $\phi(6) = 6 \left(\frac{1}{2}\right) \left(\frac{2}{3}\right)$

↳ $\phi(6) = 6 \left(\frac{2}{6}\right)$

↳ $\phi(6) = 2$

Answer : $\phi(6) = 2$

check work : ① 7, ② 8, ③ 4, ④ 5, ⑤ 6

↳ common factors : 2, 3, 4, 6 (4 values)

↳ coprimes : 1, 5 (2 values) ✓

b) $\phi(9)$

Solution :

Step 1 : List all prime factors of 9.

↳ 9 has prime factors $3 \cdot 3$

Step 2 : Use Euler's Phi Function for calculation.

Equation : $\phi(n) = n \prod_{p|n} \left(1 - \frac{1}{p}\right)$

↳ $\phi(9) = 9 \left(1 - \frac{1}{3}\right)$

↳ $\phi(9) = 9 \left(\frac{2}{3}\right)$

↳ $\phi(9) = 6$

Answer : $\phi(9) = 6$

check work : ① ② ③ ④ ⑤ ⑥ ⑦ ⑧ ⑨

↳ common factors : 3, 6, 9 (3 values)

↳ coprimes : 1, 2, 4, 5, 7, 8 (6 values) ✓

c) $\phi(15)$

Solution :

Step 1 : List all prime factors of 15.

↳ 15 has prime factors $3 \cdot 5$

Step 2 : Use Euler's Phi Function for calculation

Equation : $\phi(n) = n \prod_{p|n} \left(1 - \frac{1}{p}\right)$

↳ $\phi(15) = 15 \left(1 - \frac{1}{3}\right) \left(1 - \frac{1}{5}\right)$

↳ $\phi(15) = 15 \left(\frac{2}{3}\right) \left(\frac{4}{5}\right)$

↳ $\phi(15) = 15 \left(\frac{8}{15}\right)$

↳ $\phi(15) = 8$

Answer : $\phi(15) = 8$

check work : $\textcircled{1} \textcircled{2} \textcircled{3} \textcircled{4} \textcircled{5} \textcircled{6} \textcircled{7} \textcircled{8} \textcircled{9} \textcircled{10} \textcircled{11} \textcircled{12} \textcircled{13} \textcircled{14} \textcircled{15}$

↳ common factors : 3, 5, 6, 9, 10, 12, 15 (7 values)

↳ coprime : 1, 2, 4, 7, 8, 11, 13, 14 (8 values) ✓

$$d) \phi(17)$$

Solution :

Step 1 : List all prime factors of 17.

↳ 17 has prime factors $17 \cdot 1$

Step 2 : Use Euler's Phi Function for calculation.

$$\text{Equation : } \phi(n) = n \prod_{p|n} \left(1 - \frac{1}{p}\right)$$

$$\hookrightarrow \phi(17) = 17 \left(1 - \frac{1}{17}\right)$$

$$\hookrightarrow \phi(17) = 17 \left(\frac{16}{17}\right)$$

$$\hookrightarrow \phi(17) = 16$$

$\text{Answer : } \phi(17) = 16$

check work : (1)(2)(3)(4)(5)(6)(7)(8)(9)(10)(11),
(12)(13)(14)(15)(16) ✓

↳ common factors : 17 (1 value)

↳ coprime : 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14,
15, 16 (16 values) ✓

References

Student Discussions

- Chidanand Bangalore
 - Satya srinivas
 - lyiolowa Aromokudu
 - Naphi Tang
 - Aaron Couch
 - Matthew Holmes
 - Nahid Farady
 - Lauren Ayala
- } office hours

Tools and Resources

- Hoffsterin textbook
- Async videos (module 1)
- Youtube
- Wikipedia
- Stack Exchange
- Khan Academy
- Berkeley EECS Handout
- Wolfram Alpha