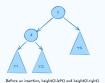
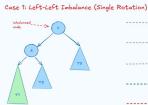
STARTING TREE for LL and LR





Then, there is an insertion somewhere in T1 that causes A's height to increase by 1. Now, height(Cleft) and height(Cright) differ by 2.

ant observations; thies in T1 are smaller than both A and C. thies in T2 are bigger than A but smaller than C thies in T3 are bigger than C.

-> Since all values in T2 are bigger than A but smaller than C, that subtree could live to the right of A or the left of C



rotateWithLeftChild(C, p A = C.left C.left = A.right A.right = C

if C == root of tree: root of tree = A else: if porentOfC.left = C: porentOfC.left = A else porentOfC.right = A

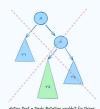


We can adjust for this imbalance with a SINSLE ROTATION by rotating the node of imbalance with its left child. Notice after the re-balancing, ${\cal A}$ has the same height as ${\bf C}$ and in the original tree.

Case 2: Left-Right Imbalance (Double Rotation)

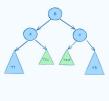








Step 1: Rotate A with Right Child B.



Step 2: Rotate C with Right Child B.

After this rotation, either T2.4 or T2.R will be as deep as T1 and T3, but not both.

T2.L foll between A and B.
could be the left child of B or the right child of A.
T2.R full between B and C.
could be the right child of B or the left child of C.

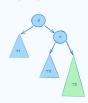
Rotate 4 with its right child (B). (same rotation RP imbalance)
Rotate C with its left child (now it is B after step 1).

STARTING TREE For RR and RL



Before an insertion, height(A.left) and height(A.right) differ by 1.

Case 4: Right Right Imbalance (Single Rotation)



Insertion into T3 causes it to be 2 levels deeper than T1, making A the node of imbalance.

Remember: All values in T2 fall between A and C. So T2 could be connected to the left of C or the Right of A.

Case 3: Right Left Imbalance (Double Rotation)



As with Case 2, an insertion into T2 caunot be solved with a single rotation. So, we need to explicitly consider the root of T2, labeled as B here.

