15-150 Fall 2015 Homework 06

Out: Wednesday, 7 October 2015 Due: Tuesday, 13 October 2015 at 23:59 EST

1 Introduction

This homework will focus on higher order functions and on continuation-passing style

1.1 Getting The Homework Assignment

The starter files for the homework assignment have been distributed through our git repository, as usual.

1.2 Submitting The Homework Assignment

Submissions will be handled through Autolab, at https://autolab.cs.cmu.edu

In preparation for submission, your hw/06 directory should contain a file named exactly hw06.pdf containing your written solutions to the homework.

To submit your solutions, run make from the hw/06 directory (that contains a code folder and a file hw06.pdf). This should produce a file hw06.tar, containing the files that should be handed in for this homework assignment. Open the Autolab web site, find the page for this assignment, and submit your hw06.tar file via the "Handin your work" link.

The Autolab handin script does some basic checks on your submission: making sure that the file names are correct; making sure that no files are missing; making sure that your code compiles cleanly. Note that the handin script is *not* a grading script—a timely submission that passes the handin script will be graded, but will not necessarily receive full credit. You can view the results of the handin script by clicking the number corresponding to the "check" section of your latest handin on the "Handin History" page. If this number is 0.0, your submission failed the check script; if it is 1.0, it passed.

Remember that your written solutions must be submitted in PDF format—we do not accept MS Word files or other formats.

All the code that you want to have graded for this assignment should be contained in hw06.sml, and must compile cleanly. If you have a function that happens to be named the

same as one of the required functions but does not have the required type, it will not be graded.

1.3 Due Date

This assignment is due on Tuesday, 13 October 2015 at 23:59 EST.

1.4 Methodology

You must use the five step methodology discussed in class for writing functions, for **every** function you write in this assignment. Recall the four step methodology:

- 1. In the first line of comments, specify the type of the function.
- 2. In the second line of comments, specify via a REQUIRES clause any assumptions about the arguments to be passed to the function.
- 3. In the third line of comments, specify via an ENSURES clause what the function computes (what it returns when applied to an argument that satisfies the assumptions in REQUIRES).
- 4. Implement the function.
- 5. Provide testcases, generally in the format val <return value> = <function> <argument value>.

For example, for the factorial function presented in lecture:

```
(* fact : int -> int
  * REQUIRES: n >= 0
  * ENSURES: fact(n) ==> n! *)
fun fact (0 : int) : int = 1
  | fact (n : int) : int = n * fact(n-1)

(* Tests: *)
val 1 = fact 0
val 6 = fact 3
val 720 = fact 6
```

1.5 Style

For this assignment, we will be grading your submissions based on your coding style. There are several ways that you can learn what is good style and what isn't:

• Your returned and graded homework submissions have been graded for style so use the markups for a reference.

- We have published solution code for the previous assignments, labs, and lectures.
- We have published a style guide at

```
https://www.cs.cmu.edu/~15150/resources/handouts/style.pdf.
```

There is also a copy in the docs subdirectory of your git clone.

• You can ask your TAs about specific examples, or post on Piazza asking general questions.

2 Polymorphic Universe

Consider the following functions:

```
fun map f [] = []
  | map f (x::L) = (f x) :: (map f) L;

fun map1 [] f = []
  | map1 (x::L) f = (f x) :: (map1 L) f;

fun map2 [] _ = []
  | map2 (x::L) f = (x f) :: (map2 L) f;

fun map3 (f, x::L) = (f x) :: map3 (f, L)
  | map3 (_, []) = [];

fun map4 (f, x::L) = (f x) :: map4 f L
  | map4 (_, []) = [];
```

Task 2.1 (4 pts). For each of the provided function definitions, what does the ML runtime system say?

Task 2.2 (2 pts). Using the typing rules, explain the reasons for the responses for map3 and map4.

Task 2.3 (4 pts). For each of the following expressions, indicate:

- If it is well typed
- The type if it is well typed
- A short informal summary of the behavior, value, or meaning of the code if it is well typed, or why it is not well typed. Note that non-termination is a behavior.

For example, for (fn (x,y) => x) is well typed, with type 'a * 'b -> 'a, and is a function that when applied, evaluates to the first element of the input tuple.

```
map (fn x => x ^{"!"}) ["Functions ", "are ", "Values"]

map (fn (x,y) => (x-1,y ^{^{^{^{^{^{^{^{^{^{^{^{}}}}}}}}}}}] [(1,"f"),("o",2),("o",3)]
```

3 Higher Order Warp Jump

Given the following definitions:

```
fun foldl f z [] = z
| foldl f z (x::L) = foldl f (f(x,z)) L;

fun foldr f z [] = z
| foldr f z (x::L) = f(x, foldr f z L);

fun flattenleft L = foldl (fn (x,A) => A@x) [] L;

fun flattenright L = foldr (fn (x,A) => A@x) [] L;
```

We want to show that flattenright and flattenleft do similar things. More specifically, they produce the same output on reversed lists.

To prove this, we will need a lemma.

Task 3.1 (13 pts).

Prove Lemma 1: for all types t1 and t2, total function f: t1 * t2 -> t2, and values z: t2, L: t1 list

You may need the following lemma:

• Lemma a: For all types t and values x: t and L,R: t list, x::(LQR) = (x::L)QR

Task 3.2 (12 pts). Prove that for all types t and values L: t list list

You may need the following lemmas:

- Lemma b: for all types t and values x: t, L: t list, rev (x::L) = (rev L)@[x] and rev [] = []
- Lemma c: @ and rev are total

Task 3.3 (8 pts). Define recursive functions

```
costleft, costright : 'a list list -> int
```

such that, for any suitably typed value L, costleft L and costright L return the number of cons operations needed to evaluate flattenleft L and flattenright L, respectively. Recall that evaluating L1@L2 uses length L1 cons operations.

Task 3.4 (2 pts).

Estimate the big-O of the number of cons operations needed to evaluate flattenleft L and flattenright L when L is a list of n lists, each of length n.

Rather than coming up with a recurrence and then solving for the closed form, use the functions from task 3.3 on different sizes of inputs.

4 Short Circuitry

It can be rather annoying that we cannot use op orelse or op andalso as the function argument for foldl or foldr. Recall that in SML, orelse uses short circuit evaluation, that is, exp2 is evaluated in exp1 orelse exp2 if and only if exp1 evaluates to false.

Task 4.1 (5 pts). Prove or disprove Theorem 1:

```
Theorem 1: For all types t1, t2 and values f : t1 -> bool, g : t2 -> bool,
    x : t1, y : t2,
    (fn (a,b) => a orelse b) (f x, g y) = f x orelse g y
```

5 Big Data

In this problem we consider the task of generating word frequencies from bodies of text.

We will generalize words to be of an arbitrary type with an associated comparison functions. You may think of them as strings, but having them be arbitrary also allows us to deal with characters, integers, etc. easily.

A histogram is a mapping from the words to their counts in a list of words. We introduce the type definition to represent histograms:

```
type 'a hist = ('a * int) list
```

If the words are of type t and with comparison function cmp, a value h of type t hist is a histogram iff it has the following properites.

- The words in h are sorted in increasing order by cmp
- Each word occurs at most once in h
- If (w,c) occurs in h, then c is positive
- If a word does not occur in the histogram, it is assumed to have count of 0

For example, for the text

["types", "are", "pretty", "useful", "types", "types", "types", "types", "types", "types"] and the standard string comparison function, the histogram will be exactly

```
[("are", 1), ("pretty", 1), ("types", 6), ("useful", 1)]
```

One method to generate a histogram over some text is to break the text into individual words, create a separate histogram for each word, and then merge them all together.

Task 5.1 (10 pts). Define the function

```
merge_histogram: ('a * 'a -> order) -> ('a hist * 'a hist) -> 'a hist such that if h1 and h2 are histograms with cmp for the comparison function, then merge_histogram cmp (h1, h2) must also be a histogram, such that if word w has count of c1 in h1 and count c2 in h2, then (w, c1+c2) occurs in the merged histogram. Notice that merge_histogram is associative and commutative.
```

Task 5.2 (10 pts). Define the function

create_histogram: ('a * 'a -> order) -> 'a list -> 'a hist

such that if w appears n times in L, then (w,n) is in create_histogram cmp L. If w doesn't appear in L, then it should not be in the histogram.

Do not write helper functions specifically for this task, other than anonymous functions to be used as arguments for higher-order functions. You should use merge_histogram

Task 5.3 (5 pts). Sometimes, the text we want to create a histogram of is too large to fit in memory. Then if we want to find the histogram for the entire text (here represented as a list of word lists), we need to generate the histogram of each sub-document and merge them together.

Define the function

corpus_histogram: ('a * 'a -> order) -> 'a list list -> 'a hist

that generates the histogram of a collection of lists. Do not flatten the list of list or write any helpers. You may use the functions from task 5.1 and 5.2.

6 The Real GridWorld

Given a 2-d grid of nodes, we want to know if there is a path, i.e. list of directions, that we can take to get to a target node from the upper left corner. Our movements are constrained thusly:

- You may only move to the node vertically down and horizontally to the right per step.
- Certain nodes are blocked, so you cannot enter them. These nodes will be labeled invalid

We will use the following datatypes:

```
datatype node = valid | invalid | target
datatype dir = down | right
```

We will represent the grids as values of type node list list, with the constraint that all of the node lists have the same length, the number of columns in the grid. The first list represents the uppermost row in the grid, and the first element of each list is in the leftmost column.

We will represent the paths as values of type dir list

Consider the following grid:

```
[[valid, valid, valid, valid],
[valid, invalid, target, valid],
[valid, valid, valid, invalid]]
```

In the grid, the only path to the target is two rights followed by a down, represented by [right, right, down]

However, in the following grid there is no path to the target because we cannot move through or around the invalid nodes blocking the target.

```
[[valid, valid, valid, valid],
[valid, valid, invalid, valid],
[valid, invalid, target, valid],
[valid, valid, valid, valid]]
```

Task 6.1 (25 pts). Define the function

such that path M s k is equivalent to s L, where L is a dir list representing a valid path to a target if such a path exists. If no such path exists, the function application is equivalent to k ().

Do not use helper functions other than continuations.

Additional information and hints:

- It may be better to think of invalid as a sinkhole where you can't move from rather than somewhere you can't move to.
- The upper left corner may be invalid. In that case, there is no path to a target.
- Think of what the recursive subproblems are.
- Think of how moving right will affect the set of potentially reachable nodes.
- [] and [[],[],[],[],[],[]] are both empty grids.
- A target may not exist in the grid. In that case there is no path to a target.
- There may be more than one target, and there may be multiple paths to the same target. If this is the case, any of the paths will suffice.
- Feel free to use the function t1, especially in conjunction with map
- Use the success continuation s to build up the result list.
- Use the failure continuation k to "go back in time" to check a different path.