

**Lecture 8:** Formal Power Series and Generating Functions

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## 1 Generating Polynomials

**Definition 1.1** (Generating polynomial). *A finite sequence  $a_0, a_1, a_2, \dots, a_n$  has the generating polynomial*

$$f(x) = a_0 + a_1x + a_2x^2 + \cdots + a_nx^n = \sum_{i=0}^n a_i x^i.$$

**Example 1.2** (Binomial coefficients). *Let  $a_k = \binom{n}{k}$  for  $k = 0, 1, \dots, n$  (note: Pascal's triangle). The corresponding generating polynomial is*

$$f(x) = \binom{n}{0} + \binom{n}{1}x + \binom{n}{2}x^2 + \cdots + \binom{n}{n}x^n = \sum_{k=0}^n \binom{n}{k} x^k.$$

**Theorem 1.3** (Binomial theorem). *For  $n \in \mathbb{N}_0$ ,*

$$\sum_{k=0}^n \binom{n}{k} x^k = (1+x)^n.$$

*In particular, this says that  $f(x) = \sum_{k=0}^n \binom{n}{k} x^k$  has a “nice” factorization.*

*Proof.* We prove the identity by induction on  $n$ .

**Base case.** For  $n = 0$ ,

$$\binom{0}{0} x^0 = 1 \cdot 1 = (1+x)^0 = 1.$$

**Inductive step.** We use the convention that  $\binom{n}{k} = 0$  when  $k < 0$  or  $k > n$ , so that

$$\sum_{k=0}^n \binom{n}{k} x^k = \sum_{k=-\infty}^{\infty} \binom{n}{k} x^k.$$

Using the identity  $\binom{n}{k} = \binom{n-1}{k-1} + \binom{n-1}{k}$ , we compute

$$\begin{aligned} \sum_k \binom{n}{k} x^k &= \sum_k \left( \binom{n-1}{k-1} + \binom{n-1}{k} \right) x^k \\ &= \sum_k \binom{n-1}{k-1} x^k + \sum_k \binom{n-1}{k} x^k \\ &= x \sum_k \binom{n-1}{k-1} x^{k-1} + \sum_k \binom{n-1}{k} x^k \\ &= x \sum_k \binom{n-1}{k} x^k + \sum_k \binom{n-1}{k} x^k \quad (\text{re-indexing}) \\ &= x(1+x)^{n-1} + (1+x)^{n-1} \\ &= (1+x)^{n-1}(1+x) \\ &= (1+x)^n, \end{aligned}$$

where we used the inductive hypothesis to replace  $\sum_k \binom{n-1}{k} x^k$  by  $(1+x)^{n-1}$ . □

*Combinatorial proof.* Expanding

$$(1+x)^n = (1+x)(1+x)\cdots(1+x) \quad (n \text{ times}),$$

the coefficient of  $x^k$  counts the number of ways to pick  $k$  of the factors from which we choose the  $x$  term (and from the remaining factors we choose 1). Hence there are  $\binom{n}{k}$  ways, so the coefficient of  $x^k$  is  $\binom{n}{k}$ .  $\square$

**Remark 1.4** (Some quick consequences). *Evaluating at  $x = 1$  gives*

$$\sum_{k=0}^n \binom{n}{k} = (1+1)^n = 2^n.$$

*Evaluating at  $x = -1$  gives*

$$\sum_{k=0}^n \binom{n}{k} (-1)^k = 0 \quad \text{for } n \geq 1.$$

*In particular,*

$$\sum_{\substack{0 \leq k \leq n \\ k \text{ even}}} \binom{n}{k} = \sum_{\substack{0 \leq k \leq n \\ k \text{ odd}}} \binom{n}{k}.$$

## 2 Generating Functions

**Definition 2.1** (Generating function). *An infinite countable sequence  $a_0, a_1, a_2, a_3, \dots$  has the generating function*

$$f(x) = a_0 + a_1x + a_2x^2 + a_3x^3 + \cdots = \sum_{k=0}^{\infty} a_k x^k.$$

*We view  $f(x)$  as a formal power series: the powers of  $x$  are placeholders, and (unlike the generating polynomial case) we do not substitute numerical values for  $x$ .*

**Example 2.2** (The all-ones sequence). *Consider  $a_0 = 1, 1, 1, 1, \dots$  (all ones). Its generating function is*

$$f(x) = \sum_{k \geq 0} x^k = 1 + x + x^2 + x^3 + \cdots.$$

*This has a “nice” representation as*

$$f(x) = \frac{1}{1-x}.$$

*Indeed,  $f(x)$  is the inverse of  $(1-x)$  since*

$$f(x)(1-x) = (1+x+x^2+x^3+\cdots)(1-x) = 1.$$

**Definition 2.3** (Coefficient extraction). *If  $A(x)$  is a generating function, we denote*

$$a_n = [x^n] A(x),$$

*to mean: extract the coefficient of  $x^n$  in  $A(x)$ .*

**Example 2.4.**

$$[x^n] \left( \sum_{k \geq 0} x^k \right) = 1 \quad \text{for all } n \geq 0.$$

**Remark 2.5.** Sometimes, a precise closed-form formula for

$$a_n = [x^n] A(x)$$

can be easily extracted from the form of  $A(x)$ .

**Remark 2.6** (Taylor series and coefficient extraction). In general, we can extract  $[x^n] A(x)$  using a Taylor series expansion of  $A(x)$  around  $x = 0$ :

$$a_n = \frac{A^{(n)}(0)}{n!}.$$

This gives the numerical value of  $a_n$  (which can be obtained computationally).

### 3 The Algebra of Formal Power Series

Before going any further, we make the formal setting explicit.

**Definition 3.1** (Formal power series over  $\mathbb{C}$ ).

$$\mathbb{C}[[x]] = \left\{ \sum_{n \geq 0} a_n x^n : a_n \in \mathbb{C} \text{ for all } n \geq 0 \right\}.$$

We refer to  $\mathbb{C}[[x]]$  as the “algebra of formal power series.” (Note:  $A(x) \in \mathbb{C}[[x]]$ .)

**Remark 3.2** (Ring operations). Let

$$A(x) = \sum_{n \geq 0} a_n x^n, \quad B(x) = \sum_{n \geq 0} b_n x^n.$$

- **Addition:**

$$A(x) + B(x) = \sum_{n \geq 0} (a_n + b_n) x^n.$$

- **Scalar multiplication:** for  $c \in \mathbb{C}$ ,

$$cA(x) = c \left( \sum_{n \geq 0} a_n x^n \right) = \sum_{n \geq 0} (c a_n) x^n.$$

- **Multiplication (convolution):**

$$A(x) \cdot B(x) = \left( \sum_{n \geq 0} a_n x^n \right) \left( \sum_{n \geq 0} b_n x^n \right) = \sum_{n \geq 0} c_n x^n =: C(x),$$

where

$$c_n = \sum_{k=0}^n a_k b_{n-k}.$$

**Remark 3.3** (Inverses and a word of caution). *Recall that an inverse means  $S(x)S^{-1}(x) = 1$  (by definition of the convolution product).*

*Some statements that are true analytically about a power series (in the sense seen in calculus) may not be true as statements in  $\mathbb{C}[[x]]$ . For example,  $\frac{1}{x}$  makes sense analytically as the inverse of  $x$  (since  $x \cdot \frac{1}{x} = 1$ ), but it does not make sense in  $\mathbb{C}[[x]]$  since the formal power series*

$$x = 0 + x + 0x^2 + 0x^3 + \dots$$

*has no inverse in  $\mathbb{C}[[x]]$ .*

*Indeed, if  $xf(x) = 1$ , then  $xf(x)$  has 0 as its constant term, but the right-hand side of  $xf(x) = 1$  is 1.*

**Fact 3.4.** *One can show that  $f(x) \in \mathbb{C}[[x]]$  has an inverse if and only if*

$$[x^0] f(x) \neq 0.$$