

1. A MIPS-like processor working at 2GHZ is available. This processor runs a program P with the following distribution of instructions.

| Tipo   | %  | CPI |
|--------|----|-----|
| load   | 25 | 1.2 |
| store  | 15 | 1   |
| add    | 40 | 1.4 |
| mult   | 10 | 1.4 |
| branch | 10 | 1.2 |

An improvement in the manufacturing process leads to an increase in the integration scales, which enables the addition of a second execution stage just after the memory access stage. This modification enriches the instruction set with arithmetic instructions having one operand in memory (`load-add`) and with combined arithmetic instructions (`add-mult`) executing two operation within a single instruction. However, adding such instructions requires a new instruction format making more complex the decoding process. Since the decoding stage is already the slowest one, this modification slows down in 5% the clock period.

Once analyzing the code generated by the compiler, it is determined that new arithmetic instructions `load-add` affect to 25% of `load` instructions. These percentage of `load` instructions and the `add` instructions following them are replaced by the new `load-add` instructions. On the other hand, `add-mult` affect to 12.5 % of `add` instructions, which are followed by a `mult` instructions. In this case the sequence of two instructions (`add` and `mult`) is replaced by the new `add-mult` instructions. Furthermore, both new arithmetic instructions, `load-add` and `add-mult`, have a CPI of 2.

- Compute the execution time of program P in the original processor. Express that time attending to the number  $n$  of executed instructions.
- Assuming the compiler uses the new instructions whenever possible, How many instructions would contain program P if executed in the new processor?
- Compute the new distribution of instructions including instructions `load-add` and `add-mult`.
- Calculate the average CPI with the modified datapath when new instructions are used.
- Justify whether the modification of the datapath and the use of the new instructions is worthy from the perspective of the execution time of program P.

### Solución:

- Compute the execution time of program P in the original processor.

$$T_{ex} = I \times CPI \times T$$

We have that  $I = n$ ,  $T = 1/2 \cdot 10^6$ . The only term missing is CPI

$$\begin{aligned}
 CPI &= 0.25 \times 1.2 + 0.15 \times 1 + 0.5 \times 1.4 + 0.1 \times 1.2 \\
 &= 0.3 + 0.15 + 0.7 + 0.12 = 1.27 \\
 T_{ex} &= n \times 1.27 \times 0.5 \text{ ns} = 0.635 \cdot n \text{ ns}
 \end{aligned}$$

- How many instructions would contain program P if executed in the new processor?  
Number of cases where a `load` instruction is followed by a `add` instruction, so they can be replaced by the new `load-add` instruction:

$$I_{load-add} = n \times 0.25 \times 0.25 = n \times 0.0625$$

Number of cases where an `add` instructions is followed by `mul` instruction, so they can be replaced by the new `mult-add`:

$$I_{add-mul} = n \times 0.4 \times 0.125 = n \times 0.05$$

In each of the aforementioned cases the number of instructions in the program is reduced in one instruction, as a result:

$$I' = n - n \times 0.0625 - n \times 0.05 = 0.8875 \cdot n$$

(c) New distribution of instructions in the new processor

$$f_{\text{load}} = \frac{0.25 - 0.25 \times 0.25}{0.8875} = \frac{0.1875}{0.8875} = 21.13\%$$

$$f_{\text{store}} = \frac{0.15}{0.8875} = 16.90\%$$

$$f_{\text{add}} = \frac{0.4 - 0.4 \times 0.125 - 0.25 \times 0.25}{0.8875} = \frac{0.2875}{0.8875} = 32.39\%$$

$$f_{\text{mult}} = \frac{0.1 - 0.4 \times 0.125}{0.8875} = \frac{0.05}{0.8875} = 5.63\%$$

$$f_{\text{branch}} = \frac{0.1}{0.8875} = 11.27\%$$

$$f_{\text{load-add}} = \frac{0.25 \times 0.25}{0.8875} = \frac{0.0625}{0.8875} = 7.04\%$$

$$f_{\text{add-mult}} = \frac{0.4 \times 0.125}{0.8875} = \frac{0.05}{0.8875} = 5.63\%$$

(d) New CPI

$$CPI' = \frac{0.1875 \times 1.2 + 0.15 \times 1 + 0.2875 \times 1.4 + 0.05 \times 1.4 + 0.1 \times 1.2 + 0.0625 \times 2 + 0.05 \times 2}{0.8875}$$

$$CPI' = \frac{0.225 + 0.15 + 0.4025 + 0.07 + 0.12 + 0.125 + 0.1}{0.8875} = \frac{1.1925}{0.8875} = 1.34$$

(e) New execution time

$$T'_{ej} = I' \times CPI' \times T'$$

$$T' = 1.05 \times 0.5 \text{ ns} = 0.525 \text{ ns}$$

$$T'_{ej} = 0.8875 \times n \times \frac{1.1925}{0.8875} \times 0.525 \text{ ns} = 1.1925 \times 0.525 \times n \text{ ns} = 0,6261 \cdot n \text{ ns}$$

$$S = \frac{T_{ej}}{T'_{ej}} = \frac{0.635 \cdot n}{0.6261 \cdot n} = 1,0142 \rightarrow 1,42\%$$

□

# Rúbricas:

| A  | B   | C  | D  |
|----|---|--|--|
| 1a | Selects the formula $t = I \times CPI \times T$ (20%).  | Computes the $CPI$ from a distribution of instructions, the CPI of each instruction, and the clock period considering the clock frequency (40%). | Applies the formula with the corresponding units (40%).  |
| 1b | Combines percentages of affected instructions (33%).  | Computes the variations in the number of instructions of each type (33%).  | Compute the total number of instructions in the new program (33%).   |
| 1c | Combines percentages of affected instructions (33%).  | Computes the variations in the number of instructions of each type (33%).  | Provides the new distribution of instructions in the program (33%).  |
| 1d | Computes the $CPI$ from a distribution of instructions and the CPI of each instructions (100%). |  |  |
| 1e | Selects the formula $t = I \times CPI \times T$ (20%).  | Computes the new clock period (20%).   | Applies the formula with the corresponding units (30%).  |
|    |   |  | Justifies correctly whether the modification is worthy and supports the decision by computing the speed up and interpreting it, or by comparing the resulting execution times (30%). |