Performance portability via Nim metaprogramming

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A DOE ECP effort from Lattice QCD

August 23 DOE COE Performance Portability Meeting 2017 Denver, Colorado "Complexity" seems to be a lot like "energy": you can transfer it from the end user to one/some of the other players, but the total amount seems to remain pretty much constant for a given task.

-Ran, 5 Mar 2000

Outline

- Benchmark code in Nim
- Application domain: Lattice Gauge Theory
- Nim programming language
- onGPU
- Benchmark result
 - P100 + XEON
 - KNL

Benchmark

$X_{ij}^l \leftarrow X_{ij}^l + Y_{ik}^l \times Z_{kj}^l$

```
3: proc test(vecLen, memLen: static[int]; N: int) =
4:
     var
       x = newColorMatrixArray(vecLen,memLen,N)
5:
                                                 Define complex
       y = newColorMatrixArray(vecLen, memLen, N)
6:
                                                 3\times3 matrix field
       z = newColorMatrixArray(vecLen, memLen, N)
7:
                 CPU threads
     threads:
32:
       x := 0
33:
                  Set diagonal elements
       y := 1
34:
       z := 2
35:
       timeit "CPU": x += y * z Benchmark
36:
37:
     timeit "GPU5":
38:
                     Run statement block on GPU
       onGpu(N, 32):
39:
         x += y * z
40:
                                                   Rerun with
     timeit "GPU6": onGpu(N, 64): x += y * z
41:
     timeit "GPU7": onGpu(N, 128): x += y * z
42:
                                                  different T/B
43:
     threads: timeit "CPU": x += y * z
44:
                                          Back to CPU threads
```

CPU threads

- Takes a block of code statements
- Wraps in a function with lexically scoped thread local objects
- References to variables outside the code block are managed by Nim
- · Runs the function under omp parallel directive
- A custom iterator over the array indices takes care of actual data parallel operations

$$x += y * z \longrightarrow x[i] += (y * z)[i]$$

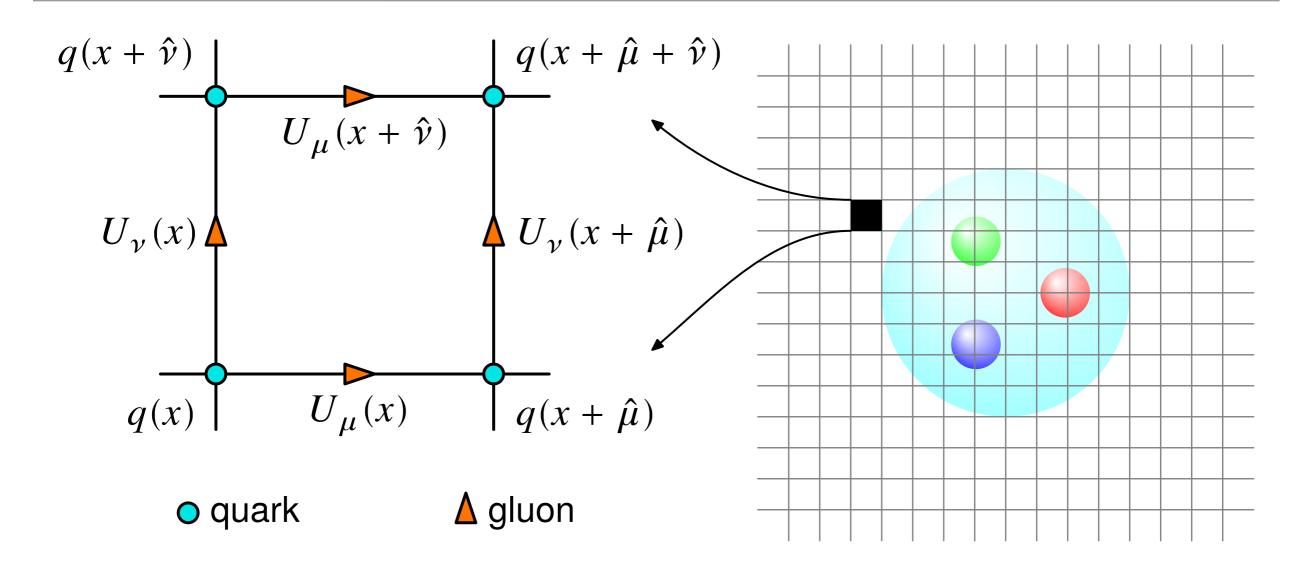
AST based overloading for data parallel ops

```
1: type ArrayIndex* = SomeInteger or ShortVectorIndex
2:
3: template indexArray*(x: ArrayObj, i: ArrayIndex): untyped =
     x.p[i]
4:
5:
6: macro indexArray*(x: ArrayObj{call}, y: ArrayIndex): untyped =
     result = newCall(ident(x[0])
8: for i in 1..<x.len:
9: let xi = x[i]
10: result.add( quote do:
         indexArray(`xi`, `y`) )
11:
12:
13: template `[]`*(x: ArrayObj, i: ArrayIndex): untyped =
     indexArray(x, i)
14:
```

• When an ArrayObj is indexed, if the object is a function call, the indexing goes inside the call

$$x[i] += (y * z)[i] \longrightarrow x[i] += y[i] * z[i]$$

Lattice gauge theory



 Large 4D (5D) grid of small vectors/matrices with homogeneous stencil operations — large sparse linear algebra

Nim

- Modern (since 2008) language
- "Efficient Expressive Elegant"
- · Statically typed systems language (full access to low-level objects & code) with type inference
- Generates C or C++ code & compile with any compiler
- · Integrated build system (no Makefile necessary): copy main program, modify, compile
- https://nim-lang.org

Nim—both low-level and high-level

- Low-level efficiency
 - Can manually manage memory instead of GC
 - Cross module inlining and constant unfolding
 - Whole program dead code elimination
- High-level wrappers & libraries
 - gmp, bignum, nimblas, linalg(LAPACK), ...
 - bindings to GTK2, the Windows API, the POSIX API, OpenGL, SDL, Cairo, Python, Lua, TCL, X11, libzip, PCRE, libcurl, mySQL, SQLite, ...
 - exportC to create static/dynamic libraries
- NimScript: shell-like scripting
 - · Used in compiler for compile-time evaluation
 - · Available to plug in to application and can interface with rest of application

Nim—metaprogramming

- Templates: in-line code substitutions, also allows overloading, completely hygienic (if desired)
- Generics: applies to types, procedures, templates, and macros also allows type-classes, concepts
- Macros: similar to Lisp: syntax tree of arguments passed to macro at compile time (type checked or untyped)
- **AST based overloading**: allows specialization based on the AST of the arguments

New framework: QEX (Quantum EXpressions)

- Data parallel library for tensor objects on a lattice including shifts, reductions
- Mostly in Nim, with USQCD SciDAC C libraries
- High level interface in development
- Available on https://github.com/jcosborn/qex
- Performance portability study: cudanim
 - Supports arrays on both CPU and GPU
 - Checkout https://github.com/jcosborn/cudanim

onGpu

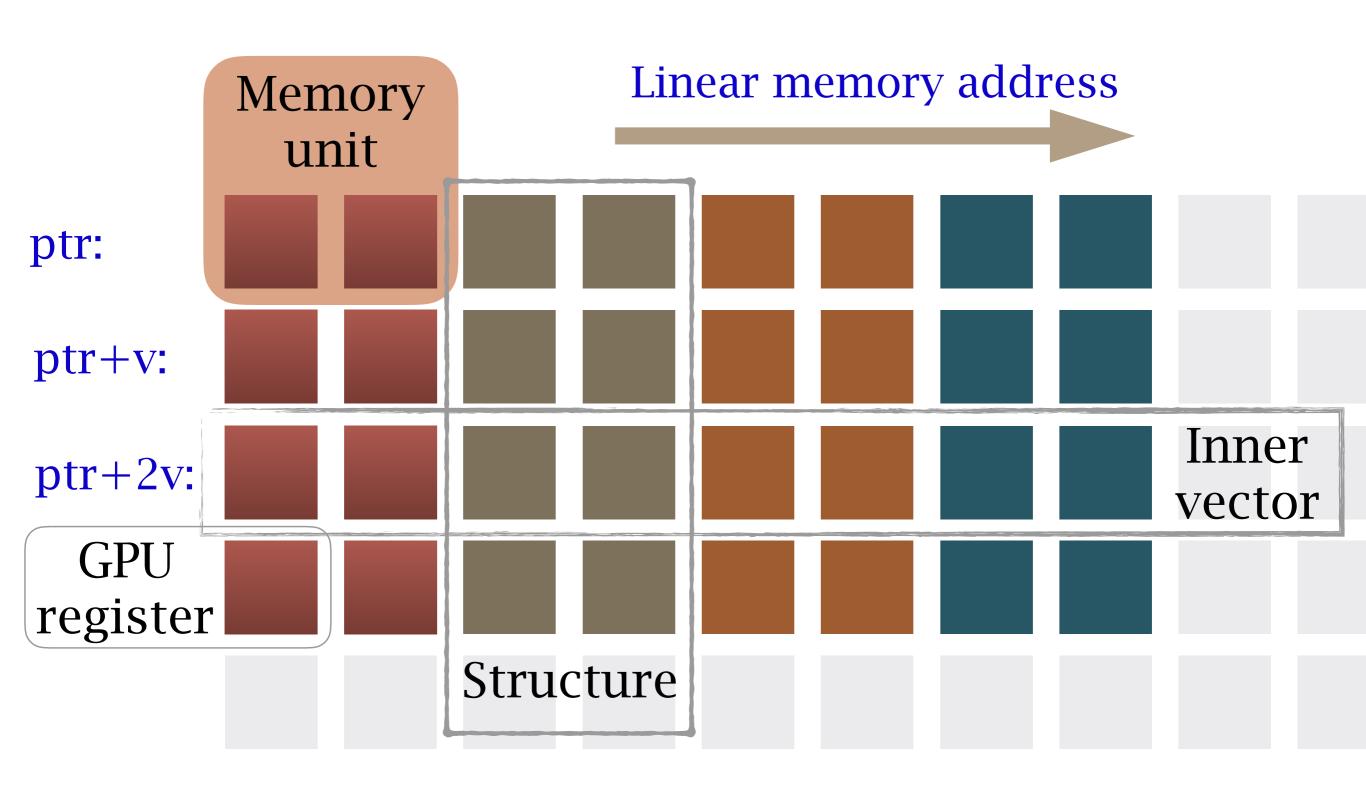
```
1: template onGpu*(nn,tpb: untyped, body: untyped): untyped =
                                             Copy memory to GPU
     block:
2:
       var v = packVars(body, getGpuPtr)
3:
                                                  if necessary
       type ByCopy {.bycopy.} [T] = object
4:
                                              collect pointers in v
 5:
         d: T
       proc kern(xx: ByCopy[type(v)]) {.cudaGlobal.} =
6:
         template deref(k: int): untyped = xx.d[k]
7:
         substVars(body, deref)
8:
                                        Generate kernel function
       let ni = nn.int32
9:
       let threadsPerBlock = tpb.int32
10:
       let blocksPerGrid = (ni+threadsPerBlock-1) div threadsPerBlock
11:
       cudaLaunch(kern, blocksPerGrid, threadsPerBlock, v)
12:
       discard cudaDeviceSynchronize()
13:
14: template onGpu*(nn: untyped, body: untyped): untyped =
     onGpu(nn, 64, body)
15:
16: template onGpu*(body: untyped): untyped =
     onGpu(512*64, 64, body)
17:
```

The generated kern

- xx is an object holding pointers to GPU memory
- substVars with the help of deref transforms

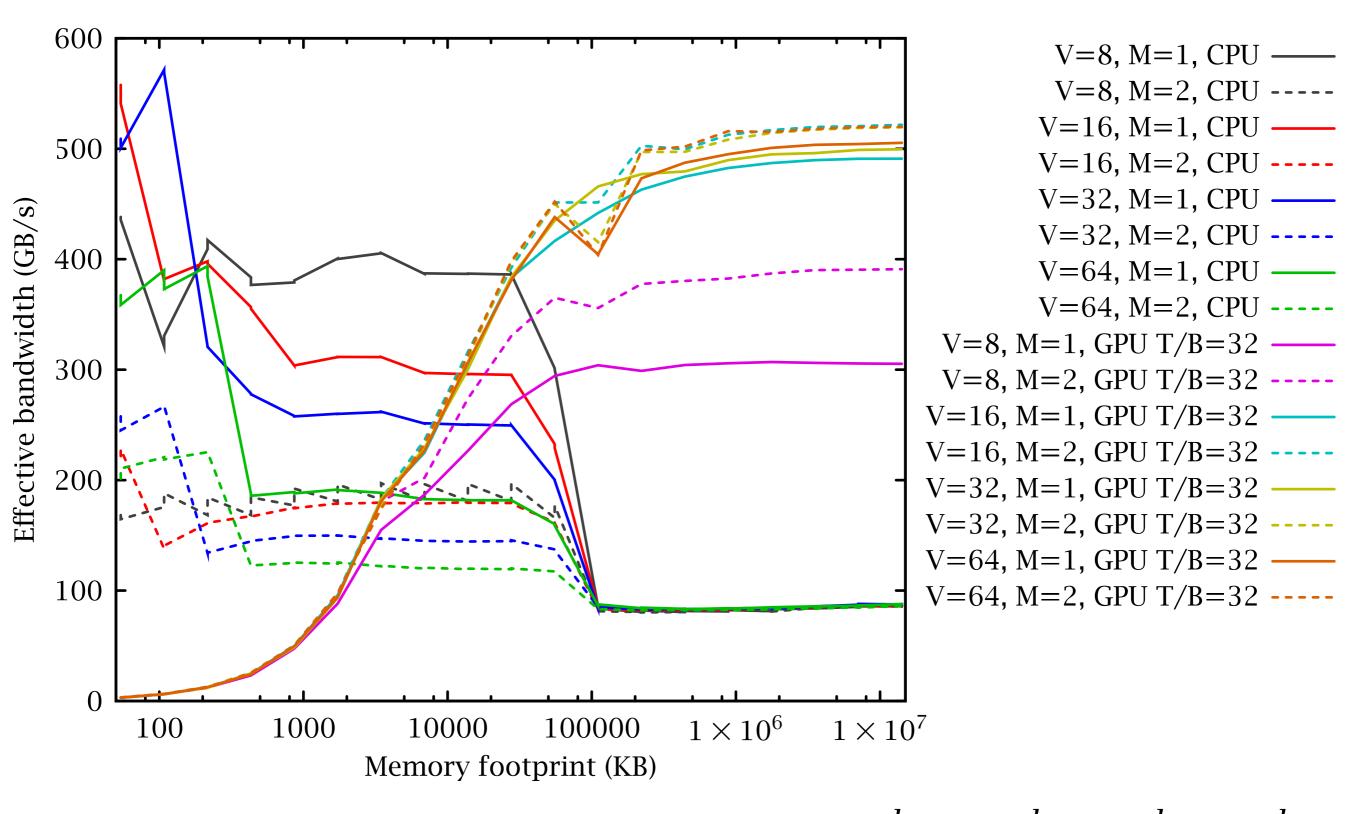
$$x += y * z \longrightarrow deref(0) += deref(1)*deref(2)$$

$$\longrightarrow xx.d[0] += xx.d[1]*xx.d[2]$$



Coalesced in-memory data layout

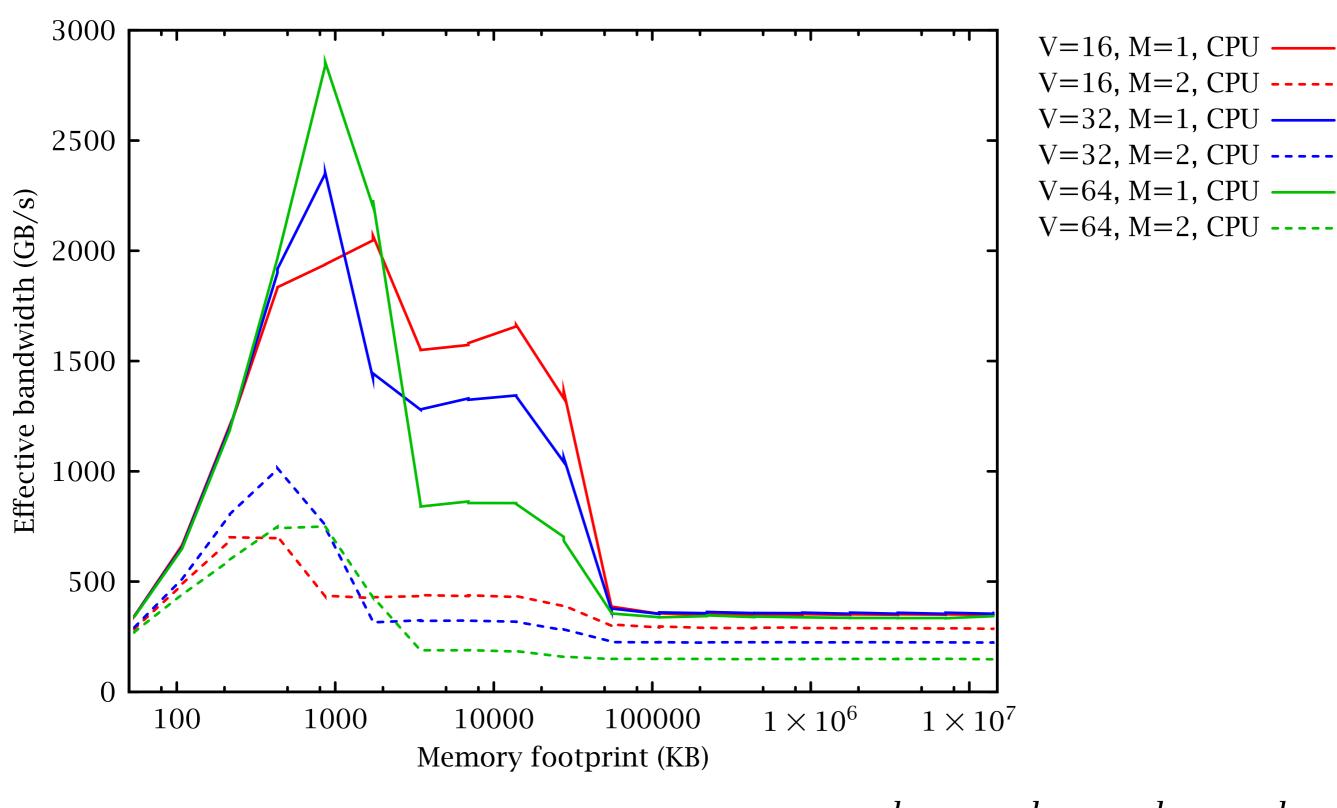
AoSoAoS



Tesla P100 + 2x Xeon E5-2687WV2

$$X_{ij}^l \leftarrow X_{ij}^l + Y_{ik}^l \times Z_{kj}^l$$

CUDA 8.0.61 + GCC 4.8.5



KNL (Xeon Phi 7210) Flat MCDRAM

$$X_{ij}^{l} \leftarrow X_{ij}^{l} + Y_{ik}^{l} \times Z_{kj}^{l}$$
GCC 7.1.0

Summary & Outlook

- Nim metaprogramming helps hiding architecture differences under a unified data parallel API
- Toy benchmark saturates GPU bandwidth
- Considering API to use both CPU & GPU in a heterogeneous setting
- Many possibilities of using AST
 - Apply AST based optimizations (inlining, loop unrolling, temporary variable elimination), across multiple statements
 - Specialization at different levels
 - Craft application specific AST transformations
 - Help general purpose compiler with application specifics

Backup slides

GPU offloading (1 of 4)

```
1: template cudaDefs(body: untyped): untyped {.dirty.} =
     var gridDim{.global,importC,noDecl.}: CudaDim3
     var blockIdx{.global,importC,noDecl.}: CudaDim3
 3:
     var blockDim{.global,importC,noDecl.}: CudaDim3
4:
     var threadIdx{.global,importC,noDecl.}: CudaDim3
 5:
     template getGridDim: untyped {.used.} = gridDim
6:
     template getBlockIdx: untyped {.used.} = blockIdx
7:
     template getBlockDim: untyped {.used.} = blockDim
8:
     template getThreadIdx: untyped {.used.} = threadIdx
9:
     template getThreadNum: untyped {.used.} = blockDim.x * blockIdx.x + threadIdx.x
10:
     template getNumThreads: untyped {.used.} = gridDim.x * blockDim.x
11:
     bind inlineProcs
12:
     inlineProcs:
13:
       body
14:
```

- Overloaded template definitions
- Inline all Nim procedures in body to prepare for generating kernel function

GPU offloading (2 of 4)

```
34: macro cuda*(s,p: untyped): auto =
35:    let ss = s.strVal
36:    p.expectKind nnkProcDef
37:    result = p
38:    result.addPragma parseExpr("{.codegenDecl:\""&ss&" $# $#$\".}")[0]
39:    result.body = getAst(cudaDefs(result.body))
40:    var sl = newStmtList()
41:    sl.add( quote do:
42:    {.push checks: off.}
43:    {.push stacktrace: off.} )
44:    sl.add result
45:    result = sl
46: template cudaGlobal*(p: untyped): auto = cuda("__global__",p)
```

- Convert a procedure definition, p, with the overloaded templates in cudaDefs
- Turn it in to a ___global___ kernel

GPU offloading (3 of 4)

```
16: template cudaLaunch*(p: proc; blocksPerGrid,threadsPerBlock: SomeInteger;
                         arg: varargs[pointer,dataAddr]) =
17:
     var pp: proc = p
18:
     var gridDim, blockDim: CudaDim3
19:
     gridDim.x = blocksPerGrid
20:
     gridDim.y = 1
21:
     qridDim.z = 1
22:
     blockDim.x = threadsPerBlock
23:
     blockDim.y = 1
24:
     blockDim.z = 1
25:
    var args: array[arg.len, pointer]
26:
    for i in 0...<arg.len: args[i] = arg[i]
27:
    #echo "really launching kernel"
28:
     let err = cudaLaunchKernel(pp, gridDim, blockDim, addr args[0])
29:
     if err:
30:
       echo err
31:
       quit cast[cint](err)
32:
```

 Calls the cuda function, cudaLaunchKernel, with a passed in procedure, p

GPU offloading (4 of 4)

```
48: template onGpu*(nn,tpb: untyped, body: untyped): untyped =
     block:
49:
       var v = packVars(body, getGpuPtr)
50:
       type ByCopy {.bycopy.} [T] = object
51:
52:
       proc kern(xx: ByCopy[type(v)]) {.cudaGlobal.} =
53:
         template deref(k: int): untyped = xx.d[k]
54:
         substVars(body, deref)
55:
       let ni = nn.int32
56:
       let threadsPerBlock = tpb.int32
57:
       let blocksPerGrid = (ni+threadsPerBlock-1) div threadsPerBlock
58:
       cudaLaunch(kern, blocksPerGrid, threadsPerBlock, v)
59:
       discard cudaDeviceSynchronize()
60:
61: template onGpu*(nn: untyped, body: untyped): untyped = onGpu(nn, 64, body)
62: template onGpu*(body: untyped): untyped = onGpu(512*64, 64, body)
```

- Take a body of code chunk and put it in a kernel definition, kern
- kern calls cudaGlobal to setup other definitions, and takes care of syncing CPU memory to GPU