

New ideas for scalable attractor detection and control of asynchronous Boolean networks

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Attractor detection and control of asynchronous Boolean networks

Two **most important** issues in Boolean networks research with many applications in various fields [Akutsu, 2018, Schwab et al., 2020].

Usually, the attractor detection issue is the preceding step of the control issue.

Unfortunately, both of them are **very challenging**, especially for **large-scale** networks.

Many methods have been proposed.

Attractor detection methods

BDD-based methods: genYsis [Garg et al., 2008],
geneFatt [Zheng et al., 2013],
[CABEAN](#) [Mizera et al., 2019, Su and Pang, 2021a].

Reduction-based methods: [AEON](#) [Benes et al., 2021].

Trap space-based methods: PyBoolNet¹ [Klarner et al., 2017],
[pystablemotifs](#) [Rozum et al., 2021b, Rozum et al., 2021a].

Feedback vertex set-based methods: FVS-ABN [Trinh et al., 2020],
iFVS-ABN [Giang and Hiraishi, 2021], [mtsNFVS](#) [Trinh et al., 2022b].

All of them have [advantages](#) and [disadvantages](#). To our best knowledge,
none of them can robustly handle [complex networks with thousands of nodes](#).

¹It is not an exact method.

Control methods

BDD-based and decomposition-based methods:

[CABEAN](#) [Su and Pang, 2021b, Su and Pang, 2021a].

Reduction-based methods: [AEON](#) [Brim et al., 2021].

Trap space-based methods:

PyBoolNet [Fontanals et al., 2020, Cifuentes-Fontanals et al., 2021],
[pystablemotifs](#) [Zañudo and Albert, 2015, Rozum et al., 2021a].

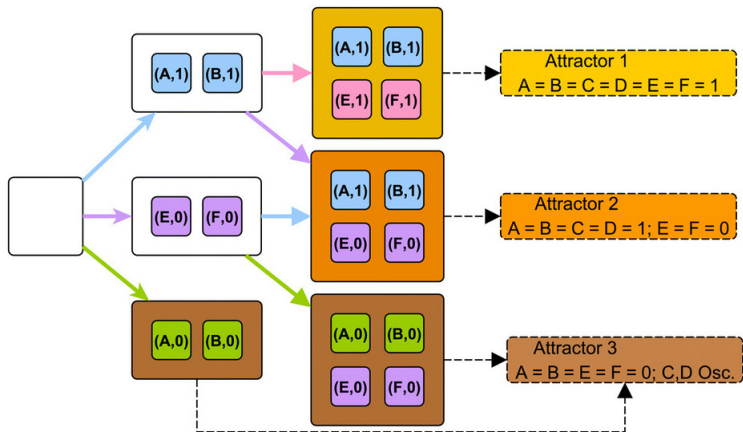
Objective

pystablemotifs and mtsNFVS have **much more potential** to reach the **genome scale**.

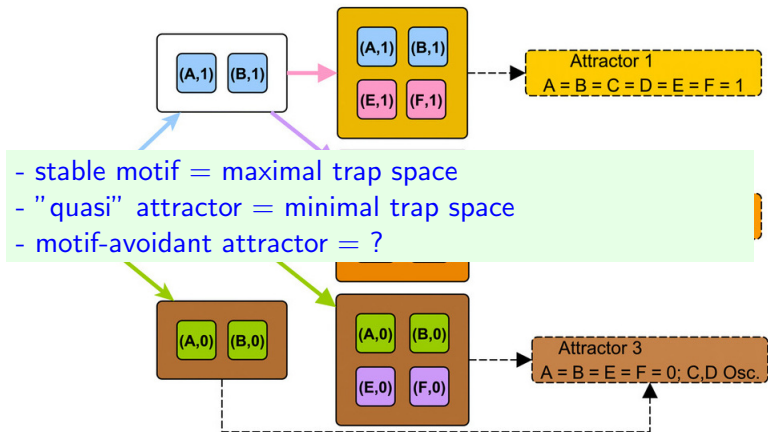
Both of them have their own advantages and disadvantages, but fortunately they can be **complementary** to each other.

Develop a **more efficient method** that can benefit both the attractor detection and control issues, and in particular can handle networks at the genome scale.

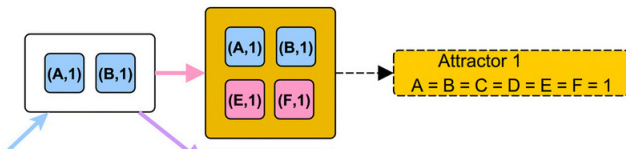
E Succession diagram and attractors



E Succession diagram and attractors

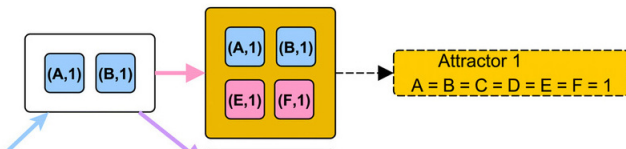


E Succession diagram and attractors

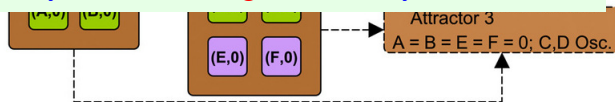


- pystablemotifs is beneficial to not only the attractor detection problem but also the control problem as it provides a **succession diagram** that can be efficiently used for control.
- If solely considering attractor identification, pystablemotifs seems to have **disadvantages** as compared to other methods such as AEON, iFVS-ABN, and mtsNFVS.

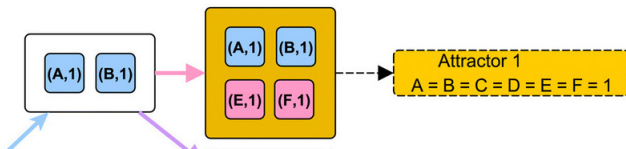
E Succession diagram and attractors



- pystablemotifs must check the existence of **motif-avoidant attractors**.
- Although it uses time reversal to prune the considered state space, the remaining part (called the terminal restriction space) may be still **too large** to be analyzed.

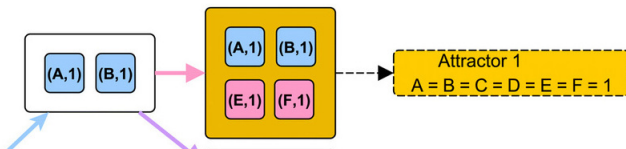


E Succession diagram and attractors

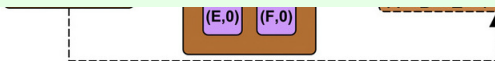


- pystablemotifs relies on PyBoolNet to compute stable motifs (i.e., maximal trap spaces). However, PyBoolNet does not work well with networks with Boolean functions having many input nodes.
- The reason for this is PyBoolNet mainly relies on prime-implicants computation, whereas the number of prime-implicants may be too huge if there are many input nodes.

E Succession diagram and attractors



- The size of the succession diagram may be too large.
- Apart from the theoretical worst case ($n_{sm}!$ where n_{sm} is the number of stable motifs), the number of stable motifs may be actually too large because for example, the network has too many source nodes (node A is a source node if $f_A = A$).

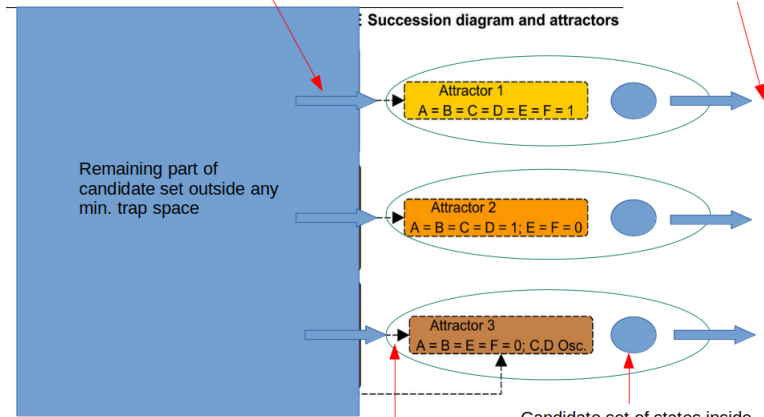


mtsNFVS

Preprocessing SSF (shrink the candidate set)

After Preprocessing SSF, if there is only one state, we do not need the reachability check. In this case, there is only one attractor inside this min. trap space.

Succession diagram and attractors



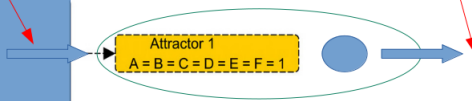
After Preprocessing SSF, if there is no state in the remaining part, we do not need the reachability check. In this case, there is no attractor outside any min. trap space.

mtsNFVS

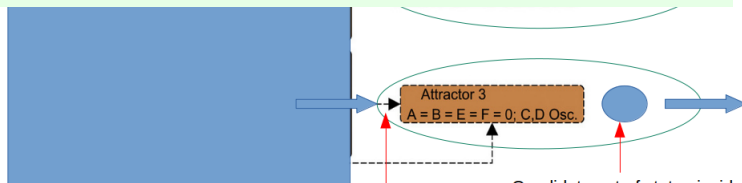
Preprocessing SSF (shrink the candidate set)

After Preprocessing SSF, if there is only one state, we do not need the reachability check. In this case, there is only one attractor inside this min. trap space.

Succession diagram and attractors



- The candidate set F (the blue part) is computed based on an negative feedback vertex set U^- of the ABN.



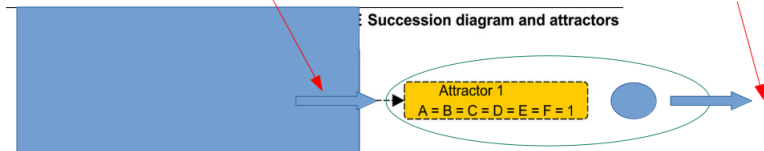
After Preprocessing SSF, if there is no state in the remaining part, we do not need the reachability check. In this case, there is no attractor outside any min. trap space.

mtsNFVS

Preprocessing SSF (shrink the candidate set)

After Preprocessing SSF, if there is only one state, we do not need the reachability check. In this case, there is only one attractor inside this min. trap space.

Succession diagram and attractors



- If the size of the negative feedback vertex set is too large, there may be too many states in the candidate set F (i.e., the set of fixed points of the reduced STG), leading to extremely longer time for the computation of F , Preprocessing SSF, and maybe the reachability analysis.
- Note that real-world models usually have **small** minimum negative feedback vertex sets.

After Preprocessing SSF, if there is no state in the remaining part, we do not need the reachability check. In this case, there is no attractor outside any min. trap space.

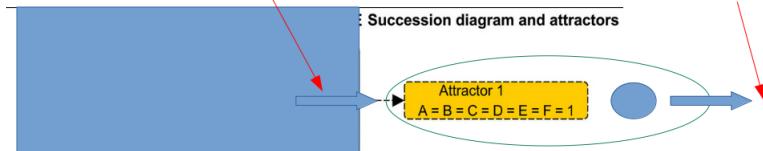
Candidate set of states inside each min. trap space

mtsNFVS

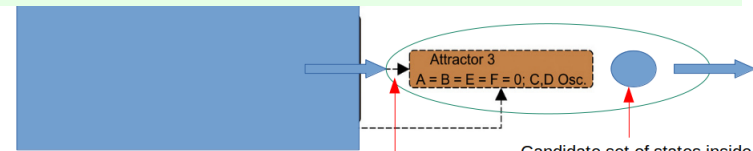
Preprocessing SSF (shrink the candidate set)

After Preprocessing SSF, if there is only one state, we do not need the reachability check. In this case, there is only one attractor inside this min. trap space.

Succession diagram and attractors



- mtsNFVS also uses PyBoolNet to compute the candidate set as well as the set of minimal trap spaces. Again, it will not work well with the case of many input nodes.



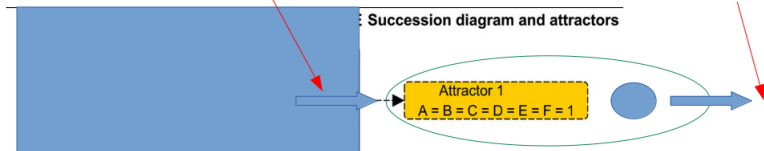
After Preprocessing SSF, if there is no state in the remaining part, we do not need the reachability check. In this case, there is no attractor outside any min. trap space.

mtsNFVS

Preprocessing SSF (shrink the candidate set)

After Preprocessing SSF, if there is only one state, we do not need the reachability check. In this case, there is only one attractor inside this min. trap space.

Succession diagram and attractors



- If some motif-avoidant attractors exist or the results of Preprocessing SSF are not good enough, mtsNFVS still must check the reachability in asynchronous BNs, which is PSPACE-complete in general.
- Moreover, the target set for the reachability analysis is maybe small (i.e., the union of all minimal trap spaces).

After Preprocessing SSF, if there is no state in the remaining part, we do not need the reachability check. In this case, there is no attractor outside any min. trap space.

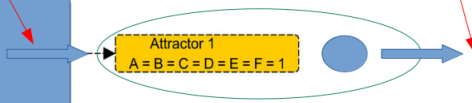
Candidate set of states inside each min. trap space

mtsNFVS

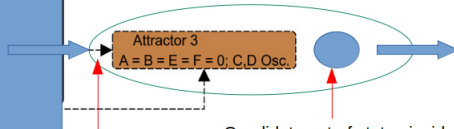
Preprocessing SSF (shrink the candidate set)

After Preprocessing SSF, if there is only one state, we do not need the reachability check. In this case, there is only one attractor inside this min. trap space.

Succession diagram and attractors



- Currently, mtsNFVS focuses solely on the attractor detection problem.



After Preprocessing SSF, if there is no state in the remaining part, we do not need the reachability check. In this case, there is no attractor outside any min. trap space.

Candidate set of states inside each min. trap space

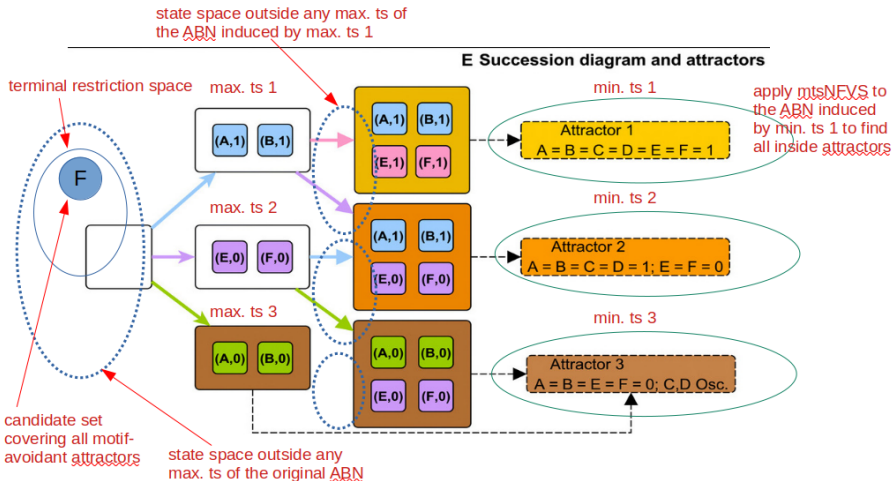
New idea

In principle, our new idea is to combine pystablemotifs and mtsNFVS to obtain a more efficient method that can benefit both the attractor detection and control problems.

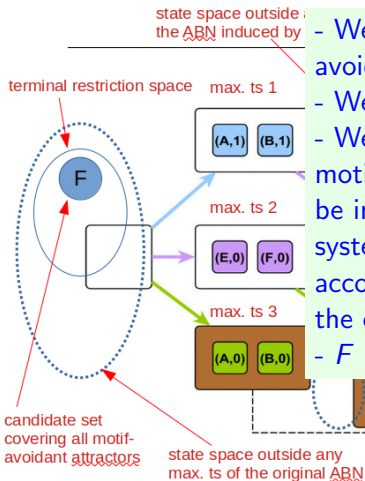
Hereafter, we will present the general approach along with many **new findings** that can speedup the whole process.

General approach

It follows the process of *pystablemotifs* (i.e., building the succession diagram).

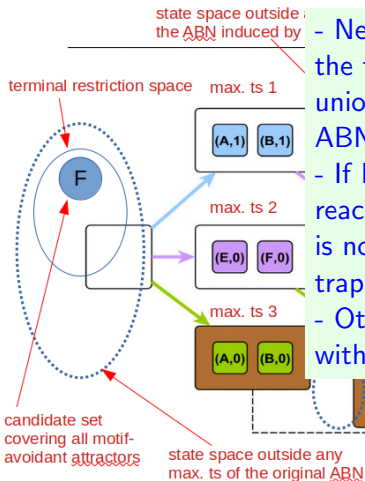


General approach



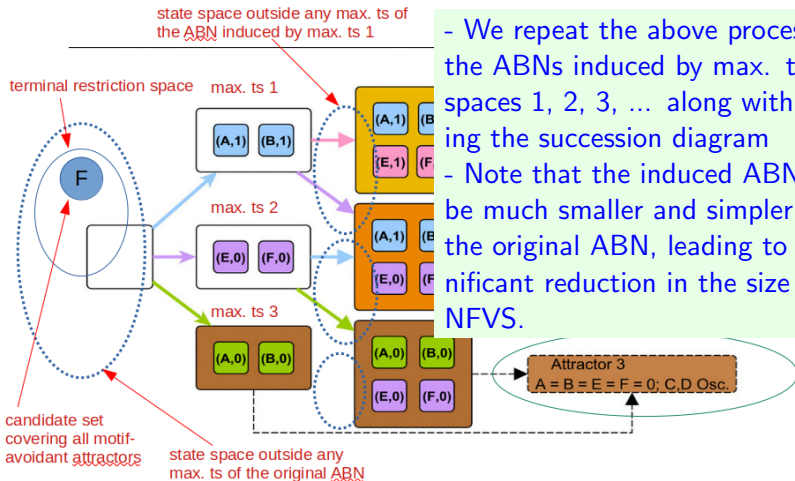
- We need to check the existence of motif-avoidant attractors.
- We compute the terminal restriction space R .
- We compute the candidate set F covering all motif-avoidant attractors. Each state in F must be in R and has no out-going successor after systematically removing some of its successors according to the negative feedback vertex set of the original ABN.
- F may be **small** even when the NFVS is large.

General approach



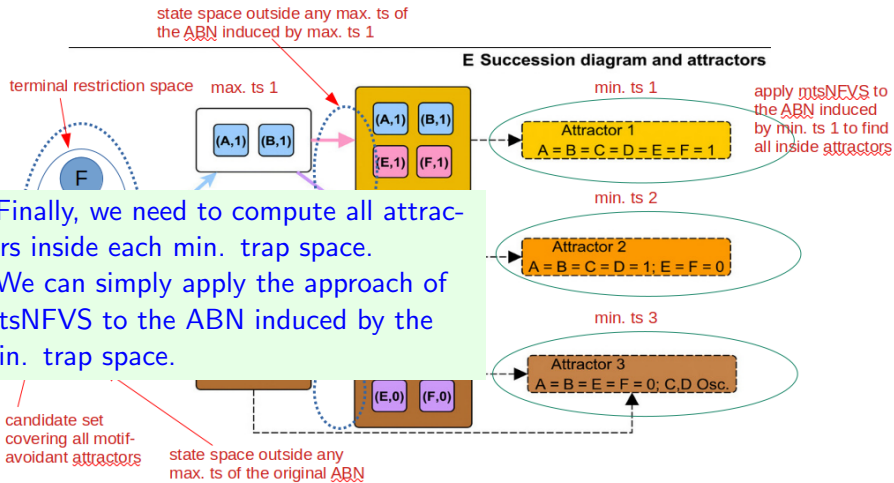
- Next, we use Preprocessing SSF for F with the target set M_{max} of states covered by the union of all max. trap spaces of the original ABN.
- If Preprocessing SSF is good enough, we can reach the **best** case where $F = \emptyset$ (i.e., there is no motif-avoidant attractor outside the max. trap spaces).
- Otherwise, we need to check the reachability with the target set including M_{max} .

General approach



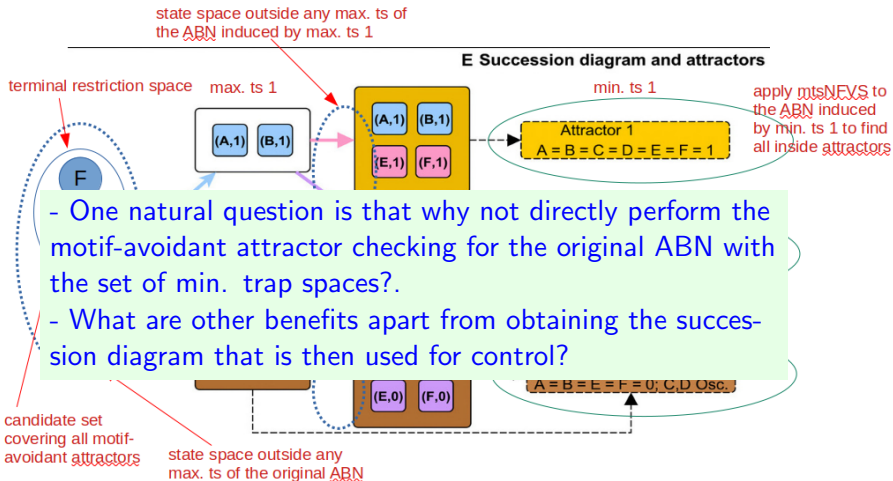
- We repeat the above process for the ABNs induced by max. trap spaces 1, 2, 3, ... along with building the succession diagram
- Note that the induced ABNs may be much smaller and simpler than the original ABN, leading to the significant reduction in the size of the NFVS.

General approach

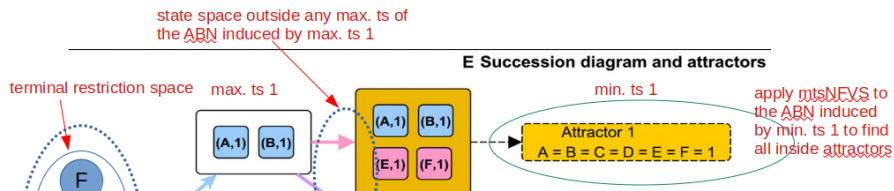


- Finally, we need to compute all attractors inside each min. trap space.
- We can simply apply the approach of mtsNFVS to the ABN induced by the min. trap space.

General approach

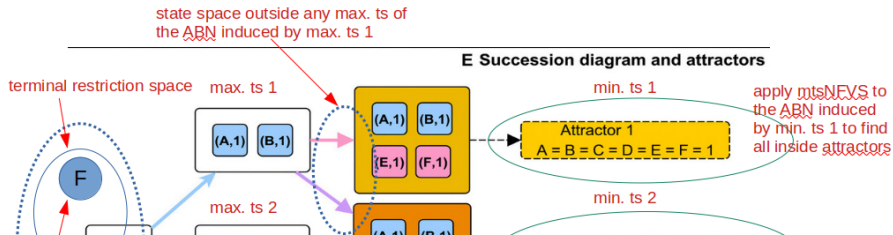


General approach



- First, the new approach of course obtains not only the set of attractors but also the succession diagram that is then used for control.
- Second, we observed that in this case the terminal restriction space may be **too huge**, leading to the candidate set F may be too large (hence, hard to be computed as well).
- Third, M_{max} is much broader than M_{min} in most cases, hence using max. trap spaces can help Preprocessing SSF to reach good cases **more easily** than using min. trap spaces.

General approach



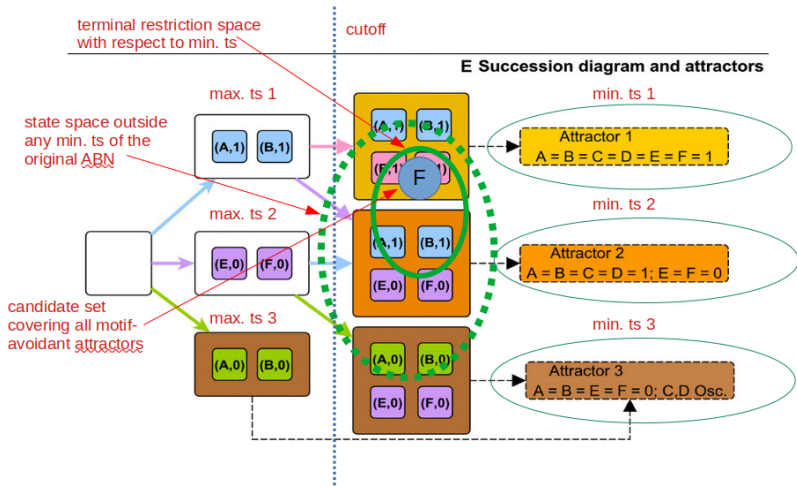
Jordan: We need to consider the percolation of a stable motif (LDOI) along with building the succession diagram. The percolation is based on PyBoolNet.

Sam: When we have already obtained the stable motif (i.e., max. trap space), we can perform the percolation easily by algebraically propagating the fixed values.

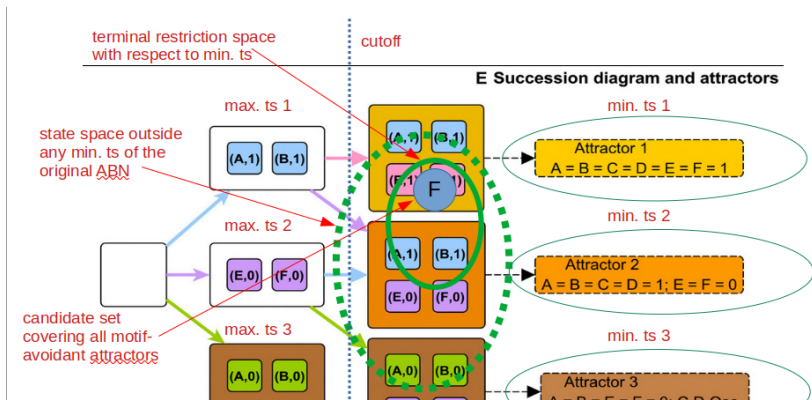
Giang: Sure. It is possible. Note that we can store Boolean functions as Boolean expressions or BDDs (not as prime-implicants).

Deal with the problem of large succession diagrams

The above approach is promising. However, how to deal with the problem of large succession diagrams?

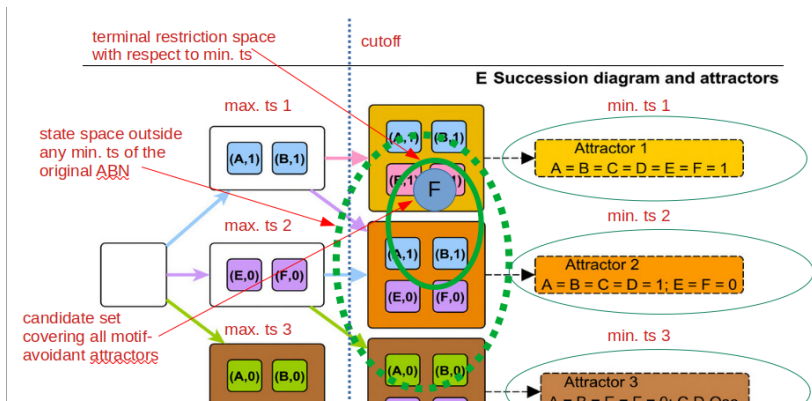


Deal with the problem of large succession diagrams



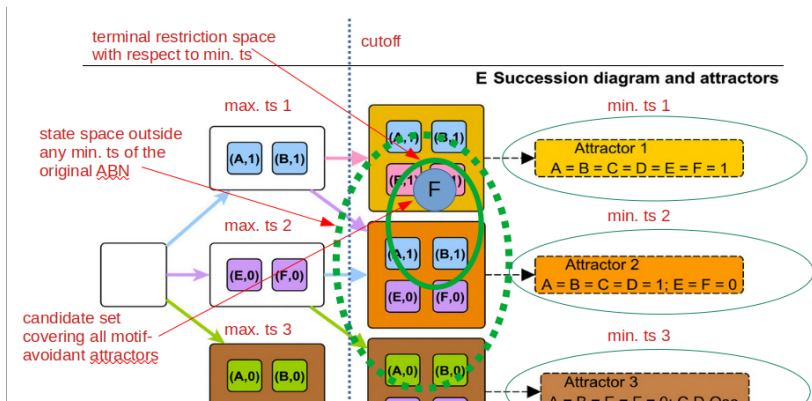
- When the size of the succession diagram becomes too large (maybe exceeding a threshold), we stop its construction.
- However, we still need to check the existence of motif-avoidant attractors.

Deal with the problem of large succession diagrams



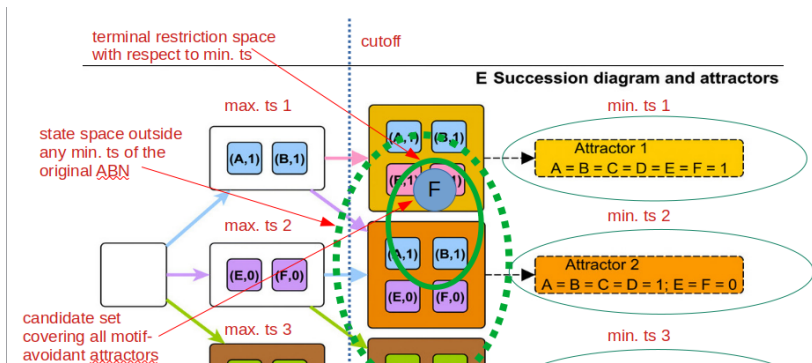
- First, we compute all min. trap spaces of the original ABN.
- Second, we compute the terminal restriction space R with respect to these min. trap spaces.
- Next, we get the intersection between R and the state spaces induced by the already computed max. trap spaces (e.g., max. ts 1, 2, 3).

Deal with the problem of large succession diagrams



- Then, we compute the candidate set F with respect to the NFVS of the original ABN.
- Each state in F must be in the intersected set obtained beforehand.
- F may be **small** even when the NFVS is large because the intersected set may be not too large.

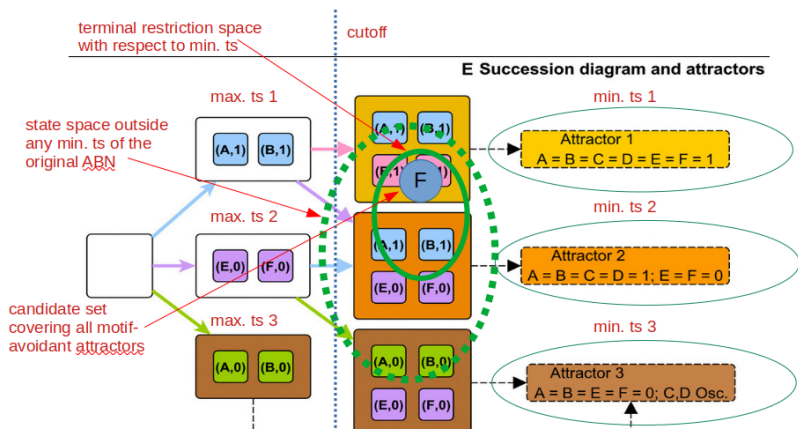
Deal with the problem of large succession diagrams



Next, we use Preprocessing SSF for F with the target set M_{min} of states covered by the union of all min. trap spaces of the original ABN.

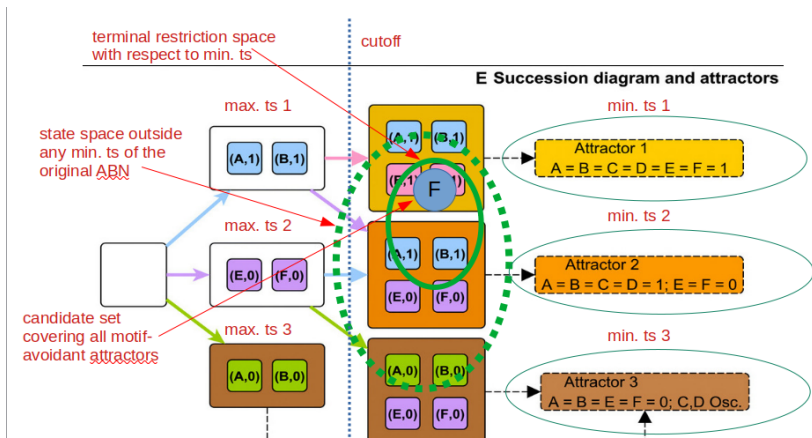
- If Preprocessing SSF is good enough, we can reach the best case where $F = \emptyset$ (i.e., there is no motif-avoidant attractor).
- Otherwise, we need to check the reachability with the target set including M_{min} .

Deal with the problem of large succession diagrams



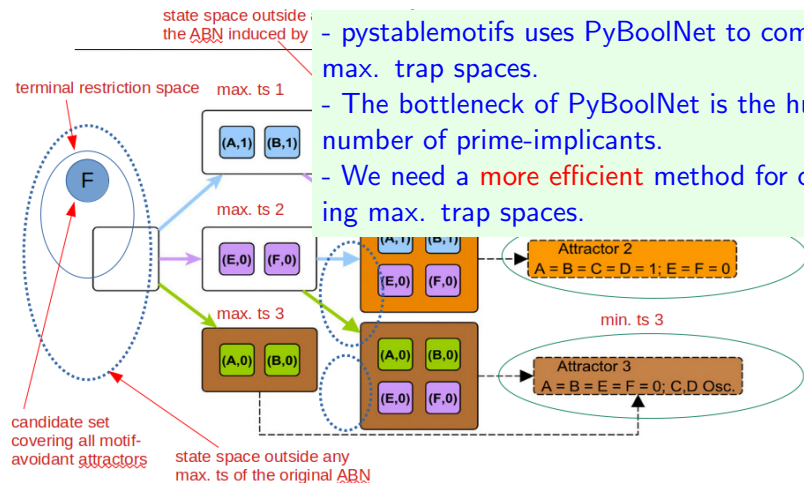
- Finally, we need to compute all attractors inside each min. trap space.
- We can simply apply the approach of mtsNFVS to the ABN induced by the min. trap space.

Deal with the problem of large succession diagrams



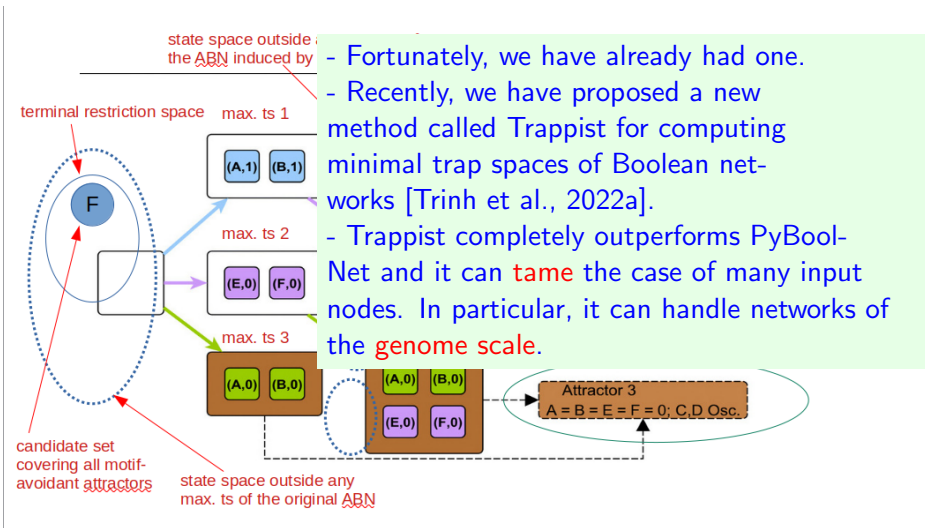
- After finishing the above process, we will obtain the whole attractor landscape and the partial succession diagram.

Computation of max. trap spaces

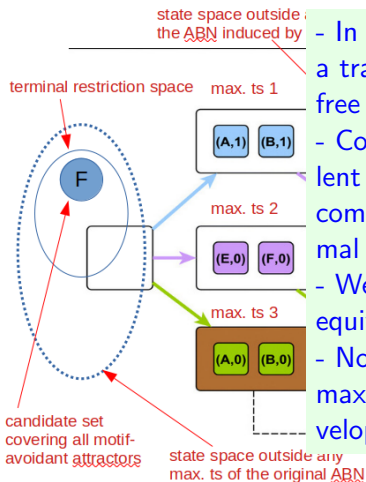


- pystablemotifs uses PyBoolNet to compute max. trap spaces.
- The bottleneck of PyBoolNet is the huge number of prime-implicants.
- We need a **more efficient** method for computing max. trap spaces.

Computation of max. trap spaces

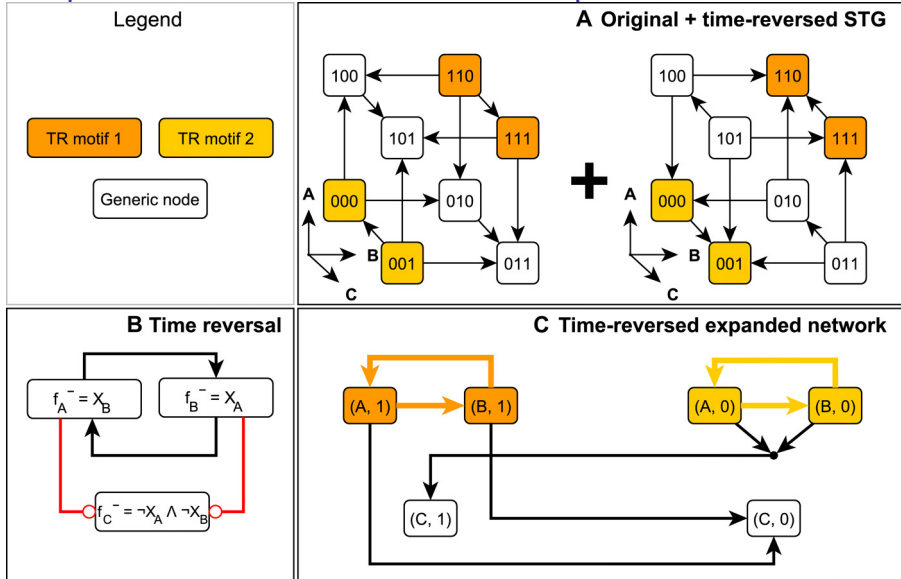


Computation of max. trap spaces

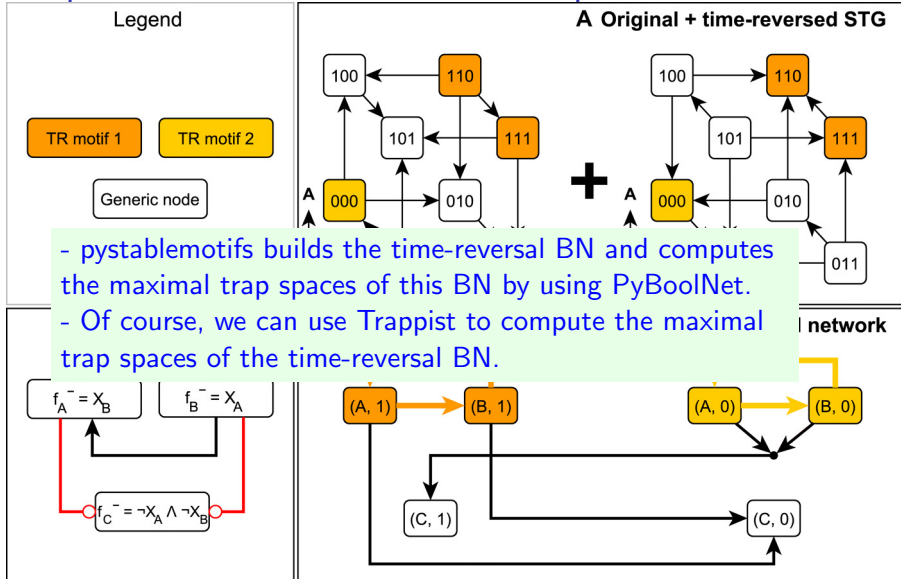


- In [Trinh et al., 2022a], we have proved that a trap space of a BN is equivalent to a conflict-free siphon of its Petri net encoding.
- Consequently, a min. trap space is equivalent to a maximal conflict-free siphon. Trappist computes min. trap spaces by computing maximal conflict-free siphons.
- We can easily prove that a **max. trap space** is equivalent to a **minimal conflict-free siphon**.
- Now, Trappist does not support computing max. trap spaces. However, we can easily develop this feature for Trappist.

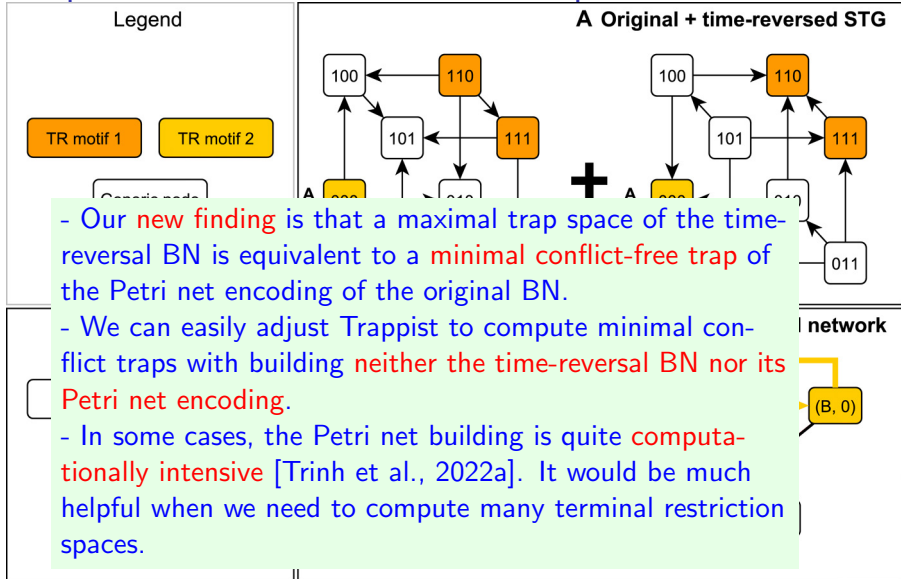
Computation of terminal restriction spaces



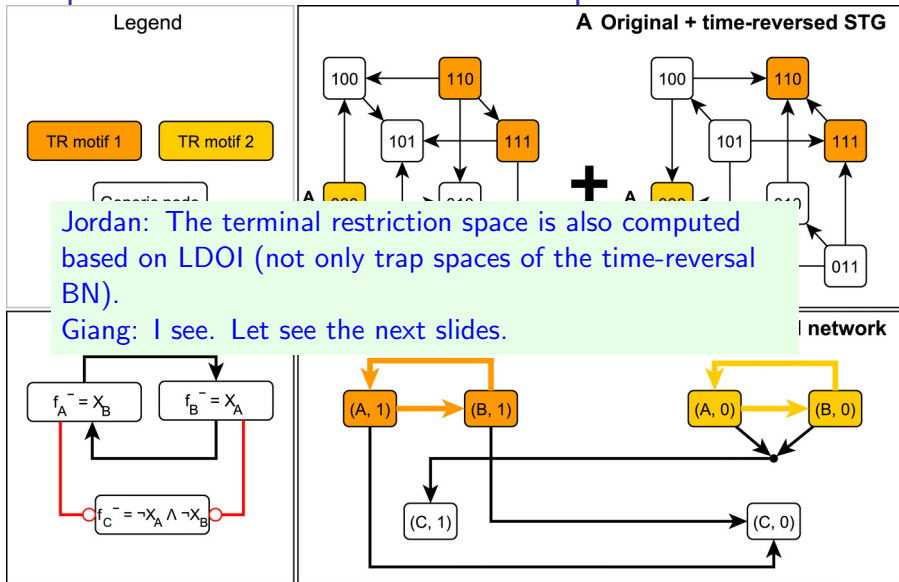
Computation of terminal restriction spaces



Computation of terminal restriction spaces



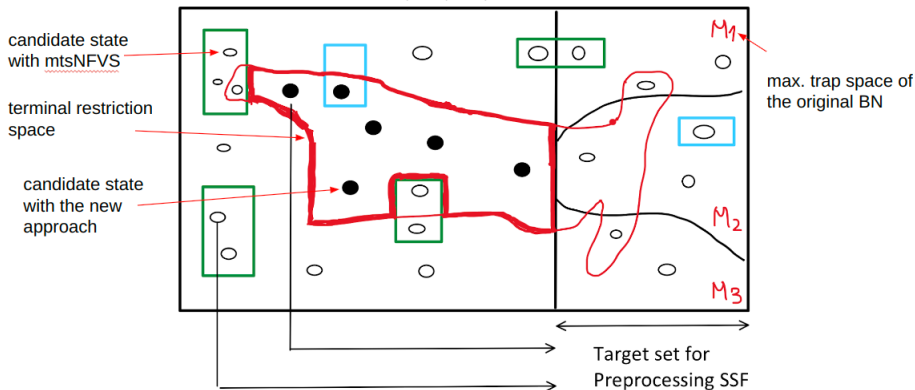
Computation of terminal restriction spaces



Computation of terminal restriction spaces (cont.)

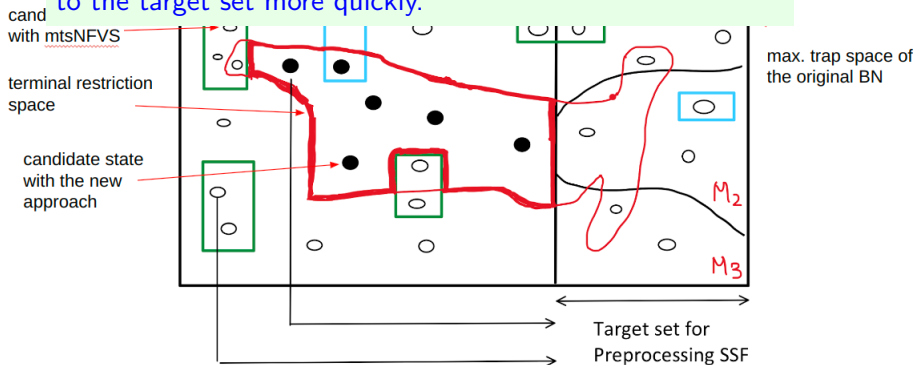
- Self-negating max. trap space of time-reversal BN
- Non-self-negating max. trap space of time-reversal BN

$R(X) = 1$ with respect to δ that is the set of all driver nodes for all M_1, M_2, M_3, \dots



Computation of terminal restriction spaces (cont.)

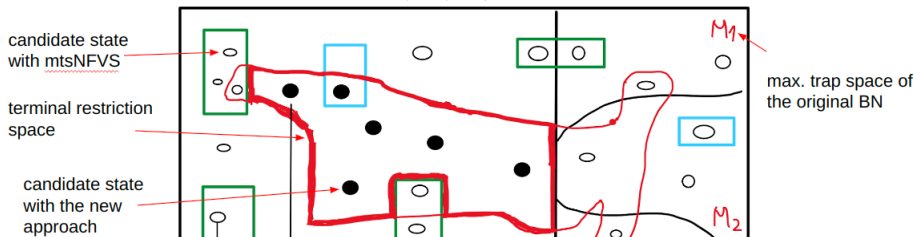
- I emphasize an **advantage** of the new approach.
- The new candidate set seems to have shorter distances to the target set.
- This is better for Preprocessing SSF as it may converge to the target set more quickly.



Computation of terminal restriction spaces (cont.)

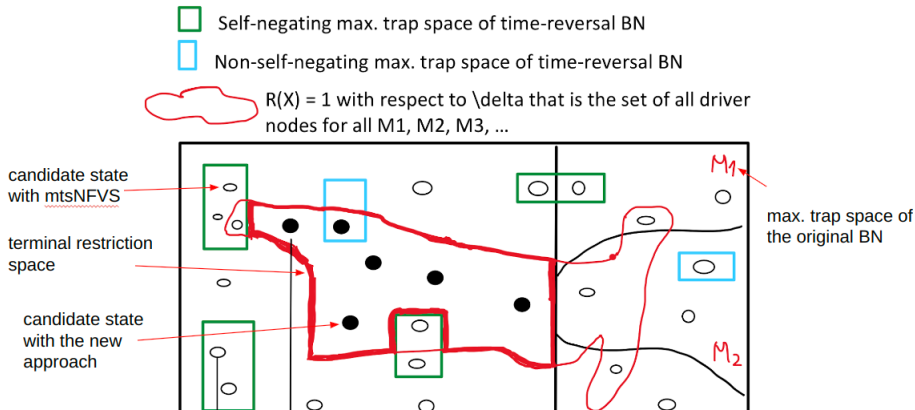
- Self-negating max. trap space of time-reversal BN
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$R(X) = 1$ with respect to δ that is the set of all driver nodes for all M_1, M_2, M_3, \dots



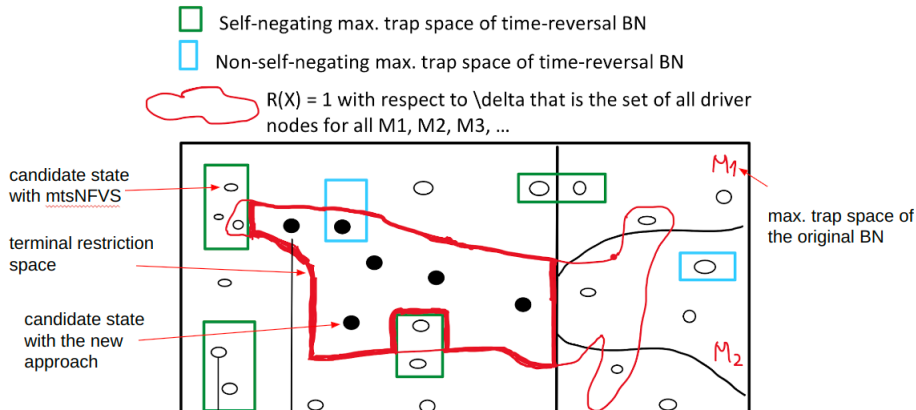
- We can efficiently compute max. trap spaces of the time-reversal BN by using Trappist.
- The self-negation of a max. trap space of the time-reversal BN can be checked by computing its LDOI in the original BN.

Computation of terminal restriction spaces (cont.)



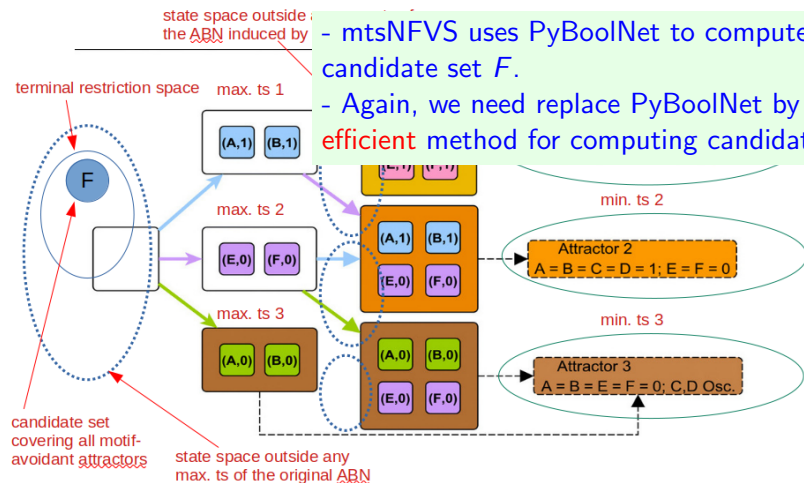
- Δ is the set of all virtual nodes that individually drive any stable motif.
- A question is how to compute Δ efficiently? Note that we need to compute Δ many times in the whole process.
- If the computation is expensive, we can ignore $R(X)$. In this case, I hope that self-negating max. trap spaces are good enough.

Computation of terminal restriction spaces (cont.)

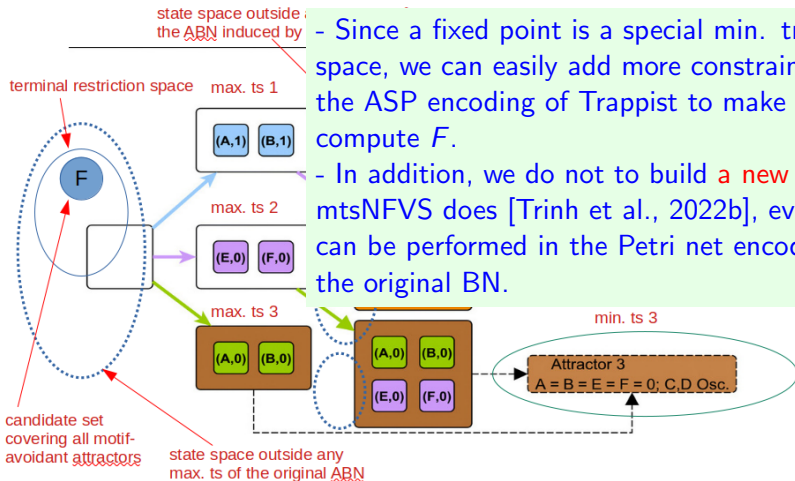


- The set of retained values B largely affects the size of the candidate set [Trinh et al., 2020].
- One interesting question is that how we can propose heuristics for setting B based on the information about the terminal restriction space to get a smaller candidate set?

Computation of candidate sets

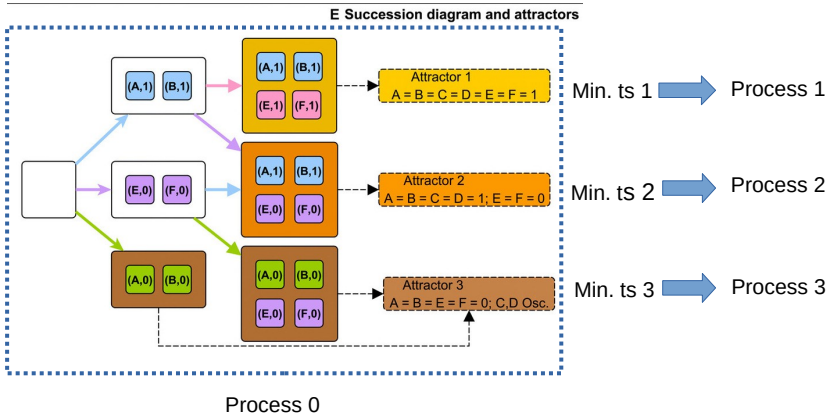


Computation of candidate sets



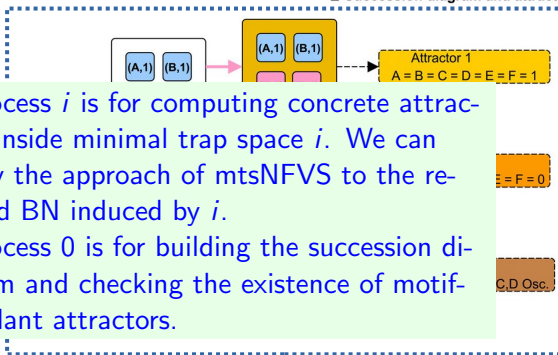
- Since a fixed point is a special min. trap space, we can easily add more constraints to the ASP encoding of Trappist to make it able to compute F .
- In addition, we do not to build a new BN as mtsNFVS does [Trinh et al., 2022b], everything can be performed in the Petri net encoding of the original BN.

Parallelization



Parallelization

E Succession diagram and attractors



- Process i is for computing concrete attractors inside minimal trap space i . We can apply the approach of mtsNFVS to the reduced BN induced by i .
- Process 0 is for building the succession diagram and checking the existence of motif-avoidant attractors.

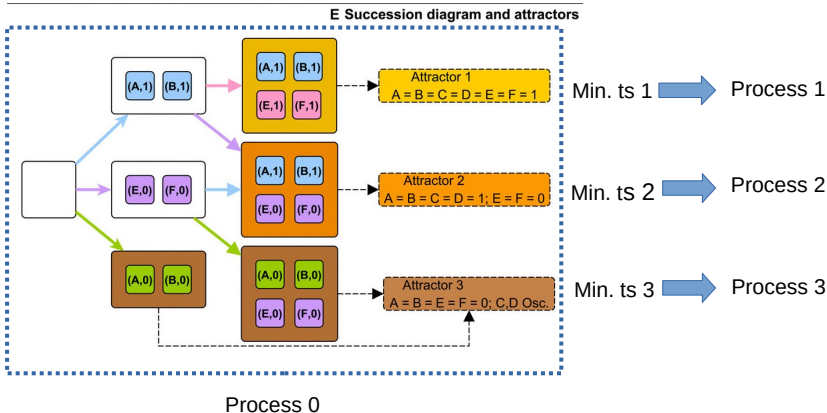
Min. ts 1 → Process 1

Min. ts 2 → Process 2

Min. ts 3 → Process 3

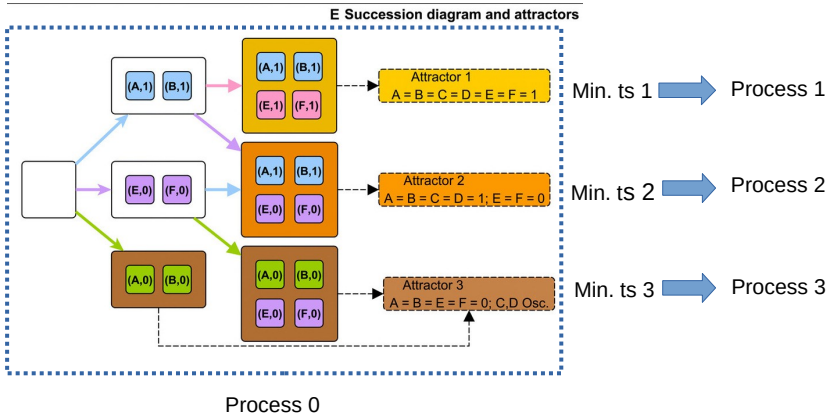
Process 0

Parallelization



- Processes 0, 1, 2, ... can be completely paralleled.

Parallelization



- Process 0 can even be divided into many paralleled sub-processes.

Target control

After finishing the attractor detection phase, we will obtain the whole attractor landscape as well as the partial/full succession diagram. If the succession diagram is full, its size should be moderate.

The next step is how to compute minimum control policies that drive the network dynamics into a target attractor/min. trap space from any initial state?

There are some existing methods for the target control of asynchronous Boolean networks: CABEAN [Su and Pang, 2021a], PyBoolNet [Fontanals et al., 2020, Cifuentes-Fontanals et al., 2021], pystablemotifs [Zañudo and Albert, 2015, Rozum et al., 2021a], AEON [Brim et al., 2021].

Target control: our observations I

CABEAN always ensures that the resulting control policies are minimum. It relies on the calculation of strong basin of attraction, which may be very computationally expensive.

PyBoolNet exploits the information on trap spaces to find control policies. However, it does not ensure that the resulting control policies are minimum. It uses model checking to improve the quality of the control policies, which may be very computationally expensive as well. Moreover, its scalability may be limited by the necessity to scan through the available perturbations.

pystablemotifs uses the succession diagram to find control policies. However, it does not ensure that the resulting control policies are minimum.

Target control: our observations II

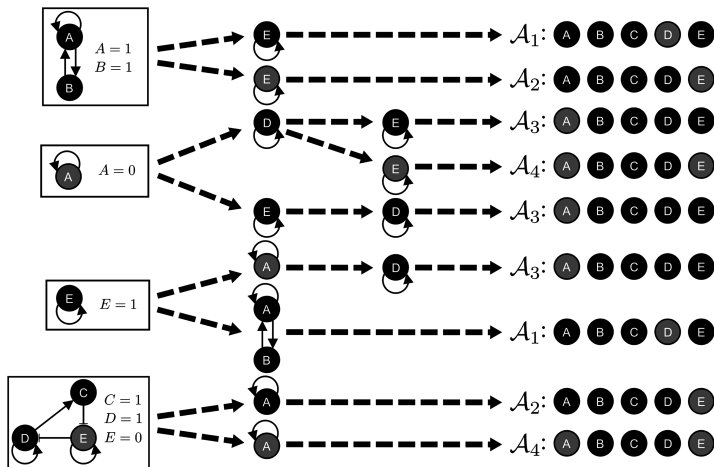
AEON uses methods based on the computation of the strong basin, but with symbolic representation to exhaustively analyse all perturbations simultaneously. A postprocessing step is then needed to extract the minimal perturbations from the final result. Alternatively, if the network has parameters, such postprocessing can also be used to extract perturbations with favourable robustness within the parameter space. The method is mainly limited by the complexity of the symbolic state space search for all available perturbations.

In the reported experiments [Su and Pang, 2021b, Su and Pang, 2021a], `pystablemotifs` is more time-efficient in most cases, though it did not return the minimum results in some cases.

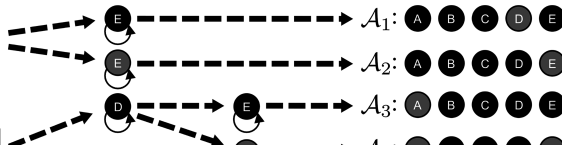
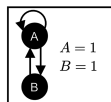
Target control: our observations III

In our opinion, pystablemotifs has much more potential than the others. We should dive into it. Of course, we need to improve both its **time-efficiency** and **accuracy**.

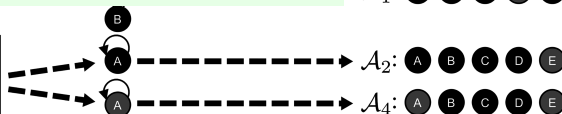
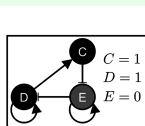
Target control with the full succession diagram



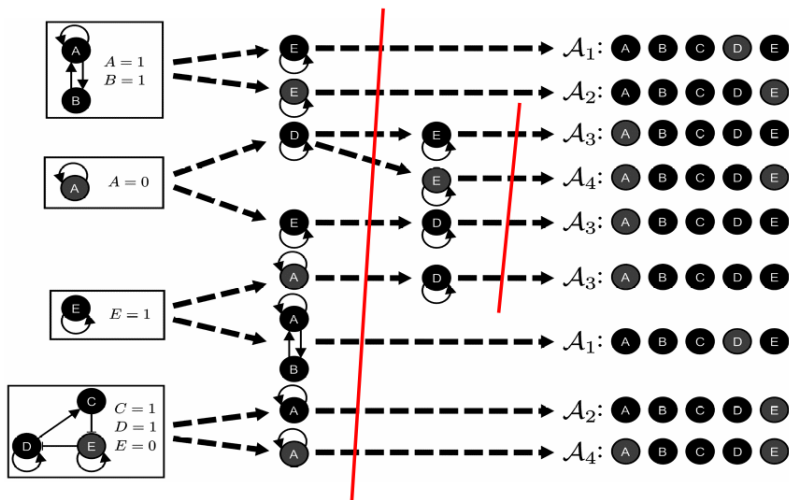
Target control with the full succession diagram



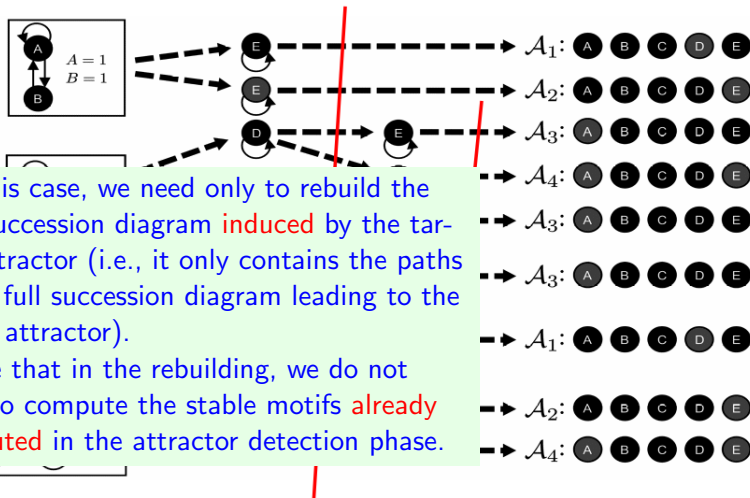
- In this case, we simply reuse the approach of `pystablemotifs`.
- Specifically, we follow all the paths in the full succession diagram leading to the target attractor to find driver nodes.



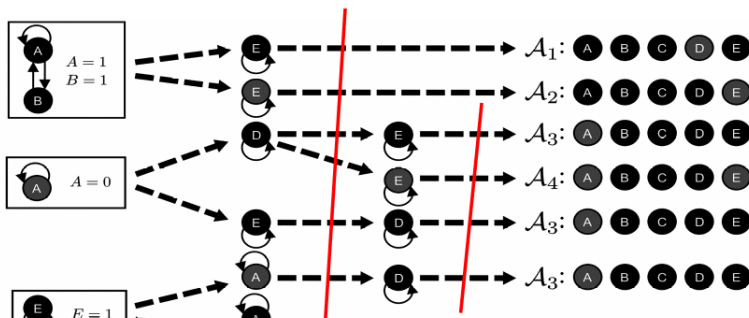
Target control with the partial succession diagram



Target control with the partial succession diagram

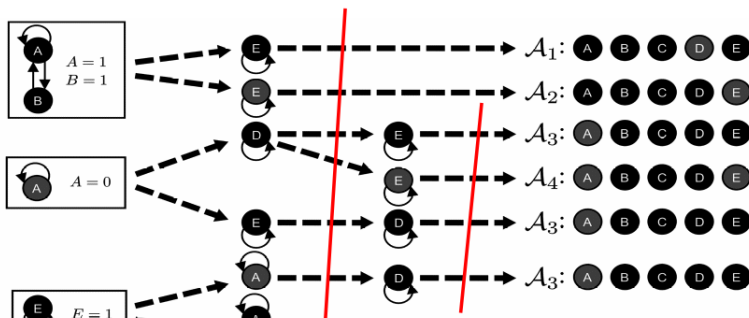


Target control with the partial succession diagram



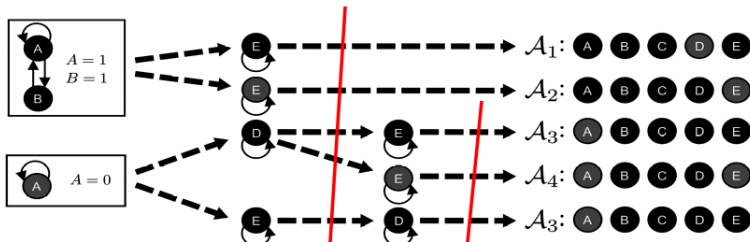
- One question may be whether we need to compute both $\{E = 1\}$ and $\{E = 0\}$. No, we **do not need**. We have already known the target attractor, hence we can add constraints to the ASP encoding to guide Traplist for searching only stable motifs that are **compatible** with the target attractor \mathcal{A}_3 .

Target control with the partial succession diagram



- Finally, we follow the **induced** succession diagram to find driver nodes (similar to the case of the full succession diagram).
- The induced succession diagram may be **much smaller** than the full one, hence the computational burden may be reduced significantly.

Target control with the partial succession diagram



Jordan: The control approach of pystablemotifs is independent to the update scheme. CABEAN is designed for asynchronous Boolean networks. Hence, the problem of not obtaining the minimum control polices is not really mattered with pystablemotifs.

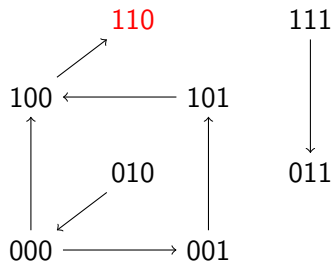
Giang: I understand. Let me recall the control example.



Improve the accuracy

Consider the example asynchronous Boolean network (Figure 2 in [Cifuentes-Fontanals et al., 2021]).

$$\begin{cases} f_1 = \neg x_2 \vee (x_1 \wedge \neg x_3) \\ f_2 = (x_1 \wedge \neg x_3) \vee (x_2 \wedge x_3) \\ f_3 = (\neg x_1 \wedge \neg x_2) \vee (x_2 \wedge x_3) \end{cases}$$



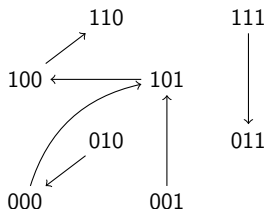
The result of `pystablemotifs` is $\{x_1 = 1, x_3 = 0\}$. The result of `CABEAN` (also of `PyBoolNet` with model checking) is $\{x_3 = 0\}$.

We have obtained a preliminary idea for improving the accuracy of `pystablemotifs`. However, we need to make it more specific as well as investigate its efficiency. In most cases, smaller control policies, longer computational time. We will present it in the next slides.

Improve the accuracy (cont.)

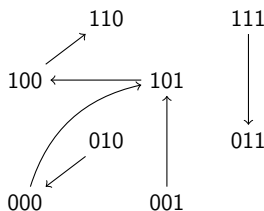
Consider the **synchronous** counterpart of the previous asynchronous Boolean network.

$$\begin{cases} f_1 = \neg x_2 \vee (x_1 \wedge \neg x_3) \\ f_2 = (x_1 \wedge \neg x_3) \vee (x_2 \wedge x_3) \\ f_3 = (\neg x_1 \wedge \neg x_2) \vee (x_2 \wedge x_3) \end{cases}$$



Improve the accuracy (cont.)

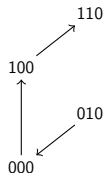
$$\begin{cases} f_1 = \neg x_2 \vee (x_1 \wedge \neg x_3) \\ f_2 = (x_1 \wedge \neg x_3) \vee (x_2 \wedge x_3) \\ f_3 = (\neg x_1 \wedge \neg x_2) \vee (x_2 \wedge x_3) \end{cases}$$



Assume that the target attractor is $\{110\}$. The result of pystablemotifs is $\{x_1 = 1, x_3 = 0\}$. The minimum result still should be $\{x_3 = 0\}$.

Improve the accuracy (cont.)

$$\begin{cases} f_1 = \neg x_2 \vee x_1 \\ f_2 = x_1 \end{cases}$$

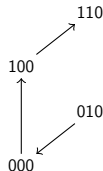


$\{x_1 = 1, x_3 = 0\}$ is a stable motif obtained by `pystablemotifs`. The function `minimal_drivers` (in the file `drivers.py`) finds the minimum subsets of stable motif's states that, when fixed in the logical model, are enough to force the state of every node in the motif into the stable motif M . This function considers all subsets of size 1, 2, ..., $|M| - 1$ and relies on `LDOI` (`implied, contradicted = logical_domain_of_influence(driver_dict, primes)`).

In the example, $\{x_3 = 0\}$ cannot be obtained by using this method.

Improve the accuracy (cont.)

$$\begin{cases} f_1 = \neg x_2 \vee x_1 \\ f_2 = x_1 \end{cases}$$



I guess if we use DOI instead of LDOI, we can get $\{x_3 = 0\}$ instead of $\{x_1 = 1, x_3 = 0\}$. However, calculating DOI can be **difficult** in general.

Research question: How to develop **an efficient algorithm** for driver node identification in a stable motif regardless of update schemes?

The **efficient** word means that the algorithm is **fast** and can return **smaller** (maybe not optimal) sets of drive nodes.

Ideas for driver node identification

From the previous example, we observed that:

- If fixing $x_1 = 1$, the resulting BN has two max. trap spaces $\{x_3 = 0\}$ and $\{x_2 = 1, x_3 = 1\}$. The latter is not compatible with the stable motif $\{x_1 = 1, x_3 = 0\}$.
- If fixing $x_3 = 0$, the resulting BN has only one max. trap space $\{x_1 = 1\}$. All max. trap spaces are compatible with the stable motif $\{x_1 = 1, x_3 = 0\}$.

We formulate a problem as follows.

Driver node identification

Let \mathcal{N} be a BN and $M = \{x_{i_1} = a_{i_1}, \dots, x_{i_k} = a_{i_k}\}$ be a max. trap space of \mathcal{N} . Find all minimum subsets S of M such that of max. trap spaces of the new BN induced by S are compatible with M .

This problem is related to **the fixed point control problem** considered in [Biane and Delaplace, 2019, Moon et al., 2022].

Ideas for driver node identification

We can solve this problem by using **answer set programming**. The main idea is as follows.

- Consider the nodes in M as controllable nodes (x_1, x_3) and the remaining nodes as uncontrollable nodes (x_2).
- Introduce an auxiliary Boolean variable d_i for each controllable node x_i (d_1, d_3). If $d_i = 1$, node x_i is fixed to a_i . If $d_i = 0$, node x_i is updated by using its original Boolean function.
- Let \mathcal{L} be the ASP characterizing the trap spaces of \mathcal{N} considering all d_i . Let C be the constraint characterizing the max. trap space M ($x_1 = 1 \wedge x_3 = 0$).
- Compute the set A of minimal answer sets (projecting only d_i not x_i) of $\mathcal{L} + C$. $A = \{\{d_1\}, \{d_3\}\}$
- Compute the set A' of minimal answer sets (projecting only d_i not x_i) of $\mathcal{L} + \neg C + A$. $A' = \{\{d_1\}\}$
- Return $A \setminus A'$. **Return** $\{\{d_3\}\}$ where $\{d_3\}$ corresponds to $\{x_3 = 0\}$

Ideas for driver node identification

This ASP-based method can avoid considering explicitly all subsets of size 1, 2, ..., $|M| - 1$. This exploits the power of ASP solvers (e.g., clingo).

However, it is still needed to test its efficiency in practice. Moreover, we also need to investigate its correctness thoroughly.

Bottlenecks of both pystablemotifs and mtsNFVS

The new approach still encounters the two **bottlenecks** of both pystablemotifs and mtsNFVS.

Bottlenecks of both pystablemotifs and mtsNFVS

The case of many source nodes where there are actually too many attractors (at least 2^k where k is the number of source nodes). mtsNFVS is less affected than pystablemotifs, but this is still problem because mtsNFVS uses the set of **explicit** (not **symbolic**) min. trap spaces.

We have found a wise way to symbolically represent the **maximal answer sets** (corresponding to **min. trap spaces**) of the encoded ASP characterizing all trap spaces.

For **max. trap spaces**, we need to adjust the current method. We are thinking about this.

Bottlenecks of both pystablemotifs and mtsNFVS

There exist some very huge attractors inside min. trap spaces. In this case, the number free variables of a min. trap space is **still too large** (even all the nodes of the original network). pystablemotifs only ends with min. trap spaces. It is **very hard** for mtsNFVS to reach the best case (i.e., $|F| = 1$) to avoid the reachability analysis.

One example is the model shown in <https://github.com/jcrozum/pystablemotifs/issues/83>. This network has no stable motifs, and we have exactly one attractor for each combination of source node values.

Our general idea is as follows.

- We have the **information** about the min. trap space. We can use this information to guide Preprocessing SSF to reach the easy cases where we can avoid or at least reduce the reachability analysis times.

• ...

Implementation

Trappist (in Python): .bnet file $\xrightarrow{\text{pyeda}}$ BDDs of Boolean functions \rightarrow Petri net (digraph) \rightarrow ASP (string) $\xrightarrow{\text{clingo}}$ min./max. trap spaces (json)

We think that we should start from scratch, i.e., we remove completely the dependency on PyBoolNet because it uses a heavy data structure (i.e, prime-implicants).

We propose to design **a completely new system** for attractor detection and control. We will gradually adapt the functions from pystablemotifs, mtsNFVS, and Trappist to the new system. Of course, we need a **good system design** first.

Note that our proposed approach contains many **constituent tasks** (e.g., computation of min./max. trap spaces, computation of LDOI, driver node identification). The new system should be easy to replace new implementations for each constituent task.

Experimentation

We plan to test the new method on models without **constant** nodes.

N-K models: expected to handle networks with $K = 2$ and up to 500 nodes.

scale-free models: expected to handle networks with $K = 10$ and up to 5000 nodes.

real-world models: expected to handle all the largest and most complex networks that we can find in the literature. The repository ² maintained by Samuel Pastva is a good source.

²<https://github.com/sybila/biodivine-boolean-models>

Conclusion

The new approach exploits advantages of both pystablemotifs and mtsNFVS.

It benefits not only the attractor detection issue but also the control issue of asynchronous Boolean networks.

We believe that it will be the **most promising** approach that can reach the genome scale.

Collaboration

We hope we will have a nice collaboration on this work. Hence, we should make clear something at the beginning.

Participants:

- Our side: Van-Giang Trinh and Samuel Pastva.
- Your side: Jordan Rozum and **Kyu Hyong Park**.

Authorship (if we can publish some papers :)):

- Two first authors, one from our side and one from your side with the equal contribution?

What do you think?

Next steps

Task 1: Complete our new approach and make it as detailed as possible.

- Find solutions for some hard problems (e.g., the case of very huge attractors).
- Prepare more specific slides (or a report) presenting technical details of the new approach.
- Add more concrete examples for illustration.

Task 2: Start the implementation of our new approach.

- Design a new system for attractor detection and control (input/output formats, components, functions, etc.).
- Programming

We should create a Github repository first. Everything will be updated there.

Our proposal

I and Samuel will focus on Task 1. Of course, Jordan and Kyu will also involve in this task to make the new approach as powerful as possible.

Jordan and Kyu will lead Task 2 because you already had the holistic view about pystacklemotifs. You can divide and assign to all of them implementation sub-tasks. I and Samuel will contribute to this process.

Should keep regular exchanges between us.

Your suggestions?

Thank you for your attention!

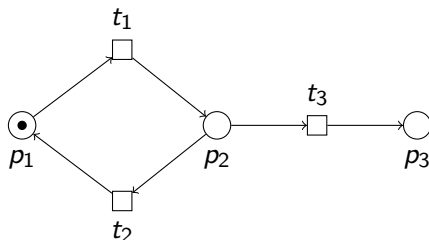
Petri nets

Bipartite graph

Places P

Transitions T

Weighted arcs W



A *marking* for a Petri net is a mapping $m : P \mapsto \mathbb{N}$ that assigns a number of tokens to each place. A place p is marked by a marking m if and only if $m(p) > 0$. For example, p_1 is marked, p_2 and p_3 are unmarked.

We shall write $\text{pred}(x)$ (resp. $\text{succ}(x)$) to represent the set of vertices that have a (non-zero weighted) arc leading to (resp. coming from) x . For example, $\text{pred}(p_2) = \{t_1\}$, $\text{succ}(p_2) = \{t_2, t_3\}$, $\text{pred}(t_2) = \{p_2\}$, and $\text{succ}(t_2) = \{p_1\}$.

Siphons

Siphon

A *siphon* of a Petri net (P, T, W) is a set of places S such that:

$$\forall t \in T, S \cap \text{succ}(t) \neq \emptyset \Rightarrow S \cap \text{pred}(t) \neq \emptyset.$$

Once a siphon is unmarked, it remains unmarked.

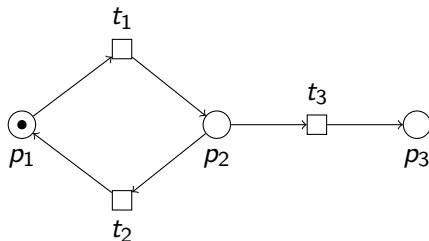
$S = \{p_1, p_3\}$ is not a siphon because

$$S \cap \text{succ}(t_3) = \{p_1, p_3\} \cap \{p_3\} =$$

$$\{p_3\} \neq \emptyset \text{ but}$$

$$S \cap \text{pred}(t_3) = \{p_1, p_3\} \cap \{p_2\} = \emptyset.$$

Here: $\emptyset, \{p_1, p_2\}, \{p_1, p_2, p_3\}$



Traps

Trap

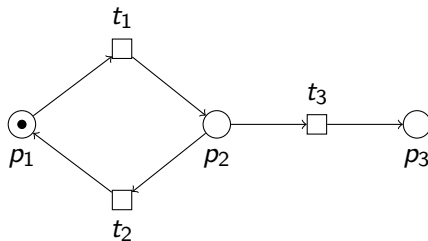
A *trap* of a Petri net (P, T, W) is a set of places S such that:

$$\forall t \in T, S \cap \text{pred}(t) \neq \emptyset \Rightarrow S \cap \text{succ}(t) \neq \emptyset.$$

Once a trap is marked (i.e., at least one of its places is marked), it remains marked.

$S = \{p_1, p_3\}$ is not a trap because
 $S \cap \text{pred}(t_1) = \{p_1, p_3\} \cap \{p_1\} = \{p_1\} \neq \emptyset$ but
 $S \cap \text{succ}(t_1) = \{p_1, p_3\} \cap \{p_2\} = \emptyset.$

Here: $\emptyset, \{p_1, p_2, p_3\}$



Petri net of a Boolean model

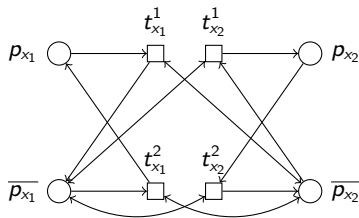
The original encoding was established in [Chaouiya et al., 2004].

Two places for each gene: $v \rightsquigarrow p_v, \overline{p_v}$

Solutions of $f_v \not\leftrightarrow v \rightsquigarrow$ transitions from p_v to $\overline{p_v}$ (and back)

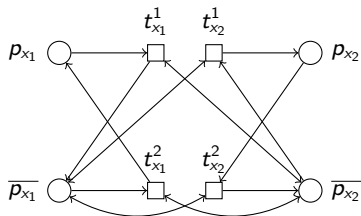
At any marking m of the Petri net encoding a Boolean model, it always holds that $m(p_v) + m(\overline{p_v}) = 1$.

$$\begin{cases} f_1 = (x_1 \wedge x_2) \vee (\neg x_1 \wedge \neg x_2) \\ f_2 = (x_1 \wedge x_2) \vee (\neg x_1 \wedge \neg x_2) \end{cases}$$



Conflict-free siphons/traps

A siphon/trap is called **conflict-free** if it does not contain both p_v and $\overline{p_v}$ for all $v \in V$.



Siphon	Conflict-free?
\emptyset	yes
$\{\overline{p_{x_1}}, \overline{p_{x_2}}\}$	yes
$\{p_{x_1}, \overline{p_{x_1}}\}$	no
$\{p_{x_2}, \overline{p_{x_2}}\}$	no

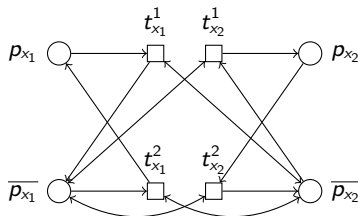
A conflict-free siphon (resp. trap) is *maximal* if it is not a subset of any other conflict-free siphon (resp. trap).

Conflict-free siphons are trap spaces

Theorem 1

Let \mathcal{M} be a Boolean model and \mathcal{P} be its Petri net encoding. There is a **one-to-one correspondence** between the set of **trap spaces** of \mathcal{M} and the set of **conflict-free siphons** of \mathcal{P} .

$$\begin{cases} f_1 = (x_1 \wedge x_2) \vee (\neg x_1 \wedge \neg x_2) \\ f_2 = (x_1 \wedge x_2) \vee (\neg x_1 \wedge \neg x_2) \end{cases}$$



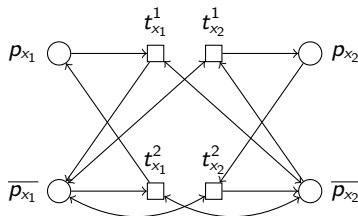
Trap space	Conflict-free siphon
★★	\emptyset
11	$\{\overline{p_{x_1}}, \overline{p_{x_2}}\}$

Maximal conflict-free siphons are minimal trap spaces

Theorem 2

Let \mathcal{M} be a Boolean model and \mathcal{P} be its Petri net encoding. There is a **one-to-one correspondence** between the set of **minimal trap spaces** of \mathcal{M} and the set of **maximal conflict-free siphons** of \mathcal{P} .

$$\begin{cases} f_1 = (x_1 \wedge x_2) \vee (\neg x_1 \wedge \neg x_2) \\ f_2 = (x_1 \wedge x_2) \vee (\neg x_1 \wedge \neg x_2) \end{cases}$$



Trap space	Conflict-free siphon
★★	\emptyset
11	$\{\overline{p_{x_1}}, \overline{p_{x_2}}\}$

Proposed method for minimal trap space computation

From Theorem 2, we propose a new method for computing minimal trap spaces of a Boolean model \mathcal{M} .

- Build the Petri net encoding \mathcal{P} of \mathcal{M} .
- Compute all maximal conflict-free siphons of \mathcal{P} .
- Convert the obtained maximal conflict-free siphons into the corresponding minimal trap spaces.

Petri net transformation

Transforming a Boolean model into its Petri net encoding can be done via computing Disjunctive Normal Forms (DNF) of each Boolean function [Chatain et al., 2014]. This transformation is implemented in the bioLQM³ library using BDDs.

Though this might appear quite computationally intensive it is important to remark first that contrary to the prime-implicants case, there is no need to find *minimal* DNFs.

We use the above transformation in our proposed method.

³<http://www.colomoto.org/biolqm/>

Maximal conflict-free siphon computation

Characterize all siphons of the encoded Petri net as a system of Boolean rules.

$$p \in S \Rightarrow \bigvee_{p' \in \text{pred}(t)} p' \in S, p \in P, t \in T, t \in \text{pred}(p)$$

Add to the system the Boolean rules representing the conflict-freeness.

$$p_v \in S \Rightarrow \overline{p_v} \notin S \wedge \overline{p_v} \in S \Rightarrow p_v \notin S, v \in V$$

Encode the system as an ASP.

Use an ASP solver (e.g., clingo [Gebser et al., 2011]) to compute all set-inclusion maximal answer sets of the ASP.

Set-maximality through “heuristics”

```
clingo --heuristic=Domain --enum-mod=domRec --dom-mod=3
```

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