New ideas for scalable attractor detection and control of asynchronous Boolean networks

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Attractor detection and control of asynchronous Boolean networks

Two most important issues in Boolean networks research with many applications in various fields [Akutsu, 2018, Schwab et al., 2020].

Usually, the attractor detection issue is the preceding step of the control issue.

Unfortunately, both of them are very challenging, especially for large-scale networks.

Many methods have been proposed.

Attractor detection methods

BDD-based methods: genYsis [Garg et al., 2008], geneFatt [Zheng et al., 2013], CABEAN [Mizera et al., 2019, Su and Pang, 2021a].

Reduction-based methods: AEON [Benes et al., 2021].

Trap space-based methods: PyBoolNet [1] [Klarner et al., 2017], pystablemotifs [Rozum et al., 2021b, Rozum et al., 2021a].

Feedback vertex set-based methods: FVS-ABN [Trinh et al., 2020], iFVS-ABN [Giang and Hiraishi, 2021], mtsNFVS [Trinh et al., 2022b].

All of them have advantages and disadvantages. To our best knowledge, none of them can robustly handle complex networks with thousands of nodes.

[1] It is not an exact method.

Control methods

BDD-based and decomposition-based methods:
CABEAN [Su and Pang, 2021b, Su and Pang, 2021a].

Reduction-based methods: AEON [Brim et al., 2021].

Trap space-based methods:
PyBoolNet [Fontanals et al., 2020, Cifuentes-Fontanals et al., 2021],
pystablemotifs [Zañudo and Albert, 2015, Rozum et al., 2021a].

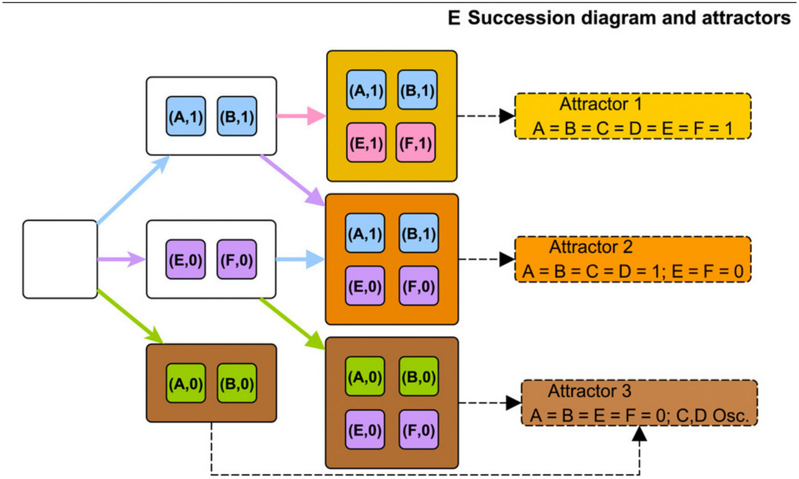
Objective

pystablemotifs and mtsNFVS have much more potential to reach the genome scale.

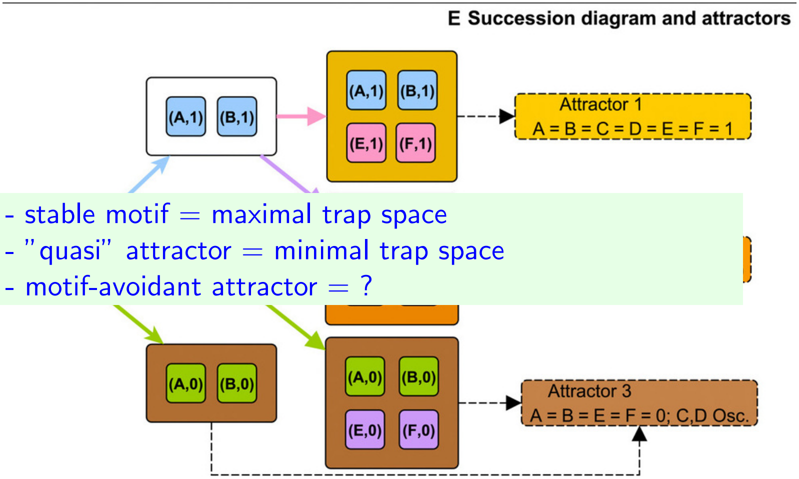
Both of them have their own advantages and disadvantages, but fortunately they can be complementary to each other.

Develop a more efficient method that can benefit both the attractor detection and control issues, and in particular can handle networks at the genome scale.

pystablemotifs

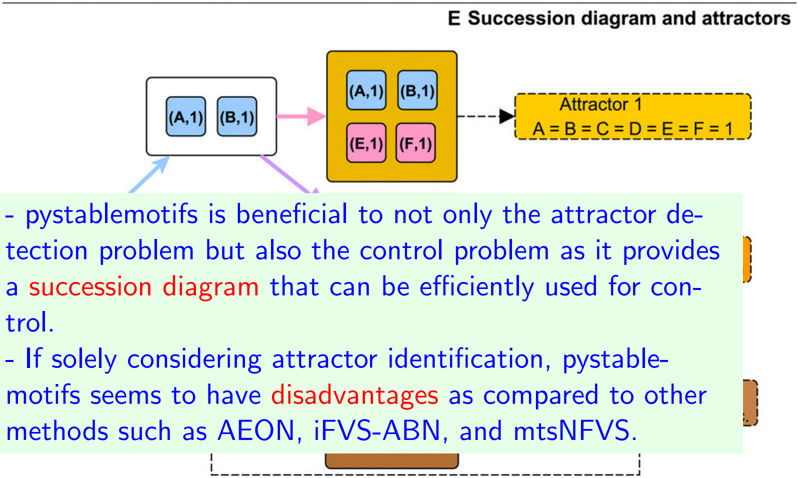


pystablemotifs


E Succession diagram and attractors
(A,1)
(B,1)
(A,1)
(B,1)
(E,1)
(F,1)
Attractor 1
A = B = C = D = E = F = 1
(A,0)
(B,0)
(A,0)
(B,0)
(E,0)
(F,0)
Attractor 3
A = B = E = F = 0; C,D Osc.

- stable motif = maximal trap space
- "quasi" attractor = minimal trap space
- motif-avoidant attractor = ?

pystablemotifs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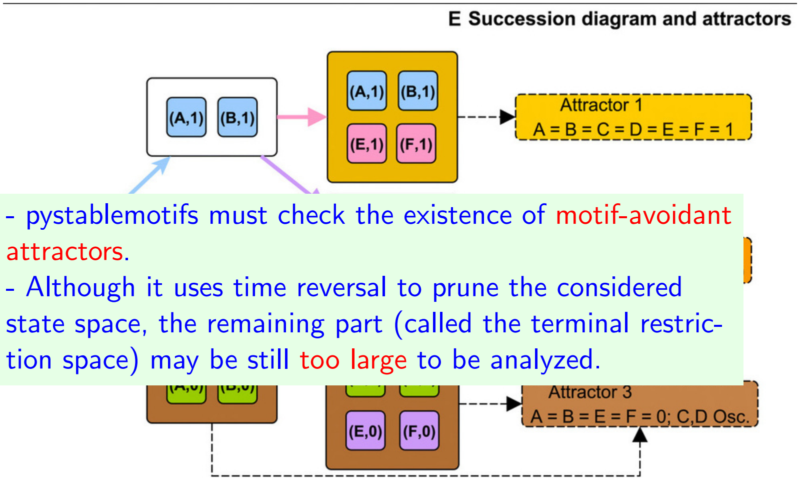
(A,1)
(B,1)
(A,1)
(B,1)
(E,1)
(F,1)
Attractor 1
A = B = C = D = E = F = 1

- pystablemotifs is beneficial to not only the attractor detection problem but also the control problem as it provides a succession diagram that can be efficiently used for control.
- If solely considering attractor identification, pystablemotifs seems to have disadvantages as compared to other methods such as AEON, iFVS-ABN, and mtsNFVS.

pystablemotifs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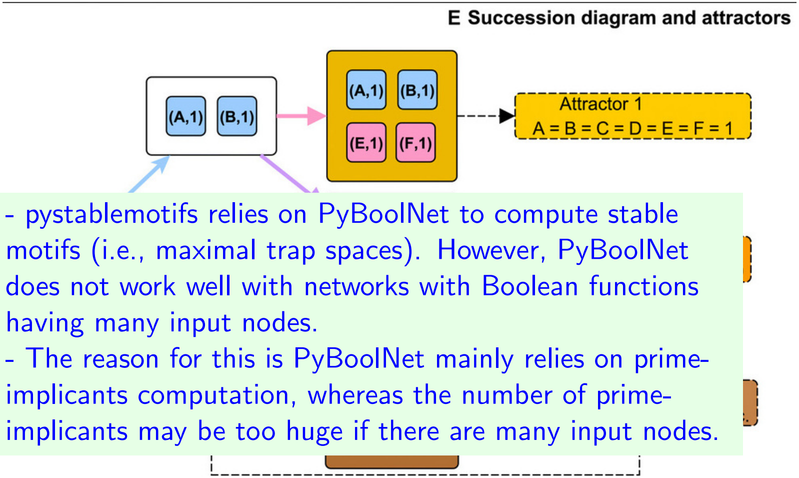
(F,1)
Attractor 1
A = B = C = D = E = F = 1
(E,0)
(F,0)
Attractor 3
A = B = E = F = 0; C,D Osc.

- pystablemotifs must check the existence of motif-avoidant attractors.
- Although it uses time reversal to prune the considered state space, the remaining part (called the terminal restriction space) may be still too large to be analyzed.

pystablemotifs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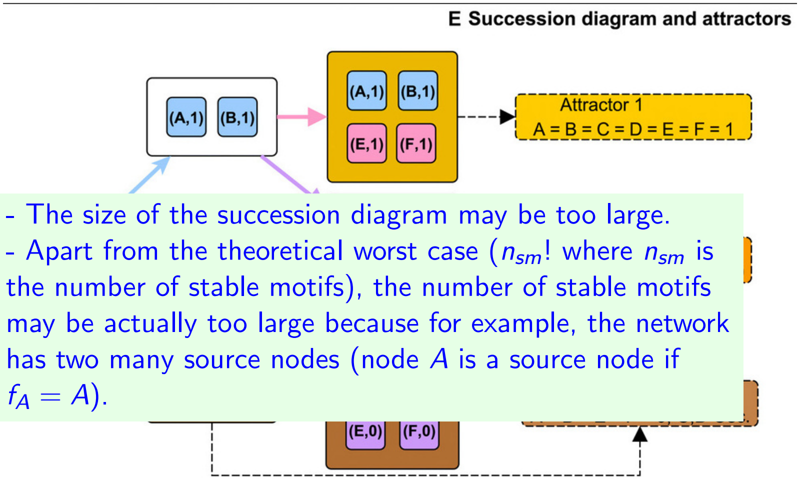
(A,1)
(B,1)
(A,1)
(B,1)
(E,1)
(F,1)
Attractor 1
A = B = C = D = E = F = 1

- pystablemotifs relies on PyBoolNet to compute stable motifs (i.e., maximal trap spaces). However, PyBoolNet does not work well with networks with Boolean functions having many input nodes.
- The reason for this is PyBoolNet mainly relies on prime-implicants computation, whereas the number of prime-implicants may be too huge if there are many input nodes.

pystablemotifs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(A,1)
(B,1)
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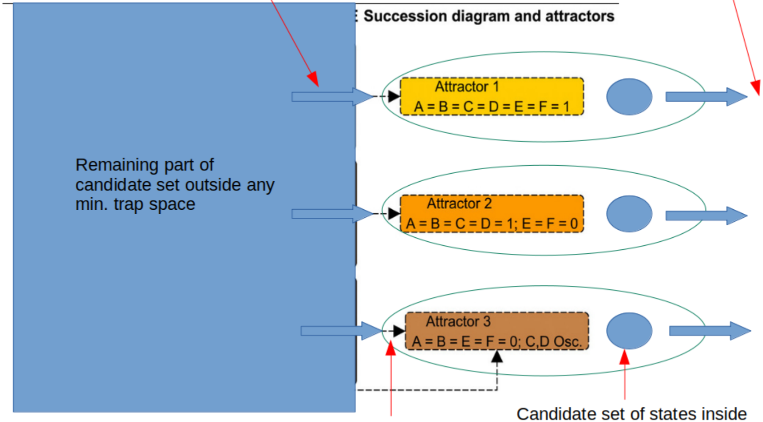
(F,1)
Attractor 1
A = B = C = D = E = F = 1
(E,0)
(F,0)

- The size of the succession diagram may be too large.
- Apart from the theoretical worst case (n_{sm}! where n_{sm} is the number of stable motifs), the number of stable motifs may be actually too large because for example, the network has two many source nodes (node A is a source node if $f_A = A$).

mtsNFVS

Preprocessing SSF (shrink the candidate set)

After Preprocessing SSF, if there is only one state, we do not need the reachability check. In this case, there is only one attractor inside this min. trap space.

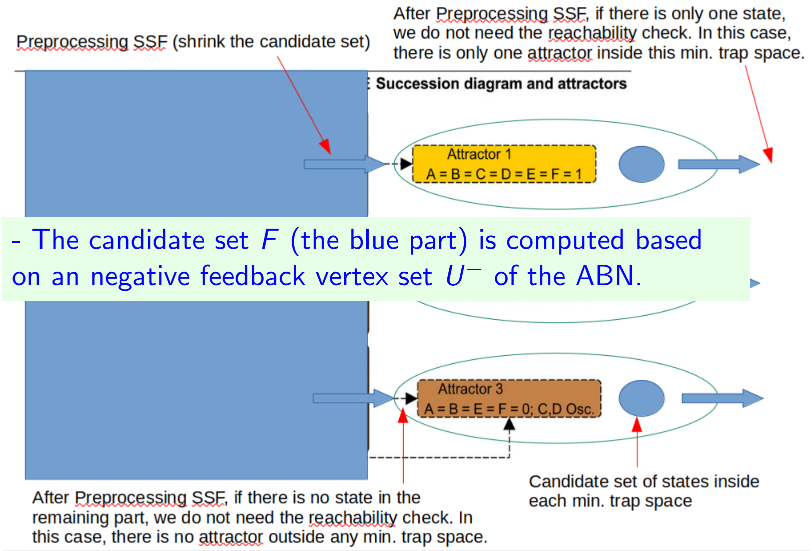


Succession diagram and attractors
Remaining part of candidate set outside any min. trap space
Attractor 1
A = B = C = D = E = F = 1
Attractor 2
A = B = C = D = 1; E = F = 0
Attractor 3
A = B = E = F = 0; C,D Osc.

After Preprocessing SSF, if there is no state in the remaining part, we do not need the reachability check. In this case, there is no attractor outside any min. trap space.

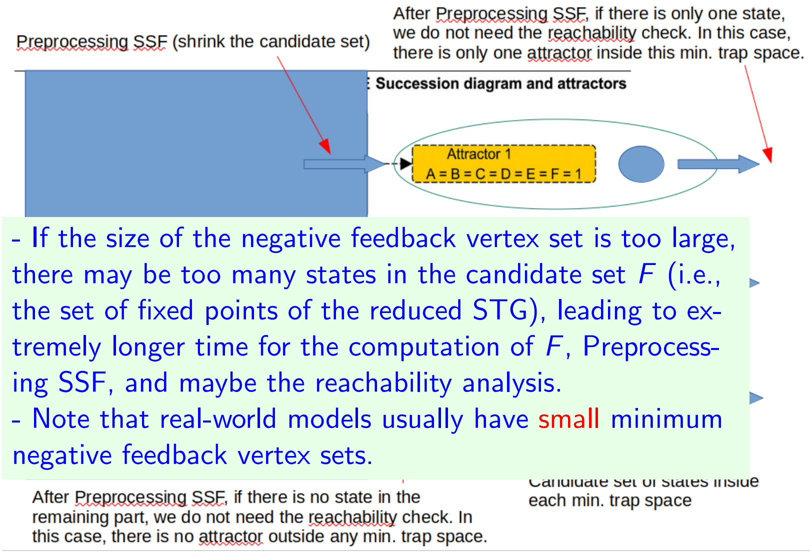
Candidate set of states inside each min. trap space

mtsNFVS


Preprocessing SSF (shrink the candidate set)
After Preprocessing SSF, if there is only one state, we do not need the reachability check. In this case, there is only one attractor inside this min. trap space.
Succession diagram and attractors
Attractor 1
A = B = C = D = E = F = 1
Attractor 3
A = B = E = F = 0; C,D Osc.
After Preprocessing SSF, if there is no state in the remaining part, we do not need the reachability check. In this case, there is no attractor outside any min. trap space.
Candidate set of states inside each min. trap space

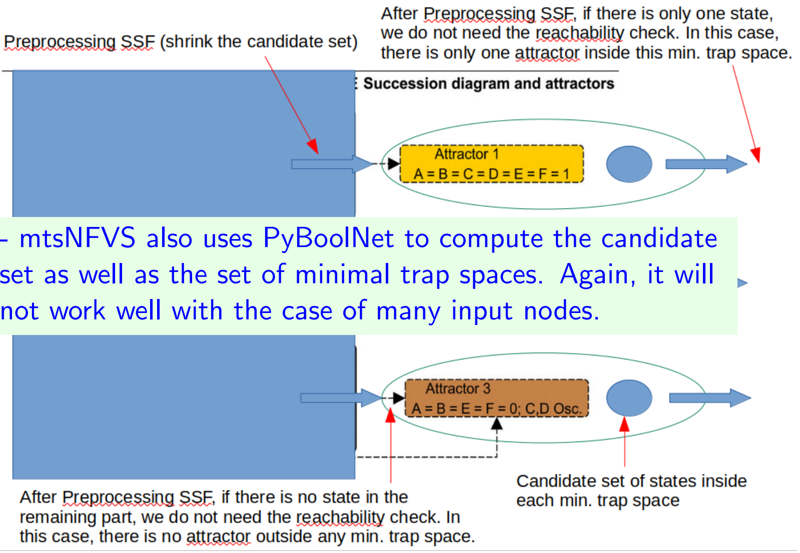
- The candidate set F (the blue part) is computed based on an negative feedback vertex set U^{-} of the ABN.

mtsNFVS


Preprocessing SSF (shrink the candidate set)
After Preprocessing SSF, if there is only one state, we do not need the reachability check. In this case, there is only one attractor inside this min. trap space.
Succession diagram and attractors
Attractor 1
A = B = C = D = E = F = 1
Candidate set of states inside each min. trap space
After Preprocessing SSF, if there is no state in the remaining part, we do not need the reachability check. In this case, there is no attractor outside any min. trap space.

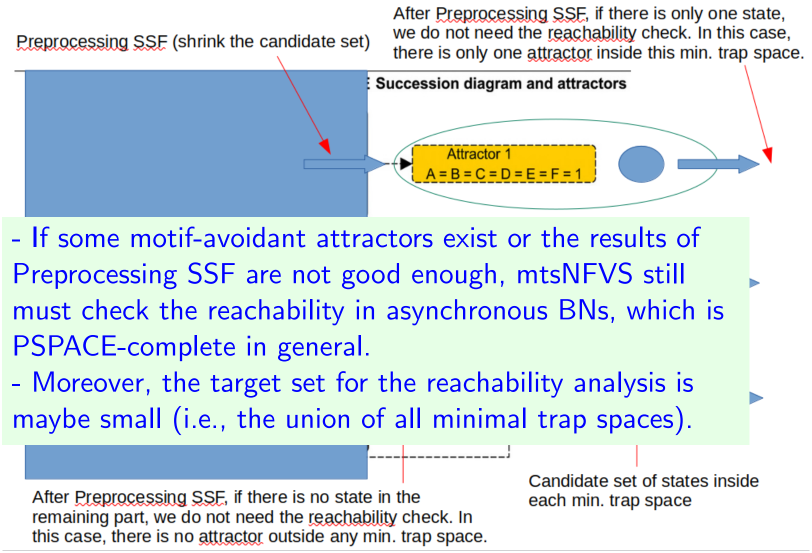
- If the size of the negative feedback vertex set is too large, there may be too many states in the candidate set F (i.e., the set of fixed points of the reduced STG), leading to extremely longer time for the computation of F, Preprocessing SSF, and maybe the reachability analysis.
- Note that real-world models usually have small minimum negative feedback vertex sets.

mtsNFVS


Preprocessing SSF (shrink the candidate set)
After Preprocessing SSF, if there is only one state, we do not need the reachability check. In this case, there is only one attractor inside this min. trap space.
Succession diagram and attractors
Attractor 1
A = B = C = D = E = F = 1
Attractor 3
A = B = E = F = 0; C,D Osc.
Candidate set of states inside each min. trap space
After Preprocessing SSF, if there is no state in the remaining part, we do not need the reachability check. In this case, there is no attractor outside any min. trap space.

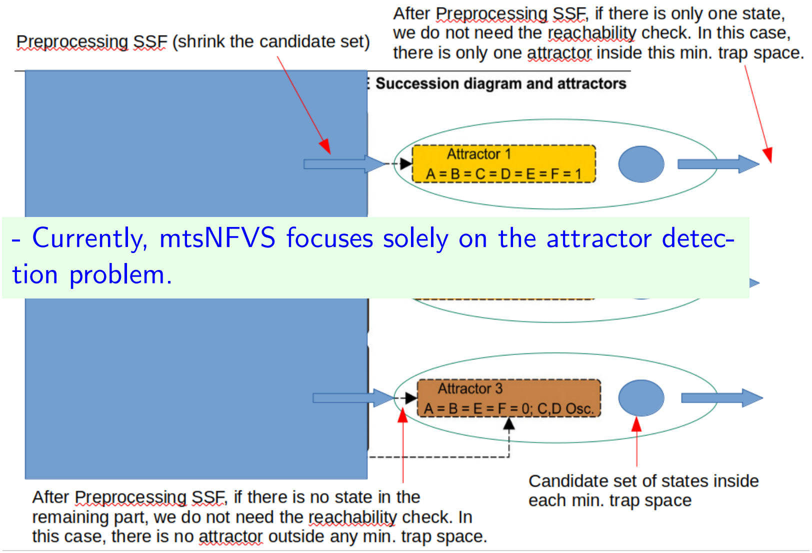
- mtsNFVS also uses PyBoolNet to compute the candidate set as well as the set of minimal trap spaces. Again, it will not work well with the case of many input nodes.

mtsNFVS


Preprocessing SSF (shrink the candidate set)
After Preprocessing SSF, if there is only one state, we do not need the reachability check. In this case, there is only one attractor inside this min. trap space.
Succession diagram and attractors
Attractor 1
A = B = C = D = E = F = 1
After Preprocessing SSF, if there is no state in the remaining part, we do not need the reachability check. In this case, there is no attractor outside any min. trap space.
Candidate set of states inside each min. trap space

- If some motif-avoidant attractors exist or the results of Preprocessing SSF are not good enough, mtsNFVS still must check the reachability in asynchronous BNs, which is PSPACE-complete in general.
- Moreover, the target set for the reachability analysis is maybe small (i.e., the union of all minimal trap spaces).

mtsNFVS



Preprocessing SSF (shrink the candidate set)
After Preprocessing SSF, if there is only one state, we do not need the reachability check. In this case, there is only one attractor inside this min. trap space.
Succession diagram and attractors
Attractor 1
A = B = C = D = E = F = 1
Attractor 3
A = B = E = F = 0; C,D Osc.
After Preprocessing SSF, if there is no state in the remaining part, we do not need the reachability check. In this case, there is no attractor outside any min. trap space.
Candidate set of states inside each min. trap space

- Currently, mtsNFVS focuses solely on the attractor detection problem.

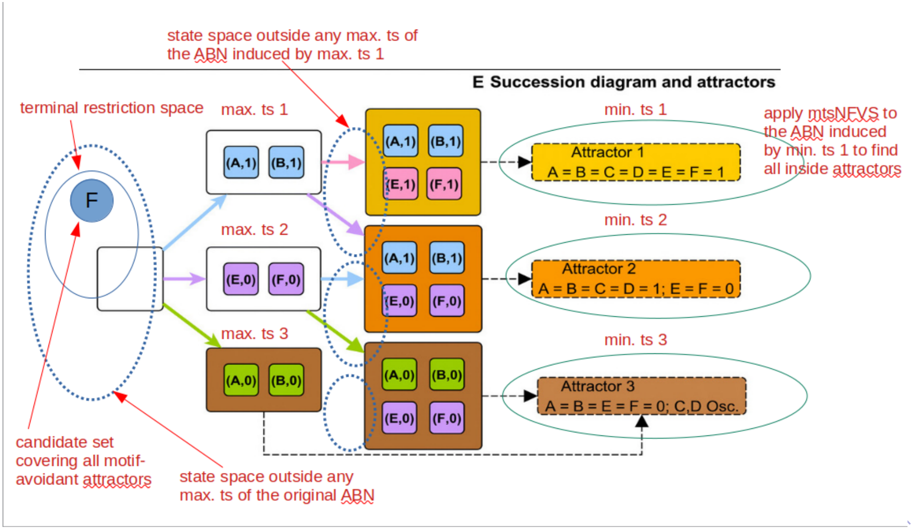
New idea

In principle, our new idea is to combine pystablemotifs and mtsNFVS to obtain a more efficient method that can benefit both the attractor detection and control problems.

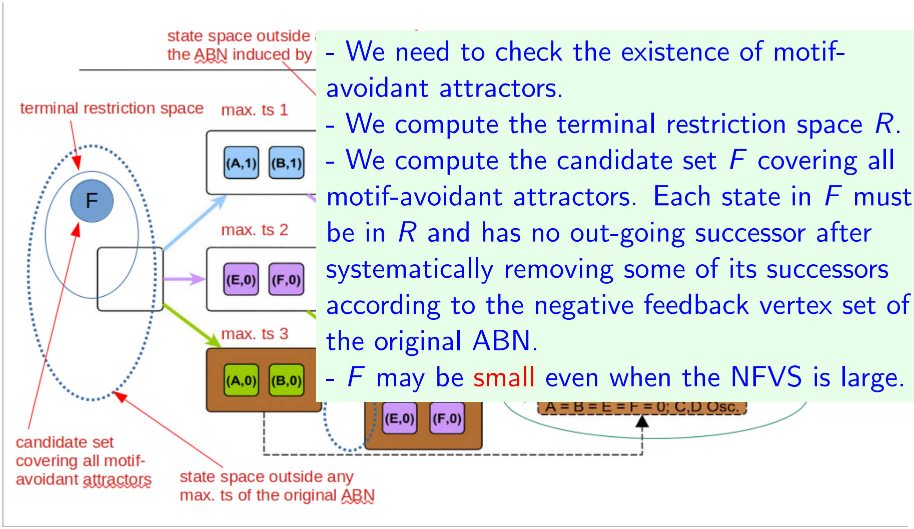
Hereafter, we will present the general approach along with many new findings that can speedup the whole process.

General approach

It follows the process of pystablemotifs (i.e., building the succession diagram).


state space outside any max. ts of the ABN induced by max. ts 1
E Succession diagram and attractors
terminal restriction space
max. ts 1
min. ts 1
apply mtsNFVS to the ABN induced by min. ts 1 to find all inside attractors
F
(A,1)
(B,1)
(E,1)
(F,1)
Attractor 1
A = B = C = D = E = F = 1
max. ts 2
min. ts 2
(E,0)
(F,0)
Attractor 2
A = B = C = D = 1; E = F = 0
max. ts 3
min. ts 3
(A,0)
(B,0)
Attractor 3
A = B = E = F = 0; C,D Osc.
candidate set covering all motif-avoidant attractors
state space outside any max. ts of the original ABN

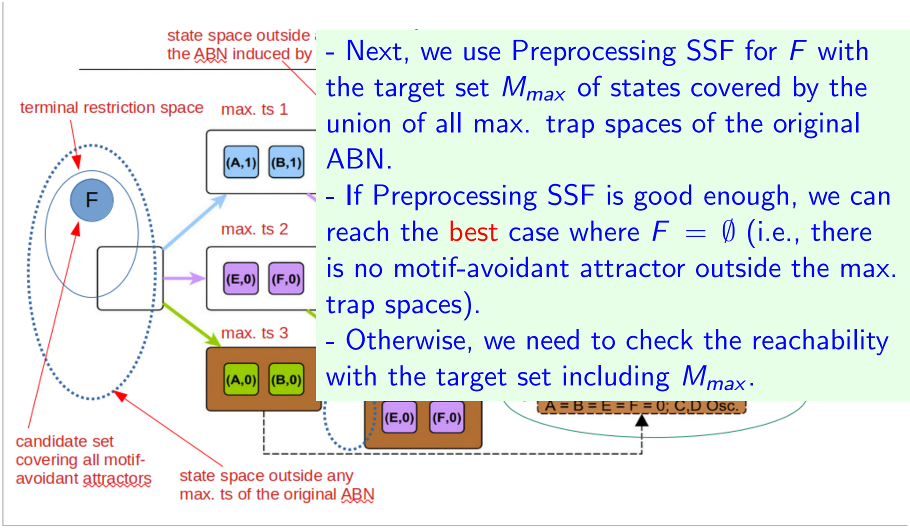
General approach



state space outside
the ABN induced by
terminal restriction space
max. ts 1
(A,1)
(B,1)
F
max. ts 2
(E,0)
(F,0)
max. ts 3
(A,0)
(B,0)
A = B = E = F = 0; C,D Osc.
(E,0)
(F,0)
candidate set covering all motif-avoidant attractors
state space outside any max. ts of the original ABN

- We need to check the existence of motif-avoidant attractors.
- We compute the terminal restriction space R.
- We compute the candidate set F covering all motif-avoidant attractors. Each state in F must be in R and has no out-going successor after systematically removing some of its successors according to the negative feedback vertex set of the original ABN.
- F may be small even when the NFVS is large.

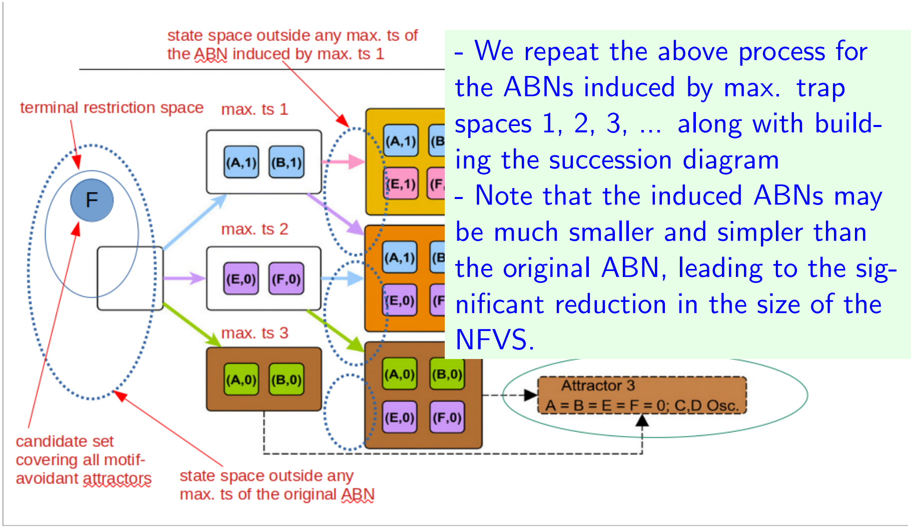
General approach



state space outside
the ABN induced by
terminal restriction space
max. ts 1
F
(A,1)
(B,1)
max. ts 2
(E,0)
(F,0)
max. ts 3
(A,0)
(B,0)
(E,0)
(F,0)
A = B = E = F = 0; C,D Osc.
candidate set covering all motif-avoidant attractors
state space outside any max. ts of the original ABN

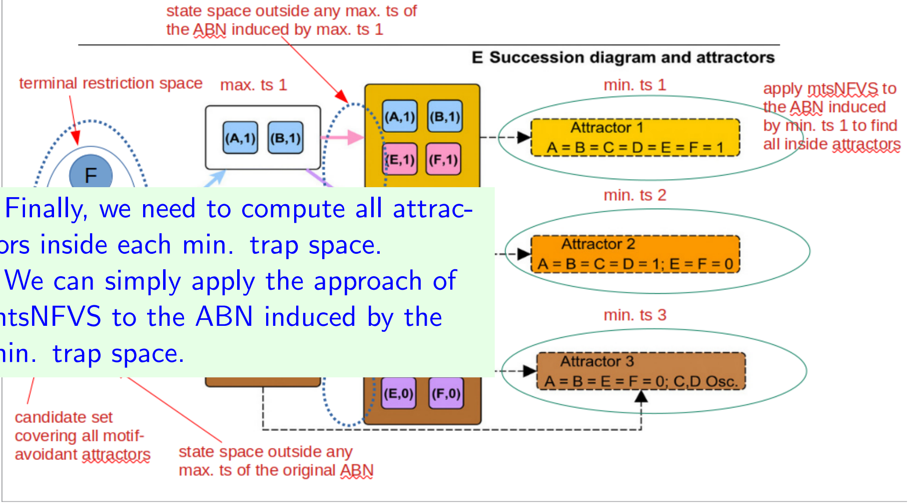
- Next, we use Preprocessing SSF for F with the target set M_{max} of states covered by the union of all max. trap spaces of the original ABN.
- If Preprocessing SSF is good enough, we can reach the best case where $F = \emptyset$ (i.e., there is no motif-avoidant attractor outside the max. trap spaces).
- Otherwise, we need to check the reachability with the target set including M_{max}.

General approach


state space outside any max. ts of the ABN induced by max. ts 1
terminal restriction space
max. ts 1
max. ts 2
max. ts 3
F
(A,1)
(B,1)
(E,0)
(F,0)
(A,0)
(B,0)
(E,1)
Attractor 3
A = B = E = F = 0; C,D Osc.
candidate set covering all motif-avoidant attractors
state space outside any max. ts of the original ABN

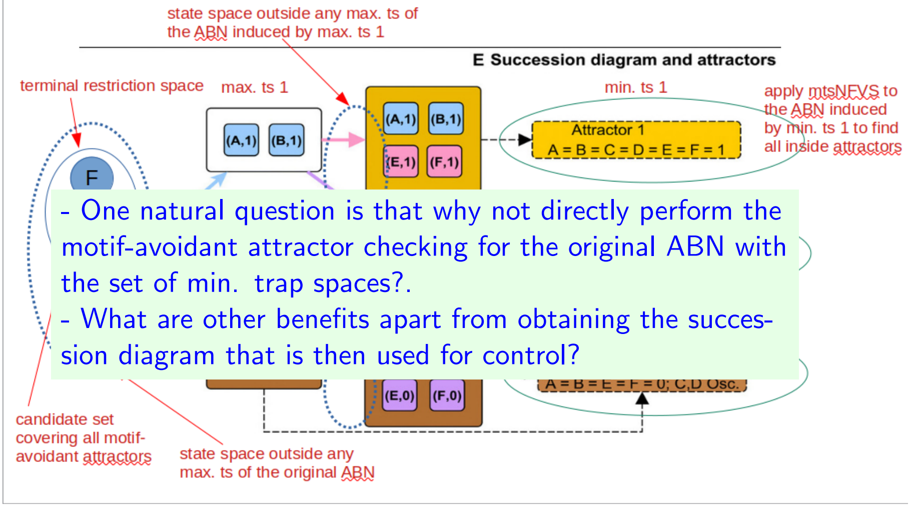
- We repeat the above process for the ABNs induced by max. trap spaces 1, 2, 3, ... along with building the succession diagram
- Note that the induced ABNs may be much smaller and simpler than the original ABN, leading to the significant reduction in the size of the NFVS.

General approach


state space outside any max. ts of the ABN induced by max. ts 1
E Succession diagram and attractors
terminal restriction space
max. ts 1
min. ts 1
apply mtsNFVS to the ABN induced by min. ts 1 to find all inside attractors
(A,1)
(B,1)
(E,1)
(F,1)
F
Attractor 1
A = B = C = D = E = F = 1
min. ts 2
Attractor 2
A = B = C = D = 1; E = F = 0
min. ts 3
Attractor 3
A = B = E = F = 0; C,D Osc.
(E,0)
(F,0)
candidate set covering all motif-avoidant attractors
state space outside any max. ts of the original ABN

- Finally, we need to compute all attractors inside each min. trap space.
- We can simply apply the approach of mtsNFVS to the ABN induced by the min. trap space.

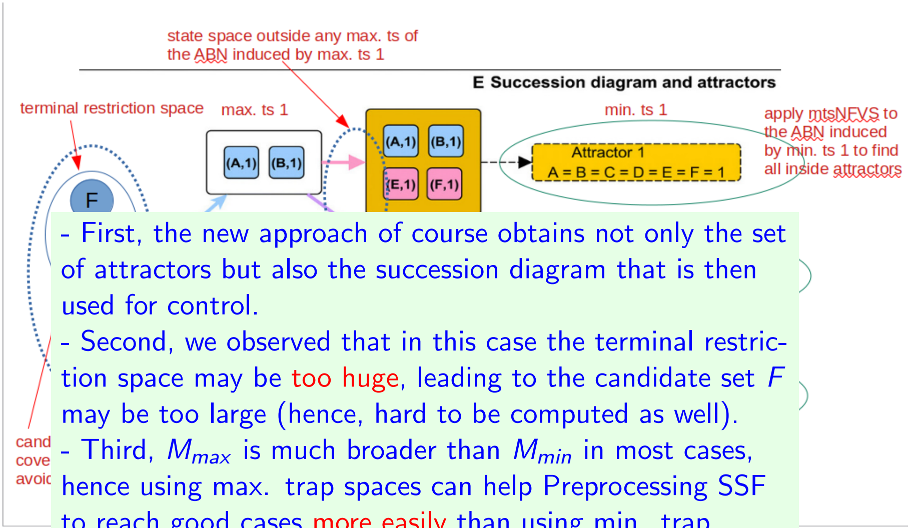
General approach



state space outside any max. ts of the ABN induced by max. ts 1
E Succession diagram and attractors
terminal restriction space
max. ts 1
min. ts 1
(A,1)
(B,1)
(E,1)
(F,1)
F
Attractor 1
A = B = C = D = E = F = 1
apply mtsNFVS to the ABN induced by min. ts 1 to find all inside attractors
A = B = E = F = 0; C,D Osc.
(E,0)
(F,0)
candidate set covering all motif-avoidant attractors
state space outside any max. ts of the original ABN

- One natural question is that why not directly perform the motif-avoidant attractor checking for the original ABN with the set of min. trap spaces?.
- What are other benefits apart from obtaining the succession diagram that is then used for control?

General approach


state space outside any max. ts of
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terminal restriction space
max. ts 1
(A,1)
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(F,1)
min. ts 1
Attractor 1
A = B = C = D = E = F = 1
apply mtsNFVS to
the ABN induced
by min. ts 1 to find
all inside attractors
F
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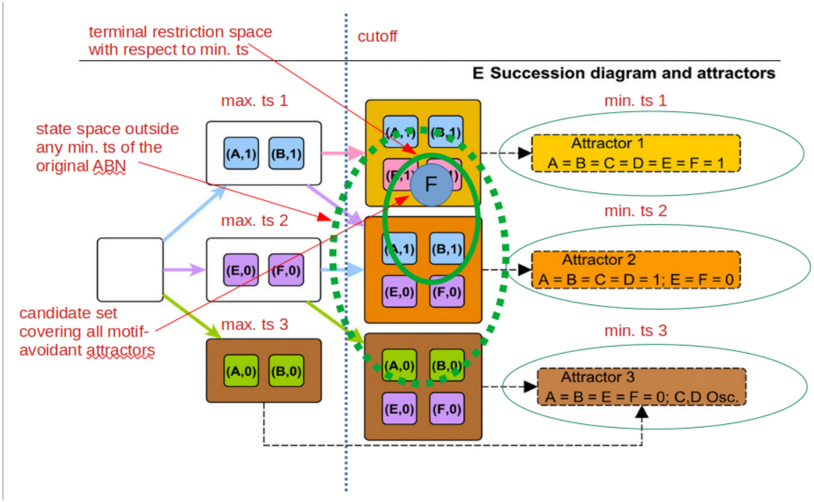
- First, the new approach of course obtains not only the set of attractors but also the succession diagram that is then used for control.
- Second, we observed that in this case the terminal restriction space may be too huge, leading to the candidate set F may be too large (hence, hard to be computed as well).
- Third, M_{max} is much broader than M_{min} in most cases, hence using max. trap spaces can help Preprocessing SSF to reach good cases more easily than using min. trap spaces.

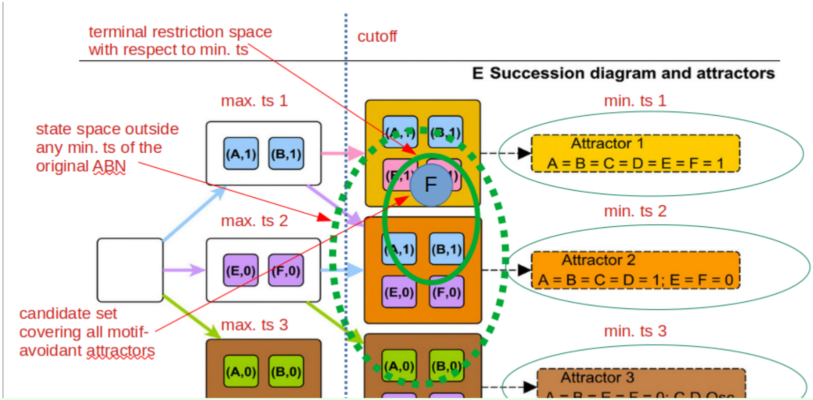
Deal with the problem of large succession diagrams

The above approach is promising. However, how to deal with the problem of large succession diagrams?.



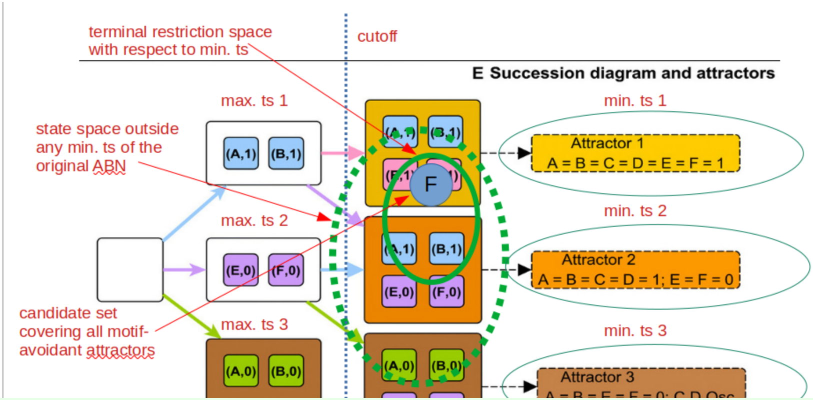
terminal restriction space with respect to min. ts
cutoff
E Succession diagram and attractors
max. ts 1
min. ts 1
state space outside any min. ts of the original ABN
(A,1)
(B,1)
(F,1)
F
Attractor 1
A = B = C = D = E = F = 1
max. ts 2
min. ts 2
(E,0)
(F,0)
Attractor 2
A = B = C = D = 1; E = F = 0
candidate set covering all motif-avoidant attractors
max. ts 3
min. ts 3
(A,0)
(B,0)
Attractor 3
A = B = E = F = 0; C,D Osc.

Deal with the problem of large succession diagrams



- When the size of the succession diagram becomes too large (maybe exceeding a threshold), we stop its construction.
- However, we still need to check the existence of motif-avoidant attractors.

Deal with the problem of large succession diagrams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F
Attractor 1
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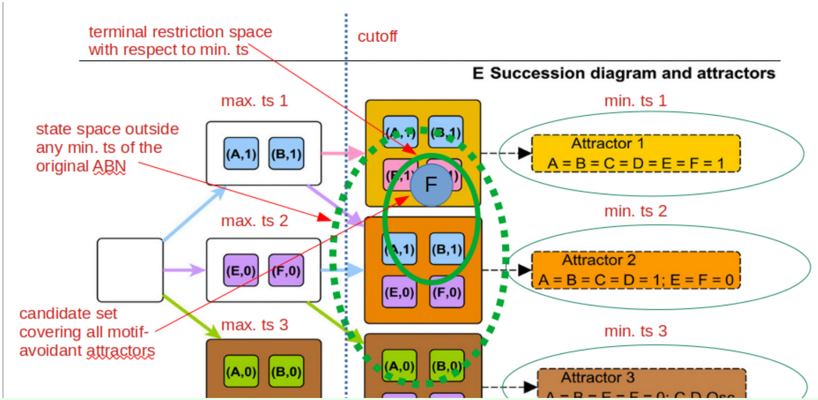
min. ts 2
(E,0)
(F,0)
Attractor 2
A = B = C = D = 1; E = F = 0
candidate set covering all motif-avoidant attractors
max. ts 3
min. ts 3
(A,0)
(B,0)
Attractor 3
A = B = E = F = 0; C,D Osc

- First, we compute all min. trap spaces of the original ABN.
- Second, we compute the terminal restriction space R with respect to these min. trap spaces.
- Next, we get the intersection between R and the state spaces induced by the already computed max. trap spaces (e.g., max. ts 1, 2, 3).

Deal with the problem of large succession diagrams



terminal restriction space with respect to min. ts
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max. ts 1
min. ts 1
state space outside any min. ts of the original ABN
(A,1)
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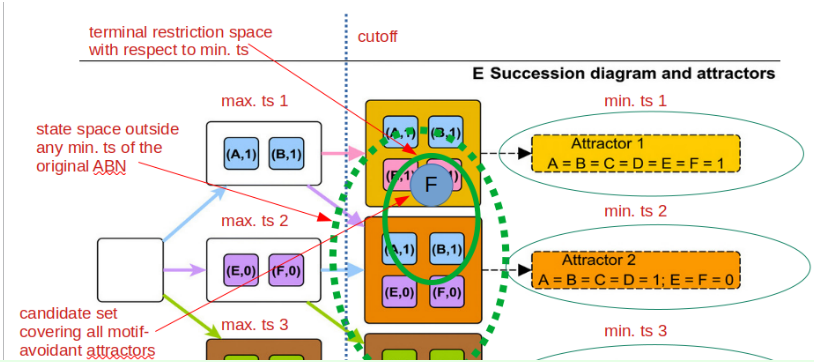
(F,1)
F
Attractor 1
A = B = C = D = E = F = 1
max. ts 2
min. ts 2
(E,0)
(F,0)
Attractor 2
A = B = C = D = 1; E = F = 0
candidate set covering all motif-avoidant attractors
max. ts 3
min. ts 3
(A,0)
(B,0)
Attractor 3

- Then, we compute the candidate set F with respect to the NFVS of the original ABN.
- Each state in F must be in the intersected set obtained beforehand.
- F may be small even when the NFVS is large because the intersected set may be not too large.

Deal with the problem of large succession diagrams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F
Attractor 1
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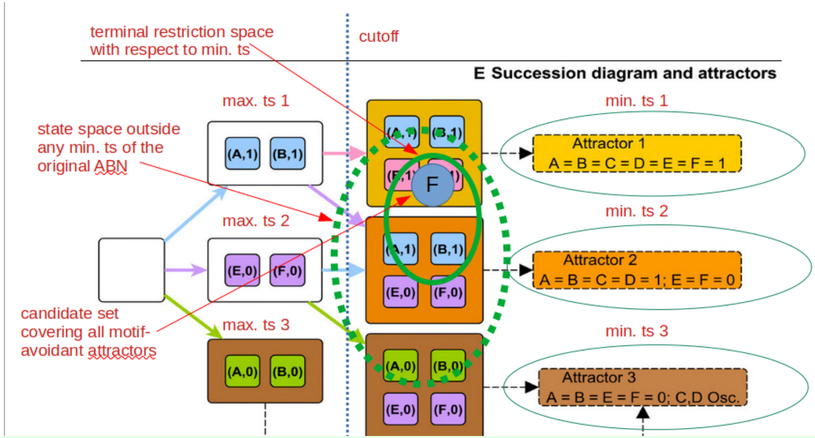
max. ts 2
min. ts 2
(E,0)
(F,0)
Attractor 2
A = B = C = D = 1; E = F = 0
candidate set covering all motif-avoidant attractors
max. ts 3
min. ts 3

Next, we use Preprocessing SSF for F with the target set M_{min} of states covered by the union of all min. trap spaces of the original ABN.
- If Preprocessing SSF is good enough, we can reach the best case where $F = \emptyset$ (i.e., there is no motif-avoidant attractor).
- Otherwise, we need to check the reachability with the target set including M_{min}.

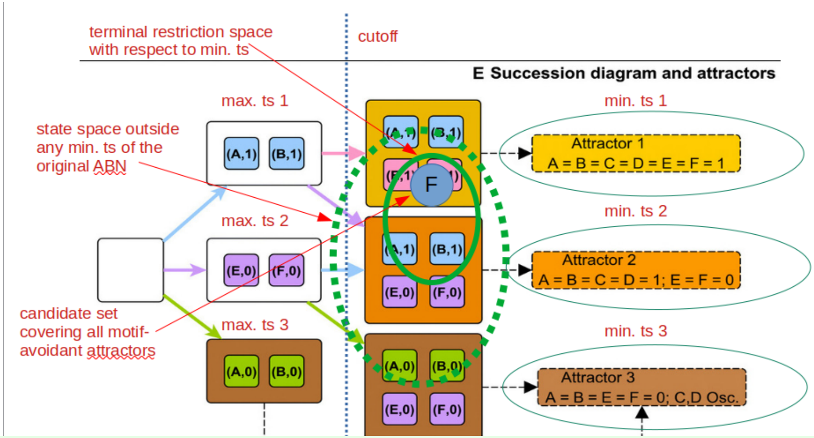
Deal with the problem of large succession diagrams



terminal restriction space with respect to min. ts
cutoff
E Succession diagram and attractors
max. ts 1
state space outside any min. ts of the original ABN
(A,1)
(B,1)
max. ts 2
(E,0)
(F,0)
candidate set covering all motif-avoidant attractors
max. ts 3
(A,0)
(B,0)
F
min. ts 1
Attractor 1
A = B = C = D = E = F = 1
min. ts 2
Attractor 2
A = B = C = D = 1; E = F = 0
min. ts 3
Attractor 3
A = B = E = F = 0; C,D Osc.

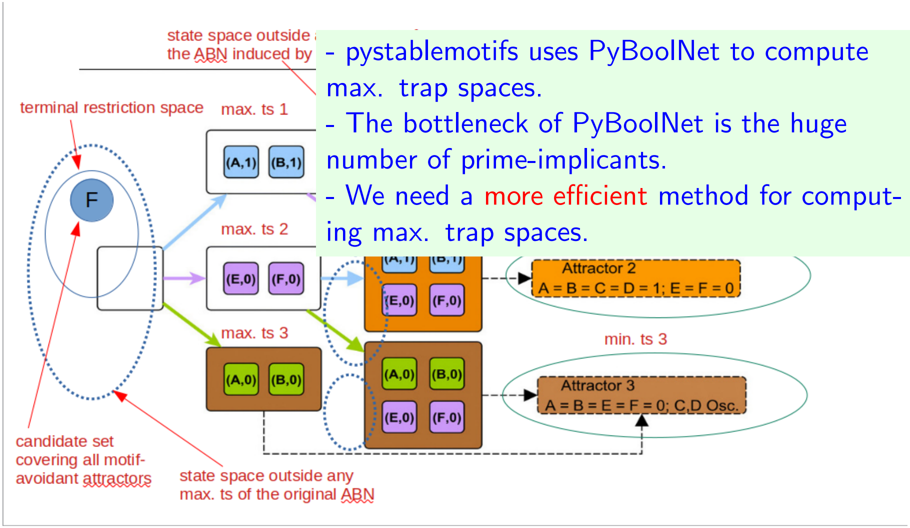
- Finally, we need to compute all attractors inside each min. trap space.
- We can simply apply the approach of mtsNFVS to the ABN induced by the min. trap space.

Deal with the problem of large succession diagrams



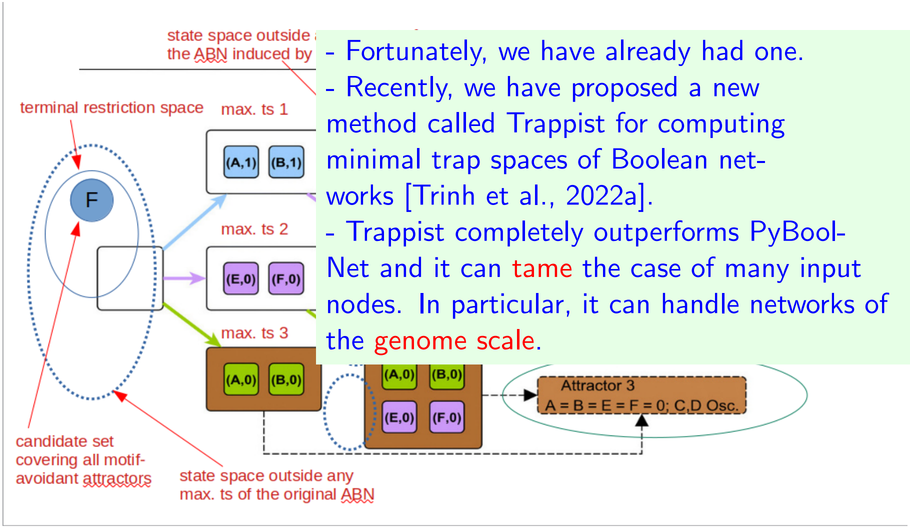
- After finishing the above process, we will obtain the whole attractor landscape and the partial succession diagram.

Computation of max. trap spaces


state space outside
the ABN induced by
terminal restriction space
max. ts 1
(A,1)
(B,1)
F
max. ts 2
(E,0)
(F,0)
max. ts 3
(A,0)
(B,0)
(A,1)
(B,1)
(E,0)
(F,0)
Attractor 2
A = B = C = D = 1; E = F = 0
min. ts 3
(A,0)
(B,0)
(E,0)
(F,0)
Attractor 3
A = B = E = F = 0; C,D Osc.
candidate set covering all motif-avoidant attractors
state space outside any max. ts of the original ABN

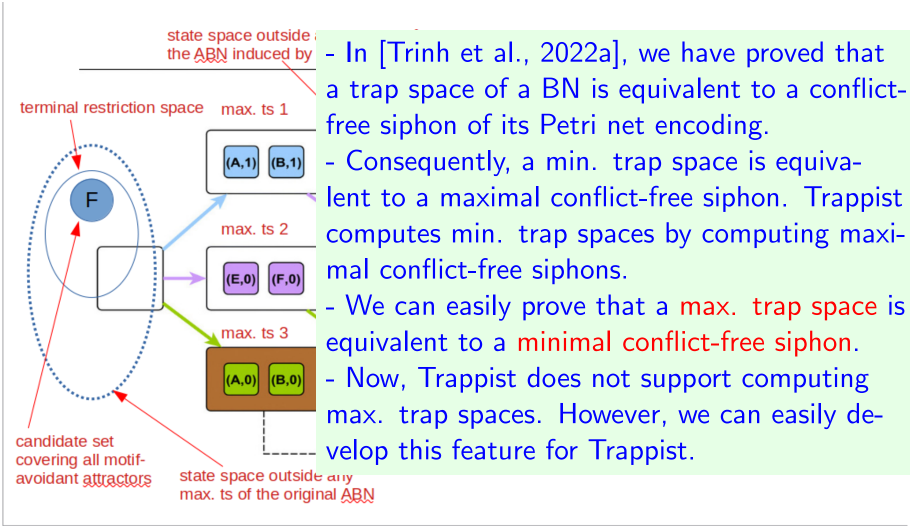
- pystablemotifs uses PyBoolNet to compute max. trap spaces.
- The bottleneck of PyBoolNet is the huge number of prime-implicants.
- We need a more efficient method for computing max. trap spaces.

Computation of max. trap spaces



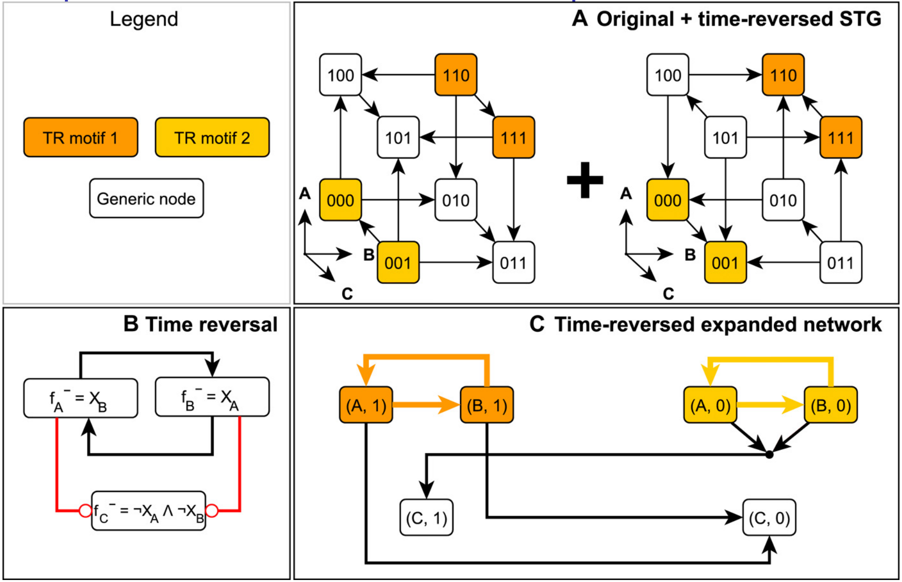
- Fortunately, we have already had one.
- Recently, we have proposed a new method called Trappist for computing minimal trap spaces of Boolean networks [Trinh et al., 2022a].
- Trappist completely outperforms PyBoolNet and it can tame the case of many input nodes. In particular, it can handle networks of the genome scale.

Computation of max. trap spaces

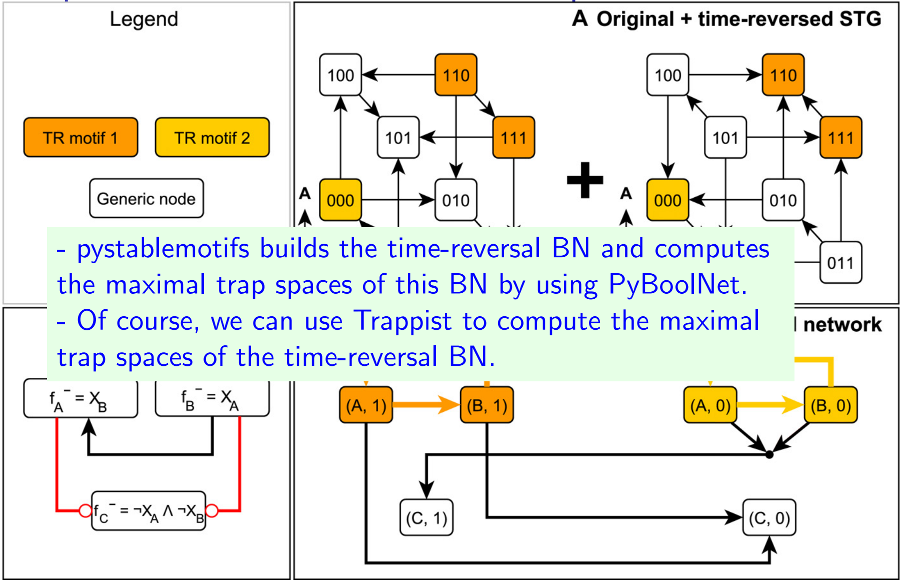

state space outside
the ABN induced by
terminal restriction space
max. ts 1
(A,1)
(B,1)
F
max. ts 2
(E,0)
(F,0)
max. ts 3
(A,0)
(B,0)
candidate set covering all motif-avoidant attractors
state space outside any max. ts of the original ABN

- In [Trinh et al., 2022a], we have proved that a trap space of a BN is equivalent to a conflict-free siphon of its Petri net encoding.
- Consequently, a min. trap space is equivalent to a maximal conflict-free siphon. Trappist computes min. trap spaces by computing maximal conflict-free siphons.
- We can easily prove that a max. trap space is equivalent to a minimal conflict-free siphon.
- Now, Trappist does not support computing max. trap spaces. However, we can easily develop this feature for Trappist.

Computation of terminal restriction spaces

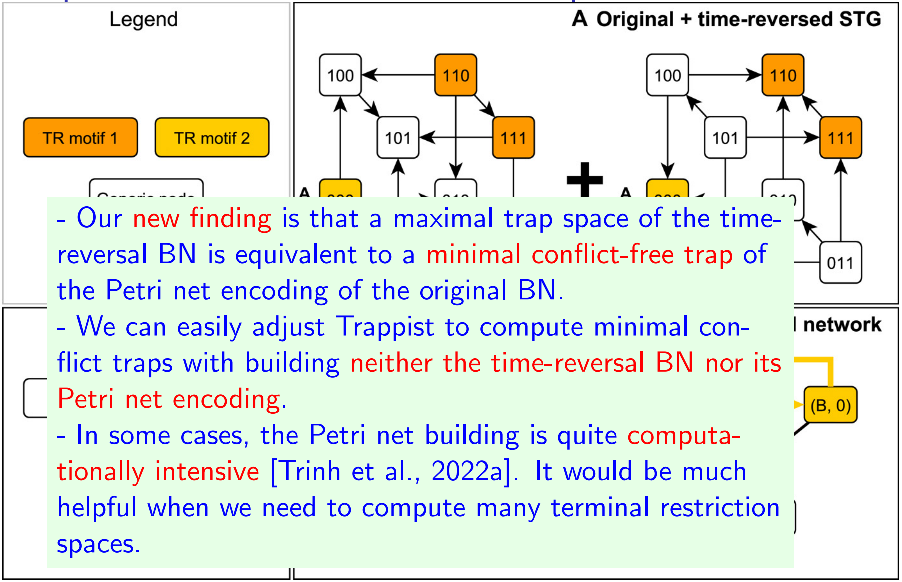


Computation of terminal restriction spaces



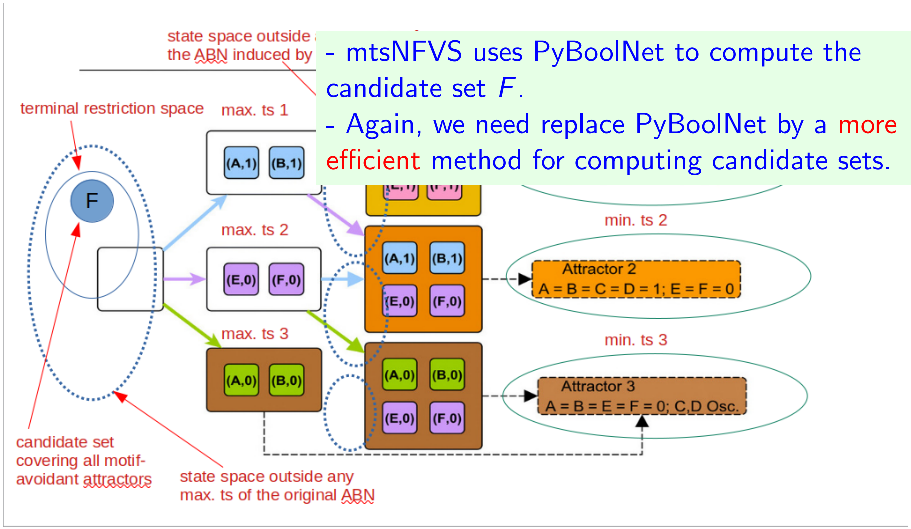
- pystablemotifs builds the time-reversal BN and computes the maximal trap spaces of this BN by using PyBoolNet.
- Of course, we can use Trappist to compute the maximal trap spaces of the time-reversal BN.

Computation of terminal restriction spaces



- Our new finding is that a maximal trap space of the time-reversal BN is equivalent to a minimal conflict-free trap of the Petri net encoding of the original BN.
- We can easily adjust Trappist to compute minimal conflict traps with building neither the time-reversal BN nor its Petri net encoding.
- In some cases, the Petri net building is quite computationally intensive [Trinh et al., 2022a]. It would be much helpful when we need to compute many terminal restriction spaces.

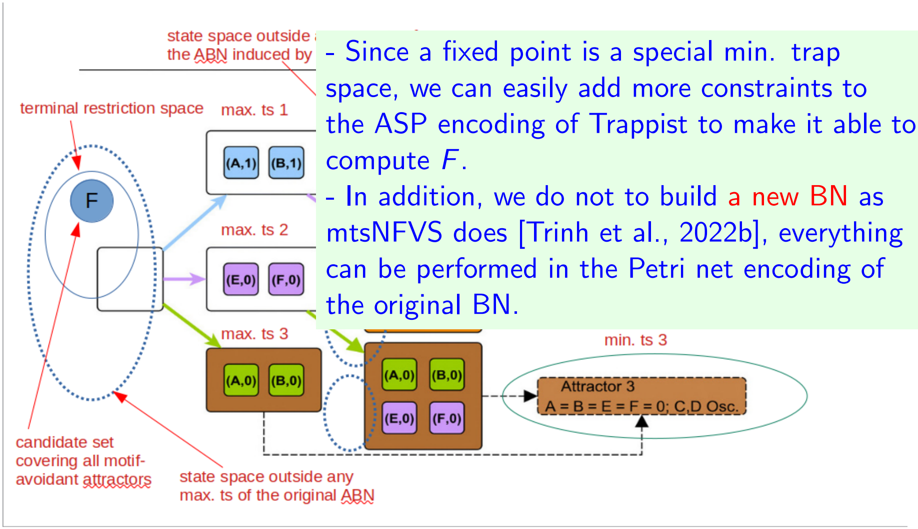
Computation of candidate sets



state space outside
the ABN induced by
terminal restriction space
max. ts 1
(A,1)
(B,1)
F
max. ts 2
(E,0)
(F,0)
max. ts 3
(A,0)
(B,0)
(A,1)
(B,1)
(E,0)
(F,0)
(A,0)
(B,0)
(E,0)
(F,0)
min. ts 2
Attractor 2
A = B = C = D = 1; E = F = 0
min. ts 3
Attractor 3
A = B = E = F = 0; C,D Osc.
candidate set covering all motif-avoidant attractors
state space outside any max. ts of the original ABN

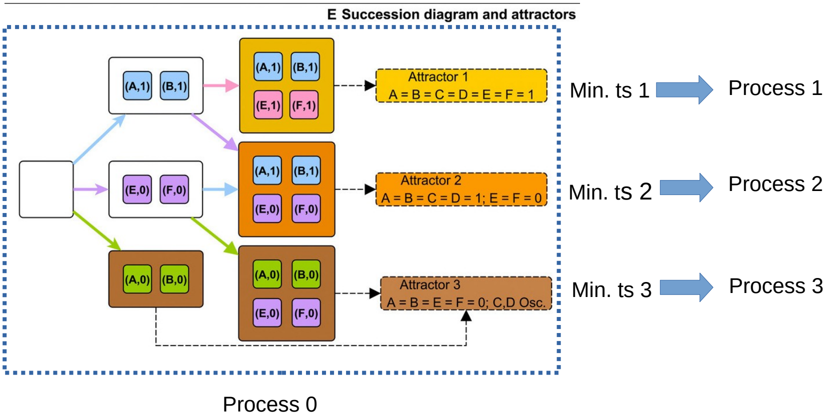
- mtsNFVS uses PyBoolNet to compute the candidate set F.
- Again, we need replace PyBoolNet by a more efficient method for computing candidate sets.

Computation of candidate sets

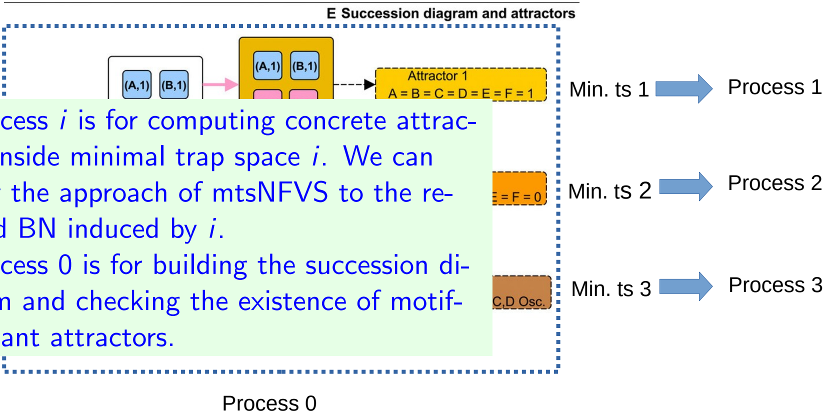

state space outside
the ABN induced by
terminal restriction space
max. ts 1
(A,1)
(B,1)
F
max. ts 2
(E,0)
(F,0)
max. ts 3
(A,0)
(B,0)
(A,0)
(B,0)
(E,0)
(F,0)
min. ts 3
Attractor 3
A = B = E = F = 0; C,D Osc.
candidate set covering all motif-avoidant attractors
state space outside any max. ts of the original ABN

- Since a fixed point is a special min. trap space, we can easily add more constraints to the ASP encoding of Trappist to make it able to compute F.
- In addition, we do not to build a new BN as mtsNFVS does [Trinh et al., 2022b], everything can be performed in the Petri net encoding of the original BN.

Parallelization

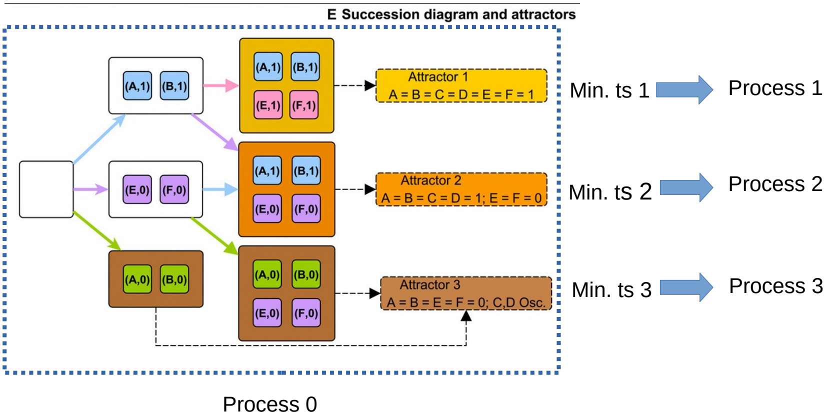


Parallelization



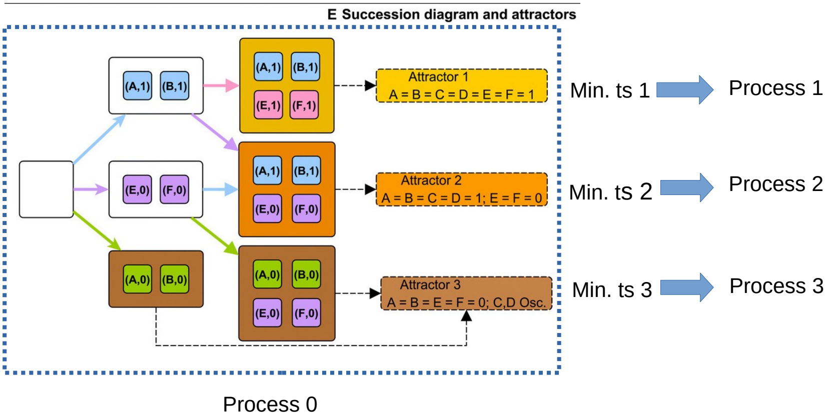
- Process *i* is for computing concrete attractors inside minimal trap space *i*. We can apply the approach of mtsNFVS to the reduced BN induced by *i*.
- Process 0 is for building the succession diagram and checking the existence of motif-avoidant attractors.

Parallelization



- Processes 0, 1, 2, ... can be completely paralleled.

Parallelization



- Process 0 can even be divided into many paralleled sub-processes.

Target control

After finishing the attractor detection phase, we will obtain the whole attractor landscape as well as the partial/full succession diagram. If the succession diagram is full, its size should be moderate.

The next step is how to compute minimum control policies that drive the network dynamics into a target attractor/min. trap space from any initial state?

There are some existing methods for the target control of asynchronous Boolean networks: CABEAN [Su and Pang, 2021a], PyBoolNet [Fontanals et al., 2020, Cifuentes-Fontanals et al., 2021], pystablemotifs [Zañudo and Albert, 2015, Rozum et al., 2021a], AEON [Brim et al., 2021].

Target control: our observations I

CABEAN always ensures that the resulting control policies are minimum. It relies on the calculation of strong basin of attraction, which may be very computationally expensive.

PyBoolNet exploits the information on trap spaces to find control policies. However, it does not ensure that the resulting control policies are minimum. It uses model checking to improve the quality of the control policies, which may be very computationally expensive as well. Moreover, its scalability may be limited by the necessity to scan through the available perturbations.

pystablemotifs uses the succession diagram to find control policies. However, it does not ensure that the resulting control policies are minimum.

Target control: our observations II

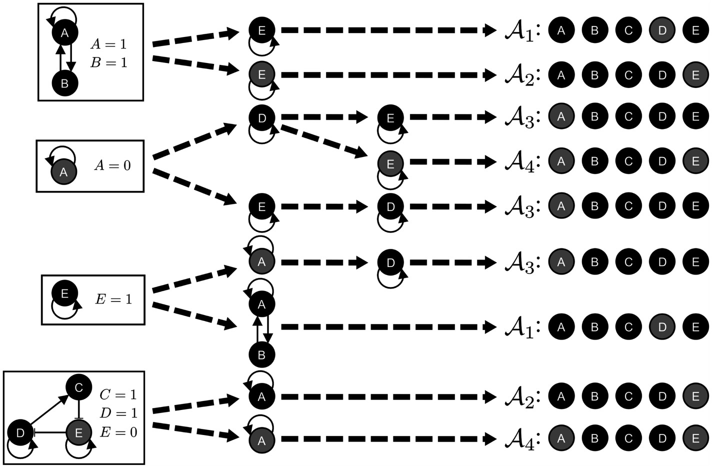
AEON uses methods based on the computation of the strong basin, but with symbolic representation to exhaustively analyse all perturbations simultaneously. A postprocessing step is then needed to extract the minimal perturbations from the final result. Alternatively, if the network has parameters, such postprocessing can also be used to extract perturbations with favourable robustness within the parameter space. The method is mainly limited by the complexity of the symbolic state space search for all available perturbations.

In the reported experiments [Su and Pang, 2021b, Su and Pang, 2021a], pystablemotifs is more time-efficient in most cases, though it did not return the minimum results in some cases.

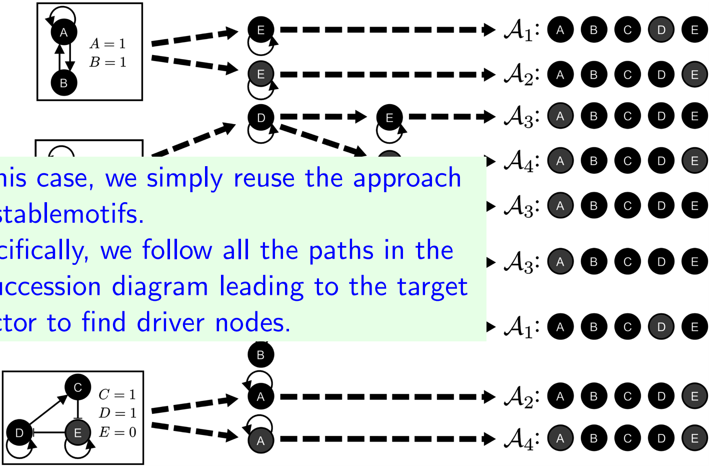
Target control: our observations III

In our opinion, pystablemotifs has much more potential than the others. We should dive into it. Of course, we need to improve both its time-efficiency and accuracy.

Target control with the full succession diagram



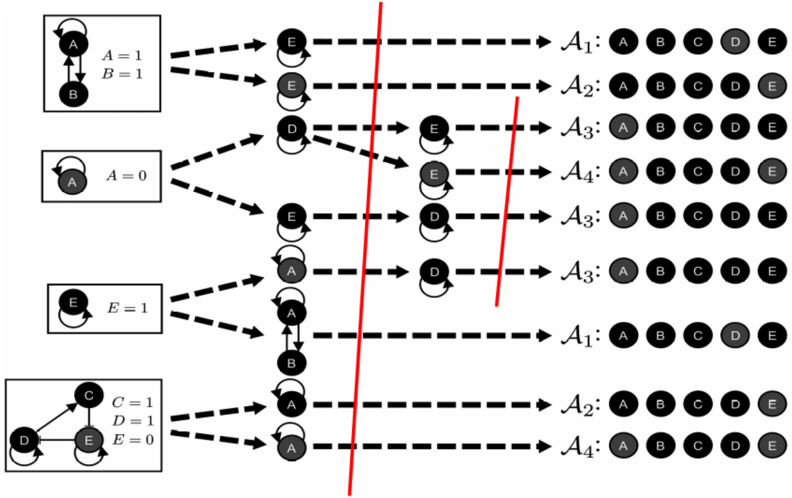
Target control with the full succession diagram



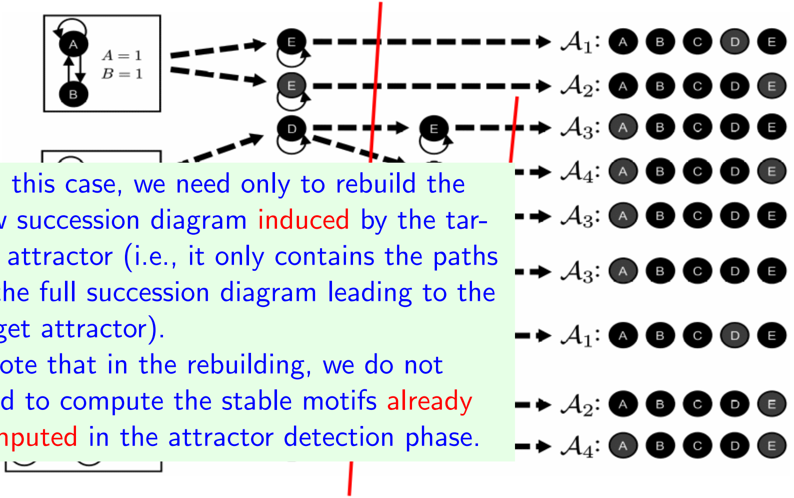
A = 1
B = 1
C = 1
D = 1
E = 0
$\mathcal{A}_1$: A B C D E
$\mathcal{A}_2$: A B C D E
$\mathcal{A}_3$: A B C D E
$\mathcal{A}_4$: A B C D E
$\mathcal{A}_3$: A B C D E
$\mathcal{A}_3$: A B C D E
$\mathcal{A}_1$: A B C D E
$\mathcal{A}_2$: A B C D E
$\mathcal{A}_4$: A B C D E

- In this case, we simply reuse the approach of pystablemotifs.
- Specifically, we follow all the paths in the full succession diagram leading to the target attractor to find driver nodes.

Target control with the partial succession diagram



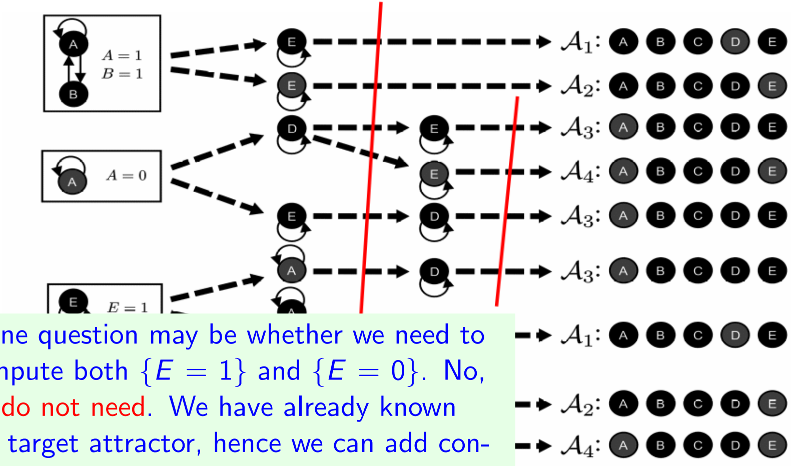
Target control with the partial succession diagram



$A = 1$
$B = 1$
A
B
E
D
$\mathcal{A}_1$: A B C D E
$\mathcal{A}_2$: A B C D E
$\mathcal{A}_3$: A B C D E
$\mathcal{A}_4$: A B C D E
$\mathcal{A}_3$: A B C D E
$\mathcal{A}_3$: A B C D E
$\mathcal{A}_1$: A B C D E
$\mathcal{A}_2$: A B C D E
$\mathcal{A}_4$: A B C D E

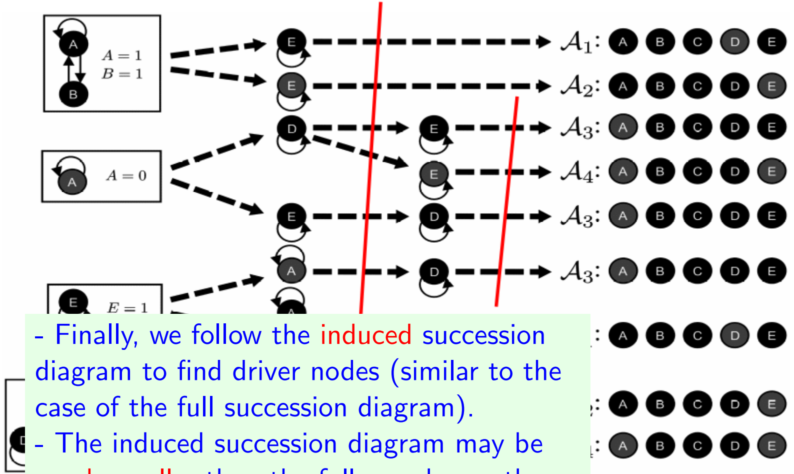
- In this case, we need only to rebuild the new succession diagram induced by the target attractor (i.e., it only contains the paths in the full succession diagram leading to the target attractor).
- Note that in the rebuilding, we do not need to compute the stable motifs already computed in the attractor detection phase.

Target control with the partial succession diagram



- One question may be whether we need to compute both $\{E = 1\}$ and $\{E = 0\}$. No, we do not need. We have already known the target attractor, hence we can add constraints to the ASP encoding to guide Trappist for searching only stable motifs that are compatible with the target attractor $\mathcal{A}_3$.

Target control with the partial succession diagram



- Finally, we follow the induced succession diagram to find driver nodes (similar to the case of the full succession diagram).
- The induced succession diagram may be much smaller than the full one, hence the computational burden may be reduced significantly.

Improve the accuracy

Consider the example asynchronous Boolean network.

$$\begin{cases} f_1 = \neg x_2 \lor (x_1 \land \neg x_3) \\ f_2 = (x_1 \land \neg x_3) \lor (x_2 \land x_3) \\ f_2 = (x_1 \land \neg x_3) \lor (x_2 \land x_3) \end{cases}$$

$$\begin{tikzcd}
 & 110 & & 111 \arrow[dd] \\
100 \arrow[ur] & & 101 \arrow[ll] & \\
 & 010 \arrow[dl] & & 011 \\
000 \arrow[uu] \arrow[rr] & & 001 \arrow[uu] &
\end{tikzcd}$$

Assume that the target attractor is $\{110\}$. The result of pystablemotifs is $\{x_1 = 1, x_3 = 0\}$. The minimum result should be $\{x_3 = 0\}$.

We have obtained a preliminary idea for improving the accuracy of pystablemotifs. However, we need to make it more specific as well as investigate its efficiency. In most cases, smaller control policies, longer computational time. We will present it to you later.

Bottlenecks of both pystablemotifs and mtsNFVS

The new approach still encounters the two bottlenecks of both pystablemotifs and mtsNFVS.

The case of many source nodes where there are actually too many attractors (at least 2^k where k is the number of source nodes). mtsNFVS is less affected than pystablemotifs, but this is still problem because mtsNFVS uses the set of explicit (not symbolic) min. trap spaces.

There exist some very huge attractors inside min. trap spaces. In this case, the number free variables of a min. trap space is still too large (even all the nodes of the original network). pystablemotifs only ends with min. trap spaces. It is very hard for mtsNFVS to reach the best case (i.e., $|F| = 1$) to avoid the reachability analysis.

We have obtained some preliminary ideas to mitigate the two above bottlenecks. We will present it to you later.

Implementation

We need to alter the code of pystablemotifs in deep.

The first thing needed to be done may be to integrate Trappist [2] to pystablemotifs. Since Trappist is written in Python, the integration should be easy.

We will discuss more about other implementation tasks.

[2] `https://github.com/soli/trap-spaces-as-siphons`

Experimentation

We plan to test the new method on models without constant nodes.

N-K models: expected to handle networks with $K = 2$ and up to 500 nodes.

scale-free models: expected to handle networks with $K = 10$ and up to 5000 nodes.

real-world models: expected to handle all the largest and most complex networks that we can find in the literature. The repository [3] maintained by Samuel Pastva is a good source.

[3] `https://github.com/sybila/biodivine-boolean-models`

Conclusion

The new approach exploits advantages of both pystablemotifs and mtsNFVS.

It benefits not only the attractor detection issue but also the control issue of asynchronous Boolean networks.

We believe that it will be the most promising approach that can reach the genome scale.

Collaboration

We hope we will have a nice collaboration on this work. Hence, we should make clear something at the beginning.

Participants:

- Our side: Van-Giang Trinh and Samuel Pastva.
- Your side: Jordan Rozum and ?

Authorship (if we can publish some papers :)):

- Two first authors, one from our side and one from your side with the equal contribution?

What do you think?

Next steps

We will prepare more specific slides (or a report) presenting technical details of the new approach.

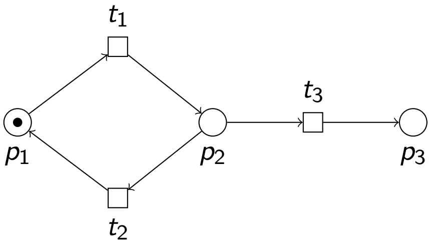
We should have more meetings to discuss in detail as well as make a specific plan.

Thank you for your attention!

Petri nets

Bipartite graph

Places P
Transitions T
Weighted arcs W



A *marking* for a Petri net is a mapping $m : P \mapsto \mathbb{N}$ that assigns a number of tokens to each place. A place p is marked by a marking m if and only if $m(p) > 0$. For example, p_1 is marked, p_2 and p_3 are unmarked.

We shall write $pred(x)$ (resp. $succ(x)$) to represent the set of vertices that have a (non-zero weighted) arc leading to (resp. coming from) x. For example, $pred(p_2) = \{t_1\}$, $succ(p_2) = \{t_2, t_3\}$, $pred(t_2) = \{p_2\}$, and $succ(t_2) = \{p_1\}$.

Siphons

Siphon

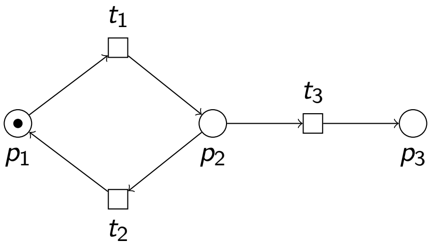
A *siphon* of a Petri net (P, T, W) is a set of places S such that:

$$\forall t \in T, S \cap succ(t) \neq \emptyset \Rightarrow S \cap pred(t) \neq \emptyset.$$

Once a siphon is unmarked, it remains unmarked.

$S = \{p_1, p_3\}$ is not a siphon because $S \cap succ(t_3) = \{p_1, p_3\} \cap \{p_3\} = \{p_3\} \neq \emptyset$ but $S \cap pred(t_3) = \{p_1, p_3\} \cap \{p_2\} = \emptyset$.

Here: $\emptyset$, $\{p_1, p_2\}$, $\{p_1, p_2, p_3\}$



Traps

Trap

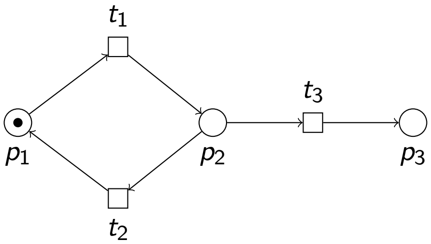
A *trap* of a Petri net (P, T, W) is a set of places S such that:

$$\forall t \in T, S \cap pred(t) \neq \emptyset \Rightarrow S \cap succ(t) \neq \emptyset.$$

Once a trap is marked (i.e., at least one of its places is marked), it remains marked.

$S = \{p_1, p_3\}$ is not a trap because $S \cap pred(t_1) = \{p_1, p_3\} \cap \{p_1\} = \{p_1\} \neq \emptyset$ but $S \cap succ(t_1) = \{p_1, p_3\} \cap \{p_2\} = \emptyset$.

Here: $\emptyset$, $\{p_1, p_2, p_3\}$



Petri net of a Boolean model

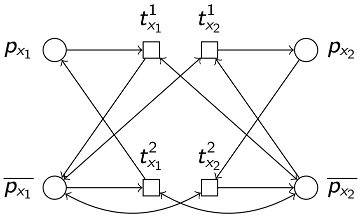
The original encoding was established in [Chaouiya et al., 2004].

Two places for each gene: $v \rightsquigarrow p_v, \overline{p_v}$

Solutions of $f_v \not\leftrightarrow v \rightsquigarrow$ transitions from p_v to $\overline{p_v}$ (and back)

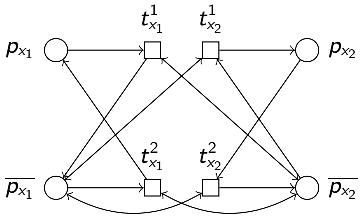
At any marking m of the Petri net encoding a Boolean model, it always holds that $m(p_v) + m(\overline{p}_v) = 1$.

$$\begin{cases} f_1 = (x_1 \wedge x_2) \vee (\neg x_1 \wedge \neg x_2) \\ f_2 = (x_1 \wedge x_2) \vee (\neg x_1 \wedge \neg x_2) \end{cases}$$



Conflict-free siphons/traps

A siphon/trap is called **conflict-free** if it does not contain both p_v and $\overline{p_v}$ for all $v \in V$.



Siphon	Conflict-free?
$\emptyset$	yes
$\{\overline{p_{x_1}}, \overline{p_{x_2}}\}$	yes
$\{p_{x_1}, \overline{p_{x_1}}\}$	no
$\{p_{x_2}, \overline{p_{x_2}}\}$	no

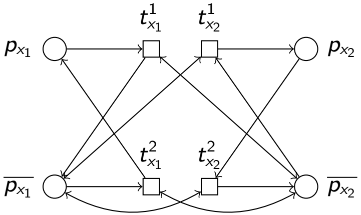
A conflict-free siphon (resp. trap) is *maximal* if it is not a subset of any other conflict-free siphon (resp. trap).

Conflict-free siphons are trap spaces

Theorem 1

Let $\mathcal{M}$ be a Boolean model and $\mathcal{P}$ be its Petri net encoding. There is a one-to-one correspondence between the set of **trap spaces** of $\mathcal{M}$ and the set of **conflict-free siphons** of $\mathcal{P}$.

$$\begin{cases} f_1 = (x_1 \wedge x_2) \vee (\neg x_1 \wedge \neg x_2) \\ f_2 = (x_1 \wedge x_2) \vee (\neg x_1 \wedge \neg x_2) \end{cases}$$



Trap space	Conflict-free siphon
$\star\star$	$\emptyset$
11	$\{\overline{p_{x_1}}, \overline{p_{x_2}}\}$

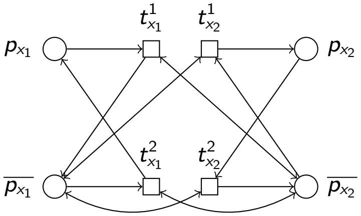
Maximal conflict-free siphons are minimal trap spaces

Theorem 2

Let $\mathcal{M}$ be a Boolean model and $\mathcal{P}$ be its Petri net encoding. There is a one-to-one correspondence between the set of **minimal trap spaces** of $\mathcal{M}$ and the set of **maximal conflict-free siphons** of $\mathcal{P}$.

$$\begin{cases} f_1 = (x_1 \land x_2) \lor (\neg x_1 \land \neg x_2) \\ f_2 = (x_1 \land x_2) \lor (\neg x_1 \land \neg x_2) \end{cases}$$

Trap space	Conflict-free siphon
$\star\star$	$\emptyset$
11	$\{\overline{p_{x_1}}, \overline{p_{x_2}}\}$



Proposed method for minimal trap space computation

From Theorem 2, we propose a new method for computing minimal trap spaces of a Boolean model $\mathcal{M}$.

- Build the Petri net encoding $\mathcal{P}$ of $\mathcal{M}$.
- Compute all maximal conflict-free siphons of $\mathcal{P}$.
- Convert the obtained maximal conflict-free siphons into the corresponding minimal trap spaces.

Petri net transformation

Transforming a Boolean model into its Petri net encoding can be done via computing Disjunctive Normal Forms (DNF) of each Boolean function [Chatain et al., 2014]. This transformation is implemented in the bioLQM[4] library using BDDs.

Though this might appear quite computationally intensive it is important to remark first that contrary to the prime-implicants case, there is no need to find *minimal* DNFs.

We use the above transformation in our proposed method.

[4] `http://www.colomoto.org/biolqm/`

Maximal conflict-free siphon computation

Characterize all siphons of the encoded Petri net as a system of Boolean rules.

$$p \in S \Rightarrow \bigvee_{p' \in pred(t)} p' \in S, p \in P, t \in T, t \in pred(p)$$

Add to the system the Boolean rules representing the conflict-freeness.

$$p_v \in S \Rightarrow \overline{p_v} \notin S \wedge \overline{p_v} \in S \Rightarrow p_v \notin S, v \in V$$

Encode the system as an ASP.

Use an ASP solver (e.g., clingo [Gebser et al., 2011]) to compute all set-inclusion maximal answer sets of the ASP.

Set-maximality through "heuristics"

```
clingo --heuristic=Domain --enum-mod=domRec --dom-mod=3
```

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