

Degrees of the logarithmic vector fields for close-to-free hyperplane arrangements

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§ Introduction

§ Line arrangement with cubic minimal log derivations.

§ Derivation Degree Sequences of Logarithmic Modules for close to Free Hyperplane Arrangements

<https://github.com/jcwjmz/LogarithmicDerivationModule>

or: <https://github.com/jcwjmz>

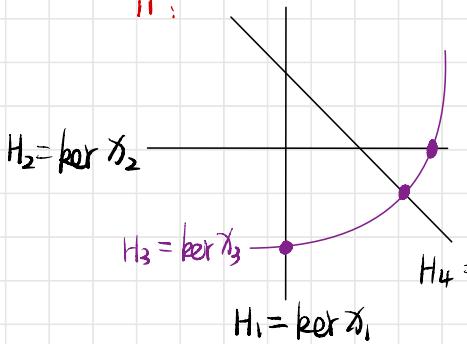
§ Introduction

1. \mathbb{K} = field. , $V = \mathbb{K}^l$ = vector space .
2. $S = \mathbb{K}[x_1, \dots, x_l]$ = the polynomial ring .
3. $\mathcal{A} := \{H_1, \dots, H_p\}$ = hyperplane arrangement
a finite set of linear hyperplanes $H_i = \ker \alpha_i \subset V$
4. $Q(\mathcal{A}) := \prod_{H \in \mathcal{A}} \alpha_H (\in S)$ = defining polynomial of \mathcal{A}

• Example :

Let $V = \underline{\mathbb{R}^3}$ and $\mathcal{A} = \{H_1, H_2, H_3, H_4\}$.
 \downarrow projectively

\mathbb{P}^2 :



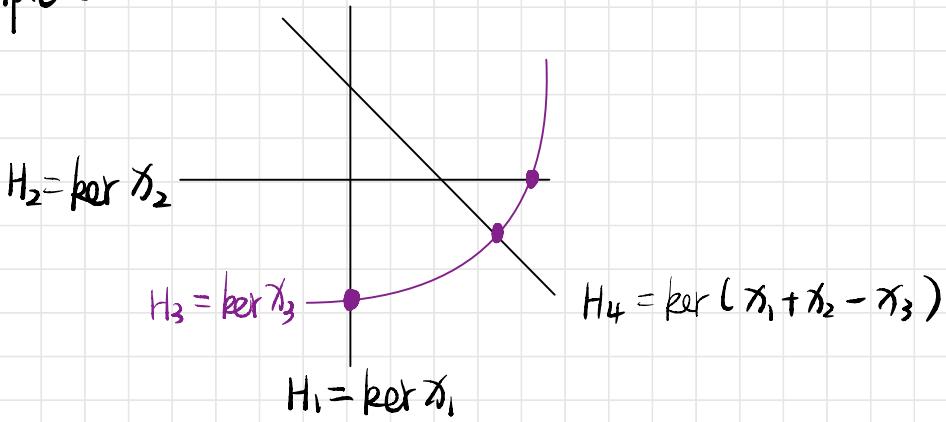
$$x_3 = 1$$

$$\left\{ \begin{array}{l} H_1 = \ker x_1 \\ H_2 = \ker x_2 \\ H_3 = \ker x_3 \\ H_4 = \ker (x_1 + x_2 - x_3) \end{array} \right.$$

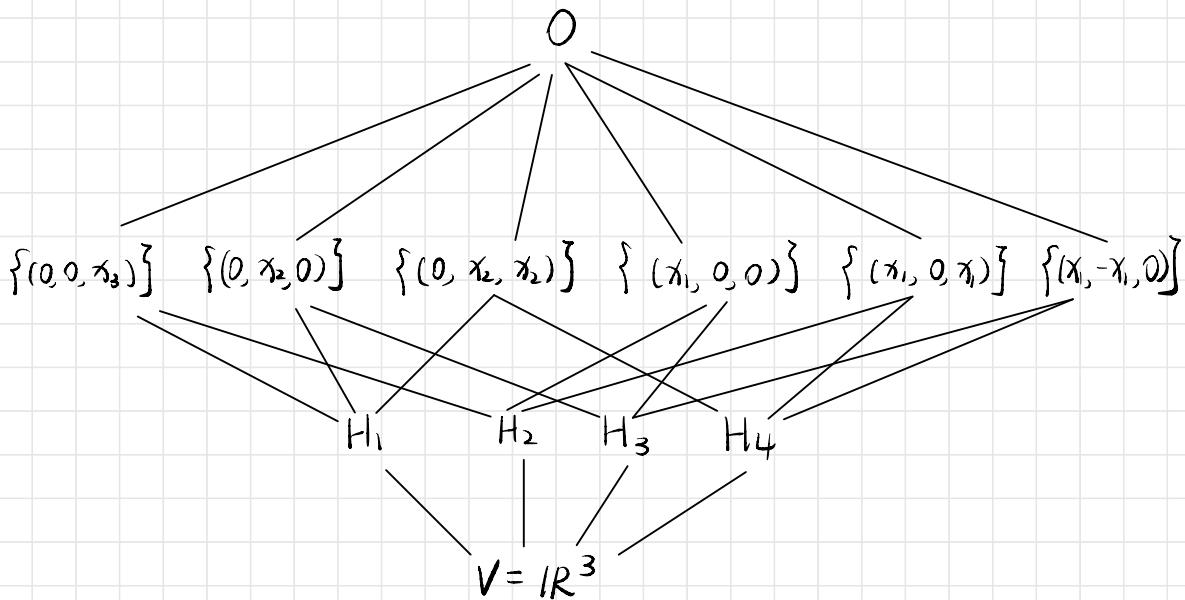
• Combinatorics of hyperplane arrangements.

- $L(\mathcal{A}) := \left\{ \bigcap_{H \in \mathcal{B}} H \mid \mathcal{B} \subset \mathcal{A} \right\}$ = intersection poset
partial order is reverse inclusion

Example:



$L(\mathcal{A})$:



- Some operations =

1. For $X \in L(\mathcal{A})$

(1) $\mathcal{A}^X := \{H \in \mathcal{A} \mid X \subset H\}$: Localization

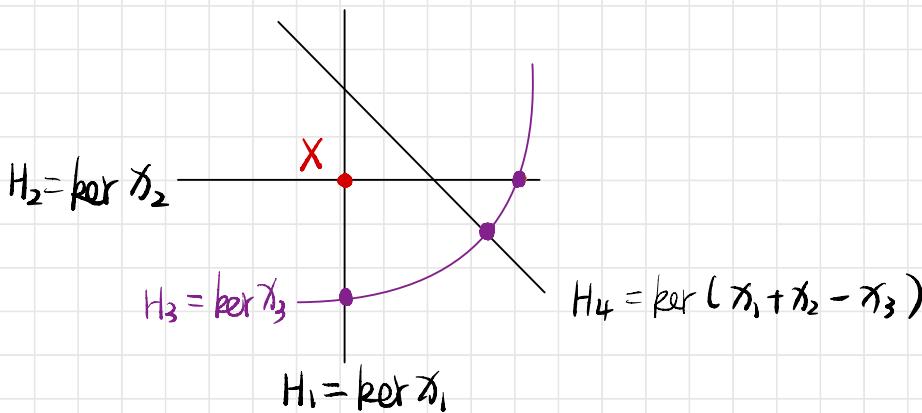
of \mathcal{A} to X . [\mathcal{A}^X is in $V/X = \mathbb{K}^{\text{codim } X}$]

(2) $\mathcal{A}^X := \{H \cap X \mid H \in \mathcal{A} \setminus \mathcal{A}^X\}$: restriction

of \mathcal{A} to X [\mathcal{A}^X is in $X = \mathbb{K}^{\dim X}$]

- Example =

Let $\ell = 3$ and $Q(\mathbf{d}) = x_1 x_2 x_3 (x_1 + x_2 - x_3)$



Then: $X = H_1 \cap H_2 \Rightarrow \mathcal{A}^X = \{H_1, H_2\}$

$$X = H_3 \Rightarrow \mathcal{A}^X = \{x_1 = 0, x_1 + x_2 = 0, x_2 = 0\}$$

● Algebraic invariant of hyperplane arrangements.

Def 1. (1) The free S -module of derivations:

$$\text{Der } S := \left\{ \sum_{i=1}^l f_i \partial_i \mid \partial_i := \frac{\partial}{\partial x_i}, f_i \in S, i=1, \dots, l \right\}$$

★ (2) The logarithmic derivation module $D(\mathbb{A})$:

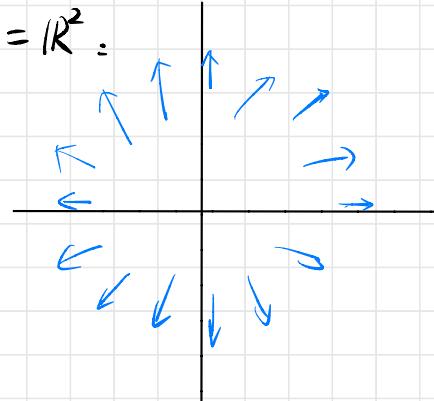
$$D(\mathbb{A}) := \{ \theta \in \text{Der } S \mid \theta(\alpha_H) \in S\alpha_H \ \forall H \in \mathbb{A} \},$$

(tangent condition)

Def 2: The Euler derivation $\theta_E \in \text{Der } S$ is:

$$\text{eg: } V = \mathbb{R}^2:$$

$$\theta_E = \sum_{i=1}^l x_i \partial_i$$



- For any hyperplane $H := \ker \alpha$ with $\alpha = \sum a_i x_i$

$$\theta_E(\alpha) = \sum a_i x_i = \alpha$$

$\Rightarrow \theta_E \in D(\mathcal{A})$ for any arrangement \mathcal{A} .

• Lemma = $D(\mathcal{A}) = S\theta_E \oplus D_H(\mathcal{A})$.

$$D_H(\mathcal{A}) = \left\{ \theta \in D(\mathcal{A}) \mid \theta(d_H) = 0 \right\}$$

\nwarrow
parallel to H

Def 3: We say \mathcal{A} is free with

$\exp(\mathcal{A}) = (d_1, d_2, \dots, d_e)$ if

$$D(\mathcal{A}) = S\theta_1 \oplus S\theta_2 \oplus \dots \oplus S\theta_e.$$

with $\deg \theta_i = d_i$ ($i=1, \dots, e$).

Examples.

coordinate planes

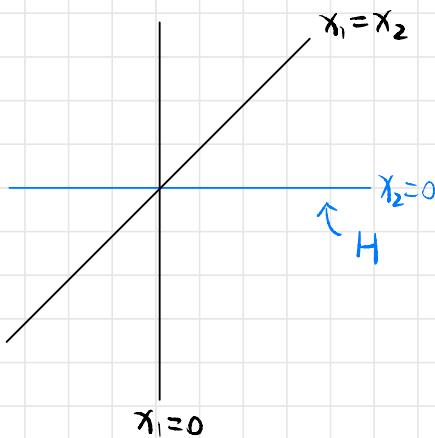
(1) Let $\mathcal{A} = \{H_i = \ker \chi_i \mid i=1, \dots, l\}$.

Then \mathcal{A} is free with $\exp(\mathcal{A}) = (1, \dots, 1)$.

In fact, $D(\mathcal{A}) = \bigoplus_{i=1}^l S\theta_i$, where $\theta_i = \chi_i d_i$.

(2) Any Δ in \mathbb{K}^2 is free with $\exp(\Delta) = (1, |\Delta| - 1)$

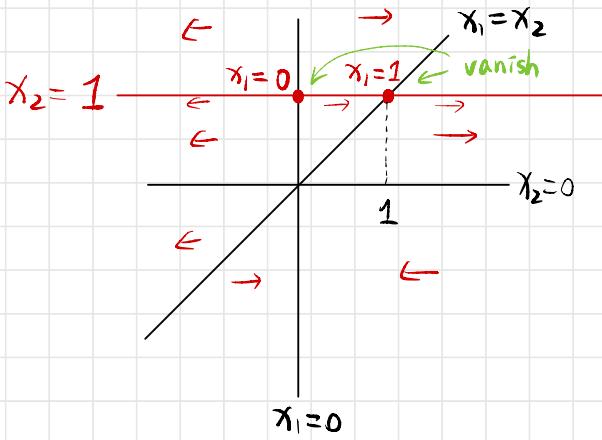
$$\text{e.g.: } Q(\Delta) = x_1 x_2 (x_1 - x_2)$$



$$\left\{ \begin{array}{l} D(\Delta) = S\theta_E \oplus D_H(\Delta) \\ \theta_E = x_1 \partial_1 + x_2 \partial_2 \end{array} \right.$$

Consider $\theta \in D_H(\Delta)$

$$\theta = f \partial_1, \quad f \in S.$$



$$(x_1 - 0)(x_1 - 1) \Big| f$$

$$\theta \in S(x_1, (x_1 - x_2) \partial_1)$$

$$\Downarrow \theta_2 := x_1(x_1 - x_2) \partial_1$$

$$D(\Delta) = S\theta_E \oplus S\theta_2$$

One of the most important conjecture in this field is Terao's Conjecture (1981).

- For arrangements \mathcal{A} and \mathcal{B} , if $L(\mathcal{A}) \cong L(\mathcal{B})$, then \mathcal{A} is free iff \mathcal{B} is free.

→ Example on SageMath

A deeper understanding of the generators of $D(\mathcal{A})$ is of utmost importance to tackle this conjecture

• Machinery to check freeness.

Theorem (Saito's criterion)

Let $\theta_1, \dots, \theta_e \in D(\mathcal{A})$ be homogeneous and linearly independent over S . Then \mathcal{A} is free with basis $\theta_1, \dots, \theta_e$ if and only if

$$\sum_{i=1}^e \deg \theta_i = |\mathcal{A}|.$$

Theorem (Addition-deletion Theorem, Terao, 1980)

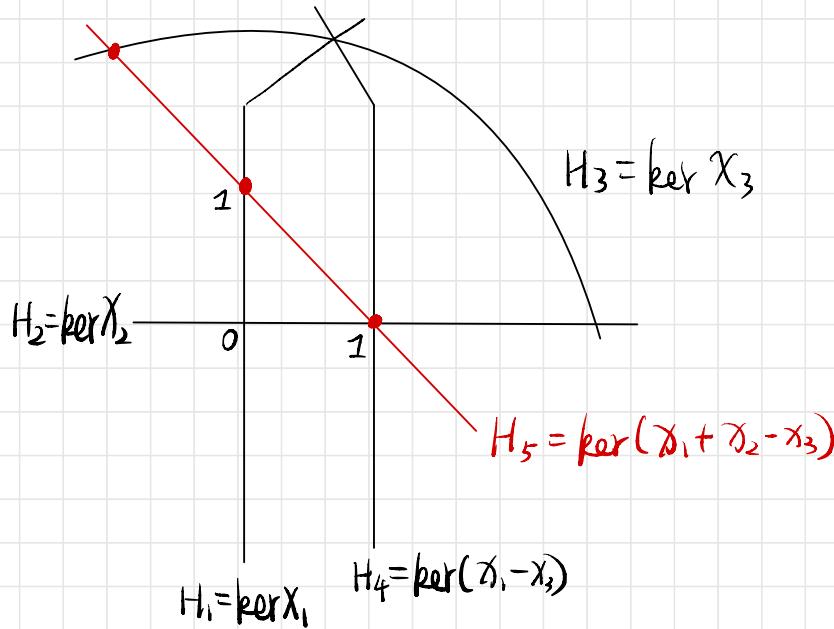
Let $H \in \mathcal{A}$, $\mathcal{A}' := \mathcal{A} \setminus \{H\}$ and $\mathcal{A}'' := \mathcal{A}^H$. Then two of the following imply the third:

- (1) \mathcal{A} is free with $\exp(\mathcal{A}) = (d_1, \dots, d_{e-1}, d_e)$
- (2) \mathcal{A}' is free with $\exp(\mathcal{A}') = (d_1, \dots, d_{e-1}, d_e - 1)$
- (3) \mathcal{A}'' is free with $\exp(\mathcal{A}'') = (d_1, \dots, d_{e-1})$

Moreover, all the three above hold true if \mathcal{A} and \mathcal{A}' are free.

Example:

Let $\ell = 3$ and $Q(\alpha) = \alpha_1 \alpha_2 \alpha_3 (\alpha_1 - \alpha_3)(\alpha_1 + \alpha_2 - \alpha_3)$



$\alpha' = \alpha \setminus \{H_5\}$ is free with $\exp(\alpha') = (1, 1, 2)$

$\alpha'' = \alpha^{H_5}$ is free with $\exp(\alpha'') = (1, 2)$

$\Rightarrow \alpha$ is free with $\exp(\alpha) = (1, 2, 2)$

§ Line arrangement with cubic minimal
log derivations.

The study of $D(\mathcal{A})$:

→ natural approach:

- the degrees of the gens.
- the min. generators.
- the min. free resolution
- :

Recall that,

$$D(\mathcal{A}) = S\theta_E \oplus D_H(\mathcal{A}), \quad \forall H \in \mathcal{A}.$$

Let $\{\theta_1, \theta_2, \dots, \theta_p \mid \theta_i \in D_A(A) \ \forall i=2, \dots, p\}$

is a minimal set of generators for $D(A)$.

The minimal degree of logarithmic derivations of arrangement A is defined as

$$\min \{ d_i \mid d_i = \deg \theta_i, i=2, \dots, p \}.$$

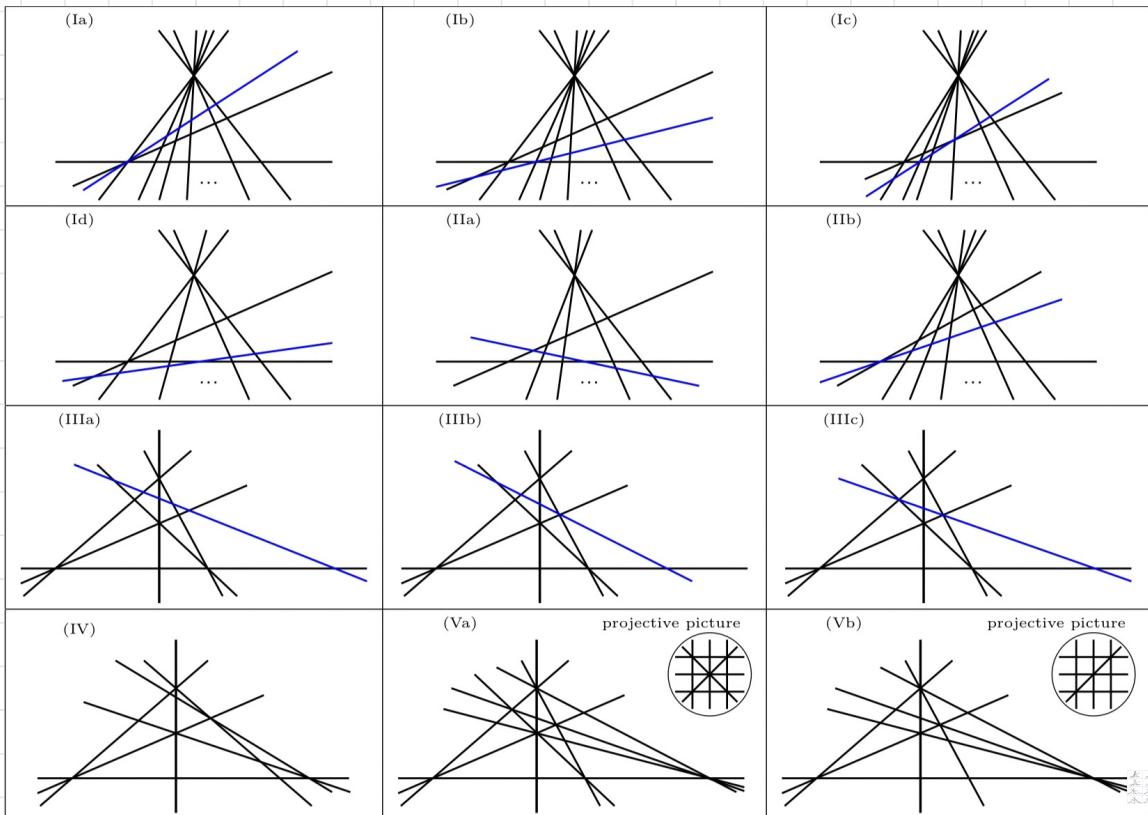
From this view,

Buriti and Tohăneanu give full classification of rank 3 line arrangements in P^2 that have a minimal log. derivation of degree 1, 2, 3.

Theorem [Burîty - Tohăneanu, 2021]:

Let $\mathcal{A} \subset \mathbb{P}^2$ be a line arrangement of rank 3, with $r(\mathcal{A}) = 3$. Then, up to a change of coordinates, \mathcal{A} has one of the following defining polynomials (and corresponding (affine) pictures):

- (Ia) $F = xyz(x+y)(bx+y) \prod_{j=6}^s (t_j y + z)$, $t_j \neq 0$, $s \geq 7$.
- (Ib) $F = xyz(x+y)(bx+z) \prod_{j=6}^s (t_j y + z)$, $t_j \neq 0$, $s \geq 6$.
- (Ic) $F = xyz(x+y)(y+z)(-x+z) \prod_{j=7}^s (t_j y + z)$, $t_j \neq 0, 1$, $s \geq 7$.
- (Id) $F = xyz(x+y)(ax+by+z) \prod_{j=6}^s (t_j y + z)$, $t_j \neq 0$, $s \geq 5$.
- (IIa) $F = xyz(x+y+z)(ax+by+z) \prod_{j=6}^s (t_j y + z)$, $t_j \neq 0$, $s \geq 5$.
- (IIb) $F = xyz(x+y+z)(bx+y) \prod_{j=6}^s (t_j y + z)$, $t_j \neq 0$, $s \geq 6$.
- (IIIa) $F = xyz(x+z)(y+z)(x+y+z)(ax+by+z)$.
- (IIIb) $F = xyz(x+z)(y+z)(x+y+z)(ax+ay+z)$.
- (IIIc) $F = xyz(x+z)(y+z)(x+y+z)(x+y+2z)$.
- (IV) $F = xy(x-z)(y-z)(x-2z)(y-2z)(x-y)$.
- (Va) $F = xyz(x^2-y^2)(x^2-z^2)(y^2-z^2)$.
- (Vb) $F = xyz(x^2-z^2)(y^2-z^2)(x-y)$.



Theorem [Chu]

Let $\mathcal{A} \subset \mathbb{P}^2$ be a line arrangement of rank 3, with $r(\mathcal{A}) = 3$. Then, up to a change of coordinates, \mathcal{A} has one of the following defining polynomials (and corresponding (affine) pictures):

- (Ia) $F = xyz(x+y)(bx+y)\prod_{j=6}^s(t_jy+z), t_j \neq 0, s \geq 7.$
- (Ib) $F = xyz(x+y)(bx+z)\prod_{j=6}^s(t_jy+z), t_j \neq 0, s \geq 6.$
- (Ic) $F = xyz(x+y)(y+z)(-x+z)\prod_{j=7}^s(t_jy+z), t_j \neq 0, 1, s \geq 7.$
- (Id) $F = xyz(x+y)(ax+by+z)\prod_{j=6}^s(t_jy+z), t_j \neq 0, s \geq 5.$
- (IIa) $F = xyz(x+y+z)(ax+by+z)\prod_{j=6}^s(t_jy+z), t_j \neq 0, s \geq 5.$
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- (Vb) $F = xyz(x^2-z^2)(y^2-z^2)(x-y).$

Table 1: Freeness of Theorem 2.1.

Cases	Freeness	Exponents
(Ia)	Free	$\{1, 3, s-4\}$
(Ib)	$t_i = -b$	$\{1, 3, s-4\}$
(Ic)	Free	$\{1, 3, s-4\}$
(Id)	$t_i = -b$ and $t_j = b-a$	$\{1, 3, s-4\}$
(IIa)	$t_i = 1$ and either $b = 1$ or $t_j = b, a = 1$ or $t_j = b, t_k = \frac{a-b}{a-1}$	$\{1, 3, s-4\}$
(IIb)	$t_i = 1$ and either $b = 1$ or $b(1-t_j) = 1$	$\{1, 3, s-4\}$
(IIIa)	either $a = 1$ or $b = 1$	$\{1, 3, 3\}$
(IIIb)	Not free	Non
(IIIc)	Free	$\{1, 3, 3\}$
(IV)	Not free	Non
(Va)	Free	$\{1, 3, 5\}$
(Vb)	Free	$\{1, 3, 4\}$

Theorem [C-]

Let $\ell=3$, $s \in \mathbb{N}$ and the defining polynomial of \mathcal{A} be:

$$Q(\mathcal{A}) = x_1 x_2 x_3 (ax_1 + bx_2 + x_3) \prod_{j=1}^s (t_j x_2 + x_3)$$

with $a, b \in \mathbb{K} \setminus 0$.

Then $[\mathcal{A} \text{ is free} \iff \exists! i \in \{1, \dots, s\} \text{ s.t. } b = t_i]$

Moreover, if \mathcal{A} is free, we have $\exp(\mathcal{A}) = (1, 2, s+1)$ and a basis is as follows:

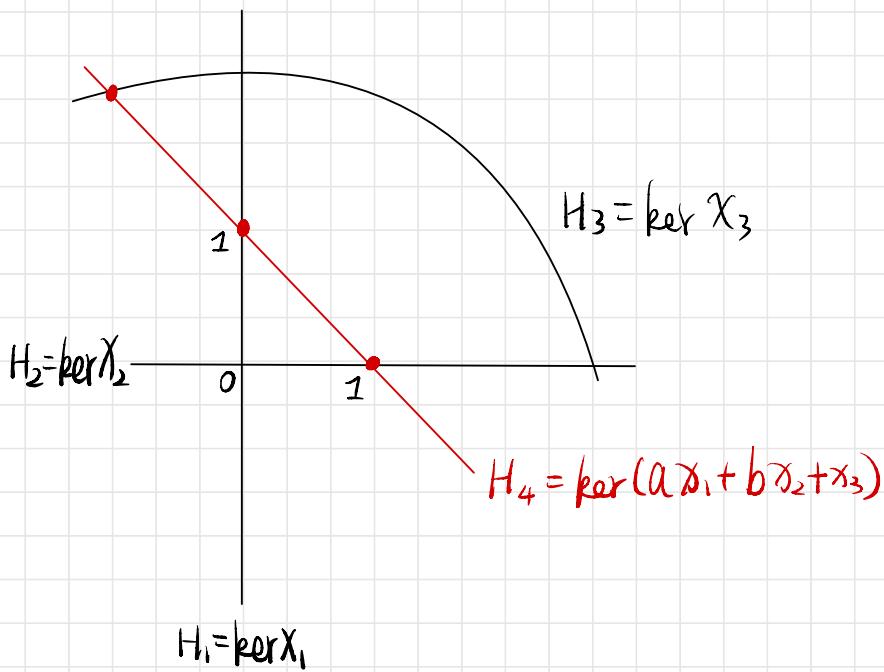
$$\mathfrak{g}_1 = \theta_E$$

$$\mathfrak{g}_2 = x_1 (ax_1 + bx_2 + x_3) \partial_{x_1}$$

$$\mathfrak{g}_3 = x_2 x_3 \prod_{\substack{j=1, \\ j \neq i}}^s (t_j x_2 + x_3) (\partial_{x_2} - t_i \partial_{x_3})$$

Example :

Let $\ell=3$ and $Q(\mathbf{A}) = \mathbf{x}_1 \mathbf{x}_2 \mathbf{x}_3 (ax_1 + bx_2 + x_3)$



- $\mathbf{A}' = \mathbf{A} \setminus \{H_4\}$ is free with $\exp(\mathbf{A}') = (1, 1, 1)$
- \mathbf{A}^{H_4} is free with $\exp(\mathbf{A}^{H_4}) = \begin{cases} (1, 2) & , a, b \in \mathbb{K} \setminus 0 \\ (1, 1) & , "a=0" \text{ or } "b=0" \end{cases}$

By A-D Theorem, we may get that

\mathbf{A} is free \Leftrightarrow either " $a=0$ " or " $b=0$ "

§ Derivation Degree Sequences of Logarithmic Modules for close to Free Hyperplane Arrangements

The study of $D(\mathcal{A})$:

→ natural approach:

- the min. generators.
- the degrees of the gens
- the min. free resolution
- :
- free arrangements
- non-free = much unexplored

Problem:

If \mathcal{A} is free, then its deletion $\mathcal{A}' = \mathcal{A} \setminus H_1$
is far from free, or not so far?

• SPOG arrangements

Def [Abe, 2021]

We say that \mathcal{A} is strictly plus-one generated (SPOG) with exponents $\text{POexp}(\mathcal{A}) = (d_1, \dots, d_e)$ and level d if $D(\mathcal{A})$ has a min. free resolution:

$$0 \rightarrow S[-d-1] \xrightarrow{(\alpha, f_1, \dots, f_e)} S[-d] \oplus \left(\bigoplus_{i=1}^e S[-d_i] \right) \longrightarrow D(\mathcal{A}) \rightarrow 0$$

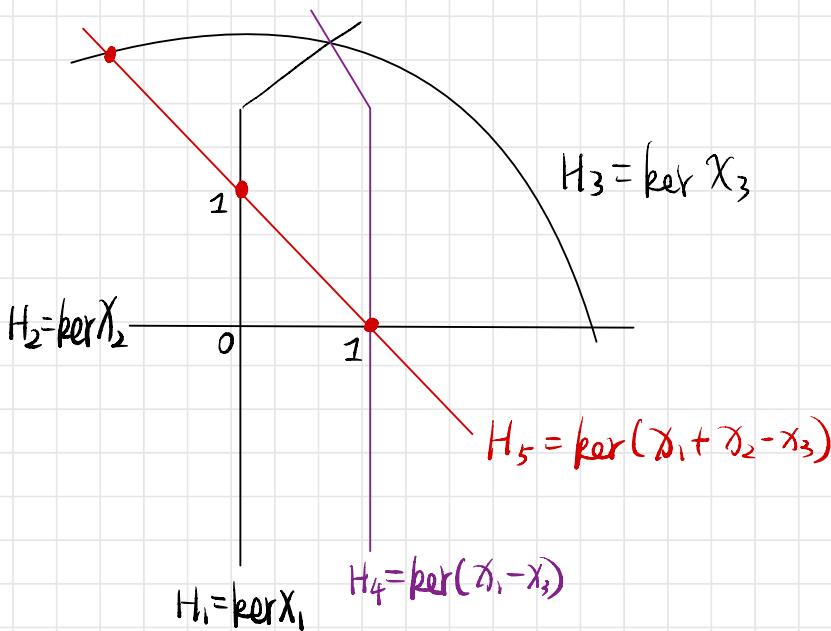
where $f_i \in S$, $\alpha \in S_1 \setminus 0$.

Theorem [A-, 2021]

Let \mathcal{A} be free with $\text{exp}(\mathcal{A}) = (d_1, \dots, d_e)$ and $H \in \mathcal{A}$. Then $\mathcal{A}' := \mathcal{A} \setminus \{H\}$ is free, or SPOG with $\text{POexp}(\mathcal{A}') = (d_1, \dots, d_e)$ and level $d = |\mathcal{A}'| - |\mathcal{A}^H|$

Example:

Let $V = \mathbb{R}^3$ and $Q(\alpha) = x_1 x_2 x_3 (x_1 - x_3)(x_1 + x_2 - x_3)$



- α is free with $\exp(\alpha) = (1, 2, 2)$
- $\alpha' = \alpha \setminus \{H_4\}$ is not free.

By the above theorem, we know that.

α' is spon with $P0\exp(\alpha') = (1, 2, 2)$ and level $|\alpha'| - |\alpha^{H_4}| = 4 - 2 = 2$.

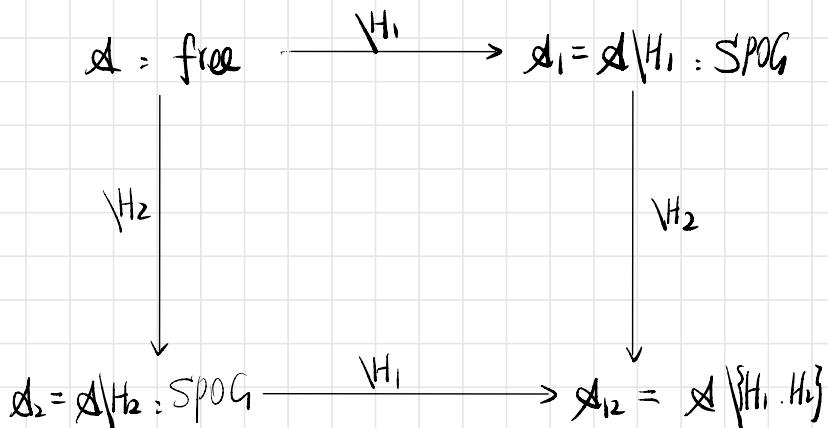
Problem :

If \mathcal{A} is free, then $\mathcal{A} \setminus \{H_1, H_2\}$ is far from free, or not so far?

Def: Assume that $\{\theta_1, \dots, \theta_p\}$ is a min. set of gen. for $D(\mathcal{A})$ with $\deg \theta_i = d_i$.

The derivation degree sequence of $D(\mathcal{A})$ defined as $DS(\mathcal{A}) = (d_1, \dots, d_p)$.
Moreover, $|DS(\mathcal{A})| = p$.

- Free minus two arrangement.
- \mathcal{A} = free with $exp(\mathcal{A}) = (d_1, d_2, \dots, d_e) \in$
- $\mathcal{A}_i := \mathcal{A} \setminus \{H_i\}$, $\mathcal{A}_{i,j} := \mathcal{A} \setminus \{H_i, H_j\}$, $H_i, H_j \in \mathcal{A}$
- \mathcal{A}_1 and \mathcal{A}_2 are SPOG with level c_1 and c_2



Theorem [C-]: Let \mathcal{A}_i be SPOG with $\text{POexp}(\mathcal{A}_i) = (d_1, d_2, \dots, d_e)$ and level C_i , where $i=1, 2$. Then $\text{DS}(\mathcal{A}_{1,2})$ is as one of following:

(I) $|\mathcal{A}_{H_1 \cap H_2}| = 2$.

In this case, $\text{DS}(\mathcal{A}_{1,2}) = (d_1, d_2, \dots, d_e, C_1, C_2)$

(II) $|\mathcal{A}_{H_1 \cap H_2}| > 2$.

(IIa) $\text{DS}(\mathcal{A}_{1,2}) = (d_1, \dots, d_e, C_1 - 1)$

(IIb) $\text{DS}(\mathcal{A}_{1,2}) = (d_1, \dots, d_e, C_1 - 1, C_2)$

(IIc) $\text{DS}(\mathcal{A}_{1,2}) \supsetneq (d_1, \dots, d_e, C_1, C_2)$

* If the relation in (IIc) is a proper subset, this shows that $|\text{DS}(\mathcal{A}_{1,2})| > e + 2$.

$H_i \in \mathcal{A}$ is in the same order as its in $Q(\mathcal{A})$.

• Example. Let $V = \mathbb{K}^4$ and

$$Q(\mathcal{A}) = x_1 x_2 x_3 x_4 (x_1 - x_2) (x_1 - x_3) (x_1 - x_4) (x_2 - x_3) \\ (x_2 - x_4) (x_3 - x_4) (x_2 - x_3 + x_4) (x_1 - x_2 + x_3 - x_4)$$

\mathcal{A} is free with $\exp(\mathcal{A}) = (1, 3, 4, 4)$.

(II) $\mathcal{A}_1, \mathcal{A}_2$: SPOG with level 4.

$\mathcal{A}_8, \mathcal{A}_{10}$: SPOG with level 5

$$(IIa) DS(\mathcal{A}_{1,2}) = (1, 3, 3, 4, 4)$$

$$(IIb) DS(\mathcal{A}_{1,8}) = (1, 3, 4, 4, 4, 5)$$

$$(IIc) DS(\mathcal{A}_{2,10}) = (1, 3, 4, 4, 4, 4)$$

• Example. Let $V = \mathbb{K}^4$ and

$$Q(\mathcal{A}) = x_1 (x_1 - x_2) (x_1 + x_2) (x_1 - x_3) (x_1 + x_3) (x_1 - x_4) \\ (x_1 + x_4) (x_2 - x_3) (x_2 + x_4) (x_3 - x_4)$$

\mathcal{A} is free with $\exp(\mathcal{A}) = (1, 3, 3, 3)$ By computer,

we have $DS(\mathcal{A}_{1,10}) = (1, 3, 3, 3, 3, 3, 3, 3) \Rightarrow |DS(\mathcal{A}_{1,10})| = 7 > 4 + 2$

* This is a counter-example to Orlik's Conjecture.
[Di Pasquale, 2023]

- Example. Let $V = \mathbb{K}^4$ and.

$$Q(\Delta) = x_1 x_2 x_3 x_4 (x_1 - x_2)(x_1 - x_3)(x_2 - x_3)(x_3 - x_4)$$

$$(x_2 - x_3 + x_4)(x_1 - x_2 + x_3 - x_4)$$

Δ is free with $\exp(\Delta) = (1, 3, 3, 3)$

$\Delta_1, \Delta_3 = \text{SPOG}$ with level 3.

By computer, we have $DS(\Delta_{1,3}) = (1, 3, 3, 3, 3, 3)$

*This is a counter-example to Orlik's Conjecture.

[Nakashima - Tsujie, 2023]

However, by computer, the first **counter example** to the Orlik's conjecture by Edelman and Reiner has at most **1+2 min. generators**.

Therefore, perhaps we cannot assert that there is a certain connection between Orlik's conjecture and the cardinality of minimal set of homogeneous generators.

3-dimensional case.

Theorem [C-]:

Let $\ell=3$ and \mathcal{A}_i be SPOG with
 $\text{POexp}(\mathcal{A}_i) = (d_1, d_2, d_3)_\leq$ and level C_i , where $i=1, 2$.

Then $\text{DS}(\mathcal{A}_{1,2})$ is as one of following:

$$(I) |\mathcal{A}_{H_1 \cap H_2}| = 2 \Leftrightarrow \text{DS}(\mathcal{A}_{1,2}) = (d_1, d_2, d_3, C_1, C_2)$$

$$(II) |\mathcal{A}_{H_1 \cap H_2}| > 2. \text{ Assume } C_1 \leq C_2.$$

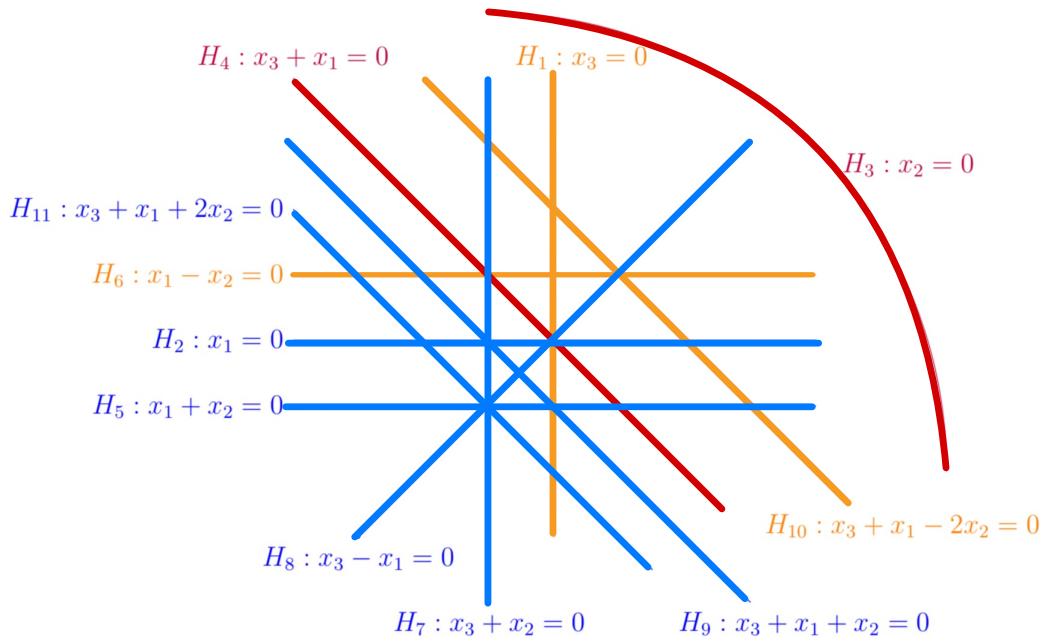
$$(IIa) C_1 = C_2 \Leftrightarrow \text{DS}(\mathcal{A}_{1,2}) = (d_1, d_2, d_3, C_1 - 1)$$

$$\text{Moreover, } C_1 = C_2 = d_3$$

$\Leftrightarrow \mathcal{A}_{1,2}$ is SPOG with level d_3

$$(IIb) C_2 > C_1 \Leftrightarrow \text{DS}(\mathcal{A}_{1,2}) = (d_1, d_2, d_3, C_1, C_2 - 1)$$

Example :

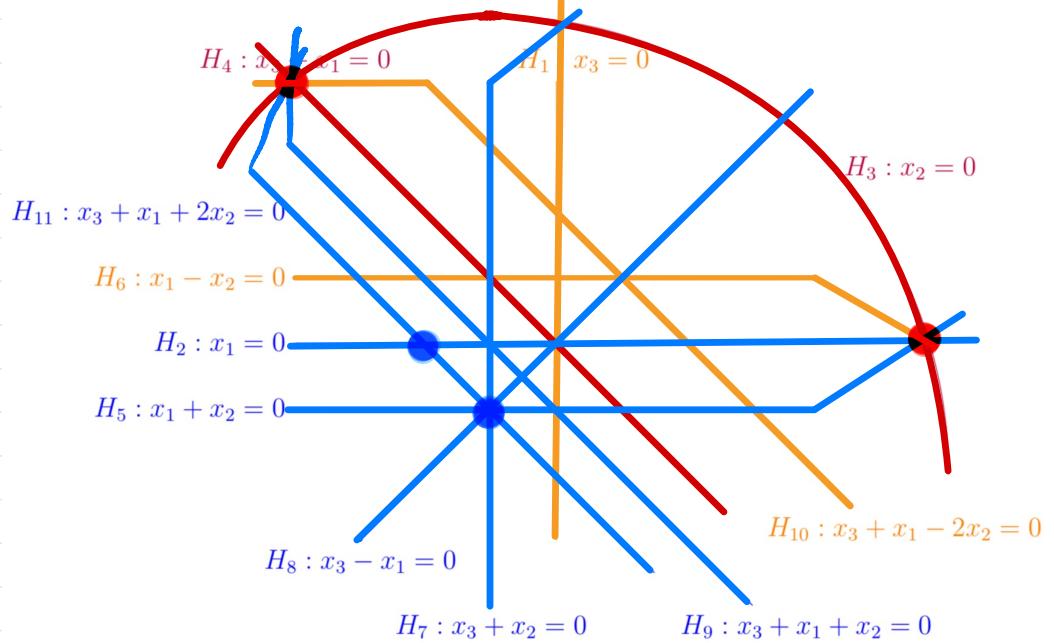


This is a free arrangement with $\exp(\Delta) = (1, 5, 5)$

$\Delta_i := \Delta \setminus H_i$: SPOG with level 5

$\Delta_j := \Delta \setminus H_j$: SPOG with level 6

$\Delta_k := \Delta \setminus H_k$: free with $(1, 4, 5)$.



By the above theorem, we may get that.

$$(1). DS(\Delta_{2,11}) = (1, 5, 5, 5, 5)$$

$$(2). DS(\Delta_{3,4}) = (1, 5, 5, 5),$$

but $\Delta_{3,4}$ is not SPOG.

(3). $\Delta_{7,8}$ is SPOG with level 5

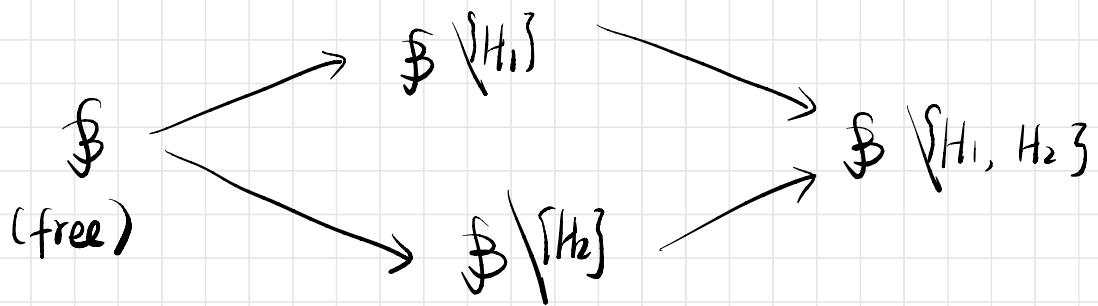
$$\text{and } DS(\Delta_{7,8}) = (1, 4, 5, 5)$$

$$(4) DS(\Delta_{2,3}) = (1, 5, 5, 5, 5)$$

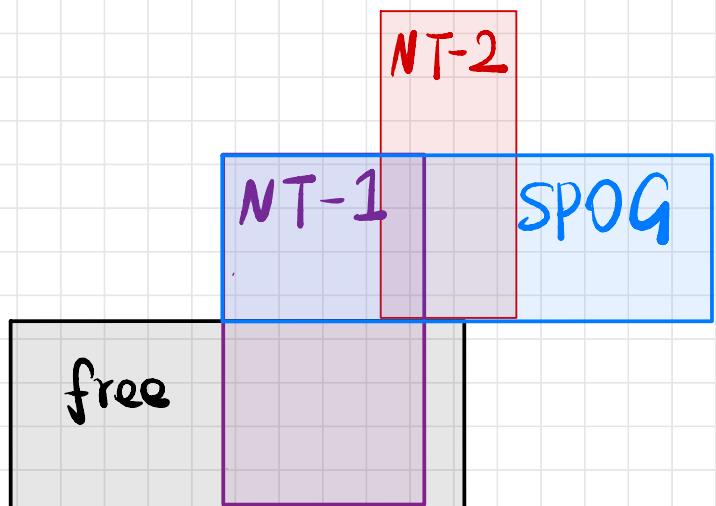
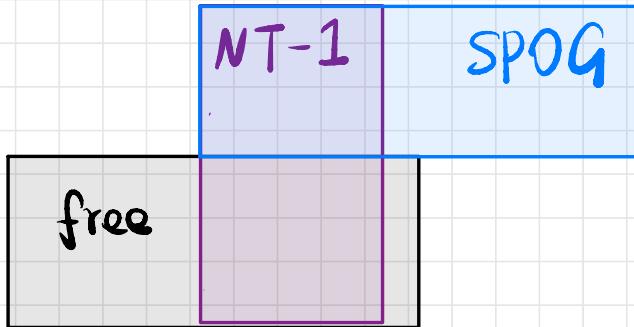
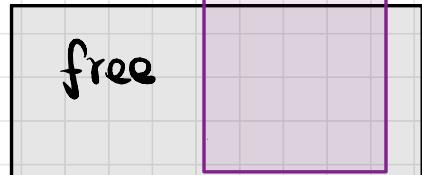
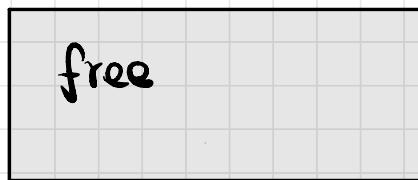
{ Conclusion & Conjecture.

Def 1 = We say that \mathcal{A} is next to free minus one ($NT-1$) if \exists free arrangement \mathbb{B} and $H \in \mathbb{B}$ s.t. $\mathcal{A} = \mathbb{B} \setminus H$

- Consider $\mathbb{B} \setminus \{H_1, H_2\}$



Def 2 = We say that \mathcal{A} is Strictly next to free minus two ($NT-2$) if \exists free arrangement \mathbb{B} and $H_1, H_2 \in \mathbb{B}$ s.t. $\mathbb{B} \setminus H_i$ is not free and $\mathcal{A} = \mathbb{B} \setminus \{H_1, H_2\}$, where $i=1, 2$



claim:

$$\bullet \{NT-2\} \cap \{\text{free}\} = \emptyset$$

$$\bullet \{NT-2\} \cap \{NT-1\} \subseteq \text{SPOG}$$

Example:

$$\text{Let } Q(\mathcal{A}) = x_1 x_2 x_3 (x_1 - x_2) (x_1 + x_2) (x_1 + 2x_2) \\ (2x_1 + x_2) (3x_1 + x_2) (x_2 + x_3) (3x_1 + x_2 + x_3)$$

\mathcal{A} is free with $\exp(\mathcal{A}) = (1, 3, 6)$

$$H_1 = \ker x_1, \quad H_2 = \ker x_2$$

$\Rightarrow \mathcal{A}_{1,2}$ is SPOG with

$$\text{POexp}(\mathcal{A}_{1,2}) = (1, 3, 5) \text{ and level 6.}$$

but $\mathcal{A}_{1,2}$ is not a NT-1. arrangement.

i.e. $\forall H \subset V, \quad \mathcal{A}_{1,2} \cup H$ is not free.

Conjecture:

1. $\{NF-1\} \cap \{NF-2\} = \emptyset$.

2. If \mathbb{A} is NF-2, then $\text{pd}_S D(\mathbb{A}) \leq 2$

Moreover, if # min. generators for $D(\mathbb{A})$ is at most $l+2$, then $\text{pd}_S D(\mathbb{A}) \leq 1$.

3. Free resolution of NF-2.

4. $L(\mathbb{A}) \cong L(\mathbb{B}) \Rightarrow \text{pd}_S D(\mathbb{A}) = \text{pd}_S D(\mathbb{B})$

Thank you

for

your attention

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Paul Mücksch.

Akiko Yazawa