

Advanced Persistent Jest

*... I wield the power of business
excellence, ten thousand senior project
managers could not manage me.*

Writeup – smartfridge1 and smartfridge2 (33C3 Part 2 of 2)

Posted on December 30, 2016 by Norman

During the last 2 days, I participated in the 33c3 CTF. This post contains the second part of my writeups for this CTF, covering the smartfridge1 and smartfridge2 challenges.

Note that the content provided for smartfridge2 (a crypto challenge!) speeds up solving smartfridge1 considerably, so I will provide the original files for both challenges together.

smartfridge1

These two challenges were presented as a binary (smartfridge1), and a pcap of someone communicating with said binary (smartfridge2), which you can download [here](#). The challenge text indicates that this is something to do with the Bluetooth LE 4.0 protocol:

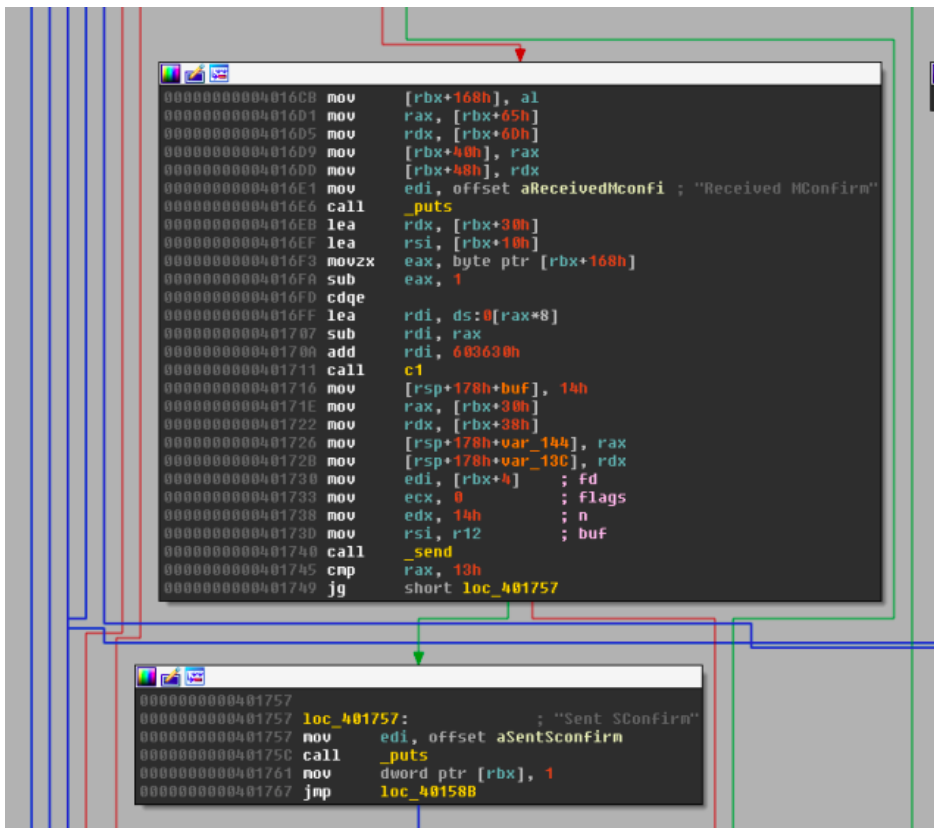


Our task appears to be to reverse engineer the provided binary, build something which can interface with it (and thus, the remote server), and send a series of commands to “retrieve” the flag.

On first inspection, we notice that the binary makes heavy use of OpenSSL AES functions, and is littered with printf debugging. As a starting point, I decided to use the pcap provided for smartfridge2, and simply send the first packet in the pcap to the binary. Instant success:

```
savra@svogthos:/ctf/33c3/smartfridge$ ./smartfridge-fdc1e0eaa6617b7bcdeea3ec9f142f3b27bd4a4
Data received on 1
Valid packet of size 21 received
Received MConfirm
Sent SConfirm
Data received on 1
Disconnecting client 1
```

We can look for these printf statements in IDA (the second SEND below is an SCONFIRM packet):



Our next step is to carefully gdb our way through the “c1” function, which appears to set an AES encryption key and encrypt *something*, to determine what’s being encrypted:

```

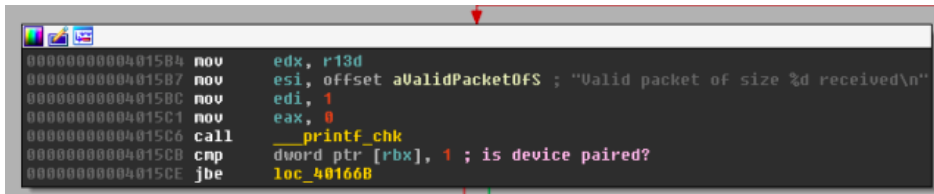
=> 0x401481 <e+56>:      call    0x400cb0 <AES_encrypt@plt>
0x401486 <e+61>:      mov     rax,QWORD PTR [rsp+0xf8]
0x40148e <e+69>:      xor     rax,QWORD PTR fs:0x28
0x401497 <e+78>:      je      0x40149e <e+85>
0x401499 <e+80>:      call    0x400d40 <__stack_chk_fail@plt>
Guessed arguments:
arg[0]: 0xe608020 --> 0xe60b7cb4bf1804a
arg[1]: 0x608040 --> 0x0
arg[2]: 0x7fffffffdec0 --> 0x1e240
[-----stack-----]
0000| 0x7fffffffdec0 --> 0x1e240
0008| 0x7fffffffdec8 --> 0x0
0016| 0x7fffffffded0 --> 0x6362812263628122
0024| 0x7fffffffded8 --> 0x6362812263628122
0032| 0x7fffffffdee0 --> 0x93fbba0ef0992b2c
0040| 0x7fffffffdee8 --> 0x93fbba0ef0992b2c
0048| 0x7fffffffdef0 --> 0xc8be8e8a5b452484
0056| 0x7fffffffdef8 --> 0xabdc0fa83827a5a6
[-----]
Legend: code, data, rodata, value

Breakpoint 1, 0x00000000401481 in e ()
qdb$

```

In this instance, the key is somewhat recognizable (the hexadecimal representation of atoi(“123456” – which appears to be a test PIN)), but the plaintext (0xe6...) is not. It doesn’t help that despite sending a static message to the server, this plaintext seems to change every time. Let’s set this aside for now.

Going back through the disassembly, we can shed some light on the structure of the packet: breakpoints at 0x40158B and 0x40159A reveal that the first DWORD is a packet length, which must be above 3 and below 84. Also, we can see a “state variable” being used to control program flow – packets must be sent in order:



```

0000000000401584  nov     edx, r13d
0000000000401587  nov     esi, offset aValidPacket0FS ; "Valid packet of size %d received\n"
000000000040158C  nov     edi, 1
00000000004015C1  nov     eax, 0
00000000004015C6  call    __printf_chk
00000000004015C8  cmp     dword ptr [rbx], 1 ; is device paired?
00000000004015CE  jbe     loc_401668

```

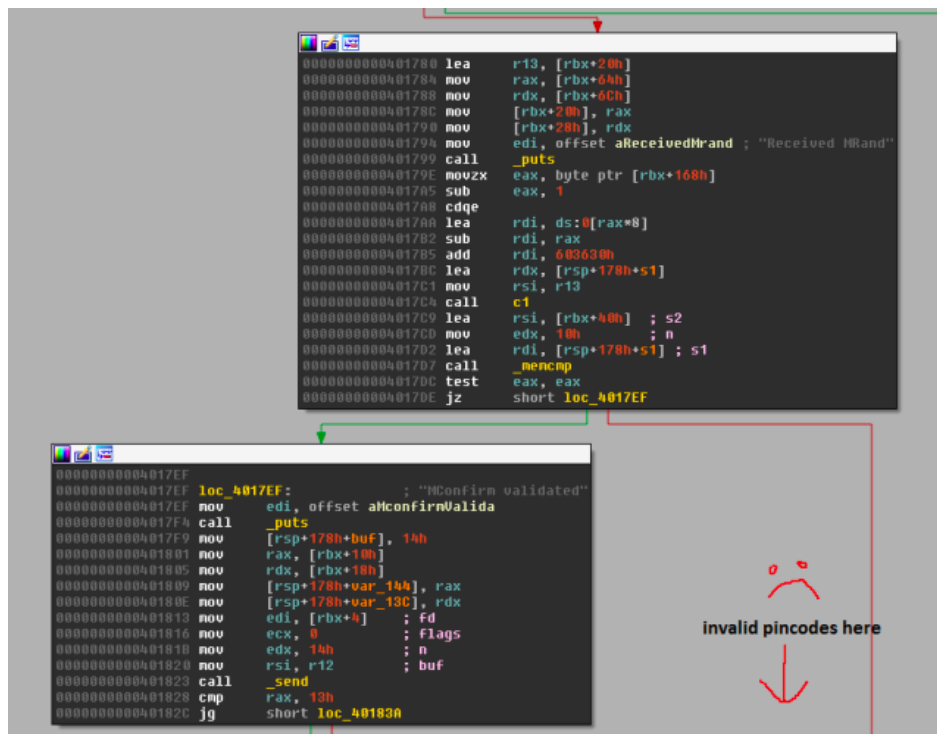
Continuing with this avenue of investigation, we add in the second packet from the pcap (the one starting with “\x14\x00\x00\x00”) to our client:

```

savra@svogthos:ctf/33c3/smartfridge$ ./smartfridge-fdcde0eaa6617b7bcdeea3ec9f142f3b27bd4a4
Data received on 1
Valid packet of size 21 received
Received MConfirm
Sent SConfirm
Data received on 1
Valid packet of size 20 received
Received MRand
Invalid pincode
Disconnecting client 1

```

This time, something more interesting comes up: once more, we return to the disassembly to find out what’s causing our “invalid pincode” message:



```

0000000000401780  lea     r13, [rbx+20h]
0000000000401784  mov     rax, [rbx+60h]
0000000000401788  mov     rdx, [rbx+60h]
000000000040178C  mov     [rbx+20h], rax
0000000000401790  mov     [rbx+28h], rdx
0000000000401794  mov     edi, offset aReceivedHrand ; "Received HRand"
0000000000401799  call    _puts
000000000040179E  movzx   eax, byte ptr [rbx+168h]
00000000004017A5  sub     eax, 1
00000000004017A8  cdq     rax
00000000004017AA  lea     rdi, ds:0[rax*8]
00000000004017B2  sub     rdi, rax
00000000004017B5  add     rdi, 603630h
00000000004017BC  lea     rdx, [rsp+178h+51]
00000000004017C1  mov     rsi, r13
00000000004017C4  call    c1
00000000004017C9  lea     rsi, [rbx+40h] ; s2
00000000004017CD  mov     edx, 10h ; n
00000000004017D2  lea     rdi, [rsp+178h+51] ; s1
00000000004017D7  call    _memcmp
00000000004017DC  test    eax, eax
00000000004017DE  jz      short loc_4017EF

```

```

00000000004017EF  loc_4017EF: ; "MConfirm validated"
00000000004017EF  mov     edi, offset aMconfirmValida
00000000004017F4  call    _puts
00000000004017F9  mov     [rsp+178h+buf], 14h
0000000000401801  mov     rax, [rbx+10h]
0000000000401805  mov     rdx, [rbx+18h]
0000000000401809  mov     [rsp+178h+var_144], rax
000000000040180E  mov     [rsp+178h+var_13C], rdx
0000000000401813  mov     edi, [rbx+4] ; fd
0000000000401816  mov     ecx, 0 ; flags
0000000000401818  mov     edx, 14h ; n
0000000000401820  mov     rsi, r12 ; buf
0000000000401823  call    _send
0000000000401828  cmp     rax, 13h
000000000040182C  jg      short loc_40183A

```

Here, we go back to dynamic analysis to find what’s being passed to c1, and what it’s being compared against. To make life easier for ourselves, we replace the “body” of each packet with clearly recognizable strings like “\x01”*16 (MConfirm, packet 1) and “\x02”*16 (MRand, packet 2).

This time, we have some success:

```

=> 0x4017c4 <handle+624>: call 0x4014a8 <c1>
0x4017c9 <handle+629>: lea rsi,[rbx+0x40]
0x4017cd <handle+633>: mov edx,0x10
0x4017d2 <handle+638>: lea rdi,[rsp+0x20]
0x4017d7 <handle+643>: call 0x400d70 <memcmp@plt>
Guessed arguments:
arg[0]: 0x603637 --> 0x363534333231 ('123456')
arg[1]: 0x608030 --> 0x2020202020202020
arg[2]: 0x7fffffff040 --> 0x7fffffff070 --> 0x7ffff763f130 --> 0xc0022000015b5

```

Here, we can see that upon receiving the “MRand” packet, “c1” is called to encrypt “\x02\x02\x02\x02...” with “123456”, our test PIN code, the result of which is “d86docfbcb7bbo343f09fad5ca6824429”. We then proceed to

the memcmp call at 0x4017D7:

```
=> 0x4017d7 <handle+643>:      call    0x400d70 <memcmp@plt>
0x4017dc <handle+648>:      test    eax,eax
0x4017de <handle+650>:      je      0x4017ef <handle+667>
0x4017e0 <handle+652>:      mov     edi,0x402037
0x4017e5 <handle+657>:      call    0x400c20 <puts@plt>
Guessed arguments:
arg[0]: 0x7fffffff040 --> 0x4303bbc7fb0c6dd8
arg[1]: 0x608050 --> 0x101010101010101
arg[2]: 0x10
```

Hang on, we know those values! It's comparing the content of our MConfirm packet with AES(key="123456",content=MRand)! Knowing this, it is trivial to pass the MRand check, by sending the correct MConfirm packet first:

```
savra@svogthos:/ctf/33c3/smartfridge$ ./smartfridge-fdc1e9eaa6617b7bcdea3ec9f142f3b27bd4a4
Data received on 1
Valid packet of size 21 received
Received MConfirm
Sent SConfirm
Data received on 1
Valid packet of size 20 received
Received MRand
MConfirm validated
Sent SRand
Paired
Data received on 1
Disconnecting client 1
```

Success! At this point, we are sending:

- Packet 1 – MConfirm: \x15\x00\x00\x00\x00\x02 + aes(key="123456",text="\x02"*16)
- Packet 2 – MRand: \x14\x00\x00\x00\x00 + "\x02"*16

Now that we're paired, let's try sending another packet. As the process logic appears to be controlled via state variable instead of any "command specifier", let's just try sending a packet containing "\x03"*16 and let's see what happens:

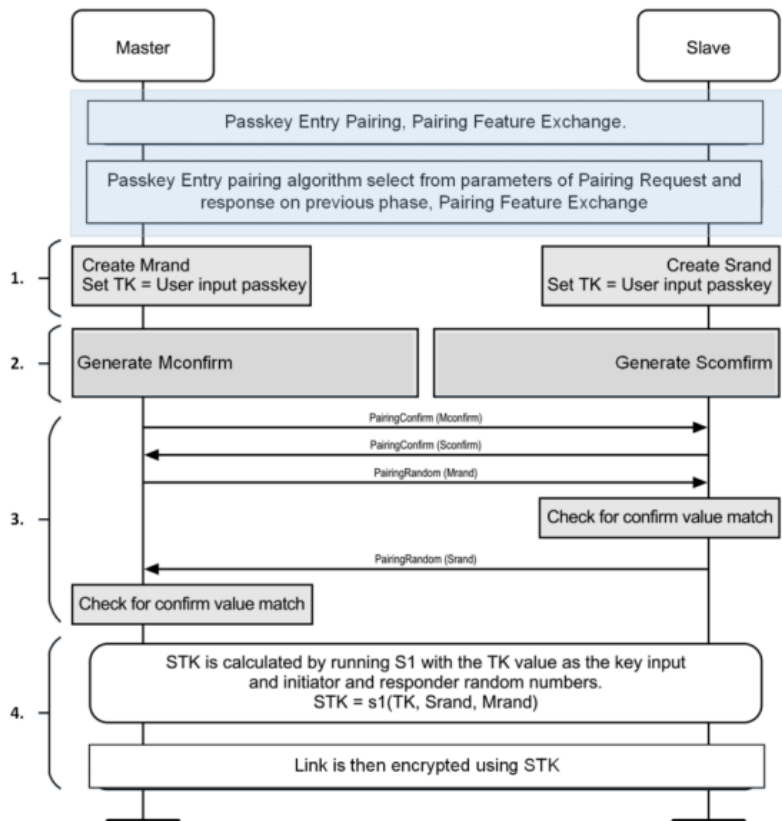
```
Paired
Data received on 1
Valid packet of size 20 received
Invalid padding
Disconnecting client 1
```

Going back through the disassembly, we can see that this comes from a decryption function at 0x4015F2:

```
00000000004015F2
00000000004015F2 loc_4015F2:
00000000004015F2 lea     r15, [rbx+0E8h]
00000000004015F9 lea     rdi, [rbx+50h]
00000000004015FD mov     rcx, r15
0000000000401600 mov     edx, ebp
0000000000401602 mov     rsi, r14
0000000000401605 call    decrypt
000000000040160A lea     eax, [r13-5]
000000000040160E movzx   eax, byte ptr [rbx+rax+0E8h]
0000000000401616 cmp     al, 10h
0000000000401618 jbe     short loc_40162E
```

At a glance, this function appears to attempt an AES decryption, and if the decryption is successful, the decrypted data gets passed to strtok (0x40162E), and from there, to command handlers which try to do string comparisons against known commands (e.g. "OPEN" at 0x4018B3), giving us some clue as to what the plaintext needs to be.

Stepping through this decryption function, I quickly noticed that the key was not the "PIN code" key. Perplexed, I turned to the Bluetooth Low Energy documentation:



from <https://blog.bluetooth.com/bluetooth-pairing-part-3-low-energy-legacy-pairing-passkey-entry>

From here, we can make a reasonable guess that the `decrypt()` function is using the STK instead of the TK / pincode. After reading through the documentation, we can determine that the STK is calculated thusly:

```
STK = AES(key=pin,text=second_half_of_srand + first_half_of_mrand)
```

Going through the disassembly, we can see that the “s1” function is present in the binary, at 0x40188F – we’ll need to emulate this function in our client, so we can use the same STK session key to encrypt command traffic to the server. In gdb:

```

=> 0x40188f <handle+827>:      call    0x40150a <s1>
0x401894 <handle+832>:      mov     DWORD PTR [rbx],0x2
0x40189a <handle+838>:      mov     edi,0x402065
0x40189f <handle+843>:      call    0x400c20 <puts@plt>
0x4018a4 <handle+848>:      jmp     0x40158b <handle+55>
Guessed arguments:
arg[0]: 0x7fffffff030 --> 0x1e240
arg[1]: 0x608020 --> 0xd1305cafc8590ca9
arg[2]: 0x608030 --> 0x2020202020202020
arg[3]: 0x608060 --> 0x0
[-----stack-----]
0000| 0x7fffffff020 --> 0x7fffffff050 --> 0xc8590ca900000014
0008| 0x7fffffff028 --> 0x7fffffff040 --> 0x4303bbc7fb0c6dd8
0016| 0x7fffffff030 --> 0x1e240
0024| 0x7fffffff038 --> 0x0
0032| 0x7fffffff040 --> 0x4303bbc7fb0c6dd8
0040| 0x7fffffff048 --> 0x294482a65cad9ff0
0048| 0x7fffffff050 --> 0xc8590ca900000014
0056| 0x7fffffff058 --> 0xf6d83fbfd1305caf
[-----]
Legend: code, data, rodata, value

Breakpoint 1, 0x000000000040188f in handle ()
gdb> x/32bx 0x608020
0x608020: 0xa9 0x0c 0x59 0xc8 0xaf 0x5c 0x30 0xd1
0x608028: 0xbf 0x3f 0xd8 0xf6 0x62 0x67 0xea 0x7f
0x608030: 0x02 0x02 0x02 0x02 0x02 0x02 0x02 0x02
0x608038: 0x02 0x02 0x02 0x02 0x02 0x02 0x02 0x02
gdb>
  
```

With the STK, we can now send correctly encrypted commands to the server. The commands “OPEN 2”, “LIST” and “TAKE 0” can be used to retrieve a flag – in our test binary, this is “33C3_NOT_FLAG_2” – the only remaining change is to modify the PIN code, and send traffic to the actual server instead:

A little trial and error later, and I realized that the failed pairing was due of the initial MConfirm message. When we prefixed the packet with “\x15\x00\x00\x00\x02”, we were selecting user 2, who seemed to have a different PIN code. Moments later (and after changing “OPEN 2” to “OPEN 1”, the flag:

You can find the complete solution script [here](#) – it’s a functional implementation of Bluetooth LE pairing over TCP!

smartfridge2

With the knowledge of the Bluetooth LE pairing protocol gained during the last challenge, this challenge was trivial. We had a PCAP of a Bluetooth pairing, for “user 2” (as indicated by “\x15\x00\x00\x00\x02” in the MConfirm packet).

With the MConfirm and MRand packets, and knowing that:


```
MConfirm = AES(key=pin,cleartext=MRand)
```

it is trivial to brute force the key. A few modifications to the client script later, and we have the second flag:

```
[!] received: 0cbc663729f9fe6991f1400b07014ae8
[!] decrypted: 5ac34b0f6ebbe57c02d5a43d5aeef0e0
[>] sending MRAND, which is 0x02 * 16
[<] recving SRAND, decrypted key in srand_decrypted...
[!] SRAND received: 5ac34b0f6ebbe57c02d5a43d5aeef0e0
[!] our unencrypted session key is [CORRECT] : 02d5a43d5aeef0e00202020202020202
[!] our session key is s1(tk,srand,rand) : d6c718ed4931b9524c7521e284e4b834
[CBK_KEY/session key IS CORRECT:BREAK AT 0x401894 and VERIFY 0x608060]
[>] trying to send OPEN 2 1, which is 5b22da7147d08d7a57449a4574c63c66
[>] trying to send LIST, which is 86e2456f53d674c9b499c82270082545
0. yoghurt
1.
2.
3.
4.
\x05\x05\x05\x05\x05
[>] trying to send TAKE 0, which is 7a64ba5f9473c1a9ecc02b08b8d2cd8f
yoghurt: 33C3_sh0rt_p4ssc0de_1s_sh0rt
```

You can find the modified solution script [here](#).

As always, I'd like to thank the organisers of 33c3's CTF, Eat Sleep Pwn Repeat, for putting together a fantastic event. While I was not able to solve as many challenges as I had hoped, I had a lot of fun and learned alot solving the ones I was able to complete, and look forward to doing the others in my own time.

If I might offer some feedback, the point values seem imbalanced, and this was reflected by other participants in this event. For example, it is clear that smartfridge2 required a miniscule fraction of the effort required by smartfridge1, yet was worth 50% of the points. Still, this by no means detracted from the excellence of the event, as it didn't make the challenges less fun.

A safe and happy New Year's to you all, and see you all in a few weeks for the Insomni'hack Teaser CTF.

About Norman

Sometimes, I write code. Occasionally, it even works.

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