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 ValarDragon Fix bolding, and add My USB writeup

6205515 8 days ago

1 contributor

68 lines (52 sloc) 4.69 KB

Raw Blame History

- 400 points
- Category: Cryptography
- Problem statement: *My friend made some sort of encryption service using the Paillier cryptosystem. Can you get him to encrypt the string `easyctf{3ncrypt_m3!}` for me? Your flag will be a base 10 integer. Access his encryption service at `paillier.tcp.easyctf.com 8570`.*
- Problem hint: [https://en.wikipedia.org/wiki/Paillier_cryptosystem]

We are not actually 'cracking' the Paillier Cryptosystem for this challenge, we are gaining the ability to encrypt our own messages already given the ability to choose the message and random number. (n,g) is the normal public key for this cryptosystem

Just trying that in the service gives us:

```
valar@valardev-Vostro-3460-mint ~/Code/CTF/ezCTF $ nc paillier.tcp.easyc.tf.com 8570
Enter a message to encrypt (int): 578781299356711768839252397261103878073419506045
Enter r (int): 1
Bad m. We only take m from 0 to 1000000000000000000.valar@valardev-Vostro-3460-mint
```

A bit of info about the Paillier Cryptosystem: Its a Homomorphic probabilistic public key cryptography system. What that means:

- Homomorphism: $\text{decryption}(\text{encryptedMessage1} * \text{encryptedMessage2}) = \text{plaintext1} + \text{plaintext2}$ (recall that plaintexts are just numbers, that are decoded in different ways)
- Probabilistic: A message can be encrypted in many different ways using the same key, so a lookup table is impossible, but all of these decrypt to the same message.

We only care about encryption here. There are 2 parameters, g and n that are fixed. g is some number less than n^2 . To encrypt a message m , you do

$$c = (q^m)(r^n) \bmod n^2$$

where r is a random number less than n

Getting g is easy! Make $m = 1$, $r = 1$, and then c is $g!$ Now that we have g , increase m by 1 and put that into the service. Lets call that c , g^2 .

$$g^2 = (g^2)(r^1) \bmod n^2$$

From the definition of modulo:

$$g^2 - g^2 = k \cdot n^2$$

If we do that for a couple of message values, $g^3, g^4, g^5 \dots$, we should get enough values of things that n factors into to find n . We can also test any n that we have by just raising r by 1, and seeing if we get the expected results. Additionally we know that n must be a perfect square.

I was about to begin coding something to check and do all that, but I decided to put $g^2 - g^2$ into factordb, to see if it was something easy.

It turns out it couldn't have been easier! n^2 is the square of a [prime](#)!

So now we have everything we need to encrypt! Set r to 1 for convenience, and then its just a single modpow. I just used the Homomorphic property in my code, since I think its cool, but you can compute it directly.

```
m5r1 = pow(g,5,n2)
goalDiv5 = goal // 5
# Now use the Homomorphic property :)
flagInt = pow(m5r1,goalDiv5,n2)
print(flagInt)
```

The flag's an integer not some string that you can hex decode, since theres of tons equally valid ways to encode this message. (Imagine all the different factorizations of the goal number you can use in the modpows, or multiply different encrypted messages with different r values)

And thats 400 more points for us! Code used to solve the problem is in `Paillier.py`

I submitted

```
44073117240618665780675193850837939995438219250244678211539041436428154743261238082817577099306521708734123
381615432054274681465095612422847370622010652215512660940106734460138798004151939831278940754163448609294265
458598883535128433424615303280599380544523443593952238464672302887846705279608801286723167548136016323776193
330983364067235836166569465230366 as our flag!
```

