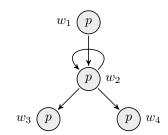
## PHIL 143 Hw4

## Jad Damaj

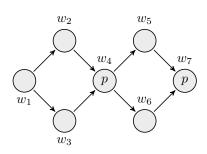
## February 14, 2024

 ${\bf Exercise}$  1. Draw the filtration of each model through the given formula.

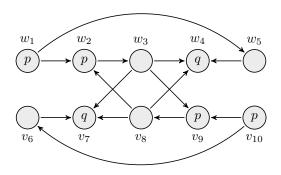
(a)  $p \to \Diamond p$ 



(b)  $\Box p \to \Box \Box p$ 

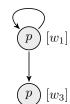


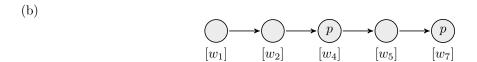
(c)  $p \to \Diamond p$ 



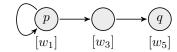
 ${\it Proof.}$ 

(a)





(c)



## Exercise 2.

- (a) Prove that if a formula  $\varphi$  contains n connectives, then  $|\operatorname{sub}(\varphi)| \leq 2n+1$ .
- (b) Show that the bound is sharp. That is, for any n, give an example of formula with n connectives and exactly 2n + 1 subformulas.

Proof.

(a) We show this by induction on the complexity of the formula.  $\varphi = p$ : If  $\varphi$  is atomic then it contains 0 connectives and  $|\mathrm{sub}(p)| = 1 = 2 \cdot 0 + 1$ .  $\varphi = \neg \psi$ : If the claim holds for  $\psi$  and  $\psi$  has n connectives then  $\varphi$  has n+1 connectives and

$$|\operatorname{sub}(\varphi)| \le |\operatorname{sub}(\psi)| + 1 \le (2n+1) + 1 \le 2(n+1) + 1$$

 $\varphi = \Box \psi$ : If the claim holds for  $\psi$  and  $\psi$  has n connectives then  $\varphi$  has n+1 connectives and

$$|\operatorname{sub}(\varphi)| \le |\operatorname{sub}(\psi)| + 1 \le (2n+1) + 1 \le 2(n+1) + 1$$

 $\varphi = \psi_1 \wedge \psi_2$ : If the claim holds for  $\psi_1$  and  $\psi_1$  and they have  $n_1$  and  $n_2$  connectives respectively then  $\varphi$  has  $n_1 + n_2 + 1$  connectives and

$$|\operatorname{sub}(\varphi)| \le |\operatorname{sub}(\psi_1)| + |\operatorname{sub}(\psi_2)| + 1 \le (2n_1 + 1) + (2n_2 + 1) + 1 = 2(n_1 + n_2 + 1) + 1$$

(b) To see that this bound is sharp inductively define a formulas by  $\varphi_0 = p_0$  and  $\varphi_{n+1} = (\varphi_n \wedge p_{n+1})$ . Then  $|\operatorname{sub}(\varphi_0)| = 1$  and  $|\operatorname{sub}(\varphi_{n+1})| = 2 + |\operatorname{sub}(\varphi_n)|$  so these witness the upper bound.

**Exercise 3.** Recursively define formulas  $\varphi_n$  as follows:

- $\bullet \ \varphi_0 = p_0$
- $\varphi_{n=1} = (\Diamond \varphi_n \wedge \Box p_{n+1})$
- (a) Find a formula for  $|\operatorname{sub}(\varphi_n)|$  in terms of n.
- (b) Find a formula for  $d(\varphi_n)$  in terms of n.
- (c) Find a formula for  $|\varphi_n|$  in terms of n.
- (d) In the limit, which method gives us a better bound on the size of the models satisfying  $\varphi_n$ ?

Proof.

- (a)  $|\operatorname{sub}(\varphi_0)| = 1$  and  $|\operatorname{sub}(\varphi_{n+1})| = |\operatorname{sub}(\varphi_n)| + 4$  and  $|\operatorname{sub}(\varphi_n)| = 4n + 1$ .
- (b)  $d(\varphi_0) = 0$  and  $d(\varphi_{n+1}) = \max\{d(\varphi_n) + 1, 1\}$  so  $d(\varphi_n) = n$ .
- (c)  $|\varphi_0| = 1$  and  $|\varphi_{n+1}| = |\varphi_n| + 6$  so  $|\varphi_n| = 6n + 1$ .

(d) Filtration gives a bound of  $2^{4n+1}$  and selection gives a bound of  $(6n+1)^n$  so filtration gives a better bound in the limit.

Exercise 4. Show that the following are theorems of K.

- 1.  $\Box(p \to \Diamond p) \to (\Diamond \Diamond p \lor \neg \Diamond p)$
- **2.**  $(\lozenge \Box (p \to q) \land \Box \lozenge p) \to \lozenge \lozenge q$

*Proof.* To show these are theorems of K we that they are valid from which their theoremhood follows by completeness.

- (a) Suppose  $\mathcal{M}, w$  is a pointed model such that  $\mathcal{M}, w \models \Box(p \to \Diamond p)$ . If there is some v such that wRv and  $\mathcal{M}, v \models p$  we have that  $\mathcal{M}, v \models p \to \Diamond p$  and so  $\mathcal{M}, v \models \Diamond p$ . Hence, it follows that  $\mathcal{M}, w \models \Diamond \Diamond p$ . Otherwise, no such v exists and we have that  $\mathcal{M}, w \models \neg \Diamond p$ . Thus,  $\mathcal{M}, w \models (\Diamond \Diamond p \lor \neg \Diamond p)$ , as desired.
- (b) Given some  $\mathcal{M}, w$  such that  $\mathcal{M}, w \models \Diamond \Box (p \to q) \land \Box \Diamond p$  we show that  $\mathcal{M}, w \models \Diamond \Diamond p$ . By assumption, there is some v such that wRv and  $\mathcal{M}, v \models \Box (p \to q)$ . Further, we must also have  $\mathcal{M}, v \models \Diamond p$  and so there is some v' such that vRv' and  $\mathcal{M}, v' \models p$ . Then, since  $\mathcal{M}, v \models \Box (p \to q), \mathcal{M}, v' \models p \to q$  and so  $\mathcal{M}, v' \models q$ . Thus, we have that  $\mathcal{M}, w \models \Diamond \Diamond q$ , as desired.