

Golf-Sport: Managing Operations A Case Study in Optimization

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1 Introduction

This problem examines a manufacturing optimization case.....

2 Background and Prior Work

Here is a good spot to insertour necessary references. We can evaluate solutions of similar problems

3 Description of Problem

Golf-Sport is a golfclub manufacturing company with three plants in Arizona. Each plant has an attached retail store and is treated as its only supplier. Each location has predicted minimum and maximum sales of each item type per period. The goal of this optimization is to maximize profit over a two-month period by determining the optimal number of each product to make at each location. In this problem, two types of constraints exist: local & global. Local constraints are unique to each plant/store combination, while the global constraints refer to company-wide resources.

4 Mathematical Formulation & Model

This problem is set up as a mixed integer programming problem, with a single objective function to be maximized (representing profit). The decision variable are set to be the production of each product item, indexed by location and period. In determining the maximum profit given the constraints, the solution will also yield the optimal production plan for Golf-Sport.

4.1 Assumptions

Two main assumptions were made in order to create an effective model for this problem. First, this model represents the company's production as if it were not a going-concern. The model is constructed to optimize profit over two months, thereby not building up any additional inventory for carryover into a third month. By default, the model will not have any inventory left at the end of the second month. If further information were to arise along with a modeling need for a third month, this could be accomplished by expanding the current model.

The second assumption used is the carryover of graphite stock between months. Advertising and graphite allotments are companywide in this problem, and it is clearly stated that the amount of money earmarked for advertising each month will not carryover into the next month. Graphite seems to have no such restriction, so it appears reasonable to incorporate the possibility of carryover of graphite stock.

4.2 Indices

Index	Location
C	Chandler
G	Glendale
T	Tucson
Index	Period Description
1	Produced & Sold in Period 1
2	Produced & Sold in Period 2
3	Produced in Period 1 & Sold in Period 2

4.3 Variables

Index	Product
s	Steel Shafts
g	Graphite Shafts
i	Forged Iron Heads
w	Metal Wood Heads
h	Titanium Insert Wood Heads
v	Set, Steel Shafts, Metal Heads
u	Set, Steel Shafts, Insert Heads
y	Set, Graphite Shafts, Metal Heads
z	Set, Graphite Shafts, Insert Heads

4.4 Objective

The objective of this model is to maximize profit, represented as the difference between sales revenue and cost. For the purposes of this function let R , C & I represent revenue, cost, and inventory cost, respectively. In this function, x represents product production. The function can be expressed as:

$$f(x) = \sum_{i=1}^3 Rx_i - \sum_{i=1}^3 Cx_i - \sum Ix_3$$

The total revenue from all products produced and sold, less the cost of manufacture and inventory, is represented in this function. While designed for a two-period optimization, the base function can be easily expanded to accomodate for more periods of production and sales.

4.5 Constraints

In construction of this model, the following constraints were constructed within the parameters set forth by the problem:

$$\sum cx_{1_i} + \sum cx_{3_i} \leq L_i \tag{1}$$

$$\sum cx_{2_i} \leq L_i \tag{2}$$

$$\sum cx_{1_i} + \sum cx_{3_i} \leq P_i \tag{3}$$

$$\sum cx_{2_i} \leq P_i \tag{4}$$

$$\sum cx_{1_i} + \sum cx_{3_i} \leq A \tag{5}$$

$$\sum Lx_{2_i} \leq A \tag{6}$$

5 Implementation

5.1 Hardware & Software

The problem was coded and solved in MATLAB R2017b on a personal computer. The code does not appear to be backward-compatible due to the use of the *optimproblem* implementation. The small scale of this optimization does not necessitate a computer beyond average specifications.

5.2 Coding

Code

5.3 Algorithm

This problem was solved using the *intlinprog* function in MATLAB. This function conducts a linear relaxation on the problem, in order to set the upper bound. Then, the algorithm iterates through a branch-and-cut approach to solve for the optimal integer solution. By default, the function will iterate through 1,000,000 iterations prior to terminating. Additionally, the function applies basic heuristics in the form of Gomory cuts. Gomory cuts

6 Solution

Explain results, compare to linear relaxation of the problem
maybe insert a section with a graph or chart to visualize production

7 Further Analysis

Conduct range sensitivity, examine extra problems in case study

8 Conclusions & Implications

8.1 Mathematical Approach

Evaluate our approach, suggest changes in future. EFFICIENCY

8.2 Company Approach

Discuss optimized solution, range sensitivity, business suggestions

References

[1]