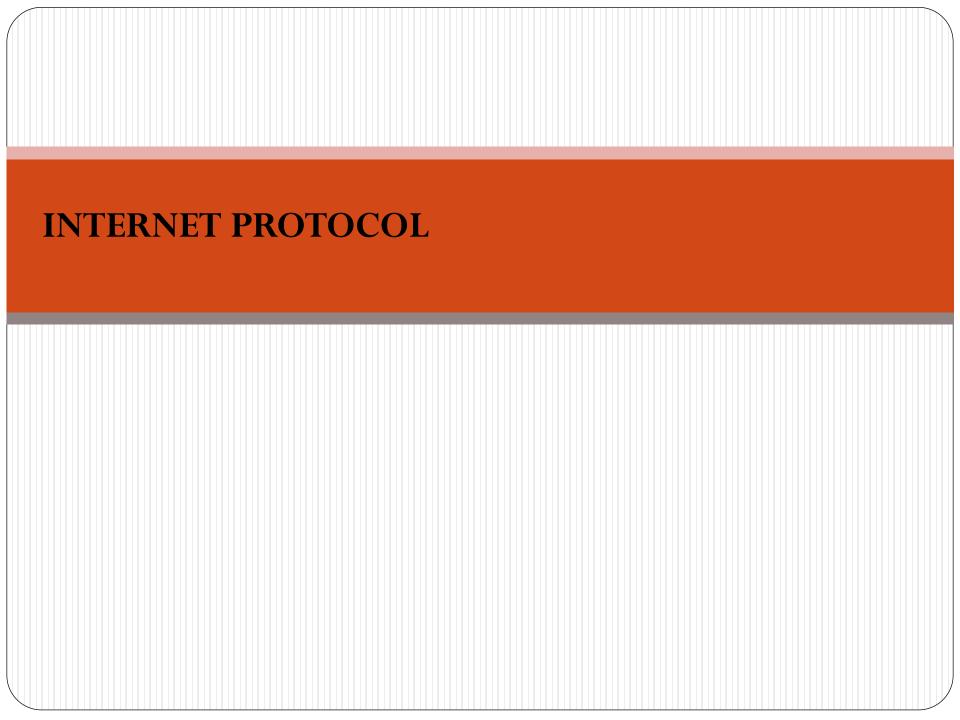
INTERNET SYSTEMS

6CCS3INS

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Outline

- The Internet Protocol (IP)
 - IPv4 header
 - IP Fragmentation & Reassembly
 - IPv6 header

Problems

- How can I send a message across the internet that will reach the right destination?
- How can I send messages of larger size than some networks along the route can carry?
- How can I know what protocols are being used to send this message (and so interpret the message correctly)?

Overview

• Fragmentation

• IPv4 header

• IPv6 header

• Path MTU

• ICMP

The Internet Protocol

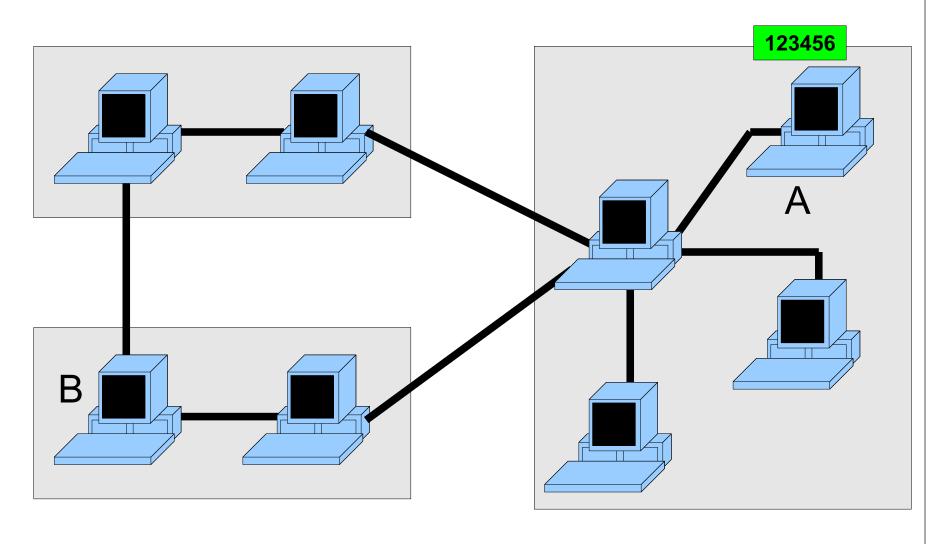
- Used to send messages across the internet
- To cope with physical network bandwidth, messages may be split into **fragments** and each fragment sent individually
- IP adds a **header** to each fragment to form an IP **datagram**

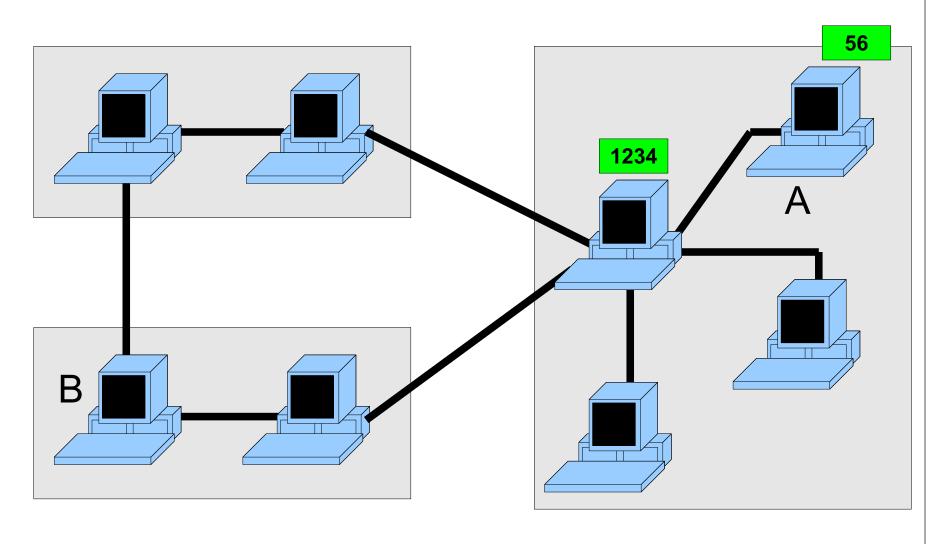
Issues not solved by IP

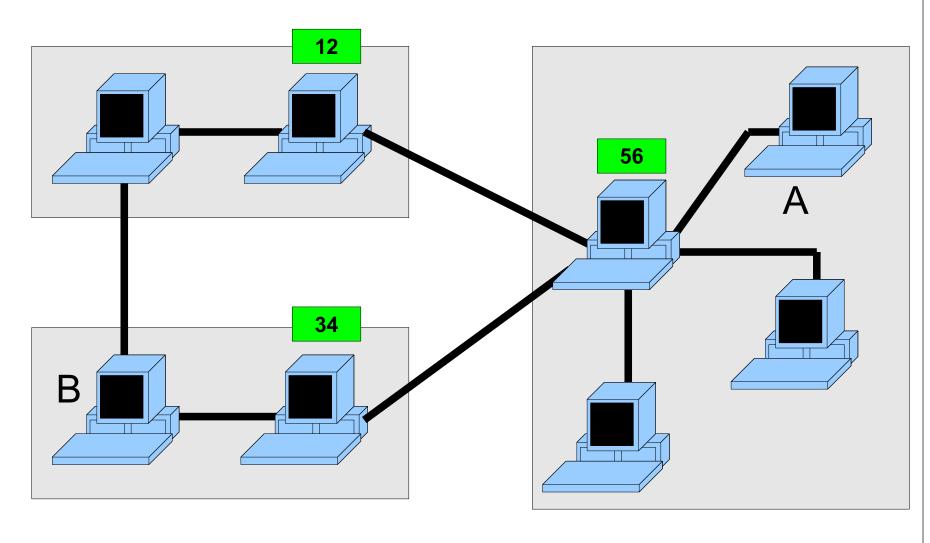
- It is unreliable, meaning that it provides no guarantees about the message successfully reaching its destination
- The message may:
 - Arrive corrupted
 - Arrive with fragments out of order
 - Arrive multiple times (duplicated)
 - Be lost entirely
- Higher level protocols, such as TCP, add reliability to IP

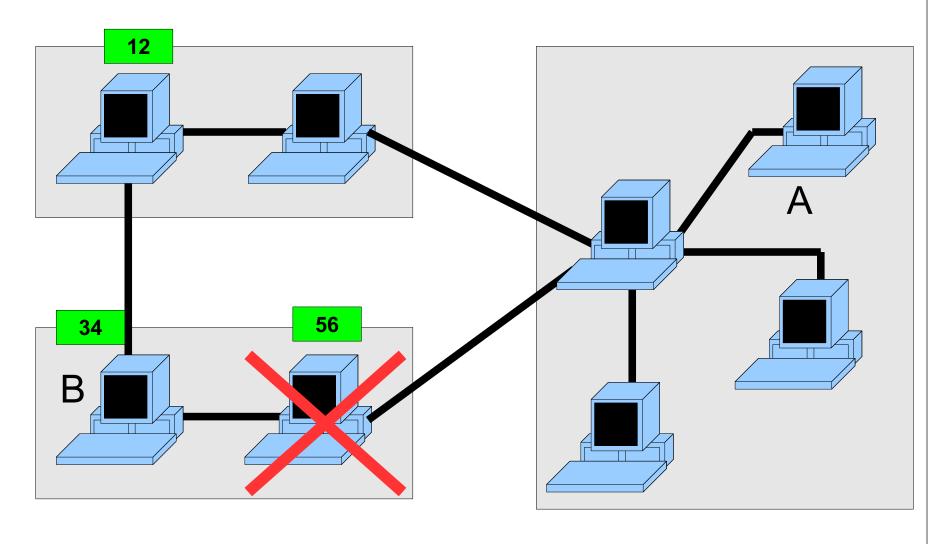
Fragmentation

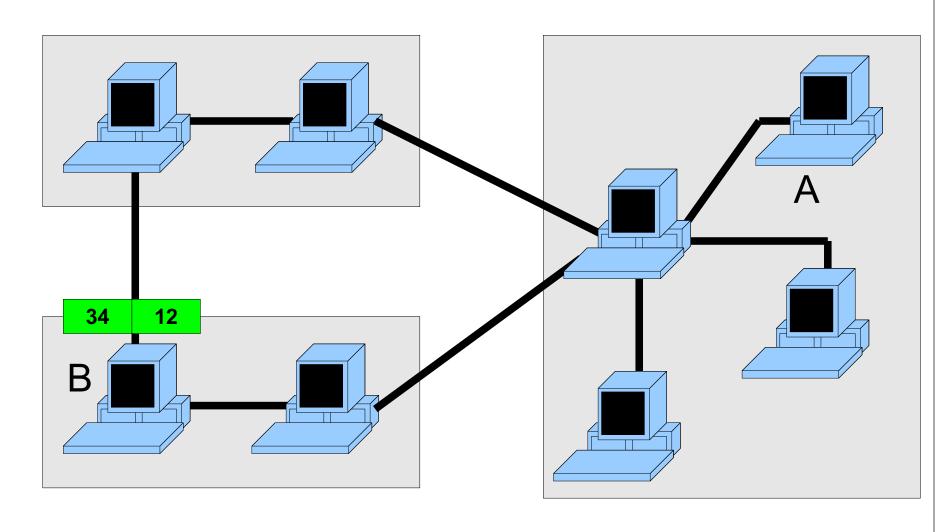
- Each network has a limit on the size of message it can carry: the **Maximum Transmission Unit (MTU)** of the network
- The sender fragments messages to fit the local MTU
- In IPv4, the fragments may be further fragmented on moving from one network to another with smaller MTU
- Fragments are reassembled at the destination
- In fragmentation, every fragment must be a multiple of 8 bytes (64 bits) except the last one



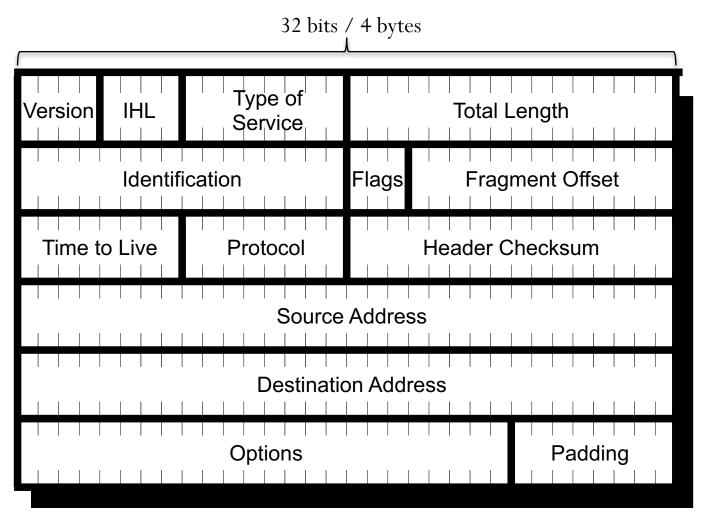








IPv4 header



An IP header with no options is 20 bytes long

Version field

- The version of IP being used
- 4 bits: 0 15
- 0100 (4) for IPv4
- 0110 (6) for IPv6

Internet header length (IHL) field

- Length of this header in 32-bit words
- 4 bits: 0 15
- Minimum value 5, to include all header fields
- Header length must be a multiple of 32 bits
- Padding field makes it up to correct length

<u>IHL</u>	<u>Header in bits</u>	
5	160	
6	192	
10	320	

•

Type of service field

- Used by message sender to specify how message should be handled in transmission
- 8 bits

- Originally used as set of flags specifying choices of the kind of service required, e.g.
 - Higher throughput vs higher reliability
- Redefined by more recent protocols handling network congestion

Total length field

- Total number of bytes in the whole datagram
- 16 bits : 0 65,535
- Header plus fragment data
- Measured in bytes/octets (multiples of 8 bits)
- Minimum: 20

Identification field

- Identifies whole of fragmented message
- 16 bits : 0 65,535
- For a receiving host
 - Every fragment with same Identification from same sending host is part of the same message
 - Can then reconstruct message from fragments

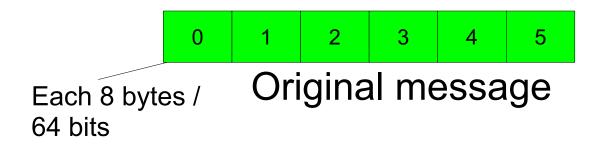
Flag fields

- A flag is an informational bit (0 = no, 1 = yes)
- 3 bits in the IPv4 header = 3 flags
- **Reserved**: First bit must always be 0
- Don't Fragment (DF): This data should not be fragmented during transmission
- More Fragments (MF): This fragment is followed by others in the message

Fragment offset field

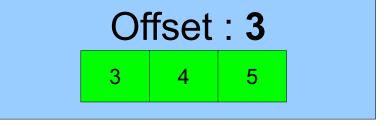
- The fragment offset (FO) field specifies where in the original message the fragment in this datagram starts
- Used to re-assemble message from fragments in the correct order
- It is a position from the start of the message counted in 8 bytes (64 bits) units, so FO=3 means this fragment starts after 24 bytes in the original message (8 x 3 = 24)
- The FO field is 13 bits long, so allows values 0 8191

Fragment Offset illustration





IP datagram 1



IP datagram 2

Time to live field

- How long this datagram is allowed to remain in the system before being deleted
- 8 bits : 0 255

- Every host handling the message must reduce Time to
 Live by 1
- When it reaches 0, the datagram should be destroyed
- Prevents data from circling round internet in an endless loop when they cannot reach their intended destination

Protocol field

- The host-to-host transport layer protocol being used by the message sent over IP
- 8 bits
- Protocols include TCP: 6, UDP: 17, ICMP: 1
- Each has a numeric identifier
- Protocol identifiers are assigned and maintained by the Internet Assigned Numbers Authority (IANA), part of ICANN
- http://www.iana.org (Protocol Numbers)

Header checksum field

- Checksum on the header, 16 bits
- Sender, receiver perform same check
- Used to ensure integrity in transmission
 - 1. One's complement sum of 16 bit words in header
 - 2. One's complement of final value

One's complement sum				
	11111111	11111100		
	00000000	00000100		
1	00000000	00000000		
		1		
	00000000	00000001		

One's complement
00110011 10101010
11001100 01010101

Source and destination fields

- IP address of sending/source host
- IP address of destination host (intended recipient)

• 32 bits each

10001001 01001001 00001001 11101000

137.73.9.232

Options and padding

- Optional arguments usable by IP processing software
- Variable number of bits
- Cover routing options, tracing of route etc.
- Not used much in practice
- See IPv4 specification (RFC 791) for more
- Followed by padding of 0s to make header up to multiple of 32 bits

Changes in fragmentation

- In IPv4, when moving from one network to another with a smaller MTU, datagrams may be re-fragmented
- Fragments must be a multiple of 8 bytes except for the last one (because the Fragment Offset is counted in 8 byte units)
- Some of the header fields will remain the same in the fragments, others will change
- For example, the Total Length field gives the length for the datagram plus header, so changes in fragmentation

More Fragments flag changes

• More Fragments Flag becomes 1 (Yes) in all but the last fragment, where it is the same as in the original datagram

MF: No

MF: Yes

MF: No

MF: Yes

MF: Yes

MF: Yes

MF: No

Fragment Offset field changes

- On fragmentation of message M, the Fragment Offset is the sum of
 - The Fragment Offset in M's header
 - The new fragment's position in M, counted in 8-byte words from the start of the fragment data

FO: 0

FO: 0

FO: 12

FO: 0

FO: 6

FO: 12

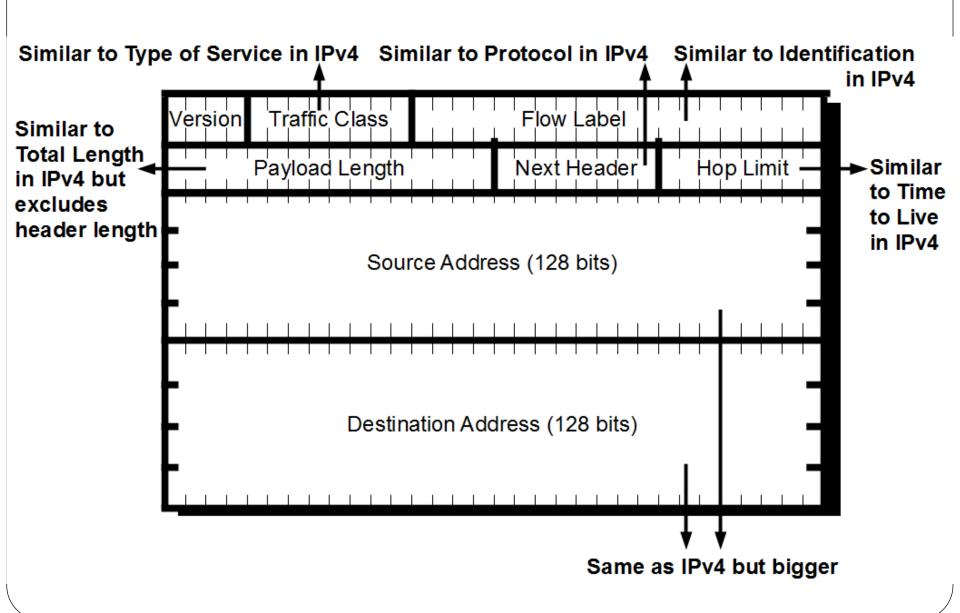
FO: 24

Header changes from IPv4 to IPv6





IPv6 header



IPv6 options

- IPv6 also allows options but, unlike IPv4, they are not all in one block inside the header
- Instead, an IPv6 header can be followed by an extension header containing optional data
- An extension header can be followed by another extension header: a chain of options
- The Next Header field is used to specify that the next block of data is an extension header, and which kind.

Extension headers

<u>IPv6 Header</u>	TCP	
Next Header	Header,	
= TCP	Body	

IPv6 Header	Routing Header	TCP
Next Header	Next Header	Header,
= Routing	= TCP	Body

IPv6 Header	Routing Header	Fragment Header	TCP
Next Header	Next Header	Next Header	Header,
= Routing	= Fragment	= TCP	Body

Fragment option

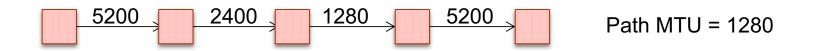
- One example of an IPv6 extension header is the fragment extension header
- If a message is fragmented (by the sender), the receiver has to be able to re-assemble it
- In order to do this, a Fragment Offset field is needed, as in IPv4
- Where a IPv6 packet contains a fragment of a message it will have a fragment extension header containing the fragment offset
- Re-assembly occurs just as in IPv4

MTU values

- On configuring a router, the MTU for each connection to a local or remote network is set
- It is set by default to the maximum value for the network access protocol used to communicate with the network, e.g. 1500 for Ethernet
- To see the MTUs obtained by IP software:
- netstat -nr to find the interface (network access protocol instantiation) of each route
- netstat -i to get the MTU of an interface

Path MTU

- IPv6 does not use fragmentation in transport
- It requires the sender to ensure messages are small enough to cross the networks **before** sending them
- The Path MTU is the minimum MTU on the path from sender to receiver



- Sender fragments message to ensure all parts are smaller than the Path MTU
- Path MTU can also be used with IPv4

Path MTU discovery algorithm

- 1. Initially assume that the Path MTU is the MTU of the first hop in the path (link to first router)
- 2. Fragments message to current assumed Path MTU, sends first fragment
- 3. If fragment reaches a link where it is too large to pass through, because the MTU is too small:
 - a) An ICMP Packet Too Big messages is return to the sender, stating that smaller MTU
 - b) Sender sets its Path MTU to MTU in ICMP message
 - c) Try again from step 2 above
- 4. If the first fragment reaches its destination, rest of message is sent, fragmented to known Path MTU

IPv4 Don't Fragment

- If the Don't Fragment Flag is set in an IPv4 message, but the message needs to be fragmented to move onto a network with a lower MTU, an error message is sent by ICMP
- ICMP code 4: fragmentation needed and DF set
- Sent back to the IP message sender

Internet Control Message Protocol

- ICMP is part of IP and used to report information on IP processing
- ICMP reports errors and other info regarding IP processing back to the datagram's source
- Examples of ICMP messages include
 - destination unreachable
 - echo request and reply (generated by ping)
 - redirect
 - time exceeded
 - router advertisement and router solicitation

RFCs

• Internet Protocol v4: 791

• Internet Procotol v6: 2460

• Internet Control Message Protocol 792

• Fragment reassembly: 815

• Path MTU Discovery: 1981

• http://www.tcpipguide.com/free

Extra Reading

- 1. W. Richard Stevens, "TCP/IP Illustrated, Volume 1: The Protocols", Addison-Wesley Professional computing series, 1994
- 2. James F. Kurose and Keith W. Ross, "Computer Networking. A Top-Down Approach Featuring the Internet", Addison-Wesley, 2000
- 3. The TCP/IP guide: http://www.tcpipguide.com/