The Envelope Theorem: Shephard's Lemma, Hotelling's Lemma, etc.

Suppose g(x, a) where a is a parameter.

Choose *x* to max the function.

In general $x^* = x(a)$.

Therefore

g((x(a), a)) is the maximum value of $g(\cdot)$ given a.

Call this a value function

$$M(a) \equiv g(x(a), a).$$

The profit function $\pi(p, w) = pf(x(p, w) - w'x(p, w))$ is an example.

Returning to the general form

Differentiate both sides of this identity with reference to a

but,

4 Substitute 3 into 2

$$\frac{dM(a)}{da} = \frac{\partial g(x(a), a)}{\partial a}$$

(This is the Envelope Theorem)

which we often write

$$\equiv \frac{\partial g(x,a)}{\partial a} \bigg|_{x=x(a)}.$$

Return to the example value function $\pi(p, w)$. By the envelope theorem

$$\frac{\partial \pi(p, w)}{\partial p} = \frac{\partial \left[pf(x) - w'x \right]}{\partial p} \Big|_{x = x(p, w)}$$

$$= f(x) \Big|_{x = x(p, w)}$$

$$= f(x(p, w))$$

which tells us that $y^s = f(x(p, w)) = \frac{\partial \pi(p, w)}{\partial p}$.

(this is Hotelling's Lemma)

Now consider the cost function

$$c = c(y, w)$$

which identifies the minimum cost of producing y given w.

Note that

$$c(y, w) = w'x(y, w)$$

conditional input demand function

So, c(y, w) is a value function (minimum not maximum).

For simplicity assume *x* is a scalar (only 1 input).

If the envelope theorem applies

$$\frac{\partial c(y, w)}{\partial w_i} = \frac{\partial \left[w'x(y, w)\right]}{\partial w_i} = \frac{\partial (w'x)}{\partial w_i} \Big|_{x = x(y, w)} = x_i((y, w))$$

which is the conditional input demand function for input i.

(This is Shephard's Lemma)

We have just used the Envelope theorem to prove Shephard's Lemma and Hotelling's Lemma. Both are special cases of the Envelope theorem.