

Boolean Algebra

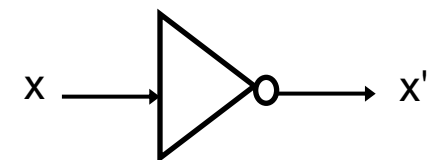
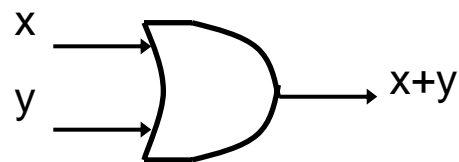
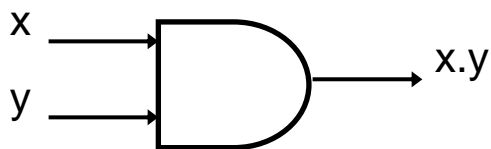
Two-valued Boolean Algebra

- Set of elements: $\{0,1\}$
- Set of operations: $\{., +, '\}$

x	y	$x \cdot y$
0	0	0
0	1	0
1	0	0
1	1	1

x	y	$x + y$
0	0	0
0	1	1
1	0	1
1	1	1

x	x'
0	1
1	0



Signals: High = 5V = 1; Low = 0V = 0

Boolean Functions

- **Boolean function** is an expression formed with binary variables, the two binary operators, OR and AND, and the unary operator, NOT, parenthesis and the equal sign.
- Its result is also a binary value.
- We usually use \cdot for AND, $+$ for OR, and $'$ for NOT. Sometimes, we may omit the \cdot if there is no ambiguity.

Complement of Boolean Functions

- Given a function, F , the **complement** of this function, F' , is obtained by **interchanging 1 with 0** in the function's output values.

Example: $F1 = x \cdot y \cdot z'$

Complement: $F1' = (x \cdot y \cdot z')'$
 $= x' + y' + (z')'$ **DeMorgan**
 $= x' + y' + z$ **Involution**

x	y	z	F1	F1'
0	0	0	0	1
0	0	1	0	1
0	1	0	0	1
0	1	1	0	1
1	0	0	0	1
1	0	1	0	1
1	1	0	1	0
1	1	1	0	1

Complement of Functions

- More general DeMorgan's theorems useful for obtaining complement functions:

$$(A + B + C + \dots + Z)' = A' \cdot B' \cdot C' \cdot \dots \cdot Z'$$

$$(A \cdot B \cdot C \dots \cdot Z)' = A' + B' + C' + \dots + Z'$$

Standard Forms

- Certain types of Boolean expressions lead to gating networks which are desirable from implementation viewpoint.
- Two Standard Forms: *Sum-of-Products* and *Product-of-Sums*
- **Literals**: a variable on its own or in its complemented form. Examples: x , x' , y , y'
- **Product Term**: a single literal or a logical product (AND) of several literals.
Examples: x , $x.y.z'$, $A'.B$, $A.B$

- **Sum Term**: a single literal or a logical sum (OR) of several literals.

Examples: x , $x+y+z'$, $A'+B$, $A+B$, $c+d+h'+j$

- **Sum-of-Products (SOP) Expression**: a product term or a logical sum (OR) of several product terms.

Examples: x , $x+y.z'$, $x.y' + x'.y.z$, $A.B + A'.B'$,
 $A + B'.C + A.C' + C.D$

- **Product-of-Sums (POS) Expression**: a sum term or a logical product (AND) of several sum terms.

Examples: x , $x.(y+z')$, $(x+y').(x'+y+z)$,
 $(A+B).(A'+B')$, $(A+B+C).D'.(B'+D+E')$

Standard Forms

- Every Boolean expression can either be expressed as sum-of-products or product-of-sums expression.

Examples:

SOP: $x'.y + x.y' + x.y.z$

POS: $(x + y').(x' + y).(x' + z')$

both: $x' + y + z$ or $x.y.z'$

neither: $x.(w' + y.z)$ or $z' + w.x'.y + v.(x.z + w')$

Minterm & Maxterm

- Consider two binary variables x, y .
- Each variable may appear as itself or in complemented form as literals (i.e. x, x' & y, y')
- For **two** variables, there are **four** possible combinations with the AND operator, namely:
$$x'.y', x'.y, x.y', x.y$$
- These product terms are called the *minterms*.
- A **minterm** of n variables is the product of n literals from the different variables.

Minterm & Maxterm

- In general, n variables can give 2^n minterms.
- In a similar fashion, a **maxterm** of n variables is the sum of n literals from the different variables.
Examples: $x'+y'$, $x'+y$, $x+y'$, $x+y$
- In general, n variables can also give 2^n maxterms.

Minterm & Maxterm

- The minterms and maxterms of 2 variables are denoted by m_0 to m_3 and M_0 to M_3 respectively:

Minterms				Maxterms	
x	y	term	notation	term	notation
0	0	$x'.y'$	m_0	$x+y$	M_0
0	1	$x'.y$	m_1	$x+y'$	M_1
1	0	$x.y'$	m_2	$x'+y$	M_2
1	1	$x.y$	m_3	$x'+y'$	M_3

Each minterm is the **complement** of the corresponding maxterm:

Example: $m_2 = x.y'$

$$m_2' = (x.y')' = x' + (y')' = x' + y = M_2$$

Canonical Form: Sum of Minterms

- What is a **canonical/normal form**?
 - ❖ A unique form for representing something.
- Minterms are product terms.
 - ❖ Can express Boolean functions using Sum-of-Minterms form.

Canonical Form: Sum of Minterms

a) Obtain the truth table. Example:

x	y	z	F1	F2	F3
0	0	0	0	0	0
0	0	1	0	1	1
0	1	0	0	0	0
0	1	1	0	0	1
1	0	0	0	1	1
1	0	1	0	1	1
1	1	0	1	1	0
1	1	1	0	1	0

Canonical Form: Sum of Minterms

- b) Obtain Sum-of-Minterms by gathering/summing the minterms of the function (where result is a 1)

$$F1 = x.y.z' = \Sigma m(6)$$

$$\begin{aligned} F2 &= x'.y'.z + x.y'.z' + \\ &\quad x.y'.z + x.y.z' + x.y.z \\ &= \Sigma m(1,4,5,6,7) \end{aligned}$$

$$\begin{aligned} F3 &= x'.y'.z + x'.y.z \\ &\quad + x.y'.z' + x.y'.z \\ &= \Sigma m(1,3,4,5) \end{aligned}$$

x	y	z	F1	F2	F3
0	0	0	0	0	0
0	0	1	0	1	1
0	1	0	0	0	0
0	1	1	0	0	1
1	0	0	0	1	1
1	0	1	0	1	1
1	1	0	1	1	0
1	1	1	0	1	0

Canonical Form: Product of Maxterms

- Maxterms are sum terms.
- For Boolean functions, the maxterms of a function are the terms for which the result is 0.
- Boolean functions can be expressed as Products-of-Maxterms.

Canonical Form: Product of Maxterms

E.g.: $F2 = \Pi M(0,2,3) = (x+y+z).(x+y'+z).(x+y'+z')$

$F3 = \Pi M(0,2,6,7)$

$= (x+y+z).(x+y'+z).(x'+y'+z).(x'+y'+z')$

x	y	z	F1	F2	F3
0	0	0	0	0	0
0	0	1	0	1	1
0	1	0	0	0	0
0	1	1	0	0	1
1	0	0	0	1	1
1	0	1	0	1	1
1	1	0	1	1	0
1	1	1	0	1	0

Canonical Form: Product of Maxterms

- Why is this so? Take F2 as an example.

$$F2 = \sum m(1,4,5,6,7)$$

- The complement function of F2 is:

$$\begin{aligned} F2' &= \sum m(0,2,3) \\ &= m0 + m2 + m3 \end{aligned}$$

(Complement functions' minterms are the opposite of their original functions, i.e. when original function = 0)

x	y	z	F2	F2'
0	0	0	0	1
0	0	1	1	0
0	1	0	0	1
0	1	1	0	1
1	0	0	1	0
1	0	1	1	0
1	1	0	1	0
1	1	1	1	0

Canonical Form: Product of Maxterms

From previous slide, $F2' = m_0 + m_2 + m_3$

Therefore:

$$\begin{aligned} F2 &= (m_0 + m_2 + m_3)' \\ &= m_0' \cdot m_2' \cdot m_3' && \text{DeMorgan} \\ &= M_0 \cdot M_2 \cdot M_3 && m_x' = M_x \\ &= \Pi M(0,2,3) \end{aligned}$$

- Every Boolean function can be expressed as either Sum-of-Minterms or Product-of-Maxterms.

Conversion of Canonical Forms

■ Sum-of-Minterms \Rightarrow Product-of-Maxterms

- ❖ Rewrite minterm shorthand using maxterm shorthand.
- ❖ Replace minterm indices with indices not already used.

$$\text{Eg: } F1(A,B,C) = \sum m(3,4,5,6,7) = \prod M(0,1,2)$$

■ Product-of-Maxterms \Rightarrow Sum-of-Minterms

- ❖ Rewrite maxterm shorthand using minterm shorthand.
- ❖ Replace maxterm indices with indices not already used.

$$\text{Eg: } F2(A,B,C) = \prod M(0,3,5,6) = \sum m(1,2,4,7)$$

Conversion of Canonical Forms

■ Sum-of-Minterms of $F \Rightarrow$ Sum-of-Minterms of F'

- ❖ In minterm shorthand form, list the indices not already used in F .

$$\text{Eg: } F1(A,B,C) = \sum m(3,4,5,6,7)$$

$$F1'(A,B,C) = \sum m(0,1,2)$$

■ Product-of-Maxterms of $F \Rightarrow$ Prod-of-Maxterms of F'

- ❖ In maxterm shorthand form, list the indices not already used in F .

$$\text{Eg: } F1(A,B,C) = \prod M(0,1,2)$$

$$F1'(A,B,C) = \prod M(3,4,5,6,7)$$

Conversion of Canonical Forms

- **Sum-of-Minterms of $F \Rightarrow$ Product-of-Maxterms of F'**
 - ❖ Rewrite in maxterm shorthand form, using the same indices as in F .
Eg: $F1(A,B,C) = \sum m(3,4,5,6,7)$
 $F1'(A,B,C) = \prod M(3,4,5,6,7)$
- **Product-of-Maxterms of $F \Rightarrow$ Sum-of-Minterms of F'**
 - ❖ Rewrite in minterm shorthand form, using the same indices as in F .
Eg: $F1(A,B,C) = \prod M(0,1,2)$
 $F1'(A,B,C) = \sum m(0,1,2)$

Binary Functions

- Given n variables, there are 2^n possible minterms.
- As each function can be expressed as sum-of-minterms, there could be 2^{2^n} different functions.
- In the case of two variables, there are $2^2=4$ possible minterms; and $2^4=16$ different possible binary functions.
- The 16 possible binary functions are shown in the next slide.

Binary Functions

x	y	F ₀	F ₁	F ₂	F ₃	F ₄	F ₅	F ₆	F ₇
0	0	0	0	0	0	0	0	0	0
0	1	0	0	0	0	1	1	1	1
1	0	0	0	1	1	0	0	1	1
1	1	0	1	0	1	0	1	0	1
Symbol			.	/		/		⊕	+
Name			AND	x, but not y		y, but not x		XOR	OR

x	y	F ₈	F ₉	F ₁₀	F ₁₁	F ₁₂	F ₁₃	F ₁₄	F ₁₅
0	0	1	1	1	1	1	1	1	1
0	1	0	0	0	0	1	1	1	1
1	0	0	0	1	1	0	0	1	1
1	1	0	1	0	1	0	1	0	1
Symbol		↓	⊙	'	⊂	'	⊃	↑	
Name		NOR	XNOR					NAND	

Logic Gates

Introduction

- A **Logic Gate** is an electronic circuit capable of making logical decisions.
- It has one output and one or more inputs.
- Basic building blocks of **digital systems**.
- 0 and 1.
- 0 means 0 to 3 V.
- 1 means 3.5 V to 5 V.
- 3 V to 3.5 V is undefined.

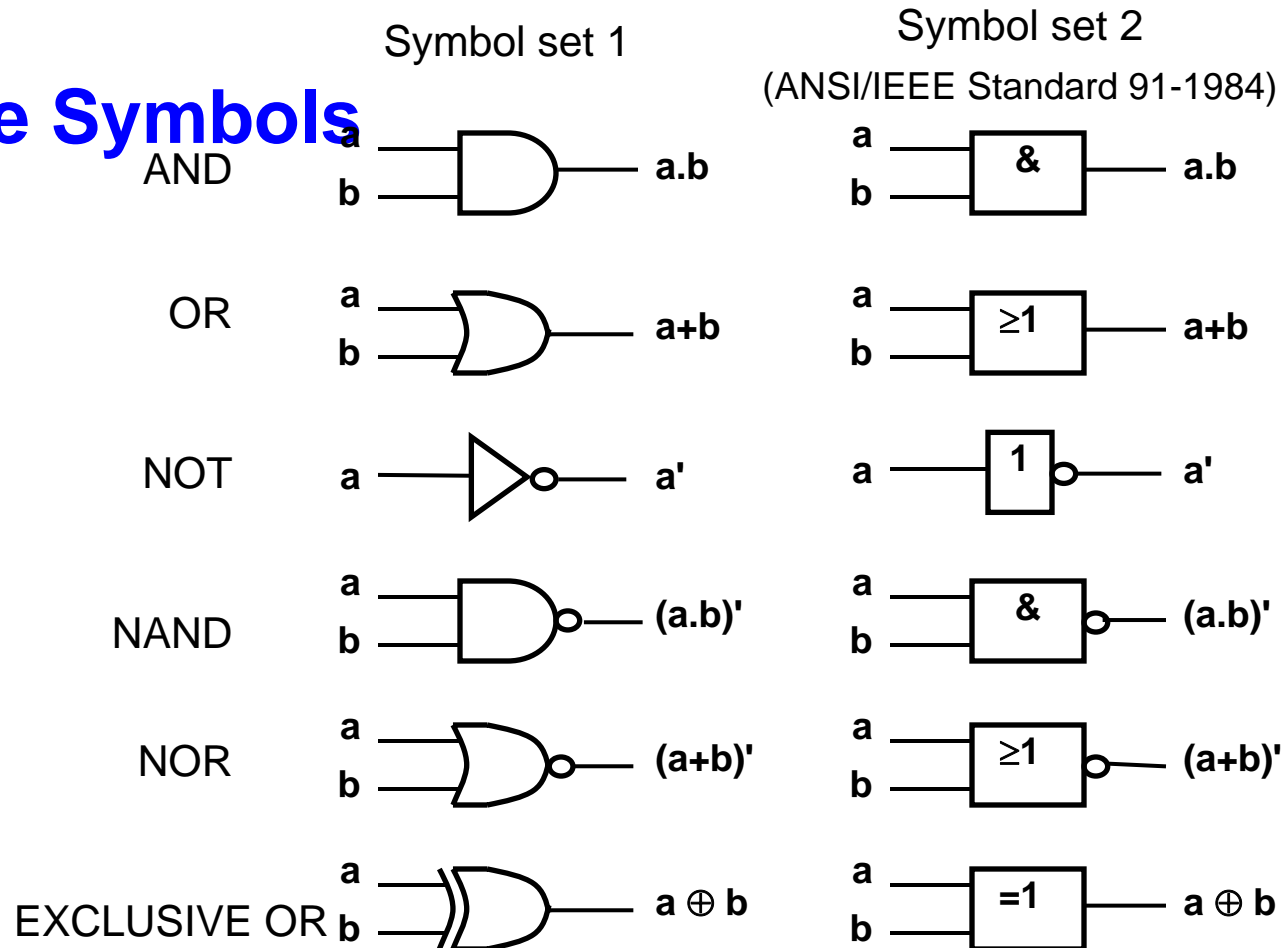
Logic Gates and Circuits

- Logic Gates
 - ❖ The Inverter
 - ❖ The AND Gate
 - ❖ The OR Gate
 - ❖ The NAND Gate
 - ❖ The NOR Gate
 - ❖ The XOR Gate
 - ❖ The XNOR Gate
- Drawing Logic Circuit
- Analysing Logic Circuit

- Universal Gates: NAND and NOR
 - ❖ NAND Gate
 - ❖ NOR Gate
- Implementation using NAND Gates
- Implementation using NOR Gates
- Implementation of SOP Expressions
- Implementation of POS Expressions
- Positive and Negative Logic
- Integrated Circuit Logic Families

Logic Gates

■ Gate Symbols



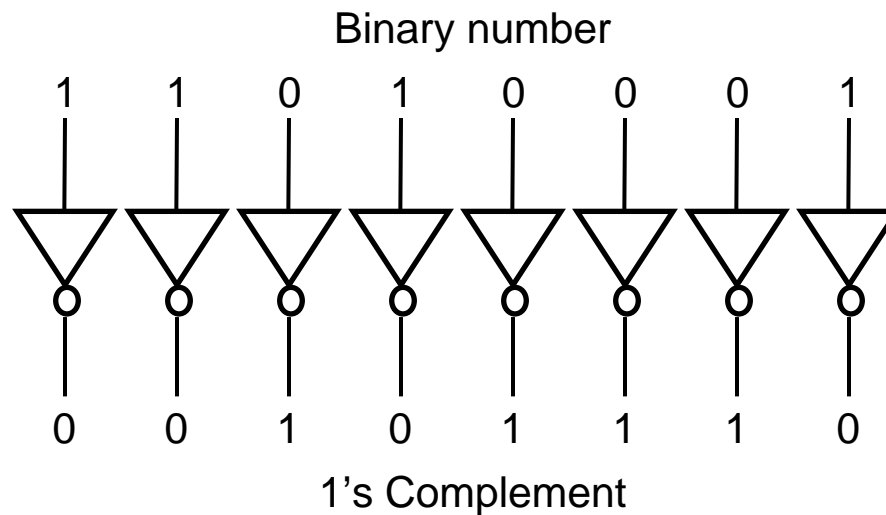
Logic Gates: The Inverter

- The **Inverter**



A	A'
0	1
1	0

- Application of the inverter: complement.



Logic Gates: The AND Gate (1/2)

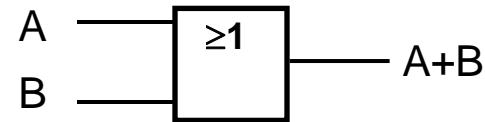
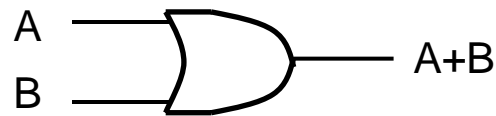
- The **AND** Gate



A	B	A . B
0	0	0
0	1	0
1	0	0
1	1	1

Logic Gates: The OR Gate

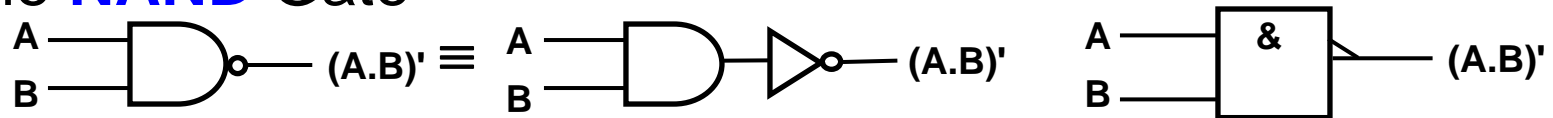
- The **OR** Gate



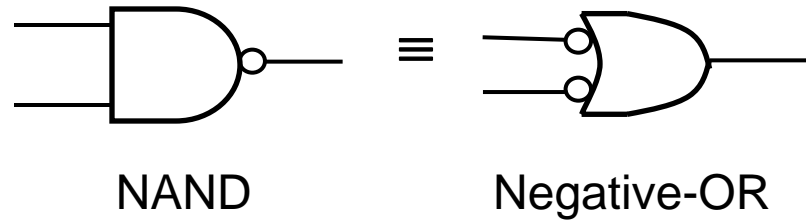
A	B	$A + B$
0	0	0
0	1	1
1	0	1
1	1	1

Logic Gates: The NAND Gate

- The **NAND** Gate

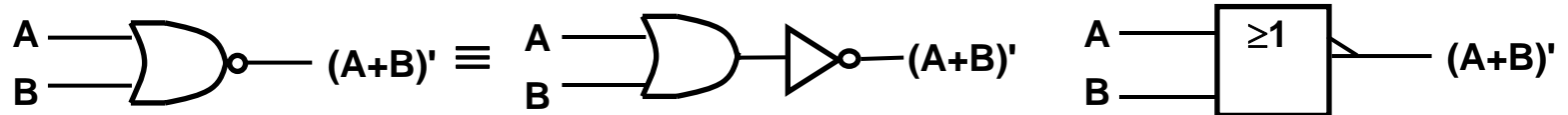


A	B	$(A.B)'$
0	0	1
0	1	1
1	0	1
1	1	0

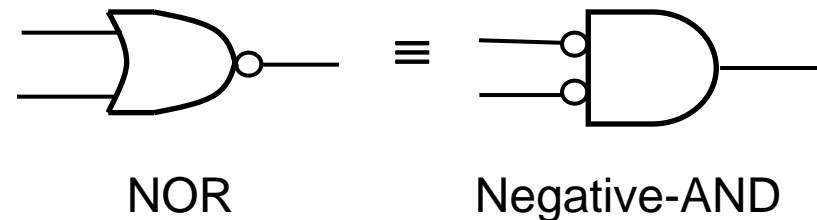


Logic Gates: The NOR Gate

- The **NOR** Gate

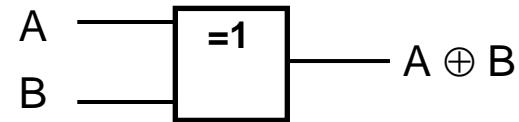
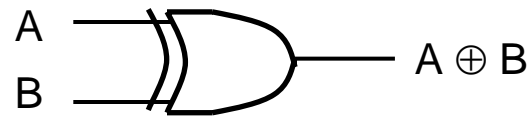


A	B	$(A+B)'$
0	0	1
0	1	0
1	0	0
1	1	0



Logic Gates: The XOR Gate

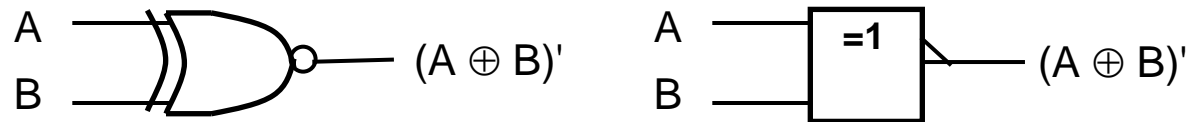
- The **XOR** Gate



A	B	$A \oplus B$
0	0	0
0	1	1
1	0	1
1	1	0

Logic Gates: The XNOR Gate

- The **XNOR** Gate



A	B	$(A \oplus B)'$
0	0	1
0	1	0
1	0	0
1	1	1