

A flexible implementation of the Harrow-Hassidim-Lloyd quantum algorithm for 2x2 symmetric matrices with equal diagonal entries

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Abstract- In recent years, due to its potential to provide an exponential speed up in many problems, Quantum Computing has become one of the most exciting and active fields of research in Physics and Computer Sciences. Among the most relevant of the tasks being studied in this field is the one concerning linear system of equations and how to solve them using Quantum Computing techniques. In this paper we develop an implementation of the quantum Harrow-Hassidim-Lloyd linear system solver algorithm for a special type of a 2 by 2 matrix. We present the circuit and test it in the IBM Quantum Computers and QASM simulators available online showing the results for two test cases and compared them to the normalize true solution computed classically.

INTRODUCTION.

For its wide range of application in various fields the problem of solving linear systems of equations is one of the most exciting tasks in which quantum computers are expected to be useful in the future. The Harrow-Hassidim-Lloyd (HHL) algorithm formulated in 2009 promises a polylogarithmic time complexity in solving systems of equations where, for whatever reason, our final goal is to find the expectation value of some operator applied to the vector x rather than x itself. A general implementation of this algorithm for arbitrary d -dimensional matrices is still currently not possible due to efficiency limitations in embedding an arbitrary system in a Hamiltonian, a step that is required for the HHL algorithm (Child & Kothari, 2010). Previously at the University of Alberta Thea Wang (2019) elaborated an implementation of the HHL algorithm for the specific case of the 2×2 linear system

$$A = \begin{bmatrix} 1.5 & 0.5 \\ 0.5 & 1.5 \end{bmatrix}; b = \begin{bmatrix} b_1 \\ b_2 \end{bmatrix}$$

Similarly, Cao, Daskin, Frankel and Kais (2012) implemented an example for a similar

fixed system of linear equations. In this project our goal

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is to go a step further and implement the quantum circuit for a more flexible case. The instance in which A is not fixed but a 2×2 symmetric matrix with equal diagonals of the form

$$A = \pi \begin{bmatrix} 2^{3-a} & 2^{3-b} \\ 2^{3-b} & 2^{3-a} \end{bmatrix}$$

Where a and b are positive integers and a is strictly less than b . We will bypass the limitations by exploiting two primary facts. The first is that as pointed out by Sadeghi (2016) the exponentials of a 2×2 antidiagonal matrix with equal entries have a friendly closed form and the second one being that:

$$e^{i[X+Y]t} = e^{iXt}e^{iYt} \leftrightarrow [X, Y] = 0$$

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1. The HHL algorithm

As explained by Dervovic et al (2018) and by Yudong Cao et al (2012) the HHL algorithm can be broken down in two key concepts: The embedding of the matrix that describe the system into a Hamiltonian and the Phase Estimation algorithm. We will walk through them in order and then we will put it all together

54 before proceeding to explain our
55 implementation

56

57 *Embedding the system*

58 The key realization of this step of the algorithm
59 is that if our system's matrix A is Hermitian and
60 has eigenvalues $|\mu_j\rangle$ and eigenvectors λ_j we can
61 make a mapping to a unitary matrix U by matrix
62 exponentiation that conserves the eigenvectors
63 and eigenvalues. This is needed because
64 Quantum gates can only perform unitary matrix
65 transformation. Under this mapping the
66 eigenvalues for the eigenvectors will be
67 transform as follows:

$$68 \quad \lambda_j \rightarrow e^{i\lambda_j t}$$

69 In the case where A does not happen to be
70 Hermitian the only relevant change we would
71 need to do would be to embed A in bigger
72 system of the form.

$$73 \quad \begin{bmatrix} 0 & A \\ A & 0 \end{bmatrix}$$

74 Which clearly is Hermitian (Harrow et al,
75 2009).

76 *Phase Estimation*

77 The phase estimation algorithm is a common
78 subroutine in several quantum algorithms in
79 which, given a target quantum register of length
80 b with all its qubits in the Hadamard positive
81 state (the $|+\rangle$ superposition state), a quantum
82 operator U , and a register in the state of one of
83 the eigenvector of the matrix A the algorithm
84 performs the transformation:

$$85 \quad |\mu_j\rangle|+\rangle \rightarrow |\mu_j\rangle|2^n \lambda_j\rangle$$

86 Where λ_j is the phase of the operator U as well
87 as the eigenvalue of the Hermitian matrix A .

88 Similarly, if the input is an arbitrary quantum
89 state b , the system would be left in a
90 superposition of the eigenvectors of the form:

$$91 \quad |\Psi\rangle = \sum_j \beta_j |2^n \lambda_j\rangle |\mu_j\rangle$$

92 Where the β s would be just the weights of the
93 linear combination of eigenvectors that spans b

94 *Putting the algorithm together.*

95 Taking this into consideration HHL works by
96 putting a quantum state $|b\rangle$ (proportional to the
97 true vector b) in the eigenvalue register of the
98 QPE while it uses the unitary matrix produced
99 by the exponentiation of A as the operator.

100 This puts the target register in a superposition
101 of the eigenvalues of A which corresponds to
102 the superposition that spans b (using the
103 eigenvectors as basis).

104 The algorithm then makes use of an additional
105 ancillary qubit to perform a controlled rotation
106 applying R_y gates conditioned on the
107 eigenvalues to the ancilla. The R_y gate is
108 defined by the following matrix transformation

$$109 \quad R_y(\theta) = \begin{bmatrix} \cos(\theta/2) & -\sin(\theta/2) \\ \sin(\theta/2) & \cos(\theta/2) \end{bmatrix}$$

110 The system is thus put in the following state as
111 result of the operation:

$$112 \quad |\Psi\rangle = \sum_j \left(\sqrt{1 - \left(\frac{C}{\lambda_j}\right)^2} |0\rangle + \frac{C}{\lambda_j} |1\rangle \right) \beta_j |2^n \lambda_j\rangle |\mu_j\rangle$$

113 This is accomplished by setting the angle of
114 rotation of the rotation to

$$115 \quad \theta = 2 \arcsin\left(\frac{C}{\lambda_j}\right)$$

116 Where C is just a proportionality constant that
117 we chose.

118 Then we clean the QPE target register by
119 applying the same gates in reverse order
120 exploiting the fact that unitary transformations
121 are reversible. This leaves the system as
122 follows:

$$123 \quad |\Psi\rangle = \sum_j \left(\sqrt{1 - \left(\frac{C}{\lambda_j}\right)^2} |0\rangle + \frac{C}{\lambda_j} |1\rangle \right) \beta_j |0\rangle |\mu_j\rangle$$

124 For reference a schematic of this circuit can be
125 found in Figure 1

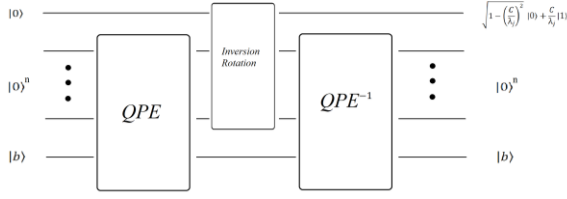


FIG 1: High level schematics of the HHL algorithm circuit

Now note that if we make measurements and post-select the states where the ancillary qubit collapses to $|1\rangle$ we are left with the state

$$|\Psi\rangle = C \sum_j \frac{\beta_j}{\lambda_j} |\mu_j\rangle |0\rangle |1\rangle$$

And if we look carefully, we will realize that we have accomplish our goal since

$$A^{-1}|b\rangle \propto C \sum_j \frac{\beta_j}{\lambda_j} |\mu_j\rangle$$

IMPLEMENTATION.

Implementing our operator

To implement the HHL algorithm for our case we rewrite our target matrix as the sum of commuting diagonal and anti-diagonal matrices.

$$\pi \begin{bmatrix} 0 & 2^{3-b} \\ 2^{3-b} & 0 \end{bmatrix} + \begin{bmatrix} 2^{3-a} & 0 \\ 0 & 2^{3-a} \end{bmatrix} = \pi \begin{bmatrix} 2^{3-a} & 2^{3-b} \\ 2^{3-b} & 2^{3-a} \end{bmatrix}$$

$$\begin{bmatrix} 0 & 2^{3-b} \\ 2^{3-b} & 0 \end{bmatrix} \begin{bmatrix} 2^{3-a} & 0 \\ 0 & 2^{3-a} \end{bmatrix} = \begin{bmatrix} 2^{3-a} & 0 \\ 0 & 2^{3-a} \end{bmatrix} \begin{bmatrix} 0 & 2^{3-b} \\ 2^{3-b} & 0 \end{bmatrix}$$

This means that if we take the matrix exponential of A, we can rewrite it as the product of the exponential of its diagonal component (X) with the exponential of antidiagonal component (Y):

$$e^{i[X+Y]t} = e^{iXt} e^{iYt}$$

Now we are going to use the fact from Sadeghi (2016) that the exponential of a 2 by 2 antidiagonal matrices with repeated entries can in general be written as.

$$e^{i\pi \begin{bmatrix} 0 & 2^{3-b} \\ 2^{3-b} & 0 \end{bmatrix} t} = \begin{bmatrix} \cosh(i\pi 2^{3-b} t) & \sinh(i\pi 2^{3-b} t) \\ \sinh(i\pi 2^{3-b} t) & \cosh(i\pi 2^{3-b} t) \end{bmatrix}$$

Which we can further simplify using the identities

$$\sinh(ix) = i\sin(x)$$

$$\cosh(ix) = \cos(x)$$

And assign

$$\theta = -\pi 2^{3-b}$$

to rewrite the matrix as:

$$e^{iYt} = \begin{bmatrix} \cos(\theta t) & -i\sin(\theta t) \\ -i\sin(\theta t) & \cos(\theta t) \end{bmatrix}$$

Note that is no more than the Rx-gate

$$R_X(2\theta t) = \begin{bmatrix} \cos(2\theta t/2) & -i\sin(2\theta t/2) \\ -i\sin(2\theta t/2) & \cos(2\theta t/2) \end{bmatrix}$$

Setting $t=1/4$ this can be easily implemented by

$$R_X\left(\frac{\theta}{2}\right) = \begin{bmatrix} \cos(\theta/4) & -i\sin(\theta/4) \\ -i\sin(\theta/4) & \cos(\theta/4) \end{bmatrix}$$

Likewise, the exponential of the diagonal matrix X

$$e^{iXt} = \begin{bmatrix} e^{i\pi 2^{3-a} t} & 0 \\ 0 & e^{i\pi 2^{3-a} t} \end{bmatrix}$$

can also be simply implemented by a combination of X-gates and R-phi

$$\begin{bmatrix} e^{i\pi 2^{3-a} t} & 0 \\ 0 & e^{i\pi 2^{3-a} t} \end{bmatrix} = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 0 & e^{i\pi 2^{3-a} t} \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 0 & e^{i\pi 2^{3-a} t} \end{bmatrix}$$

$$= X R_{(\emptyset)} X R_{(\emptyset)}$$

Where we make

$$\emptyset = i\pi 2^{3-a} t$$

The IBM Qiskit Python package conveniently has available its own implementation for this transformation.

So, our final operator for the QPE is just the controlled version of this succession of gates.

$$\exp\left(i\pi \begin{bmatrix} 2^{3-a} & 2^{3-b} \\ 2^{3-b} & 2^{3-a} \end{bmatrix} t\right) = Cnot C R_{(\emptyset)} Cnot C R_{(\emptyset)} R_y\left(\frac{\theta}{2}\right)$$

Where we are going to use $t=4$.

Implementing the inversion rotation

In the most general case for this step we would be required to implement a unitary gate U able to take the input in the register of the eigenvalues and map it to their inverses. However, in this case we are going to make use of our knowledge of the format of the input matrix and solve it in an ad-hoc manner.

It is easy to check that the eigenvalues of the antidiagonal matrix Y are b and $-b$ and thus the eigenvalues of its exponential correspond to $e^{i2\pi 2^{-b}}$ and $e^{-i2\pi 2^{-b}}$ with $t=4$. The diagonal matrix simply scales the input vector by the entry in its diagonal. Hence, the eigenvalues for the product of both matrices are

$$e^{\pm i2\pi 2^{-b}} e^{i2\pi 2^{-a}} = e^{2\pi i(2^{-a} \pm 2^{-b})}$$

In our matrix $|a| < |b|$ so to apply the correct angle of rotation to invert each eigenvalue we only need to condition the rotation in the most significant qubit of the QPE registry.

This means that we can implement the inversion by simply applying two control rotation in the most significant qubit. If it's 1 we rotate by an angle of $\theta_1 = 2 \arcsin\left(\frac{C}{2\pi(2^{-a} + 2^{-b})}\right)$ and if 0 we rotate by $\theta_2 = 2 \arcsin\left(\frac{C}{2\pi(2^{-a} - 2^{-b})}\right)$.

Therefore, we can replace the whole rotation inversion by the circuit shown in figure 2 where the control reaches to the first most significant qubit in the QPE registry.

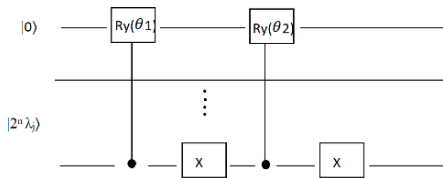


FIG 2: Circuit for the rotation inversion of the eigenvalues for our family of matrices. The rotation is condition on the first most significant qubit of the QPE registry.

TEST CASES.

We are now going to proceed to present the results we obtained for the two test cases ran in the IBM Cloud Quantum Computer and QASM simulator.

For all cases we will be using

$$C=2\pi(2^{-a} - 2^{-b})$$

Test Case 1

$$A = \begin{bmatrix} 4\pi & 2\pi \\ 2\pi & 4\pi \end{bmatrix}; b = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$$

The result of the linear system solved conventionally is:

$$x = \left(\frac{1}{\pi}\right) \begin{bmatrix} 1/3 \\ -1/6 \end{bmatrix}$$

which means that we expect the square amplitude of the corresponding quantum state to be measure to:

$$\begin{bmatrix} (C_1)^2 \\ (C_2)^2 \end{bmatrix} = \begin{bmatrix} \frac{1/9}{1/9 + 1/36} \\ \frac{1/36}{1/9 + 1/36} \end{bmatrix} = \begin{bmatrix} 0.8 \\ 0.2 \end{bmatrix}$$

So, when we read the results, we expect the states that end in 1 to have the same ratio.

The circuit generated using QISKIT as well as the results from running it in the QASM simulator can be found in Figure 3 and 4 underneath

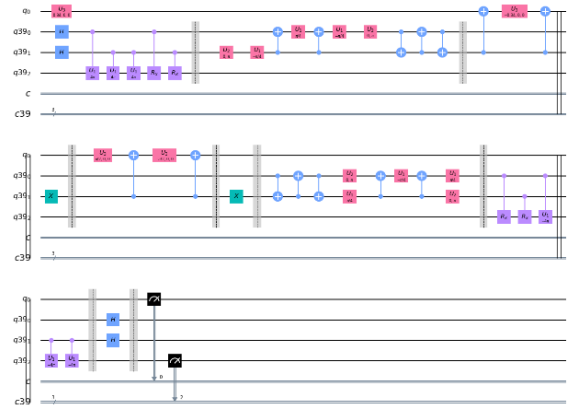


FIG 3: Quantum Circuit generated by Qiskit using the IBM subroutines for QFT and QFT⁻¹

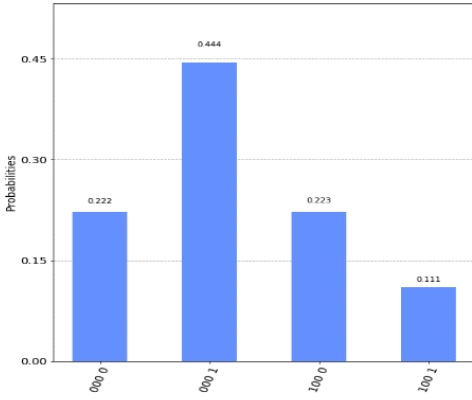


FIG 4: Histogram of results counts of the simulation of test case 1 with 8000 shots from the IBM-QASM simulator

We note from FIG 4 that

$$c_{1qasm} = \frac{0.444}{0.444 + 0.111} = 0.8$$

$$c_{2qasm} = \frac{0.111}{0.444 + 0.111} = 0.2$$

So indeed, the normalize vectors are the same. Now if we run the same circuit in the ibmq-16_melbourne quantum computer available online we get the results depicted in figure 5.

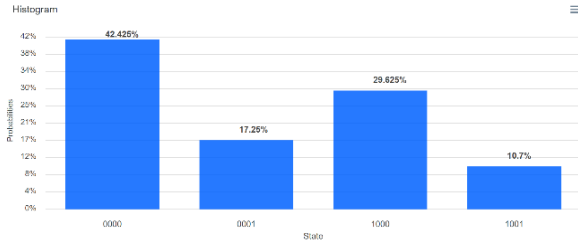


FIG 5: Histogram of counts of the simulation of test case 1 with 8000 shots from the ibmq-16_melbourne quantum computer available online

Normalizing the cases where the ancilla collapsed to 1 we get

$$c_{1melb} = \frac{29.625}{29.625 + 10.7} = 0.734$$

$$c_{2melb} = \frac{10.7}{29.625 + 10.7} = 0.265$$

Thus, the L1 norm of the difference between the amplitudes squares is

$$\|c_{qasm} - c_{melb}\|_1 = 0.0785 + 0.0784 = 0.1569$$

Test Case 2

$$A = \begin{bmatrix} 4\pi & 1\pi \\ 1\pi & 4\pi \end{bmatrix}; b = \begin{bmatrix} -\sin(\pi/8) \\ \cos(\pi/8) \end{bmatrix}$$

The solution to 4 significant digits is

$$x = \begin{bmatrix} -0.05208 \\ 0.08654 \end{bmatrix}$$

Thus, we expect

$$\begin{bmatrix} (c_1)^2 \\ (c_2)^2 \end{bmatrix} = \begin{bmatrix} \frac{0.05208^2}{0.05209^2 + 0.08654^2} \\ \frac{0.08654^2}{0.05209^2 + 0.08654^2} \end{bmatrix} = \begin{bmatrix} 0.2658 \\ 0.7341 \end{bmatrix}$$

Likewise, we present the circuit generated using Qiskit as well as the results from running it in the QASM simulator in Figure 5 and 6.

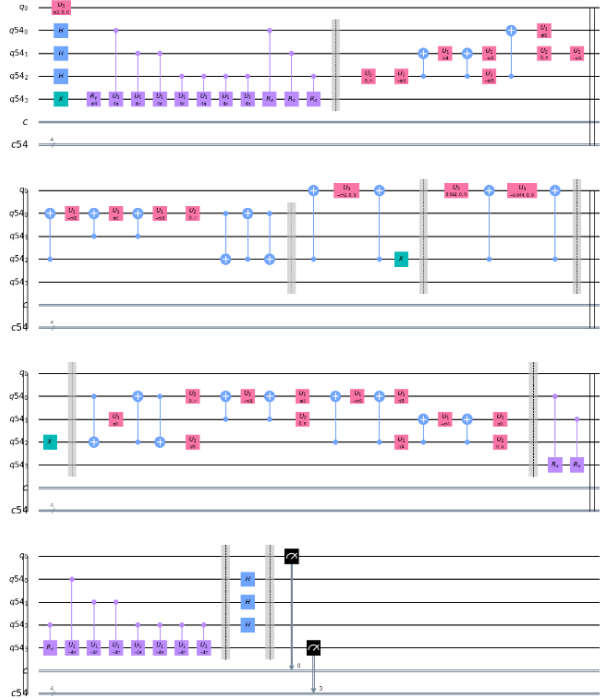
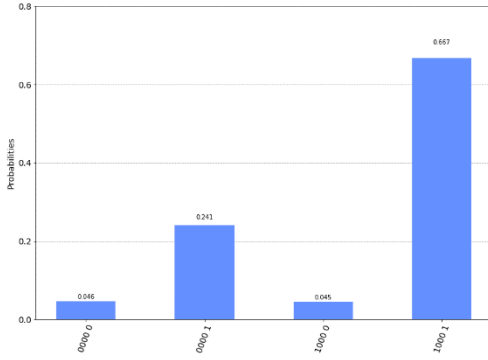


FIG 5: Quantum Circuit for test case 2 generated using Qiskit subroutines for QFT and QFT⁻¹

285



286 FIG 6: Histogram of results counts of the simulation
287 of test case 1 with 8000 shots from the IBM-QASM
288 simulator

289 And again, we note that:

$$290 \quad c_{1qasm} = \frac{0.241}{0.241 + 0.667} = 0.2654$$

$$291 \quad c_{2qasm} = \frac{0.667}{0.241 + 0.667} = 0.7345$$

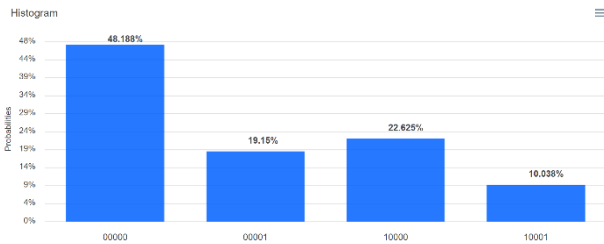
292

293 The L1 difference between the two normalized
294 vectors is

$$295 \quad ||c_{qasm} - c_{real}||_1 = 0.004 + 0.004 = 0.008$$

296 Similarly, running the same circuit with the
297 same number of shots in the Melbourne back-
298 end we get the following output:

299



300 FIG 5: Histogram of results counts of the simulation of test
301 case 1 with 8000 shots from the ibmq-16_melbourne
302 quantum computer available online

303

304 Looking at the states that end in 1 again we see
305 that

$$306 \quad c_{1melb} = \frac{10.038}{19.150 + 10.038} = 0.3439$$

307

$$c_{2melb} = \frac{19.150}{19.150 + 10.038} = 0.6561$$

308

309 And the L1 norm of the difference between the
310 amplitude square of the states gotten by the real
311 device and the QASM is

312

$$313 \quad ||c_{qasm} - c_{melb}||_1 = 0.0785 + 0.1087 = 0.1872$$

314

315 LIMITATIONS AND FURTHER WORK.

316 Even though we have accomplished our goal
317 and have been able to extract information about
318 x with reasonable precision using the currently
319 noisy quantum computers. The procedure we
320 have implemented is still no more than an
321 example with an ad-hoc solution.

322 Throughout the implementation we have relied
323 heavily in the format of our input. The
324 prescribed format of our matrix made it possible
325 to easily factorized it into unitary operators and
326 readily gave up information about its
327 eigenvalues. We relied upon this information
328 during the rotation inversion which is the key
329 step of HHL. It remains to be implemented a
330 general unitary operation that takes the result of
331 the QPE algorithm and maps directly to its
332 inverse, so we do not need to change the angle
333 of rotations with every input in our instance of
334 HHL.

335 In addition to this every other limitation of the
336 algorithm described by Cao et al (2012) and
337 Childs and Kothari (2010) such as the problem
338 of the preparation of an arbitrary state $|b\rangle$
339 applies as well to this instance

340 In conclusion even though we have been able to
341 put together a working example of the HHL
342 algorithm this is still a didactical
343 implementation.

344

345

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