

# 0/1 - Matrices

Alexandra Lassota

Consider the problem of deciding whether the polytope  $P = \{x \mid Ax = b, 0 \leq x \leq 1, x \in \mathbb{N}^{2^m}\}$  is non-empty where  $A$  is the complete 0/1-matrix of size  $m \times 2^m$ .

We denote by  $x^i$  the  $i$ -th entry of a vector  $x$ . Let  $a_i$  be a column of  $A$  for  $i \in \{1, \dots, 2^m\}$ . We define *pairs* as tuples  $(a_i, a_j)$  of columns such that  $a_i + a_j = \mathbf{1}$ . Call  $a_j$  the *complement* of  $a_i$  and vice versa. Clearly, we can partition all columns in  $A$  into disjoint pairs and this partition is unique.

Consider the decomposition  $b = k \cdot \mathbf{1} + u$  of the right-hand side  $b$  such that  $Ax = u$  is feasible for some  $0 \leq x \leq 1, x \in \mathbb{N}^{2^m}$  and  $k$  is maximal. Call such a decomposition *maximal*.

**Theorem 1.** *Given a maximal decomposition  $b = k \cdot \mathbf{1} + u$ . If  $Ax = b$  is feasible, then we can extend each support minimal solution  $x_u$  for  $Ax = u$  to a solution  $x^* = x_u + x_k$  such that  $x_u + x_k \leq \mathbf{1}$ .*

*Proof.* Suppose the opposite, i.e., there exists a solution  $x' = x_{u'} + x'_k$  such that  $Ax' = b$  but  $x_{u'}$  is **not support minimal** or  $k$  is not maximal w.r.t. the definition of a maximal decomposition, and a support minimal solution cannot be extended to solve the original problem.

Define  $X'_1$  as the set of columns taken by  $x_{u'}$ , i.e., take the  $i$ -th column if  $x_{u'}^i = 1$ . Let  $X'_2$  contain all complements of the columns chosen for  $X'_1$ . Denote by  $X'_R$  the remaining columns of  $A$ , i.e., all columns which are not present in  $X'_1$  nor  $X'_2$ . Note that  $X'_R$  can be partitioned into  $|X'_R|/2$  many pairs.

**If  $X'_1 \cup X'_2 \neq \emptyset$ , this means that  $X'_1$  contains a pair  $(a_i, a_j)$ .** In that case, we can reduce  $u'$  by  $\mathbf{1}$  and increase  $k'$  by 1 and obtain  $u''$  and  $k''$ . There exists a solution to the new decomposition, in particular, obtain  $x_{u''}$  by setting  $x_{u'}^i = x_{u'}^j = 0$ . Similarly,  $x_k''$  is obtained by setting the  $i$ -th and  $j$ -th entry to 1.

So, let us assume that  $X'_1 \cup X'_2 = \emptyset$ . If  $k' \leq |X'_R|/2$ , meaning columns present in  $X'_R$  are sufficient to define a solution for  $Ax = k' \cdot \mathbf{1}$ . In that case, some support minimal solution  $x$  would yield that  $|X_R| > |X'_R|$  and thus, we can extend to a solution for  $Ax = b$ .

Thus, we can additionally assume that  $k > |X'_R|/2$ . Thus, we necessarily have to take some column  $a_i$  in  $X'_2$  for our solution  $x'$ . For that column, there exists a complement  $a_j$  in  $X'_1$ . Now, we can show that  $u'$  is not minimal by defining a solution with smaller support to solve a smaller  $u$ . To do so, set  $x_{u'}^i = 0$ . Meaning, we delete  $a_i$  and  $a_j$  from the sets  $X'_1$  and  $X'_2$ , respectively, and add them to  $X'_R$ . This way, we expressed the same solution but with a decomposition for a smaller  $u < u'$  using less columns. We can repeat this step for every column in  $X'_2$ . This way, we obtain a solution using only columns in  $X'_R$  to satisfy the scaled  $\mathbf{1}$ -vector, and obtain a sufficient small  $u''$ . Indeed, the  $u''$  obtained this way might not be minimal w.r.t. the maximal decomposition. However, if we can solve the problem for some  $u''$  larger than a minimal  $u$  such that the remaining solution only uses

columns in the set  $X'_R$ , we already argued above that this results also holds for the maximal decomposition.  $\square$