

Typify types

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Parser and tools for type signatures of **typify**. Essentially the language is dependent type theory, with omitted lambda constructor.

1 Formal syntax definition

Terminals:

$$\frac{}{\top : \text{type}} \boxed{\text{ANY}} \quad \frac{}{\perp : \text{type}} \boxed{\text{NOTHING}} \quad \frac{}{\mathbb{1} : \text{type}} \boxed{\text{UNIT}}$$
$$\frac{}{\text{number} : \text{type}} \boxed{\text{NUMBER}} \quad \frac{}{\text{string} : \text{type}} \boxed{\text{STRING}} \quad \frac{}{\text{bool} : \text{type}} \boxed{\text{BOOLEAN}}$$

Identifiers:

$$\frac{n : \text{name}}{n : \text{type}} \boxed{\text{IDENTIFIER}}$$

Row types:

$$\frac{a : \text{type} \quad b : \text{type}}{a \wedge b : \text{type}} \boxed{\text{CONJUNCTION}} \quad \frac{a : \text{type} \quad b : \text{type}}{a \vee b : \text{type}} \boxed{\text{DISJUNCTION}} \quad \frac{a : \text{type}}{a? : \text{type}} \boxed{\text{OPTIONAL}}$$

Application:

$$\frac{a : \text{type} \quad b : \text{type}}{a \, b : \text{type}} \boxed{\text{APPLICATION}}$$

Special modifiers:

$$\frac{a : \text{type} \quad b : \text{type}}{a \times b : \text{type}} \boxed{\text{PRODUCT}} \quad \frac{a : \text{type}}{a \dots : \text{type}} \boxed{\text{VARIADIC}} \quad \frac{n : \text{type} \quad p : \text{type}}{n : p : \text{type}} \boxed{\text{NAMED}}$$

Brackets:

$$\frac{p : \text{type}}{[p] : \text{type}} \boxed{\text{SQUARE}} \quad \frac{p : \text{type}}{\{p\} : \text{type}} \boxed{\text{CURLY}}$$

Functions:

$$\frac{a : \text{type} \quad b : \text{type}}{a \rightarrow b : \text{type}} \boxed{\text{FUNCTION}}$$

Figure 1: Typify syntax rules

level	name		associativity	example
8	optional	?	postfix	$a?$
7	application	$a\ b$	left associative	$a\ b\ c \equiv (a\ b)\ c$
6	conjunction	\wedge	associative	$(a \wedge b) \wedge c \equiv a \wedge (b \wedge c)$
5	disjunction	\vee	associative	$(a \vee b) \vee c \equiv a \vee (b \vee c)$
4	ellipsis	\dots	postfix	$a\dots$
3	type	$:$	left associative	$x : a : b \equiv (x : a) : b$
2	product	\times	associative	$(a \times b) \times c \equiv a \times (b \times c)$
1	function	\rightarrow	right associative	$a \rightarrow b \rightarrow c \equiv a \rightarrow (b \rightarrow c)$

Table 1: Operator precedence

name	code variants	
any	\top	$\top\ *$
nothing	\perp	$\perp\ _ _$
unit	$\mathbb{1}$	$\mathbb{1}\ ()$
optional	$?$	$?$
conjunction	\wedge	$\wedge\ $
disjunction	\vee	$\vee\ \&$
ellipsis	\dots	\dots (either three dots, or unicode ellipsis)
type	$:$	$:$
product	\times	$\times\ ,$
function	\rightarrow	$\rightarrow\ \rightarrow$

Table 2: Special character code representations

$$\begin{aligned}
& a \times ys : b \dots \times c \rightarrow d \rightarrow e \\
& \equiv x : a \times ys : (b \dots) \times c \rightarrow d \rightarrow e \\
& \equiv (x : a) \times (ys : (b \dots)) \times c \rightarrow d \rightarrow e \\
& \equiv ((x : a) \times (ys : (b \dots)) \times c) \rightarrow d \rightarrow e \\
& \equiv ((x : a) \times (ys : (b \dots)) \times c) \rightarrow (d \rightarrow e) \\
& a \vee b? \wedge c\ d \dots \rightarrow e \\
& \equiv a \vee (b?) \wedge (c\ d) \dots \rightarrow e \\
& \equiv a \vee ((b?) \wedge (c\ d)) \dots \rightarrow e \\
& \equiv (a \vee ((b?) \wedge (c\ d))) \dots \rightarrow e \\
& \equiv ((a \vee ((b?) \wedge (c\ d))) \dots) \rightarrow e \\
& a \times y : b \rightarrow c \\
& \equiv a \times (y : b) \rightarrow c \\
& \equiv (a \times (y : b)) \rightarrow c
\end{aligned}$$

Figure 2: Examples of operator precedence

Either $a\ b \equiv \{type : "left", value : a\} \vee \{type : "right", value : b\}$
 $flatMap : (@ : Observable\ A \times f : A \rightarrow Observable\ B \vee Event\ B \vee B) \rightarrow EventStream\ B$

Figure 3: Examples of real world types