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1 Data Structures

1.1 Disjoint Set Union

```

1 struct DSU {
2     vector<int> parent, size;
3     DSU(int n) {
4         parent.resize(n);
5         size.resize(n);
6         for (int i = 0; i < n; i++) make_set(i);
7     }
8     void make_set(int v) {
9         parent[v] = v;
10        size[v] = 1;
11    }
12    bool is_same(int a, int b) { return find_set(a)
13        == find_set(b); }
14    int find_set(int v) { return v == parent[v] ? v :
15        parent[v] = find_set(parent[v]); }
16    void union_sets(int a, int b) {
17        a = find_set(a);
18        b = find_set(b);
19        if (a != b) {
20            if (size[a] < size[b]) swap(a, b);
21            parent[b] = a;
22            size[a] += size[b];
23        }
24    }
25 };

```

1.2 Minimum Queue

```

1 ll get_minimum(stack<pair<ll, ll>> &s1, stack<pair<
2     ll, ll>> &s2) {
3     if (s1.empty() || s2.empty()) {
4         return s1.empty() ? s2.top().second : s1.top().
5         second;
6     } else {
7         return min(s1.top().second, s2.top().second);
8     }
9 }
10 void add_element(ll new_element, stack<pair<ll, ll
11     >> &s1) {
12     ll minimum = s1.empty() ? new_element : min(
13         new_element, s1.top().second);
14     s1.push({new_element, minimum});
15 }
16 ll remove_element(stack<pair<ll, ll>> &s1, stack<
17     pair<ll, ll>> &s2) {
18     if (s2.empty()) {
19         while (!s1.empty()) {
20             ll element = s1.top().first;
21             s1.pop();
22             ll minimum = s2.empty() ? element : min(
23                 element, s2.top().second);
24             s2.push({element, minimum});
25         }
26     }
27     ll removed_element = s2.top().first;
28     s2.pop();
29     return removed_element;
30 }

```

1.3 Range Add Point Query

```

1 template<typename T, typename InType = T>
2 class SegTreeNode {
3 public:
4     const T IDN = 0, DEF = 0;
5     int i, j;
6     T val;
7     SegTreeNode<T, InType>* lc, * rc;
8     SegTreeNode(int i, int j) : i(i), j(j) {
9         if (j - i == 1) {
10             lc = rc = nullptr;
11             val = DEF;
12             return;
13         }
14         int k = (i + j) / 2;
15         lc = new SegTreeNode<T, InType>(i, k);
16         rc = new SegTreeNode<T, InType>(k, j);
17         val = 0;
18     }
19     SegTreeNode(const vector<InType>& a, int i, int j
20         ) : i(i), j(j) {
21         if (j - i == 1) {
22             lc = rc = nullptr;
23             val = (T) a[i];
24             return;
25         }
26         int k = (i + j) / 2;
27         lc = new SegTreeNode<T, InType>(a, i, k);
28         rc = new SegTreeNode<T, InType>(a, k, j);
29         val = 0;
30     }
31     void range_add(int l, int r, T x) {
32         if (r <= i || j <= l) return;
33         if (l <= i && j <= r) {
34             val += x;
35             return;
36         }
37         lc->range_add(l, r, x);
38         rc->range_add(l, r, x);
39     }
40     T point_query(int k) {
41         if (k < i || j <= k) return IDN;
42         if (j - i == 1) return val;
43         return val + lc->point_query(k) + rc->
44             point_query(k);
45     }
46 };
47 template<typename T, typename InType = T>
48 class SegTree {
49 public:
50     SegTreeNode<T, InType> root;
51     SegTree(int n) : root(0, n) {}
52     SegTree(const vector<InType>& a) : root(a, 0, a.
53         size()) {}
54     void range_add(int l, int r, T x) { root.
55         range_add(l, r, x); }
56     T point_query(int k) { return root.point_query(k)
57         ; }
58 };

```

1.4 Range Add Range Query

```

1 template<typename T, typename InType = T>
2 class SegTreeNode {
3 public:

```

```

4  const T IDN = 0, DEF = 0;
5  int i, j;
6  T val, to_add = 0;
7  SegTreeNode<T, InType>* lc, * rc;
8  SegTreeNode(int i, int j) : i(i), j(j) {
9      if (j - i == 1) {
10         lc = rc = nullptr;
11         val = DEF;
12         return;
13     }
14     int k = (i + j) / 2;
15     lc = new SegTreeNode<T, InType>(i, k);
16     rc = new SegTreeNode<T, InType>(k, j);
17     val = operation(lc->val, rc->val);
18 }
19 SegTreeNode(const vector<InType>& a, int i, int j
20             ) : i(i), j(j) {
21     if (j - i == 1) {
22         lc = rc = nullptr;
23         val = (T) a[i];
24         return;
25     }
26     int k = (i + j) / 2;
27     lc = new SegTreeNode<T, InType>(a, i, k);
28     rc = new SegTreeNode<T, InType>(a, k, j);
29     val = operation(lc->val, rc->val);
30 }
31 void propagate() {
32     if (to_add == 0) return;
33     val += to_add;
34     if (j - i > 1) {
35         lc->to_add += to_add;
36         rc->to_add += to_add;
37     }
38     to_add = 0;
39 }
40 void range_add(int l, int r, T delta) {
41     propagate();
42     if (r <= i || j <= l) return;
43     if (l <= i && j <= r) {
44         to_add += delta;
45         propagate();
46     } else {
47         lc->range_add(l, r, delta);
48         rc->range_add(l, r, delta);
49         val = operation(lc->val, rc->val);
50     }
51 }
52 T range_query(int l, int r) {
53     propagate();
54     if (l <= i && j <= r) return val;
55     if (j <= l || r <= i) return IDN;
56     return operation(lc->range_query(l, r), rc->
57                     range_query(l, r));
58 }
59 T operation(T x, T y) {}
60 template<typename T, typename InType = T>
61 class SegTree {
62 public:
63     SegTreeNode<T, InType> root;
64     SegTree(int n) : root(0, n) {}
65     SegTree(const vector<InType>& a) : root(a, 0, a.
66         size()) {}
67     void range_add(int l, int r, T delta) { root.
68         range_add(l, r, delta); }
69     T range_query(int l, int r) { return root.
70         range_query(l, r); }
71 };

```

1.5 Segment Tree

```

1  template<typename T, typename InType = T>
2  class SegTreeNode {
3 public:
4     const T IDN = 0, DEF = 0;
5     int i, j;
6     T val;
7     SegTreeNode<T, InType>* lc, * rc;
8     SegTreeNode(int i, int j) : i(i), j(j) {
9         if (j - i == 1) {
10             lc = rc = nullptr;
11             val = DEF;
12             return;
13         }
14         int k = (i + j) / 2;
15         lc = new SegTreeNode<T, InType>(i, k);
16         rc = new SegTreeNode<T, InType>(k, j);
17         val = op(lc->val, rc->val);
18     }
19     SegTreeNode(const vector<InType>& a, int i, int j
20                 ) : i(i), j(j) {
21         if (j - i == 1) {
22             lc = rc = nullptr;
23             val = (T) a[i];
24             return;
25         }
26         int k = (i + j) / 2;
27         lc = new SegTreeNode<T, InType>(a, i, k);
28         rc = new SegTreeNode<T, InType>(a, k, j);
29         val = op(lc->val, rc->val);
30     }
31     void set(int k, T x) {
32         if (k < i || j <= k) return;
33         if (j - i == 1) {
34             val = x;
35             return;
36         }
37         lc->set(k, x);
38         rc->set(k, x);
39         val = op(lc->val, rc->val);
40     }
41     T range_query(int l, int r) {
42         if (l <= i && j <= r) return val;
43         if (j <= l || r <= i) return IDN;
44         return op(lc->range_query(l, r), rc->
45                 range_query(l, r));
46     }
47     T op(T x, T y) {}
48 };
49 template<typename T, typename InType = T>
50 class SegTree {
51 public:
52     SegTreeNode<T, InType> root;
53     SegTree(int n) : root(0, n) {}
54     SegTree(const vector<InType>& a) : root(a, 0, a.
55         size()) {}
56     void set(int k, T x) { root.set(k, x); }
57     T range_query(int l, int r) { return root.
58         range_query(l, r); }
59 };

```

1.6 Sparse Table

```

1  ll log2_floor(ll i) {

```

```

2  return i ? __builtin_clzll(1) - __builtin_clzll(i
3  ) : -1;
4  }
5  vector<vector<ll>> build_sum(ll N, ll K, vector<ll>
6  &array) {
7  vector<vector<ll>> st(K + 1, vector<ll>(N + 1));
8  for (ll i = 0; i < N; i++) st[0][i] = array[i];
9  for (ll i = 1; i <= K; i++)
10     for (ll j = 0; j + (1 << i) <= N; j++)
11         st[i][j] = st[i - 1][j] + st[i - 1][j + (1 <<
12             (i - 1))];
13     return st;
14 }
15 ll sum_query(ll L, ll R, ll K, vector<vector<ll>> &
16 st) {
17     ll sum = 0;
18     for (ll i = K; i >= 0; i--) {
19         if ((1 << i) <= R - L + 1) {
20             sum += st[i][L];
21             L += 1 << i;
22         }
23     }
24     return sum;
25 }
26 vector<vector<ll>> build_min(ll N, ll K, vector<ll>
27 &array) {
28 vector<vector<ll>> st(K + 1, vector<ll>(N + 1));
29 for (ll i = 0; i < N; i++) st[0][i] = array[i];
30 for (ll i = 1; i <= K; i++)
31     for (ll j = 0; j + (1 << i) <= N; j++)
32         st[i][j] = min(st[i - 1][j], st[i - 1][j + (1
33             << (i - 1))]);
34     return st;
35 }
36 ll min_query(ll L, ll R, vector<vector<ll>> &st) {
37     ll i = log2_floor(R - L + 1);
38     return min(st[i][L], st[i][R - (1 << i) + 1]);
39 }

```

2 Dynamic Programming

2.1 Divide And Conquer

```

1  ll m, n;
2  vector<ll> dp_before(n), dp_cur(n);
3  ll C(ll i, ll j);
4  void compute(ll l, ll r, ll optl, ll opttr) {
5      if (l > r) {
6          return;
7      }
8      ll mid = (l + r) >> 1;
9      pair<ll, ll> best = {LLONG_MAX, -1};
10     for (ll k = optl; k <= min(mid, opttr); k++) {
11         best = min(best, {(k ? dp_before[k - 1] : 0) +
12             C(k, mid), k});
13     }
14     dp_cur[mid] = best.first;
15     ll opt = best.second;
16     compute(l, mid - 1, optl, opt);
17     compute(mid + 1, r, opt, opttr);
18 }
19 ll solve() {
20     for (ll i = 0; i < n; i++) {
21         dp_before[i] = C(0, i);
22     }
23     for (ll i = 1; i < m; i++) {

```

```

23     compute(0, n - 1, 0, n - 1);
24     dp_before = dp_cur;
25 }
26 return dp_before[n - 1];
27 }

```

2.2 Edit Distance

```

1 ll edit_distance(string x, string y, ll n, ll m) {
2     vector<vector<int>> dp(n + 1, vector<int>(m + 1,
3         INF));
4     dp[0][0] = 0;
5     for (int i = 1; i <= n; i++) {
6         dp[i][0] = i;
7     }
8     for (int j = 1; j <= m; j++) {
9         dp[0][j] = j;
10    }
11    for (int i = 1; i <= n; i++) {
12        for (int j = 1; j <= m; j++) {
13            dp[i][j] = min({dp[i - 1][j] + 1, dp[i][j -
14                1] + 1, dp[i - 1][j - 1] + (x[i - 1] !=
15                    y[j - 1])});
16        }
17    }
18    return dp[n][m];
19 }

```

2.3 Knapsack

```

1 ll knapsack(ll W, vector<ll> &wt, vector<ll> &val,
2     ll n) {
3     vector<ll> dp(W + 1, 0);
4     for (ll i = 1; i <= n; i++) {
5         for (ll w = W; w >= 0; w--) {
6             if (wt[i - 1] <= w) {
7                 dp[w] = max(dp[w], dp[w - wt[i - 1]] + val[
8                     i - 1]);
9             }
10        }
11    }
12    return dp[W];
13 }

```

2.4 Knuth Optimization

```

1 ll solve() {
2     ll N;
3     // read N and input
4     vector<vector<ll>> dp(N, vector<ll>(N)), opt(N,
5         vector<ll>(N));
6     auto C = [&](ll i, ll j) {
7         // Implement cost function C.
8     };
9     for (ll i = 0; i < N; i++) {
10        opt[i][i] = i;
11        ... // Initialize dp[i][i] according to the
12            problem
13    }
14    for (ll i = N - 2; i >= 0; i--) {
15        for (ll j = i + 1; j < N; j++) {
16            ll mn = ll_MAX, cost = C(i, j);

```

```

15     for (ll k = opt[i][j - 1]; k <= min(j - 1,
16         opt[i + 1][j]); k++) {
17         if (mn >= dp[i][k] + dp[k + 1][j] + cost) {
18             opt[i][j] = k;
19             mn = dp[i][k] + dp[k + 1][j] + cost;
20         }
21     }
22     dp[i][j] = mn;
23 }
24 cout << dp[0][N - 1] << '\n';
25 }

```

2.5 Longest Common Subsequence

```

1 ll LCS(string x, string y, ll n, ll m) {
2     vector<vector<ll>> dp(n + 1, vector<ll>(m + 1));
3     for (ll i = 0; i <= n; i++) {
4         for (ll j = 0; j <= m; j++) {
5             if (i == 0 || j == 0) {
6                 dp[i][j] = 0;
7             } else if (x[i - 1] == y[j - 1]) {
8                 dp[i][j] = dp[i - 1][j - 1] + 1;
9             } else {
10                dp[i][j] = max(dp[i - 1][j], dp[i][j - 1]);
11            }
12        }
13    }
14    ll index = dp[n][m];
15    vector<char> lcs(index + 1);
16    lcs[index] = '\0';
17    ll i = n, j = m;
18    while (i > 0 && j > 0) {
19        if (x[i - 1] == y[j - 1]) {
20            lcs[index - 1] = x[i - 1];
21            i--;
22            j--;
23            index--;
24        } else if (dp[i - 1][j] > dp[i][j - 1]) {
25            i--;
26        } else {
27            j--;
28        }
29    }
30    return dp[n][m];
31 }

```

2.6 Longest Increasing Subsequence

```

1 ll get_ceil_idx(vector<ll> &a, vector<ll> &T, ll l,
2     ll r, ll x) {
3     while (r - l > 1) {
4         ll m = l + (r - l) / 2;
5         if (a[T[m]] >= x) {
6             r = m;
7         } else {
8             l = m;
9         }
10    }
11    return r;
12 }
13 ll LIS(ll n, vector<ll> &a) {
14     ll len = 1;
15     vector<ll> T(n, 0), R(n, -1);
16     T[0] = 0;

```

```

16     for (ll i = 1; i < n; i++) {
17         if (a[i] < a[T[0]]) {
18             T[0] = i;
19         } else if (a[i] > a[T[len - 1]]) {
20             R[i] = T[len - 1];
21             T[len++] = i;
22         } else {
23             ll pos = get_ceil_idx(a, T, -1, len - 1, a[i
24                 ]);
25             R[i] = T[pos - 1];
26             T[pos] = i;
27         }
28     }
29     return len;
30 }

```

2.7 Subset Sum

```

1 bool subset_sum(ll n, vector<ll> &arr, ll sum) {
2     vector<vector<ll>> dp(n + 1, vector<ll>(sum + 1,
3         false));
4     dp[0][0] = true;
5     for (ll i = 1; i <= n; i++) {
6         for (ll j = 0; j <= sum; j++) {
7             dp[i][j] = dp[i - 1][j];
8             if (j >= arr[i]) {
9                 dp[i][j] |= dp[i - 1][j - arr[i]];
10            }
11        }
12    }
13    return dp[n][sum];
14 }

```

3 Geometry

3.1 Circle Line Intersection

```

1 double r, a, b, c; // given as input
2 double x0 = -a * c / (a * a + b * b);
3 double y0 = -b * c / (a * a + b * b);
4 if (c * c > r * r * (a * a + b * b) + EPS) {
5     puts ("no points");
6 } else if (abs (c * c - r * r * (a * a + b * b)) <
7     EPS) {
8     puts ("1 point");
9     cout << x0 << ' ' << y0 << '\n';
10 } else {
11     double d = r * r - c * c / (a * a + b * b);
12     double mult = sqrt (d / (a * a + b * b));
13     double ax, ay, bx, by;
14     ax = x0 + b * mult;
15     bx = x0 - b * mult;
16     ay = y0 + a * mult;
17     by = y0 - a * mult;
18     puts ("2 points");
19     cout << ax << ' ' << ay << '\n' << bx << ' ' <<
20         by << '\n';
21 }

```

3.2 Convex Hull

```

1 struct pt {
2     double x, y;
3 };
4 ll orientation(pt a, pt b, pt c) {
5     double v = a.x * (b.y - c.y) + b.x * (c.y - a.y)
6         + c.x * (a.y - b.y);
7     if (v < 0) {
8         return -1;
9     } else if (v > 0) {
10        return +1;
11    }
12    return 0;
13 }
14 bool cw(pt a, pt b, pt c, bool include_collinear) {
15     ll o = orientation(a, b, c);
16     return o < 0 || (include_collinear && o == 0);
17 }
18 bool collinear(pt a, pt b, pt c) {
19     return orientation(a, b, c) == 0;
20 }
21 void convex_hull(vector<pt>& a, bool
22     include_collinear = false) {
23     pt p0 = *min_element(a.begin(), a.end(), [](pt a,
24         pt b) {
25         return make_pair(a.y, a.x) < make_pair(b.y, b.x)
26             });
27     sort(a.begin(), a.end(), [&p0](const pt& a, const
28         pt& b) {
29         ll o = orientation(p0, a, b);
30         if (o == 0) {
31             return (p0.x - a.x) * (p0.x - a.x) + (p0.y -
32                 a.y) * (p0.y - a.y)
33                 < (p0.x - b.x) * (p0.x - b.x) + (p0.y -
34                     b.y) * (p0.y - b.y);
35         }
36         return o < 0;
37     });
38     if (include_collinear) {
39         ll i = (ll) a.size() - 1;
40         while (i >= 0 && collinear(p0, a[i], a.back()))
41             i--;
42         reverse(a.begin() + i + 1, a.end());
43     }
44     vector<pt> st;
45     for (ll i = 0; i < (ll) a.size(); i++) {
46         while (st.size() > 1 && !cw(st[st.size() - 2],
47             st.back(), a[i], include_collinear)) {
48             st.pop_back();
49         }
50         st.push_back(a[i]);
51     }
52     a = st;
53 }

```

3.3 Line Sweep

```

1 const double EPS = 1E-9;
2 struct pt {
3     double x, y;
4 };
5 struct seg {
6     pt p, q;
7     ll id;
8     double get_y(double x) const {
9         if (abs(p.x - q.x) < EPS) {
10            return p.y;

```

```

11        }
12        return p.y + (q.y - p.y) * (x - p.x) / (q.x - p
13            .x);
14    }
15 }
16 bool intersectld(double l1, double r1, double l2,
17     double r2) {
18     if (l1 > r1) {
19         swap(l1, r1);
20     }
21     if (l2 > r2) {
22         swap(l2, r2);
23     }
24     return max(l1, l2) <= min(r1, r2) + EPS;
25 }
26 ll vec(const pt& a, const pt& b, const pt& c) {
27     double s = (b.x - a.x) * (c.y - a.y) - (b.y - a.y
28         ) * (c.x - a.x);
29     return abs(s) < EPS ? 0 : s > 0 ? +1 : -1;
30 }
31 bool intersect(const seg& a, const seg& b) {
32     return intersectld(a.p.x, a.q.x, b.p.x, b.q.x) &&
33         intersectld(a.p.y, a.q.y, b.p.y, b.q.y) &&
34         vec(a.p, a.q, b.p) * vec(a.p, a.q, b.q) <=
35             0 &&
36             vec(b.p, b.q, a.p) * vec(b.p, b.q, a.q) <=
37                 0;
38 }
39 bool operator<(const seg& a, const seg& b) {
40     double x = max(min(a.p.x, a.q.x), min(b.p.x, b.
41         q.x));
42     return a.get_y(x) < b.get_y(x) - EPS;
43 }
44 struct event {
45     double x;
46     ll tp, id;
47     event() {}
48     event(double x, ll tp, ll id) : x(x), tp(tp), id(id) {}
49 }
50 bool operator<(const event& e) const {
51     if (abs(x - e.x) > EPS) {
52         return x < e.x;
53     }
54     return tp > e.tp;
55 }
56 }
57 pair<ll, ll> solve(const vector<seg>& a) {
58     ll n = (ll) a.size();
59     vector<event> e;
60     for (ll i = 0; i < n; ++i) {
61         e.push_back(event(min(a[i].p.x, a[i].q.x), +1,
62             i));
63         e.push_back(event(max(a[i].p.x, a[i].q.x), -1,
64             i));
65     }
66     sort(e.begin(), e.end());
67     s.clear();
68     where.resize(a.size());
69     for (size_t i = 0; i < e.size(); ++i) {
70         ll id = e[i].id;
71         if (e[i].tp == +1) {

```

```

72         set<seg>::iterator nxt = s.lower_bound(a[id])
73             , prv = prev(nxt);
74         if (nxt != s.end() && intersect(*nxt, a[id])) {
75             return make_pair(nxt->id, id);
76         }
77         if (prv != s.end() && intersect(*prv, a[id])) {
78             return make_pair(prv->id, id);
79         }
80         where[id] = s.insert(nxt, a[id]);
81     } else {
82         set<seg>::iterator nxt = next(where[id]), prv
83             = prev(where[id]);
84         if (nxt != s.end() && prv != s.end() &&
85             intersect(*nxt, *prv)) {
86             return make_pair(prv->id, nxt->id);
87         }
88         s.erase(where[id]);
89     }
90     return make_pair(-1, -1);
91 }

```

3.4 Nearest Points

```

1 struct pt {
2     ll x, y, id;
3 };
4 struct cmp_x {
5     bool operator()(const pt & a, const pt & b) const
6     {
7         return a.x < b.x || (a.x == b.x && a.y < b.y);
8     }
9 };
10 struct cmp_y {
11     bool operator()(const pt & a, const pt & b) const
12     {
13         return a.y < b.y;
14     }
15 };
16 ll n;
17 vector<pt> a;
18 double mindist;
19 pair<ll, ll> best_pair;
20 void upd_ans(const pt & a, const pt & b) {
21     double dist = sqrt((a.x - b.x) * (a.x - b.x) + (a
22         .y - b.y) * (a.y - b.y));
23     if (dist < mindist) {
24         mindist = dist;
25         best_pair = {a.id, b.id};
26     }
27 }
28 vector<pt> t;
29 void rec(ll l, ll r) {
30     if (r - l <= 3) {
31         for (ll i = l; i < r; ++i) {
32             for (ll j = i + 1; j < r; ++j) {
33                 upd_ans(a[i], a[j]);
34             }
35         }
36     }
37     sort(a.begin() + l, a.begin() + r, cmp_y());
38     return;
39 }
40 ll m = (1 + r) >> 1, midx = a[m].x;
41 rec(l, m);
42 rec(m, r);

```

```

39 merge(a.begin() + 1, a.begin() + m, a.begin() + m
    , a.begin() + r, t.begin(), cmp_y());
40 copy(t.begin(), t.begin() + r - 1, a.begin() + 1)
    ;
41 ll tsz = 0;
42 for (ll i = 1; i < r; ++i) {
43     if (abs(a[i].x - midx) < mindist) {
44         for (ll j = tsz - 1; j >= 0 && a[i].y - t[j].
            y < mindist; --j) {
45             upd_ans(a[i], t[j]);
46         }
47         t[tsz++] = a[i];
48     }
49 }
50 }
51 t.resize(n);
52 sort(a.begin(), a.end(), cmp_x());
53 mindist = 1E20;
54 rec(0, n);

```

4 Graph Theory

4.1 Articulation Point

```

1 void APUtil(vector<vector<ll>> &adj, ll u, vector<
    bool> &visited,
2 vector<ll> &disc, vector<ll> &low, ll &time, ll
    parent, vector<bool> &isAP) {
3     ll children = 0;
4     visited[u] = true;
5     disc[u] = low[u] = ++time;
6     for (auto v : adj[u]) {
7         if (!visited[v]) {
8             children++;
9             APUtil(adj, v, visited, disc, low, time, u,
                isAP);
10            low[u] = min(low[u], low[v]);
11            if (parent != -1 && low[v] >= disc[u]) {
12                isAP[u] = true;
13            }
14            } else if (v != parent) {
15                low[u] = min(low[u], disc[v]);
16            }
17        }
18        if (parent == -1 && children > 1) {
19            isAP[u] = true;
20        }
21    }
22    void AP(vector<vector<ll>> &adj, ll n) {
23        vector<ll> disc(n), low(n);
24        vector<bool> visited(n), isAP(n);
25        ll time = 0, par = -1;
26        for (ll u = 0; u < n; u++) {
27            if (!visited[u]) {
28                APUtil(adj, u, visited, disc, low, time, par,
                    isAP);
29            }
30        }
31        for (ll u = 0; u < n; u++) {
32            if (isAP[u]) {
33                cout << u << " ";
34            }
35        }
36    }

```

4.2 Bellman Ford

```

1 void bellman_ford(vector<vector<ll>> &edges, ll n,
    ll m, ll src, vector<ll> &dis) {
2     for (ll i = 0; i < n; i++) {
3         dis[i] = INF;
4     }
5     for (ll i = 0; i < n - 1; i++) {
6         for (ll j = 0; j < m; j++) {
7             ll u = edges[j][0], v = edges[j][1], w =
                edges[j][2];
8             if (dis[u] < INF) {
9                 dis[v] = min(dis[v], dis[u] + w);
10            }
11        }
12    }
13    for (ll i = 0; i < m; i++) {
14        ll u = edges[i][0], v = edges[i][1], w = edges[
            i][2];
15        if (dis[u] < INF && dis[u] + w < dis[v]) {
16            cout << "The graph contains a negative cycle.
                " << '\n';
17        }
18    }
19 }

```

4.3 Bridge

```

1 void bridge_util(vector<vector<ll>> &adj, ll u,
    vector<bool> &visited, vector<ll> &disc,
    vector<ll> &low, vector<ll> &parent) {
2     static ll time = 0;
3     visited[u] = true;
4     disc[u] = low[u] = ++time;
5     list<ll>::iterator i;
6     for (auto v : adj[u]) {
7         if (!visited[v]) {
8             parent[v] = u;
9             bridge_util(adj, v, visited, disc, low,
                parent);
10            low[u] = min(low[u], low[v]);
11            if (low[v] > disc[u]) {
12                cout << u << ' ' << v << '\n';
13            }
14        } else if (v != parent[u]) {
15            low[u] = min(low[u], disc[v]);
16        }
17    }
18 }
19 void bridge(vector<vector<ll>> &adj, ll n) {
20     vector<bool> visited(n, false);
21     vector<ll> disc(n), low(n), parent(n, -1);
22     for (ll i = 0; i < n; i++) {
23         if (!visited[i]) {
24             bridge_util(adj, i, visited, disc, low,
                parent);
25         }
26     }
27 }

```

4.4 Dijkstra

```

1 void dijkstra(ll n, vector<vector<pair<ll, ll>>> &
    adj, vector<ll> &dis) {

```

```

2     priority_queue<pair<ll, ll>, vector<pair<ll, ll
    >>, greater<pair<ll, ll>>> pq;
3     for (int i = 0; i < n; i++) {
4         dis[i] = INF;
5     }
6     dis[0] = 0;
7     pq.push({0, 0});
8     while (!pq.empty()) {
9         auto p = pq.top();
10        pq.pop();
11        ll u = p.second;
12        if (dis[u] != p.first) {
13            continue;
14        }
15        for (auto x : adj[u]) {
16            ll v = x.first, w = x.second;
17            if (dis[v] > dis[u] + w) {
18                dis[v] = dis[u] + w;
19                pq.push({dis[v], v});
20            }
21        }
22    }
23 }

```

4.5 Find Cycle

```

1 bool dfs(ll v) {
2     color[v] = 1;
3     for (ll u : adj[v]) {
4         if (color[u] == 0) {
5             parent[u] = v;
6             if (dfs(u)) {
7                 return true;
8             }
9         } else if (color[u] == 1) {
10            cycle_end = v;
11            cycle_start = u;
12            return true;
13        }
14    }
15    color[v] = 2;
16    return false;
17 }
18 void find_cycle() {
19     color.assign(n, 0);
20     parent.assign(n, -1);
21     cycle_start = -1;
22     for (ll v = 0; v < n; v++) {
23         if (color[v] == 0 && dfs(v)) {
24             break;
25         }
26     }
27     if (cycle_start == -1) {
28         cout << "Acyclic" << endl;
29     } else {
30         vector<ll> cycle;
31         cycle.push_back(cycle_start);
32         for (ll v = cycle_end; v != cycle_start; v =
            parent[v]) {
33             cycle.push_back(v);
34         }
35         cycle.push_back(cycle_start);
36         reverse(cycle.begin(), cycle.end());
37         cout << "Cycle found: ";
38         for (ll v : cycle) {
39             cout << v << ' ';
40         }

```

```

41     cout << '\n';
42 }
43 }

```

4.6 Floyd Warshall

```

1 void floyd_warshall(vector<vector<ll>> &dis, ll n)
2 {
3     for (ll i = 0; i < n; i++) {
4         for (ll j = 0; j < n; j++) {
5             dis[i][j] = (i == j ? 0 : INF);
6         }
7     }
8     for (ll k = 0; k < n; k++) {
9         for (ll i = 0; i < n; i++) {
10            for (ll j = 0; j < n; j++) {
11                if (dis[i][k] < INF && dis[k][j] < INF) {
12                    dis[i][j] = min(dis[i][j], dis[i][k] +
13                                   dis[k][j]);
14                }
15            }
16        }
17        for (ll i = 0; i < n; i++) {
18            for (ll j = 0; j < n; j++) {
19                for (ll k = 0; k < n; k++) {
20                    if (dis[k][k] < 0 && dis[i][k] < INF && dis
21                        [k][j] < INF) {
22                        dis[i][j] = -INF;
23                    }
24                }
25            }
26        }
27    }
28 }

```

4.7 Hierholzer

```

1 void print_circuit(vector<vector<ll>> &adj) {
2     map<ll, ll> edge_count;
3     for (ll i = 0; i < adj.size(); i++) {
4         edge_count[i] = adj[i].size();
5     }
6     if (!adj.size()) {
7         return;
8     }
9     stack<ll> curr_path;
10    vector<ll> circuit;
11    curr_path.push(0);
12    ll curr_v = 0;
13    while (!curr_path.empty()) {
14        if (edge_count[curr_v] {
15            curr_path.push(curr_v);
16            ll next_v = adj[curr_v].back();
17            edge_count[curr_v]--;
18            adj[curr_v].pop_back();
19            curr_v = next_v;
20        } else {
21            circuit.push_back(curr_v);
22            curr_v = curr_path.top();
23            curr_path.pop();
24        }
25    }
26    for (ll i = circuit.size() - 1; i >= 0; i--) {
27        cout << circuit[i] << ' ';
28    }
29 }

```

```

29 }

```

4.8 Is Bipartite

```

1 bool is_bipartite(vector<ll> &col, vector<vector<ll>
2 >> &adj, ll n) {
3     queue<pair<ll, ll>> q;
4     for (ll i = 0; i < n; i++) {
5         if (col[i] == -1) {
6             q.push({i, 0});
7             col[i] = 0;
8             while (!q.empty()) {
9                 pair<ll, ll> p = q.front();
10                q.pop();
11                ll v = p.first, c = p.second;
12                for (ll j : adj[v]) {
13                    if (col[j] == c) {
14                        return false;
15                    }
16                    if (col[j] == -1) {
17                        col[j] = (c ? 0 : 1);
18                        q.push({j, col[j]});
19                    }
20                }
21            }
22        }
23    }
24    return true;
25 }

```

4.9 Is Cyclic

```

1 bool is_cyclic_util(int u, vector<vector<int>> &adj
2 , vector<bool> &vis, vector<bool> &rec) {
3     vis[u] = true;
4     rec[u] = true;
5     for (auto v : adj[u]) {
6         if (!vis[v] && is_cyclic_util(v, adj, vis, rec)
7             ) {
8             return true;
9         }
10        else if (rec[v]) {
11            return true;
12        }
13    }
14    rec[u] = false;
15    return false;
16 }
17 bool is_cyclic(int n, vector<vector<int>> &adj) {
18     vector<bool> vis(n, false), rec(n, false);
19     for (int i = 0; i < n; i++) {
20         if (!vis[i] && is_cyclic_util(i, adj, vis, rec)
21             ) {
22             return true;
23         }
24     }
25    return false;
26 }

```

4.10 Kahn

```

1 void kahn(vector<vector<ll>> &adj) {
2     ll n = adj.size();

```

```

3     vector<ll> in_degree(n, 0);
4     for (ll u = 0; u < n; u++) {
5         for (ll v : adj[u]) {
6             in_degree[v]++;
7         }
8     }
9     queue<ll> q;
10    for (ll i = 0; i < n; i++) {
11        if (in_degree[i] == 0) {
12            q.push(i);
13        }
14    }
15    ll cnt = 0;
16    vector<ll> top_order;
17    while (!q.empty()) {
18        ll u = q.front();
19        q.pop();
20        top_order.push_back(u);
21        for (ll v : adj[u]) {
22            if (--in_degree[v] == 0) {
23                q.push(v);
24            }
25        }
26        cnt++;
27    }
28    if (cnt != n) {
29        cout << -1 << '\n';
30        return;
31    }
32    for (ll i = 0; i < (ll) top_order.size(); i++) {
33        cout << top_order[i] << ' ';
34    }
35    cout << '\n';
36 }

```

4.11 Kosaraju

```

1 void topo_sort(int u, vector<vector<int>>& adj,
2 vector<bool>& vis, stack<int>& stk) {
3     vis[u] = true;
4     for (int v : adj[u]) {
5         if (!vis[v]) {
6             topo_sort(v, adj, vis, stk);
7         }
8     }
9     stk.push(u);
10 }
11 vector<vector<int>> transpose(int n, vector<vector<
12 int>>& adj) {
13     vector<vector<int>> adj_t(n);
14     for (int u = 0; u < n; u++) {
15         for (int v : adj[u]) {
16             adj_t[v].push_back(u);
17         }
18     }
19    return adj_t;
20 }
21 void get_scc(int u, vector<vector<int>>& adj_t,
22 vector<bool>& vis, vector<int>& scc) {
23     vis[u] = true;
24     scc.push_back(u);
25     for (int v : adj_t[u]) {
26         if (!vis[v]) {
27             get_scc(v, adj_t, vis, scc);
28         }
29     }
30 }

```

```

28     }
29 }
30
31 void kosaraju(int n, vector<vector<int>>& adj,
32             vector<vector<int>>& sccs) {
33     vector<bool> vis(n, false);
34     stack<int> stk;
35     for (int u = 0; u < n; u++) {
36         if (!vis[u]) {
37             topo_sort(u, adj, vis, stk);
38         }
39     }
40     vector<vector<int>> adj_t = transpose(n, adj);
41     for (int u = 0; u < n; u++) {
42         vis[u] = false;
43     }
44     while (!stk.empty()) {
45         int u = stk.top();
46         stk.pop();
47         if (!vis[u]) {
48             vector<int> scc;
49             get_scc(u, adj_t, vis, scc);
50             sccs.push_back(scc);
51         }
52     }
}

```

4.12 Kruskal Mst

```

1 struct Edge {
2     ll u, v, weight;
3     bool operator<(Edge const& other) {
4         return weight < other.weight;
5     }
6 };
7 ll n;
8 vector<Edge> edges;
9 ll cost = 0;
10 vector<ll> tree_id(n);
11 vector<Edge> result;
12 for (ll i = 0; i < n; i++) {
13     tree_id[i] = i;
14 }
15 sort(edges.begin(), edges.end());
16 for (Edge e : edges) {
17     if (tree_id[e.u] != tree_id[e.v]) {
18         cost += e.weight;
19         result.push_back(e);
20         ll old_id = tree_id[e.u], new_id = tree_id[e.v];
21         for (ll i = 0; i < n; i++) {
22             if (tree_id[i] == old_id) {
23                 tree_id[i] = new_id;
24             }
25         }
26     }
27 }

```

4.13 Lowest Common Ancestor

```

1 struct LCA {
2     vector<ll> height, euler, first, segtree;
3     vector<bool> visited;
4     ll n;
5     LCA(vector<vector<ll>> &adj, ll root = 0) {

```

```

6         n = adj.size();
7         height.resize(n);
8         first.resize(n);
9         euler.reserve(n * 2);
10        visited.assign(n, false);
11        dfs(adj, root);
12        ll m = euler.size();
13        segtree.resize(m * 4);
14        build(1, 0, m - 1);
15    }
16    void dfs(vector<vector<ll>> &adj, ll node, ll h = 0) {
17        visited[node] = true;
18        height[node] = h;
19        first[node] = euler.size();
20        euler.push_back(node);
21        for (auto to : adj[node]) {
22            if (!visited[to]) {
23                dfs(adj, to, h + 1);
24                euler.push_back(node);
25            }
26        }
27    }
28    void build(ll node, ll b, ll e) {
29        if (b == e) {
30            segtree[node] = euler[b];
31        } else {
32            ll mid = (b + e) / 2;
33            build(node << 1, b, mid);
34            build(node << 1 | 1, mid + 1, e);
35            ll l = segtree[node << 1], r = segtree[node << 1 | 1];
36            segtree[node] = (height[l] < height[r]) ? l : r;
37        }
38    }
39    ll query(ll node, ll b, ll e, ll L, ll R) {
40        if (b > R || e < L) {
41            return -1;
42        }
43        if (b >= L && e <= R) {
44            return segtree[node];
45        }
46        ll mid = (b + e) >> 1;
47        ll left = query(node << 1, b, mid, L, R);
48        ll right = query(node << 1 | 1, mid + 1, e, L, R);
49        if (left == -1) return right;
50        if (right == -1) return left;
51        return height[left] < height[right] ? left : right;
52    }
53    ll lca(ll u, ll v) {
54        ll left = first[u], right = first[v];
55        if (left > right) {
56            swap(left, right);
57        }
58        return query(1, 0, euler.size() - 1, left, right);
59    }
60 };

```

4.14 Maximum Bipartite Matching

```

1 bool bpm(ll n, ll m, vector<vector<bool>> &bpGraph,
2         ll u, vector<bool> &seen, vector<ll> &matchR) {

```

```

2     for (ll v = 0; v < m; v++) {
3         if (bpGraph[u][v] && !seen[v]) {
4             seen[v] = true;
5             if (matchR[v] < 0 || bpm(n, m, bpGraph,
6                 matchR[v], seen, matchR)) {
7                 matchR[v] = u;
8                 return true;
9             }
10        }
11        return false;
12    }
13    ll maxBPM(ll n, ll m, vector<vector<bool>> &bpGraph) {
14        vector<ll> matchR(m, -1);
15        ll result = 0;
16        for (ll u = 0; u < n; u++) {
17            vector<bool> seen(m, false);
18            if (bpm(n, m, bpGraph, u, seen, matchR)) {
19                result++;
20            }
21        }
22        return result;
23    }

```

4.15 Max Flow

```

1 bool bfs(ll n, vector<vector<ll>> &r_graph, ll s,
2         ll t, vector<ll> &parent) {
3     vector<bool> visited(n, false);
4     queue<ll> q;
5     q.push(s);
6     visited[s] = true;
7     parent[s] = -1;
8     while (!q.empty()) {
9         ll u = q.front();
10        q.pop();
11        for (ll v = 0; v < n; v++) {
12            if (!visited[v] && r_graph[u][v] > 0) {
13                if (v == t) {
14                    return true;
15                }
16                q.push(v);
17                parent[v] = u;
18                visited[v] = true;
19            }
20        }
21    }
22    return false;
23 }
24 ll fordFulkerson(ll n, vector<vector<ll>> graph, ll s, ll t) {
25     ll u, v;
26     vector<vector<ll>> r_graph;
27     for (u = 0; u < n; u++) {
28         for (v = 0; v < n; v++) {
29             r_graph[u][v] = graph[u][v];
30         }
31     }
32     vector<ll> parent;
33     ll max_flow = 0;
34     while (bfs(n, r_graph, s, t, parent)) {
35         ll path_flow = INF;
36         for (v = t; v != s; v = parent[v]) {
37             u = parent[v];
38             path_flow = min(path_flow, r_graph[u][v]);

```



```

39     }
40     for (v = t; v != s; v = parent[v]) {
41         u = parent[v];
42         r_graph[u][v] -= path_flow;
43         r_graph[v][u] += path_flow;
44     }
45     max_flow += path_flow;
46 }
47 return max_flow;
48 }

```

4.16 Prim Mst

```

1  vector<ll> prim_mst(ll n, vector<vector<pair<ll, ll>>> &adj) {
2      priority_queue<pair<ll, ll>, vector<pair<ll, ll>>, greater<pair<ll, ll>>> pq;
3      ll src = 0;
4      vector<ll> key(n, INF), parent(n, -1);
5      vector<bool> in_mst(n, false);
6      pq.push(make_pair(0, src));
7      key[src] = 0;
8      while (!pq.empty()) {
9          ll u = pq.top().second;
10         pq.pop();
11         if (in_mst[u]) continue;
12         in_mst[u] = true;
13         for (auto p : adj[u]) {
14             ll v = p.first, w = p.second;
15             if (in_mst[v] == false && w < key[v]) {
16                 key[v] = w;
17                 pq.push(make_pair(key[v], v));
18                 parent[v] = u;
19             }
20         }
21     }
22     return parent;
23 }
24
25 }

```

4.17 Topological Sort

```

1  void dfs(ll v) {
2      visited[v] = true;
3      for (ll u : adj[v]) {
4          if (!visited[u]) {
5              dfs(u);
6          }
7      }
8      ans.push_back(v);
9  }
10 void topological_sort() {
11     visited.assign(n, false);
12     ans.clear();
13     for (ll i = 0; i < n; ++i) {
14         if (!visited[i]) {
15             dfs(i);
16         }
17     }
18     reverse(ans.begin(), ans.end());
19 }

```

5 Miscellaneous

5.1 Gauss

```

1  const double EPS = 1e-9;
2  const ll INF = 2;
3  ll gauss(vector<vector<double>> a, vector<double> &ans) {
4      ll n = (ll) a.size(), m = (ll) a[0].size() - 1;
5      vector<ll> where(m, -1);
6      for (ll col = 0, row = 0; col < m && row < n; ++col) {
7          ll sel = row;
8          for (ll i = row; i < n; ++i) {
9              if (abs(a[i][col]) > abs(a[sel][col])) {
10                 sel = i;
11             }
12             if (abs(a[sel][col]) < EPS) {
13                 continue;
14             }
15             for (ll i = col; i <= m; ++i) {
16                 swap(a[sel][i], a[row][i]);
17             }
18             where[col] = row;
19             for (ll i = 0; i < n; ++i) {
20                 if (i != row) {
21                     double c = a[i][col] / a[row][col];
22                     for (ll j = col; j <= m; ++j) {
23                         a[i][j] -= a[row][j] * c;
24                     }
25                 }
26             }
27             ++row;
28         }
29         ans.assign(m, 0);
30         for (ll i = 0; i < m; ++i) {
31             if (where[i] != -1) {
32                 ans[i] = a[where[i]][m] / a[where[i]][i];
33             }
34         }
35         for (ll i = 0; i < n; ++i) {
36             double sum = 0;
37             for (ll j = 0; j < m; ++j) {
38                 sum += ans[j] * a[i][j];
39             }
40             if (abs(sum - a[i][m]) > EPS) {
41                 return 0;
42             }
43         }
44         for (ll i = 0; i < m; ++i) {
45             if (where[i] == -1) {
46                 return INF;
47             }
48         }
49         return 1;
50     }
51 }

```

5.2 Ternary Search

```

1  double ternary_search(double l, double r) {
2      double eps = 1e-9;
3      while (r - l > eps) {
4          double m1 = l + (r - l) / 3;
5          double m2 = r - (r - l) / 3;

```

```

6          double f1 = f(m1);
7          double f2 = f(m2);
8          if (f1 < f2) {
9              l = m1;
10             } else {
11                 r = m2;
12             }
13     }
14     return f(l);
15 }

```

6 Number Theory

6.1 Extended Euclidean

```

1  ll gcd_extended(ll a, ll b, ll &x, ll &y) {
2      if (b == 0) {
3          x = 1;
4          y = 0;
5          return a;
6      }
7      ll x1, y1, g = gcd_extended(b, a % b, x1, y1);
8      x = y1;
9      y = x1 - (a / b) * y1;
10     return g;
11 }

```

6.2 Find All Solutions

```

1  bool find_any_solution(ll a, ll b, ll c, ll &x0, ll &y0, ll &g) {
2      g = gcd_extended(abs(a), abs(b), x0, y0);
3      if (c % g) {
4          return false;
5      }
6      x0 *= c / g;
7      y0 *= c / g;
8      if (a < 0) {
9          x0 = -x0;
10     }
11     if (b < 0) {
12         y0 = -y0;
13     }
14     return true;
15 }
16 void shift_solution(ll &x, ll &y, ll a, ll b, ll cnt) {
17     x += cnt * b;
18     y -= cnt * a;
19 }
20 ll find_all_solutions(ll a, ll b, ll c, ll minx, ll maxx, ll miny, ll maxy) {
21     ll x, y, g;
22     if (!find_any_solution(a, b, c, x, y, g)) {
23         return 0;
24     }
25     a /= g;
26     b /= g;
27     ll sign_a = a > 0 ? +1 : -1;
28     ll sign_b = b > 0 ? +1 : -1;
29     shift_solution(x, y, a, b, (minx - x) / b);
30     if (x < minx) {
31         shift_solution(x, y, a, b, sign_b);
32     }

```



```

33 if (x > maxx) {
34     return 0;
35 }
36 ll lx1 = x;
37 shift_solution(x, y, a, b, (maxx - x) / b);
38 if (x > maxx) {
39     shift_solution(x, y, a, b, -sign_b);
40 }
41 ll rx1 = x;
42 shift_solution(x, y, a, b, -(miny - y) / a);
43 if (y < miny) {
44     shift_solution(x, y, a, b, -sign_a);
45 }
46 if (y > maxy) {
47     return 0;
48 }
49 ll lx2 = x;
50 shift_solution(x, y, a, b, -(maxy - y) / a);
51 if (y > maxy) {
52     shift_solution(x, y, a, b, sign_a);
53 }
54 ll rx2 = x;
55 if (lx2 > rx2) {
56     swap(lx2, rx2);
57 }
58 ll lx = max(lx1, lx2), rx = min(rx1, rx2);
59 if (lx > rx) {
60     return 0;
61 }
62 return (rx - lx) / abs(b) + 1;
63 }

```

6.3 Linear Sieve

```

1 void linear_sieve(ll N, vector<ll> &lowest_prime,
2 vector<ll> &prime) {
3     for (ll i = 2; i <= N; i++) {
4         if (lowest_prime[i] == 0) {
5             lowest_prime[i] = i;
6             prime.push_back(i);
7         }
8         for (ll j = 0; i * prime[j] <= N; j++) {
9             lowest_prime[i * prime[j]] = prime[j];
10            if (prime[j] == lowest_prime[i]) {
11                break;
12            }
13        }
14    }

```

6.4 Miller Rabin

```

1 bool check_composite(u64 n, u64 a, u64 d, ll s) {
2     u64 x = binpower(a, d, n);
3     if (x == 1 || x == n - 1) {
4         return false;
5     }
6     for (ll r = 1; r < s; r++) {
7         x = (u128) x * x % n;
8         if (x == n - 1) {
9             return false;
10        }
11    }
12    return true;
13 }

```

```

14 bool miller_rabin(u64 n) {
15     if (n < 2) {
16         return false;
17     }
18     ll r = 0;
19     u64 d = n - 1;
20     while ((d & 1) == 0) {
21         d >>= 1;
22         r++;
23     }
24     for (ll a : {2, 3, 5, 7, 11, 13, 17, 19, 23, 29,
25                31, 37}) {
26         if (n == a) {
27             return true;
28         }
29         if (check_composite(n, a, d, r)) {
30             return false;
31         }
32     }
33     return true;

```

6.5 Modulo Inverse

```

1 ll mod_inv(ll a, ll m) {
2     if (m == 1) {
3         return 0;
4     }
5     ll m0 = m, x = 1, y = 0;
6     while (a > 1) {
7         ll q = a / m, t = m;
8         m = a % m;
9         a = t;
10        t = y;
11        y = x - q * y;
12        x = t;
13    }
14    if (x < 0) {
15        x += m0;
16    }
17    return x;
18 }

```

6.6 Pollard Rho Brent

```

1 ll mult(ll a, ll b, ll mod) {
2     return (__int128_t) a * b % mod;
3 }
4 ll f(ll x, ll c, ll mod) {
5     return (mult(x, x, mod) + c) % mod;
6 }
7 ll pollard_rho_brent(ll n, ll x0 = 2, ll c = 1) {
8     ll x = x0, g = 1, q = 1, xs, y, m = 128, l = 1;
9     while (g == 1) {
10        y = x;
11        for (ll i = 1; i < l; i++) {
12            x = f(x, c, n);
13        }
14        ll k = 0;
15        while (k < l && g == 1) {
16            xs = x;
17            for (ll i = 0; i < m && i < l - k; i++) {
18                x = f(x, c, n);
19                q = mult(q, abs(y - x), n);
20            }

```

```

21        g = __gcd(q, n);
22        k += m;
23    }
24    l *= 2;
25 }
26 if (g == n) {
27     do {
28         xs = f(xs, c, n);
29         g = __gcd(abs(xs - y), n);
30     } while (g == 1);
31 }
32 return g;
33 }

```

6.7 Range Sieve

```

1 vector<bool> range_sieve(ll l, ll r) {
2     ll n = sqrt(r);
3     vector<bool> is_prime(n + 1, true);
4     vector<ll> prime;
5     is_prime[0] = is_prime[1] = false;
6     prime.push_back(2);
7     for (ll i = 4; i <= n; i += 2) {
8         is_prime[i] = false;
9     }
10    for (ll i = 3; i <= n; i += 2) {
11        if (is_prime[i]) {
12            prime.push_back(i);
13            for (ll j = i * i; j <= n; j += i) {
14                is_prime[j] = false;
15            }
16        }
17    }
18    vector<bool> result(r - l + 1, true);
19    for (ll i : prime) {
20        for (ll j = max(i * i, (l + i - 1) / i * i); j
21              <= r; j += i) {
22            result[j - l] = false;
23        }
24    }
25    if (l == 1) {
26        result[0] = false;
27    }
28    return result;

```

6.8 Segmented Sieve

```

1 vector<ll> segmented_sieve(ll n) {
2     const ll S = 10000;
3     ll nsqrt = sqrt(n);
4     vector<char> is_prime(nsqrt + 1, true);
5     vector<ll> prime;
6     is_prime[0] = is_prime[1] = false;
7     prime.push_back(2);
8     for (ll i = 4; i <= nsqrt; i += 2) {
9         is_prime[i] = false;
10    }
11    for (ll i = 3; i <= nsqrt; i += 2) {
12        if (is_prime[i]) {
13            prime.push_back(i);
14            for (ll j = i * i; j <= nsqrt; j += i) {
15                is_prime[j] = false;
16            }
17        }

```

```

18 }
19 vector<ll> result;
20 vector<char> block(S);
21 for (ll k = 0; k * S <= n; k++) {
22     fill(block.begin(), block.end(), true);
23     for (ll p : prime) {
24         for (ll j = max((k * S + p - 1) / p, p) * p -
25                 k * S; j < S; j += p) {
26             block[j] = false;
27         }
28     }
29     if (k == 0) {
30         block[0] = block[1] = false;
31     }
32     for (ll i = 0; i < S && k * S + i <= n; i++) {
33         if (block[i]) {
34             result.push_back(k * S + i);
35         }
36     }
37     return result;
38 }

```

6.9 Tonelli Shanks

```

1 ll legendre(ll a, ll p) {
2     return bin_pow_mod(a, (p - 1) / 2, p);
3 }
4 ll tonelli_shanks(ll n, ll p) {
5     if (legendre(n, p) == p - 1) {
6         return -1;
7     }
8     if (p % 4 == 3) {
9         return bin_pow_mod(n, (p + 1) / 4, p);
10    }
11    ll Q = p - 1, S = 0;
12    while (Q % 2 == 0) {
13        Q /= 2;
14        S++;
15    }
16    ll z = 2;
17    for (; z < p; z++) {
18        if (legendre(z, p) == p - 1) {
19            break;
20        }
21    }
22    ll M = S, c = bin_pow_mod(z, Q, p), t =
23        bin_pow_mod(n, Q, p), R = bin_pow_mod(n, (Q
24            + 1) / 2, p);
25    while (t % p != 1) {
26        if (t % p == 0) {
27            return 0;
28        }
29        ll i = 1, t2 = t * t % p;
30        for (; i < M; i++) {
31            if (t2 % p == 1) {
32                break;
33            }
34            t2 = t2 * t2 % p;
35        }
36        ll b = bin_pow_mod(c, bin_pow_mod(2, M - i - 1,
37            p), p);
38        M = i;
39        c = b * b % p;
40        t = t * c % p;
41        R = R * b % p;
42    }
43 }

```

```

40     return R;
41 }

```

7 Strings

7.1 Hashing

```

1 ll compute_hash(string const& s) {
2     const ll p = 31, m = 1e9 + 9;
3     ll hash_value = 0, p_pow = 1;
4     for (char c : s) {
5         hash_value = (hash_value + (c - 'a' + 1) *
6             p_pow) % m;
7         p_pow = (p_pow * p) % m;
8     }
9     return hash_value;
10 }

```

7.2 Knuth Morris Pratt

```

1 vector<ll> prefix_function(string s) {
2     ll n = (ll) s.length();
3     vector<ll> pi(n);
4     for (ll i = 1; i < n; i++) {
5         ll j = pi[i - 1];
6         while (j > 0 && s[i] != s[j]) {
7             j = pi[j - 1];
8         }
9         if (s[i] == s[j]) {
10            j++;
11        }
12        pi[i] = j;
13    }
14    return pi;
15 }

```

7.3 Rabin Karp

```

1 vector<ll> rabin_karp(string const& s, string const
2     & t) {
3     const ll p = 31, m = 1e9 + 9;
4     ll S = s.size(), T = t.size();
5     vector<ll> p_pow(max(S, T));
6     p_pow[0] = 1;
7     for (ll i = 1; i < (ll) p_pow.size(); i++) {
8         p_pow[i] = (p_pow[i - 1] * p) % m;
9     }
10    vector<ll> h(T + 1, 0);
11    for (ll i = 0; i < T; i++) {
12        h[i + 1] = (h[i] + (t[i] - 'a' + 1) * p_pow[i])
13            % m;
14    }
15    ll h_s = 0;
16    for (ll i = 0; i < S; i++) {
17        h_s = (h_s + (s[i] - 'a' + 1) * p_pow[i]) % m;
18    }
19    vector<ll> occurrences;
20    for (ll i = 0; i + S - 1 < T; i++) {
21        ll cur_h = (h[i + S] + m - h[i]) % m;
22        if (cur_h == h_s * p_pow[i] % m) {
23            occurrences.push_back(i);
24        }
25    }
26 }

```

```

23 }
24 return occurrences;
25 }

```

7.4 Suffix Array

```

1 vector<ll> sort_cyclic_shifts(string const& s) {
2     ll n = s.size();
3     const ll alphabet = 256;
4     vector<ll> p(n), c(n), cnt(max(alphabet, n), 0);
5     for (ll i = 0; i < n; i++) {
6         cnt[s[i]]++;
7     }
8     for (ll i = 1; i < alphabet; i++) {
9         cnt[i] += cnt[i - 1];
10    }
11    for (ll i = 0; i < n; i++) {
12        p[--cnt[s[i]]] = i;
13    }
14    c[p[0]] = 0;
15    ll classes = 1;
16    for (ll i = 1; i < n; i++) {
17        if (s[p[i]] != s[p[i - 1]]) {
18            classes++;
19        }
20        c[p[i]] = classes - 1;
21    }
22    vector<ll> pn(n), cn(n);
23    for (ll h = 0; (1 << h) < n; ++h) {
24        for (ll i = 0; i < n; i++) {
25            pn[i] = p[i] - (1 << h);
26            if (pn[i] < 0) {
27                pn[i] += n;
28            }
29        }
30        fill(cnt.begin(), cnt.begin() + classes, 0);
31        for (ll i = 0; i < n; i++) {
32            cnt[c[pn[i]]]++;
33        }
34        for (ll i = 1; i < classes; i++) {
35            cnt[i] += cnt[i - 1];
36        }
37        for (ll i = n - 1; i >= 0; i--) {
38            p[--cnt[c[pn[i]]]] = pn[i];
39        }
40        cn[p[0]] = 0;
41        classes = 1;
42        for (ll i = 1; i < n; i++) {
43            pair<ll, ll> cur = {c[p[i]], c[(p[i] + (1 <<
44                h)) % n]};
45            pair<ll, ll> prev = {c[p[i - 1]], c[(p[i - 1]
46                + (1 << h)) % n]};
47            if (cur != prev) {
48                ++classes;
49            }
50            cn[p[i]] = classes - 1;
51        }
52        c.swap(cn);
53    }
54    vector<ll> build_suff_arr(string s) {
55        s += (char) 0;
56        vector<ll> sorted_shifts = sort_cyclic_shifts(s);
57        sorted_shifts.erase(sorted_shifts.begin());
58        return sorted_shifts;
59    }

```

7.5 Z Function

```
1 vector<ll> z_function(string s) {  
2     ll n = (ll) s.length();  
3     vector<ll> z(n);  
4     for (ll i = 1, l = 0, r = 0; i < n; ++i) {
```

```
5         if (i <= r) {  
6             z[i] = min (r - i + 1, z[i - l]);  
7         }  
8         while (i + z[i] < n && s[z[i]] == s[i + z[i]])  
9             ++z[i];  
10    }  
11    if (i + z[i] - 1 > r) {
```

```
12        l = i, r = i + z[i] - 1;  
13    }  
14 }  
15 return z;  
16 }
```

| | | |
|--|--|---|
| $f(n) = O(g(n))$ | iff \exists positive c, n_0 such that $0 \leq f(n) \leq cg(n) \forall n \geq n_0$. | $\sum_{i=1}^n i = \frac{n(n+1)}{2}, \quad \sum_{i=1}^n i^2 = \frac{n(n+1)(2n+1)}{6}, \quad \sum_{i=1}^n i^3 = \frac{n^2(n+1)^2}{4}.$ |
| $f(n) = \Omega(g(n))$ | iff \exists positive c, n_0 such that $f(n) \geq cg(n) \geq 0 \forall n \geq n_0$. | In general: |
| $f(n) = \Theta(g(n))$ | iff $f(n) = O(g(n))$ and $f(n) = \Omega(g(n))$. | $\sum_{i=1}^n i^m = \frac{1}{m+1} \left[(n+1)^{m+1} - 1 - \sum_{i=1}^n ((i+1)^{m+1} - i^{m+1} - (m+1)i^m) \right]$ |
| $f(n) = o(g(n))$ | iff $\lim_{n \rightarrow \infty} f(n)/g(n) = 0$. | $\sum_{i=1}^{n-1} i^m = \frac{1}{m+1} \sum_{k=0}^m \binom{m+1}{k} B_k n^{m+1-k}.$ |
| $\lim_{n \rightarrow \infty} a_n = a$ | iff $\forall \epsilon > 0, \exists n_0$ such that $ a_n - a < \epsilon, \forall n \geq n_0$. | Geometric series: |
| $\sup S$ | least $b \in \mathbb{R}$ such that $b \geq s, \forall s \in S$. | $\sum_{i=0}^n c^i = \frac{c^{n+1} - 1}{c - 1}, \quad c \neq 1, \quad \sum_{i=0}^{\infty} c^i = \frac{1}{1 - c}, \quad \sum_{i=1}^{\infty} c^i = \frac{c}{1 - c}, \quad c < 1,$ |
| $\inf S$ | greatest $b \in \mathbb{R}$ such that $b \leq s, \forall s \in S$. | $\sum_{i=0}^n ic^i = \frac{nc^{n+2} - (n+1)c^{n+1} + c}{(c-1)^2}, \quad c \neq 1, \quad \sum_{i=0}^{\infty} ic^i = \frac{c}{(1-c)^2}, \quad c < 1.$ |
| $\liminf_{n \rightarrow \infty} a_n$ | $\lim_{n \rightarrow \infty} \inf \{a_i \mid i \geq n, i \in \mathbb{N}\}.$ | Harmonic series: |
| $\limsup_{n \rightarrow \infty} a_n$ | $\lim_{n \rightarrow \infty} \sup \{a_i \mid i \geq n, i \in \mathbb{N}\}.$ | $H_n = \sum_{i=1}^n \frac{1}{i}, \quad \sum_{i=1}^n iH_i = \frac{n(n+1)}{2}H_n - \frac{n(n-1)}{4}.$ |
| $\binom{n}{k}$ | Combinations: Size k sub-sets of a size n set. | $\sum_{i=1}^n H_i = (n+1)H_n - n, \quad \sum_{i=1}^n \binom{i}{m} H_i = \binom{n+1}{m+1} \left(H_{n+1} - \frac{1}{m+1} \right).$ |
| $[n]$ | Stirling numbers (1st kind): Arrangements of an n element set into k cycles. | 1. $\binom{n}{k} = \frac{n!}{(n-k)!k!}, \quad 2. \sum_{k=0}^n \binom{n}{k} = 2^n, \quad 3. \binom{n}{k} = \binom{n}{n-k},$ |
| $\left\{ \begin{matrix} n \\ k \end{matrix} \right\}$ | Stirling numbers (2nd kind): Partitions of an n element set into k non-empty sets. | 4. $\binom{n}{k} = \frac{n}{k} \binom{n-1}{k-1}, \quad 5. \binom{n}{k} = \binom{n-1}{k} + \binom{n-1}{k-1},$ |
| $\langle \begin{matrix} n \\ k \end{matrix} \rangle$ | 1st order Eulerian numbers: Permutations $\pi_1 \pi_2 \dots \pi_n$ on $\{1, 2, \dots, n\}$ with k ascents. | 6. $\binom{n}{m} \binom{m}{k} = \binom{n}{k} \binom{n-k}{m-k}, \quad 7. \sum_{k=0}^n \binom{r+k}{k} = \binom{r+n+1}{n},$ |
| $\langle \langle \begin{matrix} n \\ k \end{matrix} \rangle \rangle$ | 2nd order Eulerian numbers. | 8. $\sum_{k=0}^n \binom{k}{m} = \binom{n+1}{m+1}, \quad 9. \sum_{k=0}^n \binom{r}{k} \binom{s}{n-k} = \binom{r+s}{n},$ |
| C_n | Catalan Numbers: Binary trees with $n+1$ vertices. | 10. $\binom{n}{k} = (-1)^k \binom{k-n-1}{k}, \quad 11. \left\{ \begin{matrix} n \\ 1 \end{matrix} \right\} = \left\{ \begin{matrix} n \\ n \end{matrix} \right\} = 1,$ |
| 14. $\begin{bmatrix} n \\ 1 \end{bmatrix} = (n-1)!,$ | 15. $\begin{bmatrix} n \\ 2 \end{bmatrix} = (n-1)!H_{n-1},$ | 16. $\begin{bmatrix} n \\ n \end{bmatrix} = 1, \quad 17. \begin{bmatrix} n \\ k \end{bmatrix} \geq \left\{ \begin{matrix} n \\ k \end{matrix} \right\},$ |
| 18. $\begin{bmatrix} n \\ k \end{bmatrix} = (n-1) \begin{bmatrix} n-1 \\ k \end{bmatrix} + \begin{bmatrix} n-1 \\ k-1 \end{bmatrix},$ | 19. $\left\{ \begin{matrix} n \\ n-1 \end{matrix} \right\} = \begin{bmatrix} n \\ n-1 \end{bmatrix} = \binom{n}{2},$ | 20. $\sum_{k=0}^n \begin{bmatrix} n \\ k \end{bmatrix} = n!, \quad 21. C_n = \frac{1}{n+1} \binom{2n}{n},$ |
| 22. $\langle \begin{matrix} n \\ 0 \end{matrix} \rangle = \langle \begin{matrix} n \\ n-1 \end{matrix} \rangle = 1,$ | 23. $\langle \begin{matrix} n \\ k \end{matrix} \rangle = \langle \begin{matrix} n \\ n-1-k \end{matrix} \rangle,$ | 24. $\langle \begin{matrix} n \\ k \end{matrix} \rangle = (k+1) \langle \begin{matrix} n-1 \\ k \end{matrix} \rangle + (n-k) \langle \begin{matrix} n-1 \\ k-1 \end{matrix} \rangle,$ |
| 25. $\langle \begin{matrix} 0 \\ k \end{matrix} \rangle = \begin{cases} 1 & \text{if } k=0, \\ 0 & \text{otherwise} \end{cases}$ | 26. $\langle \begin{matrix} n \\ 1 \end{matrix} \rangle = 2^n - n - 1,$ | 27. $\langle \begin{matrix} n \\ 2 \end{matrix} \rangle = 3^n - (n+1)2^n + \binom{n+1}{2},$ |
| 28. $x^n = \sum_{k=0}^n \langle \begin{matrix} n \\ k \end{matrix} \rangle \binom{x+k}{n},$ | 29. $\langle \begin{matrix} n \\ m \end{matrix} \rangle = \sum_{k=0}^m \binom{n+1}{k} (m+1-k)^n (-1)^k,$ | 30. $m! \left\{ \begin{matrix} n \\ m \end{matrix} \right\} = \sum_{k=0}^n \langle \begin{matrix} n \\ k \end{matrix} \rangle \binom{k}{n-m},$ |
| 31. $\langle \begin{matrix} n \\ m \end{matrix} \rangle = \sum_{k=0}^n \left\{ \begin{matrix} n \\ k \end{matrix} \right\} \binom{n-k}{m} (-1)^{n-k-m} k!,$ | 32. $\langle \langle \begin{matrix} n \\ 0 \end{matrix} \rangle \rangle = 1,$ | 33. $\langle \langle \begin{matrix} n \\ n \end{matrix} \rangle \rangle = 0 \quad \text{for } n \neq 0,$ |
| 34. $\langle \langle \begin{matrix} n \\ k \end{matrix} \rangle \rangle = (k+1) \langle \langle \begin{matrix} n-1 \\ k \end{matrix} \rangle \rangle + (2n-1-k) \langle \langle \begin{matrix} n-1 \\ k-1 \end{matrix} \rangle \rangle,$ | 35. $\sum_{k=0}^n \langle \langle \begin{matrix} n \\ k \end{matrix} \rangle \rangle = \frac{(2n)n}{2^n},$ | 36. $\left\{ \begin{matrix} x \\ x-n \end{matrix} \right\} = \sum_{k=0}^n \langle \langle \begin{matrix} n \\ k \end{matrix} \rangle \rangle \binom{x+n-1-k}{2n},$ |
| 37. $\left\{ \begin{matrix} n+1 \\ m+1 \end{matrix} \right\} = \sum_k \binom{n}{k} \left\{ \begin{matrix} k \\ m \end{matrix} \right\} = \sum_{k=0}^n \left\{ \begin{matrix} k \\ m \end{matrix} \right\} (m+1)^{n-k},$ | | |

The Chinese remainder theorem: There exists a number C such that:

$$C \equiv r_1 \pmod{m_1}$$

$$\vdots \quad \vdots \quad \vdots$$

$$C \equiv r_n \pmod{m_n}$$

if m_i and m_j are relatively prime for $i \neq j$.

Euler's function: $\phi(x)$ is the number of positive integers less than x relatively prime to x . If $\prod_{i=1}^n p_i^{e_i}$ is the prime factorization of x then

$$\phi(x) = \prod_{i=1}^n p_i^{e_i-1} (p_i - 1).$$

Euler's theorem: If a and b are relatively prime then

$$1 \equiv a^{\phi(b)} \pmod{b}.$$

Fermat's theorem:

$$1 \equiv a^{p-1} \pmod{p}.$$

The Euclidean algorithm: if $a > b$ are integers then

$$\gcd(a, b) = \gcd(a \bmod b, b).$$

If $\prod_{i=1}^n p_i^{e_i}$ is the prime factorization of x then

$$S(x) = \sum_{d|x} d = \prod_{i=1}^n \frac{p_i^{e_i+1} - 1}{p_i - 1}.$$

Perfect Numbers: x is an even perfect number iff $x = 2^{n-1}(2^n - 1)$ and $2^n - 1$ is prime.

Wilson's theorem: n is a prime iff

$$(n-1)! \equiv -1 \pmod{n}.$$

Möbius inversion:

$$\mu(i) = \begin{cases} 1 & \text{if } i = 1. \\ 0 & \text{if } i \text{ is not square-free.} \\ (-1)^r & \text{if } i \text{ is the product of } r \text{ distinct primes.} \end{cases}$$

If

$$G(a) = \sum_{d|a} F(d),$$

then

$$F(a) = \sum_{d|a} \mu(d) G\left(\frac{a}{d}\right).$$

Prime numbers:

$$p_n = n \ln n + n \ln \ln n - n + n \frac{\ln \ln n}{\ln n}$$

$$+ O\left(\frac{n}{\ln n}\right),$$

$$\pi(n) = \frac{n}{\ln n} + \frac{n}{(\ln n)^2} + \frac{2!n}{(\ln n)^3}$$

$$+ O\left(\frac{n}{(\ln n)^4}\right).$$

Definitions:

Loop An edge connecting a vertex to itself.

Directed Each edge has a direction.

Simple Graph with no loops or multi-edges.

Walk A sequence $v_0 e_1 v_1 \dots e_\ell v_\ell$.

Trail A walk with distinct edges.

Path A trail with distinct vertices.

Connected A graph where there exists a path between any two vertices.

Component A maximal connected subgraph.

Tree A connected acyclic graph.

Free tree A tree with no root.

DAG Directed acyclic graph.

Eulerian Graph with a trail visiting each edge exactly once.

Hamiltonian Graph with a cycle visiting each vertex exactly once.

Cut A set of edges whose removal increases the number of components.

Cut-set A minimal cut.

Cut edge A size 1 cut.

k-Connected A graph connected with the removal of any $k-1$ vertices.

k-Tough $\forall S \subseteq V, S \neq \emptyset$ we have $k \cdot c(G-S) \leq |S|$.

k-Regular A graph where all vertices have degree k .

k-Factor A k -regular spanning subgraph.

Matching A set of edges, no two of which are adjacent.

Clique A set of vertices, all of which are adjacent.

Ind. set A set of vertices, none of which are adjacent.

Vertex cover A set of vertices which cover all edges.

Planar graph A graph which can be embedded in the plane.

Plane graph An embedding of a planar graph.

$$\sum_{v \in V} \deg(v) = 2m.$$

If G is planar then $n - m + f = 2$, so

$$f \leq 2n - 4, \quad m \leq 3n - 6.$$

Any planar graph has a vertex with degree ≤ 5 .

Notation:

$E(G)$ Edge set

$V(G)$ Vertex set

$c(G)$ Number of components

$G[S]$ Induced subgraph

$\deg(v)$ Degree of v

$\Delta(G)$ Maximum degree

$\delta(G)$ Minimum degree

$\chi(G)$ Chromatic number

$\chi_E(G)$ Edge chromatic number

G^c Complement graph

K_n Complete graph

K_{n_1, n_2} Complete bipartite graph

$r(k, \ell)$ Ramsey number

Geometry

Projective coordinates: triples (x, y, z) , not all x, y and z zero.

$$(x, y, z) = (cx, cy, cz) \quad \forall c \neq 0.$$

Cartesian Projective

$$(x, y) \quad (x, y, 1)$$

$$y = mx + b \quad (m, -1, b)$$

$$x = c \quad (1, 0, -c)$$

Distance formula, L_p and L_∞ metric:

$$\sqrt{(x_1 - x_0)^2 + (y_1 - y_0)^2},$$

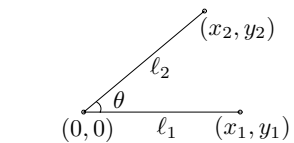
$$[|x_1 - x_0|^p + |y_1 - y_0|^p]^{1/p},$$

$$\lim_{p \rightarrow \infty} [|x_1 - x_0|^p + |y_1 - y_0|^p]^{1/p}.$$

Area of triangle $(x_0, y_0), (x_1, y_1)$ and (x_2, y_2) :

$$\frac{1}{2} \text{abs} \begin{vmatrix} x_1 - x_0 & y_1 - y_0 \\ x_2 - x_0 & y_2 - y_0 \end{vmatrix}.$$

Angle formed by three points:



$$\cos \theta = \frac{(x_1, y_1) \cdot (x_2, y_2)}{l_1 l_2}.$$

Line through two points (x_0, y_0) and (x_1, y_1) :

$$\begin{vmatrix} x & y & 1 \\ x_0 & y_0 & 1 \\ x_1 & y_1 & 1 \end{vmatrix} = 0.$$

Area of circle, volume of sphere:

$$A = \pi r^2, \quad V = \frac{4}{3} \pi r^3.$$

If I have seen farther than others, it is because I have stood on the shoulders of giants.

– Issac Newton

Taylor's series:

$$f(x) = f(a) + (x-a)f'(a) + \frac{(x-a)^2}{2}f''(a) + \dots = \sum_{i=0}^{\infty} \frac{(x-a)^i}{i!} f^{(i)}(a).$$

Expansions:

$$\begin{aligned} \frac{1}{1-x} &= 1 + x + x^2 + x^3 + x^4 + \dots = \sum_{i=0}^{\infty} x^i, \\ \frac{1}{1-cx} &= 1 + cx + c^2x^2 + c^3x^3 + \dots = \sum_{i=0}^{\infty} c^i x^i, \\ \frac{1}{1-x^n} &= 1 + x^n + x^{2n} + x^{3n} + \dots = \sum_{i=0}^{\infty} x^{ni}, \\ \frac{x}{(1-x)^2} &= x + 2x^2 + 3x^3 + 4x^4 + \dots = \sum_{i=0}^{\infty} ix^i, \\ x^k \frac{d^n}{dx^n} \left(\frac{1}{1-x} \right) &= x + 2^n x^2 + 3^n x^3 + 4^n x^4 + \dots = \sum_{i=0}^{\infty} i^n x^i, \\ e^x &= 1 + x + \frac{1}{2}x^2 + \frac{1}{6}x^3 + \dots = \sum_{i=0}^{\infty} \frac{x^i}{i!}, \\ \ln(1+x) &= x - \frac{1}{2}x^2 + \frac{1}{3}x^3 - \frac{1}{4}x^4 + \dots = \sum_{i=1}^{\infty} (-1)^{i+1} \frac{x^i}{i}, \\ \ln \frac{1}{1-x} &= x + \frac{1}{2}x^2 + \frac{1}{3}x^3 + \frac{1}{4}x^4 + \dots = \sum_{i=1}^{\infty} \frac{x^i}{i}, \\ \sin x &= x - \frac{1}{3!}x^3 + \frac{1}{5!}x^5 - \frac{1}{7!}x^7 + \dots = \sum_{i=0}^{\infty} (-1)^i \frac{x^{2i+1}}{(2i+1)!}, \\ \cos x &= 1 - \frac{1}{2!}x^2 + \frac{1}{4!}x^4 - \frac{1}{6!}x^6 + \dots = \sum_{i=0}^{\infty} (-1)^i \frac{x^{2i}}{(2i)!}, \\ \tan^{-1} x &= x - \frac{1}{3}x^3 + \frac{1}{5}x^5 - \frac{1}{7}x^7 + \dots = \sum_{i=0}^{\infty} (-1)^i \frac{x^{2i+1}}{(2i+1)}, \\ (1+x)^n &= 1 + nx + \frac{n(n-1)}{2}x^2 + \dots = \sum_{i=0}^{\infty} \binom{n}{i} x^i, \\ \frac{1}{(1-x)^{n+1}} &= 1 + (n+1)x + \binom{n+2}{2}x^2 + \dots = \sum_{i=0}^{\infty} \binom{i+n}{i} x^i, \\ \frac{x}{e^x - 1} &= 1 - \frac{1}{2}x + \frac{1}{12}x^2 - \frac{1}{720}x^4 + \dots = \sum_{i=0}^{\infty} \frac{B_i x^i}{i!}, \\ \frac{1}{2x}(1 - \sqrt{1-4x}) &= 1 + x + 2x^2 + 5x^3 + \dots = \sum_{i=0}^{\infty} \frac{1}{i+1} \binom{2i}{i} x^i, \\ \frac{1}{\sqrt{1-4x}} &= 1 + x + 2x^2 + 6x^3 + \dots = \sum_{i=0}^{\infty} \binom{2i}{i} x^i, \\ \frac{1}{\sqrt{1-4x}} \left(\frac{1 - \sqrt{1-4x}}{2x} \right)^n &= 1 + (2+n)x + \binom{4+n}{2}x^2 + \dots = \sum_{i=0}^{\infty} \binom{2i+n}{i} x^i, \\ \frac{1}{1-x} \ln \frac{1}{1-x} &= x + \frac{3}{2}x^2 + \frac{11}{6}x^3 + \frac{25}{12}x^4 + \dots = \sum_{i=1}^{\infty} H_i x^i, \\ \frac{1}{2} \left(\ln \frac{1}{1-x} \right)^2 &= \frac{1}{2}x^2 + \frac{3}{4}x^3 + \frac{11}{24}x^4 + \dots = \sum_{i=2}^{\infty} \frac{H_{i-1} x^i}{i}, \\ \frac{x}{1-x-x^2} &= x + x^2 + 2x^3 + 3x^4 + \dots = \sum_{i=0}^{\infty} F_i x^i, \\ \frac{F_n x}{1 - (F_{n-1} + F_{n+1})x - (-1)^n x^2} &= F_n x + F_{2n} x^2 + F_{3n} x^3 + \dots = \sum_{i=0}^{\infty} F_{ni} x^i. \end{aligned}$$

Ordinary power series:

$$A(x) = \sum_{i=0}^{\infty} a_i x^i.$$

Exponential power series:

$$A(x) = \sum_{i=0}^{\infty} a_i \frac{x^i}{i!}.$$

Dirichlet power series:

$$A(x) = \sum_{i=1}^{\infty} \frac{a_i}{i^x}.$$

Binomial theorem:

$$(x+y)^n = \sum_{k=0}^n \binom{n}{k} x^{n-k} y^k.$$

Difference of like powers:

$$x^n - y^n = (x-y) \sum_{k=0}^{n-1} x^{n-1-k} y^k.$$

For ordinary power series:

$$\alpha A(x) + \beta B(x) = \sum_{i=0}^{\infty} (\alpha a_i + \beta b_i) x^i,$$

$$x^k A(x) = \sum_{i=k}^{\infty} a_{i-k} x^i,$$

$$\frac{A(x) - \sum_{i=0}^{k-1} a_i x^i}{x^k} = \sum_{i=0}^{\infty} a_{i+k} x^i,$$

$$A(cx) = \sum_{i=0}^{\infty} c^i a_i x^i,$$

$$A'(x) = \sum_{i=0}^{\infty} (i+1) a_{i+1} x^i,$$

$$xA'(x) = \sum_{i=1}^{\infty} i a_i x^i,$$

$$\int A(x) dx = \sum_{i=1}^{\infty} \frac{a_{i-1}}{i} x^i,$$

$$\frac{A(x) + A(-x)}{2} = \sum_{i=0}^{\infty} a_{2i} x^{2i},$$

$$\frac{A(x) - A(-x)}{2} = \sum_{i=0}^{\infty} a_{2i+1} x^{2i+1}.$$

Summation: If $b_i = \sum_{j=0}^i a_j$ then

$$B(x) = \frac{1}{1-x} A(x).$$

Convolution:

$$A(x)B(x) = \sum_{i=0}^{\infty} \left(\sum_{j=0}^i a_j b_{i-j} \right) x^i.$$

God made the natural numbers;
all the rest is the work of man.
– Leopold Kronecker