UPLB Eliens - Pegaraw Notebook

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  1 Data Structures
1.1 Binary Trie
```

```
struct Node { struct Node* parent, child[2]; };
      struct BinaryTrie {
        Node* root;
        BinaryTrie() {
          root = new Node();
           root->parent = NULL;
          root->child[0] = NULL;
          root->child[1] = NULL;
        void insert_node(int x) {
          Node* cur = root;
          for (int place = 29; place >= 0; place--) {
            int bit = x >> place & 1;
            if (cur->child[bit] != NULL) cur = cur->child
                  [bit];
            else {
              cur->child[bit] = new Node();
              cur->child[bit]->parent = cur;
              cur = cur->child[bit];
              cur->child[0] = NULL;
              cur->child[1] = NULL;
12 24
        void remove node(int x) {
          Node* cur = root;
           for (int place = 29; place >= 0; place--) {
```

```
int bit = x >> place & 1;
             if (cur->child[bit] == NULL) return;
             cur = cur->child[bit];
13 30
           while (cur->parent != NULL && cur->child[0] ==
                NULL && cur->child[1] == NULL) {
             Node* temp = cur;
             cur = cur->parent;
14 34
             if (temp == cur->child[0]) cur->child[0] =
             else cur->child[1] = NULL;
             delete temp;
         int get_min_xor(int x) {
           Node* cur = root;
           int minXor = 0;
           for (int place = 29; place >= 0; place--) {
             int bit = x >> place & 1;
             if (cur->child[bit] != NULL) cur = cur->child
                  [bit];
             else {
               minXor ^= 1 << place;
               cur = cur->child[1 ^ bit];
           return minXor;
```

1.2 Disjoint Set Union

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```
struct DSU {
      vector<int> parent, size;
      DSU(int n) {
        parent.resize(n);
        size.resize(n);
        for (int i = 0; i < n; i++) make_set(i);</pre>
      void make_set(int v) {
        parent[v] = v;
        size[v] = 1;
      bool is_same(int a, int b) { return find_set(a)
           == find_set(b); }
      int find_set(int v) { return v == parent[v] ? v :
            parent[v] = find_set(parent[v]); }
      void union_sets(int a, int b) {
       a = find_set(a);
       b = find_set(b);
        if (a != b) {
          if (size[a] < size[b]) swap(a, b);</pre>
          parent[b] = a;
          size[a] += size[b];
23 };
```

1.3 Minimum Queue

```
1 11 get_minimum(stack<pair<11, 11>> &s1, stack<pair<</pre>
        11, 11>> &s2) {
     if (s1.empty() || s2.empty()) {
       return s1.empty() ? s2.top().second : s1.top().
            second;
```

```
} else {
        return min(s1.top().second, s2.top().second);
 7
    void add_element(ll new_element, stack<pair<ll, ll</pre>
      11 minimum = s1.empty() ? new_element : min(
           new_element, s1.top().second);
      s1.push({new_element, minimum});
11 }
12 11 remove_element(stack<pair<11, 11>> &s1, stack<
         pair<11, 11>> &s2) {
13
      if (s2.empty()) {
14
        while (!sl.empty()) {
15
          11 element = s1.top().first;
16
          s1.pop();
          11 minimum = s2.empty() ? element : min(
               element, s2.top().second);
18
          s2.push({element, minimum});
19
20
21
      11 removed_element = s2.top().first;
      s2.pop();
23
      return removed_element;
24 }
```

1.4 Mo

```
void remove(idx); // TODO: remove value at idx
         from data structure
    void add(idx);
                     // TODO: add value at idx from
         data structure
    int get_answer(); // TODO: extract the current
         answer of the data structure
    int block_size;
    struct Query {
      int 1, r, idx;
      bool operator<(Query other) const {</pre>
 8
        return make_pair(l / block_size, r) < make_pair</pre>
              (other.1 / block_size, other.r);
 9
10 };
    vector<int> mo_s_algorithm(vector<Query> queries) {
      vector<int> answers(queries.size());
13
      sort(queries.begin(), queries.end());
14
      // TODO: initialize data structure
      int cur_1 = 0, cur_r = -1;
16
      // invariant: data structure will always reflect
            the range [cur_1, cur_r]
17
      for (Query q : queries) {
18
        while (cur_1 > q.1) {
19
          cur_1--;
20
          add(cur_l);
         while (cur_r < q.r) {</pre>
23
          cur r++;
24
25
          add(cur_r);
26
27
         while (cur_1 < q.1) {
          remove(cur_l);
28
          cur_1++;
29
30
        while (cur_r > q.r) {
31
          remove(cur_r);
          cur_r--;
34
         answers[q.idx] = get_answer();
```

```
return answers;
1.5 Range Add Point Query
```

```
template<typename T, typename InType = T>
    class SegTreeNode {
      const T IDN = 0, DEF = 0;
      int i, j;
      T val;
      SegTreeNode<T, InType>* lc, * rc;
      SegTreeNode(int i, int j) : i(i), j(j) {
        if (j - i == 1) {
          lc = rc = nullptr;
          val = DEF;
          return;
        int k = (i + j) / 2;
        lc = new SegTreeNode<T, InType>(i, k);
        rc = new SegTreeNode<T, InType>(k, j);
        val = 0;
18
      SegTreeNode(const vector<InType>& a, int i, int j
          ) : i(i), j(j) {
        if (j - i == 1) {
          lc = rc = nullptr;
          val = (T) a[i];
          return;
        int k = (i + j) / 2;
        lc = new SegTreeNode<T, InType>(a, i, k);
        rc = new SegTreeNode<T, InType>(a, k, j);
        val = 0;
      void range_add(int 1, int r, T x) {
        if (r <= i || j <= 1) return;</pre>
        if (1 <= i && j <= r) {
          val += x;
          return;
        lc->range add(l, r, x);
        rc->range_add(1, r, x);
      T point_query(int k) {
40
        if (k < i \mid | j \le k) return IDN;
        if (j - i == 1) return val;
42
        return val + lc->point_query(k) + rc->
             point_query(k);
    };
    template<typename T, typename InType = T>
    class SegTree {
    public:
48
      SegTreeNode<T, InType> root;
49
      SegTree(int n) : root(0, n) {}
50
      SegTree(const vector<InType>& a) : root(a, 0, a.
           size()) {}
      void range_add(int 1, int r, T x) { root.
           range_add(1, r, x); }
52
      T point_query(int k) { return root.point_query(k)
           ; }
    };
```

1.6 Range Add Range Query

```
template<typename T, typename InType = T>
    class SegTreeNode {
3
    public:
      const T IDN = 0, DEF = 0;
      int i, j;
      T val, to_add = 0;
      SegTreeNode<T, InType>* lc, * rc;
      SeqTreeNode(int i, int j) : i(i), j(j) {
        if (j - i == 1) {
          lc = rc = nullptr;
          val = DEF:
          return;
        int k = (i + j) / 2;
        lc = new SegTreeNode<T, InType>(i, k);
        rc = new SegTreeNode<T, InType>(k, j);
        val = operation(lc->val, rc->val);
18
19
      SegTreeNode(const vector<InType>& a, int i, int j
           ) : i(i), j(j) {
        if (j - i == 1) {
          lc = rc = nullptr;
          val = (T) a[i];
          return;
        int k = (i + j) / 2;
        lc = new SegTreeNode<T, InType>(a, i, k);
        rc = new SegTreeNode<T, InType>(a, k, j);
        val = operation(lc->val, rc->val);
2.9
      void propagate() {
        if (to_add == 0) return;
        val += to_add;
        if (j - i > 1) {
34
          lc->to_add += to_add;
35
          rc->to_add += to_add;
        to\_add = 0;
38
      void range_add(int 1, int r, T delta) {
40
        propagate();
        if (r <= i | | j <= 1) return;</pre>
        if (l <= i && j <= r) {
          to_add += delta;
          propagate();
        } else {
          lc->range_add(l, r, delta);
          rc->range_add(1, r, delta);
48
          val = operation(lc->val, rc->val);
      T range_query(int 1, int r) {
        propagate();
        if (1 <= i && j <= r) return val;</pre>
        if (j <= 1 || r <= i) return IDN;</pre>
        return operation(lc->range_query(l, r), rc->
             range_query(l, r));
56
57
      T operation(T x, T y) {}
58
    template<typename T, typename InType = T>
    class SegTree {
61
62
      SegTreeNode<T, InType> root;
      SegTree(int n) : root(0, n) {}
```

1.7 Segment Tree

```
template<typename T, typename InType = T>
    class SegTreeNode {
    public:
      const T IDN = 0, DEF = 0;
      int i, j;
      T val;
      SegTreeNode<T, InType>* lc, * rc;
      SegTreeNode(int i, int j) : i(i), j(j) {
        if (j - i == 1) {
10
          lc = rc = nullptr;
11
          val = DEF;
12
          return:
13
14
        int k = (i + j) / 2;
        lc = new SegTreeNode<T, InType>(i, k);
16
         rc = new SegTreeNode<T, InType>(k, j);
17
        val = op(lc->val, rc->val);
18
19
      SegTreeNode(const vector<InType>& a, int i, int j
           ) : i(i), j(j) {
        if (j - i == 1) {
21
          lc = rc = nullptr;
22
          val = (T) a[i];
23
          return;
24
25
         int k = (i + j) / 2;
26
        lc = new SegTreeNode<T, InType>(a, i, k);
27
         rc = new SegTreeNode<T, InType>(a, k, j);
28
        val = op(lc->val, rc->val);
29
      void set(int k, T x) {
31
        if (k < i | | j <= k) return;
32
        if (j - i == 1) {
          val = x;
34
          return;
36
        lc->set(k, x);
37
        rc->set(k, x);
38
        val = op(lc->val, rc->val);
39
40
      T range_query(int 1, int r) {
41
        if (1 <= i && j <= r) return val;</pre>
42
        if (j <= 1 || r <= i) return IDN;</pre>
         return op(lc->range_query(l, r), rc->
43
             range_query(1, r));
44
45
      T \circ p(T \times, T y) \{ \}
46 };
    template<typename T, typename InType = T>
    class SegTree {
48
49
    public:
      SegTreeNode<T, InType> root;
51
      SegTree(int n) : root(0, n) {}
52
      SeqTree(const vector<InType>& a) : root(a, 0, a.
           size()) {}
      void set(int k, T x) { root.set(k, x); }
```

1.8 Segment Tree 2d

```
template<typename T, typename InType = T>
    class SegTree2dNode {
    public:
      int i, j, tree_size;
      SegTree<T, InType>* seg_tree;
      SegTree2dNode<T, InType>* lc, * rc;
      SegTree2dNode() {}
      SegTree2dNode(const vector<vector<InType>>& a,
           int i, int j) : i(i), j(j) {
        tree_size = a[0].size();
        if (j - i == 1) {
          lc = rc = nullptr;
          seg_tree = new SegTree<T, InType>(a[i]);
          return;
        int k = (i + j) / 2;
        lc = new SegTree2dNode<T, InType>(a, i, k);
        rc = new SegTree2dNode<T, InType>(a, k, j);
18
        seg_tree = new SegTree<T, InType>(vector<T>(
             tree_size));
19
        operation_2d(lc->seg_tree, rc->seg_tree);
       ~SegTree2dNode() {
        delete lc;
        delete rc:
      void set_2d(int kx, int ky, T x) {
        if (kx < i || j <= kx) return;</pre>
        if (j - i == 1) {
          seg_tree->set(ky, x);
          return;
        1c->set_2d(kx, ky, x);
        rc \rightarrow set_2d(kx, ky, x);
3.3
        operation_2d(lc->seg_tree, rc->seg_tree);
34
35
      T range_query_2d(int lx, int rx, int ly, int ry)
36
        if (lx <= i && j <= rx) return seg_tree->
             range_query(ly, ry);
37
        if (j <= lx || rx <= i) return -INF;</pre>
38
        return max(lc->range_query_2d(lx, rx, ly, ry),
             rc->range_query_2d(lx, rx, ly, ry));
39
40
      void operation_2d(SegTree<T, InType>* x, SegTree<</pre>
           T, InType>* y) {
        for (int k = 0; k < tree_size; k++) {</pre>
          seg_tree->set(k, max(x->range_query(k, k + 1)
               , y->range_query(k, k + 1)));
43
44
45
    template<typename T, typename InType = T>
    class SegTree2d {
48
    public:
      SegTree2dNode<T, InType> root;
      SegTree2d() {}
      SegTree2d(const vector<vector<InType>>& mat) :
           root(mat, 0, mat.size()) {}
      void set_2d(int kx, int ky, T x) { root.set_2d(kx
           , ky, x); }
```

```
T range_query_2d(int lx, int rx, int ly, int ry)
{ return root.range_query_2d(lx, rx, ly, ry)
; }
54 };
```

1.9 Sparse Table

```
1  11 log2_floor(ll i) {
      return i ? __builtin_clzll(1) - __builtin_clzll(i
           ): -1;
3 }
   vector<vector<ll>> build_sum(ll N, ll K, vector<ll>
          &array) {
      vector<vector<ll>>> st(K + 1, vector<ll>(N + 1));
      for (ll i = 0; i < N; i++) st[0][i] = array[i];</pre>
      for (ll i = 1; i <= K; i++)</pre>
8
        for (11 j = 0; j + (1 << i) <= N; <math>j++)
          st[i][j] = st[i - 1][j] + st[i - 1][j + (1 <<
                (i - 1))];
      return st;
    11 sum_query(11 L, 11 R, 11 K, vector<vector<11>>> &
         st) {
      11 \text{ sum} = 0;
      for (11 i = K; i >= 0; i--) {
        if ((1 << i) <= R - L + 1) {</pre>
          sum += st[i][L];
          L += 1 << i;
18
19
      return sum;
21
22
    vector<vector<ll>> build_min(ll N, ll K, vector<ll>
          &array) {
      vector<vector<ll>> st(K + 1, vector<ll>(N + 1));
      for (ll i = 0; i < N; i++) st[0][i] = array[i];</pre>
25
      for (11 i = 1; i <= K; i++)</pre>
        for (11 \ j = 0; \ j + (1 << i) <= N; \ j++)
          st[i][j] = min(st[i-1][j], st[i-1][j+(1
                 << (i - 1));
28
      return st;
2.9
    11 min_query(11 L, 11 R, vector<vector<11>>> &st) {
      ll i = log2\_floor(R - L + 1);
      return min(st[i][L], st[i][R - (1 << i) + 1]);</pre>
33 }
```

1.10 Sparse Table 2d

```
const int N = 100;
int matrix[N][N];
int table[N][N][(int)(log2(N) + 1)][(int)(log2(N) + 1)];

void build_sparse_table(int n, int m) {
   for (int i = 0; i < n; i++)
      for (int j = 0; j < m; j++)
      table[i][j][0][0] = matrix[i][j];

   for (int k = 1; k <= (int)(log2(n)); k++)
   for (int j = 0; j + (1 << k) - 1 < n; i++)
   for (int j = 0; j + (1 << k) - 1 < m; j++)
   table[i][j][k][0] = min(table[i][j][k - 1][0], table[i + (1 << (k - 1))][j][k - 1][0], table[i + (1 << (k - 1))][j][k - 1][0];

for (int k = 1; k <= (int)(log2(m)); k++)
   for (int i = 0; i < n; i++)</pre>
```

```
4
```

```
14
           for (int j = 0; j + (1 << k) - 1 < m; <math>j++)
15
             table[i][j][0][k] = min(table[i][j][0][k -
                  1], table[i][j + (1 << (k - 1))][0][k
16
       for (int k = 1; k \le (int)(log2(n)); k++)
17
         for (int 1 = 1; 1 <= (int) (log2(m)); 1++)</pre>
18
           for (int i = 0; i + (1 << k) - 1 < n; i++)
19
             for (int j = 0; j + (1 << 1) - 1 < m; <math>j++)
20
               table[i][j][k][l] = min(
21
                 min(table[i][j][k-1][l-1], table[i]
                       + (1 << (k - 1)) ] [j] [k - 1] [1 -
22
                 min(table[i][j + (1 << (l - 1))][k -
                       1][1-1], table[i + (1 << (k - 1)
                       ) ] [j + (1 << (1 - 1))] [k - 1] [1 -
23
               );
24
25 int rmq(int x1, int y1, int x2, int y2) {
26
       int k = log2(x2 - x1 + 1), l = log2(y2 - y1 + 1);
27
       return max(
28
         \max(\text{table}[x1][y1][k][1], \text{table}[x2 - (1 << k) +
              1|[v1|[k][1]),
         \max(\text{table}[x1][y2 - (1 << 1) + 1][k][1], \text{table}[
              x^2 - (1 \ll k) + 1][y^2 - (1 \ll 1) + 1][k][1
      );
31 }
```

1.11 Sqrt Decomposition

```
1 int n;
    vector<int> a (n);
    int len = (int) sqrt (n + .0) + 1; // size of the
          block and the number of blocks
    vector<int> b (len);
    for (int i = 0; i<n; ++i) b[i / len] += a[i];</pre>
    for (;;) {
       int 1, r;
       // read input data for the next query
       int sum = 0;
      for (int i = 1; i <= r; )</pre>
11
         if (i % len == 0 && i + len - 1 <= r) {</pre>
           // if the whole block starting at i belongs
                to [1, r]
           sum += b[i / len];
14
          i += len;
15
         } else {
           sum += a[i];
           ++i;
18
19
       // or
20
21
22
23
24
25
       /*
       int sum = 0;
       int c_1 = 1 / len, c_r = r / len;
       if (c_1 == c_r)
           for (int i=1; i<=r; ++i)
               sum += a[i];
26
27
           for (int i=1, end=(c_1+1)*len-1; i<=end; ++i)
28
               sum += a[i];
29
30
           for (int i=c_1+1; i<=c_r-1; ++i)
               sum += b[i];
           for (int i=c_r*len; i<=r; ++i)
               sum += a[i];
```

```
35
```

2 Dynamic Programming

2.1 Divide And Conquer

```
1 11 m, n;
   vector<ll> dp_before(n), dp_cur(n);
   11 C(11 i, 11 j);
   void compute(ll l, ll r, ll optl, ll optr) {
      if (1 > r) return;
      11 \text{ mid} = (1 + r) >> 1;
      pair<11, 11> best = {LLONG_MAX, -1};
      for (ll k = optl; k <= min(mid, optr); k++)</pre>
       best = min(best, \{(k ? dp\_before[k - 1] : 0) +
             C(k, mid), k});
      dp_cur[mid] = best.first;
      11 opt = best.second;
      compute(1, mid - 1, optl, opt);
      compute(mid + 1, r, opt, optr);
      for (11 i = 0; i < n; i++) dp_before[i] = C(0, i)</pre>
      for (11 i = 1; i < m; i++) {
18
        compute (0, n - 1, 0, n - 1);
19
        dp_before = dp_cur;
      return dp_before[n - 1];
```

2.2 Edit Distance

```
11 edit_distance(string x, string y, 11 n, 11 m) {
     vector<vector<int>> dp(n + 1, vector<int>(m + 1,
          INF));
     dp[0][0] = 0;
     for (int i = 1; i <= n; i++) {</pre>
       dp[i][0] = i;
     for (int j = 1; j \le m; j++) {
8
       dp[0][j] = j;
9
     for (int i = 1; i <= n; i++) {
       for (int j = 1; j <= m; j++) {
         dp[i][j] = min({dp[i - 1][j] + 1, dp[i][j -
              1] + 1, dp[i - 1][j - 1] + (x[i - 1] !=
              y[j - 1])));
     return dp[n][m];
```

2.3 Knapsack

2.4 Knuth Optimization

```
1 11 solve() {
     11 N;
     ... // Read input
     vector<vector<ll>> dp(N, vector<ll>(N)), opt(N,
         vector<ll>(N));
     auto C = [&](ll i, ll j) {
       ... // Implement cost function C.
     for (11 i = 0; i < N; i++) {
       opt[i][i] = i;
       ... // Initialize dp[i][i] according to the
           problem
     for (11 i = N - 2; i >= 0; i--) {
       for (11 j = i + 1; j < N; j++) {
         11 mn = 11_MAX, cost = C(i, j);
         for (11 k = opt[i][j-1]; k \le min(j-1,
              opt[i + 1][j]); k++) {
           if (mn \ge dp[i][k] + dp[k + 1][j] + cost) {
             opt[i][j] = k;
             mn = dp[i][k] + dp[k + 1][j] + cost;
         dp[i][j] = mn;
     cout << dp[0][N - 1] << '\n';
```

2.5 Longest Common Subsequence

```
1 ll LCS(string x, string y, ll n, ll m) {
     vector<vector<ll>> dp(n + 1, vector<ll>(m + 1));
     for (ll i = 0; i <= n; i++) {
       for (11 j = 0; j <= m; j++) {
         if (i == 0 || j == 0) {
            dp[i][j] = 0;
         } else if (x[i - 1] == y[j - 1]) {
            dp[i][j] = dp[i - 1][j - 1] + 1;
         } else {
           dp[i][j] = max(dp[i - 1][j], dp[i][j - 1]);
     11 \text{ index} = dp[n][m];
     vector<char> lcs(index + 1);
     lcs[index] = ' \setminus 0';
     11 i = n, j = m;
     while (i > 0 \&\& j > 0) {
       if (x[i - 1] == y[j - 1]) {
         lcs[index - 1] = x[i - 1];
         i--;
       } else if (dp[i - 1][j] > dp[i][j - 1]) {
```

2.6 Longest Increasing Subsequence

```
1  ll get_ceil_idx(vector<ll> &a, vector<ll> &T, ll l,
          11 r, 11 x) {
      while (r - 1 > 1) {
        11 m = 1 + (r - 1) / 2;
        if (a[T[m]] >= x) {
         r = m;
        } else {
          1 = m;
10
      return r;
11 }
12 11 LIS(11 n, vector<11> &a) {
13
      11 len = 1;
14
      vector<11> T(n, 0), R(n, -1);
      T[0] = 0;
      for (11 i = 1; i < n; i++) {
17
        if (a[i] < a[T[0]]) {</pre>
18
         T[0] = i;
19
        } else if (a[i] > a[T[len - 1]]) {
20
          R[i] = T[len - 1];
21
          T[len++] = i;
22
        } else {
23
         11 pos = get_ceil_idx(a, T, -1, len - 1, a[i
              1);
          R[i] = T[pos - 1];
25
          T[pos] = i;
26
27
28
      return len;
29 }
```

2.7 Max Sum

```
int max_subarray_sum(vi arr) {
  int x = 0, s = 0;
  for (int k = 0; k < n; k++) {
    s = max(arr[k], s+arr[k]);
    x = max(x,s);
  }
  return x;
}</pre>
```

2.8 Subset Sum

```
7     if (j >= arr[i]) {
8         dp[i][j] |= dp[i - 1][j - arr[i]];
9      }
10     }
11     }
12     return dp[n][sum];
13  }
```

3 Geometry

3.1 Areas

```
int signed_area_parallelogram(point2d p1, point2d
         p2, point2d p3) {
      return cross(p2 - p1, p3 - p2);
 3
   double triangle_area(point2d p1, point2d p2,
         point2d p3) {
      return abs(signed_area_parallelogram(p1, p2, p3))
            / 2.0;
   bool clockwise (point2d p1, point2d p2, point2d p3)
      return signed_area_parallelogram(p1, p2, p3) < 0;</pre>
9
   bool counter_clockwise(point2d p1, point2d p2,
         point2d p3) {
      return signed_area_parallelogram(p1, p2, p3) > 0;
12
   double area(const vector<point>& fig) {
      double res = 0;
      for (unsigned i = 0; i < fig.size(); i++) {
        point p = i ? fig[i - 1] : fig.back();
       point q = fig[i];
18
       res += (p.x - q.x) * (p.y + q.y);
19
20
      return fabs(res) / 2;
```

3.2 Basic Geometry

return *this;

```
1 struct point2d {
      ftype x, y;
      point2d() {}
     point2d(ftype x, ftype y): x(x), y(y) {}
      point2d& operator+=(const point2d &t) {
6
       x += t.x;
       y += t.y;
8
       return *this:
9
     point2d& operator-=(const point2d &t) {
       x -= t.x;
       y -= t.y;
       return *this;
14
      point2d& operator *= (ftype t) {
       x *= t;
       y *= t;
       return *this;
      point2d& operator/=(ftype t) {
       x /= t;
        y /= t;
```

```
point2d operator+(const point2d &t) const {
           return point2d(*this) += t; }
26
      point2d operator-(const point2d &t) const {
           return point2d(*this) -= t; }
      point2d operator*(ftype t) const { return point2d
           (*this) *= t; }
28
      point2d operator/(ftype t) const { return point2d
           (*this) /= t; }
30 point2d operator*(ftype a, point2d b) { return b *
31 ftype dot(point2d a, point2d b) { return a.x * b.x
         + a.y * b.y; }
32 ftype dot(point3d a, point3d b) { return a.x * b.x
         + a.y * b.y + a.z * b.z; }
33 ftype norm(point2d a) { return dot(a, a); }
    double abs(point2d a) { return sqrt(norm(a)); }
    double proj(point2d a, point2d b) { return dot(a, b
        ) / abs(b); }
    double angle(point2d a, point2d b) { return acos(
        dot(a, b) / abs(a) / abs(b)); }
   point3d cross(point3d a, point3d b) { return
        point3d(a.y * b.z - a.z * b.y, a.z * b.x - a.x
         * b.z, a.x * b.y - a.y * b.x);
38 ftype triple(point3d a, point3d b, point3d c) {
        return dot(a, cross(b, c)); }
    ftype cross(point2d a, point2d b) { return a.x * b.
        y - a.y * b.x; }
   point2d intersect(point2d a1, point2d d1, point2d
         a2, point2d d2) { return a1 + cross(a2 - a1,
         d2) / cross(d1, d2) * d1; }
   point3d intersect(point3d a1, point3d n1, point3d
         a2, point3d n2, point3d a3, point3d n3) {
      point3d x(n1.x, n2.x, n3.x);
      point3d y(n1.y, n2.y, n3.y);
44
      point3d z(n1.z, n2.z, n3.z);
      point3d d(dot(a1, n1), dot(a2, n2), dot(a3, n3));
46
      return point3d(triple(d, y, z), triple(x, d, z),
          triple(x, y, d)) / triple(n1, n2, n3);
47 }
```

3.3 Circle Line Intersection

```
double r, a, b, c; // given as input
    double x0 = -a * c / (a * a + b * b);
    double y0 = -b * c / (a * a + b * b);
   if (c * c > r * r * (a * a + b * b) + EPS) {
     puts ("no points");
   } else if (abs (c *c - r * r * (a * a + b * b)) <</pre>
         EPS) {
      puts ("1 point");
      cout << x0 << ' ' << y0 << '\n';
      double d = r * r - c * c / (a * a + b * b);
      double mult = sqrt (d / (a * a + b * b));
      double ax, ay, bx, by;
     ax = x0 + b * mult;
     bx = x0 - b * mult;
     ay = y0 - a * mult;
     bv = v0 + a * mult;
      puts ("2 points");
      cout << ax << ' ' << ay << '\n' << bx << ' ' <<
          by << '\n';
19
```

```
struct pt {
      double x, y;
 3 };
    11 orientation(pt a, pt b, pt c) {
      double v = a.x * (b.y - c.y) + b.x * (c.y - a.y)
           + c.x * (a.y - b.y);
      if (v < 0) {
        return -1;
      \} else if (v > 0) {
        return +1;
11
      return 0:
12
13 bool cw(pt a, pt b, pt c, bool include_collinear) {
14
      11 o = orientation(a, b, c);
15
      return o < 0 || (include collinear && o == 0);</pre>
16
17
    bool collinear(pt a, pt b, pt c) {
18
      return orientation(a, b, c) == 0;
19
20
    void convex_hull(vector<pt>& a, bool
         include collinear = false) {
      pt p0 = *min_element(a.begin(), a.end(), [](pt a,
        return make_pair(a.y, a.x) < make_pair(b.y, b.x</pre>
             );
24
      sort(a.begin(), a.end(), [&p0](const pt& a, const
            pt& b) {
25
        11 \circ = orientation(p0, a, b);
26
        if (o == 0) {
          return (p0.x - a.x) * (p0.x - a.x) + (p0.y - a.x)
               a.y) * (p0.y - a.y)
                < (p0.x - b.x) * (p0.x - b.x) + (p0.y -
                    b.y) * (p0.y - b.y);
29
        return o < 0;
      }):
      if (include collinear) {
        11 i = (11) a.size()-1;
34
        while (i \ge 0 \&\& collinear(p0, a[i], a.back()))
              i--;
        reverse(a.begin()+i+1, a.end());
36
      vector<pt> st;
38
      for (ll i = 0; i < (ll) a.size(); i++) {</pre>
39
        while (st.size() > 1 && !cw(st[st.size() - 2],
             st.back(), a[i], include_collinear)) {
40
          st.pop_back();
41
42
        st.push back(a[i]);
43
44
      a = st:
45 }
```

3.5 Count Lattices

3.4 Convex Hull

```
cnt += (fk * (n - 1) + 2 * fb) * n / 2;
        k -= fk;
 8
        b -= fb;
9
      auto t = k * n + b;
      auto ft = t.floor();
      if (ft \geq= 1) cnt += count_lattices(1 / k, (t - t.
           floor()) / k, t.floor());
      return cnt;
14 }
 3.6 Line Intersection
    struct pt { double x, y; };
    struct line { double a, b, c; };
    const double EPS = 1e-9;
    double det(double a, double b, double c, double d)
         { return a*d - b*c; }
    bool intersect(line m, line n, pt & res) {
      double zn = det(m.a, m.b, n.a, n.b);
      if (abs(zn) < EPS) return false;</pre>
      res.x = -det(m.c, m.b, n.c, n.b) / zn;
      res.y = -det(m.a, m.c, n.a, n.c) / zn;
      return true;
    bool parallel(line m, line n) { return abs(det(m.a,
          m.b, n.a, n.b)) < EPS; }
    bool equivalent(line m, line n) {
      return abs(det(m.a, m.b, n.a, n.b)) < EPS</pre>
15
          && abs(det(m.a, m.c, n.a, n.c)) < EPS
6
          && abs(det(m.b, m.c, n.b, n.c)) < EPS;
1.7 }
 3.7 Line Sweep
    const double EPS = 1E-9;
    struct pt { double x, y; };
    struct seq {
      pt p, q;
      11 id:
      double get v(double x) const {
        if (abs(p.x - q.x) < EPS) return p.y;</pre>
 8
        return p.y + (q.y - p.y) * (x - p.x) / (q.x - p.x)
             .x):
 9
    };
    bool intersect1d(double 11, double r1, double 12,
         double r2) {
      if (11 > r1) swap(11, r1);
      if (12 > r2) swap(12, r2);
14
      return max(11, 12) <= min(r1, r2) + EPS;</pre>
   11 vec (const pt& a, const pt& b, const pt& c) {
```

double s = (b.x - a.x) * (c.y - a.y) - (b.y - a.y)

return intersectld(a.p.x, a.q.x, b.p.x, b.q.x) &&
 intersectld(a.p.y, a.q.y, b.p.y, b.q.y) &&

vec(a.p, a.q, b.p) * vec(a.p, a.q, b.q) <=</pre>

vec(b.p, b.q, a.p) * vec(b.p, b.q, a.q) <=

return abs(s) < EPS ? 0 : s > 0 ? +1 : -1;

bool intersect(const seg& a, const seg& b) {

) * (c.x - a.x);

3.3 0

18

19

2.4

25 }

```
26 bool operator<(const seg& a, const seg& b) {</pre>
      double x = max(min(a.p.x, a.g.x), min(b.p.x, b.g.
           x));
28
      return a.get_y(x) < b.get_y(x) - EPS;</pre>
29
30 struct event {
      double v.
      11 tp, id;
      event() {}
      event (double x, 11 tp, 11 id) : x(x), tp(tp), id(
34
      bool operator<(const event& e) const {</pre>
        if (abs(x - e.x) > EPS) return x < e.x;
        return tp > e.tp;
38
39
   };
40 set<seg> s;
    vector<set<seg>::iterator> where;
    set<seg>::iterator prev(set<seg>::iterator it) {
      return it == s.begin() ? s.end() : --it;
    set<seg>::iterator next(set<seg>::iterator it) {
      return ++it;
47
48
    pair<11, 11> solve(const vector<seg>& a) {
      11 n = (11) a.size();
      vector<event> e;
      for (11 i = 0; i < n; ++i) {
        e.push_back(event(min(a[i].p.x, a[i].q.x), +1,
5.3
        e.push_back(event(max(a[i].p.x, a[i].q.x), -1,
             <u>i</u>));
54
55
      sort(e.begin(), e.end());
56
      s.clear();
      where.resize(a.size());
58
      for (size_t i = 0; i < e.size(); ++i) {</pre>
       11 id = e[i].id;
59
60
        if (e[i].tp == +1) {
          set<seq>::iterator nxt = s.lower_bound(a[id])
               , prv = prev(nxt);
          if (nxt != s.end() && intersect(*nxt, a[id]))
                return make_pair(nxt->id, id);
          if (prv != s.end() && intersect(*prv, a[id]))
                return make pair (prv->id, id);
          where[id] = s.insert(nxt, a[id]);
        } else {
          set<seg>::iterator nxt = next(where[id]), prv
                = prev(where[id]);
          if (nxt != s.end() && prv != s.end() &&
               intersect(*nxt, *prv)) return make_pair(
               prv->id, nxt->id);
68
          s.erase(where[id]);
69
      return make_pair(-1, -1);
72 }
3.8 Minkowski Sum
```

```
11 cross(const pt & p) const { return x * p.y - y
             * p.x; }
 6
    };
     void reorder_polygon(vector<pt> & P) {
       size_t pos = 0;
       for (size_t i = 1; i < P.size(); i++) {</pre>
         if (P[i].y < P[pos].y || (P[i].y == P[pos].y &&
               P[i].x < P[pos].x)) pos = i;
11
12
      rotate(P.begin(), P.begin() + pos, P.end());
13 }
14 vector<pt> minkowski(vector<pt> P, vector<pt> Q) {
15
      // the first vertex must be the lowest
      reorder_polygon(P);
17
       reorder_polygon(Q);
18
      // we must ensure cyclic indexing
19
      P.push_back(P[0]);
20
      P.push_back(P[1]);
21
       Q.push_back(Q[0]);
22
23
       Q.push_back(Q[1]);
      // main part
24
      vector<pt> result;
       size_t i = 0, j = 0;
26
       while (i < P.size() - 2 || j < Q.size() - 2){</pre>
27
        result.push_back(P[i] + Q[j]);
28
         auto cross = (P[i + 1] - P[i]).cross(Q[j + 1] -
               Q[j]);
        if (cross >= 0 && i < P.size() - 2) ++i;</pre>
        if (cross <= 0 && j < Q.size() - 2) ++j;</pre>
31
       return result;
33 }
```

3.9 Nearest Points

```
struct pt {
                        ll x, y, id;
                 struct cmp_x {
                        bool operator()(const pt & a, const pt & b) const
                                 return a.x < b.x || (a.x == b.x && a.y < b.y);</pre>
   8 };
                 struct cmp v {
                        bool operator()(const pt & a, const pt & b) const
                                                 { return a.y < b.y; }
11 };
              11 n;
             vector<pt> a;
                double mindist;
                pair<11, 11> best_pair;
16
               void upd_ans(const pt & a, const pt & b) {
                        double dist = sqrt((a.x - b.x) * (a.x - b.x) + (a.x - 
                                            .y - b.y) * (a.y - b.y);
18
                        if (dist < mindist) {</pre>
19
                                mindist = dist;
20
                                best_pair = {a.id, b.id};
21
22 }
23 vector<pt> t;
24
                void rec(ll l, ll r) {
25
                        if (r - 1 \le 3) {
26
                                 for (11 i = 1; i < r; ++i)
27
                                        for (11 j = i + 1; j < r; ++j)
28
                                                upd_ans(a[i], a[j]);
                                 sort(a.begin() + 1, a.begin() + r, cmp_y());
```

```
return;
      11 m = (1 + r) >> 1, midx = a[m].x;
      rec(1, m);
      rec(m, r);
      merge(a.begin() + 1, a.begin() + m, a.begin() + m
           , a.begin() + r, t.begin(), cmp_y());
      copy(t.begin(), t.begin() + r - 1, a.begin() + 1)
      11 \text{ tsz} = 0;
      for (ll i = 1; i < r; ++i) {</pre>
39
        if (abs(a[i].x - midx) < mindist) {</pre>
40
          for (ll j = tsz - 1; j >= 0 && a[i].y - t[j].
               y < mindist; --j)</pre>
             upd_ans(a[i], t[j]);
          t[tsz++] = a[i];
44
      }
45
46 t.resize(n);
   sort(a.begin(), a.end(), cmp_x());
48 mindist = 1E20;
49 rec(0, n);
```

3.10 Point In Convex

```
struct pt {
      long long x, y;
      pt() {}
      pt (long long \underline{x}, long long \underline{y}) : x(\underline{x}), y(\underline{y}) {}
      pt operator+(const pt &p) const { return pt(x + p
           .x, y + p.y);}
      pt operator-(const pt &p) const { return pt(x - p
           .x, y - p.y);
      long long cross(const pt &p) const { return x * p
           y - y * p.x;
      long long dot(const pt &p) const { return x * p.x
            + y * p.y; }
      long long cross(const pt &a, const pt &b) const {
            return (a - *this).cross(b - *this); }
      long long dot (const pt &a, const pt &b) const {
           return (a - *this).dot(b - *this); }
      long long sqrLen() const { return this->dot(*this
           ); }
    bool lexComp(const pt &1, const pt &r) { return 1.x
          < r.x \mid | (1.x == r.x && 1.y < r.y); }
    int sgn(long long val) { return val > 0 ? 1 : (val
         == 0 ? 0 : -1); }
    vector<pt> seq;
16
    pt translation;
    int n:
18
    bool pointInTriangle(pt a, pt b, pt c, pt point) {
      long long s1 = abs(a.cross(b, c));
      long long s2 = abs(point.cross(a, b)) + abs(point
           .cross(b, c)) + abs(point.cross(c, a));
      return s1 == s2;
22
    void prepare(vector<pt> &points) {
      n = points.size();
      int pos = 0;
      for (int i = 1; i < n; i++) {
        if (lexComp(points[i], points[pos])) pos = i;
2.8
      rotate(points.begin(), points.begin() + pos,
           points.end());
```

```
seq.resize(n);
      for (int i = 0; i < n; i++) seq[i] = points[i +</pre>
           1] - points[0];
      translation = points[0];
34
    bool pointInConvexPolygon(pt point) {
      point = point - translation;
      if (seq[0].cross(point) != 0 && sgn(seq[0].cross(
           point)) != sgn(seq[0].cross(seq[n - 1])))
        return false:
39
      if (seq[n-1].cross(point) != 0 \&\& sqn(seq[n-1])
           1].cross(point)) != sgn(seq[n - 1].cross(seq
        return false;
      if (seq[0].cross(point) == 0)
        return seq[0].sqrLen() >= point.sqrLen();
      int 1 = 0, r = n - 1;
      while (r - 1 > 1) {
        int mid = (1 + r) / 2;
        int pos = mid;
        if (seq[pos].cross(point) >= 0) 1 = mid;
        else r = mid;
      int pos = 1;
      return pointInTriangle(seq[pos], seq[pos + 1], pt
           (0, 0), point);
52
```

3.11 Segment Intersection

```
const double EPS = 1E-9;
    struct pt {
      double x, y;
      bool operator<(const pt& p) const {</pre>
        return x < p.x - EPS \mid \mid (abs(x - p.x) < EPS &&
            y < p.y - EPS);
    };
   struct line {
      double a, b, c;
      line() {}
      line(pt p, pt q) {
        a = p.y - q.y;
       b = q.x - p.x;
       c = -a * p.x - b * p.y;
        norm();
      void norm() {
18
        double z = sqrt(a * a + b * b);
19
        if (abs(z) > EPS) a \neq z, b \neq z, c \neq z;
      double dist(pt p) const { return a * p.x + b * p.
           y + c; }
    double det (double a, double b, double c, double d)
      return a * d - b * c;
   inline bool betw(double 1, double r, double x) {
      return min(1, r) \le x + EPS && x \le max(1, r) +
   inline bool intersect_ld(double a, double b, double
          c, double d) {
      if (a > b) swap(a, b);
      if (c > d) swap(c, d);
      return max(a, c) <= min(b, d) + EPS;</pre>
```

```
Pegaraw
```

```
34 bool intersect(pt a, pt b, pt c, pt d, pt& left, pt
         & right) {
      if (!intersect_ld(a.x, b.x, c.x, d.x) || !
           intersect_ld(a.y, b.y, c.y, d.y)) return
      line m(a, b);
      line n(c, d);
      double zn = det(m.a, m.b, n.a, n.b);
      if (abs(zn) < EPS) {</pre>
        if (abs(m.dist(c)) > EPS || abs(n.dist(a)) >
             EPS) return false;
41
        if (b < a) swap(a, b);
42
        if (d < c) swap(c, d);
43
        left = max(a, c);
44
         right = min(b, d);
4.5
         return true;
46
      } else {
        left.x = right.x = -det(m.c, m.b, n.c, n.b) /
48
        left.y = right.y = -det(m.a, m.c, n.a, n.c) /
49
         return betw(a.x, b.x, left.x) && betw(a.v, b.v,
              left.y) &&
50
               betw(c.x, d.x, left.x) && betw(c.y, d.y,
                     left.v);
51
52 }
```

4 Graph Theory

4.1 Articulation Point

```
void APUtil(vector<vector<ll>>> &adj, ll u, vector<</pre>
         bool> &visited,
    vector<ll> &disc, vector<ll> &low, ll &time, ll
         parent, vector<bool> &isAP) {
      11 children = 0:
      visited[u] = true;
      disc[u] = low[u] = ++time;
      for (auto v : adj[u]) {
        if (!visited[v]) {
          children++;
          APUtil(adj, v, visited, disc, low, time, u,
          low[u] = min(low[u], low[v]);
11
          if (parent != -1 && low[v] >= disc[u]) {
            isAP[u] = true;
13
        } else if (v != parent) {
15
          low[u] = min(low[u], disc[v]);
16
17
18
      if (parent == -1 && children > 1) {
19
        isAP[u] = true;
20
21
    void AP(vector<vector<ll>>> &adj, ll n) {
23
      vector<ll> disc(n), low(n);
      vector<bool> visited(n), isAP(n);
24
25
      11 time = 0, par = -1;
      for (11 u = 0; u < n; u++) {
        if (!visited[u]) {
          APUtil(adj, u, visited, disc, low, time, par,
                isAP);
2.9
```

4.2 Bellman Ford

```
struct Edge {
      int a, b, cost;
    1:
    int n, m, v;
    vector<Edge> edges;
    const int INF = 1000000000;
    void solve() {
      vector<int> d(n, INF);
      d[v] = 0;
      vector<int> p(n, -1);
      int x;
      for (int i = 0; i < n; ++i) {
       \mathbf{x} = -1;
        for (Edge e : edges)
          if (d[e.a] < INF)
            if (d[e.b] > d[e.a] + e.cost) {
              d[e.b] = max(-INF, d[e.a] + e.cost);
              p[e.b] = e.a;
              x = e.b;
      if (x == -1) cout << "No negative cycle from " <<</pre>
      else {
        int y = x;
        for (int i = 0; i < n; ++i) y = p[y];
        vector<int> path;
        for (int cur = y;; cur = p[cur]) {
          path.push_back(cur);
          if (cur == y && path.size() > 1) break;
30
        reverse(path.begin(), path.end());
32
        cout << "Negative cycle: ";</pre>
        for (int u : path) cout << u << ' ';</pre>
```

4.3 Bridge

```
int n:
vector<vector<int>> adj;
vector<bool> visited;
vector<int> tin, low;
int timer;
void dfs (int v, int p = -1) {
  visited[v] = true;
 tin[v] = low[v] = timer++;
  for (int to : adj[v]) {
   if (to == p) continue;
    if (visited[to]) {
      low[v] = min(low[v], tin[to]);
    } else {
      dfs(to, v);
      low[v] = min(low[v], low[to]);
      if (low[to] > tin[v]) IS_BRIDGE(v, to);
```

4.4 Centroid Decomposition

```
vector<vector<int>> adj;
   vector<bool> is_removed;
   vector<int> subtree_size;
   int get_subtree_size(int node, int parent = -1) {
           subtree_size[node] = 1;
           for (int child : adj[node]) {
                   if (child == parent || is_removed[
                        child]) continue;
                    subtree_size[node] +=
                        get_subtree_size(child, node);
           return subtree_size[node];
   int get_centroid(int node, int tree_size, int
        parent = -1) {
           for (int child : adj[node]) {
                   if (child == parent || is_removed[
                        child]) continue;
                    if (subtree_size[child] * 2 >
                        tree_size) return get_centroid
                         (child, tree_size, node);
           return node;
18
   void build_centroid_decomp(int node = 0) {
           int centroid = get_centroid(node,
                get_subtree_size(node));
            // do something
           is_removed[centroid] = true;
            for (int child : adj[centroid]) {
                    if (is_removed[child]) continue;
                    build_centroid_decomp(child);
27 }
```

4.5 Dijkstra

```
q.pop();
14
        if (d v != d[v]) continue;
15
        for (auto edge : adj[v]) {
16
          int to = edge.first, len = edge.second;
17
          if (d[v] + len < d[to]) {</pre>
18
              d[to] = d[v] + len;
19
              p[to] = v;
20
              q.push({d[to], to});
23
24 }
```

4.6 Dinics struct FlowEdge { int v, u; 11 cap, flow = 0;FlowEdge(int v, int u, ll cap) : v(v), u(u), cap(struct Dinic { const 11 flow_inf = 1e18; vector<FlowEdge> edges; vector<vector<int>> adj; int n, m = 0, s, t; vector<int> level, ptr; queue<int> q; 13 Dinic(int n, int s, int t) : n(n), s(s), t(t) { 14 adj.resize(n); 15 level.resize(n); 16 ptr.resize(n); 17 18 void add_edge(int v, int u, 11 cap) { 19 edges.emplace_back(v, u, cap); 20 edges.emplace_back(u, v, 0); 21 adj[v].push_back(m); $adj[u].push_back(m + 1);$ 23 m += 2;24 25 bool bfs() { while (!q.empty()) { 27 int v = q.front(); 28 q.pop(); 29 for (int id : adj[v]) { if (edges[id].cap - edges[id].flow < 1)</pre> continue; if (level[edges[id].u] != -1) continue; level[edges[id].u] = level[v] + 1; q.push(edges[id].u); 34 35 36 return level[t] != -1; 37 38 11 dfs(int v, 11 pushed) { 39 if (pushed == 0) return 0; 40 if (v == t) return pushed; 41 for (int& cid = ptr[v]; cid < (int)adj[v].size</pre> (); cid++) 42 int id = adj[v][cid], u = edges[id].u; 43 if (level[v] + 1 != level[u] || edges[id].cap - edges[id].flow < 1) continue;</pre> 11 tr = dfs(u, min(pushed, edges[id].cap edges[id].flow));

if (tr == 0) continue;

edges[id].flow += tr;

edges[id ^ 1].flow -= tr;

4.5

46

47

```
return tr;
49
50
        return 0;
      11 flow() {
        11 f = 0;
        while (true) {
          fill(level.begin(), level.end(), -1);
56
          level[s] = 0;
          q.push(s);
58
          if (!bfs()) break;
59
          fill(ptr.begin(), ptr.end(), 0);
          while (ll pushed = dfs(s, flow_inf)) f +=
               pushed;
62
        return f;
64 };
 4.7 Edmonds Karp
    int n:
   vector<vector<int>> capacity;
   vector<vector<int>> adi;
 4 int bfs(int s, int t, vector<int>& parent) {
     fill(parent.begin(), parent.end(), -1);
      parent[s] = -2;
      queue<pair<int, int>> q;
      q.push({s, INF});
      while (!q.empty()) {
10
        int cur = q.front().first, flow = q.front().
             second;
        q.pop();
        for (int next : adj[cur]) {
13
          if (parent[next] == -1 && capacity[cur][next
            parent[next] = cur;
            int new_flow = min(flow, capacity[cur][next
                 ]);
            if (next == t) return new flow;
            q.push({next, new_flow});
18
      return 0;
    int maxflow(int s, int t) {
      int flow = 0;
      vector<int> parent(n);
      int new flow;
      while (new_flow = bfs(s, t, parent)) {
       flow += new_flow;
29
        int cur = t;
30
        while (cur != s) {
          int prev = parent[cur];
          capacity[prev][cur] -= new_flow;
          capacity[cur][prev] += new_flow;
```

4.8 Fast Second Mst

return flow;

cur = prev;

```
struct edge {
        int s, e, w, id;
        bool operator<(const struct edge& other) {</pre>
             return w < other.w; }</pre>
   typedef struct edge Edge;
    const int N = 2e5 + 5;
    long long res = 0, ans = 1e18;
8 int n, m, a, b, w, id, 1 = 21;
9 vector<Edge> edges;
10 vector<int> h(N, 0), parent(N, -1), size(N, 0),
         present (N, 0);
11 vector<vector<pair<int, int>>> adj(N), dp(N, vector
         <pair<int, int>>(1));
vector<vector<int>> up(N, vector<int>(1, -1));
    pair<int, int> combine(pair<int, int> a, pair<int,</pre>
         int> b) {
      vector<int> v = {a.first, a.second, b.first, b.
          second);
      int topTwo = -3, topOne = -2;
      for (int c : v) {
        if (c > topOne) {
          topTwo = topOne;
          topOne = c;
        } else if (c > topTwo && c < topOne) topTwo = c</pre>
      return {topOne, topTwo};
23
24
    void dfs(int u, int par, int d) {
     h[u] = 1 + h[par];
      up[u][0] = par;
      dp[u][0] = {d, -1};
      for (auto v : adj[u]) {
29
        if (v.first != par) dfs(v.first, u, v.second);
    pair<int, int> lca(int u, int v) {
      pair<int, int> ans = \{-2, -3\};
      if (h[u] < h[v]) swap(u, v);</pre>
      for (int i = 1 - 1; i >= 0; i--) {
        if (h[u] - h[v] >= (1 << i)) {
          ans = combine(ans, dp[u][i]);
38
          u = up[u][i];
      if (u == v) return ans;
      for (int i = 1 - 1; i >= 0; i--) {
        if (up[u][i] != -1 && up[v][i] != -1 && up[u][i
             ] != up[v][i]) {
          ans = combine(ans, combine(dp[u][i], dp[v][i
          u = up[u][i];
          v = up[v][i];
49
      ans = combine(ans, combine(dp[u][0], dp[v][0]));
50
      return ans;
51
52
53 int main(void) {
      cin >> n >> m;
      for (int i = 1; i <= n; i++) {</pre>
        parent[i] = i;
        size[i] = 1;
59
      for (int i = 1; i <= m; i++) {</pre>
60
        cin >> a >> b >> w; // 1-indexed
```

edges.push_back($\{a, b, w, i - 1\}$);

```
62
 63
       sort(edges.begin(), edges.end());
       for (int i = 0; i <= m - 1; i++) {
 65
         a = edges[i].s;
         b = edges[i].e;
 66
 67
         w = edges[i].w;
 68
         id = edges[i].id;
         if (unite_set(a, b)) {
           adj[a].emplace_back(b, w);
           adj[b].emplace_back(a, w);
           present[id] = 1;
 73
           res += w;
 74
         }
 75
 76
       dfs(1, 0, 0);
       for (int i = 1; i \le 1 - 1; i++) {
 78
         for (int j = 1; j <= n; ++j) {</pre>
 79
           if (up[j][i - 1] != -1) {
 80
             int v = up[j][i - 1];
 81
             up[j][i] = up[v][i - 1];
 82
             dp[j][i] = combine(dp[j][i-1], dp[v][i-
 84
 85
 86
       for (int i = 0; i <= m - 1; i++) {
 87
         id = edges[i].id;
         w = edges[i].w;
 88
 89
         if (!present[id]) {
 90
           auto rem = lca(edges[i].s, edges[i].e);
 91
           if (rem.first != w) {
 92
             if (ans > res + w - rem.first) ans = res +
                  w - rem.first;
           } else if (rem.second != -1) {
             if (ans > res + w - rem.second) ans = res +
                   w - rem.second;
 95
 96
         }
 97
       cout << ans << "\n";
 99
       return 0;
100 }
```

4.9 Find Cycle

```
bool dfs(ll v) {
      color[v] = 1;
      for (ll u : adj[v]) {
        if (color[u] == 0) {
          parent[u] = v;
          if (dfs(u)) {
            return true;
        } else if (color[u] == 1) {
          cycle_end = v;
          cycle_start = u;
          return true;
13
14
15
      color[v] = 2;
16
      return false;
17 }
18 void find_cycle() {
      color.assign(n, 0);
20
      parent.assign(n, -1);
      cycle_start = -1;
      for (11 v = 0; v < n; v++) {
```

```
if (color[v] == 0 && dfs(v)) {
          break;
26
      if (cycle_start == -1) {
        cout << "Acyclic" << endl;</pre>
      } else {
30
        vector<ll> cycle;
        cycle.push_back(cycle_start);
        for (11 v = cycle_end; v != cycle_start; v =
             parent[v]) {
          cycle.push_back(v);
34
35
        cycle.push_back(cycle_start);
        reverse(cycle.begin(), cycle.end());
        cout << "Cycle found: ";</pre>
        for (11 v : cycle) {
         cout << v << ' ';
40
        cout << '\n';
43 }
```

4.10 Floyd Warshall

4.11 Ford Fulkerson

```
1 bool bfs(ll n, vector<vector<ll>>> &r_graph, ll s,
         11 t, vector<11> &parent) {
      vector<bool> visited(n, false);
      queue<11> q;
      q.push(s);
5
      visited[s] = true;
      parent[s] = -1;
      while (!q.empty()) {
       11 u = q.front();
        q.pop();
        for (11 v = 0; v < n; v++) {
          if (!visited[v] && r_graph[u][v] > 0) {
            if (v == t) {
              parent[v] = u;
              return true;
15
            q.push(v);
            parent[v] = u;
            visited[v] = true;
```

```
return false;
    11 ford_fulkerson(ll n, vector<vector<ll>>> graph,
         11 s, 11 t) {
      11 u, v;
      vector<vector<11>> r_graph;
      for (u = 0; u < n; u++)
        for (v = 0; v < n; v++)
          r_graph[u][v] = graph[u][v];
      vector<11> parent;
31
      11 \text{ max\_flow} = 0;
      while (bfs(n, r_graph, s, t, parent)) {
        11 path_flow = INF;
34
        for (v = t; v != s; v = parent[v]) {
          u = parent[v];
          path_flow = min(path_flow, r_graph[u][v]);
        for (v = t; v != s; v = parent[v]) {
          u = parent[v];
          r_graph[u][v] -= path_flow;
          r_graph[v][u] += path_flow;
        max flow += path flow;
      return max_flow;
46
```

4.12 Hierholzer

```
void print_circuit (vector<vector<ll>>> &adj) {
      map<11, 11> edge_count;
      for (ll i = 0; i < adj.size(); i++) {</pre>
        edge_count[i] = adj[i].size();
      if (!adj.size()) {
        return;
      stack<ll> curr_path;
      vector<ll> circuit;
      curr_path.push(0);
      11 curr_v = 0;
      while (!curr_path.empty()) {
        if (edge_count[curr_v]) {
          curr_path.push(curr_v);
          11 next_v = adj[curr_v].back();
          edge_count[curr_v]--;
          adj[curr_v].pop_back();
          curr_v = next_v;
          circuit.push_back(curr_v);
          curr_v = curr_path.top();
          curr_path.pop();
2.4
25
      for (ll i = circuit.size() - 1; i >= 0; i--) {
        cout << circuit[i] << ' ';
28
29 }
```

4.13 Hungarian

```
1 vector<int> u (n+1), v (m+1), p (m+1), way (m+1);
2 for (int i=1; i<=n; ++i) {
3  p[0] = i;
4 int j0 = 0;</pre>
```

```
vector<int> minv (m+1, INF);
      vector<bool> used (m+1, false);
      do {
        used[j0] = true;
         int i0 = p[j0], delta = INF, j1;
10
         for (int j=1; j<=m; ++j)</pre>
          if (!used[j]) {
             int cur = A[i0][j]-u[i0]-v[j];
             if (cur < minv[j]) minv[j] = cur, way[j] =</pre>
14
             if (minv[j] < delta) delta = minv[j], j1 =</pre>
15
16
         for (int j=0; j<=m; ++j)</pre>
17
          if (used[j]) u[p[j]] += delta, v[j] -= delta
18
           else minv[j] -= delta;
19
         j0 = j1;
20
      } while (p[j0] != 0);
21
        int j1 = way[j0];
23
        p[j0] = p[j1];
24
         j0 = j1;
25
      } while (†0);
26
27
    vector<int> ans (n+1);
28
    for (int j=1; j<=m; ++j)</pre>
      ans[p[j]] = j;
30 int cost = -v[0];
```

4.14 Is Bipartite

```
1 bool is_bipartite(vector<ll> &col, vector<vector<ll</pre>
         >> &adj, ll n) {
       queue<pair<11, 11>> q;
      for (11 i = 0; i < n; i++) {</pre>
        if (col[i] == -1) {
          q.push({i, 0});
          col[i] = 0;
          while (!q.empty()) {
             pair<11, 11> p = q.front();
             q.pop();
10
             11 v = p.first, c = p.second;
             for (11 j : adj[v]) {
              if (col[j] == c) {
13
                 return false;
15
               if (col[j] == -1) {
16
                 col[j] = (c ? 0 : 1);
17
                 q.push({j, col[j]});
18
19
20
21
23
      return true;
```

4.15 Is Cyclic

```
if (!vis[v] && is_cyclic_util(v, adj, vis, rec)
            ) return true;
 6
        else if (rec[v]) return true;
8
      rec[u] = false;
9
      return false;
    bool is_cyclic(int n, vector<vector<int>> &adj) {
12
      vector<bool> vis(n, false), rec(n, false);
      for (int i = 0; i < n; i++)
        if (!vis[i] && is_cyclic_util(i, adj, vis, rec)
            ) return true;
      return false;
16 }
```

4.16 Kahn

```
void kahn(vector<vector<ll>>> &adj) {
     11 n = adj.size();
      vector<ll> in_degree(n, 0);
      for (11 u = 0; u < n; u++)
      for (ll v: adj[u]) in_degree[v]++;
      queue<11> q;
      for (11 i = 0; i < n; i++)
       if (in_degree[i] == 0)
         q.push(i);
      11 cnt = 0;
      vector<ll> top_order;
      while (!q.empty()) {
       11 u = q.front();
        q.pop();
        top_order.push_back(u);
        for (ll v : adj[u])
         if (--in_degree[v] == 0) q.push(v);
18
        cnt++;
19
      if (cnt != n) {
        cout << -1 << '\n';
        return;
      // print top_order
```

4.17 Kosaraju

return adj_t;

```
void topo_sort(int u, vector<vector<int>>& adj,
        vector<bool>& vis, stack<int>& stk) {
     vis[u] = true;
     for (int v : adj[u]) {
       if (!vis[v]) {
         topo_sort(v, adj, vis, stk);
6
7
8
     stk.push(u);
9
   vector<vector<int>> transpose(int n, vector<vector<</pre>
        int>>& adj) {
     vector<vector<int>> adj_t(n);
     for (int u = 0; u < n; u++) {
       for (int v : adj[u]) {
         adj_t[v].push_back(u);
```

```
void get_scc(int u, vector<vector<int>>& adj_t,
         vector<bool>& vis, vector<int>& scc) {
      vis[u] = true;
23
      scc.push_back(u);
      for (int v : adj_t[u]) {
        if (!vis[v]) {
          get_scc(v, adj_t, vis, scc);
31
    void kosaraju(int n, vector<vector<int>>& adj,
         vector<vector<int>>& sccs) {
      vector<bool> vis(n, false);
      stack<int> stk;
      for (int u = 0; u < n; u++) {
       if (!vis[u]) {
          topo_sort(u, adj, vis, stk);
38
      vector<vector<int>> adj t = transpose(n, adj);
      for (int u = 0; u < n; u++) {
       vis[u] = false;
      while (!stk.empty()) {
       int u = stk.top();
        stk.pop();
       if (!vis[u]) {
          vector<int> scc;
          get_scc(u, adj_t, vis, scc);
49
          sccs.push_back(scc);
52
```

4.18 Kruskals

```
1 struct Edge {
      int u, v, weight;
      bool operator<(Edge const& other) {</pre>
        return weight < other.weight;</pre>
6
   };
   int n;
   vector<Edge> edges;
    int cost = 0;
   vector<Edge> result;
   DSU dsu = DSU(n);
    sort(edges.begin(), edges.end());
    for (Edge e : edges) {
     if (dsu.find_set(e.u) != dsu.find_set(e.v)) {
        cost += e.weight;
        result.push_back(e);
        dsu.union_sets(e.u, e.v);
19 }
```

4.19 Kuhn

```
int n, k;
vector<vector<int>> g;
vector<int> mt;
vector<bool> used;
```

```
bool try_kuhn(int v) {
      if (used[v]) return false;
      used[v] = true;
      for (int to : g[v]) {
        if (mt[to] == -1 || try_kuhn(mt[to])) {
10
          mt[to] = v;
11
          return true;
12
13
14
      return false;
15 }
16 int main() {
17
      mt.assign(k, -1);
18
        vector<bool> used1(n, false);
19
        for (int v = 0; v < n; ++v) {
20
         for (int to : q[v]) {
21
           if (mt[to] == -1) {
              mt[to] = v;
23
              used1[v] = true;
24
              break;
25
26
          }
27
28
        for (int v = 0; v < n; ++v) {
29
          if (used1[v]) continue;
          used.assign(n, false);
31
          try_kuhn(v);
        for (int i = 0; i < k; ++i)
34
          if (mt[i] != -1)
35
            printf("%d %d\n", mt[i] + 1, i + 1);
36 }
```

4.20 Lowest Common Ancestor

```
1 struct LCA {
      vector<ll> height, euler, first, segtree;
      vector<bool> visited;
      LCA(vector<vector<ll>>> &adj, ll root = 0) {
       n = adj.size();
       height.resize(n);
       first.resize(n);
        euler.reserve(n * 2);
10
        visited.assign(n, false);
11
        dfs(adj, root);
        11 m = euler.size();
13
        segtree.resize(m * 4);
14
        build(1, 0, m - 1);
15
16
      void dfs(vector<vector<ll>> &adj, ll node, ll h =
        visited[node] = true;
18
        height[node] = h;
19
         first[node] = euler.size();
20
        euler.push_back(node);
         for (auto to : adj[node]) {
          if (!visited[to]) {
23
            dfs(adj, to, h + 1);
            euler.push_back(node);
25
26
27
      void build(ll node, ll b, ll e) {
29
        if (b == e) segtree[node] = euler[b];
         11 \text{ mid} = (b + e) / 2;
```

```
build(node << 1, b, mid);</pre>
           build(node << 1 | 1, mid + 1, e);
           11 1 = segtree[node << 1], r = segtree[node</pre>
                << 1 | 1];
           segtree[node] = (height[1] < height[r]) ? 1 :</pre>
36
      11 query(11 node, 11 b, 11 e, 11 L, 11 R) {
        if (b > R | | e < L) return -1;</pre>
40
        if (b >= L && e <= R) return segtree[node];</pre>
        11 \text{ mid} = (b + e) >> 1;
         11 left = query(node << 1, b, mid, L, R);</pre>
43
         11 right = query(node << 1 | 1, mid + 1, e, L,</pre>
              R);
44
         if (left == -1) return right;
45
         if (right == -1) return left;
46
         return height[left] < height[right] ? left :</pre>
              right;
48
      ll lca(ll u, ll v) {
49
        11 left = first[u], right = first[v];
50
         if (left > right) swap(left, right);
         return query(1, 0, euler.size() - 1, left,
              right);
52
53 };
```

4.21 Maximum Bipartite Matching

```
bool bpm(ll n, ll m, vector<vector<bool>> &bpGraph,
          11 u, vector<bool> &seen, vector<11> &matchR)
      for (11 \ v = 0; \ v < m; \ v++) {
        if (bpGraph[u][v] && !seen[v]) {
          seen[v] = true;
          if (matchR[v] < 0 \mid \mid bpm(n, m, bpGraph,
               matchR[v], seen, matchR)) {
            matchR[v] = u;
            return true;
 8
9
      return false;
12
   11 maxBPM(11 n, 11 m, vector<vector<bool>> &bpGraph
      vector<11> matchR(m, -1);
      11 \text{ result} = 0;
      for (11 u = 0; u < n; u++) {
        vector<bool> seen(m, false);
        if (bpm(n, m, bpGraph, u, seen, matchR)) {
19
          result++;
20
      return result;
```

4.22 Min Cost Flow

```
1 struct Edge {
2   int from, to, capacity, cost;
3   };
4  vector<vector<int>> adj, cost, capacity;
5  const int INF = 1e9;
```

```
void shortest_paths(int n, int v0, vector<int>& d,
         vector<int>& p) {
      d.assign(n, INF);
      d[v0] = 0;
      vector<bool> inq(n, false);
      queue<int> q;
      q.push(v0);
      p.assign(n, -1);
      while (!q.empty()) {
        int u = q.front();
        q.pop();
        inq[u] = false;
        for (int v : adj[u]) {
18
          if (capacity[u][v] > 0 && d[v] > d[u] + cost[
              u][v]) {
            d[v] = d[u] + cost[u][v];
            p[v] = u;
            if (!inq[v]) {
             inq[v] = true;
              q.push(v);
28
29
    int min_cost_flow(int N, vector<Edge> edges, int K,
          int s, int t) {
      adj.assign(N, vector<int>());
      cost.assign(N, vector<int>(N, 0));
      capacity.assign(N, vector<int>(N, 0));
      for (Edge e : edges) {
34
        adj[e.from].push_back(e.to);
        adj[e.to].push_back(e.from);
        cost[e.from][e.to] = e.cost;
        cost[e.to][e.from] = -e.cost;
38
        capacity[e.from][e.to] = e.capacity;
39
      int flow = 0;
      int cost = 0;
      vector<int> d, p;
      while (flow < K) {</pre>
       shortest_paths(N, s, d, p);
        if (d[t] == INF) break;
        int f = K - flow, cur = t;
        while (cur != s) {
        f = min(f, capacity[p[cur]][cur]);
          cur = p[cur];
        flow += f;
        cost += f * d[t];
        cur = t;
        while (cur != s) {
          capacity[p[cur]][cur] -= f;
          capacity[cur][p[cur]] += f;
          cur = p[cur];
58
59
60
      if (flow < K) return -1;</pre>
      else return cost;
62 }
```

4.23 Prim

```
const int INF = 1000000000;
struct Edge {
   int w = INF, to = -1;
   bool operator((Edge const& other) const {
```

```
Pegaraw
```

```
return make_pair(w, to) < make_pair(other.w,
             other.to);
 7
    };
8 int n;
    vector<vector<Edge>> adj;
10 void prim() {
      int total_weight = 0;
      vector<Edge> min_e(n);
13
      \min e[0].w = 0;
      set < Edge > q;
      q.insert({0, 0});
      vector<bool> selected(n, false);
17
      for (int i = 0; i < n; ++i) {</pre>
18
        if (q.empty()) {
19
         cout << "No MST!" << endl;</pre>
20
          exit(0);
21
        int v = q.begin()->to;
23
         selected[v] = true;
24
        total_weight += q.begin()->w;
25
         q.erase(q.begin());
         if (min e[v].to != -1) cout << v << " " <<</pre>
             min_e[v].to << endl;</pre>
         for (Edge e : adj[v]) {
28
          if (!selected[e.to] && e.w < min_e[e.to].w) {</pre>
29
             q.erase({min_e[e.to].w, e.to});
             min_e[e.to] = \{e.w, v\};
             q.insert({e.w, e.to});
35
      cout << total_weight << endl;</pre>
```

4.24 Topological Sort

```
void dfs(ll v) {
      visited[v] = true;
      for (ll u : adj[v]) {
       if (!visited[u]) {
          dfs(u);
      ans.push_back(v);
9
10 void topological_sort() {
      visited.assign(n, false);
      ans.clear();
      for (ll i = 0; i < n; ++i) {</pre>
14
       if (!visited[i]) {
15
          dfs(i);
16
17
18
      reverse(ans.begin(), ans.end());
```

4.25 Zero One Bfs

```
1 vector<int> d(n, INF);
2 d[s] = 0;
3 deque<int> q;
4 q.push_front(s);
5 while (!q.empty()) {
    int v = q.front();
```

```
7    q.pop_front();
8    for (auto edge : adj[v]) {
9       int u = edge.first, w = edge.second;
10       if (d[v] + w < d[u]) {
11            d[u] = d[v] + w;
12            if (w == 1) q.push_back(u);
13            else q.push_front(u);
14       }
15    }
16 }</pre>
```

5 Math

5.1 Chinese Remainder Theorem

```
struct Congruence {
     11 a, m;
   };
   11 chinese_remainder_theorem(vector<Congruence>
         const& congruences) {
      11 M = 1:
      for (auto const& congruence : congruences) M *=
           congruence.m;
      11 \text{ solution} = 0:
      for (auto const& congruence : congruences) {
       11 a i = congruence.a;
        11 M_i = M / congruence.m;
        11 N_i = mod_inv(M_i, congruence.m);
        solution = (solution + a_i * M_i % M * N_i) % M
14
      return solution;
16 }
```

5.2 Extended Euclidean

```
int gcd(int a, int b, int& x, int& y) {
   if (b == 0) {
      x = 1;
      y = 0;
      return a;
   }
   int x1, y1, d = gcd(b, a % b, x1, y1);
   x = y1;
   y = x1 - y1 * (a / b);
   return d;
}
```

5.3 Factorial Modulo

```
11 return res;
12 }
```

5.4 Fast Fourier Transform

```
1 using cd = complex<double>;
    const double PI = acos(-1);
    void fft(vector<cd>& a, bool invert) {
      int n = a.size();
      if (n == 1) return;
      vector<cd> a0 (n / 2), a1 (n / 2);
      for (int i = 0; 2 * i < n; i++) {
       a0[i] = a[2 * i];
        a1[i] = a[2 * i + 1];
      fft(a0, invert);
      fft(a1, invert);
      double ang = 2 * PI / n * (invert ? -1 : 1);
      cd w(1), wn(cos(ang), sin(ang));
      for (int i = 0; 2 * i < n; i++) {
       a[i] = a0[i] + w * a1[i];
        a[i + n / 2] = a0[i] - w * a1[i];
        if (invert) {
         a[i] /= 2;
          a[i + n / 2] /= 2;
        w \star = wn;
12.4
    vector<int> multiply(vector<int> const& a, vector<</pre>
         int> const& b) {
        vector<cd> fa(a.begin(), a.end()), fb(b.begin()
            , b.end());
        int n = 1;
        while (n < a.size() + b.size()) n <<= 1;</pre>
        fa.resize(n);
        fb.resize(n);
        fft(fa, false);
        fft(fb, false);
        for (int i = 0; i < n; i++) fa[i] *= fb[i];</pre>
34
        fft(fa, true);
        vector<int> result(n);
        for (int i = 0; i < n; i++) result[i] = round(</pre>
             fa[i].real());
        return result:
```

5.5 Fibonacci

```
13 - Periodic sequence modulo p
14 - sum[i=1..n]f[i] = f[n+2] - 1
15 - sum[i=0..n-1]f[2i+1] = f[2n]
16 - sum[i=1..n]f[2i] = f[2n+1] - 1
17 - sum[i=1..n]f[i]^2 = f[n]f[n+1]
18 Fibonacci encoding:
19 1. Iterate through the Fibonacci numbers from the
         largest to the
   smallest until you find one less than or equal to n
   2. Suppose this number was F_i. Subtract F_i from
         n Łand put a 1 Ł
   in the i-2 position of the code word (indexing from
         0 from the
23 leftmost to the rightmost bit).
24 3. Repeat until there is no remainder.
25 4. Add a final 1 Lto the codeword to indicate its
   Closed-form: f[n] = (((1 + rt(5))/2)^n - ((1 - rt))^n)
        (5)) / 2) ^n)/rt(5)
28 struct matrix {
29
      11 mat[2][2];
      matrix friend operator *(const matrix &a, const
          matrix &b) {
        matrix c:
32
        for (int i = 0; i < 2; i++) {
         for (int j = 0; j < 2; j++) {
34
           c.mat[i][j] = 0;
            for (int k = 0; k < 2; k++) c.mat[i][j] +=
                a.mat[i][k] * b.mat[k][j];
36
       }
38
        return c;
39
40 };
41 matrix matpow(matrix base, 11 n) {
      matrix ans{ {
43
       {1, 0},
44
       {0, 1}
45
      } };
46
      while (n) {
47
       if (n & 1) ans = ans * base;
48
       base = base * base;
49
       n >>= 1;
51
      return ans;
52
53 11 fib(int n) {
54
      matrix base{ {
       {1, 1},
56
       {1, 0}
57
      } };
58
      return matpow(base, n).mat[0][1];
59 }
60 pair<int, int> fib (int n) {
61
    if (n == 0) return {0, 1};
62
      auto p = fib(n >> 1);
63
      int c = p.first * (2 * p.second - p.first);
64
      int d = p.first * p.first + p.second * p.second;
65
      if (n & 1) return {d, c + d};
66
      else return {c, d};
67 }
```

5.6 Find All Solutions

```
1 bool find_any_solution(ll a, ll b, ll c, ll &x0, ll | 2 Matrix exponentation:
```

```
&y0, 11 &g) {
      q = qcd \text{ extended (abs (a), abs (b), x0, y0);}
      if (c % q) return false;
      x0 *= c / g;
      y0 \star = c / q;
      if (a < 0) x0 = -x0;
      if (b < 0) y0 = -y0;
      return true;
10 void shift_solution(ll & x, ll & y, ll a, ll b, ll
        cnt) {
      x += cnt * b;
      y -= cnt * a;
    11 find_all_solutions(ll a, ll b, ll c, ll minx, ll
          maxx, 11 miny, 11 maxy) {
15
      11 x, y, g;
      if (!find_any_solution(a, b, c, x, y, g)) return
           0;
      a /= g;
      b /= q;
      11 \text{ sign}_a = a > 0 ? +1 : -1;
      11 \text{ sign } b = b > 0 ? +1 : -1;
      shift_solution(x, y, a, b, (minx - x) / b);
      if (x < minx) shift_solution(x, y, a, b, sign_b);</pre>
      if (x > maxx) return 0;
      11 1x1 = x;
      shift_solution(x, y, a, b, (maxx - x) / b);
      if (x > maxx) shift_solution(x, y, a, b, -sign_b)
      11 \text{ rx1} = x;
      shift_solution(x, y, a, b, -(miny - y) / a);
29
      if (y < miny) shift_solution(x, y, a, b, -sign_a)</pre>
30
      if (y > maxy) return 0;
      11 \ 1x2 = x;
      shift_solution(x, y, a, b, -(maxy - y) / a);
33
      if (y > maxy) shift_solution(x, y, a, b, sign_a);
      11 \text{ rx2} = x;
      if (1x2 > rx2) swap(1x2, rx2);
36
      11 1x = max(1x1, 1x2), rx = min(rx1, rx2);
      if (lx > rx) return 0;
      return (rx - 1x) / abs(b) + 1;
39 }
5.7 Linear Sieve
 void linear_sieve(ll N, vector<ll> &lowest_prime,
         vector<ll> &prime) {
      for (11 i = 2; i <= N; i++) {
        if (lowest_prime[i] == 0) {
          lowest_prime[i] = i;
          prime.push_back(i);
 6
        for (11 j = 0; i * prime[j] <= N; j++) {</pre>
         lowest_prime[i * prime[j]] = prime[j];
```

if (prime[j] == lowest_prime[i]) break;

5.8 Matrix

}

10

12 }

```
1 /*
2 Matrix exponentation
```

```
f[n] = af[n-1] + bf[n-2] + cf[n-3]
    Use:
    |f[n]| | |a|b|c||f[n-1]|
    |f[n-1]|=|1 0 0||f[n-2]|
    |f[n-2]| |0 1 0||f[n-3]|
   To get:
    |f[n] | |a b c|^(n-2)|f[2]|
   |f[n-1]|=|1 0 0| |f[1]|
   |f[n-2]| |0 1 0|
                          |f[0]|
   struct Matrix { int mat[MAX_N][MAX_N]; };
   Matrix matrix_mul(Matrix a, Matrix b) {
     Matrix ans; int i, j, k;
      for (i = 0; i < MAX_N; i++)</pre>
      for (j = 0; j < MAX_N; j++)</pre>
      for (ans.mat[i][j] = k = 0; k < MAX_N; k++)
       ans.mat[i][j] += a.mat[i][k] * b.mat[k][j];
      return ans;
   Matrix matrix_pow(Matrix base, int p) {
     Matrix ans; int i, j;
      for (i = 0; i < MAX_N; i++)</pre>
        for (j = 0; j < MAX N; j++)
         ans.mat[i][j] = (i == j);
      while (p) {
        if (p & 1) ans = matrix_mul(ans, base);
        base = matrix_mul(base, base);
        p >>= 1;
32
     return ans;
33 }
5.9 Miller Rabin
    using u64 = uint64_t;
    using u128 = __uint128_t;
```

```
u64 binpower(u64 base, u64 e, u64 mod) {
     u64 \text{ result} = 1;
     base %= mod;
     while (e) {
       if (e & 1) result = (u128) result * base % mod;
        base = (u128) base * base % mod;
        e >>= 1;
11
     return result;
    bool check_composite(u64 n, u64 a, u64 d, 11 s) {
     u64 x = binpower(a, d, n);
      if (x == 1 \mid | x == n - 1) return false;
      for (11 r = 1; r < s; r++) {</pre>
       x = (u128) x * x % n;
        if (x == n - 1) return false:
     return true:
    bool miller_rabin(u64 n) {
      if (n < 2) return false;</pre>
      11 r = 0;
      u64 d = n - 1;
      while ((d \& 1) == 0) {
27
        d >>= 1:
2.8
        r++;
      for (11 a : {2, 3, 5, 7, 11, 13, 17, 19, 23, 29,
           31, 37}) {
        if (n == a) return true;
```

if (check_composite(n, a, d, r)) return false;

```
33    }
34    return true;
35    }
```

5.10 Modulo Inverse

```
1  ll mod_inv(ll a, ll m) {
2    if (m == 1) return 0;
3    ll m0 = m, x = 1, y = 0;
4    while (a > 1) {
5        ll q = a / m, t = m;
6        m = a % m;
7        a = t;
8        t = y;
9        y = x - q * y;
10        x = t;
11    }
12    if (x < 0) x += m0;
13    return x;
14 }</pre>
```

5.11 Pollard Rho Brent

```
1 11 mult(11 a, 11 b, 11 mod) {
      return (__int128_t) a * b % mod;
 4 ll f(ll x, ll c, ll mod) {
      return (mult(x, x, mod) + c) % mod;
 6
 7 ll pollard_rho_brent(ll n, ll x0 = 2, ll c = 1) {
      11 \times = x0, g = 1, q = 1, xs, y, m = 128, 1 = 1;
      while (g == 1) {
       v = x;
11
        for (11 i = 1; i < 1; i++) x = f(x, c, n);
12
13
        while (k < 1 \&\& g == 1) {
14
        xs = x;
15
         for (11 i = 0; i < m && i < 1 - k; i++) {
16
          x = f(x, c, n);
17
            q = mult(q, abs(y - x), n);
18
19
          g = \underline{gcd}(q, n);
20
          k += m;
21
        1 *= 2;
23
24
      if (g == n) {
25
       do {
26
        xs = f(xs, c, n);
27
         g = \underline{gcd}(abs(xs - y), n);
28
        } while (g == 1);
29
      return q;
31 }
```

5.12 Range Sieve

```
1 vector<bool> range_sieve(ll 1, ll r) {
2    l1 n = sqrt(r);
3    vector<bool> is_prime(n + 1, true);
4    vector<ll> prime;
5    is_prime[0] = is_prime[1] = false;
```

```
prime.push_back(2);
      for (l1 i = 4; i <= n; i += 2) is_prime[i] =</pre>
           false:
      for (11 i = 3; i <= n; i += 2) {
9
       if (is_prime[i]) {
          prime.push_back(i);
          for (ll j = i * i; j <= n; j += i) is_prime[j</pre>
               ] = false;
12
13
     vector<bool> result(r - 1 + 1, true);
      for (ll i : prime)
    for (11 \ j = \max(i * i, (1 + i - 1) / i * i); j
           <= r; j += i)
          result[j - 1] = false;
18 if (1 == 1) result[0] = false;
      return result;
20 }
```

5.13 Segmented Sieve

```
1 vector<ll> segmented_sieve(ll n) {
      const 11 S = 10000;
     11 nsgrt = sgrt(n);
      vector<char> is_prime(nsqrt + 1, true);
      vector<11> prime;
      is_prime[0] = is_prime[1] = false;
      prime.push_back(2);
      for (11 i = 4; i <= nsgrt; i += 2) {
      is_prime[i] = false;
      for (11 i = 3; i <= nsqrt; i += 2) {</pre>
      if (is_prime[i]) {
         prime.push_back(i);
          for (ll j = i * i; j <= nsqrt; j += i) {</pre>
           is_prime[j] = false;
       }
      vector<1l> result;
      vector<char> block(S);
      for (11 k = 0; k * S \le n; k++) {
       fill(block.begin(), block.end(), true);
       for (11 p : prime) {
          for (11 j = max((k * S + p - 1) / p, p) * p -
              k * S; j < S; j += p) {
            block[j] = false;
        if (k == 0) {
         block[0] = block[1] = false;
30
        for (11 i = 0; i < S && k * S + i <= n; i++) {</pre>
         if (block[i]) {
33
            result.push_back(k * S + i);
35
       }
      return result;
38 }
```

5.14 Sum Of Divisors

```
for (int i = 2; (11) i * i <= num; i++) {</pre>
       if (num % i == 0) {
          int e = 0;
          do {
           e++;
           num /= i;
          } while (num % i == 0);
          11 \text{ sum} = 0, \text{ pow} = 1;
          do {
           sum += pow;
           pow *= i;
          } while (e-- > 0);
          total *= sum;
      if (num > 1) total *= (1 + num);
      return total;
20 }
```

5.15 Tonelli Shanks

```
1 11 legendre(ll a, ll p) {
     return bin_pow_mod(a, (p - 1) / 2, p);
4 ll tonelli_shanks(ll n, ll p) {
    if (legendre(n, p) == p - 1) {
      return -1:
    if (p % 4 == 3) {
     return bin_pow_mod(n, (p + 1) / 4, p);
     11 \ 0 = p - 1, S = 0;
      while (Q \% 2 == 0) {
      0 /= 2;
       S++;
     11 z = 2;
      for (; z < p; z++) {</pre>
      if (legendre(z, p) == p - 1) {
         break;
      11 M = S, c = bin_pow_mod(z, Q, p), t =
         bin_pow_mod(n, Q, p), R = bin_pow_mod(n, Q)
          + 1) / 2, p);
      while (t % p != 1) {
       if (t % p == 0) {
          return 0;
       11 i = 1, t2 = t * t % p;
       for (; i < M; i++) {
         if (t2 % p == 1) {
           break;
         t2 = t2 * t2 % p;
       11 b = bin_pow_mod(c, bin_pow_mod(2, M - i - 1,
             p), p);
       M = i;
       c = b * b % p;
       t = t * c % p;
       R = R * b % p;
39
     return R;
41 }
```

6 Miscellaneous

6.1 Gauss

```
const double EPS = 1e-9;
    const 11 INF = 2;
    11 gauss(vector <vector <double>> a, vector <double>
         &ans) {
      11 n = (11) a.size(), m = (11) a[0].size() - 1;
      vector<11> where (m, -1);
      for (11 col = 0, row = 0; col < m && row < n; ++</pre>
           col) {
         11 sel = row:
         for (11 i = row; i < n; ++i) {
          if (abs(a[i][col]) > abs(a[sel][col])) {
10
             sel = i;
12
13
         if (abs (a[sel][col]) < EPS) {</pre>
14
          continue;
15
16
         for (ll i = col; i <= m; ++i) {</pre>
17
          swap(a[sel][i], a[row][i]);
18
19
         where[col] = row;
20
         for (11 i = 0; i < n; ++i) {
          if (i != row) {
             double c = a[i][col] / a[row][col];
23
             for (11 j = col; j <= m; ++j) {
24
              a[i][j] = a[row][j] * c;
25
26
          }
27
28
        ++row:
29
30
      ans.assign(m, 0);
      for (11 i = 0; i < m; ++i) {
        if (where[i] != -1) {
           ans[i] = a[where[i]][m] / a[where[i]][i];
34
35
36
      for (11 i = 0; i < n; ++i) {
37
         double sum = 0;
38
         for (11 j = 0; j < m; ++j) {
39
          sum += ans[j] * a[i][j];
40
41
        if (abs (sum - a[i][m]) > EPS) {
42
43
44
45
      for (11 i = 0; i < m; ++i) {
46
        if (where[i] == -1) {
47
          return INF;
48
49
      return 1;
51
```

6.2 Ternary Search

```
double ternary_search(double 1, double r) {
   double eps = le-9;
   while (r - 1 > eps) {
    double ml = l + (r - 1) / 3;
}
```

```
5     double m2 = r - (r - 1) / 3;
6     double f1 = f(m1);
7     double f2 = f(m2);
8     if (f1 < f2) {
9         1 = m1;
10     } else {
1         r = m2;
12     }
13     }
14     return f(1);
15 }</pre>
```

7 References

7.1 Ref

```
1 st.insert(4);
 2 st.erase(4);
 3 st.empty();
 4 // permutations
 5 do {
 6 for (int num : nums) {
    cout << num << " ";
 8 }
 9 cout << endl;</pre>
10 } while (next_permutation(nums.begin(), nums.end())
        );
   // bitset
   int num = 27; // Binary representation: 11011
13 bitset<10> s(string("0010011010")); // from right
14 bitset<sizeof(int) * 8> bits(num);
15 int setBits = bits.count();
16 // sort
   sort(v.begin(), v.end());
18 sort(v.rbegin(), v.rend());
    // custom sort
   bool comp(string a, string b) {
    if (a.size() != b.size()) return a.size() < b.size</pre>
        ():
    return a < b;
24 sort(v.begin(), v.end(), comp);
   // binary search
26 int a = 0, b = n-1;
    while (a \le b) { int k = (a+b)/2; if (array[k] == x
        ) {
28 // x found at index k
   } if (array[k] > x) b = k-1; else a = k+1;}
   for (auto it = s.begin(); it != s.end(); it++) {
32 cout << *it << "\n";
33 }
    // hamming distance
    int hamming(int a, int b) {
    return __builtin_popcount(a^b);
   // custom comparator for pq
39 class Compare {
41 bool operator()(T a, T b){
42 if(cond) return true; // do not swap
    return false;
44 }
45 };
```

46 priority_queue<PII, vector<PII>, Compare> ds;

Strings

8.1 Count Unique Substrings

```
int count_unique_substrings(string const& s) {
  int n = s.size();
  const int p = 31;
  const int m = 1e9 + 9;
  vector<long long> p_pow(n);
  p_pow[0] = 1;
  for (int i = 1; i < n; i++) p_pow[i] = (p_pow[i -</pre>
        1] * p) % m;
  vector<long long> h(n + 1, 0);
  for (int i = 0; i < n; i++) h[i + 1] = (h[i] + (s
       [i] - 'a' + 1) * p_pow[i]) % m;
  int cnt = 0;
  for (int 1 = 1; 1 <= n; 1++) {
    unordered_set<long long> hs;
    for (int i = 0; i <= n - 1; i++) {</pre>
      long long cur_h = (h[i + 1] + m - h[i]) % m;
      cur_h = (cur_h * p_pow[n - i - 1]) % m;
      hs.insert(cur_h);
    cnt += hs.size();
  return cnt;
```

8.2 Finding Repetitions

```
vector<int> z_function(string const& s) {
      int n = s.size();
      vector<int> z(n);
      for (int i = 1, l = 0, r = 0; i < n; i++) {
        if (i \le r) z[i] = min(r - i + 1, z[i - 1]);
        while (i + z[i] < n \&\& s[z[i]] == s[i + z[i]])
             z[i]++;
        if (i + z[i] - 1 > r) {
          1 = i;
          \mathbf{r} = \mathbf{i} + \mathbf{z}[\mathbf{i}] - 1;
11
      return z;
    int get_z(vector<int> const& z, int i) {
      if (0 <= i && i < (int) z.size()) return z[i];</pre>
      else return 0;
    vector<pair<int, int>> repetitions;
    void convert_to_repetitions(int shift, bool left,
         int cntr, int 1, int k1, int k2) {
      for (int 11 = \max(1, 1 - k2); 11 \le \min(1, k1);
           11++) {
        if (left && 11 == 1) break;
        int 12 = 1 - 11;
        int pos = shift + (left ? cntr - 11 : cntr - 1
             - 11 + 1);
        repetitions.emplace_back(pos, pos + 2 * 1 - 1);
26
    void find_repetitions(string s, int shift = 0) {
28
      int n = s.size();
29
      if (n == 1) return;
      int nu = n / 2;
```

```
int nv = n - nu;
      string u = s.substr(0, nu);
      string v = s.substr(nu);
      string ru(u.rbegin(), u.rend());
      string rv(v.rbegin(), v.rend());
      find_repetitions(u, shift);
37
      find_repetitions(v, shift + nu);
38
      vector<int> z1 = z_function(ru);
39
      vector<int> z2 = z_function(v + '#' + u);
      vector<int> z3 = z_function(ru + '#' + rv);
41
      vector<int> z4 = z_function(v);
42
      for (int cntr = 0; cntr < n; cntr++) {</pre>
43
        int 1, k1, k2;
44
        if (cntr < nu) {</pre>
45
          1 = nu - cntr;
46
          k1 = get_z(z1, nu - cntr);
47
          k2 = get_z(z2, nv + 1 + cntr);
48
49
          1 = cntr - nu + 1;
50
          k1 = get_z(z3, nu + 1 + nv - 1 - (cntr - nu))
          k2 = get_z(z4, (cntr - nu) + 1);
        if (k1 + k2 >= 1) convert_to_repetitions(shift,
              cntr < nu, cntr, 1, k1, k2);</pre>
54
55 }
```

8.3 Group Identical Substrings

```
1 vector<vector<int>> group_identical_strings(vector<</pre>
         string> const& s) {
      int n = s.size();
      vector<pair<long long, int>> hashes(n);
      for (int i = 0; i < n; i++) hashes[i] = {</pre>
           compute_hash(s[i]), i};
      sort(hashes.begin(), hashes.end());
      vector<vector<int>> groups;
      for (int i = 0; i < n; i++) {
        if (i == 0 || hashes[i].first != hashes[i - 1].
             first) groups.emplace_back();
        groups.back().push_back(hashes[i].second);
10
11
      return groups;
12 }
```

8.4 Hashing

8.5 Knuth Morris Pratt

```
1 vector<ll> prefix_function(string s) {
```

```
ll n = (ll) s.length();
      vector<ll> pi(n);
      for (ll i = 1; i < n; i++) {
       11 j = pi[i - 1];
       while (j > 0 \&\& s[i] != s[j]) j = pi[j - 1];
       if (s[i] == s[j]) j++;
8
       pi[i] = j;
      return pi;
   // count occurences
   vector<int> ans(n + 1);
14 for (int i = 0; i < n; i++)
     ans[pi[i]]++;
16 for (int i = n-1; i > 0; i--)
     ans[pi[i-1]] += ans[i];
18 for (int i = 0; i <= n; i++)
    ans[i]++;
```

8.6 Longest Common Prefix

```
vector<int> lcp_construction(string const& s,
     vector<int> const& p) {
   int n = s.size();
  vector<int> rank(n, 0);
   for (int i = 0; i < n; i++) rank[p[i]] = i;</pre>
   int k = 0;
   vector<int> lcp(n-1, 0);
   for (int i = 0; i < n; i++) {
    if (rank[i] == n - 1) {
      k = 0:
       continue;
     int j = p[rank[i] + 1];
     while (i + k < n \&\& j + k < n \&\& s[i + k] == s[
         j + k]) k++;
     lcp[rank[i]] = k;
    if (k) k--;
  return lcp;
```

8.7 Manacher

```
vector<int> manacher_odd(string s) {
     int n = s.size();
      s = "$" + s + "^";
      vector<int> p(n + 2);
      int 1 = 1, r = 1;
      for (int i = 1; i \le n; i++) {
        p[i] = max(0, min(r - i, p[1 + (r - i)]));
        while (s[i - p[i]] == s[i + p[i]]) p[i]++;
        if(i + p[i] > r) 1 = i - p[i], r = i + p[i];
      return vector<int>(begin(p) + 1, end(p) - 1);
12
13
   vector<int> manacher(string s) {
      string t;
      for(auto c: s) t += string("#") + c;
      auto res = manacher_odd(t + "#");
      return vector<int>(begin(res) + 1, end(res) - 1);
18
```

```
8.8 Rabin Karp
```

```
vector<11> rabin_karp(string const& s, string const
     const 11 p = 31, m = 1e9 + 9;
     11 S = s.size(), T = t.size();
     vector<ll> p_pow(max(S, T));
     p_pow[0] = 1;
     for (ll i = 1; i < (ll) p_pow.size(); i++) p_pow[</pre>
          i] = (p_pow[i-1] * p) % m;
     vector<11> h(T + 1, 0);
     for (11 i = 0; i < T; i++) h[i + 1] = (h[i] + (t[
          i] - 'a' + 1) * p_pow[i]) % m;
     11 h s = 0;
     for (11 i = 0; i < S; i++) h_s = (h_s + (s[i] - '
         a' + 1) * p_pow[i]) % m;
     vector<11> occurences;
     for (11 i = 0; i + S - 1 < T; i++) {
      11 \text{ cur}_h = (h[i + S] + m - h[i]) % m;
       if (cur_h == h_s * p_pow[i] % m) occurences.
           push_back(i);
     return occurences;
```

8.9 Suffix Array

```
vector<int> sort_cyclic_shifts(string const& s) {
      int n = s.size();
      const int alphabet = 256;
      vector<int> p(n), c(n), cnt(max(alphabet, n), 0);
      for (int i = 0; i < n; i++) cnt[s[i]]++;</pre>
      for (int i = 1; i < alphabet; i++) cnt[i] += cnt[</pre>
           i - 1];
      for (int i = 0; i < n; i++) p[--cnt[s[i]]] = i;</pre>
      c[p[0]] = 0;
      int classes = 1;
      for (int i = 1; i < n; i++) {</pre>
        if (s[p[i]] != s[p[i-1]]) classes++;
        c[p[i]] = classes - 1;
      vector<int> pn(n), cn(n);
      for (int h = 0; (1 << h) < n; ++h) {
        for (int i = 0; i < n; i++) {
          pn[i] = p[i] - (1 << h);
          if (pn[i] < 0)
            pn[i] += n;
        fill(cnt.begin(), cnt.begin() + classes, 0);
        for (int i = 0; i < n; i++) cnt[c[pn[i]]]++;</pre>
        for (int i = 1; i < classes; i++) cnt[i] += cnt</pre>
             [i - 1];
        for (int i = n-1; i >= 0; i--) p[--cnt[c[pn[i
             ]]]] = pn[i];
        cn[p[0]] = 0;
        classes = 1;
        for (int i = 1; i < n; i++) {
28
          pair<int, int> cur = {c[p[i]], c[(p[i] + (1
               << h)) % n]};
          pair < int, int > prev = {c[p[i-1]], c[(p[i-1]] + }
                (1 << h)) % n]};
          if (cur != prev) ++classes;
          cn[p[i]] = classes - 1;
        c.swap(cn);
```

```
34
35
36
     }
     return p;
37
    vector<int> build_suff_arr(string s) {
38
     s += "$";
39
      vector<int> sorted_shifts = sort_cyclic_shifts(s)
40
      sorted_shifts.erase(sorted_shifts.begin());
41
      return sorted_shifts;
42 }
43 // compare two substrings
44 int compare(int i, int j, int 1, int k) {
      pair<int, int> a = {c[k][i], c[k][(i + 1 - (1 <<</pre>
```

```
k)) % n]};
46    pair<int, int> b = {c[k][j], c[k][(j + 1 - (1 << k)) % n]};
47    return a == b ? 0 : a < b ? -1 : 1;
48  }

8.10    Z Function

1    vector<int> z_function(string s) {
2        int n = s.size();
3        vector<int> z(n);
```

f(n) = O(g(n))	iff \exists positive c, n_0 such that	$n = n(n+1)$ $n = n(n+1)(2n+1)$ $n = n(2(n+1))^2$
	$0 \le f(n) \le cg(n) \ \forall n \ge n_0.$	$\sum_{i=1}^{n} i = \frac{n(n+1)}{2}, \sum_{i=1}^{n} i^2 = \frac{n(n+1)(2n+1)}{6}, \sum_{i=1}^{n} i^3 = \frac{n^2(n+1)^2}{4}.$
$f(n) = \Omega(g(n))$	iff \exists positive c, n_0 such that $f(n) \geq cg(n) \geq 0 \ \forall n \geq n_0$.	In general:
$f(n) = \Theta(g(n))$	iff $f(n) = O(g(n))$ and $f(n) = \Omega(g(n))$.	$\sum_{i=1}^{n} i^{m} = \frac{1}{m+1} \left[(n+1)^{m+1} - 1 - \sum_{i=1}^{n} \left((i+1)^{m+1} - i^{m+1} - (m+1)i^{m} \right) \right]$
f(n) = o(g(n))	iff $\lim_{n\to\infty} f(n)/g(n) = 0$.	$\sum_{i=1}^{n-1} i^m = \frac{1}{m+1} \sum_{k=0}^m \binom{m+1}{k} B_k n^{m+1-k}.$
$\lim_{n \to \infty} a_n = a$	iff $\forall \epsilon > 0$, $\exists n_0$ such that $ a_n - a < \epsilon$, $\forall n \ge n_0$.	Geometric series:
$\sup S$	least $b \in \mathbb{R}$ such that $b \geq s$, $\forall s \in S$.	$\sum_{i=0}^{n} c^{i} = \frac{c^{n+1} - 1}{c - 1}, c \neq 1, \sum_{i=0}^{\infty} c^{i} = \frac{1}{1 - c}, \sum_{i=1}^{\infty} c^{i} = \frac{c}{1 - c}, c < 1,$
$\inf S$	greatest $b \in \mathbb{R}$ such that $b \le s$, $\forall s \in S$.	$\sum_{i=0}^{n} ic^{i} = \frac{nc^{n+2} - (n+1)c^{n+1} + c}{(c-1)^{2}}, c \neq 1, \sum_{i=0}^{\infty} ic^{i} = \frac{c}{(1-c)^{2}}, c < 1.$
$ \liminf_{n \to \infty} a_n $	$\lim_{n \to \infty} \inf \{ a_i \mid i \ge n, i \in \mathbb{N} \}.$	Harmonic series: $n = n + n = n + n = n = n = n = n = n = $
$\limsup_{n \to \infty} a_n$	$\lim_{n \to \infty} \sup \{ a_i \mid i \ge n, i \in \mathbb{N} \}.$	$H_n = \sum_{i=1}^n \frac{1}{i}, \qquad \sum_{i=1}^n iH_i = \frac{n(n+1)}{2}H_n - \frac{n(n-1)}{4}.$
$\binom{n}{k}$	Combinations: Size k subsets of a size n set.	$\sum_{i=1}^{n} H_i = (n+1)H_n - n, \sum_{i=1}^{n} {i \choose m} H_i = {n+1 \choose m+1} \left(H_{n+1} - \frac{1}{m+1} \right).$
$\begin{bmatrix} n \\ k \end{bmatrix}$	Stirling numbers (1st kind): Arrangements of an n element set into k cycles.	1. $\binom{n}{k} = \frac{n!}{(n-k)!k!}$, 2. $\sum_{k=0}^{n} \binom{n}{k} = 2^{n}$, 3. $\binom{n}{k} = \binom{n}{n-k}$,
$\binom{n}{k}$	Stirling numbers (2nd kind): Partitions of an n element set into k non-empty sets.	$4. \binom{n}{k} = \frac{n}{k} \binom{n-1}{k-1}, \qquad \qquad 5. \binom{n}{k} = \binom{n-1}{k} + \binom{n-1}{k-1}, \\ 6. \binom{n}{m} \binom{m}{k} = \binom{n}{k} \binom{n-k}{m-k}, \qquad \qquad 7. \sum_{k=0}^{n} \binom{r+k}{k} = \binom{r+n+1}{n},$
$\left\langle {n\atop k}\right\rangle$	1st order Eulerian numbers: Permutations $\pi_1\pi_2\pi_n$ on $\{1,2,,n\}$ with k ascents.	8. $\sum_{k=0}^{n} {k \choose m} = {n+1 \choose m+1},$ 9. $\sum_{k=0}^{n} {r \choose k} {s \choose n-k} = {r+s \choose n},$
$\langle\!\langle {n \atop k} \rangle\!\rangle$	2nd order Eulerian numbers.	10. $\binom{n}{k} = (-1)^k \binom{k-n-1}{k}$, 11. $\binom{n}{1} = \binom{n}{n} = 1$,
C_n	Catalan Numbers: Binary trees with $n+1$ vertices.	12. $\binom{n}{2} = 2^{n-1} - 1$, 13. $\binom{n}{k} = k \binom{n-1}{k} + \binom{n-1}{k-1}$,
$14. \begin{bmatrix} n \\ 1 \end{bmatrix} = (n-1)$	15. $\begin{bmatrix} n \\ 2 \end{bmatrix} = (n - 1)^n$	$16. \begin{bmatrix} n \\ n \end{bmatrix} = 1, \qquad \qquad 17. \begin{bmatrix} n \\ k \end{bmatrix} \ge \begin{Bmatrix} n \\ k \end{Bmatrix},$
18. $\begin{bmatrix} n \\ k \end{bmatrix} = (n-1)$	$\binom{n-1}{k} + \binom{n-1}{k-1}, 19. \ \binom{n}{n-1}$	
$22. \left\langle {n \atop 0} \right\rangle = \left\langle {n \atop n-1} \right\rangle$	$\binom{n}{-1} = 1,$ 23. $\binom{n}{k} = \binom{n}{k}$	$\binom{n}{n-1-k}$, $24. \left\langle \binom{n}{k} \right\rangle = (k+1) \left\langle \binom{n-1}{k} \right\rangle + (n-k) \left\langle \binom{n-1}{k-1} \right\rangle$,
25. $\left\langle {0\atop k}\right\rangle = \left\{ {1\atop 0}\right\}$	if $k = 0$, otherwise 26. $\begin{cases} r \\ 1 \end{cases}$	$\binom{n}{2} = 2^n - n - 1,$ 27. $\binom{n}{2} = 3^n - (n+1)2^n + \binom{n+1}{2},$
28. $x^n = \sum_{k=0}^n \binom{n}{k}$	$\left\langle {x+k \choose n}, \qquad $ 29. $\left\langle {n \atop m} \right\rangle = \sum_{k=1}^m$	
		32. $\left\langle \left\langle \begin{array}{c} n \\ 0 \end{array} \right\rangle = 1,$ 33. $\left\langle \left\langle \begin{array}{c} n \\ n \end{array} \right\rangle = 0$ for $n \neq 0$,
$34. \; \left\langle \!\! \left\langle \!\! \begin{array}{c} n \\ k \end{array} \!\! \right\rangle = (k + 1)^n$	-1 $\left\langle \left\langle \left$	
$36. \left\{ \begin{array}{c} x \\ x-n \end{array} \right\} = \sum_{k}^{n} \left\{ \begin{array}{c} x \\ x \end{array} \right\}$	$\sum_{k=0}^{n} \left\langle \!\! \left\langle n \atop k \right\rangle \!\! \right\rangle \!\! \binom{x+n-1-k}{2n},$	37. $\binom{n+1}{m+1} = \sum_{k} \binom{n}{k} \binom{k}{m} = \sum_{k=0}^{n} \binom{k}{m} (m+1)^{n-k},$

The Chinese remainder theorem: There exists a number C such that:

 $C \equiv r_1 \mod m_1$

: : :

 $C \equiv r_n \bmod m_n$

if m_i and m_j are relatively prime for $i \neq j$. Euler's function: $\phi(x)$ is the number of positive integers less than x relatively prime to x. If $\prod_{i=1}^{n} p_i^{e_i}$ is the prime factorization of x then

$$\phi(x) = \prod_{i=1}^{n} p_i^{e_i - 1} (p_i - 1).$$

Euler's theorem: If a and b are relatively prime then

$$1 \equiv a^{\phi(b)} \bmod b$$
.

Fermat's theorem:

$$1 \equiv a^{p-1} \bmod p.$$

The Euclidean algorithm: if a > b are integers then

$$gcd(a, b) = gcd(a \mod b, b).$$

If $\prod_{i=1}^{n} p_i^{e_i}$ is the prime factorization of x

$$S(x) = \sum_{d|x} d = \prod_{i=1}^{n} \frac{p_i^{e_i+1} - 1}{p_i - 1}.$$

Perfect Numbers: x is an even perfect number iff $x = 2^{n-1}(2^n - 1)$ and $2^n - 1$ is prime. Wilson's theorem: n is a prime iff

$$(n-1)! \equiv -1 \mod n$$
.

$$\mu(i) = \begin{cases} (n-1)! = -1 \bmod n. \\ \text{M\"obius inversion:} \\ \mu(i) = \begin{cases} 1 & \text{if } i = 1. \\ 0 & \text{if } i \text{ is not square-free.} \\ (-1)^r & \text{if } i \text{ is the product of} \\ r & \text{distinct primes.} \end{cases}$$
 If

 If

$$G(a) = \sum_{d|a} F(d),$$

$$F(a) = \sum_{d|a} \mu(d) G\left(\frac{a}{d}\right).$$

Prime numbers:

$$p_n = n \ln n + n \ln \ln n - n + n \frac{\ln \ln n}{\ln n}$$

$$+O\left(\frac{n}{\ln n}\right),$$

$$\pi(n) = \frac{n}{\ln n} + \frac{n}{(\ln n)^2} + \frac{2!n}{(\ln n)^3} + O\left(\frac{n}{(\ln n)^4}\right).$$

T	 0	٠,	٠	
				ns

Loop An edge connecting a vertex to itself.

Directed Each edge has a direction. SimpleGraph with no loops or multi-edges.

WalkA sequence $v_0e_1v_1\ldots e_\ell v_\ell$. TrailA walk with distinct edges. Pathtrail with distinct

vertices.

ConnectedA graph where there exists a path between any two

vertices.

ComponentΑ maximal connected subgraph.

TreeA connected acyclic graph. Free tree A tree with no root. DAGDirected acyclic graph. EulerianGraph with a trail visiting each edge exactly once.

Hamiltonian Graph with a cycle visiting each vertex exactly once.

CutA set of edges whose removal increases the number of components.

Cut-setA minimal cut. Cut edge A size 1 cut.

k-Connected A graph connected with the removal of any k-1vertices.

k-Tough $\forall S \subseteq V, S \neq \emptyset$ we have $k \cdot c(G - S) \le |S|$.

A graph where all vertices k-Regular have degree k.

k-Factor Α k-regular spanning subgraph.

Matching A set of edges, no two of which are adjacent.

CliqueA set of vertices, all of which are adjacent.

Ind. set A set of vertices, none of which are adjacent.

Vertex cover A set of vertices which cover all edges.

Planar graph A graph which can be embeded in the plane.

Plane graph An embedding of a planar

$$\sum_{v \in V} \deg(v) = 2m.$$

If G is planar then n - m + f = 2, so

$$f \le 2n - 4, \quad m \le 3n - 6.$$

Any planar graph has a vertex with degree ≤ 5 .

Notation:

E(G)Edge set Vertex set V(G)

c(G)Number of components

G[S]Induced subgraph deg(v)Degree of v

Maximum degree $\Delta(G)$

 $\delta(G)$ Minimum degree $\chi(G)$ Chromatic number

 $\chi_E(G)$ Edge chromatic number G^c Complement graph K_n Complete graph

 K_{n_1,n_2} Complete bipartite graph

Ramsev number

Geometry

Projective coordinates: (x, y, z), not all x, y and z zero.

$$(x, y, z) = (cx, cy, cz) \quad \forall c \neq 0.$$

Cartesian Projective (x, y)(x, y, 1)y = mx + b(m, -1, b)

x = c(1,0,-c)Distance formula, L_p and L_{∞}

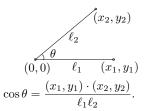
$$\sqrt{(x_1 - x_0)^2 + (y_1 - y_0)^2},$$
$$\left[|x_1 - x_0|^p + |y_1 - y_0|^p\right]^{1/p},$$

$$\lim_{p \to \infty} \left[|x_1 - x_0|^p + |y_1 - y_0|^p \right]^{1/p}.$$

Area of triangle $(x_0, y_0), (x_1, y_1)$ and (x_2, y_2) :

$$\frac{1}{2} \operatorname{abs} \begin{vmatrix} x_1 - x_0 & y_1 - y_0 \\ x_2 - x_0 & y_2 - y_0 \end{vmatrix}.$$

Angle formed by three points:



Line through two points (x_0, y_0) and (x_1, y_1) :

$$\begin{vmatrix} x & y & 1 \\ x_0 & y_0 & 1 \\ x_1 & y_1 & 1 \end{vmatrix} = 0.$$

Area of circle, volume of sphere:

$$A = \pi r^2, \qquad V = \frac{4}{3}\pi r^3.$$

If I have seen farther than others, it is because I have stood on the shoulders of giants.

- Issac Newton

Taylor's series:

$$f(x) = f(a) + (x - a)f'(a) + \frac{(x - a)^2}{2}f''(a) + \dots = \sum_{i=0}^{\infty} \frac{(x - a)^i}{i!}f^{(i)}(a).$$

Expansions:

Expansions:
$$\frac{1}{1-x} = 1 + x + x^2 + x^3 + x^4 + \cdots = \sum_{i=0}^{\infty} x^i,$$

$$\frac{1}{1-cx} = 1 + x + x^2 + x^3 + x^4 + \cdots = \sum_{i=0}^{\infty} c^i x^i,$$

$$\frac{1}{1-cx} = 1 + x + x^2 + x^3 + x^4 + \cdots = \sum_{i=0}^{\infty} c^i x^i,$$

$$\frac{1}{1-x^n} = 1 + x^n + x^{2n} + x^{3n} + \cdots = \sum_{i=0}^{\infty} ix^{ii},$$

$$\frac{x}{(1-x)^2} = x + 2x^2 + 3x^3 + 4x^4 + \cdots = \sum_{i=0}^{\infty} ix^i,$$

$$x^k \frac{d^n}{dx^n} \left(\frac{1}{1-x}\right) = x + 2^{n}x^2 + 3^n x^3 + 4^n x^4 + \cdots = \sum_{i=0}^{\infty} i^n x^i,$$

$$e^x = 1 + x + \frac{1}{2}x^2 + \frac{1}{6}x^3 + \cdots = \sum_{i=0}^{\infty} i^n x^i,$$

$$\ln(1+x) = x - \frac{1}{2}x^2 + \frac{1}{3}x^3 - \frac{1}{4}x^4 - \cdots = \sum_{i=0}^{\infty} (-1)^{i+1} \frac{x^i}{i},$$

$$\ln \frac{1}{1-x} = x + \frac{1}{2}x^2 + \frac{1}{3}x^3 + \frac{1}{4}x^4 + \cdots = \sum_{i=0}^{\infty} (-1)^{i} \frac{x^{2i+1}}{(2i+1)!},$$

$$\cos x = 1 - \frac{1}{2!}x^2 + \frac{1}{1!}x^4 - \frac{1}{6!}x^6 + \cdots = \sum_{i=0}^{\infty} (-1)^{i} \frac{x^{2i+1}}{(2i+1)!},$$

$$(1+x)^n = 1 + nx + \frac{n(n-1)}{2}x^2 + \cdots = \sum_{i=0}^{\infty} (-1)^{i} \frac{x^{2i+1}}{(2i+1)!},$$

$$(1+x)^n = 1 + nx + \frac{n(n-1)}{2}x^2 + \cdots = \sum_{i=0}^{\infty} (-1)^{i} \frac{x^{2i+1}}{(2i+1)!},$$

$$\frac{1}{(1-x)^{n+1}} = 1 + (n+1)x + (\frac{n+2}{2})x^2 + \cdots = \sum_{i=0}^{\infty} (-1)^{i} \frac{x^{2i+1}}{(i)},$$

$$\frac{x}{e^x - 1} = 1 - \frac{1}{2}x + \frac{1}{12}x^2 - \frac{1}{126}x^4 + \cdots = \sum_{i=0}^{\infty} (\frac{i+n}{i})x^i,$$

$$\frac{1}{\sqrt{1-4x}} = 1 + x + 2x^2 + 5x^3 + \cdots = \sum_{i=0}^{\infty} (\frac{2i}{i})x^i,$$

$$\frac{1}{\sqrt{1-4x}} = 1 + x + 2x^2 + 6x^3 + \cdots = \sum_{i=0}^{\infty} (\frac{2i}{i})x^i,$$

$$\frac{1}{1-x} \ln \frac{1}{1-x} = x + \frac{3}{2}x^2 + \frac{1}{16}x^3 + \frac{25}{24}x^4 + \cdots = \sum_{i=0}^{\infty} H_i x^i,$$

$$\frac{1}{1-x} \ln \frac{1}{1-x} = x + \frac{3}{2}x^2 + \frac{1}{16}x^3 + \frac{25}{24}x^4 + \cdots = \sum_{i=0}^{\infty} H_i x^i,$$

$$\frac{1}{2} \left(\ln \frac{1}{1-x} \right)^2 = \frac{1}{2}x^2 + \frac{3}{4}x^3 + \frac{11}{24}x^4 + \cdots = \sum_{i=0}^{\infty} F_{ii}x^i.$$

$$\frac{x}{1-x-x^2} = x + x^2 + 2x^3 + 3x^4 + \cdots = \sum_{i=0}^{\infty} F_{ii}x^i.$$

Ordinary power series:

$$A(x) = \sum_{i=0}^{\infty} a_i x^i.$$

Exponential power series:

$$A(x) = \sum_{i=0}^{\infty} a_i \frac{x^i}{i!}.$$

Dirichlet power series:

$$A(x) = \sum_{i=1}^{\infty} \frac{a_i}{i^x}.$$

Binomial theorem

$$(x+y)^n = \sum_{k=0}^n \binom{n}{k} x^{n-k} y^k.$$

$$x^{n} - y^{n} = (x - y) \sum_{k=0}^{n-1} x^{n-1-k} y^{k}.$$

For ordinary power se

$$\alpha A(x) + \beta B(x) = \sum_{i=0}^{\infty} (\alpha a_i + \beta b_i) x^i,$$

$$x^k A(x) = \sum_{i=k}^{\infty} a_{i-k} x^i,$$

$$\frac{A(x) - \sum_{i=0}^{k-1} a_i x^i}{x^k} = \sum_{i=0}^{\infty} a_{i+k} x^i,$$

$$A(cx) = \sum_{i=0}^{\infty} c^i a_i x^i,$$

$$A'(x) = \sum_{i=0}^{\infty} (i+1)a_{i+1}x^{i},$$

$$xA'(x) = \sum_{i=1}^{\infty} ia_i x^i,$$

$$\int A(x) dx = \sum_{i=1}^{\infty} \frac{a_{i-1}}{i} x^{i},$$

$$\frac{A(x) + A(-x)}{2} = \sum_{i=1}^{\infty} a_{2i} x^{2i},$$

$$\frac{A(x) - A(-x)}{2} = \sum_{i=0}^{\infty} a_{2i+1} x^{2i+1}.$$

Summation: If $b_i = \sum_{i=0}^i a_i$ then

$$B(x) = \frac{1}{1-x}A(x).$$

Convolution:

$$A(x)B(x) = \sum_{i=0}^{\infty} \left(\sum_{j=0}^{i} a_j b_{i-j}\right) x^i.$$

God made the natural numbers; all the rest is the work of man. Leopold Kronecker