UPLB Eliens - Pegaraw Notebook

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    1 Data Structures
```

1.1 Disjoint Set Union

```
struct DSU {
      vector<int> parent, size;
      DSU(int n) {
        parent.resize(n);
        size.resize(n);
        for (int i = 0; i < n; i++) make_set(i);</pre>
8
      void make_set(int v) {
        parent[v] = v;
        size[v] = 1;
      bool is same(int a, int b) { return find set(a)
           == find_set(b); }
      int find_set(int v) { return v == parent[v] ? v :
            parent[v] = find_set(parent[v]); }
      void union_sets(int a, int b) {
       a = find_set(a);
16
        b = find_set(b);
        if (a != b) {
          if (size[a] < size[b]) swap(a, b);</pre>
          parent[b] = a;
          size[a] += size[b];
22
23 };
```

1.2 Minimum Queue

6

```
11, 11>> &s2) {
     if (s1.empty() || s2.empty()) {
       return s1.empty() ? s2.top().second : s1.top().
       return min(s1.top().second, s2.top().second);
 6
 8 void add_element(ll new_element, stack<pair<ll, ll</pre>
        >> &s1) {
     11 minimum = s1.empty() ? new_element : min(
          new_element, s1.top().second);
     s1.push({new_element, minimum});
11
12 11 remove_element(stack<pair<11, 11>> &s1, stack<
        pair<11, 11>> &s2) {
```

```
if (s2.empty()) {
           while (!s1.emptv()) {
             11 element = s1.top().first;
             11 minimum = s2.empty() ? element : min(
                  element, s2.top().second);
             s2.push({element, minimum});
         11 removed_element = s2.top().first;
         s2.pop();
         return removed_element;
13 24 }
```

1.3 Range Add Point Query

```
template<typename T, typename InType = T>
    class SegTreeNode {
    public:
      const T IDN = 0, DEF = 0;
      int i, j;
      T val;
      SegTreeNode<T, InType>* lc, * rc;
      SegTreeNode(int i, int j) : i(i), j(j) {
        if (j - i == 1) {
          lc = rc = nullptr;
          val = DEF;
          return:
        int k = (i + j) / 2;
        lc = new SegTreeNode<T, InType>(i, k);
        rc = new SegTreeNode<T, InType>(k, j);
      SegTreeNode(const vector<InType>& a, int i, int j
          ) : i(i), j(j) {
        if (j - i == 1) {
         lc = rc = nullptr;
          val = (T) a[i];
          return;
        int k = (i + j) / 2;
        lc = new SegTreeNode<T, InType>(a, i, k);
        rc = new SegTreeNode<T, InType>(a, k, j);
28
        val = 0;
      void range_add(int 1, int r, T x) {
        if (r <= i || j <= 1) return;</pre>
        if (1 <= i && j <= r) {
          val += x;
          return;
        lc->range_add(1, r, x);
        rc->range_add(1, r, x);
      T point_query(int k) {
        if (k < i \mid | j \le k) return IDN;
        if (j - i == 1) return val;
        return val + lc->point_query(k) + rc->
            point_query(k);
   template<typename T, typename InType = T>
    class SegTree {
      SegTreeNode<T, InType> root;
```

SegTree(int n) : root(0, n) {}

1.4 Range Add Range Query

```
template<typename T, typename InType = T>
    class SegTreeNode {
    public:
      const T IDN = 0, DEF = 0;
      int i, j;
      T val, to_add = 0;
      SegTreeNode<T, InType>* lc, * rc;
      SegTreeNode(int i, int j) : i(i), j(j) {
        if (j - i == 1) {
          lc = rc = nullptr;
11
          val = DEF:
          return;
13
14
        int k = (i + j) / 2;
        lc = new SegTreeNode<T, InType>(i, k);
16
         rc = new SegTreeNode<T, InType>(k, j);
17
        val = operation(lc->val, rc->val);
18
19
      SegTreeNode(const vector<InType>& a, int i, int j
           ) : i(i), j(j) {
20
         if (j - i == 1) {
21
          lc = rc = nullptr;
          val = (T) a[i];
23
          return:
24
25
         int k = (i + j) / 2;
26
        lc = new SegTreeNode<T, InType>(a, i, k);
27
         rc = new SegTreeNode<T, InType>(a, k, j);
28
        val = operation(lc->val, rc->val);
29
      void propagate() {
31
        if (to_add == 0) return;
        val += to add;
        if (j - i > 1) {
34
          lc->to_add += to_add;
35
          rc->to_add += to_add;
36
37
        to\_add = 0;
38
39
      void range_add(int 1, int r, T delta) {
40
        propagate();
41
         if (r <= i || j <= 1) return;</pre>
42
        if (1 <= i && j <= r) {
43
          to_add += delta;
44
          propagate();
45
         } else {
46
          lc->range_add(l, r, delta);
47
           rc->range_add(l, r, delta);
48
          val = operation(lc->val, rc->val);
49
51
      T range_query(int 1, int r) {
        propagate();
53
        if (1 <= i && j <= r) return val;</pre>
54
         if (j <= 1 || r <= i) return IDN;</pre>
         return operation(lc->range_query(l, r), rc->
             range_query(1, r));
```

```
56  }
57  T operation(T x, T y) {}
58  };
59  template<typename T, typename InType = T>
60  class SegTree {
61  public:
62   SegTreeNode<T, InType> root;
63   SegTree(int n) : root(0, n) {}
64   SegTree(const vector<InType>& a) : root(a, 0, a. size()) {}
65   void range_add(int 1, int r, T delta) { root. range_add(1, r, delta); }
66   T range_query(int 1, int r) { return root. range_query(1, r); }
67  };
```

1.5 Segment Tree

```
template<typename T, typename InType = T>
    class SegTreeNode {
    public:
      const T IDN = 0, DEF = 0;
      int i, j;
      SegTreeNode<T, InType>* lc, * rc;
      SegTreeNode(int i, int j) : i(i), j(j) {
        if (j - i == 1) {
          lc = rc = nullptr;
          val = DEF;
          return;
         int k = (i + j) / 2;
         lc = new SegTreeNode<T, InType>(i, k);
         rc = new SegTreeNode<T, InType>(k, j);
         val = op(lc->val, rc->val);
18
19
      SegTreeNode(const vector<InType>& a, int i, int j
           ) : i(i), j(j) {
        if (j - i == 1) {
          lc = rc = nullptr;
          val = (T) a[i];
          return;
25
        int k = (i + j) / 2;
         lc = new SegTreeNode<T, InType>(a, i, k);
         rc = new SegTreeNode<T, InType>(a, k, j);
        val = op(lc->val, rc->val);
2.9
30
      void set(int k, T x) {
        if (k < i | | j <= k) return;</pre>
        if (j - i == 1) {
          val = x;
          return;
        1c->set(k, x);
         rc \rightarrow set(k, x);
38
        val = op(lc->val, rc->val);
39
40
      T range_query(int 1, int r) {
        if (1 <= i && j <= r) return val;</pre>
42
        if (j <= 1 || r <= i) return IDN;</pre>
4.3
        return op(lc->range_query(l, r), rc->
             range_query(l, r));
45
      T \circ p(T \times, T y) \{ \}
    };
    template<typename T, typename InType = T>
```

```
48  class SegTree {
49  public:
50    SegTreeNode<T, InType> root;
51    SegTree(int n) : root(0, n) {}
52    SegTree(const vector<InType>& a) : root(a, 0, a. size()) {}
53    void set(int k, T x) { root.set(k, x); }
54    T range_query(int 1, int r) { return root. range_query(1, r); }
55  };
```

```
1.6 Segment Tree 2d
    template<typename T, typename InType = T>
    class SegTree2dNode {
    public:
      int i, j, tree_size;
      SegTree<T, InType>* seg_tree;
      SeqTree2dNode<T, InType>* 1c, * rc;
      SegTree2dNode() {}
      SegTree2dNode(const vector<vector<InType>>& a,
           int i, int j) : i(i), j(j) {
        tree_size = a[0].size();
        if (j - i == 1) {
          lc = rc = nullptr;
          seg_tree = new SegTree<T, InType>(a[i]);
          return:
        int k = (i + j) / 2;
        lc = new SegTree2dNode<T, InType>(a, i, k);
        rc = new SegTree2dNode<T, InType>(a, k, j);
        seg_tree = new SegTree<T, InType>(vector<T>(
             tree size));
        operation_2d(1c->seq_tree, rc->seq_tree);
       ~SegTree2dNode() {
        delete 1c;
        delete rc;
24
      void set_2d(int kx, int ky, T x) {
        if (kx < i || j <= kx) return;</pre>
27
        if (j - i == 1) {
2.8
          seg_tree->set(ky, x);
          return;
        1c->set_2d(kx, ky, x);
32
        rc->set_2d(kx, ky, x);
        operation_2d(lc->seg_tree, rc->seg_tree);
34
      T range_query_2d(int lx, int rx, int ly, int ry)
36
        if (lx <= i && j <= rx) return seg_tree->
             range_query(ly, ry);
        if (j <= lx || rx <= i) return -INF;</pre>
38
        return max(lc->range_query_2d(lx, rx, ly, ry),
             rc->range_query_2d(lx, rx, ly, ry));
39
      void operation_2d(SegTree<T, InType>* x, SegTree<</pre>
           T, InType>* y) {
41
        for (int k = 0; k < tree_size; k++) {</pre>
          seg_tree->set(k, max(x->range_query(k, k + 1)
               , y->range_query(k, k + 1)));
43
4.5
    template<typename T, typename InType = T>
    class SegTree2d {
```

```
48  public:
49    SegTree2dNode<T, InType> root;
50    SegTree2d() {}
51    SegTree2d(const vector<vector<InType>>& mat):
        root(mat, 0, mat.size()) {}
52    void set_2d(int kx, int ky, T x) { root.set_2d(kx, ky, x); }
53    T range_query_2d(int lx, int rx, int ly, int ry)
        { return root.range_query_2d(lx, rx, ly, ry)
        ; }
54    };
```

1.7 Sparse Table 1 11 log2_floor(ll i) { return i ? __builtin_clzll(1) - __builtin_clzll(i vector<vector<ll>> build_sum(ll N, ll K, vector<ll> vector<vector<ll>> st(K + 1, vector<ll>(N + 1)); for (ll i = 0; i < N; i++) st[0][i] = array[i];</pre> for (ll i = 1; i <= K; i++)</pre> for $(11 \ j = 0; \ j + (1 << i) <= N; \ j++)$ st[i][j] = st[i - 1][j] + st[i - 1][j + (1 <<(i - 1))];10 return st; 12 11 sum_query(ll L, ll R, ll K, vector<vector<ll>>> & 11 sum = 0;14 for (11 i = K; i >= 0; i--) { 15 **if** ((1 << i) <= R - L + 1) { 16 sum += st[i][L]; 17 L += 1 << i; 18 19 20 return sum; 21 } vector<vector<ll>> build_min(ll N, ll K, vector<ll>

vector<vector<ll>> st(K + 1, vector<ll>(N + 1));

st[i][j] = min(st[i-1][j], st[i-1][j+(1

for (ll i = 0; i < N; i++) st[0][i] = array[i];</pre>

11 min_query(11 L, 11 R, vector<vector<11>> &st) {

return min(st[i][L], st[i][R - (1 << i) + 1]);</pre>

for (11 j = 0; j + (1 << i) <= N; <math>j++)

for (ll i = 1; i <= K; i++)</pre>

 $11 i = log2_floor(R - L + 1);$

<< (i - 1))]);

1.8 Sparse Table 2d

return st;

24

25

26

29

33 }

```
const int N = 100;
int matrix[N][N];
int table[N][N][(int)(log2(N) + 1)][(int)(log2(N) + 1)];

void build_sparse_table(int n, int m) {
   for (int i = 0; i < n; i++)
   for (int j = 0; j < m; j++)
      table[i][j][0][0] = matrix[i][j];
   for (int k = 1; k <= (int)(log2(n)); k++)</pre>
```

```
for (int i = 0; i + (1 << k) - 1 < n; i++)
           for (int j = 0; j + (1 << k) - 1 < m; <math>j++)
             table[i][j][k][0] = min(table[i][j][k -
                  1][0], table[i + (1 << (k - 1))][j][k
                  - 1][0]);
      for (int k = 1; k \le (int)(log2(m)); k++)
        for (int i = 0; i < n; i++)</pre>
           for (int j = 0; j + (1 << k) - 1 < m; <math>j++)
             table[i][j][0][k] = min(table[i][j][0][k -
                  1], table[i][j + (1 << (k - 1))][0][k
                  - 1]);
       for (int k = 1; k <= (int) (log2(n)); k++)</pre>
        for (int 1 = 1; 1 \le (int)(log2(m)); 1++)
18
           for (int i = 0; i + (1 << k) - 1 < n; i++)
19
             for (int j = 0; j + (1 << 1) - 1 < m; j++)
               table[i][j][k][l] = min(
                 min(table[i][j][k-1][l-1], table[i]
                      + (1 << (k - 1))][j][k - 1][1 -
                      1]),
                 min(table[i][j + (1 << (1 - 1))][k -
                      1] [1 - 1], table [i + (1 << (k - 1))
                      ) ] [ \dot{j} + (1 << (1 - 1)) ] [k - 1] [1 -
               );
    int rmq(int x1, int y1, int x2, int y2) {
      int k = log2(x2 - x1 + 1), 1 = log2(y2 - y1 + 1);
      return max (
        \max(table[x1][y1][k][1], table[x2 - (1 << k) +
             1][y1][k][l]),
         \max(\text{table}[x1][y2 - (1 << 1) + 1][k][1], \text{table}[
             x2 - (1 << k) + 1][y2 - (1 << 1) + 1][k][1
             ])
      );
31
```

2 Dynamic Programming

2.1 Divide And Conquer

```
11 m, n;
   vector<ll> dp_before(n), dp_cur(n);
   11 C(11 i, 11 j);
   void compute(ll 1, ll r, ll optl, ll optr) {
      if (1 > r) return;
     11 \text{ mid} = (1 + r) >> 1;
      pair<11, 11 > best = \{LLONG_MAX, -1\};
      for (11 k = opt1; k <= min(mid, optr); k++)</pre>
        best = min(best, \{(k ? dp_before[k - 1] : 0) +
             C(k, mid), k});
      dp_cur[mid] = best.first;
      11 opt = best.second;
      compute(1, mid - 1, opt1, opt);
      compute(mid + 1, r, opt, optr);
14
15 ll solve() {
      for (ll i = 0; i < n; i++) dp_before[i] = C(0, i)</pre>
      for (ll i = 1; i < m; i++) {
        compute (0, n - 1, 0, n - 1);
        dp_before = dp_cur;
      return dp_before[n - 1];
```

2.2 Edit Distance

2.3 Knapsack

2.4 Knuth Optimization

```
1 11 solve() {
     11 N:
     // read N and input
     vector<vector<ll>> dp(N, vector<ll>(N)), opt(N,
          vector<11>(N));
     auto C = [\&](11 i, 11 j) {
       // Implement cost function C.
     };
     for (11 i = 0; i < N; i++) {
       opt[i][i] = i;
       ... // Initialize dp[i][i] according to the
     for (11 i = N - 2; i >= 0; i--) {
       for (11 j = i + 1; j < N; j++) {
         11 \text{ mn} = 11\_\text{MAX}, \text{ cost} = C(i, j);
         for (ll k = opt[i][j-1]; k \le min(j-1,
              opt[i + 1][j]); k++) {
           if (mn \ge dp[i][k] + dp[k + 1][j] + cost) {
              opt[i][j] = k;
             mn = dp[i][k] + dp[k + 1][j] + cost;
         }
```

2.5 Longest Common Subsequence

```
1 11 LCS(string x, string y, 11 n, 11 m) {
      vector < vector < 11 >> dp(n + 1, vector < 11 > (m + 1));
      for (ll i = 0; i <= n; i++) {
         for (ll j = 0; j <= m; j++) {</pre>
           if (i == 0 || j == 0) {
             dp[i][j] = 0;
           } else if (x[i - 1] == y[j - 1]) {
             dp[i][j] = dp[i - 1][j - 1] + 1;
             dp[i][j] = max(dp[i - 1][j], dp[i][j - 1]);
11
        }
13
14
      ll index = dp[n][m];
15
      vector<char> lcs(index + 1);
      lcs[index] = ' \setminus 0';
17
      11 i = n, j = m;
18
      while (i > 0 \&\& j > 0) {
19
        if (x[i-1] == y[j-1]) {
20
           lcs[index - 1] = x[i - 1];
21
           i--;
22
           <del>--;</del>
           index--;
         } else if (dp[i - 1][j] > dp[i][j - 1]) {
25
          <u>i</u>--;
26
        } else {
27
           j--;
28
29
      return dp[n][m];
31 }
```

2.6 Longest Increasing Subsequence

```
1  ll get_ceil_idx(vector<ll> &a, vector<ll> &T, ll l,
          11 r, 11 x) {
      while (r - 1 > 1) {
        11 m = 1 + (r - 1) / 2;
        if (a[T[m]] >= x) {
         r = m:
        } else {
          1 = m:
      return r;
11
    11 LIS(11 n, vector<11> &a) {
      11 len = 1;
14
      vector<11> T(n, 0), R(n, -1);
15
      T[0] = 0;
      for (ll i = 1; i < n; i++) {
17
        if (a[i] < a[T[0]]) {</pre>
         T[0] = i;
19
        } else if (a[i] > a[T[len - 1]]) {
20
          R[i] = T[len - 1];
21
          T[len++] = i;
         } else {
```

2.7 Subset Sum

3 Geometry

3.1 Circle Line Intersection

```
double r, a, b, c; // given as input
    double x0 = -a * c / (a * a + b * b);
    double v0 = -b * c / (a * a + b * b);
    if (c * c > r * r * (a * a + b * b) + EPS) {
      puts ("no points");
    } else if (abs (c *c - r * r * (a * a + b * b)) <</pre>
         EPS) {
      puts ("1 point");
      cout << x0 << ' ' << y0 << '\n';
 8
   } else {
      double d = r * r - c * c / (a * a + b * b);
      double mult = sqrt (d / (a * a + b * b));
      double ax, ay, bx, by;
      ax = x0 + b * mult;
      bx = x0 - b * mult;
      ay = y0 - a * mult;
      by = y0 + a * mult;
      puts ("2 points");
      cout << ax << ' ' << ay << '\n' << bx << ' ' <<
           by << '\n';</pre>
19 }
```

3.2 Convex Hull

```
} else if (v > 0) {
        return +1;
      return 0;
12
13
    bool cw(pt a, pt b, pt c, bool include_collinear) {
      11 o = orientation(a, b, c);
      return o < 0 || (include_collinear && o == 0);</pre>
    bool collinear(pt a, pt b, pt c) {
      return orientation(a, b, c) == 0;
19
   void convex_hull(vector<pt>& a, bool
         include_collinear = false) {
      pt p0 = *min_element(a.begin(), a.end(), [](pt a,
            pt b) {
        return make_pair(a.y, a.x) < make_pair(b.y, b.x</pre>
      });
24
      sort(a.begin(), a.end(), [&p0](const pt& a, const
            pt& b) {
        11 o = orientation(p0, a, b);
        if (o == 0) {
          return (p0.x - a.x) * (p0.x - a.x) + (p0.y - a.x)
               a.y) * (p0.y - a.y)
28
               < (p0.x - b.x) * (p0.x - b.x) + (p0.y -
                    b.y) * (p0.y - b.y);
        return o < 0;
      if (include collinear) {
        11 i = (11) a.size()-1;
34
        while (i \geq= 0 && collinear(p0, a[i], a.back()))
              i --:
        reverse(a.begin()+i+1, a.end());
      vector<pt> st;
38
      for (ll i = 0; i < (ll) a.size(); i++) {</pre>
        while (st.size() > 1 && !cw(st[st.size() - 2],
             st.back(), a[i], include_collinear)) {
          st.pop_back();
        st.push_back(a[i]);
      a = st;
45 }
```

3.3 Line Sweep

```
const double EPS = 1E-9;
    struct pt {
     double x, y;
4
   };
    struct seq {
      pt p, q;
      11 id;
      double get_y (double x) const {
        if (abs(p.x - q.x) < EPS) {
          return p.y;
12
        return p.y + (q.y - p.y) * (x - p.x) / (q.x - p.x)
             .x);
   bool intersect1d(double 11, double r1, double 12,
         double r2) {
      if (l1 > r1) {
```

```
Pegarav
```

```
17
        swap(11, r1);
18
19
      if (12 > r2) {
20
        swap(12, r2);
21
22
      return max(11, 12) <= min(r1, r2) + EPS;</pre>
23
24 11 vec(const pt& a, const pt& b, const pt& c) {
      double s = (b.x - a.x) * (c.y - a.y) - (b.y - a.y)
           ) * (c.x - a.x);
      return abs(s) < EPS ? 0 : s > 0 ? +1 : -1;
27
28 bool intersect(const seg& a, const seg& b) {
29
      return intersect1d(a.p.x, a.g.x, b.p.x, b.g.x) &&
             intersect1d(a.p.y, a.q.y, b.p.y, b.q.y) &&
31
              vec(a.p, a.q, b.p) * vec(a.p, a.q, b.q) <=
                   . 4 . 0
              vec(b.p, b.q, a.p) * vec(b.p, b.q, a.q) <=
                   0:
34 bool operator<(const seg& a, const seg& b) {
         double x = max(min(a.p.x, a.q.x), min(b.p.x, b.
36
         return a.get_y(x) < b.get_y(x) - EPS;</pre>
37
38 struct event {
      double x:
40
      11 tp, id;
41
      event() {}
      event (double x, 11 tp, 11 id) : x(x), tp(tp), id(
           id) {}
43
      bool operator<(const event& e) const {</pre>
44
        if (abs(x - e.x) > EPS) {
45
          return x < e.x;
46
47
        return tp > e.tp;
48
49 };
50 set<seg> s;
51 vector<set<seg>::iterator> where;
    set<seg>::iterator prev(set<seg>::iterator it) {
53
      return it == s.begin() ? s.end() : --it;
55 set<seg>::iterator next(set<seg>::iterator it) {
     return ++it;
58 pair<11, 11> solve(const vector<seg>& a) {
     11 n = (11) a.size();
      vector<event> e;
61
      for (ll i = 0; i < n; ++i) {</pre>
        e.push_back(event(min(a[i].p.x, a[i].q.x), +1,
        e.push_back(event(max(a[i].p.x, a[i].q.x), -1,
             i));
      sort(e.begin(), e.end());
      s.clear();
67
      where.resize(a.size());
68
      for (size_t i = 0; i < e.size(); ++i) {</pre>
69
       11 \text{ id} = e[i].id;
        if (e[i].tp == +1) {
          set<seg>::iterator nxt = s.lower_bound(a[id])
               , prv = prev(nxt);
          if (nxt != s.end() && intersect(*nxt, a[id]))
            return make pair(nxt->id, id);
74
75
          if (prv != s.end() && intersect(*prv, a[id]))
```

3.4 Nearest Points

```
struct pt {
      11 x, y, id;
   };
    struct cmp_x {
      bool operator()(const pt & a, const pt & b) const
         return a.x < b.x || (a.x == b.x && a.y < b.y);</pre>
 8
    };
    struct cmp_y {
      bool operator()(const pt & a, const pt & b) const
        return a.y < b.y;</pre>
    };
    11 n;
    vector<pt> a;
    double mindist;
    pair<11, 11> best_pair;
    void upd_ans(const pt & a, const pt & b) {
      double dist = sqrt((a.x - b.x) * (a.x - b.x) + (a.x - b.x)
           .y - b.y) * (a.y - b.y);
      if (dist < mindist) {</pre>
        mindist = dist;
        best_pair = {a.id, b.id};
23
24
    vector<pt> t;
    void rec(ll 1, ll r) {
      if (r - 1 <= 3) {
        for (11 i = 1; i < r; ++i) {</pre>
          for (11 j = i + 1; j < r; ++j) {
            upd_ans(a[i], a[j]);
         sort(a.begin() + 1, a.begin() + r, cmp_y());
        return;
      11 m = (1 + r) >> 1, midx = a[m].x;
      rec(1, m);
      rec(m, r);
      merge(a.begin() + 1, a.begin() + m, a.begin() + m
           , a.begin() + r, t.begin(), cmp_y());
      copy(t.begin(), t.begin() + r - 1, a.begin() + 1)
      11 \text{ tsz} = 0;
      for (11 i = 1; i < r; ++i) {
        if (abs(a[i].x - midx) < mindist) {</pre>
```

4 Graph Theory

4.1 Articulation Point

```
void APUtil(vector<vector<ll>>> &adj, ll u, vector<</pre>
         bool> &visited,
    vector<11> &disc, vector<11> &low, 11 &time, 11
         parent, vector<bool> &isAP) {
      11 children = 0;
      visited[u] = true;
      disc[u] = low[u] = ++time;
      for (auto v : adj[u]) {
        if (!visited[v]) {
8
          children++;
          APUtil(adj, v, visited, disc, low, time, u,
              isAP);
          low[u] = min(low[u], low[v]);
          if (parent != -1 && low[v] >= disc[u]) {
            isAP[u] = true;
        } else if (v != parent) {
          low[u] = min(low[u], disc[v]);
      if (parent == -1 && children > 1) {
        isAP[u] = true;
    void AP(vector<vector<ll>> &adj, ll n) {
      vector<ll> disc(n), low(n);
      vector<bool> visited(n), isAP(n);
      11 time = 0, par = -1;
      for (11 u = 0; u < n; u++) {
        if (!visited[u]) {
2.8
          APUtil(adj, u, visited, disc, low, time, par,
                isAP);
      for (11 u = 0; u < n; u++) {
        if (isAP[u]) {
          cout << u << " ";
34
36
```

4.2 Bellman Ford

```
1 struct Edge {
2   int a, b, cost;
3   };
4   int n, m, v;
```

```
vector<Edge> edges;
    const int INF = 1000000000;
    void solve() {
      vector<int> d(n, INF);
      d[v] = 0;
      vector<int> p(n, -1);
11
      for (int i = 0; i < n; ++i) {</pre>
13
       \mathbf{x} = -1;
        for (Edge e : edges)
         if (d[e.a] < INF)</pre>
            if (d[e.b] > d[e.a] + e.cost) {
17
              d[e.b] = max(-INF, d[e.a] + e.cost);
18
              p[e.b] = e.a;
19
              x = e.b;
20
21
      if (x == -1) cout << "No negative cycle from " <<
            v;
      else {
        int y = x;
25
         for (int i = 0; i < n; ++i) y = p[y];
26
        vector<int> path;
27
        for (int cur = y;; cur = p[cur]) {
28
          path.push_back(cur);
29
          if (cur == y && path.size() > 1) break;
31
        reverse(path.begin(), path.end());
        cout << "Negative cycle: ";</pre>
         for (int u : path) cout << u << ' ';</pre>
35 }
```

4.3 Bridge

```
1 int n;
   vector<vector<int>> adi;
    vector<bool> visited;
   vector<int> tin, low;
 5 int timer;
    void dfs(int v, int p = -1) {
     visited[v] = true;
     tin[v] = low[v] = timer++;
      for (int to : adj[v]) {
       if (to == p) continue;
11
       if (visited[to]) {
        low[v] = min(low[v], tin[to]);
13
      } else {
         dfs(to, v);
         low[v] = min(low[v], low[to]);
16
         if (low[to] > tin[v]) IS_BRIDGE(v, to);
17
18
     }
19 }
20 void find_bridges() {
      timer = 0;
      visited.assign(n, false);
23
      tin.assign(n, -1);
24
      low.assign(n, -1);
      for (int i = 0; i < n; ++i) {
26
       if (!visited[i]) dfs(i);
27
```

4.4 Dijkstra

```
const int INF = 1000000000;
   vector<vector<pair<int, int>>> adj;
3 void dijkstra(int s, vector<int> & d, vector<int> &
         p) {
     int n = adj.size();
     d.assign(n, INF);
     p.assign(n, -1);
     d[s] = 0;
     using pii = pair<int, int>;
     priority queue<pii, vector<pii>, greater<pii>> g;
     q.push({0, s});
     while (!q.empty()) {
       int v = q.top().second, d_v = q.top().first;
       if (d v != d[v]) continue;
       for (auto edge : adj[v]) {
16
         int to = edge.first, len = edge.second;
         if (d[v] + len < d[to]) {</pre>
             d[to] = d[v] + len;
             p[to] = v;
             q.push({d[to], to});
```

4.5 Dinics

```
1 struct FlowEdge {
      int v, u;
      11 cap, flow = 0;
      FlowEdge(int v, int u, ll cap) : v(v), u(u), cap(
           cap) {}
5
    struct Dinic {
      const 11 flow_inf = 1e18;
     vector<FlowEdge> edges;
     vector<vector<int>> adj;
     int n, m = 0, s, t;
      vector<int> level, ptr;
      gueue<int> q;
      Dinic(int n, int s, int t) : n(n), s(s), t(t) {
        adj.resize(n);
        level.resize(n);
       ptr.resize(n);
      void add_edge(int v, int u, ll cap) {
19
        edges.emplace_back(v, u, cap);
20
        edges.emplace_back(u, v, 0);
        adj[v].push_back(m);
        adj[u].push_back(m + 1);
        m += 2;
      bool bfs() {
        while (!q.empty()) {
          int v = q.front();
          for (int id : adj[v]) {
            if (edges[id].cap - edges[id].flow < 1)</pre>
                 continue;
            if (level[edges[id].u] != -1) continue;
            level[edges[id].u] = level[v] + 1;
            g.push(edges[id].u);
34
        return level[t] != -1;
```

```
11 dfs(int v, 11 pushed) {
        if (pushed == 0) return 0;
        if (v == t) return pushed;
        for (int& cid = ptr[v]; cid < (int)adj[v].size</pre>
             (); cid++) {
          int id = adj[v][cid], u = edges[id].u;
4.3
          if (level[v] + 1 != level[u] || edges[id].cap
                - edges[id].flow < 1) continue;</pre>
          11 tr = dfs(u, min(pushed, edges[id].cap -
               edges[id].flow));
          if (tr == 0) continue;
          edges[id].flow += tr;
          edges[id ^ 1].flow -= tr;
          return tr;
        return 0:
      11 flow() {
        11 f = 0;
        while (true) {
          fill(level.begin(), level.end(), -1);
          level[s] = 0;
          q.push(s);
          if (!bfs()) break;
59
          fill(ptr.begin(), ptr.end(), 0);
60
          while (ll pushed = dfs(s, flow_inf)) f +=
               pushed;
        return f;
63
64 };
```

4.6 Edmonds Karp

```
vector<vector<int>> capacity;
   vector<vector<int>> adj;
   int bfs(int s, int t, vector<int>& parent) {
    fill(parent.begin(), parent.end(), -1);
     parent[s] = -2;
     queue<pair<int, int>> q;
      g.push({s, INF});
      while (!g.emptv()) {
        int cur = q.front().first, flow = q.front().
             second:
        q.pop();
        for (int next : adj[cur]) {
          if (parent[next] == -1 && capacity[cur][next
              ]) {
            parent[next] = cur;
            int new_flow = min(flow, capacity[cur][next
            if (next == t) return new flow;
            q.push({next, new_flow});
19
21
      return 0;
22
    int maxflow(int s, int t) {
24
     int. flow = 0:
     vector<int> parent(n);
      int new flow:
      while (new_flow = bfs(s, t, parent)) {
       flow += new flow;
29
        int cur = t;
        while (cur != s) {
```

```
int prev = parent[cur];
capacity[prev][cur] -= new_flow;
capacity[cur][prev] += new_flow;
cur = prev;
}

return flow;
}
```

```
4.7 Fast Second Mst
    struct edge {
         int s, e, w, id;
        bool operator<(const struct edge& other) {</pre>
             return w < other.w; }</pre>
    typedef struct edge Edge;
    const int N = 2e5 + 5;
    long long res = 0, ans = 1e18;
    int n, m, a, b, w, id, 1 = 21;
    vector<Edge> edges;
    vector<int> h(N, 0), parent(N, -1), size(N, 0),
         present(N, 0);
    vector<vector<pair<int, int>>> adj(N), dp(N, vector
         <pair<int, int>>(1));
   vector<vector<int>> up(N, vector<int>(1, -1));
    pair<int, int> combine(pair<int, int> a, pair<int,</pre>
         int> b) {
      vector<int> v = {a.first, a.second, b.first, b.
14
           second);
      int topTwo = -3, topOne = -2;
16
      for (int c : v) {
17
        if (c > topOne) {
18
          topTwo = topOne;
19
          topOne = c;
20
         } else if (c > topTwo && c < topOne) topTwo = c</pre>
22
      return {topOne, topTwo};
23
24
    void dfs(int u, int par, int d) {
      h[u] = 1 + h[par];
26
      up[u][0] = par;
27
      dp[u][0] = \{d, -1\};
      for (auto v : adj[u]) {
29
        if (v.first != par) dfs(v.first, u, v.second);
31
    pair<int, int> lca(int u, int v) {
      pair<int, int> ans = \{-2, -3\};
34
35
      if (h[u] < h[v]) swap(u, v);
      for (int i = 1 - 1; i >= 0; i--) {
        if (h[u] - h[v] >= (1 << i)) {
37
          ans = combine(ans, dp[u][i]);
38
          u = up[u][i];
39
40
41
      if (u == v) return ans;
42
      for (int i = 1 - 1; i >= 0; i--) {
43
        if (up[u][i] != -1 && up[v][i] != -1 && up[u][i
             ] != up[v][i]) {
          ans = combine(ans, combine(dp[u][i], dp[v][i
          u = up[u][i];
46
          v = up[v][i];
47
48
```

```
ans = combine(ans, combine(dp[u][0], dp[v][0]));
    int main(void) {
      cin >> n >> m;
       for (int i = 1; i <= n; i++) {</pre>
        parent[i] = i;
        size[i] = 1;
 58
59
       for (int i = 1; i <= m; i++) {</pre>
60
         cin >> a >> b >> w; // 1-indexed
         edges.push_back(\{a, b, w, i - 1\});
62
63
       sort(edges.begin(), edges.end());
       for (int i = 0; i \le m - 1; i++) {
        a = edges[i].s;
        b = edges[i].e;
         w = edges[i].w;
         id = edges[i].id;
         if (unite_set(a, b)) {
           adj[a].emplace_back(b, w);
           adj[b].emplace back(a, w);
           present[id] = 1;
           res += w;
       dfs(1, 0, 0);
       for (int i = 1; i \le 1 - 1; i++) {
         for (int j = 1; j \le n; ++j) {
           if (up[j][i - 1] != -1) {
             int v = up[j][i - 1];
             up[j][i] = up[v][i - 1];
             dp[j][i] = combine(dp[j][i-1], dp[v][i-
                  1]);
83
       for (int i = 0; i \le m - 1; i++) {
        id = edges[i].id;
88
         w = edges[i].w;
         if (!present[id]) {
           auto rem = lca(edges[i].s, edges[i].e);
           if (rem.first != w) {
             if (ans > res + w - rem.first) ans = res +
                  w - rem.first;
           } else if (rem.second != -1) {
             if (ans > res + w - rem.second) ans = res +
                   w - rem.second;
96
       cout << ans << "\n";
       return 0;
100 }
```

4.8 Find Cycle

```
1 bool dfs(ll v) {
2    color[v] = 1;
3    for (ll u : adj[v]) {
4        if (color[u] == 0) {
5            parent[u] = v;
6            if (dfs(u)) {
7                return true;
8        }
9        } else if (color[u] == 1) {
```

```
cycle_end = v;
          cycle start = u;
          return true;
      color[v] = 2;
      return false;
17
18 void find_cycle() {
      color.assign(n, 0);
     parent.assign(n, -1);
      cycle_start = -1;
      for (11 v = 0; v < n; v++) {
       if (color[v] == 0 && dfs(v)) {
          break;
      if (cycle_start == -1) {
        cout << "Acyclic" << endl;</pre>
      } else {
        vector<11> cycle;
        cycle.push_back(cycle_start);
        for (11 v = cycle end; v != cycle start; v =
             parent[v]) {
          cycle.push_back(v);
        cycle.push_back(cycle_start);
        reverse(cycle.begin(), cycle.end());
        cout << "Cycle found: ";</pre>
38
        for (ll v : cycle) {
          cout << v << ' ';
        cout << '\n';
```

4.9 Floyd Warshall

4.10 Ford Fulkerson

```
q.pop();
        for (11 v = 0; v < n; v++) {
11
          if (!visited[v] && r_graph[u][v] > 0) {
            if (v == t) {
13
              parent[v] = u;
14
              return true;
15
16
             q.push(v);
17
             parent[v] = u;
18
             visited[v] = true;
19
20
        }
21
22
      return false;
23
    11 ford_fulkerson(ll n, vector<vector<ll>> graph,
         ll s, ll t) {
25
      11 u, v;
26
      vector<vector<ll>>> r_graph;
27
      for (u = 0; u < n; u++)
28
        for (v = 0; v < n; v++)
29
          r_{graph[u][v]} = graph[u][v];
      vector<11> parent;
31
      11 \text{ max\_flow} = 0;
      while (bfs(n, r_graph, s, t, parent)) {
33
        11 path_flow = INF;
34
        for (v = t; v != s; v = parent[v]) {
          u = parent[v];
36
          path_flow = min(path_flow, r_graph[u][v]);
38
        for (v = t; v != s; v = parent[v]) {
39
          u = parent[v];
40
          r_graph[u][v] -= path_flow;
41
          r_graph[v][u] += path_flow;
42
43
        max_flow += path_flow;
44
45
      return max_flow;
46 }
```

4.11 Hierholzer

```
void print circuit(vector<vector<ll>> &adj) {
      map<11, 11> edge_count;
      for (ll i = 0; i < adj.size(); i++) {</pre>
         edge_count[i] = adj[i].size();
      if (!adj.size()) {
        return;
      stack<ll> curr path;
      vector<11> circuit;
11
      curr_path.push(0);
12
13
      11 curr v = 0;
      while (!curr_path.empty()) {
14
        if (edge_count[curr_v]) {
15
           curr_path.push(curr_v);
16
           11 next_v = adj[curr_v].back();
17
           edge_count[curr_v]--;
18
          adj[curr_v].pop_back();
19
          curr_v = next_v;
20
21
         } else {
          circuit.push_back(curr_v);
           curr_v = curr_path.top();
23
           curr_path.pop();
24
```

```
26     for (11 i = circuit.size() - 1; i >= 0; i--) {
27         cout << circuit[i] << ' ';
28     }
29    }</pre>
```

4.12 Hungarian

```
vector<int> u (n+1), v (m+1), p (m+1), way (m+1);
    for (int i=1; i<=n; ++i) {</pre>
      p[0] = i;
      int j0 = 0;
      vector<int> minv (m+1, INF);
      vector<bool> used (m+1, false);
        used[j0] = true;
9
        int i0 = p[j0], delta = INF, j1;
        for (int j=1; j<=m; ++j)</pre>
          if (!used[j]) {
12
            int cur = A[i0][j]-u[i0]-v[j];
            if (cur < minv[j]) minv[j] = cur, way[j] =</pre>
14
            if (minv[j] < delta) delta = minv[j], j1 =</pre>
                  j;
16
         for (int j=0; j<=m; ++j)</pre>
          if (used[j]) u[p[j]] += delta, v[j] -= delta
          else minv[j] -= delta;
         j0 = j1;
      } while (p[j0] != 0);
        int j1 = way[j0];
        p[j0] = p[j1];
         j0 = j1;
      } while (j0);
26
    vector<int> ans (n+1);
    for (int j=1; j<=m; ++j)</pre>
      ans[p[j]] = j;
30 int cost = -v[0];
```

4.13 Is Bipartite

```
bool is_bipartite(vector<ll> &col, vector<vector<ll</pre>
         >> &adj, ll n) {
      queue<pair<ll, ll>> q;
      for (11 i = 0; i < n; i++) {
        if (col[i] == -1) {
          q.push({i, 0});
 6
          col[i] = 0;
          while (!q.empty()) {
            pair<11, 11> p = q.front();
            q.pop();
            11 v = p.first, c = p.second;
            for (ll j : adj[v]) {
              if (col[j] == c) {
                return false;
15
              if (col[j] == -1) {
16
17
                col[j] = (c ? 0 : 1);
                q.push({j, col[j]});
            }
```

```
22 }
23 return true;
24 }
```

4.14 Is Cyclic

```
bool is cyclic util(int u, vector<vector<int>> &adj
        , vector<bool> &vis, vector<bool> &rec) {
     vis[u] = true;
     rec[u] = true;
     for(auto v : adj[u]) {
       if (!vis[v] && is_cyclic_util(v, adj, vis, rec)
             ) return true;
6
       else if (rec[v]) return true;
      rec[u] = false;
      return false;
   bool is_cyclic(int n, vector<vector<int>> &adj) {
      vector<bool> vis(n, false), rec(n, false);
      for (int i = 0; i < n; i++)</pre>
       if (!vis[i] && is_cyclic_util(i, adj, vis, rec)
            ) return true;
     return false;
16 }
```

4.15 Kahn

```
void kahn(vector<vector<11>> &adj) {
      11 n = adj.size();
      vector<ll> in_degree(n, 0);
      for (11 u = 0; u < n; u++)
        for (ll v: adj[u]) in_degree[v]++;
      queue<11> q;
      for (ll i = 0; i < n; i++)</pre>
       if (in_degree[i] == 0)
          q.push(i);
      11 \text{ cnt} = 0;
      vector<ll> top_order;
      while (!q.empty()) {
       11 u = q.front();
       q.pop();
       top_order.push_back(u);
        for (11 v : adj[u])
          if (--in_degree[v] == 0) q.push(v);
18
      if (cnt != n) {
       cout << -1 << '\n';
       return;
      // print top_order
```

4.16 Kosaraju

```
9
```

```
void kosaraju(int n, vector<vector<int>>& adj,
         vector<vector<int>>& sccs) {
      vector<bool> vis(n, false);
      stack<int> stk;
      for (int u = 0; u < n; u++) {</pre>
35
       if (!vis[u]) {
36
          topo_sort(u, adj, vis, stk);
38
39
      vector<vector<int>> adj_t = transpose(n, adj);
40
      for (int u = 0; u < n; u++) {
41
       vis[u] = false;
42
43
      while (!stk.empty()) {
44
        int u = stk.top();
45
        stk.pop();
46
        if (!vis[u]) {
47
          vector<int> scc;
48
          get_scc(u, adj_t, vis, scc);
49
          sccs.push_back(scc);
51
52 }
 4.17 Kruskals
    struct Edge {
      int u, v, weight;
      bool operator<(Edge const& other) {</pre>
         return weight < other.weight;</pre>
    };
    int n;
   vector<Edge> edges;
    int cost = 0;
10 vector<Edge> result;
   DSU dsu = DSU(n);
    sort(edges.begin(), edges.end());
    for (Edge e : edges) {
      if (dsu.find_set(e.u) != dsu.find_set(e.v)) {
```

stk.push(u);

return adj_t;

vis[u] = true;

scc.push_back(u);

if (!vis[v]) {

for (int v : adj_t[u]) {

int>>% adi) {

vector<vector<int>> adj_t(n);

for (int u = 0; u < n; u++) {

adj_t[v].push_back(u);

for (int v : adj[u]) {

11 vector<vector<int>>> transpose(int n, vector<vector<</pre>

void get_scc(int u, vector<vector<int>>& adj_t,

get_scc(v, adj_t, vis, scc);

vector<bool>& vis, vector<int>& scc) {

10

13

14

15

16

17

18

22

23

27

28

29

19 }

```
cost += e.weight;
        result.push back(e);
        dsu.union_sets(e.u, e.v);
18
19 }
                                                         34
 4.18 Kruskal Mst
    struct Edge {
      ll u, v, weight;
      bool operator<(Edge const& other) {</pre>
        return weight < other.weight;</pre>
 6
    };
   11 n;
    vector<Edge> edges;
 9 11 cost = 0;
10 vector<11> tree_id(n);
11 vector<Edge> result;
12 for (11 i = 0; i < n; i++) {
    tree_id[i] = i;
14 }
15 sort(edges.begin(), edges.end());
16 for (Edge e : edges) {
      if (tree_id[e.u] != tree_id[e.v]) {
18
        cost += e.weight;
        result.push_back(e);
20
        ll old_id = tree_id[e.u], new_id = tree_id[e.v
        for (ll i = 0; i < n; i++) {
          if (tree_id[i] == old_id) {
            tree_id[i] = new_id;
 4.19 Kuhn
 1 int n, k;
 2 vector<vector<int>> g;
 3 vector<int> mt;
 4 vector<bool> used;
 5 bool try_kuhn(int v) {
     if (used[v]) return false;
      used[v] = true;
 8
      for (int to : g[v]) {
 9
       if (mt[to] == -1 || try_kuhn(mt[to])) {
          mt[to] = v;
          return true;
12
      return false;
15
    int main() {
      mt.assign(k, -1);
        vector<bool> used1(n, false);
        for (int v = 0; v < n; ++v) {
20
          for (int to : g[v]) {
```

if (mt[to] == -1) {

used1[v] = true;

mt[to] = v;

break;

```
for (int v = 0; v < n; ++v) {
          if (used1[v]) continue;
          used.assign(n, false);
          try_kuhn(v);
        for (int i = 0; i < k; ++i)
          if (mt[i] != -1)
            printf("%d %d\n", mt[i] + 1, i + 1);
36 }
4.20 Lowest Common Ancestor
    struct LCA {
      vector<ll> height, euler, first, segtree;
      vector<bool> visited;
      LCA(vector<vector<ll>> &adj, ll root = 0) {
        n = adj.size();
        height.resize(n);
        first.resize(n);
        euler.reserve(n * 2);
        visited.assign(n, false);
        dfs(adj, root);
        11 m = euler.size();
        segtree.resize(m * 4);
        build(1, 0, m - 1);
      void dfs(vector<vector<11>> &adj, 11 node, 11 h =
        visited[node] = true;
        height[node] = h;
        first[node] = euler.size();
        euler.push_back(node);
        for (auto to : adj[node]) {
          if (!visited[to]) {
            dfs(adj, to, h + 1);
            euler.push_back(node);
      void build(ll node, ll b, ll e) {
        if (b == e) segtree[node] = euler[b];
        else {
          11 \text{ mid} = (b + e) / 2;
          build(node << 1, b, mid);</pre>
          build(node << 1 | 1, mid + 1, e);</pre>
          11 1 = segtree[node << 1], r = segtree[node</pre>
                << 1 | 1];
          segtree[node] = (height[1] < height[r]) ? 1 :</pre>
      11 query(11 node, 11 b, 11 e, 11 L, 11 R) {
        if (b > R | | e < L) return -1;</pre>
        if (b >= L && e <= R) return segtree[node];</pre>
        11 \text{ mid} = (b + e) >> 1;
        11 left = query(node << 1, b, mid, L, R);</pre>
        11 right = query(node << 1 | 1, mid + 1, e, L,</pre>
             R);
        if (left == -1) return right;
        if (right == -1) return left;
        return height[left] < height[right] ? left :</pre>
             right;
48
      ll lca(ll u, ll v) {
        11 left = first[u], right = first[v];
```

4.21 Maximum Bipartite Matching

```
bool bpm(ll n, ll m, vector<vector<bool>> &bpGraph,
          11 u, vector<bool> &seen, vector<11> &matchR)
      for (11 \ v = 0; \ v < m; \ v++)  {
        if (bpGraph[u][v] && !seen[v]) {
          seen[v] = true;
          if (matchR[v] < 0 || bpm(n, m, bpGraph,</pre>
               matchR[v], seen, matchR)) {
             matchR[v] = u;
            return true;
        }
10
11
      return false;
12
13 11 maxBPM(11 n, 11 m, vector<vector<bool>> &bpGraph
      vector<11> matchR(m, -1);
      11 \text{ result} = 0;
      for (11 u = 0; u < n; u++) {
17
        vector<bool> seen(m, false);
18
        if (bpm(n, m, bpGraph, u, seen, matchR)) {
19
          result++;
21
22
      return result;
```

4.22 Min Cost Flow

```
struct Edge {
      int from, to, capacity, cost;
 4 vector<vector<int>> adj, cost, capacity;
    const int INF = 1e9;
    void shortest_paths(int n, int v0, vector<int>& d,
         vector<int>& p) {
      d.assign(n, INF);
      d[v0] = 0;
      vector<bool> inq(n, false);
      queue<int> q;
      q.push(v0);
      p.assign(n, -1);
      while (!q.empty()) {
        int u = q.front();
15
        q.pop();
16
        inq[u] = false;
17
        for (int v : adj[u]) {
18
          if (capacity[u][v] > 0 && d[v] > d[u] + cost[
               u][v]) {
            d[v] = d[u] + cost[u][v];
20
            p[v] = u;
            if (!ing[v]) {
22
              ing[v] = true;
23
              q.push(v);
24
          }
```

```
28
   int min_cost_flow(int N, vector<Edge> edges, int K,
          int s, int t) {
      adj.assign(N, vector<int>());
      cost.assign(N, vector<int>(N, 0));
      capacity.assign(N, vector<int>(N, 0));
      for (Edge e : edges) {
        adj[e.from].push back(e.to);
        adj[e.to].push_back(e.from);
        cost[e.from][e.to] = e.cost;
        cost[e.to][e.from] = -e.cost;
38
        capacity[e.from][e.to] = e.capacity;
39
      int flow = 0;
      int cost = 0;
      vector<int> d, p;
      while (flow < K) {
        shortest_paths(N, s, d, p);
        if (d[t] == INF) break;
        int f = K - flow, cur = t;
        while (cur != s) {
         f = min(f, capacity[p[cur]][cur]);
          cur = p[cur];
        flow += f;
        cost += f * d[t];
        cur = t;
        while (cur != s) {
         capacity[p[cur]][cur] -= f;
          capacity[cur][p[cur]] += f;
          cur = p[cur];
58
59
60
      if (flow < K) return -1;</pre>
      else return cost;
62
```

4.23 Prim

```
1 const int INF = 10000000000;
 2 struct Edge {
      int w = INF, to = -1;
      bool operator<(Edge const& other) const {</pre>
        return make_pair(w, to) < make_pair(other.w,</pre>
             other.to);
 6
    };
8
    int n;
    vector<vector<Edge>> adj;
    void prim() {
      int total_weight = 0;
      vector<Edge> min_e(n);
      \min_{e[0].w = 0;}
      set<Edge> q;
      q.insert({0, 0});
      vector<bool> selected(n, false);
      for (int i = 0; i < n; ++i) {
        if (q.empty()) {
19
          cout << "No MST!" << endl;</pre>
2.0
          exit(0);
        int v = q.begin()->to;
        selected[v] = true;
        total_weight += q.begin()->w;
        q.erase(q.begin());
```

4.24 Topological Sort

```
1 void dfs(11 v) {
     visited[v] = true;
     for (ll u : adj[v]) {
       if (!visited[u]) {
         dfs(u);
8
     ans.push_back(v);
10 void topological_sort() {
     visited.assign(n, false);
     ans.clear();
     for (11 i = 0; i < n; ++i) {
       if (!visited[i]) {
         dfs(i);
18
     reverse(ans.begin(), ans.end());
19 1
```

4.25 Zero One Bfs

```
1  vector<int> d(n, INF);
2  d[s] = 0;
3  deque<int> q;
4  q.push_front(s);
5  while (!q.empty()) {
6   int v = q.front();
7   q.pop_front();
8   for (auto edge : adj[v]) {
9    int u = edge.first, w = edge.second;
10   if (d[v] + w < d[u]) {
11    d[u] = d[v] + w;
12   if (w == 1) q.push_back(u);
13   else q.push_front(u);
14  }
15  }
16 }</pre>
```

5 Miscellaneous

5.1 Gauss

```
11 n = (11) a.size(), m = (11) a[0].size() - 1;
      vector<11> where (m, -1);
      for (11 col = 0, row = 0; col < m && row < n; ++
           col) {
        11 sel = row:
        for (11 i = row; i < n; ++i) {</pre>
          if (abs(a[i][col]) > abs(a[sel][col])) {
11
12
13
        if (abs (a[sel][col]) < EPS) {</pre>
14
          continue;
15
16
         for (ll i = col; i <= m; ++i) {</pre>
17
          swap(a[sel][i], a[row][i]);
18
19
        where[col] = row;
20
         for (11 i = 0; i < n; ++i) {
         if (i != row) {
            double c = a[i][col] / a[row][col];
23
            for (11 j = col; j <= m; ++j) {</pre>
24
              a[i][j] = a[row][j] * c;
25
26
27
28
        ++row;
29
      ans.assign(m, 0);
      for (11 i = 0; i < m; ++i) {
        if (where[i] != -1) {
33
          ans[i] = a[where[i]][m] / a[where[i]][i];
34
35
      for (11 i = 0; i < n; ++i) {
37
        double sum = 0;
38
        for (11 j = 0; j < m; ++j) {
39
         sum += ans[j] * a[i][j];
40
41
        if (abs (sum - a[i][m]) > EPS) {
          return 0:
43
44
45
      for (11 i = 0; i < m; ++i) {
46
        if (where[i] == -1) {
47
          return INF;
48
49
      return 1;
51 }
```

5.2 Ternary Search

```
1  double ternary_search(double 1, double r) {
2   double eps = 1e-9;
3   while (r - 1 > eps) {
4    double m1 = 1 + (r - 1) / 3;
5    double m2 = r - (r - 1) / 3;
6    double f1 = f(m1);
7    double f2 = f(m2);
8    if (f1 < f2) {
9        1 = m1;
10    } else {
1        r = m2;
12    }
13    }
14    return f(1);
15 }</pre>
```

6 Number Theory

6.1 Extended Euclidean

```
1 ll gcd_extended(ll a, ll b, ll &x, ll &y) {
2    if (b == 0) {
3         x = 1;
4         y = 0;
5         return a;
6    }
7    ll x1, y1, g = gcd_extended(b, a % b, x1, y1);
8    x = y1;
9    y = x1 - (a / b) * y1;
10    return g;
11 }
```

6.2 Find All Solutions

```
1 bool find_any_solution(ll a, ll b, ll c, ll &x0, ll
          &v0, 11 &g) {
      g = gcd_extended(abs(a), abs(b), x0, y0);
      if (c % q) {
        return false;
 6
      x0 \star = c / q;
      y0 \star = c / q;
 8
      if (a < 0) {
       x0 = -x0;
      if (b < 0) {
       y0 = -y0;
14
      return true;
15
    void shift_solution(ll & x, ll & y, ll a, ll b, ll
        cnt) {
      x += cnt * b;
18
     y -= cnt * a;
19
    11 find all solutions (11 a, 11 b, 11 c, 11 minx, 11
          maxx, 11 miny, 11 maxy) {
      11 x, y, g;
      if (!find_any_solution(a, b, c, x, y, g)) {
23
        return 0:
2.4
25
      a /= g;
      b /= q;
      11 \text{ sign}_a = a > 0 ? +1 : -1;
      11 \text{ sign } b = b > 0 ? +1 : -1;
      shift_solution(x, y, a, b, (minx - x) / b);
      if (x < minx) {</pre>
        shift_solution(x, y, a, b, sign_b);
33
      if (x > maxx) {
34
        return 0;
3.5
36
      11 \ 1x1 = x;
      shift_solution(x, y, a, b, (maxx - x) / b);
      if (x > maxx) {
        shift_solution(x, y, a, b, -sign_b);
39
40
      shift_solution(x, y, a, b, -(miny - y) / a);
      if (y < miny) {
```

```
shift_solution(x, y, a, b, -sign_a);
      if (y > maxy) {
        return 0;
48
      11 \ 1x2 = x;
      shift_solution(x, y, a, b, -(maxy - y) / a);
      if (y > maxy) {
        shift_solution(x, y, a, b, sign_a);
      11 \text{ rx2} = x;
      if (1x2 > rx2) {
        swap(1x2, rx2);
58
      11 1x = max(1x1, 1x2), rx = min(rx1, rx2);
      if (1x > rx) {
60
       return 0;
61
      return (rx - 1x) / abs(b) + 1;
63
```

6.3 Linear Sieve

6.4 Miller Rabin

```
1 bool check_composite(u64 n, u64 a, u64 d, 11 s) {
     u64 x = binpower(a, d, n);
      if (x == 1 | | x == n - 1) {
       return false;
      for (11 r = 1; r < s; r++) {
      x = (u128) x * x % n;
       if (x == n - 1) {
          return false:
      return true;
   bool miller_rabin(u64 n) {
     if (n < 2) {
       return false:
18
     11 r = 0:
      u64 d = n - 1:
      while ((d & 1) == 0) {
       d >>= 1;
22
        r++;
```

```
Pegaraw
```

6.5 Modulo Inverse

```
1  11 mod_inv(11 a, 11 m) {
       if (m == 1) {
         return 0;
      11 \text{ m0} = \text{m}, \text{ x} = 1, \text{ y} = 0;
      while (a > 1) {
        11 q = a / m, t = m;
       m = a % m;
        a = t;
        t = y;
        y = x - q * y;
        x = t;
      if (x < 0) {
        x += m0;
16
17
       return x;
18 }
```

6.6 Pollard Rho Brent

```
1 11 mult(11 a, 11 b, 11 mod) {
       return (__int128_t) a * b % mod;
 3
 4 11 f(11 x, 11 c, 11 mod) {
      return (mult(x, x, mod) + c) % mod;
 6
    ll pollard_rho_brent(ll n, ll x0 = 2, ll c = 1) {
      11 \times = \times 0, q = 1, q = 1, \times s, y, m = 128, 1 = 1;
      while (g == 1) {
        y = x;
         for (ll i = 1; i < 1; i++) {</pre>
          x = f(x, c, n);
13
         11 k = 0:
         while (k < 1 \&\& q == 1) {
          xs = x;
           for (ll i = 0; i < m && i < 1 - k; i++) {
            x = f(x, c, n);
19
             q = mult(q, abs(y - x), n);
20
           g = \underline{gcd}(q, n);
22
           k += m;
24
25
26
         1 *= 2;
       if (q == n) {
27
         do {
28
           xs = f(xs, c, n);
29
           g = \underline{gcd}(abs(xs - y), n);
         } while (g == 1);
```

```
return q;
 6.7 Range Sieve
    vector<bool> range_sieve(ll 1, ll r) {
      11 n = sqrt(r);
      vector<bool> is_prime(n + 1, true);
     vector<ll> prime;
     is_prime[0] = is_prime[1] = false;
      prime.push_back(2);
      for (11 i = 4; i <= n; i += 2) {
       is_prime[i] = false;
9
10
      for (11 i = 3; i <= n; i += 2) {
       if (is_prime[i]) {
          prime.push_back(i);
          for (11 j = i * i; j <= n; j += i) {
           is_prime[j] = false;
          }
16
       }
18
      vector<bool> result(r - 1 + 1, true);
```

for (11 j = max(i * i, (1 + i - 1) / i * i); j

6.8 Segmented Sieve

if (1 == 1) {

return result;

23

26

for (ll i : prime) {

result[0] = false;

<= r; j += i) {

result[j - l] = false;

```
vector<ll> segmented_sieve(ll n) {
     const 11 S = 10000;
     11 nsgrt = sgrt(n);
     vector<char> is_prime(nsqrt + 1, true);
     vector<1l> prime;
     is_prime[0] = is_prime[1] = false;
     prime.push_back(2);
8
     for (11 i = 4; i <= nsqrt; i += 2) {</pre>
9
      is_prime[i] = false;
     for (11 i = 3; i <= nsqrt; i += 2) {</pre>
      if (is_prime[i]) {
         prime.push_back(i);
         for (ll j = i * i; j \le nsqrt; j += i) {
           is_prime[j] = false;
     vector<ll> result;
     vector<char> block(S);
     for (11 k = 0; k * S \le n; k++) {
      fill(block.begin(), block.end(), true);
       for (ll p : prime) {
         for (11 j = max((k * S + p - 1) / p, p) * p -
               k * S; j < S; j += p) {
           block[j] = false;
```

```
27     }
28     if (k == 0) {
29         block[0] = block[1] = false;
30     }
31     for (ll i = 0; i < S && k * S + i <= n; i++) {
32         if (block[i]) {
33             result.push_back(k * S + i);
34         }
35     }
36     }
37     return result;
38 }</pre>
```

6.9 Tonelli Shanks

```
1  ll legendre(ll a, ll p) {
      return bin_pow_mod(a, (p - 1) / 2, p);
    11 tonelli_shanks(ll n, ll p) {
      if (legendre(n, p) == p - 1) {
        return -1;
      if (p % 4 == 3) {
        return bin_pow_mod(n, (p + 1) / 4, p);
      11 Q = p - 1, S = 0;
      while (Q \% 2 == 0) {
        0 /= 2;
        S++:
      11 z = 2;
      for (; z < p; z++) {</pre>
        if (legendre(z, p) == p - 1) {
19
          break;
20
21
      11 M = S, c = bin_pow_mod(z, Q, p), t =
           bin_pow_mod(n, Q, p), R = bin_pow_mod(n, Q)
           + 1) / 2, p);
      while (t % p != 1) {
        if (t % p == 0) {
          return 0;
        11 i = 1, t2 = t * t % p;
        for (; i < M; i++) {</pre>
          if (t2 % p == 1) {
            break;
          t2 = t2 * t2 % p;
        11 b = bin_pow_mod(c, bin_pow_mod(2, M - i - 1,
              p), p);
        M = i;
        c = b * b % p;
        t = t * c % p;
38
        R = R * b % p;
40
      return R;
41
```

7 Strings

7.1 Count Unique Substrings

```
int count_unique_substrings(string const& s) {
       int n = s.size();
       const int p = 31;
       const int m = 1e9 + 9;
       vector<long long> p_pow(n);
       p_pow[0] = 1;
       for (int i = 1; i < n; i++) p_pow[i] = (p_pow[i -</pre>
             1] * p) % m;
       vector<long long> h(n + 1, 0);
       for (int i = 0; i < n; i++) h[i + 1] = (h[i] + (s)
            [i] - 'a' + 1) * p_pow[i]) % m;
       int cnt = 0:
11
       for (int 1 = 1; 1 <= n; 1++) {
12
        unordered_set<long long> hs;
13
         for (int i = 0; i \le n - 1; i++) {
14
           long long cur_h = (h[i + 1] + m - h[i]) % m;
15
           \operatorname{cur}_h = (\operatorname{cur}_h * \operatorname{p_pow}[n - i - 1]) % m;
16
           hs.insert(cur_h);
17
18
        cnt += hs.size();
19
20
       return cnt;
21
```

7.2 Finding Repetitions

```
vector<int> z_function(string const& s) {
      int n = s.size();
      vector<int> z(n);
      for (int i = 1, l = 0, r = 0; i < n; i++) {
        if (i \le r) z[i] = min(r - i + 1, z[i - 1]);
        while (i + z[i] < n \&\& s[z[i]] == s[i + z[i]])
             z[i]++;
        if (i + z[i] - 1 > r) {
          1 = i;
          r = i + z[i] - 1;
10
12
      return z;
13
    int get_z(vector<int> const& z, int i) {
      if (0 <= i && i < (int) z.size()) return z[i];</pre>
16
      else return 0;
17 }
18 vector<pair<int, int>> repetitions;
    void convert_to_repetitions(int shift, bool left,
         int cntr, int 1, int k1, int k2) {
      for (int 11 = \max(1, 1 - k2); 11 \le \min(1, k1);
           11++) {
        if (left && l1 == 1) break;
22
        int 12 = 1 - 11;
        int pos = shift + (left ? cntr - 11 : cntr - 1
             -11+1):
        repetitions.emplace_back(pos, pos + 2 * 1 - 1);
25
26
27
    void find_repetitions(string s, int shift = 0) {
28
      int n = s.size();
29
      if (n == 1) return;
     int nu = n / 2;
31
      int nv = n - nu;
32
      string u = s.substr(0, nu);
     string v = s.substr(nu);
      string ru(u.rbegin(), u.rend());
35
      string rv(v.rbegin(), v.rend());
      find_repetitions(u, shift);
      find_repetitions(v, shift + nu);
```

```
vector<int> z1 = z_function(ru);
      vector<int> z2 = z_function(v + '#' + u);
40
      vector<int> z3 = z_function(ru + '#' + rv);
      vector<int> z4 = z_function(v);
      for (int cntr = 0; cntr < n; cntr++) {</pre>
        int 1, k1, k2;
        if (cntr < nu) {</pre>
         1 = nu - cntr;
          k1 = get_z(z1, nu - cntr);
          k2 = get_z(z2, nv + 1 + cntr);
        } else {
          1 = cntr - nu + 1;
          k1 = get_z(z3, nu + 1 + nv - 1 - (cntr - nu))
          k2 = get_z(z4, (cntr - nu) + 1);
        if (k1 + k2 >= 1) convert_to_repetitions(shift,
              cntr < nu, cntr, 1, k1, k2);</pre>
54
55 }
```

7.3 Group Identical Substrings

```
vector<vector<int>>> group_identical_strings(vector
string> const& s) {

int n = s.size();

vector<pair<long long, int>> hashes(n);

for (int i = 0; i < n; i++) hashes[i] = {
    compute_hash(s[i]), i};

sort(hashes.begin(), hashes.end());

vector<vector<int>>> groups;

for (int i = 0; i < n; i++) {

if (i == 0 || hashes[i].first != hashes[i - 1].
    first) groups.emplace_back();

groups.back().push_back(hashes[i].second);

return groups;

return groups;

}
</pre>
```

7.4 Hashing

7.5 Knuth Morris Pratt

```
vector<ll> prefix_function(string s) {
    11 n = (11) s.length();
    vector<ll> pi(n);

for (11 i = 1; i < n; i++) {
    11 j = pi[i - 1];
    while (j > 0 && s[i] != s[j]) j = pi[j - 1];
    if (s[i] == s[j]) j++;
    pi[i] = j;
```

```
9  }
10  return pi;
11  }
12  // count occurences
13  vector<int> ans(n + 1);
14  for (int i = 0; i < n; i++)
15   ans[pi[i]]++;
16  for (int i = n-1; i > 0; i--)
17  ans[pi[i-1]] += ans[i];
18  for (int i = 0; i <= n; i++)
19  ans[i]++;</pre>
```

7.6 Longest Common Prefix

```
vector<int> lcp_construction(string const& s,
         vector<int> const& p) {
      int n = s.size();
      vector<int> rank(n, 0);
      for (int i = 0; i < n; i++) rank[p[i]] = i;</pre>
      int k = 0;
      vector<int> lcp(n-1, 0);
      for (int i = 0; i < n; i++) {
        if (rank[i] == n - 1) {
          k = 0:
          continue;
        int j = p[rank[i] + 1];
        while (i + k < n \&\& j + k < n \&\& s[i + k] == s[
             \frac{1}{1} + k]) k++;
        lcp[rank[i]] = k;
        if (k) k--;
      return lcp;
18 }
```

7.7 Manacher

```
vector<int> manacher_odd(string s) {
      int n = s.size();
      s = "$" + s + "^";
      vector<int> p(n + 2);
      int 1 = 1, r = 1;
      for (int i = 1; i \le n; i++) {
        p[i] = max(0, min(r - i, p[l + (r - i)]));
        while (s[i - p[i]] == s[i + p[i]]) p[i]++;
        if(i + p[i] > r) 1 = i - p[i], r = i + p[i];
11
      return vector<int>(begin(p) + 1, end(p) - 1);
12
   vector<int> manacher(string s) {
14
      string t;
      for(auto c: s) t += string("#") + c;
      auto res = manacher_odd(t + "#");
      return vector<int>(begin(res) + 1, end(res) - 1);
18 }
```

7.8 Rabin Karp

```
vector<11> p_pow(max(S, T));
      p pow[0] = 1;
      for (ll i = 1; i < (ll) p_pow.size(); i++) p_pow[</pre>
           i] = (p_pow[i-1] * p) % m;
      vector<11> h(T + 1, 0);
      for (ll i = 0; i < T; i++) h[i + 1] = (h[i] + (t[
           i] - 'a' + 1) * p_pow[i]) % m;
      11 h_s = 0;
      for (11 i = 0; i < S; i++) h_s = (h_s + (s[i] - '
           a' + 1) * p pow[i]) % m;
11
      vector<11> occurences;
      for (11 i = 0; i + S - 1 < T; i++) {
13
        11 \text{ cur}_h = (h[i + S] + m - h[i]) % m;
14
        if (cur_h == h_s * p_pow[i] % m) occurences.
             push back(i);
15
16
      return occurences;
17 }
```

7.9 Suffix Array

```
if (s[p[i]] != s[p[i-1]]) classes++;
    c[p[i]] = classes - 1;
  vector<int> pn(n), cn(n);
  for (int h = 0; (1 << h) < n; ++h) {
    for (int i = 0; i < n; i++) {
      pn[i] = p[i] - (1 << h);
      if (pn[i] < 0)
        pn[i] += n;
    fill(cnt.begin(), cnt.begin() + classes, 0);
    for (int i = 0; i < n; i++) cnt[c[pn[i]]]++;</pre>
    for (int i = 1; i < classes; i++) cnt[i] += cnt</pre>
         [i - 1];
    for (int i = n-1; i >= 0; i--) p[--cnt[c[pn[i
         ]]]] = pn[i];
    cn[p[0]] = 0;
    classes = 1;
    for (int i = 1; i < n; i++) {
      pair<int, int> cur = {c[p[i]], c[(p[i] + (1)
           << h)) % n]};
      pair<int, int> prev = \{c[p[i-1]], c[(p[i-1]] +
            (1 << h)) % n];
      if (cur != prev) ++classes;
      cn[p[i]] = classes - 1;
    c.swap(cn);
  return p;
vector<int> build_suff_arr(string s) {
  s += "$";
  vector<int> sorted_shifts = sort_cyclic_shifts(s)
       ;
```

24

25

26

28

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7.10 Z Function

Pegaraw Pegaraw

4() 0(())	100 7							
f(n) = O(g(n))	iff \exists positive c, n_0 such that $0 \le f(n) \le cg(n) \ \forall n \ge n_0$.	$\sum_{i=1}^{n} i = \frac{n(n+1)}{2}, \sum_{i=1}^{n} i^2 = \frac{n(n+1)(2n+1)}{6}, \sum_{i=1}^{n} i^3 = \frac{n^2(n+1)^2}{4}.$						
$f(n) = \Omega(g(n))$	iff \exists positive c, n_0 such that $f(n) \ge cg(n) \ge 0 \ \forall n \ge n_0$.	In general: $ \begin{array}{cccc} $						
$f(n) = \Theta(g(n))$	iff $f(n) = O(g(n))$ and $f(n) = \Omega(g(n))$.	$\sum_{i=1}^{n} i^{m} = \frac{1}{m+1} \left[(n+1)^{m+1} - 1 - \sum_{i=1}^{n} \left((i+1)^{m+1} - i^{m+1} - (m+1)i^{m} \right) \right]$						
f(n) = o(g(n))	iff $\lim_{n\to\infty} f(n)/g(n) = 0$.	$\sum_{i=1}^{n-1} i^m = \frac{1}{m+1} \sum_{k=0}^m {m+1 \choose k} B_k n^{m+1-k}.$						
$\lim_{n \to \infty} a_n = a$	iff $\forall \epsilon > 0$, $\exists n_0$ such that $ a_n - a < \epsilon$, $\forall n \ge n_0$.	Geometric series:						
$\sup S$	least $b \in \mathbb{R}$ such that $b \ge s$, $\forall s \in S$.	$\sum_{i=0}^{n} c^{i} = \frac{c^{n+1} - 1}{c - 1}, c \neq 1, \sum_{i=0}^{\infty} c^{i} = \frac{1}{1 - c}, \sum_{i=1}^{\infty} c^{i} = \frac{c}{1 - c}, c < 1,$						
$\inf S$	greatest $b \in \mathbb{R}$ such that $b \le s$, $\forall s \in S$.	$\sum_{i=0}^{n} ic^{i} = \frac{nc^{n+2} - (n+1)c^{n+1} + c}{(c-1)^{2}}, c \neq 1, \sum_{i=0}^{\infty} ic^{i} = \frac{c}{(1-c)^{2}}, c < 1.$						
$ \liminf_{n \to \infty} a_n $	$\lim_{n \to \infty} \inf \{ a_i \mid i \ge n, i \in \mathbb{N} \}.$	Harmonic series: $n + n = n + n = n = n = n = n = n = n = $						
$\limsup_{n \to \infty} a_n$	$\lim_{n \to \infty} \sup \{ a_i \mid i \ge n, i \in \mathbb{N} \}.$	$H_n = \sum_{i=1}^n \frac{1}{i}, \qquad \sum_{i=1}^n iH_i = \frac{n(n+1)}{2}H_n - \frac{n(n-1)}{4}.$						
$\binom{n}{k}$	Combinations: Size k subsets of a size n set.	$\sum_{i=1}^{n} H_i = (n+1)H_n - n, \sum_{i=1}^{n} {i \choose m} H_i = {n+1 \choose m+1} \left(H_{n+1} - \frac{1}{m+1} \right).$						
$\begin{bmatrix} n \\ k \end{bmatrix}$	Stirling numbers (1st kind): Arrangements of an n element set into k cycles.	1. $\binom{n}{k} = \frac{n!}{(n-k)!k!}$, 2. $\sum_{k=0}^{n} \binom{n}{k} = 2^{n}$, 3. $\binom{n}{k} = \binom{n}{n-k}$,						
$\left\{ egin{array}{c} n \\ k \end{array} \right\}$	Stirling numbers (2nd kind): Partitions of an <i>n</i> element	4. $\binom{n}{k} = \frac{n}{k} \binom{n-1}{k-1}$, $5.$ $\binom{n}{k} = \binom{n-1}{k} + \binom{n-1}{k-1}$,						
	set into k non-empty sets.	6. $\binom{n}{m}\binom{m}{k} = \binom{n}{k}\binom{n-k}{m-k}$, 7. $\sum_{k=0}^{n} \binom{r+k}{k} = \binom{r+n+1}{n}$,						
$\left\langle {n\atop k}\right\rangle$	1st order Eulerian numbers: Permutations $\pi_1\pi_2\pi_n$ on $\{1,2,,n\}$ with k ascents.	$8. \ \sum_{k=0}^{n} \binom{k}{m} = \binom{n+1}{m+1}, \qquad \qquad 9. \ \sum_{k=0}^{n} \binom{r}{k} \binom{s}{n-k} = \binom{r+s}{n},$						
$\left\langle\!\!\left\langle {n\atop k}\right\rangle\!\!\right\rangle$	2nd order Eulerian numbers.	$10. \begin{pmatrix} n \\ k \end{pmatrix} = (-1)^k \binom{k-n-1}{k}, \qquad 11. \begin{Bmatrix} n \\ 1 \end{Bmatrix} = \begin{Bmatrix} n \\ n \end{Bmatrix} = 1,$						
C_n	Catalan Numbers: Binary trees with $n+1$ vertices.	12. $\binom{n}{2} = 2^{n-1} - 1$, 13. $\binom{n}{k} = k \binom{n-1}{k} + \binom{n-1}{k-1}$,						
14. $\begin{bmatrix} n \\ 1 \end{bmatrix} = (n-1)$)!, $ 15. \begin{bmatrix} n \\ 2 \end{bmatrix} = (n - 1) $	$16. \begin{bmatrix} n \\ n \end{bmatrix} = 1, \qquad \qquad 17. \begin{bmatrix} n \\ k \end{bmatrix} \ge \begin{Bmatrix} n \\ k \end{Bmatrix},$						
18. $\begin{bmatrix} n \\ k \end{bmatrix} = (n-1)$	$\binom{n-1}{k} + \binom{n-1}{k-1}, 19. \ \binom{n}{n-1}$							
$22. \left\langle {n \atop 0} \right\rangle = \left\langle {n \atop n-1} \right\rangle$	$ 22. \ \left\langle {n \atop 0} \right\rangle = \left\langle {n \atop n-1} \right\rangle = 1, $ $ 23. \ \left\langle {n \atop k} \right\rangle = \left\langle {n \atop n-1-k} \right\rangle, $ $ 24. \ \left\langle {n \atop k} \right\rangle = (k+1) \left\langle {n-1 \atop k} \right\rangle + (n-k) \left\langle {n-1 \atop k-1} \right\rangle, $							
$25. \ \left\langle \begin{array}{c} 0 \\ k \end{array} \right\rangle = \left\{ \begin{array}{c} 1 & \text{if } k = 0, \\ 0 & \text{otherwise} \end{array} \right. $ $26. \ \left\langle \begin{array}{c} n \\ 1 \end{array} \right\rangle = 2^n - n - 1, $ $27. \ \left\langle \begin{array}{c} n \\ 2 \end{array} \right\rangle = 3^n - (n+1)2^n + \left(\begin{array}{c} n+1 \\ 2 \end{array} \right)$								
$25. \ \left\langle {0 \atop k} \right\rangle = \left\{ {1 \atop 0 \text{ otherwise}} \right. $ $26. \ \left\langle {n \atop 1} \right\rangle = 2^n - n - 1, $ $27. \ \left\langle {n \atop 2} \right\rangle = 3^n - (n+1)2^n + $ $28. \ x^n = \sum_{k=0}^n \left\langle {n \atop k} \right\rangle {x+k \choose n}, $ $29. \ \left\langle {n \atop m} \right\rangle = \sum_{k=0}^m {n+1 \choose k} (m+1-k)^n (-1)^k, $ $30. \ m! \left\{ {n \atop m} \right\} = \sum_{k=0}^n \left\langle {n \atop k} \right\rangle {x+k \choose n} $								
		32. $\left\langle \left\langle {n\atop 0} \right\rangle \right\rangle = 1,$ 33. $\left\langle \left\langle {n\atop n} \right\rangle \right\rangle = 0$ for $n \neq 0,$						
$34. \left\langle\!\!\left\langle {n\atop k} \right\rangle\!\!\right\rangle = (k + 1)^n$	(-1) $\binom{n-1}{k}$ $+ (2n-1-k)$ $\binom{n-1}{k}$							
$36. \left\{ \begin{array}{c} x \\ x-n \end{array} \right\} = \frac{1}{2}$	$\sum_{k=0}^{n} \left\langle \!\! \left\langle n \right\rangle \!\! \right\rangle \left(x + n - 1 - k \right), $	37. ${n+1 \choose m+1} = \sum_{k} {n \choose k} {k \choose m} = \sum_{k=0}^{n} {k \choose m} (m+1)^{n-k},$						

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The Chinese remainder theorem: There exists a number C such that:

 $C \equiv r_1 \mod m_1$

: : :

 $C \equiv r_n \mod m_n$

if m_i and m_j are relatively prime for $i \neq j$. Euler's function: $\phi(x)$ is the number of positive integers less than x relatively prime to x. If $\prod_{i=1}^{n} p_i^{e_i}$ is the prime factorization of x then

$$\phi(x) = \prod_{i=1}^{n} p_i^{e_i - 1} (p_i - 1).$$

Euler's theorem: If a and b are relatively prime then

$$1 \equiv a^{\phi(b)} \bmod b.$$

Fermat's theorem:

$$1 \equiv a^{p-1} \bmod p.$$

The Euclidean algorithm: if a > b are integers then

$$gcd(a, b) = gcd(a \mod b, b).$$

If $\prod_{i=1}^{n} p_i^{e_i}$ is the prime factorization of x

$$S(x) = \sum_{d|x} d = \prod_{i=1}^{n} \frac{p_i^{e_i+1} - 1}{p_i - 1}.$$

Perfect Numbers: x is an even perfect number iff $x = 2^{n-1}(2^n - 1)$ and $2^n - 1$ is prime. Wilson's theorem: n is a prime iff

$$(n-1)! \equiv -1 \bmod n.$$

$$\mu(i) = \begin{cases} (n-1)! = -1 \bmod n. \\ \text{M\"obius inversion:} \\ \mu(i) = \begin{cases} 1 & \text{if } i = 1. \\ 0 & \text{if } i \text{ is not square-free.} \\ (-1)^r & \text{if } i \text{ is the product of} \\ r & \text{distinct primes.} \end{cases}$$
 If

 If

$$G(a) = \sum_{d|a} F(d),$$

$$F(a) = \sum_{d|a} \mu(d) G\left(\frac{a}{d}\right).$$

Prime numbers:

$$p_n = n \ln n + n \ln \ln n - n + n \frac{\ln \ln n}{\ln n}$$

$$+O\left(\frac{n}{\ln n}\right),$$

$$\pi(n) = \frac{n}{\ln n} + \frac{n}{(\ln n)^2} + \frac{2!n}{(\ln n)^3} + O\left(\frac{n}{(\ln n)^4}\right).$$

)ef			

Loop An edge connecting a vertex to itself.

Directed Each edge has a direction. SimpleGraph with no loops or multi-edges.

WalkA sequence $v_0e_1v_1\ldots e_\ell v_\ell$. TrailA walk with distinct edges. Pathtrail with distinct

vertices.

ConnectedA graph where there exists a path between any two

vertices.

ComponentΑ maximal connected

subgraph.

TreeA connected acyclic graph. Free tree A tree with no root. DAGDirected acyclic graph. EulerianGraph with a trail visiting each edge exactly once.

Hamiltonian Graph with a cycle visiting each vertex exactly once.

CutA set of edges whose removal increases the number of components.

Cut-setA minimal cut. Cut edge A size 1 cut.

k-Connected A graph connected with the removal of any k-1vertices.

k-Tough $\forall S \subseteq V, S \neq \emptyset$ we have $k \cdot c(G - S) \le |S|$.

A graph where all vertices k-Regular have degree k.

k-Factor Α k-regular spanning subgraph.

Matching A set of edges, no two of which are adjacent.

CliqueA set of vertices, all of which are adjacent.

Ind. set A set of vertices, none of which are adjacent.

Vertex cover A set of vertices which cover all edges.

Planar graph A graph which can be embeded in the plane.

Plane graph An embedding of a planar

$$\sum_{v \in V} \deg(v) = 2m.$$

If G is planar then n - m + f = 2, so

$$f \le 2n - 4, \quad m \le 3n - 6.$$

Any planar graph has a vertex with degree ≤ 5 .

Notation:

E(G)Edge set Vertex set V(G)

c(G)Number of components

G[S]Induced subgraph deg(v)Degree of v

Maximum degree $\Delta(G)$

 $\delta(G)$ Minimum degree $\chi(G)$ Chromatic number

 $\chi_E(G)$ Edge chromatic number G^c Complement graph K_n Complete graph

 K_{n_1,n_2} Complete bipartite graph

Ramsev number

Geometry

Projective coordinates: (x, y, z), not all x, y and z zero.

$$(x, y, z) = (cx, cy, cz) \quad \forall c \neq 0.$$

Cartesian Projective (x, y)(x, y, 1)

y = mx + b(m, -1, b)x = c(1,0,-c)

Distance formula, L_p and L_{∞}

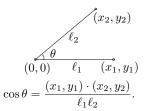
$$\sqrt{(x_1 - x_0)^2 + (y_1 - y_0)^2},$$
$$[|x_1 - x_0|^p + |y_1 - y_0|^p]^{1/p},$$

$$\lim_{x \to \infty} \left[|x_1 - x_0|^p + |y_1 - y_0|^p \right]^{1/p}.$$

Area of triangle $(x_0, y_0), (x_1, y_1)$ and (x_2, y_2) :

$$\frac{1}{2} \operatorname{abs} \begin{vmatrix} x_1 - x_0 & y_1 - y_0 \\ x_2 - x_0 & y_2 - y_0 \end{vmatrix}.$$

Angle formed by three points:



Line through two points (x_0, y_0) and (x_1, y_1) :

$$\begin{vmatrix} x & y & 1 \\ x_0 & y_0 & 1 \\ x_1 & y_1 & 1 \end{vmatrix} = 0.$$

Area of circle, volume of sphere:

$$A = \pi r^2, \qquad V = \frac{4}{3}\pi r^3.$$

If I have seen farther than others, it is because I have stood on the shoulders of giants.

- Issac Newton

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Taylor's series:

$$f(x) = f(a) + (x - a)f'(a) + \frac{(x - a)^2}{2}f''(a) + \dots = \sum_{i=0}^{\infty} \frac{(x - a)^i}{i!}f^{(i)}(a).$$

Expansions:

Expansions:
$$\frac{1}{1-x} = 1 + x + x^2 + x^3 + x^4 + \cdots = \sum_{i=0}^{\infty} x^i,$$

$$\frac{1}{1-cx} = 1 + x + x^2 + x^3 + x^4 + \cdots = \sum_{i=0}^{\infty} c^i x^i,$$

$$\frac{1}{1-cx} = 1 + x + x^2 + x^3 + x^4 + \cdots = \sum_{i=0}^{\infty} c^i x^i,$$

$$\frac{1}{1-x^n} = 1 + x^n + x^{2n} + x^{3n} + \cdots = \sum_{i=0}^{\infty} ix^{ii},$$

$$\frac{x}{(1-x)^2} = x + 2x^2 + 3x^3 + 4x^4 + \cdots = \sum_{i=0}^{\infty} ix^i,$$

$$x^k \frac{d^n}{dx^n} \left(\frac{1}{1-x}\right) = x + 2^{n}x^2 + 3^n x^3 + 4^n x^4 + \cdots = \sum_{i=0}^{\infty} i^n x^i,$$

$$e^x = 1 + x + \frac{1}{2}x^2 + \frac{1}{6}x^3 + \cdots = \sum_{i=0}^{\infty} i^n x^i,$$

$$\ln(1+x) = x - \frac{1}{2}x^2 + \frac{1}{3}x^3 - \frac{1}{4}x^4 - \cdots = \sum_{i=0}^{\infty} (-1)^{i+1} \frac{x^i}{i},$$

$$\ln \frac{1}{1-x} = x + \frac{1}{2}x^2 + \frac{1}{3}x^3 + \frac{1}{4}x^4 + \cdots = \sum_{i=0}^{\infty} (-1)^{i} \frac{x^{2i+1}}{(2i+1)!},$$

$$\cos x = 1 - \frac{1}{2!}x^2 + \frac{1}{1!}x^4 - \frac{1}{6!}x^6 + \cdots = \sum_{i=0}^{\infty} (-1)^{i} \frac{x^{2i+1}}{(2i+1)!},$$

$$(1+x)^n = 1 + nx + \frac{n(n-1)}{2}x^2 + \cdots = \sum_{i=0}^{\infty} (-1)^{i} \frac{x^{2i+1}}{(2i+1)!},$$

$$(1+x)^n = 1 + nx + \frac{n(n-1)}{2}x^2 + \cdots = \sum_{i=0}^{\infty} (-1)^{i} \frac{x^{2i+1}}{(2i+1)!},$$

$$\frac{1}{(1-x)^{n+1}} = 1 + (n+1)x + (\frac{n+2}{2})x^2 + \cdots = \sum_{i=0}^{\infty} (-1)^{i} \frac{x^{2i+1}}{(i)},$$

$$\frac{x}{e^x - 1} = 1 - \frac{1}{2}x + \frac{1}{12}x^2 - \frac{1}{126}x^4 + \cdots = \sum_{i=0}^{\infty} (\frac{i+n}{i})x^i,$$

$$\frac{1}{\sqrt{1-4x}} = 1 + x + 2x^2 + 5x^3 + \cdots = \sum_{i=0}^{\infty} (\frac{2i}{i})x^i,$$

$$\frac{1}{\sqrt{1-4x}} = 1 + x + 2x^2 + 6x^3 + \cdots = \sum_{i=0}^{\infty} (\frac{2i}{i})x^i,$$

$$\frac{1}{1-x} \ln \frac{1}{1-x} = x + \frac{3}{2}x^2 + \frac{1}{16}x^3 + \frac{25}{24}x^4 + \cdots = \sum_{i=0}^{\infty} H_i x^i,$$

$$\frac{1}{1-x} \ln \frac{1}{1-x} = x + \frac{3}{2}x^2 + \frac{1}{16}x^3 + \frac{25}{24}x^4 + \cdots = \sum_{i=0}^{\infty} H_i x^i,$$

$$\frac{1}{2} \left(\ln \frac{1}{1-x} \right)^2 = \frac{1}{2}x^2 + \frac{3}{4}x^3 + \frac{11}{24}x^4 + \cdots = \sum_{i=0}^{\infty} F_{ii}x^i.$$

$$\frac{x}{1-x-x^2} = x + x^2 + 2x^3 + 3x^4 + \cdots = \sum_{i=0}^{\infty} F_{ii}x^i.$$

Ordinary power series:

$$A(x) = \sum_{i=0}^{\infty} a_i x^i.$$

Exponential power series:

$$A(x) = \sum_{i=0}^{\infty} a_i \frac{x^i}{i!}.$$

Dirichlet power series:

$$A(x) = \sum_{i=1}^{\infty} \frac{a_i}{i^x}.$$

Binomial theorem

$$(x+y)^n = \sum_{k=0}^n \binom{n}{k} x^{n-k} y^k.$$

$$x^{n} - y^{n} = (x - y) \sum_{k=0}^{n-1} x^{n-1-k} y^{k}.$$

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$$\alpha A(x) + \beta B(x) = \sum_{i=0}^{\infty} (\alpha a_i + \beta b_i) x^i,$$

$$x^k A(x) = \sum_{i=k}^{\infty} a_{i-k} x^i,$$

$$\frac{A(x) - \sum_{i=0}^{k-1} a_i x^i}{x^k} = \sum_{i=0}^{\infty} a_{i+k} x^i,$$

$$A(cx) = \sum_{i=0}^{\infty} c^i a_i x^i,$$

$$A'(x) = \sum_{i=0}^{\infty} (i+1)a_{i+1}x^{i},$$

$$xA'(x) = \sum_{i=1}^{\infty} ia_i x^i,$$

$$\int A(x) dx = \sum_{i=1}^{\infty} \frac{a_{i-1}}{i} x^{i},$$

$$\frac{A(x) + A(-x)}{2} = \sum_{i=1}^{\infty} a_{2i} x^{2i},$$

$$\frac{A(x) - A(-x)}{2} = \sum_{i=0}^{\infty} a_{2i+1} x^{2i+1}.$$

Summation: If $b_i = \sum_{i=0}^i a_i$ then

$$B(x) = \frac{1}{1-x}A(x).$$

Convolution:

$$A(x)B(x) = \sum_{i=0}^{\infty} \left(\sum_{j=0}^{i} a_j b_{i-j}\right) x^i.$$

God made the natural numbers; all the rest is the work of man. Leopold Kronecker