# UPLB Eliens - Pegaraw Notebook

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1 Data Structures
1.1 Disjoint Set Union
```

```
struct DSU {
         vector<int> parent, size;
         DSU(int n) {
           parent.resize(n);
           size.resize(n);
           for (int i = 0; i < n; i++) make_set(i);</pre>
         void make_set(int v) {
           parent[v] = v;
           size[v] = 1;
  11
         bool is_same(int a, int b) { return find_set(a)
              == find_set(b); }
9
         int find_set(int v) { return v == parent[v] ? v :
               parent[v] = find_set(parent[v]); }
         void union_sets(int a, int b) {
           a = find_set(a);
           b = find_set(b);
           if (a != b) {
             if (size[a] < size[b]) swap(a, b);</pre>
10
  19
             parent[b] = a;
10
             size[a] += size[b];
10
11
11
   23 };
```

# 1.2 Minimum Queue

```
11
 11
      11, 11>> &s2) {
12
    if (s1.empty() || s2.empty()) {
```

```
return s1.empty() ? s2.top().second : s1.top().
    return min(s1.top().second, s2.top().second);
void add_element(ll new_element, stack<pair<ll, ll</pre>
     >> &s1) {
  11 minimum = s1.empty() ? new_element : min(
       new element, s1.top().second);
  s1.push({new_element, minimum});
11 remove_element(stack<pair<11, 11>> &s1, stack
     pair<11, 11>> &s2) {
  if (s2.empty()) {
    while (!sl.empty()) {
      11 element = s1.top().first;
      s1.pop();
      11 minimum = s2.empty() ? element : min(
           element, s2.top().second);
      s2.push({element, minimum});
  11 removed_element = s2.top().first;
  return removed_element;
```

# 1.3 Range Add Point Query

```
template<typename T, typename InType = T>
class SegTreeNode {
public:
  const T IDN = 0, DEF = 0;
  int i, j;
  T val;
  SegTreeNode<T, InType>* lc, * rc;
  SegTreeNode(int i, int j) : i(i), j(j) {
   if (j - i == 1) {
     lc = rc = nullptr;
      val = DEF;
      return:
    int k = (i + j) / 2;
    lc = new SegTreeNode<T, InType>(i, k);
    rc = new SegTreeNode<T, InType>(k, j);
    val = 0;
  SegTreeNode(const vector<InType>& a, int i, int j
      ) : i(i), j(j) {
    if (j - i == 1) {
     lc = rc = nullptr;
      val = (T) a[i];
      return;
    int k = (i + j) / 2;
    lc = new SegTreeNode<T, InType>(a, i, k);
    rc = new SegTreeNode<T, InType>(a, k, j);
    val = 0;
  void range_add(int 1, int r, T x) {
   if (r <= i || j <= 1) return;</pre>
    if (1 <= i && j <= r) {</pre>
      val += x;
      return;
    lc->range_add(l, r, x);
```

```
rc->range_add(1, r, x);
38
39
      T point_query(int k) {
40
        if (k < i \mid | j \le k) return IDN;
41
        if (j - i == 1) return val;
42
        return val + lc->point_query(k) + rc->
             point_query(k);
43
44 };
45 template<typename T, typename InType = T>
    class SegTree {
47
48
      SegTreeNode<T, InType> root;
49
      SegTree(int n) : root(0, n) {}
50
      SegTree(const vector<InType>& a) : root(a, 0, a.
           size()) {}
51
      void range_add(int 1, int r, T x) { root.
           range_add(l, r, x); }
      T point_query(int k) { return root.point_query(k)
           ; }
53 };
```

# 1.4 Range Add Range Query

```
template<typename T, typename InType = T>
    class SegTreeNode {
    public:
      const T IDN = 0, DEF = 0;
      int i, j;
      T val, to_add = 0;
      SegTreeNode<T, InType>* lc, * rc;
      SegTreeNode(int i, int j) : i(i), j(j) {
        if (j - i == 1) {
10
          lc = rc = nullptr;
          val = DEF;
12
          return;
13
14
         int k = (i + j) / 2;
        lc = new SegTreeNode<T, InType>(i, k);
16
         rc = new SegTreeNode<T, InType>(k, j);
17
        val = operation(lc->val, rc->val);
18
19
      SegTreeNode(const vector<InType>& a, int i, int j
           ) : i(i), j(j) {
20
         if (j - i == 1) {
21
          lc = rc = nullptr;
          val = (T) a[i];
23
          return;
24
25
        int k = (i + j) / 2;
        lc = new SegTreeNode<T, InType>(a, i, k);
        rc = new SegTreeNode<T, InType>(a, k, j);
28
        val = operation(lc->val, rc->val);
29
30
      void propagate() {
31
        if (to_add == 0) return;
32
33
        val += to_add;
        if (j - i > 1) {
34
          lc->to_add += to_add;
          rc->to_add += to_add;
36
37
        to\_add = 0;
38
39
      void range_add(int 1, int r, T delta) {
40
        propagate();
41
        if (r <= i || j <= 1) return;</pre>
        if (1 <= i && j <= r) {
```

```
to_add += delta;
          propagate();
        } else {
          lc->range_add(l, r, delta);
          rc->range_add(1, r, delta);
48
          val = operation(lc->val, rc->val);
      T range_query(int 1, int r) {
        propagate();
        if (1 <= i && j <= r) return val;</pre>
        if (j <= 1 || r <= i) return IDN;</pre>
        return operation(lc->range_query(l, r), rc->
             range_query(1, r));
      T operation(T x, T y) {}
58
    template<typename T, typename InType = T>
    class SegTree {
    public:
      SegTreeNode<T, InType> root;
      SegTree(int n) : root(0, n) {}
      SegTree(const vector<InType>& a) : root(a, 0, a.
           size()) {}
      void range_add(int 1, int r, T delta) { root.
           range_add(l, r, delta); }
66
      T range_query(int 1, int r) { return root.
           range_query(1, r); }
67 };
```

# 1.5 Segment Tree

return;

```
template<typename T, typename InType = T>
    class SegTreeNode {
      const T IDN = 0, DEF = 0;
      int i, j;
      T val;
      SegTreeNode<T, InType>* lc, * rc;
      SegTreeNode(int i, int j) : i(i), j(j) {
        if (j - i == 1) {
          lc = rc = nullptr;
          val = DEF;
          return;
        int k = (i + j) / 2;
        lc = new SegTreeNode<T, InType>(i, k);
        rc = new SegTreeNode<T, InType>(k, j);
        val = op(lc->val, rc->val);
18
      SegTreeNode(const vector<InType>& a, int i, int j
          ) : i(i), j(j) {
        if (j - i == 1) {
          lc = rc = nullptr;
          val = (T) a[i];
          return;
        int k = (i + j) / 2;
        lc = new SegTreeNode<T, InType>(a, i, k);
        rc = new SegTreeNode<T, InType>(a, k, j);
28
        val = op(lc->val, rc->val);
      void set(int k, T x) {
        if (k < i | | j <= k) return;</pre>
        if (j - i == 1) {
          val = x;
```

```
1c->set(k, x);
        rc \rightarrow set(k, x);
38
        val = op(lc->val, rc->val);
39
      T range_query(int 1, int r) {
        if (1 <= i && j <= r) return val;</pre>
        if (j <= 1 || r <= i) return IDN;</pre>
        return op(lc->range_query(l, r), rc->
              range_query(1, r));
      T \circ p(T \times, T y) \{ \}
46
    template<typename T, typename InType = T>
    class SegTree {
    public:
50
      SegTreeNode<T, InType> root;
      SegTree(int n) : root(0, n) {}
      SegTree(const vector<InType>& a) : root(a, 0, a.
           size()) {}
      void set(int k, T x) { root.set(k, x); }
      T range_query(int 1, int r) { return root.
           range_query(1, r); }
```

# 1.6 Segment Tree 2d

```
template<typename T, typename InType = T>
    class SegTree2dNode {
    public:
      int i, j, tree_size;
      SegTree<T, InType>* seg_tree;
      SegTree2dNode<T, InType>* lc, * rc;
      SegTree2dNode() {}
      SegTree2dNode(const vector<vector<InType>>& a,
           int i, int j) : i(i), j(j) {
        tree_size = a[0].size();
        if (j - i == 1) {
         lc = rc = nullptr;
          seg_tree = new SegTree<T, InType>(a[i]);
          return;
        int k = (i + j) / 2;
        lc = new SegTree2dNode<T, InType>(a, i, k);
        rc = new SegTree2dNode<T, InType>(a, k, j);
18
        seg_tree = new SegTree<T, InType>(vector<T>(
             tree_size));
19
        operation_2d(lc->seg_tree, rc->seg_tree);
      ~SegTree2dNode() {
        delete lc;
        delete rc:
      void set_2d(int kx, int ky, T x) {
        if (kx < i | | j <= kx) return;</pre>
        if (j - i == 1) {
          seg_tree->set(ky, x);
          return;
        1c->set_2d(kx, ky, x);
        rc->set_2d(kx, ky, x);
        operation_2d(lc->seg_tree, rc->seg_tree);
34
      T range_query_2d(int lx, int rx, int ly, int ry)
        if (lx <= i && j <= rx) return seg_tree->
             range_query(ly, ry);
```

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Pegaraw
```

```
if (j <= lx || rx <= i) return -INF;</pre>
38
        return max(lc->range_query_2d(lx, rx, ly, ry),
             rc->range_query_2d(lx, rx, ly, ry));
39
40
      void operation_2d(SegTree<T, InType>* x, SegTree<</pre>
           T, InType>* y) {
41
        for (int k = 0; k < tree_size; k++) {</pre>
          seg_tree->set(k, max(x->range_query(k, k + 1)
               , y->range_query(k, k + 1)));
43
44
45 };
46 template<typename T, typename InType = T>
47 class SegTree2d {
48 public:
49
      SegTree2dNode<T, InType> root;
50
      SegTree2d() {}
51
      SegTree2d(const vector<vector<InType>>& mat) :
           root(mat, 0, mat.size()) {}
      void set_2d(int kx, int ky, T x) { root.set_2d(kx
           , ky, x); }
      T range_query_2d(int lx, int rx, int ly, int ry)
           { return root.range_query_2d(lx, rx, ly, ry)
54 };
```

# 1.7 Sparse Table

```
1 11 log2_floor(ll i) {
      return i ? __builtin_clzll(1) - __builtin_clzll(i
 4 vector<vector<ll>> build_sum(ll N, ll K, vector<ll>>
       vector<vector<ll>> st(K + 1, vector<ll>(N + 1));
      for (ll i = 0; i < N; i++) st[0][i] = array[i];</pre>
      for (11 i = 1; i <= K; i++)</pre>
        for (11 j = 0; j + (1 << i) <= N; <math>j++)
           st[i][j] = st[i - 1][j] + st[i - 1][j + (1 <<
                 (i - 1));
      return st;
11
   11 sum_query(11 L, 11 R, 11 K, vector<vector<11>>> &
      11 \text{ sum} = 0:
14
      for (11 i = K; i >= 0; i--) {
        if ((1 << i) <= R - L + 1) {</pre>
16
          sum += st[i][L];
17
           L += 1 << i;
18
19
      return sum;
21
    vector<vector<ll>> build_min(ll N, ll K, vector<ll>
      vector<vector<ll>> st(K + 1, vector<ll>(N + 1));
      for (ll i = 0; i < N; i++) st[0][i] = array[i];</pre>
25
      for (11 i = 1; i <= K; i++)
26
        for (11 \ j = 0; \ j + (1 << i) <= N; \ j++)
           st[i][j] = min(st[i-1][j], st[i-1][j+(1
                 << (i - 1))]);
2.8
      return st;
29 }
30 ll min_query(ll L, ll R, vector<vector<ll>>> &st) {
      11 i = log2\_floor(R - L + 1);
32
      return min(st[i][L], st[i][R - (1 << i) + 1]);</pre>
33 }
```

```
1.8 Sparse Table 2d
```

```
const int N = 100;
    int matrix[N][N];
    int table[N][N][(int)(log2(N) + 1)][(int)(log2(N) +
    void build_sparse_table(int n, int m) {
      for (int i = 0; i < n; i++)</pre>
        for (int j = 0; j < m; j++)
          table[i][j][0][0] = matrix[i][j];
      for (int k = 1; k \le (int)(log2(n)); k++)
        for (int i = 0; i + (1 << k) - 1 < n; i++)
          for (int j = 0; j + (1 << k) - 1 < m; <math>j++)
            table[i][j][k][0] = min(table[i][j][k -
                 1][0], table[i + (1 << (k - 1))][j][k
      for (int k = 1; k \le (int)(log2(m)); k++)
        for (int i = 0; i < n; i++)
          for (int j = 0; j + (1 << k) - 1 < m; j++)
15
            table[i][j][0][k] = min(table[i][j][0][k -
                 1], table[i][j + (1 << (k - 1))][0][k
                 - 11);
      for (int k = 1; k \le (int)(log2(n)); k++)
        for (int 1 = 1; 1 <= (int) (log2(m)); 1++)</pre>
          for (int i = 0; i + (1 << k) - 1 < n; i++)
19
            for (int j = 0; j + (1 << 1) - 1 < m; <math>j++)
              table[i][j][k][1] = min(
21
                 min(table[i][j][k-1][l-1], table[i]
                     + (1 << (k - 1)) [j] [k - 1] [1 -
                     1]),
                min(table[i][j + (1 << (1 - 1))][k -
                     1] [1 - 1], table [i + (1 << (k - 1))
                     )][j + (1 << (1 - 1))][k - 1][1 -
24
    int rmq(int x1, int y1, int x2, int y2) {
      int k = log2(x2 - x1 + 1), l = log2(y2 - y1 + 1);
      return max (
        \max(table[x1][y1][k][1], table[x2 - (1 << k) +
             1][y1][k][l]),
        \max(table[x1][y2 - (1 << 1) + 1][k][1], table[
             x^2 - (1 \ll k) + 1][y^2 - (1 \ll 1) + 1][k][1
             ])
     );
31 }
```

# 2 Dynamic Programming

# 2.1 Divide And Conquer

#### 2.2 Edit Distance

#### 2.3 Knapsack

# 2.4 Knuth Optimization

```
... // Initialize dp[i][i] according to the
              problem
11
12
      for (11 i = N - 2; i >= 0; i--) {
13
         for (11 j = i + 1; j < N; j++) {
14
           11 \text{ mn} = 11\_\text{MAX}, \text{ cost} = C(i, j);
15
           for (l1 k = opt[i][j - 1]; k <= min(j - 1,</pre>
                opt[i + 1][j]); k++) {
             if (mn \ge dp[i][k] + dp[k + 1][j] + cost) {
17
               opt[i][j] = k;
               mn = dp[i][k] + dp[k + 1][j] + cost;
19
20
21
           dp[i][j] = mn;
23
24
      cout << dp[0][N - 1] << '\n';
25 }
```

# 2.5 Longest Common Subsequence

```
1  11 LCS(string x, string y, 11 n, 11 m) {
      vector < vector < 11 >> dp(n + 1, vector < 11 > (m + 1));
      for (ll i = 0; i <= n; i++) {
         for (11 j = 0; j \le m; j++) {
          if (i == 0 || j == 0) {
             dp[i][j] = 0;
           } else if (x[i - 1] == y[j - 1]) {
             dp[i][j] = dp[i - 1][j - 1] + 1;
             dp[i][j] = max(dp[i - 1][j], dp[i][j - 1]);
11
        }
13
14
      ll index = dp[n][m];
15
      vector<char> lcs(index + 1);
      lcs[index] = ' \setminus 0';
17
      11 i = n, j = m;
18
      while (i > 0 \&\& j > 0) {
19
        if (x[i-1] == y[j-1]) {
20
          lcs[index - 1] = x[i - 1];
21
          i--;
          j--;
23
          index--;
         } else if (dp[i - 1][j] > dp[i][j - 1]) {
25
          i--:
        } else {
          j--;
28
29
       return dp[n][m];
31 }
```

# 2.6 Longest Increasing Subsequence

```
return r;
    11 LIS(11 n, vector<11> &a) {
13
      11 len = 1;
      vector<11> T(n, 0), R(n, -1);
      T[0] = 0;
      for (ll i = 1; i < n; i++) {</pre>
        if (a[i] < a[T[0]]) {</pre>
          T[0] = i;
        } else if (a[i] > a[T[len - 1]]) {
          R[i] = T[len - 1];
          T[len++] = i;
        } else {
          ll pos = get_ceil_idx(a, T, -1, len - 1, a[i
               1);
          R[i] = T[pos - 1];
25
          T[pos] = i;
26
8.5
      return len:
29
```

#### 2.7 Subset Sum

#### 2.8 Maximum Subarray Sum

```
int max_subarray_sum(vi arr) {
   int x = 0, s = 0;
   for (int k = 0; k < n; k++) {
      s = max(arr[k], s+arr[k]);
      x = max(x,s);
   }
   return x;
}</pre>
```

# 3 Geometry

#### 3.1 Basic Geometry

```
struct point2d {
ftype x, y;
point2d() {}
point2d(ftype x, ftype y): x(x), y(y) {}
point2d& operator+=(const point2d &t) {
    x += t.x;
    y += t.y;
return *this;
```

```
point2d& operator-=(const point2d &t) {
       x -= t.x;
       y -= t.y;
       return *this;
      point2d& operator*=(ftype t) {
       x *= †:
       y *= t;
       return *this;
      point2d& operator/=(ftype t) {
       x /= t;
       y /= t;
        return *this;
      point2d operator+(const point2d &t) const {
           return point2d(*this) += t; }
      point2d operator-(const point2d &t) const {
          return point2d(*this) -= t; }
      point2d operator*(ftype t) const { return point2d
           (*this) *= t; }
      point2d operator/(ftype t) const { return point2d
           (*this) /= t; }
2.9
   point2d operator*(ftype a, point2d b) { return b *
    ftype dot(point2d a, point2d b) { return a.x * b.x
         + a.y * b.y; }
    ftype dot(point3d a, point3d b) { return a.x * b.x
        + a.y * b.y + a.z * b.z; }
    ftype norm(point2d a) { return dot(a, a); }
    double abs(point2d a) { return sqrt(norm(a)); }
   double proj(point2d a, point2d b) { return dot(a, b
        ) / abs(b); }
   double angle(point2d a, point2d b) { return acos(
         dot(a, b) / abs(a) / abs(b)); }
   point3d cross(point3d a, point3d b) { return
        point3d(a.y * b.z - a.z * b.y, a.z * b.x - a.x
         * b.z, a.x * b.y - a.y * b.x); }
38 ftype triple(point3d a, point3d b, point3d c) {
         return dot(a, cross(b, c)); }
   ftype cross(point2d a, point2d b) { return a.x * b.
         y - a.y * b.x; }
40 point2d intersect(point2d a1, point2d d1, point2d
         a2, point2d d2) { return a1 + cross(a2 - a1,
         d2) / cross(d1, d2) * d1; }
41 point3d intersect (point3d al, point3d nl, point3d
         a2, point3d n2, point3d a3, point3d n3) {
      point3d x(n1.x, n2.x, n3.x);
     point3d y(n1.y, n2.y, n3.y);
     point3d z(n1.z, n2.z, n3.z);
      point3d d(dot(a1, n1), dot(a2, n2), dot(a3, n3));
      return point3d(triple(d, y, z), triple(x, d, z),
          triple(x, y, d)) / triple(n1, n2, n3);
47 }
```

#### 3.2 Circle Line Intersection

```
cout << x0 << ' ' << y0 << '\n';
    } else {
10
      double d = r * r - c * c / (a * a + b * b);
      double mult = sqrt (d / (a * a + b * b));
      double ax, ay, bx, by;
13
      ax = x0 + b * mult;
      bx = x0 - b * mult;
      ay = y0 - a * mult;
16
      by = y0 + a * mult;
17
      puts ("2 points");
      cout << ax << ' ' << ay << '\n' << bx << ' ' <<
           by << '\n';
19 }
```

#### 3.3 Convex Hull

```
struct pt {
      double x, y;
 3
    11 orientation(pt a, pt b, pt c) {
      double v = a.x * (b.y - c.y) + b.x * (c.y - a.y)
           + c.x * (a.y - b.y);
      if (v < 0) {
         return -1;
      } else if (v > 0) {
        return +1;
      return 0;
12
13 bool cw(pt a, pt b, pt c, bool include_collinear) {
14
      11 o = orientation(a, b, c);
15
      return o < 0 || (include collinear && o == 0);</pre>
16
17
    bool collinear(pt a, pt b, pt c) {
18
      return orientation(a, b, c) == 0;
19
20 void convex_hull(vector<pt>& a, bool
         include_collinear = false) {
      pt p0 = *min_element(a.begin(), a.end(), [](pt a,
        return make_pair(a.y, a.x) < make_pair(b.y, b.x</pre>
             );
23
      sort(a.begin(), a.end(), [&p0](const pt& a, const
            pt& b) {
25
        11 o = orientation(p0, a, b);
        if (o == 0) {
          return (p0.x - a.x) * (p0.x - a.x) + (p0.y - a.x)
               a.y) * (p0.y - a.y)
               < (p0.x - b.x) * (p0.x - b.x) + (p0.y -
                    b.y) * (p0.y - b.y);
29
        return o < 0;
      if (include_collinear) {
33
        11 i = (11) a.size()-1;
34
         while (i \ge 0 \&\& collinear(p0, a[i], a.back()))
              i--;
35
        reverse(a.begin()+i+1, a.end());
36
      vector<pt> st;
      for (11 i = 0; i < (11) a.size(); i++) {</pre>
38
        while (st.size() > 1 && !cw(st[st.size() - 2],
             st.back(), a[i], include_collinear)) {
40
          st.pop_back();
41
42.
         st.push_back(a[i]);
```

```
a = st;
45 }
 3.4 Line Intersection
    struct pt { double x, y; };
    struct line { double a, b, c; };
    const double EPS = 1e-9;
    double det (double a, double b, double c, double d)
         { return a*d - b*c; }
    bool intersect(line m, line n, pt & res) {
      double zn = det(m.a, m.b, n.a, n.b);
      if (abs(zn) < EPS) return false;</pre>
      res.x = -det(m.c, m.b, n.c, n.b) / zn;
      res.y = -det(m.a, m.c, n.a, n.c) / zn;
      return true:
11 }
12 bool parallel(line m, line n) { return abs(det(m.a,
          m.b, n.a, n.b)) < EPS; }
    bool equivalent(line m, line n) {
      return abs(det(m.a, m.b, n.a, n.b)) < EPS</pre>
          && abs(det(m.a, m.c, n.a, n.c)) < EPS
          && abs(det(m.b, m.c, n.b, n.c)) < EPS;
17 }
 3.5 Line Sweep
    const double EPS = 1E-9;
    struct pt { double x, y; };
    struct seg {
      pt p, q;
      11 id;
      double get_y(double x) const {
        if (abs(p.x - q.x) < EPS) return p.y;</pre>
         return p.y + (q.y - p.y) * (x - p.x) / (q.x - p
             .x);
1.0
    bool intersect1d(double 11, double r1, double 12,
         double r2) {
      if (11 > r1) swap(11, r1);
      if (12 > r2) swap(12, r2);
      return max(11, 12) <= min(r1, r2) + EPS;
16 11 vec(const pt& a, const pt& b, const pt& c) {
      double s = (b.x - a.x) * (c.y - a.y) - (b.y - a.y)
           ) * (c.x - a.x);
18
      return abs(s) < EPS ? 0 : s > 0 ? +1 : -1;
19
    bool intersect(const seg& a, const seg& b) {
      return intersect1d(a.p.x, a.q.x, b.p.x, b.q.x) &&
             intersect1d(a.p.y, a.q.y, b.p.y, b.q.y) &&
             vec(a.p, a.q, b.p) * vec(a.p, a.q, b.q) <=
                   330
             vec(b.p, b.q, a.p) * vec(b.p, b.q, a.q) <=</pre>
                   0:
25
    bool operator<(const seg& a, const seg& b) {</pre>
      double x = max(min(a.p.x, a.q.x), min(b.p.x, b.q.
           x));
      return a.get_y(x) < b.get_y(x) - EPS;</pre>
29 1
30 struct event {
```

double x;

```
11 tp, id;
      event() {}
34
      event (double x, 11 tp, 11 id) : x(x), tp(tp), id(
      bool operator<(const event& e) const {</pre>
        if (abs(x - e.x) > EPS) return x < e.x;
        return tp > e.tp;
38
39 };
40 set<seg> s;
    vector<set<seg>::iterator> where;
    set<seg>::iterator prev(set<seg>::iterator it) {
      return it == s.begin() ? s.end() : --it;
44
45
    set<seg>::iterator next(set<seg>::iterator it) {
     return ++it;
47
48
   pair<11, 11> solve(const vector<seg>& a) {
      11 n = (11) a.size();
      vector<event> e;
      for (11 i = 0; i < n; ++i) {
        e.push_back(event(min(a[i].p.x, a[i].q.x), +1,
        e.push_back(event(max(a[i].p.x, a[i].q.x), -1,
             <u>i</u>));
      sort(e.begin(), e.end());
56
      s.clear();
      where.resize(a.size());
58
      for (size_t i = 0; i < e.size(); ++i) {</pre>
        11 \text{ id} = e[i].id;
        if (e[i].tp == +1) {
61
          set<seg>::iterator nxt = s.lower_bound(a[id])
               , prv = prev(nxt);
62
          if (nxt != s.end() && intersect(*nxt, a[id]))
                return make_pair(nxt->id, id);
          if (prv != s.end() && intersect(*prv, a[id]))
                return make_pair(prv->id, id);
          where[id] = s.insert(nxt, a[id]);
        } else {
          set<seq>::iterator nxt = next(where[id]), prv
                = prev(where[id]);
          if (nxt != s.end() && prv != s.end() &&
               intersect(*nxt, *prv)) return make pair(
               prv->id, nxt->id);
68
          s.erase(where[id]);
      return make_pair(-1, -1);
72
```

# 3.6 Nearest Points

```
гезага
```

```
13 vector<pt> a;
    double mindist;
    pair<11, 11> best_pair;
    void upd_ans(const pt & a, const pt & b) {
      double dist = sqrt((a.x - b.x) * (a.x - b.x) + (a.x - b.x)
           .y - b.y) * (a.y - b.y);
      if (dist < mindist) {</pre>
        mindist = dist;
20
        best_pair = {a.id, b.id};
21
22 }
23 vector<pt> t;
24 void rec(11 1, 11 r) {
25
     if (r - 1 <= 3) {
26
        for (ll i = 1; i < r; ++i)
27
         for (11 \ j = i + 1; \ j < r; ++j)
28
           upd_ans(a[i], a[j]);
29
        sort(a.begin() + 1, a.begin() + r, cmp_y());
        return:
      11 m = (1 + r) >> 1, midx = a[m].x;
      rec(1, m);
      rec(m, r);
      merge(a.begin() + 1, a.begin() + m, a.begin() + m
           , a.begin() + r, t.begin(), cmp_y());
      copy(t.begin(), t.begin() + r - 1, a.begin() + 1)
37
      11 \text{ tsz} = 0;
      for (11 i = 1; i < r; ++i) {
39
       if (abs(a[i].x - midx) < mindist) {</pre>
        for (ll j = tsz - 1; j >= 0 && a[i].y - t[j].
            y < mindist; --j)
            upd_ans(a[i], t[j]);
          t[tsz++] = a[i];
43
44
     }
45 }
46 t.resize(n);
47 sort(a.begin(), a.end(), cmp_x());
48 mindist = 1E20;
49 rec(0, n);
```

# 4 Graph Theory

#### 4.1 Articulation Point

```
void APUtil(vector<vector<ll>>> &adj, ll u, vector
         bool> &visited,
   vector<ll> &disc, vector<ll> &low, ll &time, ll
        parent, vector<bool> &isAP) {
      11 children = 0:
      visited[u] = true;
      disc[u] = low[u] = ++time;
      for (auto v : adj[u]) {
       if (!visited[v]) {
          children++;
          APUtil(adj, v, visited, disc, low, time, u,
          low[u] = min(low[u], low[v]);
11
          if (parent != -1 && low[v] >= disc[u]) {
           isAP[u] = true;
        } else if (v != parent) {
15
          low[u] = min(low[u], disc[v]);
```

```
if (parent == -1 && children > 1) {
       isAP[u] = true;
20
   void AP(vector<vector<ll>> &adj, ll n) {
     vector<ll> disc(n), low(n);
     vector<bool> visited(n), isAP(n);
     11 time = 0, par = -1;
      for (11 u = 0; u < n; u++) {
       if (!visited[u]) {
          APUtil(adj, u, visited, disc, low, time, par,
29
30
      for (11 u = 0; u < n; u++) {
      if (isAP[u]) {
33
         cout << u << " ";
34
     }
36 }
```

# 4.2 Bellman Ford

```
struct Edge {
     int a, b, cost;
    int n, m, v;
    vector<Edge> edges;
    const int INF = 1000000000;
    void solve() {
     vector<int> d(n, INF);
     d[v] = 0;
     vector<int> p(n, -1);
      int x;
      for (int i = 0; i < n; ++i) {
       \mathbf{x} = -1;
       for (Edge e : edges)
         if (d[e.a] < INF)</pre>
            if (d[e.b] > d[e.a] + e.cost) {
              d[e.b] = max(-INF, d[e.a] + e.cost);
              p[e.b] = e.a;
19
              x = e.b;
20
      if (x == -1) cout << "No negative cycle from " <<
      else {
        int y = x;
        for (int i = 0; i < n; ++i) y = p[y];
        vector<int> path;
        for (int cur = y;; cur = p[cur]) {
         path.push_back(cur);
          if (cur == y && path.size() > 1) break;
3.0
        reverse(path.begin(), path.end());
        cout << "Negative cycle: ";</pre>
        for (int u : path) cout << u << ' ';</pre>
35 }
```

#### 4.3 Bridge

```
1 int n;
2 vector<vector<int>> adj;
3 vector<bool> visited;
```

```
vector<int> tin, low;
   int timer;
   void dfs (int v, int p = -1) {
     visited[v] = true;
     tin[v] = low[v] = timer++;
      for (int to : adj[v]) {
      if (to == p) continue;
       if (visited[to]) {
         low[v] = min(low[v], tin[to]);
          dfs(to, v);
          low[v] = min(low[v], low[to]);
          if (low[to] > tin[v]) IS_BRIDGE(v, to);
18
     }
19
20 void find_bridges() {
     timer = 0;
     visited.assign(n, false);
      tin.assign(n, -1);
      low.assign(n, -1);
      for (int i = 0; i < n; ++i) {</pre>
       if (!visited[i]) dfs(i);
28 }
```

# 4.4 Centroid Decomposition

```
1 vector<vector<int>> adj;
2 vector<bool> is removed;
3 vector<int> subtree size;
4 int get_subtree_size(int node, int parent = -1) {
           subtree_size[node] = 1;
           for (int child : adj[node]) {
                   if (child == parent || is_removed[
                        child]) continue;
                   subtree size[node] +=
                        get_subtree_size(child, node);
           return subtree size[node];
   int get_centroid(int node, int tree_size, int
        parent = -1) {
           for (int child : adj[node]) {
                   if (child == parent || is_removed[
                        childl) continue;
                   if (subtree_size[child] * 2 >
                        tree size) return get centroid
                        (child, tree size, node);
           return node;
18
   void build centroid decomp(int node = 0) {
           int centroid = get_centroid(node,
                get_subtree_size(node));
           // do something
           is_removed[centroid] = true;
           for (int child : adj[centroid]) {
                   if (is removed[child]) continue;
                   build_centroid_decomp(child);
```

#### 4.5 Dijkstra

```
const int INF = 1000000000;
    vector<vector<pair<int, int>>> adj;
    void dijkstra(int s, vector<int> & d, vector<int> &
      int n = adj.size();
      d.assign(n, INF);
      p.assign(n, -1);
      d[s] = 0;
      using pii = pair<int, int>;
      priority_queue<pii, vector<pii>, greater<pii>> q;
      q.push({0, s});
      while (!q.empty()) {
12
        int v = g.top().second, d_v = g.top().first;
13
14
        if (d_v != d[v]) continue;
15
        for (auto edge : adj[v]) {
16
          int to = edge.first, len = edge.second;
17
          if (d[v] + len < d[to]) {</pre>
18
              d[to] = d[v] + len;
19
              p[to] = v;
20
              q.push({d[to], to});
21
23
24 }
```

#### 4.6 Dinics

```
struct FlowEdge {
      int v, u;
      11 \text{ cap, flow} = 0;
      FlowEdge(int v, int u, ll cap) : v(v), u(u), cap( | 4.7 Edmonds Karp
           cap) {}
    struct Dinic {
      const 11 flow_inf = 1e18;
      vector<FlowEdge> edges;
      vector<vector<int>> adj;
     int n, m = 0, s, t;
      vector<int> level, ptr;
      queue<int> q;
      Dinic(int n, int s, int t) : n(n), s(s), t(t) {
        adj.resize(n);
15
        level.resize(n):
16
        ptr.resize(n);
17
      void add_edge(int v, int u, 11 cap) {
19
        edges.emplace_back(v, u, cap);
20
        edges.emplace_back(u, v, 0);
21
        adj[v].push_back(m);
22
23
        adj[u].push_back(m + 1);
        m += 2:
24
25
26
      bool bfs() {
         while (!q.empty()) {
27
          int v = q.front();
28
          q.pop();
29
          for (int id : adj[v]) {
30
            if (edges[id].cap - edges[id].flow < 1)</pre>
                  continue:
             if (level[edges[id].u] != -1) continue;
             level[edges[id].u] = level[v] + 1;
             q.push(edges[id].u);
34
35
36
         return level[t] != -1;
```

```
11 dfs(int v, 11 pushed) {
        if (pushed == 0) return 0;
        if (v == t) return pushed;
         for (int& cid = ptr[v]; cid < (int)adj[v].size</pre>
             (); cid++) {
42
          int id = adj[v][cid], u = edges[id].u;
          if (level[v] + 1 != level[u] || edges[id].cap
4.3
                - edges[id].flow < 1) continue;</pre>
          11 tr = dfs(u, min(pushed, edges[id].cap -
               edges[id].flow));
          if (tr == 0) continue;
          edges[id].flow += tr;
          edges[id ^ 1].flow -= tr;
48
          return tr;
49
50
        return 0;
      11 flow() {
        11 f = 0;
        while (true) {
          fill(level.begin(), level.end(), -1);
          level[s] = 0;
          q.push(s);
58
          if (!bfs()) break;
59
          fill(ptr.begin(), ptr.end(), 0);
60
          while (ll pushed = dfs(s, flow_inf)) f +=
               pushed;
62
        return f;
63
64 };
```

```
vector<vector<int>> capacity;
   vector<vector<int>> adj;
 4 int bfs(int s, int t, vector<int>& parent) {
     fill(parent.begin(), parent.end(), -1);
      parent[s] = -2;
      queue<pair<int, int>> q;
      g.push({s, INF});
      while (!q.emptv()) {
        int cur = q.front().first, flow = q.front().
             second:
        q.pop();
        for (int next : adj[cur]) {
13
          if (parent[next] == -1 && capacity[cur][next
               ]) {
14
            parent[next] = cur;
            int new_flow = min(flow, capacity[cur][next
            if (next == t) return new_flow;
            q.push({next, new_flow});
19
       }
      return 0;
    int maxflow(int s, int t) {
      int flow = 0;
      vector<int> parent(n);
      int new flow:
      while (new_flow = bfs(s, t, parent)) {
       flow += new flow;
29
        int cur = t;
```

while (cur != s) {

```
int prev = parent[cur];
    capacity[prev][cur] -= new flow;
    capacity[cur][prev] += new_flow;
    cur = prev;
return flow;
```

```
4.8 Fast Second Mst
    struct edge {
        int s, e, w, id;
        bool operator<(const struct edge& other) {</pre>
             return w < other.w; }</pre>
    typedef struct edge Edge;
    const int N = 2e5 + 5;
    long long res = 0, ans = 1e18;
8 int n, m, a, b, w, id, 1 = 21;
   vector<Edge> edges;
10 vector<int> h(N, 0), parent(N, -1), size(N, 0),
         present (N, 0);
11 vector<vector<pair<int, int>>> adj(N), dp(N, vector
         <pair<int, int>>(1));
    vector<vector<int>> up(N, vector<int>(1, -1));
    pair<int, int> combine(pair<int, int> a, pair<int,</pre>
         int> b) {
      vector<int> v = {a.first, a.second, b.first, b.
          second);
      int topTwo = -3, topOne = -2;
      for (int c : v) {
        if (c > topOne) {
          topTwo = topOne;
          topOne = c;
20
        } else if (c > topTwo && c < topOne) topTwo = c
22
      return {topOne, topTwo};
2.3
2.4
    void dfs(int u, int par, int d) {
     h[u] = 1 + h[par];
      up[u][0] = par;
      dp[u][0] = \{d, -1\};
28
      for (auto v : adj[u]) {
29
        if (v.first != par) dfs(v.first, u, v.second);
    pair<int, int> lca(int u, int v) {
      pair<int, int> ans = \{-2, -3\};
      if (h[u] < h[v]) swap(u, v);</pre>
      for (int i = 1 - 1; i >= 0; i--) {
        if (h[u] - h[v] >= (1 << i)) {
          ans = combine(ans, dp[u][i]);
38
          u = up[u][i];
39
      if (u == v) return ans;
      for (int i = 1 - 1; i >= 0; i--) {
        if (up[u][i] != -1 && up[v][i] != -1 && up[u][i
             ] != up[v][i]) {
          ans = combine(ans, combine(dp[u][i], dp[v][i
          u = up[u][i];
          v = up[v][i];
```

```
49
       ans = combine(ans, combine(dp[u][0], dp[v][0]));
 50
       return ans;
 51
 52
     int main(void) {
 54
       cin >> n >> m;
       for (int i = 1; i <= n; i++) {
 55
 56
         parent[i] = i;
 57
         size[i] = 1;
 58
 59
       for (int i = 1; i <= m; i++) {</pre>
 60
         cin >> a >> b >> w; // 1-indexed
 61
         edges.push_back(\{a, b, w, i - 1\});
 62
 63
       sort(edges.begin(), edges.end());
 64
       for (int i = 0; i <= m - 1; i++) {
 65
        a = edges[i].s;
         b = edges[i].e;
 66
 67
         w = edges[i].w;
 68
         id = edges[i].id;
 69
         if (unite_set(a, b)) {
           adj[a].emplace_back(b, w);
           adj[b].emplace back(a, w);
           present[id] = 1;
 73
           res += w;
 74
 76
       dfs(1, 0, 0);
       for (int i = 1; i <= 1 - 1; i++) {</pre>
78
         for (int j = 1; j \le n; ++j) {
 79
           if (up[j][i - 1] != -1) {
80
             int v = up[j][i - 1];
 81
             up[j][i] = up[v][i - 1];
 82
             dp[j][i] = combine(dp[j][i-1], dp[v][i-
                  11);
83
 84
         }
 85
 86
       for (int i = 0; i <= m - 1; i++) {
 87
         id = edges[i].id;
 88
         w = edges[i].w;
 89
         if (!present[id]) {
 90
           auto rem = lca(edges[i].s, edges[i].e);
 91
           if (rem.first != w) {
 92
             if (ans > res + w - rem.first) ans = res +
                  w - rem.first;
 93
           } else if (rem.second != -1) {
 94
             if (ans > res + w - rem.second) ans = res +
                   w - rem.second;
 95
 96
 97
 98
       cout << ans << "\n";
 99
       return 0;
100 }
```

# 4.9 Find Cycle

```
bool dfs(ll v) {
  color[v] = 1;
  for (ll u : adj[v]) {
    if (color[u] == 0) {
      parent[u] = v;
      if (dfs(u)) {
        return true;
    } else if (color[u] == 1) {
```

```
cycle_end = v;
          cvcle start = u;
          return true;
14
      color[v] = 2;
16
      return false;
18 void find_cycle() {
      color.assign(n, 0);
      parent.assign(n, -1);
      cycle_start = -1;
      for (11 v = 0; v < n; v++) {
        if (color[v] == 0 && dfs(v)) {
          break:
26
      if (cycle_start == -1) {
        cout << "Acyclic" << endl;</pre>
      } else {
        vector<ll> cycle;
        cycle.push_back(cycle_start);
        for (11 v = cycle end; v != cycle start; v =
             parent[v]) {
          cycle.push_back(v);
34
        cycle.push_back(cycle_start);
        reverse(cycle.begin(), cycle.end());
        cout << "Cycle found: ";</pre>
        for (11 v : cycle) {
          cout << v << ' ';
40
        cout << '\n';</pre>
 4.10 Floyd Warshall
```

```
void floyd_warshall(vector<vector<ll>>> &dis, ll n)
      for (11 k = 0; k < n; k++)
        for (11 i = 0; i < n; i++)
          for (11 j = 0; j < n; j++)
            if (dis[i][k] < INF && dis[k][j] < INF)</pre>
              dis[i][j] = min(dis[i][j], dis[i][k] +
                   dis[k][j]);
      for (ll i = 0; i < n; i++)</pre>
        for (11 j = 0; j < n; j++)
9
          for (11 k = 0; k < n; k++)
            if (dis[k][k] < 0 && dis[i][k] < INF && dis</pre>
                 [k][j] < INF
              dis[i][j] = -INF;
12 }
```

#### 4.11 Ford Fulkerson

```
1 bool bfs(ll n, vector<vector<ll>>> &r_graph, ll s,
        11 t, vector<11> &parent) {
     vector<bool> visited(n, false);
     queue<11> q;
     q.push(s);
     visited[s] = true;
     parent[s] = -1;
     while (!q.empty()) {
      ll u = q.front();
```

```
q.pop();
        for (11 v = 0; v < n; v++) {
          if (!visited[v] && r_graph[u][v] > 0) {
            if (v == t) {
              parent[v] = u;
              return true;
            q.push(v);
            parent[v] = u;
            visited[v] = true;
22
      return false;
    11 ford_fulkerson(ll n, vector<vector<ll>> graph,
         11 s, 11 t) {
      11 u. v:
      vector<vector<11>>> r_graph;
      for (u = 0; u < n; u++)
        for (v = 0; v < n; v++)
         r_{graph[u][v]} = graph[u][v];
      vector<11> parent;
      11 \text{ max\_flow} = 0;
      while (bfs(n, r_graph, s, t, parent)) {
        11 path_flow = INF;
34
        for (v = t; v != s; v = parent[v]) {
          u = parent[v];
          path_flow = min(path_flow, r_graph[u][v]);
38
        for (v = t; v != s; v = parent[v]) {
          u = parent[v];
40
          r_graph[u][v] -= path_flow;
          r_graph[v][u] += path_flow;
42
43
        max_flow += path_flow;
44
45
      return max_flow;
```

# 4.12 Hierholzer

```
void print circuit(vector<vector<ll>> &adj) {
      map<11, 11> edge_count;
      for (ll i = 0; i < adj.size(); i++) {</pre>
        edge_count[i] = adj[i].size();
      if (!adj.size()) {
        return;
      stack<ll> curr path;
      vector<ll> circuit:
      curr_path.push(0);
      11 \text{ curr } v = 0;
      while (!curr_path.empty()) {
        if (edge_count[curr_v]) {
          curr_path.push(curr_v);
          11 next_v = adj[curr_v].back();
          edge_count[curr_v]--;
          adj[curr_v].pop_back();
          curr_v = next_v;
        } else {
          circuit.push_back(curr_v);
          curr_v = curr_path.top();
          curr_path.pop();
24
```

```
9
```

```
Pegars
```

#### 4.13 Hungarian

```
1 vector<int> u (n+1), v (m+1), p (m+1), way (m+1);
    for (int i=1; i<=n; ++i) {</pre>
      p[0] = i;
      int j0 = 0;
      vector<int> minv (m+1, INF);
      vector<bool> used (m+1, false);
        used[j0] = true;
        int i0 = p[j0], delta = INF, j1;
         for (int j=1; j<=m; ++j)</pre>
11
          if (!used[j]) {
12
             int cur = A[i0][j]-u[i0]-v[j];
             if (cur < minv[j]) minv[j] = cur, way[j] =</pre>
             if (minv[j] < delta) delta = minv[j], j1 =</pre>
                   j;
15
16
         for (int j=0; j<=m; ++j)</pre>
17
          if (used[j]) u[p[j]] += delta, v[j] -= delta
           else minv[j] -= delta;
19
         j0 = j1;
20
      } while (p[j0] != 0);
        int j1 = way[j0];
23
        p[j0] = p[j1];
         i0 = j1;
24
25
      } while (†0);
26
    vector<int> ans (n+1);
    for (int j=1; j<=m; ++j)</pre>
      ans[p[j]] = j;
30 int cost = -v[0];
```

# 4.14 Is Bipartite

```
bool is_bipartite(vector<ll> &col, vector<vector<ll</pre>
         >> &adj, ll n) {
      queue<pair<ll, ll>> q;
      for (11 i = 0; i < n; i++) {
        if (col[i] == -1) {
          q.push({i, 0});
          col[i] = 0;
          while (!q.empty()) {
            pair<11, 11> p = q.front();
            q.pop();
            11 v = p.first, c = p.second;
            for (ll j : adj[v]) {
              if (col[j] == c) {
                return false;
15
              if (col[j] == -1) {
16
                col[j] = (c ? 0 : 1);
17
                q.push({j, col[j]});
18
19
20
```

```
22 }
23 return true;
24 }
```

# 4.15 Is Cyclic

```
bool is cyclic util(int u, vector<vector<int>> &adj
        , vector<bool> &vis, vector<bool> &rec) {
     vis[u] = true;
     rec[u] = true;
     for(auto v : adj[u]) {
       if (!vis[v] && is_cyclic_util(v, adj, vis, rec)
            ) return true;
6
       else if (rec[v]) return true;
     rec[u] = false;
     return false;
   bool is_cyclic(int n, vector<vector<int>> &adj) {
     vector<bool> vis(n, false), rec(n, false);
     for (int i = 0; i < n; i++)</pre>
       if (!vis[i] && is_cyclic_util(i, adj, vis, rec)
            ) return true;
     return false;
```

#### 4.16 Kahn

```
void kahn(vector<vector<ll>> &adj) {
      ll n = adj.size();
      vector<1l> in_degree(n, 0);
      for (11 u = 0; u < n; u++)
        for (ll v: adj[u]) in_degree[v]++;
      queue<11> q;
      for (11 i = 0; i < n; i++)
        if (in_degree[i] == 0)
9
          q.push(i);
      11 \text{ cnt} = 0;
      vector<ll> top_order;
      while (!q.empty()) {
       11 u = q.front();
        q.pop();
        top_order.push_back(u);
        for (ll v : adj[u])
          if (--in_degree[v] == 0) q.push(v);
18
19
      if (cnt != n) {
        cout << -1 << '\n';
        return;
      // print top_order
```

# 4.17 Kosaraju

```
stk.push(u);
    vector<vector<int>> transpose(int n, vector<vector<</pre>
         int>>% adi) {
      vector<vector<int>> adj_t(n);
      for (int u = 0; u < n; u++) {
        for (int v : adj[u]) {
          adj_t[v].push_back(u);
      return adj_t;
19
    void get_scc(int u, vector<vector<int>>& adj_t,
         vector<bool>& vis, vector<int>& scc) {
      vis[u] = true;
      scc.push back(u);
      for (int v : adj_t[u]) {
        if (!vis[v]) {
          get_scc(v, adj_t, vis, scc);
2.9
    void kosaraju(int n, vector<vector<int>>& adj,
         vector<vector<int>>& sccs) {
      vector<bool> vis(n, false);
      stack<int> stk;
34
      for (int u = 0; u < n; u++) {
        if (!vis[u]) {
          topo_sort(u, adj, vis, stk);
38
      vector<vector<int>> adj_t = transpose(n, adj);
      for (int u = 0; u < n; u++) {
       vis[u] = false;
43
      while (!stk.empty()) {
        int u = stk.top();
        stk.pop();
        if (!vis[u]) {
          vector<int> scc;
          get_scc(u, adj_t, vis, scc);
          sccs.push_back(scc);
52 }
```

# 4.18 Kruskals

```
1  struct Edge {
2    int u, v, weight;
3    bool operator<(Edge const& other) {
4       return weight < other.weight;
5    }
6    };
7    int n;
8    vector<Edge> edges;
9    int cost = 0;
10    vector<Edge> result;
11    DSU dsu = DSU(n);
12    sort(edges.begin(), edges.end());
13    for (Edge e : edges) {
14       if (dsu.find_set(e.u) != dsu.find_set(e.v)) {
```

#### 4.19 Kruskal Mst

```
struct Edge {
      ll u, v, weight;
      bool operator<(Edge const& other) {</pre>
        return weight < other.weight;</pre>
 6 };
 7 11 n;
   vector<Edge> edges;
 9 11 cost = 0;
10 vector<ll> tree_id(n);
11 vector<Edge> result;
    for (11 i = 0; i < n; i++) {
13
     tree_id[i] = i;
14 }
15 sort(edges.begin(), edges.end());
16 for (Edge e : edges) {
17
      if (tree_id[e.u] != tree_id[e.v]) {
        cost += e.weight;
18
19
        result.push_back(e);
20
        ll old_id = tree_id[e.u], new_id = tree_id[e.v
        for (11 i = 0; i < n; i++) {</pre>
          if (tree_id[i] == old_id) {
            tree_id[i] = new_id;
2.4
25
26
```

# 4.20 Kuhn

```
1 int n, k;
    vector<vector<int>> q;
 3 vector<int> mt;
 4 vector<bool> used;
    bool try_kuhn(int v) {
     if (used[v]) return false;
      used[v] = true;
      for (int to : g[v]) {
       if (mt[to] == -1 || try_kuhn(mt[to])) {
         mt[to] = v;
          return true;
13
14
      return false;
15
    int main() {
17
      mt.assign(k, -1);
18
        vector<bool> used1(n, false);
19
        for (int v = 0; v < n; ++v) {
20
          for (int to : g[v]) {
21
           if (mt[to] == -1) {
              mt[to] = v;
23
              used1[v] = true;
24
              break;
25
          }
```

#### 4.21 Lowest Common Ancestor

```
struct LCA {
      vector<ll> height, euler, first, segtree;
       vector<bool> visited;
       LCA(vector<vector<ll>> &adj, ll root = 0) {
       n = adj.size();
        height.resize(n);
        first.resize(n);
        euler.reserve(n * 2);
        visited.assign(n, false);
        dfs(adj, root);
         11 m = euler.size();
         segtree.resize(m * 4);
        build(1, 0, m - 1);
15
       void dfs(vector<vector<ll>>> &adj, ll node, ll h =
         visited[node] = true;
         height[node] = h;
         first[node] = euler.size();
         euler.push_back(node);
         for (auto to : adj[node]) {
           if (!visited[to]) {
             dfs(adj, to, h + 1);
             euler.push_back(node);
2.6
       void build(ll node, ll b, ll e) {
        if (b == e) segtree[node] = euler[b];
30
           11 \text{ mid} = (b + e) / 2;
           build(node << 1, b, mid);</pre>
           build(node << 1 | 1, mid + 1, e);</pre>
           11 1 = segtree[node << 1], r = segtree[node</pre>
                << 1 | 1];
           segtree[node] = (height[1] < height[r]) ? 1 :</pre>
                 r;
36
       11 query(11 node, 11 b, 11 e, 11 L, 11 R) {
        if (b > R | | e < L) return -1;</pre>
40
         if (b >= L && e <= R) return segtree[node];</pre>
         11 \text{ mid} = (b + e) >> 1;
         11 left = query(node << 1, b, mid, L, R);</pre>
         11 right = query(node << 1 | 1, mid + 1, e, L,</pre>
              R);
         if (left == -1) return right;
         if (right == -1) return left;
         return height[left] < height[right] ? left :</pre>
              right;
48
       11 lca(ll u, ll v) {
         11 left = first[u], right = first[v];
```

# 4.22 Maximum Bipartite Matching

```
bool bpm(ll n, ll m, vector<vector<bool>>> &bpGraph,
          11 u, vector<bool> &seen, vector<ll> &matchR)
      for (11 v = 0; v < m; v++) {
        if (bpGraph[u][v] && !seen[v]) {
          seen[v] = true;
          if (matchR[v] < 0 || bpm(n, m, bpGraph,</pre>
               matchR[v], seen, matchR)) {
            matchR[v] = u;
            return true;
        }
      return false;
    11 maxBPM(11 n, 11 m, vector<vector<bool>>> &bpGraph
      vector<11> matchR(m, -1);
      11 \text{ result} = 0;
      for (11 u = 0; u < n; u++) {
        vector<bool> seen(m, false);
        if (bpm(n, m, bpGraph, u, seen, matchR)) {
          result++;
      return result;
23
```

# 4.23 Min Cost Flow

```
1 struct Edge {
     int from, to, capacity, cost;
   vector<vector<int>> adj, cost, capacity;
    const int INF = 1e9;
    void shortest_paths(int n, int v0, vector<int>& d,
         vector<int>& p) {
      d.assign(n, INF);
      d[v0] = 0;
      vector<bool> inq(n, false);
      queue<int> q;
      q.push(v0);
      p.assign(n, -1);
      while (!q.empty()) {
       int u = q.front();
        q.pop();
        inq[u] = false;
        for (int v : adj[u]) {
         if (capacity[u][v] > 0 && d[v] > d[u] + cost[
              u][v]) {
            d[v] = d[u] + cost[u][v];
            p[v] = u;
            if (!inq[v]) {
             inq[v] = true;
              q.push(v);
24
          }
```

```
27
28
29 int min_cost_flow(int N, vector<Edge> edges, int K,
          int s, int t) {
      adj.assign(N, vector<int>());
31
      cost.assign(N, vector<int>(N, 0));
      capacity.assign(N, vector<int>(N, 0));
33
      for (Edge e : edges) {
34
        adj[e.from].push back(e.to);
        adj[e.to].push_back(e.from);
36
        cost[e.from][e.to] = e.cost;
37
        cost[e.to][e.from] = -e.cost;
38
        capacity[e.from][e.to] = e.capacity;
39
40
      int flow = 0;
41
      int cost = 0;
42
      vector<int> d, p;
43
      while (flow < K) {
       shortest_paths(N, s, d, p);
45
        if (d[t] == INF) break;
46
        int f = K - flow, cur = t;
47
        while (cur != s) {
48
          f = min(f, capacity[p[cur]][cur]);
49
          cur = p[cur];
50
51
        flow += f;
        cost += f * d[t];
        cur = t;
        while (cur != s) {
          capacity[p[cur]][cur] -= f;
          capacity[cur][p[cur]] += f;
57
          cur = p[cur];
58
59
60
      if (flow < K) return -1;</pre>
61
      else return cost;
62 F
```

# 4.24 Prim

```
1 const int INF = 10000000000;
    struct Edge {
      int w = INF, to = -1;
      bool operator<(Edge const& other) const {</pre>
         return make_pair(w, to) < make_pair(other.w,</pre>
             other.to);
 6
 7
    };
 8
    int n;
    vector<vector<Edge>> adj;
10 void prim() {
      int total_weight = 0;
      vector<Edge> min_e(n);
      \min_{e}[0].w = 0;
14
      set < Edge > q;
15
      q.insert({0, 0});
      vector<bool> selected(n, false);
17
      for (int i = 0; i < n; ++i) {
18
       if (q.empty()) {
19
          cout << "No MST!" << endl;</pre>
20
          exit(0);
22
        int v = q.begin()->to;
23
         selected[v] = true;
24
         total_weight += q.begin()->w;
25
         q.erase(q.begin());
```

# 4.25 Topological Sort

```
1 void dfs(11 v) {
     visited[v] = true;
      for (ll u : adj[v]) {
       if (!visited[u]) {
         dfs(u);
6
 7
     }
8
     ans.push_back(v);
9 }
10 void topological_sort() {
     visited.assign(n, false);
     ans.clear();
     for (11 i = 0; i < n; ++i) {
      if (!visited[i]) {
         dfs(i);
16
18
     reverse(ans.begin(), ans.end());
```

# 4.26 Zero One Bfs

```
1 vector<int> d(n, INF);
 2 	 d[s] = 0;
3 deque<int> q;
4 q.push_front(s);
   while (!q.empty()) {
     int v = q.front();
     q.pop_front();
      for (auto edge : adj[v]) {
      int u = edge.first, w = edge.second;
10
       if (d[v] + w < d[u]) {
        d[u] = d[v] + w;
12
          if (w == 1) q.push_back(u);
13
          else q.push_front(u);
14
     }
16 }
```

# 5 Math

# 5.1 Chinese Remainder Theorem

```
1 struct Congruence {
2     11 a, m;
3     };
4
```

#### 5.2 Extended Euclidean

```
1 int gcd(int a, int b, int& x, int& y) {
2    if (b == 0) {
3         x = 1;
4         y = 0;
5         return a;
6    }
7    int x1, y1, d = gcd(b, a % b, x1, y1);
8         x = y1;
9         y = x1 - y1 * (a / b);
10    return d;
11 }
```

# 5.3 Factorial Modulo

#### 5.4 Fast Fourier Transform

```
1  using cd = complex<double>;
2  const double PI = acos(-1);
3  void fft(vector<cd>& a, bool invert) {
4   int n = a.size();
5   if (n == 1) return;
6  vector<cd> a0 (n / 2), a1 (n / 2);
7  for (int i = 0; 2 * i < n; i++) {
8   a0[i] = a[2 * i];
9   a1[i] = a[2 * i + 1];
10  }
11  fft(a0, invert);
12  fft(a1, invert);
13  double ang = 2 * PI / n * (invert ? -1 : 1);
14  cd w(1), wn(cos(ang), sin(ang));</pre>
```

```
for (int i = 0; 2 * i < n; i++) {
16
        a[i] = a0[i] + w * a1[i];
17
        a[i + n / 2] = a0[i] - w * a1[i];
18
        if (invert) {
19
         a[i] /= 2;
20
          a[i + n / 2] /= 2;
21
        w \star = wn:
23
     }
24 }
25 vector<int> multiply(vector<int> const& a, vector<
         int> const& b) {
         vector<cd> fa(a.begin(), a.end()), fb(b.begin()
27
        int n = 1;
28
        while (n < a.size() + b.size()) n <<= 1;</pre>
29
        fa.resize(n);
         fb.resize(n);
31
        fft(fa, false);
        fft(fb, false);
         for (int i = 0; i < n; i++) fa[i] *= fb[i];</pre>
34
         fft(fa, true);
35
         vector<int> result(n);
36
         for (int i = 0; i < n; i++) result[i] = round(</pre>
             fa[i].real());
         return result;
38 }
```

#### 5.5 Fibonacci

```
1 /*
   Properties:
    - Cassini's identity: f[n-1]f[n+1] - f[n]^2 = (-1)^2
    - d'Ocagne's identity: f[m]f[n+1] - f[m+1]f[n] =
         (-1)^n f[m-n]
   - Addition rule: f[n+k] = f[k]f[n+1] + f[k-1]f[n]
 6 - k = n case: f[2n] = f[n](f[n+1] + f[n-1])
 7 - f[n] \mid f[nk]
 8 - f[n] / f[m] => n / m
   - GCD rule: gcd(f[m], f[n]) = f[gcd(m, n)]
10 - [[1 \ 1], [1 \ 0]]^n = [[f[n+1] \ f[n]], [f[n], f[n]]
         -11111
   -f[2k+1] = f[k+1]^2 + f[k]^2
    -f[2k] = f[k](f[k+1] + f[k-1]) = f[k](2f[k+1] - f[
13 - Periodic sequence modulo p
14 - sum[i=1..n]f[i] = f[n+2] - 1
15 - sum[i=0..n-1]f[2i+1] = f[2n]
16 - sum[i=1..n]f[2i] = f[2n+1] - 1
17 - sum[i=1..n]f[i]^2 = f[n]f[n+1]
18 Fibonacci encoding:
   1. Iterate through the Fibonacci numbers from the
         largest to the
   smallest until you find one less than or equal to n
21 2. Suppose this number was F_i. Subtract F_i from
          n and put a 1
22 in the i-2 position of the code word (indexing from
          O from the
   leftmost to the rightmost bit).
24 3. Repeat until there is no remainder.
    4. Add a final 1 to the codeword to indicate its
   Closed-form: f[n] = (((1 + rt(5))/2)^n - ((1 - rt))/2)^n
         (5)) / 2) ^n) /rt (5)
```

```
struct matrix {
     11 mat[2][2];
      matrix friend operator *(const matrix &a, const
           matrix &b) {
        matrix c;
32
        for (int i = 0; i < 2; i++) {
3.3
          for (int j = 0; j < 2; j++) {
            c.mat[i][j] = 0;
            for (int k = 0; k < 2; k++) c.mat[i][j] +=</pre>
                 a.mat[i][k] * b.mat[k][j];
        return c:
39
40
    matrix matpow (matrix base, 11 n) {
     matrix ans{ {
      {1, 0},
       {0, 1}
44
      } }:
      while (n) {
        if (n & 1) ans = ans * base;
       base = base * base;
       n >>= 1;
      return ans;
52
5.3
    11 fib(int n) {
      matrix base{ {
        {1, 1},
56
       {1, 0}
      } };
      return matpow(base, n).mat[0][1];
59
    pair<int, int> fib (int n) {
      if (n == 0) return {0, 1};
      auto p = fib(n >> 1);
63
      int c = p.first * (2 * p.second - p.first);
      int d = p.first * p.first + p.second * p.second;
      if (n & 1) return {d, c + d};
66
      else return {c, d};
```

#### 5.6 Find All Solutions

```
1 bool find_any_solution(ll a, ll b, ll c, ll &x0, ll
      &y0, 11 &g) {
     g = gcd_extended(abs(a), abs(b), x0, y0);
     if (c % g) return false;
    x0 *= c / g;
    y0 \star = c / g;
     if (a < 0) x0 = -x0;
     if (b < 0) y0 = -y0;
     return true;
  void shift_solution(ll & x, ll & y, ll a, ll b, ll
      cnt) {
     x += cnt * b;
     y -= cnt * a;
  11 find_all_solutions(ll a, ll b, ll c, ll minx, ll
        maxx, 11 miny, 11 maxy) {
     11 x, y, g;
     if (!find_any_solution(a, b, c, x, y, g)) return
         0;
     a /= g;
     b /= g;
```

```
11 \text{ sign}_a = a > 0 ? +1 : -1;
      11 \text{ sign } b = b > 0 ? +1 : -1;
      shift_solution(x, y, a, b, (minx - x) / b);
      if (x < minx) shift_solution(x, y, a, b, sign_b);</pre>
23
      if (x > maxx) return 0;
2.4
      11 1x1 = x:
      shift_solution(x, y, a, b, (maxx - x) / b);
      if (x > maxx) shift_solution(x, y, a, b, -sign_b)
      11 \text{ rx1} = x;
      shift_solution(x, y, a, b, -(miny - y) / a);
29
      if (y < miny) shift_solution(x, y, a, b, -sign_a)</pre>
          ;
      if (y > maxy) return 0;
      11 \ 1x2 = x;
      shift_solution(x, y, a, b, -(maxy - y) / a);
      if (y > maxy) shift_solution(x, y, a, b, sign_a);
      11 \text{ rx2} = x;
      if (1x2 > rx2) swap(1x2, rx2);
      11 1x = max(1x1, 1x2), rx = min(rx1, rx2);
      if (1x > rx) return 0;
38
      return (rx - 1x) / abs(b) + 1;
39 }
```

#### 5.7 Linear Sieve

#### 5.8 Matrix

```
Matrix exponentation:
3 	 f[n] = af[n-1] + bf[n-2] + cf[n-3]
   |f[n]| | |a|b|c||f[n-1]|
   |f[n-1]| = |1 \ 0 \ 0| |f[n-2]|
   |f[n-2]| |0 1 0||f[n-3]|
   To get:
   |f[n] | |a b c|^(n-2)|f[2]|
   |f[n-1]|=|1 0 0| |f[1]|
   |f[n-2]| |0 1 0|
                         |f[0]|
   struct Matrix { int mat[MAX_N][MAX_N]; };
   Matrix matrix_mul(Matrix a, Matrix b) {
     Matrix ans; int i, j, k;
     for (i = 0; i < MAX_N; i++)</pre>
     for (j = 0; j < MAX_N; j++)
     for (ans.mat[i][j] = k = 0; k < MAX_N; k++)</pre>
      ans.mat[i][j] += a.mat[i][k] * b.mat[k][j];
     return ans;
  Matrix matrix_pow(Matrix base, int p) {
     Matrix ans; int i, j;
```

#### 5.9 Miller Rabin

```
1 using u64 = uint64 t;
    using u128 = uint128 t;
    u64 binpower(u64 base, u64 e, u64 mod) {
    u64 result = 1;
     base %= mod:
     while (e) {
       if (e & 1) result = (u128) result * base % mod;
      base = (u128) base * base % mod;
       e >>= 1:
    return result:
12
13 bool check_composite(u64 n, u64 a, u64 d, 11 s) {
    u64 x = binpower(a, d, n);
    if (x == 1 | | x == n - 1) return false;
    for (11 r = 1; r < s; r++) {
17
       x = (u128) x * x % n;
18
       if (x == n - 1) return false;
19
20
    return true;
21 }
22 bool miller rabin(u64 n) {
    if (n < 2) return false;</pre>
    11 r = 0;
u64 d = n - 1;
     while ((d & 1) == 0) {
     d >>= 1;
28
29
     for (11 a : {2, 3, 5, 7, 11, 13, 17, 19, 23, 29,
       31, 37}) {
       if (n == a) return true;
       if (check_composite(n, a, d, r)) return false;
      return true;
35 }
```

#### 5.10 Modulo Inverse

```
1  11 mod_inv(11 a, 11 m) {
2    if (m == 1) return 0;
3    11 m0 = m, x = 1, y = 0;
4    while (a > 1) {
5        11 q = a / m, t = m;
6        m = a % m;
7        a = t;
8        t = y;
9        y = x - q * y;
10        x = t;
11    }
12    if (x < 0) x += m0;
13    return x;</pre>
```

# 5.11 Pollard Rho Brent

14 }

```
1 11 mult(11 a, 11 b, 11 mod) {
     2 return (__int128_t) a * b % mod;
     3 }
     4 11 f(11 x, 11 c, 11 mod) {
                               return (mult(x, x, mod) + c) % mod;
      7 ll pollard_rho_brent(ll n, ll x0 = 2, ll c = 1) {
      8 11 \times 10^{-2} \times 10^{-2}
                                     while (g == 1) {
                                        v = x:
                                              for (11 i = 1; i < 1; i++) x = f(x, c, n);
                                               11 k = 0;
                                                 while (k < 1 \&\& q == 1) {
                                                         xs = x;
                                                       for (11 i = 0; i < m && i < 1 - k; i++) {
                                                             x = f(x, c, n);
                                                                       q = mult(q, abs(y - x), n);
                                                             g = \underline{gcd}(q, n);
                                                             k += m;
                                                 1 *= 2;
                                     if (q == n) {
                                                do {
                                                  xs = f(xs, c, n);
                                                       g = \underline{gcd}(abs(xs - y), n);
2.8
                                             } while (g == 1);
 29
3.0
                                 return q;
```

# 5.12 Range Sieve

```
1 vector<bool> range_sieve(ll 1, ll r) {
    11 n = sqrt(r);
      vector<bool> is_prime(n + 1, true);
      vector<ll> prime;
      is prime[0] = is prime[1] = false;
      prime.push_back(2);
      for (11 i = 4; i <= n; i += 2) is_prime[i] =</pre>
          false;
      for (11 i = 3; i <= n; i += 2) {
 9
      if (is prime[i]) {
        prime.push_back(i);
          for (ll j = i * i; j <= n; j += i) is_prime[j</pre>
              ] = false;
      vector<bool> result(r - 1 + 1, true);
      for (ll i : prime)
      for (11 j = max(i * i, (1 + i - 1) / i * i); j
         <= r; j += i)
         result[j - 1] = false;
   if (1 == 1) result[0] = false;
      return result;
20 }
```

# 5.13 Segmented Sieve

```
vector<ll> segmented_sieve(ll n) {
     const 11 S = 10000;
     11 nsqrt = sqrt(n);
    vector<char> is_prime(nsqrt + 1, true);
     vector<11> prime;
     is_prime[0] = is_prime[1] = false;
      prime.push_back(2);
      for (11 i = 4; i <= nsqrt; i += 2) {
      is prime[i] = false;
      for (11 i = 3; i <= nsqrt; i += 2) {</pre>
       if (is_prime[i]) {
         prime.push_back(i);
         for (11 j = i * i; j <= nsqrt; j += i) {</pre>
          is prime[j] = false;
       }
18
      }
      vector<11> result;
      vector<char> block(S);
      for (11 k = 0; k * S \le n; k++) {
       fill(block.begin(), block.end(), true);
       for (11 p : prime) {
        for (11 \ j = max((k * S + p - 1) / p, p) * p -
              k * S; j < S; j += p) {
           block[i] = false;
         }
        if (k == 0) {
        block[0] = block[1] = false;
        for (11 i = 0; i < S && k * S + i <= n; i++) {
        if (block[i]) {
           result.push back(k * S + i);
34
       }
      return result;
```

#### 5.14 Sum Of Divisors

```
1     11     sum_of_divisors(11     num) {
    11 total = 1;
   for (int i = 2; (11) i * i <= num; i++) {</pre>
     if (num % i == 0) {
         int e = 0;
         do {
          e++;
          num /= i;
         } while (num % i == 0);
         11 \text{ sum} = 0, \text{ pow} = 1;
         do {
          sum += pow;
          pow *= i;
         } while (e-- > 0);
         total *= sum;
     if (num > 1) total *= (1 + num);
     return total;
```

#### 5.15 Tonelli Shanks

```
11 legendre(ll a, ll p) {
      return bin pow mod(a, (p-1) / 2, p);
 3
    11 tonelli_shanks(ll n, ll p) {
      if (legendre(n, p) == p - 1) {
        return -1;
      if (p % 4 == 3) {
        return bin_pow_mod(n, (p + 1) / 4, p);
11
      11 Q = p - 1, S = 0;
      while (Q \% 2 == 0) {
13
       Q /= 2;
14
        S++;
15
16
      11 z = 2;
17
      for (; z < p; z++) {</pre>
18
       if (legendre(z, p) == p - 1) {
19
          break:
20
21
22
      11 M = S, c = bin_pow_mod(z, Q, p), t =
           bin_pow_mod(n, Q, p), R = bin_pow_mod(n, Q)
           + 1) / 2, p);
      while (t % p != 1) {
24
        if (t % p == 0) {
25
          return 0;
26
27
        11 i = 1, t2 = t * t % p;
28
        for (; i < M; i++) {</pre>
29
          if (t2 % p == 1) {
            break;
          t2 = t2 * t2 % p;
34
        11 b = bin_pow_mod(c, bin_pow_mod(2, M - i - 1,
              p), p);
        M = i;
36
        c = b * b % p;
       t = t * c % p;
38
       R = R * b % p;
39
40
      return R;
41
```

# 6 Miscellaneous

# 6.1 Gauss

```
const double EPS = 1e-9;
    const 11 INF = 2;
    11 gauss (vector <vector <double>> a, vector <double>
      ll n = (ll) a.size(), m = (ll) a[0].size() - 1;
      vector<ll> where (m, -1);
      for (11 col = 0, row = 0; col < m && row < n; ++
           col) {
         11 sel = row;
         for (ll i = row; i < n; ++i) {</pre>
          if (abs(a[i][col]) > abs(a[sel][col])) {
            sel = i;
11
12
         if (abs (a[sel][col]) < EPS) {</pre>
14
           continue;
15
```

```
for (ll i = col; i <= m; ++i) {</pre>
          swap(a[sel][i], a[row][i]);
18
19
        where[col] = row;
20
        for (ll i = 0; i < n; ++i) {
          if (i != row) {
            double c = a[i][col] / a[row][col];
            for (11 j = col; j <= m; ++j) {</pre>
              a[i][j] -= a[row][j] * c;
          }
28
29
30
      ans.assign(m, 0);
      for (11 i = 0; i < m; ++i) {
       if (where[i] != -1) {
33
          ans[i] = a[where[i]][m] / a[where[i]][i];
34
      for (11 i = 0; i < n; ++i) {
        double sum = 0;
        for (11 j = 0; j < m; ++j) {
          sum += ans[j] * a[i][j];
40
        if (abs (sum - a[i][m]) > EPS) {
42
          return 0:
43
45
      for (11 i = 0; i < m; ++i) {
        if (where[i] == -1) {
          return INF;
48
49
      return 1;
```

# 6.2 Ternary Search

# 7 Strings

# 7.1 Count Unique Substrings

```
int count_unique_substrings(string const& s) {
  int n = s.size();
  const int p = 31;
  const int m = 1e9 + 9;
  vector<long long> p_pow(n);
}
```

```
p_pow[0] = 1;
      for (int i = 1; i < n; i++) p pow[i] = (p pow[i -
            1] * p) % m;
      vector<long long> h(n + 1, 0);
      for (int i = 0; i < n; i++) h[i + 1] = (h[i] + (s)
           [i] - 'a' + 1) * p_pow[i]) % m;
      int cnt = 0;
      for (int 1 = 1; 1 <= n; 1++) {
        unordered_set<long long> hs;
        for (int i = 0; i <= n - 1; i++) {</pre>
          long long cur_h = (h[i + 1] + m - h[i]) % m;
          cur_h = (cur_h * p_pow[n - i - 1]) % m;
          hs.insert(cur_h);
18
        cnt += hs.size();
19
      return cnt;
21
```

# 7.2 Finding Repetitions

```
vector<int> z_function(string const& s) {
      int n = s.size();
      vector<int> z(n);
      for (int i = 1, l = 0, r = 0; i < n; i++) {
        if (i \le r) z[i] = min(r - i + 1, z[i - 1]);
        while (i + z[i] < n \&\& s[z[i]] == s[i + z[i]])
            z[i]++;
        if (i + z[i] - 1 > r) {
         1 = i;
          r = i + z[i] - 1;
      return z;
    int get_z(vector<int> const& z, int i) {
      if (0 <= i && i < (int) z.size()) return z[i];</pre>
      else return 0;
18
    vector<pair<int, int>> repetitions;
    void convert_to_repetitions(int shift, bool left,
         int cntr, int 1, int k1, int k2) {
      for (int 11 = \max(1, 1 - k2); 11 \le \min(1, k1);
           11++) {
        if (left && 11 == 1) break;
22
        int 12 = 1 - 11;
        int pos = shift + (left ? cntr - 11 : cntr - 1
             -11+1);
2.4
        repetitions.emplace_back(pos, pos + 2 * 1 - 1);
25
    void find_repetitions(string s, int shift = 0) {
      int n = s.size();
      if (n == 1) return;
     int nu = n / 2;
      int nv = n - nu;
      string u = s.substr(0, nu);
      string v = s.substr(nu);
      string ru(u.rbegin(), u.rend());
      string rv(v.rbegin(), v.rend());
      find_repetitions(u, shift);
      find_repetitions(v, shift + nu);
      vector<int> z1 = z_function(ru);
      vector<int> z2 = z_function(v + '#' + u);
      vector<int> z3 = z function(ru + '#' + rv);
      vector<int> z4 = z_function(v);
      for (int cntr = 0; cntr < n; cntr++) {</pre>
```

```
43
         int 1, k1, k2;
44
         if (cntr < nu) {</pre>
45
          1 = nu - cntr;
46
          k1 = get_z(z1, nu - cntr);
47
          k2 = get_z(z2, nv + 1 + cntr);
48
         l else (
49
          1 = cntr - nu + 1;
          k1 = \text{get}_z(z3, nu + 1 + nv - 1 - (cntr - nu))
           k2 = get_z(z4, (cntr - nu) + 1);
52
53
         if (k1 + k2 >= 1) convert_to_repetitions(shift,
               cntr < nu, cntr, 1, k1, k2);</pre>
54
55 }
```

# 7.3 Group Identical Substrings

```
1 vector<vector<int>> group_identical_strings(vector<</pre>
         string> const& s) {
      int n = s.size();
      vector<pair<long long, int>> hashes(n);
      for (int i = 0; i < n; i++) hashes[i] = {</pre>
           compute_hash(s[i]), i);
      sort(hashes.begin(), hashes.end());
      vector<vector<int>> groups;
      for (int i = 0; i < n; i++) {
        if (i == 0 || hashes[i].first != hashes[i - 1].
             first) groups.emplace_back();
9
        groups.back().push_back(hashes[i].second);
11
      return groups;
12 }
```

# 7.4 Hashing

# 7.5 Knuth Morris Pratt

```
vector<11> prefix_function(string s) {
    ll n = (ll) s.length();
    vector<1l> pi(n);
    for (ll i = 1; i < n; i++) {
        ll j = pi[i - 1];
        while (j > 0 && s[i] != s[j]) j = pi[j - 1];
        if (s[i] == s[j]) j++;
        pi[i] = j;
    }
    return pi;
}
vector<int> ans(n + 1);
```

```
14 for (int i = 0; i < n; i++)
15 ans[pi[i]]++;
16 for (int i = n-1; i > 0; i--)
17 ans[pi[i-1]] += ans[i];
18 for (int i = 0; i <= n; i++)
19 ans[i]++;
```

# 7.6 Longest Common Prefix

```
vector<int> lcp_construction(string const& s,
         vector<int> const& p) {
      int n = s.size();
      vector<int> rank(n, 0);
      for (int i = 0; i < n; i++) rank[p[i]] = i;</pre>
 4
 5
      int k = 0;
      vector<int> lcp(n-1, 0);
      for (int i = 0; i < n; i++) {
        if (rank[i] == n - 1) {
         k = 0:
          continue;
12
        int j = p[rank[i] + 1];
        while (i + k < n \&\& j + k < n \&\& s[i + k] == s[
            j + k]) k++;
        lcp[rank[i]] = k;
        if (k) k--;
      return lcp;
```

# 7.7 Manacher

```
vector<int> manacher odd(string s) {
      int n = s.size();
      s = "$" + s + "^";
      vector<int> p(n + 2);
      int 1 = 1, r = 1;
      for(int i = 1; i <= n; i++) {</pre>
       p[i] = max(0, min(r - i, p[1 + (r - i)]));
        while(s[i - p[i]] == s[i + p[i]]) p[i]++;
       if(i + p[i] > r) l = i - p[i], r = i + p[i];
10
11
      return vector<int>(begin(p) + 1, end(p) - 1);
12
   vector<int> manacher(string s) {
14
15
      for(auto c: s) t += string("#") + c;
      auto res = manacher_odd(t + "#");
      return vector<int>(begin(res) + 1, end(res) - 1);
18 }
```

# 7.8 Rabin Karp

# 7.9 Suffix Array

```
vector<int> sort_cyclic_shifts(string const& s) {
      int n = s.size();
      const int alphabet = 256;
      vector<int> p(n), c(n), cnt(max(alphabet, n), 0);
      for (int i = 0; i < n; i++) cnt[s[i]]++;</pre>
      for (int i = 1; i < alphabet; i++) cnt[i] += cnt[</pre>
           i - 11;
      for (int i = 0; i < n; i++) p[--cnt[s[i]]] = i;</pre>
      c[p[0]] = 0;
      int classes = 1;
      for (int i = 1; i < n; i++) {</pre>
        if (s[p[i]] != s[p[i-1]]) classes++;
        c[p[i]] = classes - 1;
      vector<int> pn(n), cn(n);
      for (int h = 0; (1 << h) < n; ++h) {
        for (int i = 0; i < n; i++) {</pre>
          pn[i] = p[i] - (1 << h);
          if (pn[i] < 0)
19
            pn[i] += n;
21
        fill(cnt.begin(), cnt.begin() + classes, 0);
2.2.
        for (int i = 0; i < n; i++) cnt[c[pn[i]]]++;</pre>
2.3
        for (int i = 1; i < classes; i++) cnt[i] += cnt</pre>
             [i - 1];
        for (int i = n-1; i >= 0; i--) p[--cnt[c[pn[i
             ]]]] = pn[i];
        cn[p[0]] = 0;
        classes = 1;
        for (int i = 1; i < n; i++) {</pre>
2.8
          pair<int, int> cur = {c[p[i]], c[(p[i] + (1
               << h)) % n]};
          pair < int, int > prev = {c[p[i-1]], c[(p[i-1]] +
                (1 << h)) % n]};
          if (cur != prev) ++classes;
          cn[p[i]] = classes - 1;
        c.swap(cn);
34
      return p;
    vector<int> build_suff_arr(string s) {
      s += "$";
39
      vector<int> sorted_shifts = sort_cyclic_shifts(s)
      sorted_shifts.erase(sorted_shifts.begin());
41
      return sorted shifts;
   // compare two substrings
44 int compare(int i, int j, int 1, int k) {
```

# 7.10 Z Function

# 8 References

# 8.1 Cheatsheet

();

```
1 st.insert(4);
2 st.erase(4);
    st.empty();
    // permutations
    do {
    for (int num : nums) {
    cout << num << " ";
 8
    }
    cout << endl;
   } while (next_permutation(nums.begin(), nums.end())
11 // bitset
12 int num = 27; // Binary representation: 11011
13 bitset<10> s(string("0010011010")); // from right
14 bitset<sizeof(int) * 8> bits(num);
15 int setBits = bits.count();
16 // sort
17 sort(v.begin(), v.end());
18 sort(v.rbegin(),v.rend());
19 // custom sort
20 bool comp(string a, string b) {
21 if (a.size() != b.size()) return a.size() < b.size
```

```
22 return a < b;
24 sort(v.begin(), v.end(), comp);
   // binary search
26 int a = 0, b = n-1;
27 while (a <= b) { int k = (a+b)/2; if (array[k] == x)
       ) {
28 // x found at index k
29 } if (array[k] > x) b = k-1; else a = k+1;}
30 // iterator
31 for (auto it = s.begin(); it != s.end(); it++) {
32 cout << *it << "\n";
33 }
34 // hamming distance
35 int hamming(int a, int b) {
36 return __builtin_popcount(a^b);
37 }
38 // custom comparator for pq
39 class Compare {
   public:
    bool operator()(T a, T b){
   if(cond) return true; // do not swap
    return false;
45 };
46 priority_queue<PII, vector<PII>, Compare> ds;
```

Pegaraw Pegaraw

a( ) O( ( ))	1.07			
f(n) = O(g(n))	iff $\exists$ positive $c, n_0$ such that $0 \le f(n) \le cg(n) \ \forall n \ge n_0$ .	$\sum_{i=1}^{n} i = \frac{n(n+1)}{2},  \sum_{i=1}^{n} i^2 = \frac{n(n+1)(2n+1)}{6},  \sum_{i=1}^{n} i^3 = \frac{n^2(n+1)^2}{4}.$		
$f(n) = \Omega(g(n))$	iff $\exists$ positive $c, n_0$ such that $f(n) \ge cg(n) \ge 0 \ \forall n \ge n_0$ .	$ \begin{array}{ccc}                                   $		
$f(n) = \Theta(g(n))$	iff $f(n) = O(g(n))$ and $f(n) = \Omega(g(n))$ .	$\sum_{i=1}^{n} i^{m} = \frac{1}{m+1} \left[ (n+1)^{m+1} - 1 - \sum_{i=1}^{n} \left( (i+1)^{m+1} - i^{m+1} - (m+1)i^{m} \right) \right]$		
f(n) = o(g(n))	iff $\lim_{n\to\infty} f(n)/g(n) = 0$ .	$\sum_{i=1}^{n-1} i^m = \frac{1}{m+1} \sum_{k=0}^m {m+1 \choose k} B_k n^{m+1-k}.$		
$\lim_{n \to \infty} a_n = a$	iff $\forall \epsilon > 0$ , $\exists n_0$ such that $ a_n - a  < \epsilon$ , $\forall n \ge n_0$ .	Geometric series:		
$\sup S$	least $b \in \mathbb{R}$ such that $b \ge s$ , $\forall s \in S$ .	$\sum_{i=0}^{n} c^{i} = \frac{c^{n+1} - 1}{c - 1},  c \neq 1,  \sum_{i=0}^{\infty} c^{i} = \frac{1}{1 - c},  \sum_{i=1}^{\infty} c^{i} = \frac{c}{1 - c},   c  < 1,$		
$\inf S$	greatest $b \in \mathbb{R}$ such that $b \le s$ , $\forall s \in S$ .	$\sum_{i=0}^{n} ic^{i} = \frac{nc^{n+2} - (n+1)c^{n+1} + c}{(c-1)^{2}},  c \neq 1,  \sum_{i=0}^{\infty} ic^{i} = \frac{c}{(1-c)^{2}},   c  < 1.$		
$ \liminf_{n \to \infty} a_n $	$\lim_{n \to \infty} \inf \{ a_i \mid i \ge n, i \in \mathbb{N} \}.$	Harmonic series: $n   1   \sum_{n=1}^{n} 1   \sum_{n=1}^{n} n(n+1)   n(n-1)$		
$\limsup_{n \to \infty} a_n$	$\lim_{n \to \infty} \sup \{ a_i \mid i \ge n, i \in \mathbb{N} \}.$	$H_n = \sum_{i=1}^n \frac{1}{i}, \qquad \sum_{i=1}^n iH_i = \frac{n(n+1)}{2}H_n - \frac{n(n-1)}{4}.$		
$\binom{n}{k}$	Combinations: Size $k$ subsets of a size $n$ set.	$\sum_{i=1}^{n} H_i = (n+1)H_n - n,  \sum_{i=1}^{n} {i \choose m} H_i = {n+1 \choose m+1} \left( H_{n+1} - \frac{1}{m+1} \right).$		
$\begin{bmatrix} n \\ k \end{bmatrix}$	Stirling numbers (1st kind): Arrangements of an $n$ element set into $k$ cycles.	<b>1.</b> $\binom{n}{k} = \frac{n!}{(n-k)!k!}$ , <b>2.</b> $\sum_{k=0}^{n} \binom{n}{k} = 2^{n}$ , <b>3.</b> $\binom{n}{k} = \binom{n}{n-k}$ ,		
$\binom{n}{k}$	Stirling numbers (2nd kind): Partitions of an <i>n</i> element	$4.  \binom{n}{k} = \frac{n}{k} \binom{n-1}{k-1}, \qquad \qquad 5.  \binom{n}{k} = \binom{n-1}{k} + \binom{n-1}{k-1},$		
	set into $k$ non-empty sets.	<b>6.</b> $\binom{n}{m}\binom{m}{k} = \binom{n}{k}\binom{n-k}{m-k},$ <b>7.</b> $\sum_{k=0}^{n} \binom{r+k}{k} = \binom{r+n+1}{n},$		
$\left\langle {n\atop k} \right\rangle$	1st order Eulerian numbers: Permutations $\pi_1\pi_2\pi_n$ on $\{1, 2,, n\}$ with $k$ ascents.	<b>8.</b> $\sum_{k=0}^{n} {k \choose m} = {n+1 \choose m+1},$ <b>9.</b> $\sum_{k=0}^{n} {r \choose k} {s \choose n-k} = {r+s \choose n},$		
$\left\langle\!\left\langle {n\atop k}\right\rangle\!\right\rangle$	2nd order Eulerian numbers.	<b>10.</b> $\binom{n}{k} = (-1)^k \binom{k-n-1}{k}$ , <b>11.</b> $\binom{n}{1} = \binom{n}{n} = 1$ ,		
$C_n$	Catalan Numbers: Binary trees with $n+1$ vertices.	<b>12.</b> $\binom{n}{2} = 2^{n-1} - 1,$ <b>13.</b> $\binom{n}{k} = k \binom{n-1}{k} + \binom{n-1}{k-1},$		
<b>14.</b> $\begin{bmatrix} n \\ 1 \end{bmatrix} = (n-1)!,$ <b>15.</b> $\begin{bmatrix} n \\ 2 \end{bmatrix} = (n-1)!H_{n-1},$ <b>16.</b> $\begin{bmatrix} n \\ n \end{bmatrix} = 1,$ <b>17.</b> $\begin{bmatrix} n \\ k \end{bmatrix} \ge \begin{Bmatrix} n \\ k \end{Bmatrix},$				
<b>18.</b> $\begin{bmatrix} n \\ k \end{bmatrix} = (n-1) \begin{bmatrix} n-1 \\ k \end{bmatrix} + \begin{bmatrix} n-1 \\ k-1 \end{bmatrix},$ <b>19.</b> $\begin{Bmatrix} n \\ n-1 \end{Bmatrix} = \begin{bmatrix} n \\ n-1 \end{bmatrix} = \binom{n}{2},$ <b>20.</b> $\sum_{k=0}^{n} \binom{n}{k} = n!,$ <b>21.</b> $C_n = \frac{1}{n+1} \binom{2n}{n},$				
$22. \  \left\langle {n \atop 0} \right\rangle = \left\langle {n \atop n-1} \right\rangle = 1, \qquad 23. \  \left\langle {n \atop k} \right\rangle = \left\langle {n \atop n-1-k} \right\rangle, \qquad 24. \  \left\langle {n \atop k} \right\rangle = (k+1) \left\langle {n-1 \atop k} \right\rangle + (n-k) \left\langle {n-1 \atop k-1} \right\rangle,$				
$25. \  \left\langle \begin{array}{c} 0 \\ k \end{array} \right\rangle = \left\{ \begin{array}{c} 1 & \text{if } k = 0, \\ 0 & \text{otherwise} \end{array} \right. $ $26. \ \left\langle \begin{array}{c} n \\ 1 \end{array} \right\rangle = 2^n - n - 1, $ $27. \ \left\langle \begin{array}{c} n \\ 2 \end{array} \right\rangle = 3^n - (n+1)2^n + \binom{n+1}{2}$				
$25. \  \left\langle \begin{array}{c} 0 \\ k \end{array} \right\rangle = \left\{ \begin{array}{c} 1 & \text{if } k = 0, \\ 0 & \text{otherwise} \end{array} \right. $ $26. \  \left\langle \begin{array}{c} n \\ 1 \end{array} \right\rangle = 2^n - n - 1, $ $27. \  \left\langle \begin{array}{c} n \\ 2 \end{array} \right\rangle = 3^n - (n+1)2^n + \binom{n+1}{2} $ $28. \  x^n = \sum_{k=0}^n \left\langle \begin{array}{c} n \\ k \end{array} \right\rangle \binom{x+k}{n}, $ $29. \  \left\langle \begin{array}{c} n \\ m \end{array} \right\rangle = \sum_{k=0}^m \binom{n+1}{k} (m+1-k)^n (-1)^k, $ $30. \  m! \left\{ \begin{array}{c} n \\ m \end{array} \right\} = \sum_{k=0}^n \left\langle \begin{array}{c} n \\ k \end{array} \right\rangle \binom{k}{n-m} $				
$31. \ \left\langle {n \atop m} \right\rangle = \sum_{k=0}^n \left\{ {n \atop k} \right\} {n-k \choose m} (-1)^{n-k-m} k!, \qquad \qquad 32. \ \left\langle {n \atop 0} \right\rangle = 1, \qquad \qquad 33. \ \left\langle {n \atop n} \right\rangle = 0  \text{for } n \neq 0$				
$34.  \left\langle\!\!\left\langle {n\atop k} \right\rangle\!\!\right\rangle = (k + 1)^n$	$+1$ $\binom{n-1}{k}$ $+(2n-1-k)$ $\binom{n-1}{k}$			
$36. \left\{ \begin{array}{c} x \\ x-n \end{array} \right\} = \sum_{k}^{n} \left\{ \begin{array}{c} x \\ x \end{array} \right\}$	$\sum_{k=0}^{n} \left\langle \!\! \left\langle \!\! \left\langle n \right\rangle \!\! \right\rangle \!\! \left( \!\! \left( \!\! \begin{array}{c} x+n-1-k \\ 2n \end{array} \!\! \right) \!\! \right. \!\! \right.$	37. ${n+1 \choose m+1} = \sum_{k} {n \choose k} {k \choose m} = \sum_{k=0}^{n} {k \choose m} (m+1)^{n-k},$		

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The Chinese remainder theorem: There exists a number C such that:

 $C \equiv r_1 \mod m_1$ 

: : :

 $C \equiv r_n \mod m_n$ 

if  $m_i$  and  $m_j$  are relatively prime for  $i \neq j$ . Euler's function:  $\phi(x)$  is the number of positive integers less than x relatively prime to x. If  $\prod_{i=1}^{n} p_i^{e_i}$  is the prime factorization of x then

$$\phi(x) = \prod_{i=1}^{n} p_i^{e_i - 1} (p_i - 1).$$

Euler's theorem: If a and b are relatively prime then

$$1 \equiv a^{\phi(b)} \bmod b$$
.

Fermat's theorem:

$$1 \equiv a^{p-1} \bmod p.$$

The Euclidean algorithm: if a > b are integers then

$$gcd(a, b) = gcd(a \mod b, b).$$

If  $\prod_{i=1}^{n} p_i^{e_i}$  is the prime factorization of x

$$S(x) = \sum_{d|x} d = \prod_{i=1}^{n} \frac{p_i^{e_i+1} - 1}{p_i - 1}.$$

Perfect Numbers: x is an even perfect number iff  $x = 2^{n-1}(2^n - 1)$  and  $2^n - 1$  is prime. Wilson's theorem: n is a prime iff

$$(n-1)! \equiv -1 \mod n$$
.

$$\mu(i) = \begin{cases} (n-1)! = -1 \bmod n. \\ \text{M\"obius inversion:} \\ \mu(i) = \begin{cases} 1 & \text{if } i = 1. \\ 0 & \text{if } i \text{ is not square-free.} \\ (-1)^r & \text{if } i \text{ is the product of} \\ r & \text{distinct primes.} \end{cases}$$
 If

 $\operatorname{If}$ 

$$G(a) = \sum_{d|a} F(d),$$

$$F(a) = \sum_{d|a} \mu(d) G\left(\frac{a}{d}\right).$$

Prime numbers:

$$p_n = n \ln n + n \ln \ln n - n + n \frac{\ln \ln n}{\ln n}$$

$$+O\left(\frac{n}{\ln n}\right),$$

$$\pi(n) = \frac{n}{\ln n} + \frac{n}{(\ln n)^2} + \frac{2!n}{(\ln n)^3} + O\left(\frac{n}{(\ln n)^4}\right).$$

# Definitions:

Loop An edge connecting a vertex to itself.

Directed Each edge has a direction. SimpleGraph with no loops or multi-edges.

WalkA sequence  $v_0e_1v_1\ldots e_\ell v_\ell$ . TrailA walk with distinct edges. Pathtrail with distinct

vertices.

ConnectedA graph where there exists a path between any two

vertices.

connected ComponentΑ maximal subgraph.

TreeA connected acyclic graph. Free tree A tree with no root. DAGDirected acyclic graph. EulerianGraph with a trail visiting each edge exactly once.

Hamiltonian Graph with a cycle visiting each vertex exactly once.

CutA set of edges whose removal increases the number of components.

Cut-setA minimal cut. Cut edge A size 1 cut.

k-Connected A graph connected with the removal of any k-1vertices.

k-Tough  $\forall S \subseteq V, S \neq \emptyset$  we have  $k \cdot c(G - S) \le |S|$ .

A graph where all vertices k-Regular have degree k.

k-Factor Α k-regular spanning subgraph.

Matching A set of edges, no two of which are adjacent.

CliqueA set of vertices, all of which are adjacent.

Ind. set A set of vertices, none of which are adjacent.

Vertex cover A set of vertices which cover all edges.

Planar graph A graph which can be embeded in the plane.

Plane graph An embedding of a planar

$$\sum_{v \in V} \deg(v) = 2m.$$

If G is planar then n - m + f = 2, so

$$f \le 2n - 4, \quad m \le 3n - 6.$$

Any planar graph has a vertex with degree  $\leq 5$ .

# Notation:

E(G)Edge set Vertex set V(G)

c(G)Number of components

G[S]Induced subgraph

deg(v)Degree of v

Maximum degree  $\Delta(G)$  $\delta(G)$ Minimum degree  $\chi(G)$ Chromatic number

 $\chi_E(G)$ Edge chromatic number  $G^c$ Complement graph

 $K_n$ Complete graph  $K_{n_1,n_2}$ Complete bipartite graph

Ramsev number

# Geometry

Projective coordinates: (x, y, z), not all x, y and z zero.

$$(x, y, z) = (cx, cy, cz) \quad \forall c \neq 0.$$

Cartesian Projective (x, y)(x, y, 1)y = mx + b(m, -1, b)x = c(1,0,-c)

Distance formula,  $L_p$  and  $L_{\infty}$ 

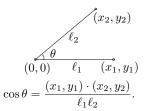
$$\sqrt{(x_1 - x_0)^2 + (y_1 - y_0)^2},$$
$$[|x_1 - x_0|^p + |y_1 - y_0|^p]^{1/p},$$

$$\lim_{p \to \infty} \left[ |x_1 - x_0|^p + |y_1 - y_0|^p \right]^{1/p}.$$

Area of triangle  $(x_0, y_0), (x_1, y_1)$ and  $(x_2, y_2)$ :

$$\frac{1}{2} \operatorname{abs} \begin{vmatrix} x_1 - x_0 & y_1 - y_0 \\ x_2 - x_0 & y_2 - y_0 \end{vmatrix}.$$

Angle formed by three points:



Line through two points  $(x_0, y_0)$ and  $(x_1, y_1)$ :

$$\begin{vmatrix} x & y & 1 \\ x_0 & y_0 & 1 \\ x_1 & y_1 & 1 \end{vmatrix} = 0.$$

Area of circle, volume of sphere:

$$A = \pi r^2, \qquad V = \frac{4}{3}\pi r^3.$$

If I have seen farther than others, it is because I have stood on the shoulders of giants.

- Issac Newton

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Taylor's series:

$$f(x) = f(a) + (x - a)f'(a) + \frac{(x - a)^2}{2}f''(a) + \dots = \sum_{i=0}^{\infty} \frac{(x - a)^i}{i!}f^{(i)}(a).$$

Expansions:

Expansions: 
$$\frac{1}{1-x} = 1 + x + x^2 + x^3 + x^4 + \cdots = \sum_{i=0}^{\infty} x^i,$$

$$\frac{1}{1-cx} = 1 + x + x^2 + x^3 + x^4 + \cdots = \sum_{i=0}^{\infty} c^i x^i,$$

$$\frac{1}{1-cx} = 1 + x + x^2 + x^3 + x^4 + \cdots = \sum_{i=0}^{\infty} c^i x^i,$$

$$\frac{1}{1-x^n} = 1 + x^n + x^{2n} + x^{3n} + \cdots = \sum_{i=0}^{\infty} ix^{ii},$$

$$\frac{x}{(1-x)^2} = x + 2x^2 + 3x^3 + 4x^4 + \cdots = \sum_{i=0}^{\infty} ix^i,$$

$$x^k \frac{d^n}{dx^n} \left(\frac{1}{1-x}\right) = x + 2^{n}x^2 + 3^n x^3 + 4^n x^4 + \cdots = \sum_{i=0}^{\infty} i^n x^i,$$

$$e^x = 1 + x + \frac{1}{2}x^2 + \frac{1}{6}x^3 + \cdots = \sum_{i=0}^{\infty} i^n x^i,$$

$$\ln(1+x) = x - \frac{1}{2}x^2 + \frac{1}{3}x^3 - \frac{1}{4}x^4 - \cdots = \sum_{i=0}^{\infty} (-1)^{i+1} \frac{x^i}{i},$$

$$\ln \frac{1}{1-x} = x + \frac{1}{2}x^2 + \frac{1}{3}x^3 + \frac{1}{4}x^4 + \cdots = \sum_{i=0}^{\infty} (-1)^{i} \frac{x^{2i+1}}{(2i+1)!},$$

$$\cos x = 1 - \frac{1}{2!}x^2 + \frac{1}{1!}x^4 - \frac{1}{6!}x^6 + \cdots = \sum_{i=0}^{\infty} (-1)^{i} \frac{x^{2i+1}}{(2i+1)!},$$

$$(1+x)^n = 1 + nx + \frac{n(n-1)}{2}x^2 + \cdots = \sum_{i=0}^{\infty} (-1)^{i} \frac{x^{2i+1}}{(2i+1)!},$$

$$(1+x)^n = 1 + nx + \frac{n(n-1)}{2}x^2 + \cdots = \sum_{i=0}^{\infty} (-1)^{i} \frac{x^{2i+1}}{(2i+1)!},$$

$$\frac{1}{(1-x)^{n+1}} = 1 + (n+1)x + (\frac{n+2}{2})x^2 + \cdots = \sum_{i=0}^{\infty} (-1)^{i} \frac{x^{2i+1}}{(i)},$$

$$\frac{x}{e^x - 1} = 1 - \frac{1}{2}x + \frac{1}{12}x^2 - \frac{1}{126}x^4 + \cdots = \sum_{i=0}^{\infty} (\frac{i+n}{i})x^i,$$

$$\frac{1}{\sqrt{1-4x}} = 1 + x + 2x^2 + 5x^3 + \cdots = \sum_{i=0}^{\infty} (\frac{2i}{i})x^i,$$

$$\frac{1}{\sqrt{1-4x}} = 1 + x + 2x^2 + 6x^3 + \cdots = \sum_{i=0}^{\infty} (\frac{2i}{i})x^i,$$

$$\frac{1}{1-x} \ln \frac{1}{1-x} = x + \frac{3}{2}x^2 + \frac{1}{16}x^3 + \frac{25}{24}x^4 + \cdots = \sum_{i=0}^{\infty} H_i x^i,$$

$$\frac{1}{1-x} \ln \frac{1}{1-x} = x + \frac{3}{2}x^2 + \frac{1}{16}x^3 + \frac{25}{24}x^4 + \cdots = \sum_{i=0}^{\infty} H_i x^i,$$

$$\frac{1}{2} \left( \ln \frac{1}{1-x} \right)^2 = \frac{1}{2}x^2 + \frac{3}{4}x^3 + \frac{11}{24}x^4 + \cdots = \sum_{i=0}^{\infty} F_{ii}x^i.$$

$$\frac{x}{1-x-x^2} = x + x^2 + 2x^3 + 3x^4 + \cdots = \sum_{i=0}^{\infty} F_{ii}x^i.$$

Ordinary power series:

$$A(x) = \sum_{i=0}^{\infty} a_i x^i.$$

Exponential power series:

$$A(x) = \sum_{i=0}^{\infty} a_i \frac{x^i}{i!}.$$

Dirichlet power series:

$$A(x) = \sum_{i=1}^{\infty} \frac{a_i}{i^x}.$$

Binomial theorem

$$(x+y)^n = \sum_{k=0}^n \binom{n}{k} x^{n-k} y^k.$$

$$x^{n} - y^{n} = (x - y) \sum_{k=0}^{n-1} x^{n-1-k} y^{k}.$$

For ordinary power se

$$\alpha A(x) + \beta B(x) = \sum_{i=0}^{\infty} (\alpha a_i + \beta b_i) x^i,$$

$$x^k A(x) = \sum_{i=k}^{\infty} a_{i-k} x^i,$$

$$\frac{A(x) - \sum_{i=0}^{k-1} a_i x^i}{x^k} = \sum_{i=0}^{\infty} a_{i+k} x^i,$$

$$A(cx) = \sum_{i=0}^{\infty} c^i a_i x^i,$$

$$A'(x) = \sum_{i=0}^{\infty} (i+1)a_{i+1}x^{i},$$

$$xA'(x) = \sum_{i=1}^{\infty} ia_i x^i,$$

$$\int A(x) dx = \sum_{i=1}^{\infty} \frac{a_{i-1}}{i} x^i,$$

$$\frac{A(x) + A(-x)}{2} = \sum_{i=1}^{\infty} a_{2i} x^{2i},$$

$$\frac{A(x) - A(-x)}{2} = \sum_{i=0}^{\infty} a_{2i+1} x^{2i+1}.$$

Summation: If  $b_i = \sum_{i=0}^i a_i$  then

$$B(x) = \frac{1}{1-x}A(x).$$

Convolution:

$$A(x)B(x) = \sum_{i=0}^{\infty} \left(\sum_{j=0}^{i} a_j b_{i-j}\right) x^i.$$

God made the natural numbers; all the rest is the work of man. Leopold Kronecker