UPLB - Pegaraw Notebook

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1 Strings
1.1 Manacher
1 vector<int> manacher_odd(string s) {
  int n = s.size();
  s = "$" + s + "^";
  vector<int> p(n + 2);
  int 1 = 1, r = 1;
   for(int i = 1; i <= n; i++) {</pre>
   p[i] = max(0, min(r - i, p[1 + (r - i)]));
   while (s[i - p[i]] == s[i + p[i]]) p[i]++;
9
   if(i + p[i] > r) l = i - p[i], r = i + p[i];
  return vector<int>(begin(p) + 1, end(p) - 1);
12
13 vector<int> manacher(string s) {
  string t:
   for(auto c: s) t += string("#") + c;
  auto res = manacher_odd(t + "#");
  return vector<int>(begin(res) + 1, end(res) - 1);
```

18 }

1.2 Hashing

1 11 compute_hash(string const& s) {

const 11 p = 31, m = 1e9 + 9;

```
11 hash_value = 0, p_pow = 1;
11
11
          for (char c : s) {
11
           hash\_value = (hash\_value + (c - 'a' + 1) *
12
                p_pow) % m;
12
           p_pow = (p_pow * p) % m;
          return hash_value;
13
13
13
    1.3 Rabin Karp
13
      vector<ll> rabin_karp(string const& s, string const
13
            & t) {
          const 11 p = 31, m = 1e9 + 9;
          11 S = s.size(), T = t.size();
14
          vector<ll> p_pow(max(S, T));
14
          p pow[0] = 1;
14
          for (11 i = 1; i < (11) p_pow.size(); i++) p_pow[</pre>
               i] = (p_pow[i-1] * p) % m;
          vector<ll> h(T + 1, 0);
15
          for (11 i = 0; i < T; i++) h[i + 1] = (h[i] + (t[
15
               i] - 'a' + 1) * p_pow[i]) % m;
15
          11 h_s = 0;
          for (11 i = 0; i < S; i++) h_s = (h_s + (s[i] - '
16
              a' + 1) * p_pow[i]) % m;
16
          vector<11> occurences;
          for (11 i = 0; i + S - 1 < T; i++) {
            11 cur h = (h[i + S] + m - h[i]) % m;
17
            if (cur_h == h_s * p_pow[i] % m) occurences.
17
                push_back(i);
          return occurences;
18
    1.4 Z Function
    1 vector<int> z_function(string s) {
          int n = s.size();
          vector<int> z(n);
          for (int i = 1, l = 0, r = 0; i < n; i++) {
           if (i < r) z[i] = min(r - i, z[i - 1]);
           while (i + z[i] < n \&\& s[z[i]] == s[i + z[i]])
                z[i]++;
           if (i + z[i] > r) {
             1 = i;
              r = i + z[i];
          return z;
   13 }
    1.5 Finding Repetitions
       vector<int> z_function(string const& s) {
         int n = s.size();
          vector<int> z(n);
          for (int i = 1, l = 0, r = 0; i < n; i++) {
           if (i \le r) z[i] = min(r - i + 1, z[i - 1]);
            while (i + z[i] < n \&\& s[z[i]] == s[i + z[i]])
                z[i]++;
            if (i + z[i] - 1 > r) {
             1 = i;
```

```
r = i + z[i] - 1;
        }
11
      return z;
13 }
14 int get_z(vector<int> const& z, int i) {
15
      if (0 <= i && i < (int) z.size()) return z[i];</pre>
      else return 0:
17 }
18 vector<pair<int, int>> repetitions;
    void convert_to_repetitions(int shift, bool left,
         int cntr, int 1, int k1, int k2) {
      for (int 11 = \max(1, 1 - k2); 11 \le \min(1, k1);
        if (left && 11 == 1) break;
22
        int 12 = 1 - 11;
23
        int pos = shift + (left ? cntr - 11 : cntr - 1
             -11+1);
24
        repetitions.emplace_back(pos, pos + 2 * 1 - 1);
25
26 }
27 void find_repetitions(string s, int shift = 0) {
      int n = s.size();
29
      if (n == 1) return;
     int nu = n / 2;
      int nv = n - nu;
      string u = s.substr(0, nu);
33
      string v = s.substr(nu);
      string ru(u.rbegin(), u.rend());
      string rv(v.rbegin(), v.rend());
      find repetitions (u, shift);
      find_repetitions(v, shift + nu);
38
      vector<int> z1 = z_function(ru);
39
      vector<int> z2 = z function(v + '#' + u);
40
      vector<int> z3 = z_function(ru + '#' + rv);
41
      vector<int> z4 = z_function(v);
42
      for (int cntr = 0; cntr < n; cntr++) {</pre>
43
       int 1, k1, k2;
44
        if (cntr < nu) {</pre>
45
         1 = nu - cntr;
46
          k1 = get_z(z1, nu - cntr);
47
          k2 = get_z(z2, nv + 1 + cntr);
48
         } else {
49
          1 = cntr - nu + 1;
50
          k1 = get_z(z3, nu + 1 + nv - 1 - (cntr - nu))
51
          k2 = get_z(z4, (cntr - nu) + 1);
52
        if (k1 + k2 >= 1) convert_to_repetitions(shift,
              cntr < nu, cntr, 1, k1, k2);</pre>
54
55 }
```

1.6 Longest Common Prefix

vector<int> sort_cyclic_shifts(string const& s) {

1.7 Suffix Array

int n = s.size();

```
const int alphabet = 256;
      vector<int> p(n), c(n), cnt(max(alphabet, n), 0);
      for (int i = 0; i < n; i++) cnt[s[i]]++;</pre>
      for (int i = 1; i < alphabet; i++) cnt[i] += cnt[</pre>
           i - 1];
      for (int i = 0; i < n; i++) p[--cnt[s[i]]] = i;</pre>
      c[p[0]] = 0;
      int classes = 1;
      for (int i = 1; i < n; i++) {</pre>
        if (s[p[i]] != s[p[i-1]]) classes++;
        c[p[i]] = classes - 1;
13
      vector<int> pn(n), cn(n);
      for (int h = 0; (1 << h) < n; ++h) {
        for (int i = 0; i < n; i++) {</pre>
          pn[i] = p[i] - (1 << h);
18
          if (pn[i] < 0)
            pn[i] += n;
        fill(cnt.begin(), cnt.begin() + classes, 0);
        for (int i = 0; i < n; i++) cnt[c[pn[i]]]++;</pre>
        for (int i = 1; i < classes; i++) cnt[i] += cnt</pre>
             [i - 1];
        for (int i = n-1; i >= 0; i--) p[--cnt[c[pn[i
             ]]]] = pn[i];
        cn[p[0]] = 0;
2.6
        classes = 1;
        for (int i = 1; i < n; i++) {
          pair<int, int> cur = {c[p[i]], c[(p[i] + (1)
               << h)) % n]};
          pair < int, int > prev = {c[p[i-1]], c[(p[i-1]] + }
                (1 << h)) % n]};
          if (cur != prev) ++classes;
          cn[p[i]] = classes - 1;
32
33
        c.swap(cn);
34
      }
      return p;
    vector<int> build_suff_arr(string s) {
      vector<int> sorted_shifts = sort_cyclic_shifts(s)
      sorted_shifts.erase(sorted_shifts.begin());
41
      return sorted_shifts;
    // compare two substrings
44 int compare(int i, int j, int l, int k) {
      pair<int, int> a = {c[k][i], c[k][(i + 1 - (1 <<
           k)) % n]};
      pair<int, int> b = \{c[k][j], c[k][(j + 1 - (1 <<
           k)) % n]};
```

return a == b ? 0 : a < b ? -1 : 1;

48 }

1.8 Count Unique Substrings

```
int count_unique_substrings(string const& s) {
      int n = s.size();
      const int p = 31;
      const int m = 1e9 + 9;
      vector<long long> p_pow(n);
      p_pow[0] = 1;
      for (int i = 1; i < n; i++) p_pow[i] = (p_pow[i -</pre>
            1] * p) % m;
      vector<long long> h(n + 1, 0);
      for (int i = 0; i < n; i++) h[i + 1] = (h[i] + (s)
           [i] - 'a' + 1) * p_pow[i]) % m;
      int cnt = 0;
      for (int 1 = 1; 1 <= n; 1++) {</pre>
        unordered_set<long long> hs;
        for (int i = 0; i <= n - 1; i++) {</pre>
          long long cur_h = (h[i + 1] + m - h[i]) % m;
          cur_h = (cur_h * p_pow[n - i - 1]) % m;
          hs.insert(cur_h);
        cnt += hs.size();
      return cnt;
21
```

1.9 Knuth Morris Pratt

```
1 vector<ll> prefix_function(string s) {
    11 n = (11) s.length();
      vector<ll> pi(n);
      for (11 i = 1; i < n; i++) {
        11 j = pi[i - 1];
        while (j > 0 \&\& s[i] != s[j]) j = pi[j - 1];
        if (s[i] == s[j]) j++;
        pi[i] = j;
9
      return pi;
    // count occurences
13 vector < int > ans(n + 1);
    for (int i = 0; i < n; i++)</pre>
     ans[pi[i]]++;
    for (int i = n-1; i > 0; i--)
     ans[pi[i-1]] += ans[i];
    for (int i = 0; i <= n; i++)</pre>
      ans[i]++;
```

1.10 Group Identical Substrings

```
vector<vector<int>> group_identical_strings(vector
string> const& s) {

int n = s.size();

vector<pair<long long, int>> hashes(n);

for (int i = 0; i < n; i++) hashes[i] = {
    compute_hash(s[i]), i};

sort(hashes.begin(), hashes.end());

vector<vector<int>> groups;

for (int i = 0; i < n; i++) {

if (i == 0 || hashes[i].first != hashes[i - 1].
    first) groups.emplace_back();

groups.back().push_back(hashes[i].second);
}
</pre>
```

```
11 return groups;
12 }
```

2 Geometry

2.1 Nearest Points

```
struct pt {
      11 x, y, id;
    struct cmp_x {
      bool operator()(const pt & a, const pt & b) const
        return a.x < b.x || (a.x == b.x && a.y < b.y);</pre>
 8 };
    struct cmp_y {
      bool operator()(const pt & a, const pt & b) const
            { return a.y < b.y; }</pre>
11 };
12
    11 n;
13 vector<pt> a;
14 double mindist;
15 pair<11, 11> best_pair;
16 void upd_ans(const pt & a, const pt & b) {
17
      double dist = sqrt((a.x - b.x) * (a.x - b.x) + (a.x - b.x)
           .y - b.y) * (a.y - b.y);
      if (dist < mindist) {</pre>
        mindist = dist;
20
        best_pair = {a.id, b.id};
21
23
    vector<pt> t;
24
    void rec(ll l, ll r) {
25
      if (r - 1 <= 3) {
26
        for (ll i = l; i < r; ++i)
27
          for (11 j = i + 1; j < r; ++j)
            upd_ans(a[i], a[j]);
        sort(a.begin() + 1, a.begin() + r, cmp_y());
        return:
31
      11 m = (1 + r) >> 1, midx = a[m].x;
33
      rec(1, m);
34
      merge(a.begin() + 1, a.begin() + m, a.begin() + m
           , a.begin() + r, t.begin(), cmp_y());
      copy(t.begin(), t.begin() + r - 1, a.begin() + 1)
      11 \text{ tsz} = 0;
38
      for (11 i = 1; i < r; ++i) {
39
        if (abs(a[i].x - midx) < mindist) {</pre>
40
          for (11 j = tsz - 1; j >= 0 && a[i].y - t[j].
               y < mindist; --j)
             upd_ans(a[i], t[j]);
42
          t[tsz++] = a[i];
43
44
45 }
46 t.resize(n);
    sort(a.begin(), a.end(), cmp_x());
48 mindist = 1E20;
49 rec(0, n);
```

```
struct pt {
      11 x, y;
      pt operator + (const pt & p) const { return pt {x}
            + p.x, y + p.y; }
      pt operator - (const pt & p) const { return pt{x
           - p.x, y - p.y; }
      11 cross(const pt & p) const { return x * p.y - y
            * p.x; }
6 };
    void reorder_polygon(vector<pt> & P) {
8
      size_t pos = 0;
      for (size_t i = 1; i < P.size(); i++){</pre>
        if (P[i].y < P[pos].y || (P[i].y == P[pos].y &&
              P[i].x < P[pos].x)) pos = i;
      rotate(P.begin(), P.begin() + pos, P.end());
13
    vector<pt> minkowski(vector<pt> P, vector<pt> Q) {
      // the first vertex must be the lowest
      reorder_polygon(P);
      reorder_polygon(Q);
      // we must ensure cyclic indexing
      P.push back(P[0]);
      P.push_back(P[1]);
      Q.push_back(Q[0]);
      Q.push_back(Q[1]);
      // main part
      vector<pt> result;
      size_t i = 0, j = 0;
      while (i < P.size() - 2 || j < Q.size() - 2){</pre>
        result.push back(P[i] + O[j]);
        auto cross = (P[i + 1] - P[i]).cross(Q[j + 1] -
              Q[j]);
        if (cross >= 0 && i < P.size() - 2) ++i;</pre>
30
        if (cross <= 0 && j < Q.size() - 2) ++j;</pre>
31
      return result;
33
```

2.3 Point In Convex

```
struct pt {
      long long x, y;
      pt() {}
      pt (long long \underline{x}, long long \underline{y}) : x(\underline{x}), y(\underline{y}) {}
      pt operator+(const pt &p) const { return pt(x + p
            .x, y + p.y);
      pt operator-(const pt &p) const { return pt(x - p
            .x, y - p.y);
      long long cross(const pt &p) const { return x * p
           y - y * p.x;
      long long dot(const pt &p) const { return x * p.x
            + y * p.y; }
      long long cross(const pt &a, const pt &b) const {
            return (a - *this).cross(b - *this); }
      long long dot (const pt &a, const pt &b) const {
           return (a - *this).dot(b - *this); }
      long long sqrLen() const { return this->dot(*this
12 };
13 bool lexComp(const pt &1, const pt &r) { return l.x
          < r.x \mid | (1.x == r.x && 1.y < r.y); }
14 int sgn(long long val) { return val > 0 ? 1 : (val
         == 0 ? 0 : -1); }
    vector<pt> seq;
16 pt translation;
```

```
bool pointInTriangle(pt a, pt b, pt c, pt point) {
      long long s1 = abs(a.cross(b, c));
      long long s2 = abs(point.cross(a, b)) + abs(point
           .cross(b, c)) + abs(point.cross(c, a));
      return s1 == s2;
    void prepare(vector<pt> &points) {
     n = points.size();
25
      int pos = 0;
      for (int i = 1; i < n; i++) {
        if (lexComp(points[i], points[pos])) pos = i;
      rotate(points.begin(), points.begin() + pos,
      n--;
      seq.resize(n);
      for (int i = 0; i < n; i++) seq[i] = points[i +</pre>
          1] - points[0];
      translation = points[0];
    bool pointInConvexPolygon(pt point) {
      point = point - translation;
      if (seq[0].cross(point) != 0 && sgn(seq[0].cross(
           point)) != sqn(seq[0].cross(seq[n-1])))
38
        return false;
39
      if (seq[n-1].cross(point) != 0 \&\& sgn(seq[n-1])
           1].cross(point)) != sqn(seq[n - 1].cross(seq
           [0])))
        return false;
      if (seq[0].cross(point) == 0)
        return seq[0].sqrLen() >= point.sqrLen();
      int 1 = 0, r = n - 1;
      while (r - 1 > 1) {
       int mid = (1 + r) / 2;
        int pos = mid;
        if (seq[pos].cross(point) >= 0) 1 = mid;
        else r = mid;
49
50
      int pos = 1;
      return pointInTriangle(seq[pos], seq[pos + 1], pt
           (0, 0), point);
```

2.4 Line Sweep

```
const double EPS = 1E-9;
    struct pt { double x, y; };
    struct seg {
      pt p, q;
      11 id:
      double get_y (double x) const {
        if (abs(p.x - q.x) < EPS) return p.y;</pre>
        return p.y + (q.y - p.y) * (x - p.x) / (q.x - p.x)
             .x):
    bool intersect1d(double 11, double r1, double 12,
         double r2) {
      if (11 > r1) swap(11, r1);
      if (12 > r2) swap(12, r2);
      return max(11, 12) <= min(r1, r2) + EPS;</pre>
16 11 vec(const pt& a, const pt& b, const pt& c) {
      double s = (b.x - a.x) * (c.y - a.y) - (b.y - a.y)
           ) * (c.x - a.x);
      return abs(s) < EPS ? 0 : s > 0 ? +1 : -1;
19
```

```
20 bool intersect(const seg& a, const seg& b) {
21
      return intersect1d(a.p.x, a.g.x, b.p.x, b.g.x) &&
             intersect1d(a.p.y, a.g.y, b.p.y, b.g.y) &&
23
             vec(a.p, a.q, b.p) * vec(a.p, a.q, b.q) <=
                   33 0
24
             vec(b.p, b.q, a.p) * vec(b.p, b.q, a.q) <=
                   0:
25 1
26 bool operator<(const seg& a, const seg& b) {</pre>
      double x = max(min(a.p.x, a.q.x), min(b.p.x, b.q.
      return a.get_y(x) < b.get_y(x) - EPS;</pre>
29 }
30 struct event {
      double x;
      11 tp, id;
      event() {}
      event (double x, 11 tp, 11 id) : x(x), tp(tp), id(
           id) {}
      bool operator<(const event& e) const {</pre>
        if (abs(x - e.x) > EPS) return x < e.x;
37
        return tp > e.tp;
38
39 };
40 set<seg> s;
41 vector<set<seg>::iterator> where;
42 set<seg>::iterator prev(set<seg>::iterator it) {
      return it == s.begin() ? s.end() : --it;
44 }
45 set<seg>::iterator next(set<seg>::iterator it) {
     return ++it;
47 }
48 pair<11, 11> solve(const vector<seg>& a) {
     11 n = (11) a.size();
      vector<event> e;
51
      for (11 i = 0; i < n; ++i) {
        e.push_back(event(min(a[i].p.x, a[i].q.x), +1,
        e.push_back(event(max(a[i].p.x, a[i].q.x), -1,
             i));
      sort(e.begin(), e.end());
      s.clear();
      where.resize(a.size());
      for (size t i = 0; i < e.size(); ++i) {</pre>
59
       11 id = e[i].id;
60
        if (e[i].tp == +1) {
61
          set<seg>::iterator nxt = s.lower_bound(a[id])
               , prv = prev(nxt);
          if (nxt != s.end() && intersect(*nxt, a[id]))
                return make_pair(nxt->id, id);
          if (prv != s.end() && intersect(*prv, a[id]))
                return make pair(prv->id, id);
64
          where[id] = s.insert(nxt, a[id]);
65
          set<seq>::iterator nxt = next(where[id]), prv
                = prev(where[id]);
67
          if (nxt != s.end() && prv != s.end() &&
               intersect(*nxt, *prv)) return make_pair(
               prv->id, nxt->id);
68
          s.erase(where[id]);
69
71
      return make_pair(-1, -1);
```

2.5 Line Intersection

```
struct pt { double x, y; };
struct line { double a, b, c; };
const double EPS = 1e-9;
double det (double a, double b, double c, double d)
    { return a*d - b*c; }
bool intersect(line m, line n, pt & res) {
  double zn = det(m.a, m.b, n.a, n.b);
  if (abs(zn) < EPS) return false;</pre>
  res.x = -det(m.c, m.b, n.c, n.b) / zn;
  res.y = -det(m.a, m.c, n.a, n.c) / zn;
  return true;
bool parallel(line m, line n) { return abs(det(m.a,
     m.b, n.a, n.b)) < EPS; }
bool equivalent(line m, line n) {
  return abs(det(m.a, m.b, n.a, n.b)) < EPS
      && abs(det(m.a, m.c, n.a, n.c)) < EPS
      && abs(det(m.b, m.c, n.b, n.c)) < EPS;
```

2.6 Basic Geometry

```
struct point2d {
      ftype x, y;
      point2d() {}
      point2d(ftype x, ftype y): x(x), y(y) {}
      point2d& operator+=(const point2d &t) {
       x += t.x;
       v += t.v:
       return *this;
      point2d& operator-=(const point2d &t) {
       x -= t.x;
       y -= t.y;
        return *this;
      point2d& operator*=(ftype t) {
       x *= t;
       v *= t;
       return *this:
      point2d& operator/=(ftype t) {
       x /= t;
        y /= t;
        return *this;
      point2d operator+(const point2d &t) const {
           return point2d(*this) += t; }
      point2d operator-(const point2d &t) const {
           return point2d(*this) -= t; }
      point2d operator*(ftype t) const { return point2d
           (*this) *= t; }
      point2d operator/(ftype t) const { return point2d
           (*this) /= t; }
30 point2d operator*(ftype a, point2d b) { return b *
31 ftype dot(point2d a, point2d b) { return a.x * b.x
         + a.y * b.y; }
32 ftype dot(point3d a, point3d b) { return a.x * b.x
         + a.y * b.y + a.z * b.z; }
33 ftype norm(point2d a) { return dot(a, a); }
34 double abs(point2d a) { return sgrt(norm(a)); }
35 double proj(point2d a, point2d b) { return dot(a, b)
         ) / abs(b); }
```

```
36 double angle(point2d a, point2d b) { return acos(
         dot(a, b) / abs(a) / abs(b)); }
37 point3d cross(point3d a, point3d b) { return
        point3d(a.y * b.z - a.z * b.y, a.z * b.x - a.x
         * b.z, a.x * b.y - a.y * b.x); }
38 ftype triple(point3d a, point3d b, point3d c) {
         return dot(a, cross(b, c)); }
39 ftype cross(point2d a, point2d b) { return a.x * b.
         y - a.y * b.x; }
40 point2d intersect(point2d al, point2d dl, point2d
         a2, point2d d2) { return a1 + cross(a2 - a1,
         d2) / cross(d1, d2) * d1; }
41 point3d intersect(point3d al, point3d nl, point3d
         a2, point3d n2, point3d a3, point3d n3) {
      point3d x(n1.x, n2.x, n3.x);
      point3d y(n1.y, n2.y, n3.y);
44
      point3d z(n1.z, n2.z, n3.z);
      point3d d(dot(a1, n1), dot(a2, n2), dot(a3, n3));
      return point3d(triple(d, y, z), triple(x, d, z),
          triple(x, y, d)) / triple(n1, n2, n3);
```

2.7 Circle Line Intersection

```
double r, a, b, c; // given as input
    double x0 = -a * c / (a * a + b * b);
    double y0 = -b * c / (a * a + b * b);
   if (c * c > r * r * (a * a + b * b) + EPS) {
     puts ("no points");
   } else if (abs (c *c - r * r * (a * a + b * b)) <</pre>
        EPS) {
      puts ("1 point");
     cout << x0 << ' ' << y0 << '\n';
   } else {
     double d = r * r - c * c / (a * a + b * b);
     double mult = sqrt (d / (a * a + b * b));
     double ax, ay, bx, by;
    ax = x0 + b * mult;
    bx = x0 - b * mult;
     ay = y0 - a * mult;
     by = y0 + a * mult;
     puts ("2 points");
      cout << ax << ' ' << ay << '\n' << bx << ' ' <<
          by << '\n';
19 1
```

2.8 Convex Hull

```
17 bool collinear(pt a, pt b, pt c) {
18
      return orientation(a, b, c) == 0;
19 }
20 void convex_hull(vector<pt>& a, bool
         include_collinear = false) {
      pt p0 = *min_element(a.begin(), a.end(), [](pt a,
        return make_pair(a.y, a.x) < make_pair(b.y, b.x</pre>
             ):
      });
      sort(a.begin(), a.end(), [&p0](const pt& a, const
25
        11 o = orientation(p0, a, b);
26
        if (o == 0) {
27
          return (p0.x - a.x) * (p0.x - a.x) + (p0.y - a.x)
               a.y) * (p0.y - a.y)
               < (p0.x - b.x) * (p0.x - b.x) + (p0.y -
                    b.y) * (p0.y - b.y);
        return o < 0;
31
32
      if (include_collinear) {
        11 i = (11) a.size()-1;
        while (i \ge 0 \&\& collinear(p0, a[i], a.back()))
        reverse(a.begin()+i+1, a.end());
36
      vector<pt> st;
      for (ll i = 0; i < (ll) a.size(); i++) {</pre>
        while (st.size() > 1 && !cw(st[st.size() - 2],
             st.back(), a[i], include collinear)) {
40
          st.pop_back();
41
42
        st.push back(a[i]);
43
44
45 }
```

2.9 Count Lattices

```
1 int count_lattices(Fraction k, Fraction b, long
         long n) {
      auto fk = k.floor();
      auto fb = b.floor();
      auto cnt = 0LL;
      if (k >= 1 || b >= 1) {
       cnt += (fk * (n - 1) + 2 * fb) * n / 2;
        k -= fk:
        b -= fb;
      auto t = k * n + b;
      auto ft = t.floor();
      if (ft >= 1) cnt += count_lattices(1 / k, (t - t.
           floor()) / k, t.floor());
1.3
      return cnt;
14 }
```

2.10 Segment Intersection

```
};
    struct line {
      double a, b, c;
      line() {}
      line(pt p, pt q) {
        a = p.y - q.y;
        b = q.x - p.x;
        c = -a * p.x - b * p.y;
        norm();
      void norm() {
18
        double z = sqrt(a * a + b * b);
19
        if (abs(z) > EPS) a /= z, b /= z, c /= z;
      double dist(pt p) const { return a * p.x + b * p.
          y + c; }
    double det(double a, double b, double c, double d)
      return a * d - b * c;
    inline bool betw(double 1, double r, double x) {
      return min(l, r) \le x + EPS \&\& x \le max(l, r) +
28
    inline bool intersect_1d(double a, double b, double
         c, double d) {
      if (a > b) swap(a, b);
      if (c > d) swap(c, d);
      return max(a, c) <= min(b, d) + EPS;</pre>
33
34
    bool intersect(pt a, pt b, pt c, pt d, pt& left, pt
         & right) {
      if (!intersect_ld(a.x, b.x, c.x, d.x) || !
           intersect_ld(a.y, b.y, c.y, d.y)) return
           false:
36
      line m(a, b);
      line n(c, d);
      double zn = det(m.a, m.b, n.a, n.b);
      if (abs(zn) < EPS)  {
        if (abs(m.dist(c)) > EPS || abs(n.dist(a)) >
             EPS) return false;
        if (b < a) swap(a, b);</pre>
        if (d < c) swap(c, d);
        left = max(a, c);
        right = min(b, d);
45
        return true;
      } else {
        left.x = right.x = -det(m.c, m.b, n.c, n.b) /
48
        left.y = right.y = -det(m.a, m.c, n.a, n.c) /
             zn:
49
        return betw(a.x, b.x, left.x) && betw(a.y, b.y,
               betw(c.x, d.x, left.x) && betw(c.y, d.y,
                     left.y);
52 }
```

2.11 Areas

```
double triangle_area (point2d p1, point2d p2,
         point2d p3) {
      return abs(signed_area_parallelogram(p1, p2, p3))
           / 2.0:
   bool clockwise(point2d p1, point2d p2, point2d p3)
      return signed_area_parallelogram(p1, p2, p3) < 0;</pre>
9
10 bool counter clockwise(point2d p1, point2d p2,
         point2d p3) {
      return signed_area_parallelogram(p1, p2, p3) > 0;
12
    double area(const vector<point>& fig) {
      double res = 0;
      for (unsigned i = 0; i < fig.size(); i++) {</pre>
       point p = i ? fig[i - 1] : fig.back();
        point q = fig[i];
        res += (p.x - q.x) * (p.y + q.y);
      return fabs(res) / 2;
21
```

3 Dynamic Programming

3.1 Knuth Optimization

```
1 11 solve() {
     11 N:
      ... // Read input
      vector<vector<ll>> dp(N, vector<ll>(N)), opt(N,
           vector<ll>(N));
      auto C = [&](11 i, 11 j) {
        ... // Implement cost function C.
8
      for (ll i = 0; i < N; i++) {
        opt[i][i] = i;
        ... // Initialize dp[i][i] according to the
      for (11 i = N - 2; i >= 0; i--) {
        for (11 j = i + 1; j < N; j++) {
          11 \text{ mn} = 11\_\text{MAX}, \text{ cost} = C(i, j);
          for (11 k = opt[i][j-1]; k \le min(j-1,
               opt[i + 1][j]); k++) {
            if (mn \ge dp[i][k] + dp[k + 1][j] + cost) {
              opt[i][j] = k;
              mn = dp[i][k] + dp[k + 1][j] + cost;
          dp[i][j] = mn;
      cout << dp[0][N - 1] << '\n';
25
```

3.2 Knapsack

3.3 Divide And Conquer

```
1 11 m, n;
   vector<11> dp_before(n), dp_cur(n);
 3 11 C(11 i, 11 j);
    void compute(l1 1, 11 r, 11 opt1, 11 optr) {
      if (1 > r) return;
      11 \text{ mid} = (1 + r) >> 1;
      pair<11, 11> best = {LLONG_MAX, -1};
      for (ll k = optl; k <= min(mid, optr); k++)</pre>
       best = min(best, \{(k ? dp_before[k - 1] : 0) +
             C(k, mid), k});
10
      dp_cur[mid] = best.first;
      11 opt = best.second;
11
      compute(1, mid - 1, optl, opt);
      compute(mid + 1, r, opt, optr);
14 }
15 11 solve() {
      for (ll i = 0; i < n; i++) dp_before[i] = C(0, i)</pre>
      for (ll i = 1; i < m; i++) {
18
       compute (0, n - 1, 0, n - 1);
19
        dp_before = dp_cur;
21
      return dp_before[n - 1];
```

3.4 Digit Dp

```
1 vector<vector<vector<ll>>>> dp(K + 1, vector
         <vector<vector<11>>>(9 * K + 1, vector<vector</pre>
         11>> (9 * K + 1, vector<11>(9 * K, 0)));
    for (11 n = 1; n \leq 9 * K; n++) dp[0][n][0][0] = 1;
    11 pow10 = 1;
    for (11 k = 1; k <= K; k++) {
      for (11 n = 1; n \le 9 * K; n++) {
        for (11 s = 0; s <= 9 * K; s++) {
          for (11 m = 0; m < n; m++) {
            for (11 y = 0; y \le 9; y++) {
                if (s \ge y) dp[k][n][s][m] += dp[k -
                     1][n][s - y][((m - y * pow10) % n]
                     + n) % n];
      pow10 *= 10;
15 }
16 string N;
17 cin >> N;
18  11  n = N.length(), ans = 0;
19 vector<11> g(9 * K + 1, 0);
20 for (11 s = 1; s <= 9 * K; s++) {
     string substring = "";
      11 pow10 = 1;
      for (11 i = 0; i < n - 1; i++) pow10 *= 10;
      for (ll i = 0; i < n; i++) {
```

```
substring += '0';
        for (11 j = 0; j < N[i] - '0'; j++) {
          11 digit_sum = j;
28
          for (11 k = 0; k < i; k++) digit_sum +=</pre>
               substring[k] - '0';
29
          if (s \ge digit\_sum) g[s] += dp[n - 1 - i][s][
               s - digit_sum][((-pow10 * stoll(
               substring)) % s + s) % s];
          substring[i]++;
        pow10 /= 10;
34
      ans += q[s];
35 }
36 auto is_good = [&](string s) -> bool {
     11 digit_sum = 0;
      for (ll i = 0; i < (ll) s.length(); i++)</pre>
          digit_sum += s[i] - '0';
      return stoll(s) % digit_sum == 0;
40 };
41 if (is_good(N)) ans++;
42 cout << ans << "\n";
```

3.5 Subset Sum

3.6 Longest Increasing Subsequence

```
1  ll get_ceil_idx(vector<ll> &a, vector<ll> &T, ll l,
        11 r, 11 x) {
     while (r - 1 > 1) {
     11 m = 1 + (r - 1) / 2;
      if (a[T[m]] >= x) {
       r = m:
       } else {
         1 = m:
8
     return r;
  11 LIS(11 n, vector<11> &a) {
     11 len = 1;
     vector<ll> T(n, 0), R(n, -1);
    \mathbf{T}[0] = 0;
     for (ll i = 1; i < n; i++) {
     if (a[i] < a[T[0]]) {</pre>
       T[0] = i;
      } else if (a[i] > a[T[len - 1]]) {
        R[i] = T[len - 1];
         T[len++] = i;
       } else {
```

3.7 Longest Common Subsequence

```
1 11 LCS(string x, string y, 11 n, 11 m) {
      vector<vector<ll>> dp(n + 1, vector<ll>(m + 1));
      for (ll i = 0; i <= n; i++) {</pre>
        for (11 j = 0; j <= m; j++) {
          if (i == 0 || j == 0) {
            dp[i][j] = 0;
          } else if (x[i - 1] == y[j - 1]) {
            dp[i][j] = dp[i - 1][j - 1] + 1;
          } else {
            dp[i][j] = max(dp[i - 1][j], dp[i][j - 1]);
      11 \text{ index} = dp[n][m];
      vector<char> lcs(index + 1);
      lcs[index] = ' \setminus 0';
      11 i = n, j = m;
      while (i > 0 \&\& j > 0) {
       if (x[i-1] == y[j-1]) {
         lcs[index - 1] = x[i - 1];
         i--;
          j--;
          index--;
        } else if (dp[i - 1][j] > dp[i][j - 1]) {
         i --:
        } else {
          j--;
28
      return dp[n][m];
```

3.8 Max Sum

```
int max_subarray_sum(vi arr) {
   int x = 0, s = 0;
   for (int k = 0; k < n; k++) {
      s = max(arr[k], s+arr[k]);
      x = max(x,s);
   }
   return x;
}</pre>
```

3.9 Edit Distance

```
for (int j = 1; j <= m; j++) {
8
        dp[0][j] = j;
9
10
      for (int i = 1; i <= n; i++) {</pre>
11
        for (int j = 1; j <= m; j++) {
12
          dp[i][j] = min({dp[i-1][j] + 1, dp[i][j-1]}
               1] + 1, dp[i - 1][j - 1] + (x[i - 1] !=
               y[j - 1])));
14
15
      return dp[n][m];
16 }
```

4 Math

4.1 Chinese Remainder Theorem

```
struct Congruence {
      11 a, m;
 3
    };
   11 chinese_remainder_theorem(vector<Congruence>
         const& congruences) {
      11 M = 1;
      for (auto const& congruence : congruences) M *=
           congruence.m;
      11 \text{ solution} = 0;
      for (auto const& congruence : congruences) {
       11 a_i = congruence.a;
11
       11 M_i = M / congruence.m;
        11 N_i = mod_inv(M_i, congruence.m);
        solution = (solution + a_i * M_i % M * N_i) % M
      return solution:
16 }
```

4.2 Extended Euclidean

```
int gcd(int a, int b, int& x, int& y) {
   if (b == 0) {
       x = 1;
       y = 0;
       return a;
   }
   int x1, y1, d = gcd(b, a % b, x1, y1);
   x = y1;
   y = x1 - y1 * (a / b);
   return d;
}
```

4.3 Modulo Inverse

```
1  11 mod_inv(11 a, 11 m) {
2    if (m == 1) return 0;
3    11 m0 = m, x = 1, y = 0;
4    while (a > 1) {
5         11 q = a / m, t = m;
6         m = a % m;
7         a = t;
8    t = y;
```

4.4 Sum Of Divisors

```
1 11 sum_of_divisors(11 num) {
 2
     11 total = 1;
      for (int i = 2; (11)i * i <= num; i++) {</pre>
        if (num % i == 0) {
 5
           int e = 0;
 6
           do {
            e++;
            num /= i;
           } while (num % i == 0);
           11 \text{ sum} = 0, \text{ pow} = 1;
           do {
           sum += pow;
            pow *= i;
           } while (e-- > 0);
           total *= sum;
16
      if (num > 1) total *= (1 + num);
19
      return total;
20 }
```

4.5 Range Sieve

```
vector<bool> range_sieve(ll 1, ll r) {
      11 n = sart(r):
      vector<bool> is_prime(n + 1, true);
      vector<ll> prime;
      is_prime[0] = is_prime[1] = false;
 6
      prime.push_back(2);
      for (11 i = 4; i <= n; i += 2) is_prime[i] =</pre>
           false:
      for (11 i = 3; i <= n; i += 2) {
9
       if (is_prime[i]) {
10
          prime.push_back(i);
11
          for (ll j = i * i; j <= n; j += i) is_prime[j</pre>
               | = false;
14
      vector<bool> result(r - 1 + 1, true);
15
      for (ll i : prime)
       for (11 j = max(i * i, (1 + i - 1) / i * i); j
            <= r; j += i)
          result[j - 1] = false;
      if (1 == 1) result[0] = false;
19
      return result;
20 }
```

4.6 Pollard Rho Brent

```
1 ll mult(11 a, 11 b, 11 mod) {
2 return (_int128_t) a * b % mod;
3 }
4 ll f(11 x, 11 c, 11 mod) {
```

```
return (mult(x, x, mod) + c) % mod;
    11 pollard_rho_brent(11 n, 11 x0 = 2, 11 c = 1) {
      11 x = x0, g = 1, q = 1, xs, y, m = 128, 1 = 1;
      while (q == 1) {
       y = x;
        for (11 i = 1; i < 1; i++) x = f(x, c, n);
        11 k = 0:
        while (k < 1 \&\& g == 1) {
          xs = x;
          for (ll i = 0; i < m \&\& i < l - k; i++) {
           x = f(x, c, n);
            q = mult(q, abs(y - x), n);
          g = \underline{gcd}(q, n);
          k += m;
        1 *= 2:
      if (g == n) {
        do {
         xs = f(xs, c, n);
          q = \underline{gcd(abs(xs - y), n)};
        } while (q == 1);
2.9
      return q;
```

4.7 Factorial Modulo

4.8 Matrix

```
Matrix exponentation:
    f[n] = af[n-1] + bf[n-2] + cf[n-3]
    Use:
    |f[n]| | |a|b|c||f[n-1]|
    |f[n-1]|=|1 0 0||f[n-2]|
    |f[n-2]| |0 1 0||f[n-3]|
    To get:
    |f[n]| | |a|b|c|^{(n-2)}|f[2]|
    |f[n-1]|=|1 0 0| |f[1]|
    |f[n-2]| |0 1 0|
                          |f[0]|
12
13 struct Matrix { int mat[MAX_N][MAX_N]; };
   Matrix matrix_mul(Matrix a, Matrix b) {
     Matrix ans; int i, j, k;
      for (i = 0; i < MAX_N; i++)</pre>
      for (j = 0; j < MAX_N; j++)</pre>
18
      for (ans.mat[i][\dot{\eta}] = k = 0; k < MAX_N; k++)
        ans.mat[i][j] += a.mat[i][k] * b.mat[k][j];
```

```
return ans;
21
22 Matrix matrix_pow(Matrix base, int p) {
23
      Matrix ans; int i, j;
24
      for (i = 0; i < MAX_N; i++)</pre>
25
      for (j = 0; j < MAX_N; j++)
26
        ans.mat[i][j] = (i == j);
      while (p) {
       if (p & 1) ans = matrix_mul(ans, base);
       base = matrix mul(base, base);
32
      return ans;
33 }
```

4.9 Find All Solutions

```
bool find_any_solution(ll a, ll b, ll c, ll &x0, ll
          &v0, 11 &q) {
      g = gcd_extended(abs(a), abs(b), x0, y0);
      if (c % g) return false;
      x0 \star = c / q;
      y0 \star = c / q;
      if (a < 0) \times 0 = -x0;
      if (b < 0) y0 = -y0;
      return true;
9 }
10 void shift_solution(ll & x, ll & y, ll a, ll b, ll
11
      x += cnt * b;
     y -= cnt * a;
12
13 }
14 ll find_all_solutions(ll a, ll b, ll c, ll minx, ll
          maxx, 11 miny, 11 maxy) {
      11 x, y, g;
      if (!find_any_solution(a, b, c, x, y, g)) return
           0;
      a /= q;
      b /= q;
      11 \text{ sign}_a = a > 0 ? +1 : -1;
      11 \text{ sign}_b = b > 0 ? +1 : -1;
      shift_solution(x, y, a, b, (minx - x) / b);
      if (x < minx) shift_solution(x, y, a, b, sign_b);</pre>
      if (x > maxx) return 0;
24
      11 1x1 = x;
25
      shift_solution(x, y, a, b, (maxx - x) / b);
      if (x > maxx) shift_solution(x, y, a, b, -sign_b)
      shift_solution(x, y, a, b, -(miny - y) / a);
      if (y < miny) shift_solution(x, y, a, b, -sign_a)</pre>
      if (y > maxy) return 0;
      11 \ 1x2 = x;
      shift_solution(x, y, a, b, -(maxy - y) / a);
      if (y > maxy) shift_solution(x, y, a, b, sign_a);
      11 \text{ rx2} = x;
      if (1x2 > rx2) swap(1x2, rx2);
      11 1x = max(1x1, 1x2), rx = min(rx1, rx2);
      if (1x > rx) return 0;
38
      return (rx - lx) / abs(b) + 1;
39 }
```

```
4.10 Miller Rabin
```

```
using u64 = uint64_t;
    using u128 = uint128 t;
   u64 binpower(u64 base, u64 e, u64 mod) {
     u64 result = 1;
   base %= mod;
     while (e) {
      if (e & 1) result = (u128) result * base % mod;
       base = (u128) base * base % mod;
       e >>= 1;
      return result;
13 bool check_composite(u64 n, u64 a, u64 d, 11 s) {
     u64 x = binpower(a, d, n);
15
     if (x == 1 \mid | x == n - 1) return false;
      for (11 r = 1; r < s; r++) {
      x = (u128) x * x % n;
       if (x == n - 1) return false;
      return true;
   bool miller_rabin(u64 n) {
     if (n < 2) return false;</pre>
     11 r = 0;
      u64 d = n - 1;
      while ((d & 1) == 0) {
       d >>= 1;
       r++;
      for (11 a : {2, 3, 5, 7, 11, 13, 17, 19, 23, 29,
         31, 37}) {
        if (n == a) return true;
       if (check_composite(n, a, d, r)) return false;
33
      return true;
```

4.11 Fibonacci

and put a 1

```
1 /*
 2 Properties:
   - Cassini's identity: f[n-1]f[n+1] - f[n]^2 = (-1)^
 4 - d'Ocagne's identity: f[m]f[n+1] - f[m+1]f[n] =
 5 - Addition rule: f[n+k] = f[k]f[n+1] + f[k-1]f[n]
 6 - k = n \text{ case: } f[2n] = f[n](f[n+1] + f[n-1])
 7 - f[n] / f[nk]
 8 - f[n] / f[m] => n / m
9 - GCD rule: gcd(f[m], f[n]) = f[gcd(m, n)]
10 - [[1 1], [1 0]]^n = [[f[n+1] f[n]], [f[n], f[n]]
        -1]]]
   - f(2k+1) = f(k+1)^2 + f(k)^2
   -f[2k] = f[k](f[k+1] + f[k-1]) = f[k](2f[k+1] - f[k])
   - Periodic sequence modulo p
   - sum[i=1..n]f[i] = f[n+2] - 1
15 - sum[i=0..n-1]f[2i+1] = f[2n]
16 - sum[i=1..n]f[2i] = f[2n+1] - 1
17 - sum[i=1..n]f[i]^2 = f[n]f[n+1]
18 Fibonacci encoding:
19 1. Iterate through the Fibonacci numbers from the
        largest to the
20 smallest until you find one less than or equal to n
21 2. Suppose this number was F_i. Subtract F_i from n
```

```
22 in the i-2 position of the code word (indexing from
          0 from the
23 leftmost to the rightmost bit).
24 3. Repeat until there is no remainder.
25 4. Add a final 1 to the codeword to indicate its
26 Closed-form: f[n] = (((1 + rt(5))/2)^n - ((1 - rt))^n
         (5)) / 2) ^n)/rt(5)
28
29 struct matrix {
        11 mat[2][2];
        matrix friend operator *(const matrix &a, const
31
              matrix &b) {
          matrix c;
          for (int i = 0; i < 2; i++) {</pre>
          for (int j = 0; j < 2; j++) {
             c.mat[i][j] = 0;
              for (int k = 0; k < 2; k++) c.mat[i][j]</pre>
                  += a.mat[i][k] * b.mat[k][j];
38
39
          return c;
      matrix matpow(matrix base, 11 n) {
       matrix ans{ {
         {1, 0},
          {0, 1}
        } };
        while (n) {
          if (n & 1) ans = ans * base;
          base = base * base;
          n >>= 1;
52
        return ans;
53
      11 fib(int n) {
        matrix base{ {
         {1, 1},
         {1, 0}
        } };
        return matpow(base, n).mat[0][1];
      pair<int, int> fib (int n) {
        if (n == 0) return {0, 1};
        auto p = fib(n >> 1);
        int c = p.first * (2 * p.second - p.first);
        int d = p.first * p.first + p.second * p.second
        if (n & 1) return {d, c + d};
        else return {c, d};
69 }
```

4.12 Fast Fourier Transform

```
1  using cd = complex<double>;
2  const double PI = acos(-1);
3  void fft(vector<cd>& a, bool invert) {
4   int n = a.size();
5   if (n == 1) return;
6  vector<cd> a0 (n / 2), a1 (n / 2);
7  for (int i = 0; 2 * i < n; i++) {
8   a0[i] = a[2 * i];
9   a1[i] = a[2 * i + 1];
</pre>
```

```
fft(a0, invert);
12
       fft(al, invert);
13
      double ang = 2 * PI / n * (invert ? -1 : 1);
      cd w(1), wn(cos(ang), sin(ang));
      for (int i = 0; 2 * i < n; i++) {
16
        a[i] = a0[i] + w * a1[i];
17
        a[i + n / 2] = a0[i] - w * a1[i];
18
        if (invert) {
19
          a[i] /= 2;
20
          a[i + n / 2] /= 2;
         w \star = wn;
23
      }
24 }
25 vector<int> multiply(vector<int> const& a, vector<
         int> const& b) {
26
         vector<cd> fa(a.begin(), a.end()), fb(b.begin()
             , b.end());
27
         int n = 1;
28
         while (n < a.size() + b.size()) n <<= 1;</pre>
29
         fa.resize(n);
         fb.resize(n);
         fft(fa, false);
         fft(fb, false);
33
         for (int i = 0; i < n; i++) fa[i] *= fb[i];</pre>
34
         fft(fa, true);
         vector<int> result(n);
36
         for (int i = 0; i < n; i++) result[i] = round(</pre>
             fa[i].real());
         return result;
38 }
```

4.13 Segmented Sieve

```
vector<ll> segmented_sieve(ll n) {
      const 11 S = 10000;
      11 nsqrt = sqrt(n);
      vector<char> is_prime(nsqrt + 1, true);
      vector<ll> prime;
      is_prime[0] = is_prime[1] = false;
      prime.push_back(2);
      for (11 i = 4; i <= nsqrt; i += 2) {</pre>
       is_prime[i] = false;
10
11
      for (11 i = 3; i <= nsqrt; i += 2) {
       if (is_prime[i]) {
13
          prime.push_back(i);
14
          for (11 j = i * i; j <= nsqrt; j += i) {</pre>
15
            is_prime[j] = false;
16
17
        }
18
19
      vector<ll> result:
20
      vector<char> block(S);
      for (11 k = 0; k * S <= n; k++) {
        fill(block.begin(), block.end(), true);
23
        for (ll p : prime) {
24
          for (11 j = max((k * S + p - 1) / p, p) * p -
                k * S; j < S; j += p) {
            block[j] = false;
26
27
28
        if (k == 0) {
29
          block[0] = block[1] = false;
31
         for (11 i = 0; i < S && k * S + i <= n; i++) {
          if (block[i]) {
```

4.14 Linear Sieve

4.15 Tonelli Shanks

```
1 ll legendre(ll a, ll p) {
     return bin_pow_mod(a, (p - 1) / 2, p);
3
   11 tonelli_shanks(ll n, ll p) {
     if (legendre(n, p) == p - 1) {
       return -1;
     if (p % 4 == 3) {
       return bin_pow_mod(n, (p + 1) / 4, p);
      11 Q = p - 1, S = 0;
      while (0 % 2 == 0) {
       0 /= 2;
       S++;
      for (; z < p; z++) {
       if (legendre(z, p) == p - 1) {
19
      11 M = S, c = bin_pow_mod(z, Q, p), t =
          bin_pow_mod(n, Q, p), R = bin_pow_mod(n, Q)
          + 1) / 2, p);
      while (t % p != 1) {
        if (t % p == 0) {
          return 0;
        11 i = 1, t2 = t * t % p;
        for (; i < M; i++) {
         if (t2 % p == 1) {
           break;
          t2 = t2 * t2 % p;
        11 b = bin_pow_mod(c, bin_pow_mod(2, M - i - 1,
             p), p);
        M = i;
        c = b * b % p;
        t = t * c % p;
```

5 Miscellaneous

5.1 Techniques

```
Dynamic Programming
3 - Bitmask
4 - Range
5 - Digit
6 - Knapsack
   Graph Theory
8 - Tree diameter
   - Reversing edges
10 - Tree re-rooting
   - DP on trees
   - DFS tree
   - Euler tour
   - Binary Jumping
   - Centroid
   - DAG
   - Condense
18 Data Structures
19 - Multiple information
20 - Binary searching on the tree
21 - 2D range query
22 - SORT decomposition
   - Small-to-large
   Sorting and searching
25 - Sliding window
   - Two pointers
   - Binary search on the answer
```

5.2 Gauss

```
1 const double EPS = 1e-9;
    const 11 INF = 2;
    11 gauss(vector <vector <double>> a, vector <double>
          &ans) {
      11 n = (11) a.size(), m = (11) a[0].size() - 1;
      vector<ll> where (m, -1);
      for (11 col = 0, row = 0; col < m && row < n; ++</pre>
           col) {
        11 \text{ sel} = \text{row};
        for (ll i = row; i < n; ++i) {</pre>
           if (abs(a[i][col]) > abs(a[sel][col])) {
             sel = i;
        if (abs (a[sel][col]) < EPS) {</pre>
           continue:
        for (ll i = col; i <= m; ++i) {</pre>
           swap(a[sel][i], a[row][i]);
        where[col] = row;
        for (11 i = 0; i < n; ++i) {
21
           if (i != row) {
             double c = a[i][col] / a[row][col];
```

```
for (ll j = col; j <= m; ++j) {</pre>
24
               a[i][j] -= a[row][j] * c;
27
28
         ++row;
29
30
31
32
33
       ans.assign(m, 0);
       for (ll i = 0; i < m; ++i) {</pre>
         if (where[i] != -1) {
           ans[i] = a[where[i]][m] / a[where[i]][i];
34
36
       for (11 i = 0; i < n; ++i) {
37
         double sum = 0;
38
         for (11 j = 0; j < m; ++j) {
39
          sum += ans[j] * a[i][j];
40
41
         if (abs (sum - a[i][m]) > EPS) {
42
           return 0;
43
44
45
       for (11 i = 0; i < m; ++i) {
46
        if (where[i] == -1) {
47
           return INF;
48
49
50
       return 1;
```

5.3 Ternary Search

6 Data Structures

6.1 Segment Tree 2d

```
seg_tree = new SegTree<T, InType>(a[i]);
          return;
15
        int k = (i + j) / 2;
16
        lc = new SegTree2dNode<T, InType>(a, i, k);
        rc = new SegTree2dNode<T, InType>(a, k, j);
        seg_tree = new SegTree<T, InType>(vector<T>(
             tree_size));
        operation_2d(lc->seg_tree, rc->seg_tree);
       ~SegTree2dNode() {
        delete lc;
        delete rc;
      void set_2d(int kx, int ky, T x) {
        if (kx < i || j <= kx) return;</pre>
        if (j - i == 1) {
          seg_tree->set(ky, x);
          return:
        1c->set_2d(kx, ky, x);
        rc \rightarrow set_2d(kx, ky, x);
        operation_2d(lc->seg_tree, rc->seg_tree);
35
      T range_query_2d(int lx, int rx, int ly, int ry)
36
        if (lx <= i && j <= rx) return seq_tree->
             range_query(ly, ry);
37
        if (j <= lx || rx <= i) return -INF;</pre>
        return max(lc->range_query_2d(lx, rx, ly, ry),
             rc->range_query_2d(lx, rx, ly, ry));
39
40
      void operation_2d(SegTree<T, InType>* x, SegTree<</pre>
           T, InType>* v) {
        for (int k = 0; k < tree_size; k++) {</pre>
          seg_tree->set(k, max(x->range_query(k, k + 1)
               , y->range_query(k, k + 1)));
43
44
     }
45
    };
46
    template<typename T, typename InType = T>
    class SegTree2d {
    public:
49
      SegTree2dNode<T, InType> root;
50
      SegTree2d() {}
      SegTree2d(const vector<vector<InType>>& mat) :
            root(mat, 0, mat.size()) {}
      void set_2d(int kx, int ky, T x) { root.set_2d(kx
      T range_query_2d(int lx, int rx, int ly, int ry)
           { return root.range_query_2d(lx, rx, ly, ry)
           ; }
54 };
```

6.2 Range Add Point Query

```
return;
        int k = (i + j) / 2;
        lc = new SegTreeNode<T, InType>(i, k);
        rc = new SegTreeNode<T, InType>(k, j);
17
        val = 0;
18
19
      SegTreeNode(const vector<InType>& a, int i, int j
           ) : i(i), j(j) {
        if (i - i == 1) {
          lc = rc = nullptr;
          val = (T) a[i];
          return;
        int k = (i + j) / 2;
        lc = new SegTreeNode<T, InType>(a, i, k);
        rc = new SegTreeNode<T, InType>(a, k, j);
        val = 0;
      void range_add(int 1, int r, T x) {
        if (r <= i || j <= 1) return;</pre>
        if (1 <= i && j <= r) {</pre>
          val += x;
          return;
        lc->range_add(1, r, x);
        rc->range_add(l, r, x);
38
      T point_query(int k) {
        if (k < i || j <= k) return IDN;</pre>
        if (j - i == 1) return val;
        return val + lc->point_query(k) + rc->
             point_query(k);
    template<typename T, typename InType = T>
    class SegTree {
47
    public:
      SegTreeNode<T, InType> root;
      SegTree(int n) : root(0, n) {}
      SegTree(const vector<InType>& a) : root(a, 0, a.
           size()) {}
      void range_add(int 1, int r, T x) { root.
           range_add(1, r, x); }
      T point_query(int k) { return root.point_query(k)
           ; }
53 };
```

6.3 Disjoint Set Union

```
16     b = find_set(b);
17     if (a != b) {
18         if (size[a] < size[b]) swap(a, b);
19         parent[b] = a;
20         size[a] += size[b];
21     }
22     }
23     };</pre>
```

6.4 Sparse Table 2d

```
const int N = 100;
    int matrix[N][N];
    int table[N][N][(int)(log2(N) + 1)][(int)(log2(N) +
     void build_sparse_table(int n, int m) {
       for (int i = 0; i < n; i++)</pre>
         for (int j = 0; j < m; j++)
           table[i][j][0][0] = matrix[i][j];
      for (int k = 1; k \le (int)(log2(n)); k++)
         for (int i = 0; i + (1 << k) - 1 < n; i++)
           for (int j = 0; j + (1 << k) - 1 < m; j++)
             table[i][j][k][0] = min(table[i][j][k -
                 1][0], table[i + (1 << (k - 1))][j][k
                  - 1][0]);
      for (int k = 1; k \le (int)(log2(m)); k++)
13
         for (int i = 0; i < n; i++)
14
           for (int j = 0; j + (1 << k) - 1 < m; <math>j++)
15
             table[i][j][0][k] = min(table[i][j][0][k -
                  1], table[i][j + (1 << (k - 1))][0][k
                  - 1]);
16
      for (int k = 1; k \le (int)(log2(n)); k++)
17
         for (int 1 = 1; 1 <= (int) (log2(m)); 1++)</pre>
18
           for (int i = 0; i + (1 << k) - 1 < n; i++)
19
             for (int j = 0; j + (1 << 1) - 1 < m; <math>j++)
20
               table[i][j][k][l] = min(
                 min(table[i][j][k-1][l-1], table[i]
                      + (1 << (k - 1))][j][k - 1][1 -
                      1]),
                 min(table[i][j + (1 << (1 - 1))][k -
                      1][1-1], table[i + (1 << (k - 1)
                      ) ] [j + (1 << (1 - 1))] [k - 1] [1 -
                      11)
23
               );
25 int rmq(int x1, int y1, int x2, int y2) {
26
      int k = log2(x2 - x1 + 1), 1 = log2(y2 - y1 + 1);
27
28
        \max(\text{table}[x1][y1][k][1], \text{table}[x2 - (1 << k) +
              1][y1][k][l]),
29
         \max(\text{table}[x1][y2 - (1 << 1) + 1][k][1], \text{ table}[
              x2 - (1 << k) + 1][y2 - (1 << 1) + 1][k][1
              1)
      );
31 }
```

6.5 Mo

```
struct Query {
      int 1, r, idx;
      bool operator<(Query other) const {</pre>
 8
        return make_pair(l / block_size, r) < make_pair</pre>
             (other.l / block_size, other.r);
 9
    };
11
    vector<int> mo_s_algorithm(vector<Query> queries) {
12
      vector<int> answers(queries.size());
      sort(queries.begin(), queries.end());
      // TODO: initialize data structure
      int cur_1 = 0, cur_r = -1;
      // invariant: data structure will always reflect
           the range [cur_1, cur_r]
      for (Query q : queries) {
        while (cur_1 > q.1) {
19
          cur_1--;
          add(cur_l);
        while (cur_r < q.r) {</pre>
          cur_r++;
          add(cur_r);
        while (cur_1 < q.1) {
          remove(cur_l);
          cur 1++;
3.0
        while (cur_r > q.r) {
          remove(cur_r);
          cur_r--;
        answers[q.idx] = get_answer();
      return answers;
```

6.6 Sparse Table

```
1  11 log2_floor(l1 i) {
      return i ? __builtin_clzll(1) - __builtin_clzll(i
           ): -1;
    vector<vector<ll>> build_sum(ll N, ll K, vector<ll>
      vector<vector<ll>> st(K + 1, vector<ll>(N + 1));
      for (ll i = 0; i < N; i++) st[0][i] = array[i];</pre>
      for (11 i = 1; i <= K; i++)
        for (11 j = 0; j + (1 << i) <= N; <math>j++)
          st[i][j] = st[i - 1][j] + st[i - 1][j + (1 <<
                 (i - 1))];
      return st;
11
    11 sum_query(11 L, 11 R, 11 K, vector<vector<11>>> &
         st) {
      11 \text{ sum} = 0;
      for (11 i = K; i >= 0; i--) {
        if ((1 << i) <= R - L + 1) {
          sum += st[i][L];
          L += 1 << i;
18
19
      return sum;
    vector<vector<ll>> build_min(ll N, ll K, vector<ll>
          &arrav) {
      vector<vector<ll>> st(K + 1, vector<ll>(N + 1));
```

for (ll i = 0; i < N; i++) st[0][i] = array[i];</pre>

```
6.7 Binary Trie
    struct Node { struct Node* parent, child[2]; };
    struct BinaryTrie {
      Node* root;
      BinaryTrie() {
        root = new Node();
        root->parent = NULL;
        root->child[0] = NULL:
        root->child[1] = NULL:
      void insert_node(int x) {
       Node* cur = root;
        for (int place = 29; place >= 0; place--) {
          int bit = x >> place & 1;
          if (cur->child[bit] != NULL) cur = cur->child
               [bit]:
          else {
            cur->child[bit] = new Node();
            cur->child[bit]->parent = cur;
            cur = cur->child[bit];
            cur->child[0] = NULL;
            cur->child[1] = NULL;
23
      void remove_node(int x) {
        Node* cur = root;
        for (int place = 29; place >= 0; place--) {
          int bit = x >> place & 1;
          if (cur->child[bit] == NULL) return;
          cur = cur->child[bit];
        while (cur->parent != NULL && cur->child[0] ==
            NULL && cur->child[1] == NULL) {
          Node* temp = cur;
          cur = cur->parent;
          if (temp == cur->child[0]) cur->child[0] =
               NULL;
          else cur->child[1] = NULL;
          delete temp;
      int get_min_xor(int x) {
        Node* cur = root;
        int minXor = 0;
        for (int place = 29; place >= 0; place--) {
          int bit = x >> place & 1;
44
          if (cur->child[bit] != NULL) cur = cur->child
               [bit];
          else {
            minXor ^= 1 << place:
            cur = cur->child[1 ^ bit];
        return minXor;
```

```
51 };
52 };
```

6.8 Segment Tree

```
template<typename T, typename InType = T>
    class SegTreeNode {
    public:
      const T IDN = 0, DEF = 0;
      int i, j;
      T val;
      SegTreeNode<T, InType>* lc, * rc;
      SegTreeNode(int i, int j) : i(i), j(j) {
        if (j - i == 1) {
          lc = rc = nullptr;
11
          val = DEF;
          return;
13
14
         int k = (i + j) / 2;
        lc = new SegTreeNode<T, InType>(i, k);
         rc = new SegTreeNode<T, InType>(k, j);
17
         val = op(lc->val, rc->val);
18
19
      SegTreeNode(const vector<InType>& a, int i, int j
           ) : i(i), j(j) {
20
         if (j - i == 1) {
21
         lc = rc = nullptr;
          val = (T) a[i];
23
          return;
24
25
26
         int k = (i + j) / 2;
        lc = new SegTreeNode<T, InType>(a, i, k);
         rc = new SegTreeNode<T, InType>(a, k, j);
28
        val = op(lc -> val, rc -> val);
29
30
31
      void set(int k, T x) {
        if (k < i || j <= k) return;</pre>
32
        if (i - i == 1) {
33
34
          vai = x;
          return;
35
36
        1c->set(k, x);
         rc \rightarrow set(k, x);
38
        val = op(lc -> val, rc -> val);
39
40
     T range_query(int 1, int r) {
41
        if (1 <= i && j <= r) return val;</pre>
42
         if (j <= 1 || r <= i) return IDN;</pre>
4.3
         return op(lc->range_query(l, r), rc->
             range_query(1, r));
44
45
     T \circ p(T \times, T y) \{ \}
46 };
47
    template<typename T, typename InType = T>
    class SegTree {
49
50
      SegTreeNode<T, InType> root;
51
      SegTree(int n) : root(0, n) {}
52
      SegTree(const vector<InType>& a) : root(a, 0, a.
            size()) {}
       void set(int k, T x) { root.set(k, x); }
54
      T range_query(int 1, int r) { return root.
            range_query(1, r); }
55 };
```

6.9 Sqrt Decomposition

```
1 int n:
   vector<int> a (n);
   int len = (int) sqrt (n + .0) + 1; // size of the
        block and the number of blocks
    vector<int> b (len);
    for (int i = 0; i<n; ++i) b[i / len] += a[i];</pre>
   for (;;) {
     int 1, r;
 8
      // read input data for the next query
      int sum = 0;
      for (int i = 1; i <= r; )</pre>
       if (i % len == 0 && i + len - 1 <= r) {</pre>
        // if the whole block starting at i belongs
              to [1, r]
          sum += b[i / len];
          i += len;
        } else {
         sum += a[i];
          ++i;
18
19
      // or
      int sum = 0;
      if (c_1 == c_r)
          for (int i=1; i<=r; ++i)
25
             sum += a[i];
          for (int i=1, end=(c_1+1)*len-1; i<=end; ++i)
             sum += a[i];
29
          for (int i=c_1+1; i<=c_r-1; ++i)
             sum += b[i];
          for (int i=c_r*len; i<=r; ++i)
             sum += a[i];
34
35 }
```

6.10 Minimum Queue

```
1 11 get_minimum(stack<pair<11, 11>> &s1, stack<pair<</pre>
         11, 11>> &s2) {
      if (s1.empty() || s2.empty()) {
        return s1.empty() ? s2.top().second : s1.top().
        return min(s1.top().second, s2.top().second);
 6
    void add_element(ll new_element, stack<pair<ll, ll</pre>
      11 minimum = s1.empty() ? new_element : min(
           new_element, s1.top().second);
      s1.push({new_element, minimum});
11
    11 remove_element(stack<pair<11, 11>> &s1, stack
         pair<11, 11>> &s2) {
      if (s2.empty()) {
        while (!s1.empty()) {
          11 element = s1.top().first;
16
          11 minimum = s2.empty() ? element : min(
               element, s2.top().second);
          s2.push({element, minimum});
```

```
19     }
20     }
21     11 removed_element = s2.top().first;
22     s2.pop();
23     return removed_element;
24  }
```

6.11 Range Add Range Query

```
template<typename T, typename InType = T>
    class SegTreeNode {
    public:
      const T IDN = 0, DEF = 0;
      int i, j;
      T val, to_add = 0;
      SegTreeNode<T, InType>* lc, * rc;
      SegTreeNode(int i, int j) : i(i), j(j) {
        if (j - i == 1) {
          lc = rc = nullptr;
          val = DEF:
          return;
        int k = (i + j) / 2;
        lc = new SegTreeNode<T, InType>(i, k);
        rc = new SegTreeNode<T, InType>(k, j);
        val = operation(lc->val, rc->val);
18
      SegTreeNode(const vector<InType>& a, int i, int j
          ) : i(i), j(j) {
        if (j - i == 1) {
          lc = rc = nullptr;
          val = (T) a[i];
          return;
        int k = (i + j) / 2;
        lc = new SegTreeNode<T, InType>(a, i, k);
        rc = new SegTreeNode<T, InType>(a, k, j);
2.8
        val = operation(lc->val, rc->val);
29
      void propagate() {
        if (to_add == 0) return;
        val += to add:
        if (i - i > 1) {
          lc->to_add += to_add;
          rc->to_add += to_add;
        to\_add = 0;
38
39
      void range_add(int 1, int r, T delta) {
        propagate();
        if (r <= i || j <= 1) return;</pre>
        if (1 <= i && j <= r) {
          to add += delta;
          propagate();
        } else {
          lc->range_add(1, r, delta);
          rc->range_add(1, r, delta);
          val = operation(lc->val, rc->val);
50
51
      T range_query(int 1, int r) {
        propagate();
        if (1 <= i && j <= r) return val;</pre>
54
        if (j <= 1 || r <= i) return IDN;</pre>
55
        return operation(lc->range_query(l, r), rc->
             range_query(1, r));
```

```
T operation(T x, T y) {}
58
    };
    template<typename T, typename InType = T>
60
    class SegTree {
61
    public:
62
      SegTreeNode<T, InType> root;
63
      SegTree(int n) : root(0, n) {}
      SeqTree(const vector<InType>& a) : root(a, 0, a.
           size()) {}
      void range_add(int 1, int r, T delta) { root.
           range_add(l, r, delta); }
      T range_query(int 1, int r) { return root.
           range_query(1, r); }
67 };
```

7 Graph Theory

7.1 Bridge

```
int n:
    vector<vector<int>> adj;
    vector<bool> visited;
    vector<int> tin, low;
    int timer;
    void dfs(int v, int p = -1) {
     visited[v] = true;
      tin[v] = low[v] = timer++;
      for (int to : adj[v]) {
       if (to == p) continue;
11
        if (visited[to]) {
12
         low[v] = min(low[v], tin[to]);
13
       } else {
14
          dfs(to, v);
15
          low[v] = min(low[v], low[to]);
          if (low[to] > tin[v]) IS_BRIDGE(v, to);
17
18
     }
19 }
20 void find_bridges() {
21
      timer = 0;
      visited.assign(n, false);
23
      tin.assign(n, -1);
      low.assign(n, -1);
25
      for (int i = 0; i < n; ++i) {
       if (!visited[i]) dfs(i);
27
28 1
```

7.2 Dijkstra

```
const int INF = 1000000000;
    vector<vector<pair<int, int>>> adj;
    void dijkstra(int s, vector<int> & d, vector<int> &
      int n = adj.size();
      d.assign(n, INF);
      p.assign(n, -1);
      d[s] = 0;
      using pii = pair<int, int>;
      priority_queue<pii, vector<pii>, greater<pii>> q;
      q.push({0, s});
11
      while (!q.empty()) {
12
        int v = q.top().second, d_v = q.top().first;
13
        q.pop();
```

```
if (d_v != d[v]) continue;
for (auto edge : adj[v]) {
    int to = edge.first, len = edge.second;
    if (d[v] + len < d[to]) {
        d[to] = d[v] + len;
        p[to] = v;
        q.push({d[to], to});
}

21    }
}
23  }
}</pre>
```

7.3 Zero One Bfs

```
1  vector<int> d(n, INF);
2  d(s] = 0;
3  deque<int> q;
4  q.push_front(s);
5  while (!q.empty()) {
6   int v = q.front();
7  q.pop_front();
8  for (auto edge : adj[v]) {
9   int u = edge.first, w = edge.second;
10  if (d[v] + w < d[u]) {
11  d[u] = d[v] + w;
12  if (w == 1) q.push_back(u);
13  else q.push_front(u);
14  }
15  }
16 }</pre>
```

7.4 Hungarian

```
vector<int> u (n+1), v (m+1), p (m+1), way (m+1);
 2 for (int i=1; i<=n; ++i) {</pre>
      p[0] = i;
      int j0 = 0;
      vector<int> minv (m+1, INF);
 6
      vector<bool> used (m+1, false);
      do √
        used[j0] = true;
        int i0 = p[j0], delta = INF, j1;
10
        for (int j=1; j<=m; ++j)</pre>
          if (!used[j]) {
12
            int cur = A[i0][j]-u[i0]-v[j];
1.3
            if (cur < minv[j]) minv[j] = cur, way[j] =</pre>
             if (minv[j] < delta) delta = minv[j], j1 =</pre>
                   j;
16
         for (int j=0; j \le m; ++j)
          if (used[j]) u[p[j]] += delta, v[j] -= delta
          else minv[j] -= delta;
         j0 = j1;
      } while (p[j0] != 0);
        int j1 = way[j0];
        p[j0] = p[j1];
         j0 = j1;
      } while (†0);
    vector<int> ans (n+1);
    for (int j=1; j<=m; ++j)</pre>
      ans[p[j]] = j;
```

```
30 int cost = -v[0];
7.5 Ford Fulkerson
bool bfs(ll n, vector<vector<ll>>> &r_graph, ll s,
         11 t, vector<11> &parent) {
      vector<bool> visited(n, false);
      queue<11> q;
      q.push(s);
      visited[s] = true;
      parent[s] = -1;
      while (!q.empty()) {
       11 u = q.front();
        q.pop();
        for (11 \ v = 0; \ v < n; \ v++) {
          if (!visited[v] && r_graph[u][v] > 0) {
            if (v == t) {
              parent[v] = u;
              return true;
            q.push(v);
            parent[v] = u;
            visited[v] = true;
      return false;
24
    11 ford_fulkerson(ll n, vector<vector<ll>>> graph,
         11 s, 11 t) {
      11 u, v;
      vector<vector<11>> r_graph;
      for (u = 0; u < n; u++)
2.8
        for (v = 0; v < n; v++)
          r_{graph[u][v]} = graph[u][v];
      vector<11> parent;
      11 \text{ max\_flow} = 0;
      while (bfs(n, r_graph, s, t, parent)) {
        11 path_flow = INF;
34
        for (v = t; v != s; v = parent[v]) {
         u = parent[v];
          path_flow = min(path_flow, r_graph[u][v]);
        for (v = t; v != s; v = parent[v]) {
         u = parent[v];
          r_graph[u][v] -= path_flow;
          r_graph[v][u] += path_flow;
```

7.6 Prim

46

```
const int INF = 1000000000;
struct Edge {
   int w = INF, to = -1;
   bool operator<(Edge const& other) const {
      return make_pair(w, to) < make_pair(other.w, other.to);
};
int n;
vector<vector<Edge>> adj;
```

max_flow += path_flow;

return max_flow;

```
10 void prim() {
11
      int total weight = 0;
      vector<Edge> min_e(n);
      min_e[0].w = 0;
      set < Edge > q;
15
      g.insert({0, 0});
16
      vector<bool> selected(n, false);
17
      for (int i = 0; i < n; ++i) {</pre>
18
         if (q.empty()) {
19
           cout << "No MST!" << endl;</pre>
20
           exit(0);
21
22
         int v = q.begin()->to;
23
         selected[v] = true;
24
         total weight += g.begin()->w;
25
         q.erase(q.begin());
2.6
         if (min_e[v].to != -1) cout << v << " " <<</pre>
             min_e[v].to << endl;</pre>
         for (Edge e : adj[v]) {
28
           if (!selected[e.to] && e.w < min_e[e.to].w) {</pre>
29
             q.erase({min_e[e.to].w, e.to});
             min_e[e.to] = \{e.w, v\};
31
             g.insert({e.w, e.to});
33
34
      cout << total_weight << endl;</pre>
```

7.7 Centroid Decomposition

```
1 vector<vector<int>> adj;
    vector<bool> is removed;
    vector<int> subtree size;
    int get_subtree_size(int node, int parent = -1) {
            subtree_size[node] = 1;
            for (int child : adj[node]) {
                    if (child == parent || is_removed[
                         child]) continue;
 8
                    subtree_size[node] +=
                        get_subtree_size(child, node);
            return subtree size[node];
11
12
    int get_centroid(int node, int tree_size, int
         parent = -1) {
13
            for (int child : adj[node]) {
14
                    if (child == parent || is_removed[
                         child]) continue;
15
                    if (subtree_size[child] * 2 >
                         tree_size) return get_centroid
                         (child, tree_size, node);
16
17
            return node;
18 }
    void build centroid decomp(int node = 0) {
20
            int centroid = get centroid(node,
                 get_subtree_size(node));
21
            // do something
            is_removed[centroid] = true;
23
            for (int child : adj[centroid]) {
24
                    if (is_removed[child]) continue;
25
                    build_centroid_decomp(child);
26
27 }
```

7.8 Kahn void kahn(vector<vector<ll>> &adi) { 11 n = adj.size();vector<ll> in_degree(n, 0); for (11 u = 0; u < n; u++) for (ll v: adj[u]) in_degree[v]++; queue<11> q; for (ll i = 0; i < n; i++)</pre> if (in_degree[i] == 0) q.push(i); 11 cnt = 0;vector<ll> top_order; while (!q.empty()) { 11 u = q.front();q.pop(); top_order.push_back(u); for (ll v : adj[u]) if (--in_degree[v] == 0) q.push(v); cnt++; 19 **if** (cnt != n) { cout << -1 << '\n'; return; // print top_order

7.9 Dinics

```
struct FlowEdge {
      int v, u;
      11 \text{ cap, flow} = 0;
      FlowEdge(int v, int u, ll cap) : v(v), u(u), cap(
           cap) {}
 5
    struct Dinic {
      const 11 flow_inf = 1e18;
      vector<FlowEdge> edges;
      vector<vector<int>> adj;
      int n, m = 0, s, t;
      vector<int> level, ptr;
       queue<int> q;
      Dinic(int n, int s, int t) : n(n), s(s), t(t) {
        adj.resize(n);
        level.resize(n);
16
        ptr.resize(n);
18
      void add_edge(int v, int u, ll cap) {
        edges.emplace_back(v, u, cap);
        edges.emplace_back(u, v, 0);
        adj[v].push_back(m);
         adj[u].push_back(m + 1);
        m += 2;
24
      bool bfs() {
        while (!q.empty()) {
          int v = q.front();
          q.pop();
          for (int id : adj[v]) {
            if (edges[id].cap - edges[id].flow < 1)</pre>
             if (level[edges[id].u] != -1) continue;
            level[edges[id].u] = level[v] + 1;
            g.push(edges[id].u);
```

```
36
        return level[t] != -1;
38
      11 dfs(int v, 11 pushed) {
39
        if (pushed == 0) return 0;
40
        if (v == t) return pushed;
        for (int& cid = ptr[v]; cid < (int)adj[v].size</pre>
41
             (); cid++) {
          int id = adj[v][cid], u = edges[id].u;
43
          if (level[v] + 1 != level[u] || edges[id].cap
                - edges[id].flow < 1) continue;</pre>
44
          11 tr = dfs(u, min(pushed, edges[id].cap -
               edges[id].flow));
          if (tr == 0) continue;
          edges[id].flow += tr;
          edges[id ^ 1].flow -= tr;
          return tr;
        return 0:
      11 flow() {
        11 f = 0;
        while (true) {
          fill(level.begin(), level.end(), -1);
56
          level[s] = 0;
          q.push(s);
58
          if (!bfs()) break;
59
          fill(ptr.begin(), ptr.end(), 0);
          while (ll pushed = dfs(s, flow_inf)) f +=
               pushed;
        return f;
63
64 };
```

7.10 Floyd Warshall

7.11 Kosaraju

```
11
    vector<vector<int>> transpose(int n, vector<vector<</pre>
         int>>& adj) {
      vector<vector<int>> adj_t(n);
13
      for (int u = 0; u < n; u++) {
14
        for (int v : adj[u]) {
15
          adj_t[v].push_back(u);
16
17
18
      return adj t;
19 }
20
    void get_scc(int u, vector<vector<int>>& adj_t,
         vector<bool>& vis, vector<int>& scc) {
      vis[u] = true;
23
      scc.push_back(u);
24
      for (int v : adj_t[u]) {
25
       if (!vis[v]) {
26
          get_scc(v, adj_t, vis, scc);
27
28
      }
29
    void kosaraju(int n, vector<vector<int>>& adj,
         vector<vector<int>>& sccs) {
      vector<bool> vis(n, false);
      stack<int> stk;
      for (int u = 0; u < n; u++) {</pre>
        if (!vis[u]) {
          topo_sort(u, adj, vis, stk);
38
39
      vector<vector<int>> adj_t = transpose(n, adj);
      for (int u = 0; u < n; u++) {
41
       vis[u] = false;
42
43
      while (!stk.empty()) {
44
       int u = stk.top();
45
       stk.pop();
46
        if (!vis[u]) {
47
        vector<int> scc;
          get_scc(u, adj_t, vis, scc);
49
          sccs.push_back(scc);
50
51
52 }
```

7.12 Maximum Bipartite Matching

7.13 Kruskals

```
1 struct Edge {
      int u, v, weight;
      bool operator<(Edge const& other) {</pre>
        return weight < other.weight;</pre>
 6
    };
    int n;
    vector<Edge> edges;
    int cost = 0;
10 vector<Edge> result;
11 DSU dsu = DSU(n);
12 sort(edges.begin(), edges.end());
13 for (Edge e : edges) {
   if (dsu.find_set(e.u) != dsu.find_set(e.v)) {
        cost += e.weight;
        result.push back(e);
        dsu.union_sets(e.u, e.v);
18
19 }
```

7.14 Is Cyclic

```
bool is_cyclic_util(int u, vector<vector<int>> &adj
        , vector<bool> &vis, vector<bool> &rec) {
      vis[u] = true;
      rec[u] = true;
      for(auto v : adj[u]) {
        if (!vis[v] && is_cyclic_util(v, adj, vis, rec)
            ) return true;
 6
        else if (rec[v]) return true;
 8
      rec[u] = false;
 9
      return false;
    bool is_cyclic(int n, vector<vector<int>> &adj) {
      vector<bool> vis(n, false), rec(n, false);
      for (int i = 0; i < n; i++)
14
       if (!vis[i] && is_cyclic_util(i, adj, vis, rec)
            ) return true;
      return false;
```

7.15 Find Cycle

```
1 bool dfs(ll v) {
2   color[v] = 1;
3   for (ll u : adj[v]) {
4     if (color[u] == 0) {
5       parent[u] = v;
6     if (dfs(u)) {
```

```
return true:
        } else if (color[u] == 1) {
          cycle_end = v;
          cycle_start = u;
          return true;
14
      color[v] = 2;
      return false;
18
    void find_cycle() {
      color.assign(n, 0);
      parent.assign(n, -1);
      cycle start = -1;
      for (11 v = 0; v < n; v++) {
       if (color[v] == 0 && dfs(v)) {
          break:
      if (cycle_start == -1) {
        cout << "Acyclic" << endl;</pre>
      } else {
        vector<11> cycle;
        cycle.push_back(cycle_start);
        for (ll v = cycle_end; v != cycle_start; v =
             parent[v]) {
          cycle.push_back(v);
        cycle.push_back(cycle_start);
        reverse(cycle.begin(), cycle.end());
        cout << "Cycle found: ";</pre>
38
        for (11 v : cycle) {
39
          cout << v << ' ';
        cout << '\n';
43 }
```

7.16 Topological Sort

```
void dfs(ll v) {
    visited[v] = true;
    for (ll u : adj[v]) {
        if (!visited[u]) {
            dfs(u);
        }
    }
    ans.push_back(v);
    }

void topological_sort() {
    visited.assign(n, false);
    ans.clear();
    for (ll i = 0; i < n; ++i) {
        if (!visited[i]) {
            dfs(i);
        }
    }
    reverse(ans.begin(), ans.end());
}</pre>
```

7.17 Min Cost Flow

```
1 struct Edge {
2 int from, to, capacity, cost;
```

```
3 };
    vector<vector<int>> adj, cost, capacity;
    const int INF = 1e9;
    void shortest_paths(int n, int v0, vector<int>& d,
         vector<int>& p) {
      d.assign(n, INF);
      d[v0] = 0;
      vector<bool> inq(n, false);
      queue<int> q;
11
      q.push(v0);
      p.assign(n, -1);
13
      while (!q.empty()) {
        int u = q.front();
14
15
16
        ing[u] = false;
17
        for (int v : adj[u]) {
18
         if (capacity[u][v] > 0 && d[v] > d[u] + cost[
               u][v]) {
            d[v] = d[u] + cost[u][v];
20
            p[v] = u;
21
            if (!ing[v]) {
22
              inq[v] = true;
23
              q.push(v);
24
25
26
27
28
    int min_cost_flow(int N, vector<Edge> edges, int K,
          int s, int t) {
      adj.assign(N, vector<int>());
      cost.assign(N, vector<int>(N, 0));
      capacity.assign(N, vector<int>(N, 0));
      for (Edge e : edges) {
34
        adj[e.from].push_back(e.to);
35
        adj[e.to].push_back(e.from);
        cost[e.from][e.to] = e.cost;
        cost[e.to][e.from] = -e.cost;
38
        capacity[e.from][e.to] = e.capacity;
39
40
      int flow = 0;
      int cost = 0;
      vector<int> d, p;
43
      while (flow < K) {</pre>
44
        shortest paths(N, s, d, p);
45
        if (d[t] == INF) break;
46
        int f = K - flow, cur = t;
47
        while (cur != s) {
48
          f = min(f, capacity[p[cur]][cur]);
49
          cur = p[cur];
51
        flow += f;
        cost += f * d[t];
        cur = t;
54
        while (cur != s) {
          capacity[p[cur]][cur] -= f;
56
          capacity[cur][p[cur]] += f;
57
          cur = p[cur];
58
59
60
      if (flow < K) return -1;</pre>
61
      else return cost;
```

```
7.18 Kuhn
```

```
1 int n, k;
```

```
vector<vector<int>> g;
     vector<int> mt;
    vector<bool> used;
    bool try_kuhn(int v) {
      if (used[v]) return false;
       used[v] = true;
       for (int to : q[v]) {
        if (mt[to] == -1 || try_kuhn(mt[to])) {
           mt[to] = v;
           return true;
 13
 14
       return false;
 15 }
 16 int main() {
      mt.assign(k, -1);
18
        vector<bool> used1(n, false);
         for (int v = 0; v < n; ++v) {
           for (int to : q[v]) {
            if (mt[to] == -1) {
               mt[to] = v;
               used1[v] = true;
               break;
 26
         for (int v = 0; v < n; ++v) {
           if (used1[v]) continue;
           used.assign(n, false);
           try_kuhn(v);
 32
         for (int i = 0; i < k; ++i)
 34
           if (mt[i] != −1)
 35
             printf("%d %d\n", mt[i] + 1, i + 1);
```

7.19 Articulation Point

```
void APUtil(vector<vector<ll>> &adj, ll u, vector
         bool> &visited.
 2 vector<11> &disc, vector<11> &low, 11 &time, 11
         parent, vector<bool> &isAP) {
      11 children = 0;
      visited[u] = true;
5
      disc[u] = low[u] = ++time;
      for (auto v : adj[u]) {
        if (!visited[v]) {
8
          children++;
9
          APUtil(adj, v, visited, disc, low, time, u,
              isAP):
          low[u] = min(low[u], low[v]);
          if (parent != -1 && low[v] >= disc[u]) {
           isAP[u] = true;
        } else if (v != parent) {
          low[u] = min(low[u], disc[v]);
18
      if (parent == -1 && children > 1) {
19
       isAP[u] = true;
20
2.1
   void AP(vector<vector<ll>>> &adj, ll n) {
      vector<ll> disc(n), low(n);
      vector<bool> visited(n), isAP(n);
      11 time = 0, par = -1;
```

for (11 u = 0; u < n; u++) {

7.20 Hierholzer

```
void print circuit(vector<vector<11>> &adj) {
      map<11, 11> edge_count;
      for (ll i = 0; i < adj.size(); i++) {</pre>
        edge_count[i] = adj[i].size();
6
      if (!adj.size()) {
       return;
8
     stack<ll> curr_path;
      vector<ll> circuit;
      curr_path.push(0);
      11 \text{ curr } v = 0;
      while (!curr_path.empty()) {
        if (edge_count[curr_v]) {
          curr_path.push(curr_v);
          11 next_v = adj[curr_v].back();
          edge_count[curr_v]--;
          adj[curr_v].pop_back();
19
          curr_v = next_v;
        } else {
          circuit.push_back(curr_v);
          curr_v = curr_path.top();
23
          curr_path.pop();
24
      for (ll i = circuit.size() - 1; i >= 0; i--) {
2.7
        cout << circuit[i] << ' ';
28
```

7.21 Lowest Common Ancestor

```
1 struct LCA {
     vector<ll> height, euler, first, segtree;
     vector<bool> visited;
     11 n:
     LCA(vector<vector<11>> &adj, 11 root = 0) {
       n = adj.size();
       height.resize(n);
       first.resize(n);
       euler.reserve(n * 2);
      visited.assign(n, false);
       dfs(adj, root);
      11 m = euler.size();
       segtree.resize(m * 4);
       build(1, 0, m - 1);
     void dfs(vector<vector<11>> &adj, 11 node, 11 h =
           0) {
       visited[node] = true;
       height[node] = h;
```

```
19
         first[node] = euler.size();
20
         euler.push back(node);
21
         for (auto to : adj[node]) {
           if (!visited[to]) {
23
             dfs(adj, to, h + 1);
24
              euler.push_back(node);
25
26
27
28
29
       void build(ll node, ll b, ll e) {
         if (b == e) segtree[node] = euler[b];
30
31
           11 \text{ mid} = (b + e) / 2;
32
           build(node << 1, b, mid);</pre>
33
           build(node << 1 | 1, mid + 1, e);</pre>
34
           11 1 = segtree[node << 1], r = segtree[node</pre>
                << 1 | 1];
           segtree[node] = (height[1] < height[r]) ? 1 :</pre>
36
37
38
       11 query(11 node, 11 b, 11 e, 11 L, 11 R) {
39
         if (b > R | | e < L) return -1;</pre>
40
         if (b >= L && e <= R) return segtree[node];</pre>
41
         11 \text{ mid} = (b + e) >> 1;
42
         11 left = query(node << 1, b, mid, L, R);</pre>
43
         11 right = query(node << 1 | 1, mid + 1, e, L,</pre>
44
         if (left == -1) return right;
45
         if (right == -1) return left;
46
         return height[left] < height[right] ? left :</pre>
47
48
       ll lca(ll u, ll v) {
49
         11 left = first[u], right = first[v];
50
         if (left > right) swap(left, right);
         return query(1, 0, euler.size() - 1, left,
              right);
52 }
53 };
```

7.22 Bellman Ford

```
struct Edge {
      int a, b, cost;
 3 };
 4 int n, m, v;
 5 vector<Edge> edges;
    const int INF = 1000000000;
    void solve() {
      vector<int> d(n, INF);
      d[v] = 0;
      vector<int> p(n, -1);
      int x;
      for (int i = 0; i < n; ++i) {
13
       \mathbf{x} = -1;
14
        for (Edge e : edges)
15
         if (d[e.a] < INF)
16
            if (d[e.b] > d[e.a] + e.cost) {
17
              d[e.b] = max(-INF, d[e.a] + e.cost);
18
              p[e.b] = e.a;
19
              x = e.b;
20
21
      if (x == -1) cout << "No negative cycle from " <<
      else {
```

7.23 Edmonds Karp

```
vector<vector<int>> capacity;
    vector<vector<int>> adj;
    int bfs(int s, int t, vector<int>& parent) {
     fill(parent.begin(), parent.end(), -1);
      parent[s] = -2;
      queue<pair<int, int>> q;
 8
      q.push({s, INF});
      while (!q.empty()) {
        int cur = q.front().first, flow = q.front().
             second;
        q.pop();
        for (int next : adj[cur]) {
          if (parent[next] == -1 && capacity[cur][next
14
            parent[next] = cur;
            int new_flow = min(flow, capacity[cur][next
            if (next == t) return new_flow;
            q.push({next, new_flow});
18
19
        }
20
      }
      return 0;
    int maxflow(int s, int t) {
      int flow = 0;
      vector<int> parent(n);
26
27
      int new_flow;
      while (new_flow = bfs(s, t, parent)) {
       flow += new_flow;
        int cur = t;
        while (cur != s) {
          int prev = parent[cur];
          capacity[prev][cur] -= new_flow;
          capacity[cur][prev] += new_flow;
          cur = prev;
35
36
      return flow;
38 1
```

7.24 Is Bipartite

7.25 Fast Second Mst

```
struct edge {
        int s, e, w, id;
 3
        bool operator<(const struct edge& other) {</pre>
             return w < other.w; }</pre>
    typedef struct edge Edge;
 6 const int N = 2e5 + 5;
    long long res = 0, ans = 1e18;
    int n, m, a, b, w, id, 1 = 21;
    vector<Edge> edges;
    vector<int> h(N, 0), parent(N, -1), size(N, 0),
         present (N, 0);
    vector<vector<pair<int, int>>> adj(N), dp(N, vector
         <pair<int, int>>(1));
    vector<vector<int>> up(N, vector<int>(1, -1));
    pair<int, int> combine(pair<int, int> a, pair<int,</pre>
         int> b) {
      vector<int> v = {a.first, a.second, b.first, b.
           second);
      int topTwo = -3, topOne = -2;
      for (int c : v) {
       if (c > topOne) {
18
          topTwo = topOne;
19
          topOne = c;
        } else if (c > topTwo && c < topOne) topTwo = c
22
      return {topOne, topTwo};
23
    void dfs(int u, int par, int d) {
      h[u] = 1 + h[par];
      up[u][0] = par;
      dp[u][0] = \{d, -1\};
2.8
      for (auto v : adj[u]) {
2.9
        if (v.first != par) dfs(v.first, u, v.second);
32
    pair<int, int> lca(int u, int v) {
      pair<int, int> ans = \{-2, -3\};
      if (h[u] < h[v]) swap(u, v);</pre>
      for (int i = 1 - 1; i >= 0; i--) {
        if (h[u] - h[v] >= (1 << i)) {
          ans = combine(ans, dp[u][i]);
38
          u = up[u][i];
39
```

```
40
 41
       if (u == v) return ans;
       for (int i = 1 - 1; i >= 0; i--) {
       if (up[u][i] != -1 && up[v][i] != -1 && up[u][i
             ] != up[v][i]) {
           ans = combine(ans, combine(dp[u][i], dp[v][i
               1));
          u = up[u][i];
          v = up[v][i];
 49
       ans = combine(ans, combine(dp[u][0], dp[v][0]));
       return ans:
 51 }
 52
 53 int main(void) {
 54
     cin >> n >> m;
       for (int i = 1; i <= n; i++) {</pre>
       parent[i] = i;
        size[i] = 1;
 58
 59
       for (int i = 1; i <= m; i++) {</pre>
 60
        cin >> a >> b >> w; // 1-indexed
 61
        edges.push_back(\{a, b, w, i - 1\});
 62
 63
       sort(edges.begin(), edges.end());
 64
       for (int i = 0; i <= m - 1; i++) {
       a = edges[i].s;
      b = edges[i].e;
 67
      w = edges[i].w;
       id = edges[i].id;
      if (unite_set(a, b)) {
        adj[a].emplace_back(b, w);
        adj[b].emplace back(a, w);
        present[id] = 1;
         res += w;
 74
       dfs(1, 0, 0);
       for (int i = 1; i <= 1 - 1; i++) {
        for (int j = 1; j <= n; ++j) {</pre>
 79
         if (up[j][i - 1] != -1) {
 80
            int v = up[j][i - 1];
 81
             up[j][i] = up[v][i - 1];
             dp[j][i] = combine(dp[j][i-1], dp[v][i-
                 1]);
 84
 85
       for (int i = 0; i <= m - 1; i++) {
 87
        id = edges[i].id;
 88
        w = edges[i].w;
 89
        if (!present[id]) {
 90
         auto rem = lca(edges[i].s, edges[i].e);
 91
         if (rem.first != w) {
            if (ans > res + w - rem.first) ans = res +
                 w - rem.first;
         } else if (rem.second != -1) {
            if (ans > res + w - rem.second) ans = res +
                  w - rem.second;
 96
        }
 97
 98
       cout << ans << "\n";
99
       return 0;
100 }
```

8 References

8.1 Ref

```
1 // vector
   push back()
3 pop_back()
   size()
 5 clear()
 6 erase()
 7 empty()
 8  Iterator lower_bound(Iterator first, Iterator last,
         const val)
   Iterator upper_bound(Iterator first, Iterator last,
     const val)
11 push()
12 pop()
13 top()
14 empty()
15 size()
16 // queue
   push()
18 pop()
19 front()
20 empty()
21 back()
22 size()
23 // priority_queue
24 push()
25 pop()
26 size()
27 empty()
28 top()
29 // set
30 insert()
31 begin()
32 end()
33 size()
34 find()
35 count()
36 empty()
37 // multiset
38 begin()
39 end()
40 size()
41 max size()
42 emptv()
43 insert(x) // O(log n)
44 clear()
45 erase(x)
46 // map
47 begin()
48 end()
49 size()
50 max size()
51 emptv()
52 pair insert(keyvalue, mapvalue)
53 erase(iterator position)
54 erase(const q)
55 clear()
56 // ordered set
57 find_by_order(k)
58 order_of_key(k)
```

59 **#include** <ext/pb_ds/assoc_container.hpp>

```
#include <ext/pb_ds/tree_policy.hpp>
     using namespace gnu pbds;
     #define ordered_set
        tree<int, null_type, less<int>, rb_tree_tag, \
             tree_order_statistics_node_update>
 66 // tuple
 67 get<i>(tuple)
 68 make_tuple(a1, a2, ...)
 69 tuple size < decltype (tuple) >:: value
 70 tuple1.swap(tuple2)
 71 tie(a1, a2, ...) = tuple
 72 tuple_cat(tuple1, tuple2)
 73 // iterator
 74 for (auto it = s.begin(); it != s.end(); it++) cout
         << *it << "\n";
 75 begin()
 76 end()
    advance(ptr, k)
 78 next(ptr, k)
 79 prev(ptr, k)
 80 // permutations
 81 do {} while (next_permutation(nums.begin(), nums.
         end()));
 82 // bitset
 83 int num = 27; // Binary representation: 11011
 84 bitset<10> s(string("0010011010")); // from right
         to left
 85 bitset<sizeof(int) * 8> bits(num);
 86 int setBits = bits.count();
 87 bits.set(index, val);
 88 bits.reset();
 89 bits.flip();
 90 bits.all();
 91 bits.any();
 92 bits.none();
 93 bits.test();
 94 to string();
 95 to ulong();
 96 to_ullong();
 97 [], &, |, !, >>=, <<=, &=, |=, ^=, ~;
    // sort
 100 sort(v.rbegin(), v.rend());
101 // custom sort
102 bool comp(string a, string b) {
103 if (a.size() != b.size()) return a.size() < b.
          size();
    return a < b; }
105 sort(v.begin(), v.end(), comp);
106 // hamming distance
107 int hamming(int a, int b) { return
         __builtin_popcount(a ^ b); }
108 // custom comparator for pq
109 class Compare {
110 public:
111 bool operator() (T a, T b) {
112 if(cond) return true; // do not swap
113 return false; } };
114 priority_queue<PII, vector<PII>, Compare> ds;
115 // gcc compiler
116 __builtin_popcount(x)
    __builtin_parity(x)
     __builtin_clz(x) // leading
119 builtin_ctz(x) // trailing
```

f(n) = O(g(n))	iff \exists positive c, n_0 such that	n n						
	$0 \le f(n) \le cg(n) \ \forall n \ge n_0.$	$\sum_{i=1}^{n} i = \frac{n(n+1)}{2}, \sum_{i=1}^{n} i^2 = \frac{n(n+1)(2n+1)}{6}, \sum_{i=1}^{n} i^3 = \frac{n^2(n+1)^2}{4}.$						
$f(n) = \Omega(g(n))$	iff \exists positive c, n_0 such that $f(n) \ge cg(n) \ge 0 \ \forall n \ge n_0$.	In general:						
$f(n) = \Theta(g(n))$	iff $f(n) = O(g(n))$ and $f(n) = \Omega(g(n))$.	$\sum_{i=1}^{n} i^{m} = \frac{1}{m+1} \left[(n+1)^{m+1} - 1 - \sum_{i=1}^{n} \left((i+1)^{m+1} - i^{m+1} - (m+1)i^{m} \right) \right]$						
f(n) = o(g(n))	iff $\lim_{n\to\infty} f(n)/g(n) = 0$.	$\sum_{i=1}^{n-1} i^m = \frac{1}{m+1} \sum_{k=0}^m \binom{m+1}{k} B_k n^{m+1-k}.$						
$\lim_{n \to \infty} a_n = a$	iff $\forall \epsilon > 0$, $\exists n_0$ such that $ a_n - a < \epsilon, \forall n \ge n_0$.	Geometric series:						
$\sup S$	least $b \in \mathbb{R}$ such that $b \geq s$, $\forall s \in S$.	$\sum_{i=0}^{n} c^{i} = \frac{c^{n+1} - 1}{c - 1}, c \neq 1, \sum_{i=0}^{\infty} c^{i} = \frac{1}{1 - c}, \sum_{i=1}^{\infty} c^{i} = \frac{c}{1 - c}, c < 1,$						
$\inf S$	greatest $b \in \mathbb{R}$ such that $b \le s$, $\forall s \in S$.	$\sum_{i=0}^{n} ic^{i} = \frac{nc^{n+2} - (n+1)c^{n+1} + c}{(c-1)^{2}}, c \neq 1, \sum_{i=0}^{\infty} ic^{i} = \frac{c}{(1-c)^{2}}, c < 1.$						
$ \liminf_{n \to \infty} a_n $	$\lim_{n\to\infty}\inf\{a_i\mid i\geq n, i\in\mathbb{N}\}.$	Harmonic series: $n + n = n + n = n = n = n = n = n = n = $						
$ \limsup_{n \to \infty} a_n $	$\lim_{n \to \infty} \sup \{ a_i \mid i \ge n, i \in \mathbb{N} \}.$	$H_n = \sum_{i=1}^n \frac{1}{i}, \qquad \sum_{i=1}^n iH_i = \frac{n(n+1)}{2}H_n - \frac{n(n-1)}{4}.$						
$\binom{n}{k}$	Combinations: Size k subsets of a size n set.	$\sum_{i=1}^{n} H_i = (n+1)H_n - n, \sum_{i=1}^{n} {i \choose m} H_i = {n+1 \choose m+1} \left(H_{n+1} - \frac{1}{m+1} \right).$						
$\begin{bmatrix} n \\ k \end{bmatrix}$	Stirling numbers (1st kind): Arrangements of an n element set into k cycles.	$1. \binom{n}{k} = \frac{n!}{(n-k)!k!}, \qquad 2. \sum_{k=0}^{n} \binom{n}{k} = 2^n, \qquad 3. \binom{n}{k} = \binom{n}{n-k},$						
$\left\{ egin{array}{c} n \\ k \end{array} \right\}$	Stirling numbers (2nd kind): Partitions of an n element set into k non-empty sets.	$4. \binom{n}{k} = \frac{n}{k} \binom{n-1}{k-1}, \qquad \qquad 5. \binom{n}{k} = \binom{n-1}{k} + \binom{n-1}{k-1}, \\ 6. \binom{n}{m} \binom{m}{k} = \binom{n}{k} \binom{n-k}{m-k}, \qquad \qquad 7. \sum_{k=0}^{n} \binom{r+k}{k} = \binom{r+n+1}{n},$						
$\left\langle {n\atop k}\right\rangle$	1st order Eulerian numbers: Permutations $\pi_1\pi_2\pi_n$ on $\{1, 2,, n\}$ with k ascents.	8. $\sum_{k=0}^{n} {k \choose m} = {n+1 \choose m+1},$ 9. $\sum_{k=0}^{n} {r \choose k} {s \choose n-k} = {r+s \choose n},$						
$\langle\!\langle {n \atop k} \rangle\!\rangle$	2nd order Eulerian numbers.	10. $\binom{n}{k} = (-1)^k \binom{k-n-1}{k}$, 11. $\binom{n}{1} = \binom{n}{n} = 1$,						
C_n	Catalan Numbers: Binary trees with $n+1$ vertices.	12. $\binom{n}{2} = 2^{n-1} - 1$, 13. $\binom{n}{k} = k \binom{n-1}{k} + \binom{n-1}{k-1}$,						
$14. \begin{bmatrix} n \\ 1 \end{bmatrix} = (n-1)$	15. $\begin{bmatrix} n \\ 2 \end{bmatrix} = (n - 1)^n$	16. $\begin{bmatrix} n \\ n \end{bmatrix} = 1,$ 17. $\begin{bmatrix} n \\ k \end{bmatrix} \ge \begin{Bmatrix} n \\ k \end{Bmatrix},$						
$18. \begin{bmatrix} n \\ k \end{bmatrix} = (n-1) \begin{bmatrix} n-1 \\ k \end{bmatrix} + \begin{bmatrix} n-1 \\ k-1 \end{bmatrix}, 19. \begin{bmatrix} n \\ n-1 \end{bmatrix} = \begin{bmatrix} n \\ n-1 \end{bmatrix} = \begin{pmatrix} n \\ 2 \end{pmatrix}, 20. \sum_{k=0}^{n} \begin{bmatrix} n \\ k \end{bmatrix} = n!, 21. \ C_n = \frac{1}{n+1} \binom{2n}{n},$								
$22. \binom{n}{0} = \binom{n}{n}$	$\begin{pmatrix} n \\ -1 \end{pmatrix} = 1,$ 23. $\begin{pmatrix} n \\ k \end{pmatrix} = \langle$	$\binom{n}{n-1-k}$, $24. \left\langle \binom{n}{k} \right\rangle = (k+1) \left\langle \binom{n-1}{k} \right\rangle + (n-k) \left\langle \binom{n-1}{k-1} \right\rangle$,						
25. $\binom{0}{k} = \begin{Bmatrix} 1 \\ 0 \end{Bmatrix}$	if $k = 0$, otherwise 26. $\left\langle {}^{n}\right\rangle$	$\binom{n}{2} = 2^n - n - 1,$ 27. $\binom{n}{2} = 3^n - (n+1)2^n + \binom{n+1}{2},$						
$28. \ \ x^n = \sum_{k=0}^n \binom{n}{k} \binom{x+k}{n}, \qquad 29. \ \ \binom{n}{m} = \sum_{k=0}^m \binom{n+1}{k} (m+1-k)^n (-1)^k, \qquad 30. \ \ m! \begin{Bmatrix} n \\ m \end{Bmatrix} = \sum_{k=0}^n \binom{n}{k} \binom{k}{n-m},$								
$31. \ \left\langle \begin{array}{c} n \\ m \end{array} \right\rangle = \sum_{k=0}^{n} \left\{ \begin{array}{c} n \\ k \end{array} \right\} \binom{n-k}{m} (-1)^{n-k-m} k!, \qquad \qquad 32. \ \left\langle \begin{array}{c} n \\ 0 \end{array} \right\rangle = 1, \qquad \qquad 33. \ \left\langle \begin{array}{c} n \\ n \end{array} \right\rangle = 0 \text{for } n \neq 0,$								
$34. \; \left\langle \!\! \left\langle \!\! \begin{array}{c} n \\ k \end{array} \!\! \right\rangle = (k + 1)^n$	$+1$) $\left\langle \left\langle \begin{array}{c} n-1 \\ k \end{array} \right\rangle + (2n-1-k) \left\langle \left\langle \begin{array}{c} n-1 \\ k \end{array} \right\rangle \right\rangle$							
$36. \left\{ \begin{array}{c} x \\ x-n \end{array} \right\} = \frac{1}{n}$	$\sum_{k=0}^{n} \left\langle \!\! \left\langle n \right\rangle \!\! \right\rangle \left(x + n - 1 - k \right), $ $2n$	37. ${n+1 \choose m+1} = \sum_{k} {n \choose k} {k \choose m} = \sum_{k=0}^{n-1} {k \choose m} (m+1)^{n-k},$						

The Chinese remainder theorem: There exists a number C such that:

 $C \equiv r_1 \mod m_1$

: : :

 $C \equiv r_n \mod m_n$

if m_i and m_j are relatively prime for $i \neq j$. Euler's function: $\phi(x)$ is the number of positive integers less than x relatively prime to x. If $\prod_{i=1}^{n} p_i^{e_i}$ is the prime factorization of x then

$$\phi(x) = \prod_{i=1}^{n} p_i^{e_i - 1} (p_i - 1).$$

Euler's theorem: If a and b are relatively prime then

$$1 \equiv a^{\phi(b)} \bmod b.$$

Fermat's theorem:

$$1 \equiv a^{p-1} \bmod p.$$

The Euclidean algorithm: if a > b are integers then

$$gcd(a, b) = gcd(a \mod b, b).$$

If $\prod_{i=1}^{n} p_i^{e_i}$ is the prime factorization of x

$$S(x) = \sum_{d|x} d = \prod_{i=1}^{n} \frac{p_i^{e_i+1} - 1}{p_i - 1}.$$

Perfect Numbers: x is an even perfect number iff $x = 2^{n-1}(2^n - 1)$ and $2^n - 1$ is prime. Wilson's theorem: n is a prime iff $(n-1)! \equiv -1 \mod n$.

$$\mu(i) = \begin{cases} (n-1)! = -1 \bmod n. \\ \text{M\"obius inversion:} \\ \mu(i) = \begin{cases} 1 & \text{if } i = 1. \\ 0 & \text{if } i \text{ is not square-free.} \\ (-1)^r & \text{if } i \text{ is the product of} \\ r & \text{distinct primes.} \end{cases}$$
 If

 If

$$G(a) = \sum_{d|a} F(d),$$

$$F(a) = \sum_{d \mid a} \mu(d) G\left(\frac{a}{d}\right).$$

Prime numbers:

$$p_n = n \ln n + n \ln \ln n - n + n \frac{\ln \ln n}{\ln n}$$

$$+O\left(\frac{n}{\ln n}\right),$$

$$\pi(n) = \frac{n}{\ln n} + \frac{n}{(\ln n)^2} + \frac{2!n}{(\ln n)^3} + O\left(\frac{n}{(\ln n)^4}\right).$$

efir		

An edge connecting a ver-Looptex to itself.

Directed Each edge has a direction. SimpleGraph with no loops or multi-edges.

WalkA sequence $v_0e_1v_1\ldots e_\ell v_\ell$. TrailA walk with distinct edges. Pathtrail with distinct

vertices.

ConnectedA graph where there exists a path between any two

vertices.

ComponentΑ maximal connected subgraph.

TreeA connected acyclic graph. Free tree A tree with no root. DAGDirected acyclic graph. EulerianGraph with a trail visiting each edge exactly once.

Hamiltonian Graph with a cycle visiting each vertex exactly once.

CutA set of edges whose removal increases the number of components.

Cut-setA minimal cut. Cut edge A size 1 cut.

k-Connected A graph connected with the removal of any k-1vertices.

k-Tough $\forall S \subseteq V, S \neq \emptyset$ we have $k \cdot c(G - S) \le |S|$.

A graph where all vertices k-Regular have degree k.

k-Factor Α k-regular spanning subgraph.

Matching A set of edges, no two of which are adjacent.

CliqueA set of vertices, all of which are adjacent.

Ind. set A set of vertices, none of which are adjacent.

Vertex cover A set of vertices which cover all edges.

Planar graph A graph which can be embeded in the plane.

Plane graph An embedding of a planar

$$\sum_{v \in V} \deg(v) = 2m.$$

If G is planar then n - m + f = 2, so

$$f \le 2n - 4, \quad m \le 3n - 6.$$

Any planar graph has a vertex with degree ≤ 5 .

Notation:

E(G)Edge set Vertex set V(G)

c(G)Number of components

G[S]Induced subgraph

deg(v)Degree of v

Maximum degree $\Delta(G)$ $\delta(G)$ Minimum degree Chromatic number

 $\chi(G)$ $\chi_E(G)$ Edge chromatic number G^c Complement graph

 K_n Complete graph

 K_{n_1,n_2} Complete bipartite graph

Ramsev number

Geometry

Projective coordinates: (x, y, z), not all x, y and z zero.

$$(x, y, z) = (cx, cy, cz) \quad \forall c \neq 0.$$

Cartesian Projective (x, y)(x, y, 1)y = mx + b(m, -1, b)x = c(1,0,-c)

Distance formula, L_p and L_{∞}

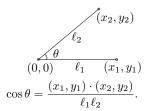
$$\sqrt{(x_1 - x_0)^2 + (y_1 - y_0)^2},$$
$$[|x_1 - x_0|^p + |y_1 - y_0|^p]^{1/p},$$

$$\lim_{n \to \infty} \left[|x_1 - x_0|^p + |y_1 - y_0|^p \right]^{1/p}.$$

Area of triangle $(x_0, y_0), (x_1, y_1)$ and (x_2, y_2) :

$$\frac{1}{2} \operatorname{abs} \begin{vmatrix} x_1 - x_0 & y_1 - y_0 \\ x_2 - x_0 & y_2 - y_0 \end{vmatrix}.$$

Angle formed by three points:



Line through two points (x_0, y_0) and (x_1, y_1) :

$$\begin{vmatrix} x & y & 1 \\ x_0 & y_0 & 1 \\ x_1 & y_1 & 1 \end{vmatrix} = 0.$$

Area of circle, volume of sphere:

$$A = \pi r^2, \qquad V = \frac{4}{3}\pi r^3.$$

If I have seen farther than others, it is because I have stood on the shoulders of giants.

- Issac Newton