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 1.1 Manacher
```

### 1 Strings

2

```
vector<int> manacher odd(string s) {
     int n = s.size();
     s = "$" + s + "^n;
     vector<int> p(n + 2);
     int 1 = 1, r = 1;
     for (int i = 1; i \le n; i++) {
       p[i] = max(0, min(r - i, p[1 + (r - i)]));
       while(s[i - p[i]] == s[i + p[i]]) p[i]++;
       if(i + p[i] > r) l = i - p[i], r = i + p[i];
     return vector<int>(begin(p) + 1, end(p) - 1);
   vector<int> manacher(string s) {
     string t;
     for(auto c: s) t += string("#") + c;
     auto res = manacher_odd(t + "#");
     return vector<int>(begin(res) + 1, end(res) - 1);
. 8
```

```
1.2 Hashing
```

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9

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11 11

```
11 compute_hash(string const& s) {
11
11
          const 11 p = 31, m = 1e9 + 9;
12
          11 hash_value = 0, p_pow = 1;
12
          for (char c : s) {
12
           hash\_value = (hash\_value + (c - 'a' + 1) *
                 p_pow) % m;
           p_pow = (p_pow * p) % m;
13
          return hash_value;
13
    9 }
13
```

### 1.3 Rabin Karp

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```
vector<ll> rabin_karp(string const& s, string const
        & t) {
     const 11 p = 31, m = 1e9 + 9;
     11 S = s.size(), T = t.size();
     vector<ll> p_pow(max(S, T));
     p_pow[0] = 1;
     for (ll i = 1; i < (ll) p_pow.size(); i++) p_pow[</pre>
          i] = (p_pow[i-1] * p) % m;
     vector<ll> h(T + 1, 0);
     for (ll i = 0; i < T; i++) h[i + 1] = (h[i] + (t[
          i] - 'a' + 1) * p_pow[i]) % m;
     11 h s = 0;
     for (11 i = 0; i < S; i++) h_s = (h_s + (s[i] - '
          a' + 1) * p_pow[i]) % m;
     vector<11> occurences;
     for (11 i = 0; i + S - 1 < T; i++) {
       11 \text{ cur}_h = (h[i + S] + m - h[i]) % m;
       if (cur_h == h_s * p_pow[i] % m) occurences.
           push_back(i);
     return occurences;
```

### 1.4 Z Function

```
vector<int> z_function(string s) {
  int n = s.size();
  vector<int> z(n);
  for (int i = 1, l = 0, r = 0; i < n; i++) {
    if (i < r) z[i] = min(r - i, z[i - 1]);
    while (i + z[i] < n \&\& s[z[i]] == s[i + z[i]])
        z[i]++;
    if (i + z[i] > r) {
     1 = i;
      r = i + z[i];
  return z;
```

### 1.5 Finding Repetitions

```
vector<int> z_function(string const& s) {
  int n = s.size();
  vector<int> z(n);
  for (int i = 1, l = 0, r = 0; i < n; i++) {
    if (i \le r) z[i] = min(r - i + 1, z[i - 1]);
    while (i + z[i] < n \&\& s[z[i]] == s[i + z[i]])
         z[i]++;
```

```
if (i + z[i] - 1 > r) {
         1 = i;
          r = i + z[i] - 1;
11
12
      return z;
13 }
14 int get_z(vector<int> const& z, int i) {
      if (0 <= i && i < (int) z.size()) return z[i];</pre>
16
      else return 0;
17 }
18 vector<pair<int, int>> repetitions;
19 void convert_to_repetitions(int shift, bool left,
         int cntr, int 1, int k1, int k2) {
      for (int 11 = max(1, 1 - k2); 11 <= min(1, k1);</pre>
           11++) {
21
        if (left && 11 == 1) break;
22
        int 12 = 1 - 11;
        int pos = shift + (left ? cntr - 11 : cntr - 1
             -11+1);
        repetitions.emplace_back(pos, pos + 2 * 1 - 1);
25
26
void find_repetitions(string s, int shift = 0) {
28
     int n = s.size();
29
     if (n == 1) return;
     int nu = n / 2;
     int nv = n - nu;
     string u = s.substr(0, nu);
     string v = s.substr(nu);
     string ru(u.rbegin(), u.rend());
      string rv(v.rbegin(), v.rend());
      find_repetitions(u, shift);
      find repetitions (v, shift + nu);
      vector<int> z1 = z_function(ru);
      vector<int> z2 = z_function(v + '#' + u);
39
40
      vector<int> z3 = z_function(ru + '#' + rv);
41
      vector<int> z4 = z_function(v);
42
      for (int cntr = 0; cntr < n; cntr++) {</pre>
43
       int 1, k1, k2;
44
        if (cntr < nu) {</pre>
45
         1 = nu - cntr;
          k1 = get_z(z1, nu - cntr);
47
          k2 = get_z(z2, nv + 1 + cntr);
48
49
          1 = cntr - nu + 1;
50
          k1 = get_z(z3, nu + 1 + nv - 1 - (cntr - nu))
51
          k2 = get_z(z4, (cntr - nu) + 1);
        if (k1 + k2 >= 1) convert_to_repetitions(shift,
              cntr < nu, cntr, 1, k1, k2);</pre>
55 }
```

### 1.6 Longest Common Prefix

### 1.7 Suffix Array

```
vector<int> sort_cyclic_shifts(string const& s) {
      int n = s.size();
      const int alphabet = 256;
      vector<int> p(n), c(n), cnt(max(alphabet, n), 0);
      for (int i = 0; i < n; i++) cnt[s[i]]++;</pre>
      for (int i = 1; i < alphabet; i++) cnt[i] += cnt[</pre>
           i - 11;
      for (int i = 0; i < n; i++) p[--cnt[s[i]]] = i;</pre>
      c[p[0]] = 0;
      int classes = 1;
      for (int i = 1; i < n; i++) {
       if (s[p[i]] != s[p[i-1]]) classes++;
        c[p[i]] = classes - 1;
13
14
      vector<int> pn(n), cn(n);
      for (int h = 0; (1 << h) < n; ++h) {
        for (int i = 0; i < n; i++) {</pre>
          pn[i] = p[i] - (1 << h);
          if (pn[i] < 0)
            pn[i] += n;
        fill(cnt.begin(), cnt.begin() + classes, 0);
        for (int i = 0; i < n; i++) cnt[c[pn[i]]]++;</pre>
        for (int i = 1; i < classes; i++) cnt[i] += cnt</pre>
             [i - 1];
        for (int i = n-1; i >= 0; i--) p[--cnt[c[pn[i
             ]]]] = pn[i];
        cn[p[0]] = 0;
        classes = 1;
        for (int i = 1; i < n; i++) {</pre>
          pair<int, int> cur = {c[p[i]], c[(p[i] + (1)
               << h)) % n]};
          pair < int, int > prev = {c[p[i-1]], c[(p[i-1]] + }
                 (1 << h)) % n]};
          if (cur != prev) ++classes;
          cn[p[i]] = classes - 1;
32
33
        c.swap(cn);
34
      }
35
      return p;
36
    vector<int> build_suff_arr(string s) {
38
      s += "$";
      vector<int> sorted_shifts = sort_cyclic_shifts(s)
40
      sorted_shifts.erase(sorted_shifts.begin());
41
      return sorted_shifts;
43 // compare two substrings
44 int compare(int i, int j, int l, int k) {
      pair<int, int> a = \{c[k][i], c[k][(i + 1 - (1 <<
           k)) % n]};
      pair<int, int> b = \{c[k][j], c[k][(j + 1 - (1 <<
```

```
47 return a == b ? 0 : a < b ? -1 : 1;
48 }
```

### 1.8 Count Unique Substrings

```
int count unique substrings(string const& s) {
      int n = s.size();
      const int p = 31;
      const int m = 1e9 + 9;
      vector<long long> p_pow(n);
      p_pow[0] = 1;
      for (int i = 1; i < n; i++) p_pow[i] = (p_pow[i -</pre>
            1] * p) % m;
      vector<long long> h(n + 1, 0);
      for (int i = 0; i < n; i++) h[i + 1] = (h[i] + (s
            [i] - 'a' + 1) * p_pow[i]) % m;
      int cnt = 0;
      for (int 1 = 1; 1 <= n; 1++) {</pre>
        unordered_set<long long> hs;
        for (int i = 0; i <= n - 1; i++) {</pre>
           long long cur_h = (h[i + 1] + m - h[i]) % m;
           \operatorname{cur}_h = (\operatorname{cur}_h * \operatorname{p_pow}[n - i - 1]) % m;
           hs.insert(cur_h);
        cnt += hs.size();
      return cnt;
21 }
```

#### 1.9 Knuth Morris Pratt

```
1 vector<11> prefix_function(string s) {
      11 n = (11) s.length();
      vector<ll> pi(n);
      for (ll i = 1; i < n; i++) {
       11 j = pi[i - 1];
       while (j > 0 \&\& s[i] != s[j]) j = pi[j - 1];
       if (s[i] == s[j]) j++;
       pi[i] = j;
     return pi;
11 }
12 // count occurences
13 vector < int > ans(n + 1);
14 for (int i = 0; i < n; i++)
    ans[pi[i]]++;
16 for (int i = n-1; i > 0; i--)
    ans[pi[i-1]] += ans[i];
18 for (int i = 0; i <= n; i++)
    ans[i]++;
```

### 1.10 Group Identical Substrings

```
vector<vector<int>> group_identical_strings(vector
string> const& s) {
int n = s.size();
vector<pair<long long, int>> hashes(n);
for (int i = 0; i < n; i++) hashes[i] = {
    compute_hash(s[i]), i};
sort(hashes.begin(), hashes.end());
vector<vector<int>> groups;
for (int i = 0; i < n; i++) {</pre>
```

### 2 Geometry

#### 2.1 Nearest Points

```
struct pt {
                11 x, y, id;
            struct cmp_x {
                 bool operator()(const pt & a, const pt & b) const
                       return a.x < b.x || (a.x == b.x && a.y < b.y);</pre>
   8
          };
            struct cmp v {
                 bool operator()(const pt & a, const pt & b) const
                               { return a.y < b.y; }
11 };
12 11 n;
13 vector<pt> a;
14 double mindist;
15 pair<11, 11> best_pair;
void upd_ans(const pt & a, const pt & b) {
                 double dist = sqrt((a.x - b.x) * (a.x - b.x) + (a.x - 
                             .y - b.y) * (a.y - b.y));
                 if (dist < mindist) {</pre>
19
                      mindist = dist:
20
                      best_pair = {a.id, b.id};
21
22
23
          vector<pt> t;
24 void rec(ll 1, ll r) {
25
                if (r - 1 \le 3) {
26
                      for (11 i = 1; i < r; ++i)</pre>
27
                           for (11 \ j = i + 1; \ j < r; ++j)
28
                                 upd_ans(a[i], a[j]);
29
                       sort(a.begin() + 1, a.begin() + r, cmp_y());
                11 m = (1 + r) >> 1, midx = a[m].x;
33
                rec(1, m);
34
                rec(m, r);
                 merge(a.begin() + 1, a.begin() + m, a.begin() + m
                              , a.begin() + r, t.begin(), cmp_y());
                 copy(t.begin(), t.begin() + r - 1, a.begin() + 1)
37
                 11 \text{ tsz} = 0;
38
                 for (11 i = 1; i < r; ++i) {
39
                     if (abs(a[i].x - midx) < mindist) {</pre>
40
                            for (11 j = tsz - 1; j >= 0 && a[i].y - t[j].
                                       y < mindist; --j)
                                 upd_ans(a[i], t[j]);
42
                            t[tsz++] = a[i];
43
               }
44
45 }
46 t.resize(n);
48 mindist = 1E20;
49 rec(0, n);
```

# 2.2 Minkowski Sum

```
struct pt {
       11 x, y;
       pt operator + (const pt & p) const { return pt {x}
            + p.x, y + p.y; }
       pt operator - (const pt & p) const { return pt {x}
            -p.x, y - p.y; }
       11 cross(const pt & p) const { return x * p.y - y
             * p.x; }
 6
    };
    void reorder_polygon(vector<pt> & P) {
 8
       size_t pos = 0;
       for (size_t i = 1; i < P.size(); i++) {</pre>
         if (P[i].y < P[pos].y || (P[i].y == P[pos].y &&
               P[i].x < P[pos].x)) pos = i;
       rotate(P.begin(), P.begin() + pos, P.end());
    vector<pt> minkowski(vector<pt> P, vector<pt> Q) {
      // the first vertex must be the lowest
       reorder_polygon(P);
      reorder_polygon(Q);
       // we must ensure cyclic indexing
      P.push_back(P[0]);
      P.push_back(P[1]);
       0.push back(0[0]);
       Q.push_back(Q[1]);
       // main part
       vector<pt> result;
25
       size_t i = 0, j = 0;
       while (i < P.size() - 2 || j < Q.size() - 2){</pre>
        result.push_back(P[i] + Q[j]);
         auto cross = (P[i + 1] - P[i]).cross(Q[j + 1] -
               ([i]O
         if (cross >= 0 && i < P.size() - 2) ++i;</pre>
         if (cross <= 0 && j < Q.size() - 2) ++j;</pre>
32
      return result;
33 }
```

### 2.3 Point In Convex

```
struct pt {
      long long x, y;
      pt() {}
      pt (long long \underline{x}, long long \underline{y}) : x(\underline{x}), y(\underline{y}) {}
      pt operator+(const pt &p) const { return pt(x + p
            .x, y + p.y); }
      pt operator-(const pt &p) const { return pt(x - p
            .x, y - p.y);
      long long cross(const pt &p) const { return x * p
            y - y * p.x;
      long long dot(const pt &p) const { return x * p.x
            + y * p.y; }
      long long cross(const pt &a, const pt &b) const {
            return (a - *this).cross(b - *this); }
10
      long long dot(const pt &a, const pt &b) const {
            return (a - *this).dot(b - *this); }
11
      long long sqrLen() const { return this->dot(*this
           ); }
12
    bool lexComp(const pt &1, const pt &r) { return 1.x
           < r.x \mid | (1.x == r.x && 1.y < r.y); }
```

```
4 int sgn(long long val) { return val > 0 ? 1 : (val
         == 0 ? 0 : -1); }
   vector<pt> seq;
   pt translation;
   int n:
18
   bool pointInTriangle(pt a, pt b, pt c, pt point) {
      long long s1 = abs(a.cross(b, c));
      long long s2 = abs(point.cross(a, b)) + abs(point
           .cross(b, c)) + abs(point.cross(c, a));
      return s1 == s2;
    void prepare(vector<pt> &points) {
      n = points.size();
      int pos = 0;
      for (int i = 1; i < n; i++) {
        if (lexComp(points[i], points[pos])) pos = i;
      rotate(points.begin(), points.begin() + pos,
          points.end());
      seq.resize(n);
      for (int i = 0; i < n; i++) seq[i] = points[i +</pre>
          1] - points[0];
      translation = points[0];
34
    bool pointInConvexPolygon(pt point) {
      point = point - translation;
      if (seq[0].cross(point) != 0 && sgn(seq[0].cross(
          point)) != sgn(seq[0].cross(seq[n - 1])))
3.8
        return false;
      if (seg[n-1].cross(point) != 0 && sgn(seg[n-1])
           1].cross(point)) != sqn(seq[n - 1].cross(seq
        return false;
      if (seq[0].cross(point) == 0)
        return seq[0].sqrLen() >= point.sqrLen();
      int 1 = 0, r = n - 1;
      while (r - 1 > 1) {
44
       int mid = (1 + r) / 2;
        int pos = mid;
        if (seq[pos].cross(point) >= 0) 1 = mid;
        else r = mid;
50
      int pos = 1;
      return pointInTriangle(seg[pos], seg[pos + 1], pt
           (0, 0), point);
52 }
```

### 2.4 Line Sweep

```
1  const double EPS = 1E-9;
2  struct pt { double x, y; };
3  struct seg {
4   pt p, q;
5   11 id;
6   double get_y(double x) const {
7   if (abs(p.x - q.x) < EPS) return p.y;
8   return p.y + (q.y - p.y) * (x - p.x) / (q.x - p.x);
9  }
10  };
11  bool intersect1d(double 11, double r1, double 12, double r2) {
12   if (11 > r1) swap(11, r1);
13   if (12 > r2) swap(12, r2);
14   return max(11, 12) <= min(r1, r2) + EPS;
15  }</pre>
```

```
16 11 vec(const pt& a, const pt& b, const pt& c) {
      double s = (b.x - a.x) * (c.y - a.y) - (b.y - a.y)
           ) * (c.x - a.x);
      return abs(s) < EPS ? 0 : s > 0 ? +1 : -1;
19 1
20 bool intersect(const seg& a, const seg& b) {
21
      return intersect1d(a.p.x, a.q.x, b.p.x, b.q.x) &&
             intersect1d(a.p.y, a.q.y, b.p.y, b.q.y) &&
23
             vec(a.p, a.q, b.p) * vec(a.p, a.q, b.q) <=</pre>
                   3.3 0
24
              vec(b.p, b.q, a.p) * vec(b.p, b.q, a.q) <=
25 }
26 bool operator<(const seq& a, const seq& b) {
      double x = max(min(a.p.x, a.g.x), min(b.p.x, b.g.
28
      return a.get_y(x) < b.get_y(x) - EPS;</pre>
29 }
30 struct event {
      double x;
      11 tp, id;
      event() {}
34
      event (double x, 11 tp, 11 id) : x(x), tp(tp), id(
      bool operator<(const event& e) const {</pre>
36
        if (abs(x - e.x) > EPS) return x < e.x;
        return tp > e.tp;
38
39 };
40 set<seg> s;
41 vector<set<seg>::iterator> where;
    set<seg>::iterator prev(set<seg>::iterator it) {
      return it == s.begin() ? s.end() : --it;
44 }
45 set<seg>::iterator next(set<seg>::iterator it) {
46
     return ++it;
47 }
48 pair<11, 11> solve(const vector<seg>& a) {
49
      11 n = (11) a.size();
      vector<event> e;
      for (11 i = 0; i < n; ++i) {
        e.push_back(event(min(a[i].p.x, a[i].q.x), +1,
        e.push_back(event(max(a[i].p.x, a[i].q.x), -1,
54
      sort(e.begin(), e.end());
      s.clear():
      where.resize(a.size());
      for (size_t i = 0; i < e.size(); ++i) {</pre>
59
       11 \text{ id} = e[i].id;
60
        if (e[i].tp == +1) {
          set<seq>::iterator nxt = s.lower_bound(a[id])
61
               , prv = prev(nxt);
62
          if (nxt != s.end() && intersect(*nxt, a[id]))
                return make_pair(nxt->id, id);
63
          if (prv != s.end() && intersect(*prv, a[id]))
                return make pair (prv->id, id);
64
          where[id] = s.insert(nxt, a[id]);
65
         } else {
66
          set<seg>::iterator nxt = next(where[id]), prv
                = prev(where[id]);
          if (nxt != s.end() && prv != s.end() &&
               intersect(*nxt, *prv)) return make_pair(
               prv->id, nxt->id);
68
          s.erase(where[id]);
69
71
      return make_pair(-1, -1);
```

```
17 }
```

```
2.5 Line Intersection
    struct pt { double x, y; };
    struct line { double a, b, c; };
    const double EPS = 1e-9;
    double det (double a, double b, double c, double d)
         { return a*d - b*c; }
   bool intersect(line m, line n, pt & res) {
      double zn = det(m.a, m.b, n.a, n.b);
      if (abs(zn) < EPS) return false;</pre>
      res.x = -det(m.c, m.b, n.c, n.b) / zn;
      res.y = -det(m.a, m.c, n.a, n.c) / zn;
     return true;
12 bool parallel(line m, line n) { return abs(det(m.a,
         m.b, n.a, n.b)) < EPS; }
13 bool equivalent(line m, line n) {
      return abs(det(m.a, m.b, n.a, n.b)) < EPS
          && abs(det(m.a, m.c, n.a, n.c)) < EPS
          && abs(det(m.b, m.c, n.b, n.c)) < EPS;
2.6 Basic Geometry
```

### struct point2d { ftype x, y; point2d() {} point2d(ftype x, ftype y): x(x), y(y) {} point2d& operator+=(const point2d &t) { x += t.x: y += t.y; return \*this; point2d& operator-=(const point2d &t) { x -= t.x: y -= t.y; return \*this; point2d& operator\*=(ftype t) { y \*= t;

return \*this;

return \*this;

x /= t;

y /= t;

point2d& operator/=(ftype t) {

(\*this) \*= t; }

(\*this) /= t; }

+ a.y \* b.y + a.z \* b.z; }

+ a.v \* b.v;

return point2d(\*this) += t; }

return point2d(\*this) -= t; }

```
point2d operator+(const point2d &t) const {
      point2d operator-(const point2d &t) const {
      point2d operator*(ftype t) const { return point2d
      point2d operator/(ftype t) const { return point2d
   point2d operator*(ftype a, point2d b) { return b *
31 ftype dot(point2d a, point2d b) { return a.x * b.x
32 ftype dot(point3d a, point3d b) { return a.x * b.x
```

```
33 ftype norm(point2d a) { return dot(a, a); }
   double abs(point2d a) { return sgrt(norm(a)); }
   double proj(point2d a, point2d b) { return dot(a, b
        ) / abs(b); }
36 double angle(point2d a, point2d b) { return acos(
         dot(a, b) / abs(a) / abs(b)); }
37 point3d cross(point3d a, point3d b) { return
         point3d(a.y \star b.z - a.z \star b.y, a.z \star b.x - a.x
          * b.z, a.x * b.y - a.y * b.x); }
38 ftype triple(point3d a, point3d b, point3d c) {
         return dot(a, cross(b, c)); }
39 ftype cross(point2d a, point2d b) { return a.x * b.
         y - a.y * b.x; }
40 point2d intersect(point2d al, point2d dl, point2d
         a2, point2d d2) { return a1 + cross(a2 - a1,
         d2) / cross(d1, d2) * d1; }
41 point3d intersect(point3d a1, point3d n1, point3d
         a2, point3d n2, point3d a3, point3d n3) {
      point3d x(n1.x, n2.x, n3.x);
      point3d y(n1.y, n2.y, n3.y);
      point3d z(n1.z, n2.z, n3.z);
      point3d d(dot(a1, n1), dot(a2, n2), dot(a3, n3));
      return point3d(triple(d, v, z), triple(x, d, z),
          triple(x, y, d)) / triple(n1, n2, n3);
```

### 2.7 Circle Line Intersection

```
double r, a, b, c; // given as input
   double x0 = -a * c / (a * a + b * b);
   double y0 = -b * c / (a * a + b * b);
   if (c * c > r * r * (a * a + b * b) + EPS) {
     puts ("no points");
   } else if (abs (c *c - r * r * (a * a + b * b)) <
        EPS) {
     puts ("1 point");
     cout << x0 << ' ' << y0 << '\n';
     double d = r * r - c * c / (a * a + b * b);
     double mult = sqrt (d / (a * a + b * b));
     double ax, ay, bx, by;
     ax = x0 + b * mult;
     bx = x0 - b * mult;
     ay = y0 - a * mult;
     by = y0 + a * mult;
     puts ("2 points");
     cout << ax << ' ' << ay << '\n' << bx << ' ' <<
          by << '\n';
19 }
```

#### 2.8 Convex Hull

```
1 struct pt {
     double x, y;
   11 orientation(pt a, pt b, pt c) {
     double v = a.x * (b.y - c.y) + b.x * (c.y - a.y)
         + c.x * (a.y - b.y);
     if (v < 0) {
       return -1:
     \} else if (v > 0) {
       return +1;
     return 0;
```

```
13 bool cw(pt a, pt b, pt c, bool include_collinear) {
14
      11 o = orientation(a, b, c);
15
      return o < 0 || (include_collinear && o == 0);</pre>
16
17
    bool collinear(pt a, pt b, pt c) {
1.8
      return orientation(a, b, c) == 0;
19 }
20 void convex_hull(vector<pt>& a, bool
         include_collinear = false) {
      pt p0 = *min_element(a.begin(), a.end(), [](pt a,
        return make_pair(a.y, a.x) < make_pair(b.y, b.x</pre>
23
      sort(a.begin(), a.end(), [&p0](const pt& a, const
25
        11 o = orientation(p0, a, b);
26
        if (o == 0) {
          return (p0.x - a.x) * (p0.x - a.x) + (p0.y - a.x)
               a.y) * (p0.y - a.y)
               < (p0.x - b.x) * (p0.x - b.x) + (p0.y -
                    b.y) * (p0.y - b.y);
        return o < 0;
31
      if (include_collinear) {
        11 i = (11) a.size()-1;
34
        while (i \geq= 0 && collinear(p0, a[i], a.back()))
        reverse(a.begin()+i+1, a.end());
36
      vector<pt> st;
38
      for (ll i = 0; i < (ll) a.size(); i++) {</pre>
39
        while (st.size() > 1 && !cw(st[st.size() - 2],
             st.back(), a[i], include_collinear)) {
40
          st.pop_back();
41
42
        st.push_back(a[i]);
43
44
      a = st;
45 }
```

### 2.9 Count Lattices

```
1 int count_lattices(Fraction k, Fraction b, long
         long n) {
      auto fk = k.floor();
      auto fb = b.floor();
      auto cnt = 0LL;
      if (k >= 1 || b >= 1) {
        cnt += (fk * (n - 1) + 2 * fb) * n / 2;
        k = fk;
        b -= fb;
      auto t = k * n + b;
      auto ft = t.floor();
      if (ft >= 1) cnt += count_lattices(1 / k, (t - t.
           floor()) / k, t.floor());
13
      return cnt;
14 }
```

#### 2.10 Segment Intersection

```
1 const double EPS = 1E-9;
2 struct pt {
```

```
double x, y;
      bool operator<(const pt& p) const {</pre>
        return x < p.x - EPS \mid \mid (abs(x - p.x) < EPS &&
             y < p.y - EPS);
 6
 7
    };
    struct line {
      double a, b, c;
      line() {}
      line(pt p, pt q) {
        a = p.y - q.y;
        b = q.x - p.x;
        c = -a * p.x - b * p.y;
      void norm() {
18
        double z = sqrt(a * a + b * b);
        if (abs(z) > EPS) a /= z, b /= z, c /= z;
      double dist(pt p) const { return a * p.x + b * p.
           y + c; }
    };
    double det (double a, double b, double c, double d)
      return a * d - b * c;
    inline bool betw(double 1, double r, double x) {
      return min(1, r) \le x + EPS \&\& x \le max(1, r) +
28
    inline bool intersect_1d(double a, double b, double
          c, double d) {
      if (a > b) swap(a, b);
      if (c > d) swap(c, d);
      return max(a, c) <= min(b, d) + EPS;</pre>
34 bool intersect(pt a, pt b, pt c, pt d, pt& left, pt
         & right) {
      if (!intersect_ld(a.x, b.x, c.x, d.x) || !
           intersect_ld(a.y, b.y, c.y, d.y)) return
           false:
      line m(a, b);
      line n(c, d);
      double zn = det(m.a, m.b, n.a, n.b);
      if (abs(zn) < EPS) {
        if (abs(m.dist(c)) > EPS || abs(n.dist(a)) >
             EPS) return false;
        if (b < a) swap(a, b);
        if (d < c) swap(c, d);
        left = max(a, c);
        right = min(b, d);
45
        return true;
      } else {
        left.x = right.x = -det(m.c, m.b, n.c, n.b) /
48
        left.y = right.y = -det(m.a, m.c, n.a, n.c) /
49
        return betw(a.x, b.x, left.x) && betw(a.y, b.y,
              left.y) &&
               betw(c.x, d.x, left.x) && betw(c.y, d.y,
                     left.y);
52 }
```

#### 2.11 Areas

```
int signed_area_parallelogram(point2d p1, point2d
```

```
p2, point2d p3) {
      return cross (p2 - p1, p3 - p2);
    double triangle_area (point2d p1, point2d p2,
         point2d p3) {
      return abs(signed_area_parallelogram(p1, p2, p3))
            / 2.0;
    bool clockwise (point2d p1, point2d p2, point2d p3)
      return signed_area_parallelogram(p1, p2, p3) < 0;</pre>
10 bool counter_clockwise(point2d p1, point2d p2,
         point2d p3) {
      return signed_area_parallelogram(p1, p2, p3) > 0;
    double area(const vector<point>& fig) {
      double res = 0;
      for (unsigned i = 0; i < fig.size(); i++) {</pre>
        point p = i ? fig[i - 1] : fig.back();
        point q = fig[i];
        res += (p.x - q.x) * (p.y + q.y);
      return fabs(res) / 2;
21 }
```

### 3 Dynamic Programming

### 3.1 Knuth Optimization

```
1 11 solve() {
      11 N:
      ... // Read input
      vector<vector<ll>> dp(N, vector<ll>(N)), opt(N,
           vector<ll>(N));
      auto C = [\&](11 i, 11 j) {
        ... // Implement cost function C.
      };
8
      for (11 i = 0; i < N; i++) {
        opt[i][i] = i;
        ... // Initialize dp[i][i] according to the
             problem
      for (11 i = N - 2; i >= 0; i--) {
        for (11 j = i + 1; j < N; j++) {
          11 \text{ mn} = 11\_\text{MAX}, \text{ cost} = C(i, j);
          for (11 k = opt[i][j-1]; k \le min(j-1,
               opt[i + 1][j]); k++) {
            if (mn \ge dp[i][k] + dp[k + 1][j] + cost) {
              opt[i][j] = k;
              mn = dp[i][k] + dp[k + 1][j] + cost;
          dp[i][j] = mn;
      cout << dp[0][N - 1] << '\n';
25
```

#### 3.2 Knapsack

### 3.3 Divide And Conquer

```
1 11 m, n;
    vector<ll> dp_before(n), dp_cur(n);
    11 C(11 i, 11 j);
    void compute(ll 1, ll r, ll optl, ll optr) {
      if (1 > r) return;
      11 \text{ mid} = (1 + r) >> 1;
      pair<11, 11> best = {LLONG_MAX, -1};
      for (ll k = optl; k <= min(mid, optr); k++)</pre>
      best = min(best, \{(k ? dp\_before[k - 1] : 0) +
             C(k, mid), k});
      dp_cur[mid] = best.first;
11
      11 opt = best.second;
      compute(1, mid - 1, optl, opt);
13
      compute(mid + 1, r, opt, optr);
14
15 11 solve() {
      for (ll i = 0; i < n; i++) dp_before[i] = C(0, i)</pre>
      for (11 i = 1; i < m; i++) {
        compute(0, n - 1, 0, n - 1);
18
19
        dp_before = dp_cur;
20
21
      return dp_before[n - 1];
```

### 3.4 Digit Dp

```
1 vector<vector<vector<ll>>>> dp(K + 1, vector
         <vector<vector<11>>>(9 * K + 1, vector<vector<</pre>
         11>> (9 * K + 1, vector<11>(9 * K, 0)));
    for (11 n = 1; n \le 9 * K; n++) dp[0][n][0][0] = 1;
    11 pow10 = 1;
    for (11 k = 1; k <= K; k++) {
      for (11 n = 1; n \le 9 * K; n++) {
        for (11 s = 0; s <= 9 * K; s++) {
          for (11 m = 0; m < n; m++) {
            for (11 y = 0; y \le 9; y++) {
                if (s \ge y) dp[k][n][s][m] += dp[k -
                     1][n][s - y][((m - y * pow10) % n]
                     + n) % n];
13
14
      pow10 *= 10;
15 }
16 string N;
17 cin \gg N;
18 11 n = N.length(), ans = 0;
19 vector<11> g(9 * K + 1, 0);
20 for (11 s = 1; s \leq 9 * K; s++) {
```

string substring = "";

```
11 pow10 = 1;
      for (11 i = 0; i < n - 1; i++) pow10 *= 10;
      for (ll i = 0; i < n; i++) {
         substring += '0';
         for (11 \ \dot{j} = 0; \ \dot{j} < N[i] - '0'; \ \dot{j}++)  {
          11 digit_sum = j;
           for (11 k = 0; k < i; k++) digit_sum +=</pre>
                substring[k] - '0';
           if (s \ge digit\_sum) g[s] += dp[n - 1 - i][s][
                s - digit_sum][((-pow10 * stoll(
                substring)) % s + s) % s];
           substring[i]++;
        pow10 /= 10;
34
      ans += q[s];
35 }
36
    auto is_good = [&](string s) -> bool {
      11 \text{ digit\_sum} = 0;
      for (ll i = 0; i < (ll) s.length(); i++)</pre>
           digit_sum += s[i] - '0';
      return stoll(s) % digit_sum == 0;
    if (is_good(N)) ans++;
    cout << ans << "\n";
```

#### 3.5 Subset Sum

#### 3.6 Longest Increasing Subsequence

```
1  ll get_ceil_idx(vector<ll> &a, vector<ll> &T, ll l,
         11 r, 11 x) {
     while (r - 1 > 1) {
      11 m = 1 + (r - 1) / 2;
       if (a[T[m]] >= x) {
       r = m;
       } else {
         1 = m;
8
     return r;
  11 LIS(ll n, vector<ll> &a) {
    11 len = 1;
     vector<ll> T(n, 0), R(n, -1);
     T[0] = 0;
     for (ll i = 1; i < n; i++) {
       if (a[i] < a[T[0]]) {</pre>
        T[0] = i;
```

} else if (a[i] > a[T[len - 1]]) {

#### 3.7 Longest Common Subsequence

```
1 11 LCS(string x, string y, 11 n, 11 m) {
      vector<vector<ll>> dp(n + 1, vector<ll>(m + 1));
      for (11 i = 0; i <= n; i++) {</pre>
        for (11 j = 0; j \le m; j++) {
          if (i == 0 || j == 0) {
            dp[i][j] = 0;
          } else if (x[i - 1] == y[j - 1]) {
            dp[i][j] = dp[i - 1][j - 1] + 1;
          } else {
            dp[i][j] = max(dp[i - 1][j], dp[i][j - 1]);
      11 \text{ index} = dp[n][m];
      vector<char> lcs(index + 1);
      lcs[index] = ' \setminus 0';
      11 i = n, j = m;
      while (i > 0 \&\& j > 0) {
       if (x[i-1] == y[j-1]) {
          lcs[index - 1] = x[i - 1];
          i--;
          j--;
          index--;
        } else if (dp[i - 1][j] > dp[i][j - 1]) {
        } else {
28
29
      return dp[n][m];
31
```

#### 3.8 Max Sum

```
int max_subarray_sum(vi arr) {
   int x = 0, s = 0;
   for (int k = 0; k < n; k++) {
      s = max(arr[k], s+arr[k]);
      x = max(x, s);
   }
   return x;
}</pre>
```

### 3.9 Bitmask Weights

```
1 vector<pair<11, 11>> dp(1 << n, {INF, 0});
2 dp[0] = {1, 0};
3 for (11 mask = 1; mask < (1 << n); mask++)</pre>
```

#### 3.10 Edit Distance

```
1  ll edit_distance(string x, string y, ll n, ll m) {
      vector<vector<int>> dp(n + 1, vector<int>(m + 1,
       dp[0][0] = 0;
       for (int i = 1; i <= n; i++) {</pre>
        dp[i][0] = i;
 6
      for (int j = 1; j <= m; j++) {</pre>
 8
        dp[0][j] = j;
 9
10
      for (int i = 1; i <= n; i++) {</pre>
11
         for (int j = 1; j <= m; j++) {</pre>
           dp[i][j] = min({dp[i-1][j] + 1, dp[i][j-1]}
                1] + 1, dp[i - 1][j - 1] + (x[i - 1] !=
               y[j - 1])));
13
14
15
      return dp[n][m];
16 }
```

### 4 Math

#### 4.1 Chinese Remainder Theorem

```
struct Congruence {
      11 a, m;
 3 };
 5 11 chinese_remainder_theorem(vector<Congruence>
         const& congruences) {
      11. M = 1:
      for (auto const& congruence : congruences) M *=
           congruence.m:
      11 \text{ solution} = 0;
      for (auto const& congruence : congruences) {
       11 a_i = congruence.a;
        11 M_i = M / congruence.m;
        11 N_i = mod_inv(M_i, congruence.m);
        solution = (solution + a_i * M_i % M * N_i) % M
14
15
      return solution;
16 }
```

### 4.2 Extended Euclidean

```
1 int gcd(int a, int b, int& x, int& y) {
2    if (b == 0) {
3         x = 1;
4         y = 0;
5         return a;
6    }
7    int x1, y1, d = gcd(b, a % b, x1, y1);
8         x = y1;
9         y = x1 - y1 * (a / b);
10    return d;
11 }
```

#### 4.3 Modulo Inverse

```
1  11 mod_inv(11 a, 11 m) {
2    if (m == 1) return 0;
3    11 m0 = m, x = 1, y = 0;
4    while (a > 1) {
5        11 q = a / m, t = m;
6        m = a % m;
7        a = t;
8        t = y;
9        y = x - q * y;
10        x = t;
11    }
12    if (x < 0) x += m0;
13    return x;
14 }</pre>
```

#### 4.4 Sum Of Divisors

```
1 ll sum of divisors(ll num) {
      11 total = 1;
      for (int i = 2; (11) i * i <= num; i++) {</pre>
        if (num % i == 0) {
           int e = 0;
           do {
            e++;
            num /= i;
9
           } while (num % i == 0);
          11 \text{ sum} = 0, \text{ pow} = 1;
           do {
            sum += pow;
            pow *= i;
          } while (e-- > 0);
15
          total *= sum;
16
18
      if (num > 1) total *= (1 + num);
      return total:
20 }
```

### 4.5 Range Sieve

```
vector<bool> range_sieve(11 1, 11 r) {
    11 n = sqrt(r);
    vector<bool> is_prime(n + 1, true);
    vector<11> prime;
    is_prime[0] = is_prime[1] = false;
    prime.push_back(2);
    for (11 i = 4; i <= n; i += 2) is_prime[i] =
        false;</pre>
```

### 4.6 Pollard Rho Brent

```
1 ll mult(ll a, ll b, ll mod) {
    return (__int128_t) a * b % mod;
3
4 11 f(11 x, 11 c, 11 mod) {
     return (mult(x, x, mod) + c) % mod;
    ll pollard_rho_brent(ll n, ll x0 = 2, ll c = 1) {
     11 \times = x0, g = 1, q = 1, xs, y, m = 128, 1 = 1;
      while (q == 1) {
        y = x;
        for (11 i = 1; i < 1; i++) x = f(x, c, n);
        11 k = 0;
        while (k < 1 \&\& g == 1) {
          xs = x;
          for (ll i = 0; i < m && i < l - k; i++) {
           x = f(x, c, n);
            q = mult(q, abs(y - x), n);
          g = \underline{gcd}(q, n);
20
          k += m;
21
        1 *= 2:
23
24
      if (g == n) {
        do {
         xs = f(xs, c, n);
          g = \underline{gcd(abs(xs - y), n)};
        } while (q == 1);
      return g;
31 }
```

### 4.7 Factorial Modulo

#### 4.8 Matrix

```
1 /*
    Matrix exponentation:
   f[n] = af[n-1] + bf[n-2] + cf[n-3]
   |f[n] | |a b c||f[n-1]|
   |f[n-1]|=|1 0 0||f[n-2]|
    |f[n-2]| |0 1 0||f[n-3]|
    |f[n] | |a b c|^(n-2)|f[2]|
10 |f[n-1]| = |1 \ 0 \ 0| |f[1]|
11 |f[n-2]| |0 1 0|
                          | f [ 0 ] |
12 */
13 struct Matrix { int mat[MAX_N][MAX_N]; };
14 Matrix matrix mul(Matrix a, Matrix b) {
15
    Matrix ans; int i, j, k;
16
    for (i = 0; i < MAX_N; i++)</pre>
      for (j = 0; j < MAX_N; j++)</pre>
      for (ans.mat[i][j] = k = 0; k < MAX_N; k++)
19
       ans.mat[i][i] += a.mat[i][k] * b.mat[k][i];
20
      return ans:
21
22 Matrix matrix_pow(Matrix base, int p) {
23
      Matrix ans; int i, j;
24
      for (i = 0; i < MAX_N; i++)</pre>
25
       for (j = 0; j < MAX_N; j++)</pre>
         ans.mat[i][j] = (i == j);
      while (p) {
       if (p & 1) ans = matrix mul(ans, base);
29
       base = matrix_mul(base, base);
       p >>= 1;
31
32
      return ans;
33 }
```

#### 4.9 Find All Solutions

```
bool find_any_solution(ll a, ll b, ll c, ll &x0, ll
          &v0, 11 &g) {
      q = qcd_{extended(abs(a), abs(b), x0, y0)};
      if (c % q) return false;
      x0 \star = c / q;
      y0 \star = c / g;
      if (a < 0) x0 = -x0;
      if (b < 0) y0 = -y0;
      return true;
 9
10 void shift_solution(11 & x, 11 & y, 11 a, 11 b, 11
         cnt) {
      x += cnt * b;
     y -= cnt * a;
13
   11 find_all_solutions(11 a, 11 b, 11 c, 11 minx, 11
          maxx, 11 miny, 11 maxy) {
      11 x, y, q;
      if (!find_any_solution(a, b, c, x, y, g)) return
           0;
      a /= g;
      b /= g;
      11 \text{ sign}_a = a > 0 ? +1 : -1;
      11 \text{ sign } b = b > 0 ? +1 : -1;
      shift_solution(x, y, a, b, (minx - x) / b);
      if (x < minx) shift_solution(x, y, a, b, sign_b);</pre>
      if (x > maxx) return 0;
```

```
11 \ 1x1 = x;
      shift solution (x, y, a, b, (maxx - x) / b);
      if (x > maxx) shift_solution(x, y, a, b, -sign_b)
      11 \text{ rx1} = x;
28
      shift_solution(x, y, a, b, -(miny - y) / a);
      if (y < miny) shift_solution(x, y, a, b, -sign_a)</pre>
      if (v > maxy) return 0;
      11 \ 1x2 = x;
      shift_solution(x, y, a, b, -(maxy - y) / a);
      if (y > maxy) shift_solution(x, y, a, b, sign_a);
     11 \text{ rx2} = x;
35
      if (1x2 > rx2) swap(1x2, rx2);
36
      11 1x = max(1x1, 1x2), rx = min(rx1, rx2);
      if (1x > rx) return 0;
      return (rx - 1x) / abs(b) + 1;
39 }
```

#### 4.10 Miller Rabin

```
1 using u64 = uint64 t;
   using u128 = __uint128_t;
   u64 binpower(u64 base, u64 e, u64 mod) {
     u64 \text{ result} = 1;
     base %= mod;
     while (e) {
      if (e & 1) result = (u128) result * base % mod;
       base = (u128) base * base % mod;
9
       e >>= 1;
      return result;
12
   bool check_composite(u64 n, u64 a, u64 d, l1 s) {
     u64 x = binpower(a, d, n);
      if (x == 1 \mid | x == n - 1) return false:
      for (11 r = 1; r < s; r++) {
       x = (u128) x * x % n;
       if (x == n - 1) return false;
      return true;
   bool miller_rabin(u64 n) {
     if (n < 2) return false;</pre>
      11 r = 0;
      u64 d = n - 1;
      while ((d \& 1) == 0) {
       d >>= 1;
       r++;
      for (11 a : {2, 3, 5, 7, 11, 13, 17, 19, 23, 29,
         31, 37}) {
        if (n == a) return true;
32
       if (check_composite(n, a, d, r)) return false;
      return true:
```

### 4.11 Fibonacci

```
4 - d'Ocagne's identity: f[m]f[n+1] - f[m+1]f[n] =
         (-1) ^n f[m-n]
5 - Addition rule: f[n+k] = f[k]f[n+1] + f[k-1]f[n]
6 - k = n case: f[2n] = f[n](f[n+1] + f[n-1])
7 - f[n] \mid f[nk]
8 - f[n] | f[m] => n | m
9 - GCD rule: gcd(f[m], f[n]) = f[gcd(m, n)]
10 - [[1 \ 1], [1 \ 0]]^n = [[f[n+1] \ f[n]], [f[n], f[n]]
         -1]]]
11 - f[2k+1] = f[k+1]^2 + f[k]^2
12 - f[2k] = f[k](f[k+1] + f[k-1]) = f[k](2f[k+1] - f[k])
        k])
13 - Periodic sequence modulo p
14 - sum[i=1..n]f[i] = f[n+2] - 1
15 - sum[i=0..n-1]f[2i+1] = f[2n]
16 - sum[i=1..n]f[2i] = f[2n+1] - 1
17 - sum[i=1..n]f[i]^2 = f[n]f[n+1]
18 Fibonacci encoding:
19 1. Iterate through the Fibonacci numbers from the
        largest to the
20 smallest until you find one less than or equal to n
    2. Suppose this number was F i. Subtract F i from n
         and put a 1
    in the i-2 position of the code word (indexing from
         0 from the
   leftmost to the rightmost bit).
24 3. Repeat until there is no remainder.
   4. Add a final 1 to the codeword to indicate its
    Closed-form: f[n] = (((1 + rt(5))/2)^n - ((1 - rt))^n)
         (5)) / 2)^n/rt(5)
28
29
    struct matrix {
       11 mat[2][2];
        matrix friend operator *(const matrix &a, const
             matrix &b) {
          matrix c;
          for (int i = 0; i < 2; i++) {</pre>
           for (int j = 0; j < 2; j++) {
             c.mat[i][j] = 0;
              for (int k = 0; k < 2; k++) c.mat[i][j]</pre>
                   += a.mat[i][k] * b.mat[k][j];
39
          return c;
40
41
      matrix matpow(matrix base, 11 n) {
       matrix ans{ {
        {1, 0},
         {0, 1}
        } };
        while (n) {
         if (n & 1) ans = ans * base;
          base = base * base;
         n >>= 1;
51
52
        return ans;
53
      11 fib(int n) {
        matrix base{ {
         {1, 1},
         {1, 0}
        return matpow(base, n).mat[0][1];
      pair<int, int> fib (int n) {
```

### 4.12 Fast Fourier Transform

```
1 using cd = complex<double>;
    const double PI = acos(-1);
    void fft(vector<cd>& a, bool invert) {
      int n = a.size();
      if (n == 1) return;
      vector<cd> a0 (n / 2), a1 (n / 2);
      for (int i = 0; 2 * i < n; i++) {
       a0[i] = a[2 * i];
       a1[i] = a[2 * i + 1];
      fft(a0, invert);
      fft(a1, invert);
13
      double ang = 2 * PI / n * (invert ? -1 : 1);
      cd w(1), wn(cos(ang), sin(ang));
      for (int i = 0; 2 * i < n; i++) {
       a[i] = a0[i] + w * a1[i];
17
        a[i + n / 2] = a0[i] - w * a1[i];
18
        if (invert) {
19
          a[i] /= 2;
          a[i + n / 2] /= 2;
21
        w \star = wn;
23
24 }
    vector<int> multiply(vector<int> const& a, vector<</pre>
         int> const& b) {
         vector<cd> fa(a.begin(), a.end()), fb(b.begin()
            , b.end());
27
        int n = 1;
28
        while (n < a.size() + b.size()) n <<= 1;</pre>
29
        fa.resize(n);
        fb.resize(n);
         fft(fa, false);
         fft(fb, false);
         for (int i = 0; i < n; i++) fa[i] *= fb[i];</pre>
         fft(fa, true);
35
        vector<int> result(n);
36
         for (int i = 0; i < n; i++) result[i] = round(</pre>
             fa[i].real());
37
         return result;
38 }
```

### 4.13 Segmented Sieve

```
vector<11> segmented_sieve(11 n) {
   const 11 S = 10000;
   11 nsqrt = sqrt(n);
   vector<char> is_prime(nsqrt + 1, true);
   vector<11> prime;
   is_prime[0] = is_prime[1] = false;
   prime.push_back(2);
   for (11 i = 4; i <= nsqrt; i += 2) {
    is_prime[i] = false;
}</pre>
```

```
for (11 i = 3; i <= nsgrt; i += 2) {</pre>
        if (is_prime[i]) {
          prime.push_back(i);
          for (11 j = i * i; j <= nsqrt; j += i) {
            is_prime[j] = false;
18
19
      vector<11> result;
      vector<char> block(S);
      for (11 k = 0; k * S \le n; k++) {
       fill(block.begin(), block.end(), true);
        for (11 p : prime) {
          for (11 j = max((k * S + p - 1) / p, p) * p -
               k * S; j < S; j += p) {
            block[j] = false;
          }
        if (k == 0) {
          block[0] = block[1] = false;
        for (11 i = 0; i < S && k * S + i <= n; i++) {
          if (block[i]) {
            result.push_back(k * S + i);
36
      return result;
```

### 4.14 Linear Sieve

### 4.15 Tonelli Shanks

```
1 ll legendre(ll a, ll p) {
2    return bin_pow_mod(a, (p - 1) / 2, p);
3 }
4 ll tonelli_shanks(ll n, ll p) {
5    if (legendre(n, p) == p - 1) {
6       return -1;
7    }
8    if (p % 4 == 3) {
7       return bin_pow_mod(n, (p + 1) / 4, p);
10    }
11    Q = p - 1, S = 0;
12    while (Q % 2 == 0) {
13       Q /= 2;
14       S++;
15    }
16    ll z = 2;
```

```
for (; z < p; z++) {</pre>
        if (legendre(z, p) == p - 1) {
19
          break;
21
22
      11 M = S, c = bin_pow_mod(z, Q, p), t =
          bin_pow_mod(n, Q, p), R = bin_pow_mod(n, Q)
           + 1) / 2, p);
      while (t % p != 1) {
        if (t % p == 0) {
          return 0;
        11 i = 1, t2 = t * t % p;
        for (; i < M; i++) {
         if (t2 % p == 1) {
            break;
          t2 = t2 * t2 % p;
        11 b = bin_pow_mod(c, bin_pow_mod(2, M - i - 1,
              p), p);
        M = i;
        c = b * b % p;
        t = t * c % p;
38
        R = R * b % p;
39
40
      return R:
41
```

### 5 Miscellaneous

### 5.1 Techniques

```
Dynamic Programming
   - Bitmask
   - Range
   - Digit
   - Knapsack
 7 Graph Theory
8 - Tree diameter
9 - Reversing edges
10 - Tree re-rooting
11 - DP on trees
12 - DFS tree
13 - Euler tour
14 - Binary Jumping
15 - Centroid
16 - DAG
17 - Condense
18 Data Structures
19 - Multiple information
20 - Binary searching on the tree
21 - 2D range query
22 - SQRT decomposition
   - Small-to-large
   Sorting and searching
   - Sliding window
   - Two pointers
   - Binary search on the answer
28 */
```

### 5.2 Gauss

```
const double EPS = 1e-9;
     const 11 INF = 2;
    11 gauss(vector <vector <double>> a, vector <double>
       11 n = (11) a.size(), m = (11) a[0].size() - 1;
       vector<11> where (m, -1);
       for (11 col = 0, row = 0; col < m && row < n; ++</pre>
            col) {
         11 sel = row;
         for (11 i = row; i < n; ++i) {</pre>
           if (abs(a[i][col]) > abs(a[sel][col])) {
11
12
         if (abs (a[sel][col]) < EPS) {</pre>
14
           continue:
15
16
         for (ll i = col; i <= m; ++i) {</pre>
17
           swap(a[sel][i], a[row][i]);
18
19
         where[col] = row;
20
         for (ll i = 0; i < n; ++i) {</pre>
21
           if (i != row) {
22
             double c = a[i][col] / a[row][col];
23
             for (ll j = col; j <= m; ++j) {</pre>
24
               a[i][j] = a[row][j] * c;
25
26
27
28
         ++row;
29
      ans.assign(m, 0);
31
       for (ll i = 0; i < m; ++i) {</pre>
         if (where[i] != -1) {
           ans[i] = a[where[i]][m] / a[where[i]][i];
34
36
      for (11 i = 0; i < n; ++i) {
37
         double sum = 0;
38
         for (11 j = 0; j < m; ++j) {
39
           sum += ans[j] * a[i][j];
40
41
         if (abs (sum - a[i][m]) > EPS) {
42
           return 0;
43
44
45
      for (ll i = 0; i < m; ++i) {
46
        if (where[i] == -1) {
47
           return INF;
48
49
50
      return 1;
51 }
```

### 5.3 Ternary Search

```
double ternary_search(double 1, double r) {
      double eps = 1e-9;
      while (r - 1 > eps) {
        double m1 = 1 + (r - 1) / 3;
        double m2 = r - (r - 1) / 3;
        double f1 = f(m1);
        double f2 = f(m2):
        if (f1 < f2) {
         1 = m1;
10
        } else {
11
          r = m2:
```

```
return f(1);
.5 }
```

### Data Structures

SegTree2d() {}

### 6.1 Segment Tree 2d

```
template<typename T, typename InType = T>
    class SegTree2dNode {
    public:
      int i, j, tree_size;
      SegTree<T, InType>* seg_tree;
      SegTree2dNode<T, InType>* lc, * rc;
      SegTree2dNode() {}
      SegTree2dNode(const vector<vector<InType>>& a,
           int i, int j) : i(i), j(j) {
        tree_size = a[0].size();
        if († - i == 1) {
          lc = rc = nullptr;
          seq_tree = new SegTree<T, InType>(a[i]);
        int k = (i + j) / 2;
16
        lc = new SegTree2dNode<T, InType>(a, i, k);
        rc = new SegTree2dNode<T, InType>(a, k, j);
        seg_tree = new SegTree<T, InType>(vector<T>(
             tree_size));
        operation_2d(lc->seg_tree, rc->seg_tree);
       ~SeqTree2dNode() {
        delete lc:
        delete rc;
      void set_2d(int kx, int ky, T x) {
        if (kx < i || j <= kx) return;</pre>
        if (j - i == 1) {
          seg_tree->set(ky, x);
          return:
        1c - \sec_2 d(kx, ky, x);
        rc \rightarrow set_2d(kx, ky, x);
33
        operation_2d(lc->seg_tree, rc->seg_tree);
34
35
      T range_query_2d(int lx, int rx, int ly, int ry)
36
        if (lx <= i && j <= rx) return seg_tree->
             range_query(ly, ry);
        if (j <= lx || rx <= i) return -INF;</pre>
        return max(lc->range_query_2d(lx, rx, ly, ry),
             rc->range_query_2d(lx, rx, ly, ry));
39
      void operation_2d(SegTree<T, InType>* x, SegTree<</pre>
           T, InType>* y) {
        for (int k = 0; k < tree_size; k++) {</pre>
          seg_tree->set(k, max(x->range_query(k, k + 1)
               , y->range_query(k, k + 1)));
43
44
4.5
    template<typename T, typename InType = T>
    class SegTree2d {
48 public:
49
      SegTree2dNode<T, InType> root;
```

```
SegTree2d(const vector<vector<InType>>& mat) :
           root(mat, 0, mat.size()) {}
      void set_2d(int kx, int ky, T x) { root.set_2d(kx
           , ky, x); }
      T range_query_2d(int lx, int rx, int ly, int ry)
           { return root.range_query_2d(lx, rx, ly, ry)
54 };
```

### 6.2 Range Add Point Query

```
template<typename T, typename InType = T>
    class SegTreeNode {
    public:
      const T IDN = 0, DEF = 0;
      int i, j;
      T val:
      SegTreeNode<T, InType>* lc, * rc;
      SegTreeNode(int i, int j) : i(i), j(j) {
        if (j - i == 1) {
          lc = rc = nullptr;
          val = DEF;
          return;
        int k = (i + j) / 2;
        lc = new SegTreeNode<T, InType>(i, k);
16
        rc = new SegTreeNode<T, InType>(k, j);
        val = 0:
18
      SegTreeNode(const vector<InType>& a, int i, int j
          ) : i(i), j(j) {
        if (j - i == 1) {
          lc = rc = nullptr;
          val = (T) a[i];
          return;
        int k = (i + j) / 2;
        lc = new SegTreeNode<T, InType>(a, i, k);
        rc = new SegTreeNode<T, InType>(a, k, j);
2.8
        val = 0;
      void range_add(int 1, int r, T x) {
        if (r <= i || j <= 1) return;</pre>
        if (l <= i && j <= r) {
          val += x;
          return;
        lc->range_add(1, r, x);
        rc->range_add(1, r, x);
38
      T point_query(int k) {
        if (k < i \mid | j \le k) return IDN;
        if (j - i == 1) return val;
        return val + lc->point_query(k) + rc->
             point_query(k);
    template<typename T, typename InType = T>
    class SegTree {
    public:
      SegTreeNode<T, InType> root;
      SegTree(int n) : root(0, n) {}
      SegTree(const vector<InType>& a) : root(a, 0, a.
      void range_add(int 1, int r, T x) { root.
           range_add(l, r, x); }
```

```
52  T point_query(int k) { return root.point_query(k)
    ; }
53 };
```

### 6.3 Disjoint Set Union

```
struct DSU {
      vector<int> parent, size;
      DSU(int n) {
        parent.resize(n);
        size.resize(n);
        for (int i = 0; i < n; i++) make_set(i);</pre>
      void make set(int v) {
        parent[v] = v;
        size[v] = 1;
11
      bool is same(int a, int b) { return find set(a)
           == find_set(b); }
      int find_set(int v) { return v == parent[v] ? v :
            parent[v] = find_set(parent[v]); }
      void union_sets(int a, int b) {
15
       a = find_set(a);
16
        b = find set(b);
        if (a != b) {
          if (size[a] < size[b]) swap(a, b);</pre>
          parent[b] = a;
20
          size[a] += size[b];
23 };
```

### 6.4 Sparse Table 2d

```
const int N = 100;
    int matrix[N][N];
    int table[N][N][(int)(log2(N) + 1)][(int)(log2(N) +
          1)];
    void build_sparse_table(int n, int m) {
      for (int i = 0; i < n; i++)
         for (int j = 0; j < m; j++)
          table[i][j][0][0] = matrix[i][j];
       for (int k = 1; k \le (int)(log2(n)); k++)
         for (int i = 0; i + (1 << k) - 1 < n; i++)
           for (int j = 0; j + (1 << k) - 1 < m; j++)
             table[i][j][k][0] = min(table[i][j][k -
                 1][0], table[i + (1 << (k - 1))][j][k
                 - 1][0]);
      for (int k = 1; k \le (int)(log2(m)); k++)
13
         for (int i = 0; i < n; i++)</pre>
14
           for (int j = 0; j + (1 << k) - 1 < m; j++)
             table[i][j][0][k] = min(table[i][j][0][k -
                 1], table[i][j + (1 << (k - 1))][0][k
                 - 11):
16
      for (int k = 1; k \le (int)(log2(n)); k++)
17
         for (int 1 = 1; 1 <= (int) (log2(m)); 1++)</pre>
18
          for (int i = 0; i + (1 << k) - 1 < n; i++)
19
             for (int j = 0; j + (1 << 1) - 1 < m; <math>j++)
               table[i][j][k][l] = min(
                 min(table[i][j][k-1][l-1], table[i]
                      + (1 << (k - 1)) ] [j] [k - 1] [1 -
22
                 min(table[i][j + (1 << (1 - 1))][k -
                      1] [1 - 1], table [i + (1 << (k - 1))
```

### 6.5 Mo

```
1 void remove(idx); // TODO: remove value at idx
         from data structure
    void add(idx);
                       // TODO: add value at idx from
         data structure
    int get_answer(); // TODO: extract the current
         answer of the data structure
    int block size;
    struct Query {
      int 1, r, idx;
      bool operator<(Query other) const {</pre>
 8
        return make_pair(l / block_size, r) < make_pair</pre>
             (other.l / block_size, other.r);
9
    };
    vector<int> mo_s_algorithm(vector<Query> queries) {
      vector<int> answers(queries.size());
      sort(queries.begin(), queries.end());
      // TODO: initialize data structure
      int cur_1 = 0, cur_r = -1;
      // invariant: data structure will always reflect
           the range [cur_1, cur_r]
      for (Query q : queries) {
        while (cur_1 > q.1) {
19
          cur_1--;
          add(cur_l);
        while (cur_r < q.r) {</pre>
          cur r++;
          add(cur_r);
        while (cur_1 < q.1) {
          remove(cur_l);
          cur_1++;
        while (cur_r > q.r) {
          remove(cur_r);
          cur_r--;
3.3
        answers[q.idx] = get_answer();
      return answers;
37 }
```

### 6.6 Sparse Table

```
vector<vector<ll>> build sum(ll N, ll K, vector<ll>
          &array) {
      vector<vector<ll>> st(K + 1, vector<ll>(N + 1));
      for (ll i = 0; i < N; i++) st[0][i] = array[i];</pre>
      for (11 i = 1; i <= K; i++)
        for (11 j = 0; j + (1 << i) <= N; j++)
          st[i][j] = st[i - 1][j] + st[i - 1][j + (1 <<
                (i - 1));
      return st;
    11 sum_query(11 L, 11 R, 11 K, vector<vector<11>> &
         st) {
      11 \text{ sum} = 0;
      for (11 i = K; i >= 0; i--) {
        if ((1 << i) <= R - L + 1) {
          sum += st[i][L];
          L += 1 << i;
      return sum;
    vector<vector<ll>> build min(ll N, ll K, vector<ll>>
      vector<vector<ll>>> st(K + 1, vector<ll>(N + 1));
      for (ll i = 0; i < N; i++) st[0][i] = array[i];</pre>
      for (ll i = 1; i <= K; i++)</pre>
        for (11 j = 0; j + (1 << i) <= N; j++)
          st[i][j] = min(st[i-1][j], st[i-1][j+(1
                << (i - 1))];
      return st;
29
    11 min_query(11 L, 11 R, vector<vector<11>>> &st) {
      11 i = log2 floor(R - L + 1);
32
      return min(st[i][L], st[i][R - (1 << i) + 1]);</pre>
33 }
```

### 6.7 Binary Trie

```
struct Node { struct Node* parent, child[2]; };
   struct BinaryTrie {
     Node* root:
     BinarvTrie() {
       root = new Node();
       root->parent = NULL;
       root->child[0] = NULL;
       root->child[1] = NULL;
9
     void insert_node(int x) {
       Node* cur = root;
       for (int place = 29; place >= 0; place--) {
         int bit = x >> place & 1;
         if (cur->child[bit] != NULL) cur = cur->child
              [bit];
         else {
           cur->child[bit] = new Node();
           cur->child[bit]->parent = cur;
           cur = cur->child[bit];
           cur->child[0] = NULL;
           cur->child[1] = NULL;
     void remove node(int x) {
       Node* cur = root;
       for (int place = 29; place >= 0; place--) {
         int bit = x >> place & 1;
```

```
28
          if (cur->child[bit] == NULL) return;
29
          cur = cur->child[bit];
31
        while (cur->parent != NULL && cur->child[0] ==
             NULL && cur->child[1] == NULL) {
          Node* temp = cur;
          cur = cur->parent;
34
          if (temp == cur->child[0]) cur->child[0] =
          else cur->child[1] = NULL;
36
          delete temp;
37
38
39
      int get_min_xor(int x) {
40
        Node* cur = root;
41
        int minXor = 0;
42
        for (int place = 29; place >= 0; place--) {
          int bit = x >> place & 1;
43
44
          if (cur->child[bit] != NULL) cur = cur->child
               [bit1:
45
          else (
46
            minXor ^= 1 << place;
47
            cur = cur->child[1 ^ bit];
48
49
50
        return minXor;
51
52 };
```

### 6.8 Segment Tree

```
template<typename T, typename InType = T>
    class SegTreeNode {
    public:
      const T IDN = 0, DEF = 0;
      int i, j;
      T val;
      SegTreeNode<T, InType>* lc, * rc;
      SegTreeNode(int i, int j) : i(i), j(j) {
       if (i - i == 1) {
         lc = rc = nullptr;
11
          val = DEF:
          return;
13
14
        int k = (i + j) / 2;
15
        lc = new SegTreeNode<T, InType>(i, k);
         rc = new SegTreeNode<T, InType>(k, j);
17
        val = op(lc->val, rc->val);
18
19
      SegTreeNode(const vector<InType>& a, int i, int j
          ) : i(i), j(j) {
         if (j - i == 1) {
          lc = rc = nullptr;
          val = (T) a[i];
23
          return;
24
         int k = (i + j) / 2;
        lc = new SegTreeNode<T, InType>(a, i, k);
27
         rc = new SegTreeNode<T, InType>(a, k, j);
28
        val = op(lc->val, rc->val);
29
30
31
      void set(int k, T x) {
        if (k < i || j <= k) return;</pre>
32
        if (j - i == 1) {
33
          val = x;
34
          return;
35
```

```
1c->set(k, x);
         rc \rightarrow set(k, x);
38
        val = op(lc->val, rc->val);
39
40
      T range_query(int 1, int r) {
        if (1 <= i && j <= r) return val;</pre>
41
         if (j <= 1 || r <= i) return IDN;</pre>
        return op(lc->range_query(1, r), rc->
             range_query(1, r));
45
      T \circ p(T \times, T y) \{ \}
46
   };
47 template<typename T, typename InType = T>
48 class SegTree {
49
    public:
50
      SegTreeNode<T, InType> root;
      SegTree(int n) : root(0, n) {}
      SegTree(const vector<InType>& a) : root(a, 0, a.
           size()) {}
      void set(int k, T x) { root.set(k, x); }
      T range_query(int 1, int r) { return root.
            range_query(1, r); }
55 };
```

### 6.9 Sqrt Decomposition

```
int n;
   vector<int> a (n);
    int len = (int) sqrt (n + .0) + 1; // size of the
         block and the number of blocks
    vector<int> b (len);
    for (int i = 0; i<n; ++i) b[i / len] += a[i];</pre>
    for (;;) {
      int 1, r:
      // read input data for the next query
      int sum = 0;
      for (int i = 1; i <= r; )</pre>
       if (i % len == 0 && i + len - 1 <= r) {</pre>
          // if the whole block starting at i belongs
               to [1, r]
          sum += b[i / len];
14
          i += len:
        } else {
          sum += a[i];
          ++i;
18
19
      // or
20
      /*
      int sum = 0;
      int c_1 = 1 / len, c_r = r / len;
      if (c_1 == c_r)
          for (int i=1; i<=r; ++i)
              sum += a[i];
26
      else {
          for (int i=1, end=(c_1+1)*len-1; i<=end; ++i)
              sum += a[i];
          for (int i=c_1+1; i<=c_r-1; ++i)
              sum += b[i]:
          for (int i=c_r*len; i<=r; ++i)
             sum += a[i];
34
      */
35 }
```

### 6.10 Minimum Queue

```
1 11 get_minimum(stack<pair<11, 11>> &s1, stack<pair<</pre>
         11, 11>> &s2) {
      if (s1.empty() || s2.empty()) {
        return s1.empty() ? s2.top().second : s1.top().
             second:
      } else {
        return min(s1.top().second, s2.top().second);
 7
   void add element(ll new element, stack<pair<11, ll</pre>
      11 minimum = s1.empty() ? new_element : min(
           new_element, s1.top().second);
      s1.push({new_element, minimum});
    11 remove_element(stack<pair<11, 11>> &s1, stack
         pair<11, 11>> &s2) {
      if (s2.empty()) {
        while (!s1.empty()) {
          11 element = s1.top().first;
          s1.pop();
          11 minimum = s2.empty() ? element : min(
               element, s2.top().second);
          s2.push({element, minimum});
21
      11 removed_element = s2.top().first;
      s2.pop();
      return removed_element;
24
```

### 6.11 Range Add Range Query

```
template<typename T, typename InType = T>
    class SegTreeNode {
    public:
      const T IDN = 0, DEF = 0;
      int i, j;
      T val, to_add = 0;
      SegTreeNode<T, InType>* lc, * rc;
      SegTreeNode(int i, int j) : i(i), j(j) {
       if (i - i == 1) {
         lc = rc = nullptr;
          val = DEF:
          return;
        int k = (i + j) / 2;
        lc = new SegTreeNode<T, InType>(i, k);
        rc = new SegTreeNode<T, InType>(k, j);
        val = operation(lc->val, rc->val);
18
      SegTreeNode(const vector<InType>& a, int i, int j
          ) : i(i), j(j) {
        if (j - i == 1) {
          lc = rc = nullptr;
          val = (T) a[i];
          return;
        int k = (i + j) / 2;
        lc = new SegTreeNode<T, InType>(a, i, k);
        rc = new SegTreeNode<T, InType>(a, k, j);
2.8
        val = operation(lc->val, rc->val);
29
      void propagate() {
        if (to add == 0) return;
        val += to_add;
        if (j - i > 1) {
```

```
lc->to_add += to_add;
35
          rc->to add += to add;
36
37
        to\_add = 0;
38
39
      void range_add(int 1, int r, T delta) {
40
        propagate();
41
        if (r <= i | | j <= 1) return;</pre>
42
        if (1 <= i && j <= r) {
43
          to add += delta;
44
          propagate();
45
        } else {
46
          lc->range_add(l, r, delta);
47
           rc->range_add(1, r, delta);
48
          val = operation(lc->val, rc->val);
49
50
51
      T range_query(int 1, int r) {
52
        propagate();
53
         if (1 <= i && j <= r) return val;</pre>
         if (j <= 1 || r <= i) return IDN;</pre>
         return operation(lc->range_query(l, r), rc->
             range query(1, r));
56
57
      T operation(T x, T y) {}
58 };
59
    template<typename T, typename InType = T>
    class SegTree {
61
    public:
62
      SegTreeNode<T, InType> root;
63
      SegTree(int n) : root(0, n) {}
      SegTree(const vector<InType>& a) : root(a, 0, a.
           size()) {}
65
       void range add(int 1, int r, T delta) { root.
           range_add(l, r, delta); }
66
      T range_query(int 1, int r) { return root.
           range_query(1, r); }
67 };
```

### 7 Graph Theory

### 7.1 Bridge

```
vector<vector<int>> adj;
    vector<bool> visited;
 4 vector<int> tin, low;
 5 int timer;
    void dfs(int v, int p = -1) {
      visited[v] = true;
      tin[v] = low[v] = timer++;
      for (int to : adj[v]) {
       if (to == p) continue;
11
        if (visited[to]) {
          low[v] = min(low[v], tin[to]);
        } else {
14
          dfs(to, v);
15
          low[v] = min(low[v], low[to]);
          if (low[to] > tin[v]) IS_BRIDGE(v, to);
17
18
    }
19 }
20 void find_bridges() {
21
      timer = 0;
      visited.assign(n, false);
      tin.assign(n, -1);
```

```
low.assign(n, -1);
      for (int i = 0; i < n; ++i) {
        if (!visited[i]) dfs(i);
28 }
 7.2 Dijkstra
    const int INF = 1000000000;
   vector<vector<pair<int, int>>> adj;
 3 void dijkstra(int s, vector<int> & d, vector<int> &
          p) {
      int n = adj.size();
      d.assign(n, INF);
 6
      p.assign(n, -1);
      d[s] = 0;
      using pii = pair<int, int>;
      priority_queue<pii, vector<pii>, greater<pii>> g;
      q.push({0, s});
      while (!q.empty()) {
        int v = q.top().second, d_v = q.top().first;
        q.pop();
        if (d_v != d[v]) continue;
        for (auto edge : adj[v]) {
          int to = edge.first, len = edge.second;
          if (d[v] + len < d[to]) {</pre>
              d[to] = d[v] + len;
              p[to] = v;
              q.push({d[to], to});
      }
24 }
```

### 7.3 Zero One Bfs

```
vector<int> d(n, INF);
2 d[s] = 0;
3 deque<int> q;
4 q.push_front(s);
5 while (!q.empty()) {
    int v = q.front();
     q.pop_front();
      for (auto edge : adj[v]) {
      int u = edge.first, w = edge.second;
10
       if (d[v] + w < d[u]) {
         d[u] = d[v] + w;
12
         if (w == 1) q.push_back(u);
         else q.push_front(u);
14
     }
16 }
```

### 7.4 Hungarian

```
1  vector<int> u (n+1), v (m+1), p (m+1), way (m+1);
2  for (int i=1; i<=n; ++i) {
3    p[0] = i;
4    int j0 = 0;
5    vector<int> minv (m+1, INF);
6    vector<bod> used (m+1, false);
7    do (
```

```
used[j0] = true;
        int i0 = p[j0], delta = INF, j1;
        for (int j=1; j<=m; ++j)</pre>
          if (!used[j]) {
12
            int cur = A[i0][j]-u[i0]-v[j];
13
             if (cur < minv[j]) minv[j] = cur, way[j] =</pre>
             if (minv[j] < delta) delta = minv[j], j1 =</pre>
        for (int j=0; j<=m; ++j)</pre>
          if (used[j]) u[p[j]] += delta, v[j] -= delta
           else minv[j] -= delta;
         i0 = i1;
      } while (p[j0] != 0);
        int j1 = way[j0];
        p[j0] = p[j1];
         j0 = j1;
      } while (†0);
    vector<int> ans (n+1);
    for (int j=1; j<=m; ++j)</pre>
      ans[p[j]] = j;
30 int cost = -v[0];
```

#### 7.5 Ford Fulkerson

```
bool bfs(ll n, vector<vector<ll>>> &r_graph, ll s,
         11 t, vector<11> &parent) {
      vector<bool> visited(n, false);
      queue<11> q;
      q.push(s);
      visited[s] = true;
      parent[s] = -1;
      while (!q.empty()) {
        11 u = q.front();
        q.pop();
        for (11 \ v = 0; \ v < n; \ v++)  {
          if (!visited[v] && r_graph[u][v] > 0) {
            if (v == t) {
              parent[v] = u;
              return true;
            q.push(v);
            parent[v] = u;
            visited[v] = true;
        }
      return false:
    11 ford_fulkerson(ll n, vector<vector<ll>>> graph,
         11 s, 11 t) {
      11 u, v;
      vector<vector<ll>> r_graph;
      for (u = 0; u < n; u++)
        for (v = 0; v < n; v++)
          r_graph[u][v] = graph[u][v];
      vector<11> parent;
      11 \text{ max flow} = 0;
      while (bfs(n, r_graph, s, t, parent)) {
        11 path_flow = INF;
34
        for (v = t; v != s; v = parent[v]) {
          u = parent[v];
          path_flow = min(path_flow, r_graph[u][v]);
```

```
37     }
38     for (v = t; v != s; v = parent[v]) {
39         u = parent[v];
40         r_graph[u][v] -= path_flow;
41         r_graph[v][u] += path_flow;
42     }
43     max_flow += path_flow;
44     }
45     return max_flow;
46  }
```

#### 7.6 Prim

```
const int INF = 1000000000;
    struct Edge {
       int w = INF, to = -1;
      bool operator<(Edge const& other) const {</pre>
         return make_pair(w, to) < make_pair(other.w,</pre>
             other to):
    };
8
    int n;
    vector<vector<Edge>> adj;
10 void prim() {
      int total_weight = 0;
      vector<Edge> min_e(n);
      min_e[0].w = 0;
      set < Edge > q;
      q.insert({0, 0});
      vector<bool> selected(n, false);
17
      for (int i = 0; i < n; ++i) {</pre>
18
        if (q.empty()) {
19
           cout << "No MST!" << endl;</pre>
20
          exit(0);
         int v = q.begin()->to;
23
         selected[v] = true;
24
         total_weight += q.begin()->w;
25
         q.erase(q.begin());
         if (min_e[v].to != -1) cout << v << " " <<</pre>
              min_e[v].to << endl;</pre>
         for (Edge e : adj[v]) {
28
          if (!selected[e.to] && e.w < min_e[e.to].w) {</pre>
29
             q.erase({min_e[e.to].w, e.to});
             min_e[e.to] = \{e.w, v\};
             q.insert({e.w, e.to});
34
      cout << total_weight << endl;</pre>
```

#### 7.7 Centroid Decomposition

```
return subtree_size[node];
   int get_centroid(int node, int tree_size, int
        parent = -1) {
            for (int child : adj[node]) {
14
                    if (child == parent || is_removed[
                         child]) continue;
                    if (subtree_size[child] * 2 >
                         tree_size) return get_centroid
                         (child, tree size, node);
            return node;
18
19
   void build_centroid_decomp(int node = 0) {
            int centroid = get_centroid(node,
20
                 get_subtree_size(node));
            // do something
            is_removed[centroid] = true;
            for (int child : adj[centroid]) {
                    if (is_removed[child]) continue;
                    build_centroid_decomp(child);
```

#### 7.8 Kahn

```
void kahn(vector<vector<ll>> &adj) {
      11 n = adj.size();
      vector<11> in_degree(n, 0);
      for (11 u = 0; u < n; u++)
       for (ll v: adj[u]) in_degree[v]++;
      queue<11> q;
      for (11 i = 0; i < n; i++)
       if (in_degree[i] == 0)
          q.push(i);
      11 \text{ cnt} = 0;
      vector<11> top_order;
      while (!q.empty()) {
       11 u = q.front();
        q.pop();
15
        top_order.push_back(u);
16
        for (11 v : adj[u])
          if (--in_degree[v] == 0) g.push(v);
18
        cnt++;
      if (cnt != n) {
        cout << -1 << '\n';
        return:
      // print top_order
```

### 7.9 Dinics

```
vector<int> level, ptr;
      queue<int> q;
      Dinic(int n, int s, int t) : n(n), s(s), t(t) {
        adj.resize(n);
        level.resize(n);
        ptr.resize(n);
18
      void add_edge(int v, int u, ll cap) {
        edges.emplace_back(v, u, cap);
        edges.emplace_back(u, v, 0);
        adj[v].push_back(m);
        adj[u].push_back(m + 1);
        m += 2:
24
      bool bfs() {
        while (!q.empty()) {
          int v = q.front();
          q.pop();
          for (int id : adj[v]) {
            if (edges[id].cap - edges[id].flow < 1)</pre>
                 continue;
            if (level[edges[id].u] != -1) continue;
            level[edges[id].u] = level[v] + 1;
            g.push(edges[id].u);
34
        return level[t] != -1;
38
      11 dfs(int v, 11 pushed) {
        if (pushed == 0) return 0;
40
        if (v == t) return pushed;
        for (int& cid = ptr[v]; cid < (int)adj[v].size</pre>
             (); cid++) {
          int id = adj[v][cid], u = edges[id].u;
43
          if (level[v] + 1 != level[u] || edges[id].cap
                - edges[id].flow < 1) continue;</pre>
44
          11 tr = dfs(u, min(pushed, edges[id].cap -
               edges[id].flow));
          if (tr == 0) continue;
          edges[id].flow += tr;
          edges[id ^ 1].flow -= tr;
          return tr;
        return 0;
      11 flow() {
        11 f = 0;
        while (true) {
          fill(level.begin(), level.end(), -1);
          level[s] = 0;
          q.push(s);
58
          if (!bfs()) break;
59
          fill(ptr.begin(), ptr.end(), 0);
60
          while (11 pushed = dfs(s, flow_inf)) f +=
               pushed;
62
        return f;
63
64 };
```

#### 7.10 Floyd Warshall

```
1  void floyd_warshall(vector<vector<ll>>> &dis, ll n)
{
2   for (ll k = 0; k < n; k++)
3    for (ll i = 0; i < n; i++)
4   for (ll j = 0; j < n; j++)</pre>
```

### 7.11 Kosaraju

```
void topo_sort(int u, vector<vector<int>>& adj,
         vector<bool>& vis, stack<int>& stk) {
      vis[u] = true;
      for (int v : adj[u]) {
        if (!vis[v]) {
          topo_sort(v, adj, vis, stk);
      stk.push(u);
 9
   vector<vector<int>>> transpose(int n, vector<vector<</pre>
         int>>& adj) {
      vector<vector<int>> adj_t(n);
13
      for (int u = 0; u < n; u++) {
14
        for (int v : adj[u]) {
15
          adj_t[v].push_back(u);
16
17
18
      return adj_t;
19
20
    void get_scc(int u, vector<vector<int>>& adj_t,
         vector<bool>& vis, vector<int>& scc) {
      vis[u] = true;
23
      scc.push_back(u);
      for (int v : adj_t[u]) {
24
        if (!vis[v]) {
          get_scc(v, adj_t, vis, scc);
27
29
    void kosaraju(int n, vector<vector<int>>& adj,
         vector<vector<int>>& sccs) {
      vector<bool> vis(n, false);
      stack<int> stk;
      for (int u = 0; u < n; u++) {
       if (!vis[u]) {
36
          topo_sort(u, adj, vis, stk);
37
39
      vector<vector<int>> adj_t = transpose(n, adj);
40
      for (int u = 0; u < n; u++) {
41
       vis[u] = false;
42
43
      while (!stk.empty()) {
44
        int u = stk.top();
45
        stk.pop();
46
        if (!vis[u]) {
47
          vector<int> scc;
48
          get_scc(u, adj_t, vis, scc);
49
          sccs.push_back(scc);
50
```

```
51 }
52 }
```

### 7.12 Maximum Bipartite Matching

```
bool bpm(ll n, ll m, vector<vector<bool>> &bpGraph,
          11 u, vector<bool> &seen, vector<11> &matchR)
      for (11 v = 0; v < m; v++) {
        if (bpGraph[u][v] && !seen[v]) {
 4
          seen[v] = true;
          if (matchR[v] < 0 \mid \mid bpm(n, m, bpGraph,
               matchR[v], seen, matchR)) {
            matchR[v] = u;
            return true;
      return false:
    11 maxBPM(11 n, 11 m, vector<vector<bool>> &bpGraph
      vector<11> matchR(m, -1);
      11 \text{ result} = 0;
      for (11 u = 0; u < n; u++) {
        vector<bool> seen(m, false);
        if (bpm(n, m, bpGraph, u, seen, matchR)) {
          result++;
20
      return result;
```

#### 7.13 Kruskals

```
struct Edge {
      int u, v, weight;
      bool operator<(Edge const& other) {</pre>
        return weight < other.weight;</pre>
 6
   };
   int n;
   vector<Edge> edges;
 9 int cost = 0;
10 vector<Edge> result;
11 DSU dsu = DSU(n);
12 sort(edges.begin(), edges.end());
13 for (Edge e : edges) {
14 if (dsu.find_set(e.u) != dsu.find_set(e.v)) {
      cost += e.weight;
16
       result.push_back(e);
        dsu.union_sets(e.u, e.v);
19 }
```

#### 7.14 Is Cyclic

### 7.15 Find Cycle

```
1 bool dfs(ll v) {
      color[v] = 1;
      for (ll u : adj[v]) {
        if (color[u] == 0) {
          parent[u] = v;
          if (dfs(u)) {
            return true;
        } else if (color[u] == 1) {
          cycle_end = v;
          cycle_start = u;
          return true;
      color[v] = 2;
16
      return false;
    void find_cycle() {
      color.assign(n, 0);
      parent.assign(n, -1);
      cycle_start = -1;
      for (11 v = 0; v < n; v++) {
        if (color[v] == 0 && dfs(v)) {
          break;
      if (cycle_start == -1) {
        cout << "Acyclic" << endl;</pre>
      } else {
        vector<ll> cycle;
        cycle.push_back(cycle_start);
        for (11 v = cycle_end; v != cycle_start; v =
             parent[v]) {
          cycle.push_back(v);
        cycle.push_back(cycle_start);
        reverse(cycle.begin(), cycle.end());
        cout << "Cycle found: ";</pre>
        for (11 v : cycle) {
39
          cout << v << ' ';
        cout << '\n';
43 }
```

### 7.16 Topological Sort

```
1 void dfs(ll v) {
```

```
visited[v] = true;
      for (ll u : adj[v]) {
        if (!visited[u]) {
          dfs(u);
      ans.push back(v);
9
10 void topological_sort() {
11
      visited.assign(n, false);
      ans.clear();
13
      for (11 i = 0; i < n; ++i) {
14
       if (!visited[i]) {
15
         dfs(i);
16
17
18
     reverse(ans.begin(), ans.end());
19 }
```

#### 7.17 Min Cost Flow

```
struct Edge {
      int from, to, capacity, cost;
    vector<vector<int>> adj, cost, capacity;
    const int INF = 1e9;
    void shortest_paths(int n, int v0, vector<int>& d,
         vector<int>& p) {
      d.assign(n, INF);
      d[v0] = 0;
      vector<bool> ing(n, false);
      queue<int> q;
      q.push(v0);
      p.assign(n, -1);
13
      while (!q.empty()) {
14
        int u = q.front();
15
        q.pop();
16
        inq[u] = false;
17
        for (int v : adj[u]) {
18
          if (capacity[u][v] > 0 && d[v] > d[u] + cost[
               u][v]) {
19
            d[v] = d[u] + cost[u][v];
20
            p[v] = u;
21
            if (!ing[v]) {
22
              inq[v] = true;
23
              q.push(v);
24
25
          }
26
        }
27
28
    int min_cost_flow(int N, vector<Edge> edges, int K,
          int s, int t) {
      adj.assign(N, vector<int>());
      cost.assign(N, vector<int>(N, 0));
      capacity.assign(N, vector<int>(N, 0));
      for (Edge e : edges) {
34
       adj[e.from].push_back(e.to);
35
        adj[e.to].push_back(e.from);
        cost[e.from][e.to] = e.cost;
37
        cost[e.to][e.from] = -e.cost;
38
        capacity[e.from][e.to] = e.capacity;
39
40
      int flow = 0;
41
      int cost = 0;
      vector<int> d, p;
      while (flow < K) {</pre>
```

```
shortest_paths(N, s, d, p);
        if (d[t] == INF) break;
        int f = K - flow, cur = t;
        while (cur != s) {
         f = min(f, capacity[p[cur]][cur]);
49
          cur = p[cur];
50
        flow += f:
        cost += f * d[t];
        cur = t;
        while (cur != s) {
          capacity[p[cur]][cur] -= f;
          capacity[cur][p[cur]] += f;
          cur = p[cur];
58
59
60
     if (flow < K) return -1;</pre>
      else return cost;
```

### 7.18 Kuhn

```
1 int n, k;
   vector<vector<int>> q;
   vector<int> mt;
    vector<bool> used;
   bool try_kuhn(int v) {
     if (used[v]) return false;
      used[v] = true;
      for (int to : g[v]) {
       if (mt[to] == -1 || try_kuhn(mt[to])) {
          mt[to] = v;
11
          return true:
12
      return false;
15
   int main() {
      mt.assign(k, -1);
        vector<bool> used1(n, false);
19
        for (int v = 0; v < n; ++v) {
          for (int to : g[v]) {
            if (mt[to] == -1) {
              mt[to] = v;
              used1[v] = true;
              break:
        for (int v = 0; v < n; ++v) {
          if (used1[v]) continue;
          used.assign(n, false);
          try_kuhn(v);
32
3.3
        for (int i = 0; i < k; ++i)
          if (mt[i] != −1)
            printf("%d %d\n", mt[i] + 1, i + 1);
```

#### 7.19 Articulation Point

```
11 children = 0;
      visited[u] = true;
      disc[u] = low[u] = ++time;
      for (auto v : adj[u]) {
       if (!visited[v]) {
          children++;
          APUtil(adj, v, visited, disc, low, time, u,
          low[u] = min(low[u], low[v]);
          if (parent != -1 && low[v] >= disc[u]) {
            isAP[u] = true;
       } else if (v != parent) {
         low[u] = min(low[u], disc[v]);
      if (parent == -1 && children > 1) {
       isAP[u] = true;
20
21
    void AP(vector<vector<11>> &adj, 11 n) {
      vector<ll> disc(n), low(n);
      vector<bool> visited(n), isAP(n);
      11 time = 0, par = -1;
      for (11 u = 0; u < n; u++) {
       if (!visited[u]) {
28
          APUtil(adj, u, visited, disc, low, time, par,
               isAP):
      for (11 u = 0; u < n; u++) {
       if (isAP[u]) {
          cout << u << " ";
34
```

#### 7.20 Hierholzer

```
void print_circuit (vector<vector<ll>>> &adj) {
      map<11, 11> edge_count;
      for (11 i = 0; i < adj.size(); i++) {</pre>
        edge_count[i] = adj[i].size();
5
      if (!adj.size()) {
        return;
      stack<ll> curr_path;
      vector<ll> circuit;
      curr_path.push(0);
      11 curr_v = 0;
      while (!curr_path.empty()) {
        if (edge_count[curr_v]) {
          curr path.push(curr v);
          11 next_v = adj[curr_v].back();
          edge_count[curr_v]--;
          adj[curr_v].pop_back();
          curr_v = next_v;
        } else {
          circuit.push_back(curr_v);
          curr_v = curr_path.top();
          curr_path.pop();
26
      for (ll i = circuit.size() - 1; i >= 0; i--) {
27
        cout << circuit[i] << ' ';
```

### 7.21 Lowest Common Ancestor

```
struct LCA {
       vector<ll> height, euler, first, segtree;
       vector<bool> visited;
      LCA(vector<vector<11>>> &adj, 11 root = 0) {
        n = adi.size():
        height.resize(n);
         first.resize(n);
         euler.reserve(n * 2);
         visited.assign(n, false);
         dfs(adj, root);
12
         11 m = euler.size();
13
         segtree.resize(m * 4);
14
         build(1, 0, m - 1);
15
      void dfs(vector<vector<ll>>> &adj, ll node, ll h = |22
16
         visited[node] = true;
         height[node] = h;
19
         first[node] = euler.size();
20
         euler.push back(node);
         for (auto to : adj[node]) {
          if (!visited[to]) {
23
             dfs(adj, to, h + 1);
24
             euler.push_back(node);
25
26
27
28
29
      void build(ll node, ll b, ll e) {
         if (b == e) segtree[node] = euler[b];
          11 \text{ mid} = (b + e) / 2;
32
          build(node << 1, b, mid);</pre>
          build(node << 1 | 1, mid + 1, e);</pre>
34
           11 1 = segtree[node << 1], r = segtree[node</pre>
                << 1 | 11;
           segtree[node] = (height[1] < height[r]) ? 1 :</pre>
36
38
      11 query(11 node, 11 b, 11 e, 11 L, 11 R) {
39
        if (b > R | | e < L) return -1;</pre>
40
         if (b >= L && e <= R) return segtree[node];</pre>
41
         11 \text{ mid} = (b + e) >> 1;
42
         11 left = query(node << 1, b, mid, L, R);</pre>
         11 right = query(node << 1 | 1, mid + 1, e, L,</pre>
         if (left == -1) return right;
         if (right == -1) return left;
45
         return height[left] < height[right] ? left :</pre>
47
48
      ll lca(ll u, ll v) {
49
         11 left = first[u], right = first[v];
50
         if (left > right) swap(left, right);
51
         return query(1, 0, euler.size() - 1, left,
              right);
52
```

```
struct Edge {
     int a, b, cost;
    };
    int n, m, v;
    vector<Edge> edges;
 6 const int INF = 1000000000;
    void solve() {
     vector<int> d(n, INF);
      d[v] = 0;
      vector<int> p(n, -1);
      for (int i = 0; i < n; ++i) {</pre>
       \mathbf{x} = -1;
       for (Edge e : edges)
         if (d[e.a] < INF)</pre>
            if (d[e.b] > d[e.a] + e.cost) {
              d[e.b] = max(-INF, d[e.a] + e.cost);
18
              p[e.b] = e.a;
              x = e.b;
      if (x == -1) cout << "No negative cycle from " <<</pre>
      else {
        int y = x;
        for (int i = 0; i < n; ++i) y = p[y];</pre>
        vector<int> path;
        for (int cur = y;; cur = p[cur]) {
         path.push_back(cur);
          if (cur == y && path.size() > 1) break;
        reverse(path.begin(), path.end());
32
        cout << "Negative cycle: ";</pre>
        for (int u : path) cout << u << ' ';</pre>
34
 7.23 Edmonds Karp
 1 int n;
   vector<vector<int>> capacity;
    vector<vector<int>> adi:
```

```
int bfs(int s, int t, vector<int>& parent) {
      fill(parent.begin(), parent.end(), -1);
      parent[s] = -2;
      queue<pair<int, int>> q;
      q.push({s, INF});
      while (!q.empty()) {
       int cur = q.front().first, flow = q.front().
             second:
        for (int next : adj[cur]) {
          if (parent[next] == -1 && capacity[cur][next
               1) {
            parent[next] = cur;
            int new_flow = min(flow, capacity[cur][next
            if (next == t) return new_flow;
            q.push({next, new_flow});
19
       }
2.0
      return 0:
23 int maxflow(int s, int t) {
      int flow = 0;
```

vector<int> parent(n);

```
int new_flow;
      while (new flow = bfs(s, t, parent)) {
       flow += new_flow;
        int cur = t;
        while (cur != s) {
          int prev = parent[cur];
          capacity[prev][cur] -= new_flow;
          capacity[cur][prev] += new_flow;
          cur = prev;
36
      return flow;
38 }
```

### 7.24 Is Bipartite

```
1 bool is_bipartite(vector<11> &col, vector<vector<11</pre>
         >> &adj, 11 n) {
      queue<pair<11, 11>> q;
      for (11 i = 0; i < n; i++) {
        if (col[i] == -1) {
          q.push({i, 0});
          col[i] = 0;
          while (!q.empty()) {
            pair<11, 11> p = q.front();
            q.pop();
            11 v = p.first, c = p.second;
            for (11 j : adj[v]) {
              if (col[j] == c) {
                return false;
              if (col[j] == -1) {
                col[j] = (c ? 0 : 1);
                q.push({j, col[j]});
21
      return true;
24
```

### 7.25 Fast Second Mst

```
struct edge {
       int s, e, w, id;
        bool operator<(const struct edge& other) {</pre>
             return w < other.w; }</pre>
   typedef struct edge Edge;
   const int N = 2e5 + 5;
    long long res = 0, ans = 1e18;
   int n, m, a, b, w, id, 1 = 21;
   vector<Edge> edges;
   vector<int> h(N, 0), parent(N, -1), size(N, 0),
         present (N, 0);
vector<vector<pair<int, int>>> adj(N), dp(N, vector
         <pair<int, int>>(1));
12 vector<vector<int>> up(N, vector<int>(1, -1));
   pair<int, int> combine(pair<int, int> a, pair<int,</pre>
      vector<int> v = {a.first, a.second, b.first, b.
           second);
      int topTwo = -3, topOne = -2;
      for (int c : v) {
```

```
17
        if (c > topOne) {
18
          topTwo = topOne;
19
          topOne = c;
20
         } else if (c > topTwo && c < topOne) topTwo = c</pre>
21
      return {topOne, topTwo};
23 }
24 void dfs(int u, int par, int d) {
25
      h[u] = 1 + h[par];
26
      up[u][0] = par;
27
      dp[u][0] = \{d, -1\};
28
      for (auto v : adj[u]) {
29
        if (v.first != par) dfs(v.first, u, v.second);
31 }
32 pair<int, int> lca(int u, int v) {
33
      pair<int, int> ans = \{-2, -3\};
      if (h[u] < h[v]) swap(u, v);</pre>
34
35
      for (int i = 1 - 1; i >= 0; i--) {
36
        if (h[u] - h[v] >= (1 << i)) {
          ans = combine(ans, dp[u][i]);
38
          u = up[u][i];
39
40
41
      if (u == v) return ans;
42
      for (int i = 1 - 1; i >= 0; i--) {
        if (up[u][i] != -1 && up[v][i] != -1 && up[u][i
43
             ] != up[v][i]) {
44
          ans = combine(ans, combine(dp[u][i], dp[v][i
               1));
45
          u = up[u][i];
46
          v = up[v][i];
47
48
49
      ans = combine(ans, combine(dp[u][0], dp[v][0]));
      return ans;
51
52
53 int main(void) {
54
      cin >> n >> m;
      for (int i = 1; i <= n; i++) {</pre>
55
       parent[i] = i;
57
        size[i] = 1;
58
59
      for (int i = 1; i <= m; i++) {</pre>
60
        cin >> a >> b >> w; // 1-indexed
61
        edges.push_back({a, b, w, i - 1});
62
63
      sort(edges.begin(), edges.end());
64
      for (int i = 0; i <= m - 1; i++) {
65
       a = edges[i].s;
66
        b = edges[i].e;
67
        w = edges[i].w;
68
        id = edges[i].id;
69
        if (unite_set(a, b)) {
70
          adj[a].emplace_back(b, w);
71
          adj[b].emplace_back(a, w);
72
          present[id] = 1;
73
          res += w;
74
75
76
      dfs(1, 0, 0);
      for (int i = 1; i \le 1 - 1; i++) {
78
        for (int j = 1; j \le n; ++j) {
79
          if (up[j][i - 1] != -1) {
80
            int v = up[j][i - 1];
81
             up[j][i] = up[v][i - 1];
82
             dp[j][i] = combine(dp[j][i-1], dp[v][i-
```

```
1]);
      for (int i = 0; i \le m - 1; i++) {
        id = edges[i].id;
        w = edges[i].w;
        if (!present[id]) {
          auto rem = lca(edges[i].s, edges[i].e);
          if (rem.first != w) {
            if (ans > res + w - rem.first) ans = res +
                 w - rem.first;
          } else if (rem.second != -1) {
            if (ans > res + w - rem.second) ans = res +
                  w - rem.second;
96
       }
      }
98
     cout << ans << "\n";
      return 0;
```

### References

### 8.1 Stack

```
// declaration
  stack<T> stk;
3 // functions
4 stk.empty();
5 stk.size();
6 stk.top();
7 stk.push(x);
8 stk.pop();
```

### 8.2 Queue

```
1 // declaration
  queue<T> q;
3
  // functions
4 q.empty();
  q.size();
  q.front();
   q.back();
  q.push(x);
  q.pop();
```

#### 8.3 Set

```
1 // declaration
   set<T> st;
   set<T, greater<T>> st;
   // custom comparator
5 class Compare {
6 public:
     bool operator() (T a, T b) {
       if (cond) return true; // do not swap
       return false;
10
   set<T, Compare> st;
```

```
14 st.insert(x);
   st.erase(x);
   st.size();
   st.empty();
18 st.begin();
19 st.end();
20 st.clear();
21 st.count(x);
22 st.find(x);
23 st.lower_bound(x); // first element not less than x
24 st.upper_bound(x); // first element greater than x
```

# 8.4 Syntax

```
1 // multiset
  ms.insert(x)
  ms.begin()
4 ms.end()
5 ms.clear()
6 ms.erase(x)
7 ms.size()
8 ms.empty()
   // map
10 begin()
11 end()
12 size()
13 max_size()
14 empty()
15 pair insert(keyvalue, mapvalue)
16 erase(iterator position)
  erase(const q)
18 clear()
    // ordered_set
   find_by_order(k)
    order_of_key(k)
    #include <ext/pb_ds/assoc_container.hpp>
    #include <ext/pb_ds/tree_policy.hpp>
    using namespace __gnu_pbds;
    #define ordered_set
       tree<int, null_type, less<int>, rb_tree_tag, \
28
            tree_order_statistics_node_update>
   // tuple
30 get<i>(tuple)
31 make_tuple(a1, a2, ...)
32 tuple_size<decltype(tuple)>::value
   tuple1.swap(tuple2)
   tie(a1, a2, ...) = tuple
   tuple_cat(tuple1, tuple2)
   // iterator
    for (auto it = s.begin(); it != s.end(); it++) cout
         << *it << "\n";
38 begin()
   end()
   advance(ptr, k)
   next(ptr, k)
    prev(ptr, k)
    // permutations
   do {} while (next_permutation(nums.begin(), nums.
        end()));
   // bitset
46 int num = 27; // Binary representation: 11011
47 bitset<10> s(string("0010011010")); // from right
        to left
48 bitset<sizeof(int) * 8> bits(num);
49 int set_bits = bits.count();
```

50 bits.set(index, val);

```
51 bits.reset();
52 bits.flip();
53 bits.all();
54 bits.any();
55 bits.none();
56 bits.test();
57 to_string();
58 to_ulong();
59 to_ullong();
60 [], &, |, !, >>=, <<=, &=, |=, ^=, ~;
61 // hamming distance
62 int hamming(int a, int b) {
63
      return __builtin_popcount(a ^ b);
64 }
65 // gcc compiler
66 __builtin_popcount(x)
67 __builtin_parity(x)
68 __builtin_clz(x) // leading
69 __builtin_ctz(x) // trailing
```

### 8.5 Priority Queue

```
1 // declaration
  priority_queue<T> pq;
  priority_queue<T> pq(v.begin(), v.end());
4 priority_queue<T, vector<T>, greater<T>> pq;
5 // custom comparator
```

```
6 class Compare {
    public:
     bool operator() (T a, T b) {
       if (cond) return true; // do not swap
10
       return false;
11
12 };
13 priority_queue<T, vector<T>, Compare> pq;
14 // functions
15 pq.empty();
16 pq.size();
17 pq.top();
18 pq.push(x);
19 pq.pop();
 8.6 Vector
```

```
1 // declaration
 2 vector<T> v:
 3 vector<T> v = \{v0, v1, v2, ...\};
 4 vector<T> v(size, initial);
 5 // functions
 6 v.begin();
 7 v.end();
 8 v.size();
 9 v.empty();
10 v.push_back(x);
```

```
v.pop_back();
   v.insert();
   v.erase(x);
14 v.clear();
15 // algorithms
16 lower_bound(v.begin(), v.end(), x);
   upper_bound(v.begin(), v.end(), x);
18 binary_search(v.begin(), v.end(), x);
19 // sort
20     sort(v.begin(), v.end());
21 sort(v.rbegin(), v.rend()); // reverse iterators
22 sort(v.begin(), v.end(), greater<T>()); // using
        functor
23 bool comp(T a, T b) {
    if (cond) {
       return true; // do not swap
26
27
     return false;
28
   sort(v.begin(), v.end(), comp); // using custom
30 sort(v.begin(), v.end(), [](const T a, const T b) {
     if (cond) {
       return true;
34
     return false;
35 }); // using lambda function
```

f(n) = O(g(n))	iff $\exists$ positive $c, n_0$ such that $0 \le f(n) \le cg(n) \ \forall n \ge n_0$ .	$\sum_{i=1}^{n} i = \frac{n(n+1)}{2},  \sum_{i=1}^{n} i^{2} = \frac{n(n+1)(2n+1)}{6},  \sum_{i=1}^{n} i^{3} = \frac{n^{2}(n+1)^{2}}{4}.$					
$f(n) = \Omega(g(n))$	iff $\exists$ positive $c, n_0$ such that $f(n) \geq cg(n) \geq 0 \ \forall n \geq n_0$ .	i=1 $i=1$ $i=1$ In general:					
$f(n) = \Theta(g(n))$	iff $f(n) = O(g(n))$ and $f(n) = \Omega(g(n))$ .	$\sum_{i=1}^{n} i^{m} = \frac{1}{m+1} \left[ (n+1)^{m+1} - 1 - \sum_{i=1}^{n} \left( (i+1)^{m+1} - i^{m+1} - (m+1)i^{m} \right) \right]$					
f(n) = o(g(n))	iff $\lim_{n\to\infty} f(n)/g(n) = 0$ .	$\sum_{i=1}^{n-1} i^m = \frac{1}{m+1} \sum_{k=0}^m \binom{m+1}{k} B_k n^{m+1-k}.$					
$ \lim_{n \to \infty} a_n = a $	iff $\forall \epsilon > 0$ , $\exists n_0$ such that $ a_n - a  < \epsilon$ , $\forall n \ge n_0$ .	Geometric series:					
$\sup S$	least $b \in \mathbb{R}$ such that $b \geq s$ , $\forall s \in S$ .	$\sum_{i=0}^{n} c^{i} = \frac{c^{n+1} - 1}{c - 1},  c \neq 1,  \sum_{i=0}^{\infty} c^{i} = \frac{1}{1 - c},  \sum_{i=1}^{\infty} c^{i} = \frac{c}{1 - c},   c  < 1,$					
$\inf S$	greatest $b \in \mathbb{R}$ such that $b \le s$ , $\forall s \in S$ .	$\sum_{i=0}^{n} ic^{i} = \frac{nc^{n+2} - (n+1)c^{n+1} + c}{(c-1)^{2}},  c \neq 1,  \sum_{i=0}^{\infty} ic^{i} = \frac{c}{(1-c)^{2}},   c  < 1.$					
$ \liminf_{n \to \infty} a_n $	$\lim_{n\to\infty}\inf\{a_i\mid i\geq n, i\in\mathbb{N}\}.$	Harmonic series: $n = \sum_{i=1}^{n} 1 \qquad \sum_{i=1}^{n} n(n+1) \qquad n(n-1)$					
$\limsup_{n \to \infty} a_n$	$\lim_{n \to \infty} \sup \{ a_i \mid i \ge n, i \in \mathbb{N} \}.$	$H_n = \sum_{i=1}^n \frac{1}{i}, \qquad \sum_{i=1}^n iH_i = \frac{n(n+1)}{2}H_n - \frac{n(n-1)}{4}.$					
$\binom{n}{k}$	Combinations: Size $k$ subsets of a size $n$ set.	$\sum_{i=1}^{n} H_i = (n+1)H_n - n,  \sum_{i=1}^{n} {i \choose m} H_i = {n+1 \choose m+1} \left( H_{n+1} - \frac{1}{m+1} \right).$					
$\begin{bmatrix} n \\ k \end{bmatrix}$	Stirling numbers (1st kind): Arrangements of an $n$ element set into $k$ cycles.	1. $\binom{n}{k} = \frac{n!}{(n-k)!k!}$ , 2. $\sum_{k=0}^{n} \binom{n}{k} = 2^n$ , 3. $\binom{n}{k} = \binom{n}{n-k}$ ,					
$\left\{ egin{array}{c} n \\ k \end{array} \right\}$	Stirling numbers (2nd kind): Partitions of an n element	$4. \binom{n}{k} = \frac{n}{k} \binom{n-1}{k-1}, \qquad 5. \binom{n}{k} = \binom{n-1}{k} + \binom{n-1}{k-1},$					
(m)	set into $k$ non-empty sets.	<b>6.</b> $\binom{n}{m}\binom{m}{k} = \binom{n}{k}\binom{n-k}{m-k},$ <b>7.</b> $\sum_{k=0}^{n} \binom{r+k}{k} = \binom{r+n+1}{n},$					
$\left\langle {n\atop k} \right\rangle$	1st order Eulerian numbers: Permutations $\pi_1\pi_2\pi_n$ on $\{1,2,,n\}$ with $k$ ascents.	8. $\sum_{k=0}^{n} {k \choose m} = {n+1 \choose m+1},$ 9. $\sum_{k=0}^{n} {r \choose k} {s \choose n-k} = {r+s \choose n},$					
$\left\langle\!\left\langle {n\atop k}\right\rangle\!\right\rangle$	2nd order Eulerian numbers.	<b>10.</b> $\binom{n}{k} = (-1)^k \binom{k-n-1}{k},$ <b>11.</b> $\binom{n}{1} = \binom{n}{n} = 1,$					
$C_n$	Catalan Numbers: Binary trees with $n+1$ vertices.	<b>12.</b> $\binom{n}{2} = 2^{n-1} - 1,$ <b>13.</b> $\binom{n}{k} = k \binom{n-1}{k} + \binom{n-1}{k-1},$					
<b>14.</b> $\begin{bmatrix} n \\ 1 \end{bmatrix} = (n-1)!,$ <b>15.</b> $\begin{bmatrix} n \\ 2 \end{bmatrix} = (n-1)!H_{n-1},$ <b>16.</b> $\begin{bmatrix} n \\ n \end{bmatrix} = 1,$ <b>17.</b> $\begin{bmatrix} n \\ k \end{bmatrix} \ge \begin{Bmatrix} n \\ k \end{Bmatrix},$							
$18. \begin{bmatrix} n \\ k \end{bmatrix} = (n-1) \begin{bmatrix} n-1 \\ k \end{bmatrix} + \begin{bmatrix} n-1 \\ k-1 \end{bmatrix},  19. \begin{bmatrix} n \\ n-1 \end{bmatrix} = \begin{bmatrix} n \\ n-1 \end{bmatrix} = \begin{pmatrix} n \\ 2 \end{pmatrix},  20. \sum_{k=0}^{n} \begin{bmatrix} n \\ k \end{bmatrix} = n!,  21. \ C_n = \frac{1}{n+1} \binom{2n}{n},$							
$22. \left\langle {n \atop 0} \right\rangle = \left\langle {n \atop n-1} \right\rangle = 1, \qquad 23. \left\langle {n \atop k} \right\rangle = \left\langle {n \atop n-1-k} \right\rangle, \qquad 24. \left\langle {n \atop k} \right\rangle = (k+1) \left\langle {n-1 \atop k} \right\rangle + (n-k) \left\langle {n-1 \atop k-1} \right\rangle,$							
$25. \  \left\langle \begin{matrix} 0 \\ k \end{matrix} \right\rangle = \left\{ \begin{matrix} 1 & \text{if } k = 0, \\ 0 & \text{otherwise} \end{matrix} \right. $ $26. \  \left\langle \begin{matrix} n \\ 1 \end{matrix} \right\rangle = 2^n - n - 1, $ $27. \  \left\langle \begin{matrix} n \\ 2 \end{matrix} \right\rangle = 3^n - (n+1)2^n + \binom{n+1}{2}, $							
<b>28.</b> $x^n = \sum_{k=0}^n \binom{n}{k} \binom{x+k}{n}$ , <b>29.</b> $\binom{n}{m} = \sum_{k=0}^m \binom{n+1}{k} (m+1-k)^n (-1)^k$ , <b>30.</b> $m! \binom{n}{m} = \sum_{k=0}^n \binom{n}{k} \binom{k}{n-m}$ ,							
		<b>32.</b> $\left\langle \left\langle {n\atop 0} \right\rangle \right\rangle = 1,$ <b>33.</b> $\left\langle \left\langle {n\atop n} \right\rangle \right\rangle = 0$ for $n \neq 0,$					
$34. \; \left\langle \!\! \left\langle \!\! \begin{array}{c} n \\ k \end{array} \!\! \right\rangle \!\! \right\rangle = (k + 1)^n$	$-1$ ) $\left\langle \left\langle {n-1 \atop k} \right\rangle \right\rangle + (2n-1-k)\left\langle \left\langle {n-1 \atop k} \right\rangle \right\rangle$						
$36. \left\{ \begin{array}{c} x \\ x-n \end{array} \right\} = \sum_{k}^{n} \left\{ \begin{array}{c} x \\ x \end{array} \right\}$	$\sum_{k=0}^{n} \left\langle \!\! \left\langle n \atop k \right\rangle \!\! \right\rangle \left( \begin{matrix} x+n-1-k \\ 2n \end{matrix} \right),$	<b>37.</b> ${n+1 \choose m+1} = \sum_{k} {n \choose k} {k \choose m} = \sum_{k=0}^{n} {k \choose m} (m+1)^{n-k},$					

The Chinese remainder theorem: There exists a number C such that:

 $C \equiv r_1 \mod m_1$ 

: : :

 $C \equiv r_n \bmod m_n$ 

if  $m_i$  and  $m_j$  are relatively prime for  $i \neq j$ . Euler's function:  $\phi(x)$  is the number of positive integers less than x relatively prime to x. If  $\prod_{i=1}^{n} p_i^{e_i}$  is the prime factorization of x then

$$\phi(x) = \prod_{i=1}^{n} p_i^{e_i - 1} (p_i - 1).$$

Euler's theorem: If a and b are relatively prime then

$$1 \equiv a^{\phi(b)} \bmod b.$$

Fermat's theorem:

$$1 \equiv a^{p-1} \bmod p.$$

The Euclidean algorithm: if a > b are integers then

$$gcd(a, b) = gcd(a \mod b, b).$$

If  $\prod_{i=1}^{n} p_i^{e_i}$  is the prime factorization of x

$$S(x) = \sum_{d|x} d = \prod_{i=1}^{n} \frac{p_i^{e_i+1} - 1}{p_i - 1}.$$

Perfect Numbers: x is an even perfect number iff  $x = 2^{n-1}(2^n - 1)$  and  $2^n - 1$  is prime. Wilson's theorem: n is a prime iff

$$(n-1)! \equiv -1 \mod n$$
.

$$\mu(i) = \begin{cases} (n-1)! = -1 \bmod n. \\ \text{M\"obius inversion:} \\ \mu(i) = \begin{cases} 1 & \text{if } i = 1. \\ 0 & \text{if } i \text{ is not square-free.} \\ (-1)^r & \text{if } i \text{ is the product of} \\ r & \text{distinct primes.} \end{cases}$$
 If

 $\operatorname{If}$ 

$$G(a) = \sum_{d|a} F(d),$$

$$F(a) = \sum_{d \mid a} \mu(d) G\left(\frac{a}{d}\right).$$

Prime numbers:

$$p_n = n \ln n + n \ln \ln n - n + n \frac{\ln \ln n}{\ln n}$$

$$+O\left(\frac{n}{\ln n}\right),$$

$$\pi(n) = \frac{n}{\ln n} + \frac{n}{(\ln n)^2} + \frac{2!n}{(\ln n)^3} + O\left(\frac{n}{(\ln n)^4}\right).$$

-	_		
- 11	Otir	nit:	ions
-	cm	II U.	min

An edge connecting a ver-Looptex to itself. Directed Each edge has a direction.

SimpleGraph with no loops or multi-edges.

WalkA sequence  $v_0e_1v_1\ldots e_\ell v_\ell$ . TrailA walk with distinct edges. Pathtrail with distinct

vertices.

ConnectedA graph where there exists a path between any two

vertices.

ComponentΑ maximal connected subgraph.

TreeA connected acyclic graph. Free tree A tree with no root. DAGDirected acyclic graph. EulerianGraph with a trail visiting each edge exactly once.

Hamiltonian Graph with a cycle visiting each vertex exactly once.

CutA set of edges whose removal increases the number of components.

Cut-setA minimal cut. Cut edge A size 1 cut.

k-Connected A graph connected with the removal of any k-1vertices.

k-Tough  $\forall S \subseteq V, S \neq \emptyset$  we have  $k \cdot c(G - S) \le |S|$ .

A graph where all vertices k-Regular have degree k.

k-Factor Α k-regular spanning subgraph.

Matching A set of edges, no two of which are adjacent.

CliqueA set of vertices, all of which are adjacent.

Ind. set A set of vertices, none of which are adjacent.

Vertex cover A set of vertices which cover all edges.

Planar graph A graph which can be embeded in the plane.

Plane graph An embedding of a planar

$$\sum_{v \in V} \deg(v) = 2m.$$

If G is planar then n - m + f = 2, so

$$f \le 2n - 4, \quad m \le 3n - 6.$$

Any planar graph has a vertex with degree  $\leq 5$ .

### Notation:

E(G)Edge set Vertex set V(G)

c(G)Number of components

G[S]Induced subgraph deg(v)Degree of v

Maximum degree  $\Delta(G)$ 

 $\delta(G)$ Minimum degree  $\chi(G)$ Chromatic number

 $\chi_E(G)$ Edge chromatic number  $G^c$ Complement graph  $K_n$ Complete graph

 $K_{n_1,n_2}$ Complete bipartite graph

Ramsev number

### Geometry

Projective coordinates: (x, y, z), not all x, y and z zero.

 $(x, y, z) = (cx, cy, cz) \quad \forall c \neq 0.$ 

Cartesian Projective (x, y)(x, y, 1)y = mx + b(m, -1, b)x = c(1,0,-c)

Distance formula,  $L_p$  and  $L_{\infty}$ 

$$\sqrt{(x_1 - x_0)^2 + (y_1 - y_0)^2},$$
$$[|x_1 - x_0|^p + |y_1 - y_0|^p]^{1/p},$$

$$\lim_{p \to \infty} \left[ |x_1 - x_0|^p + |y_1 - y_0|^p \right]^{1/p}.$$

Area of triangle  $(x_0, y_0), (x_1, y_1)$ and  $(x_2, y_2)$ :

$$\frac{1}{2} \operatorname{abs} \begin{vmatrix} x_1 - x_0 & y_1 - y_0 \\ x_2 - x_0 & y_2 - y_0 \end{vmatrix}.$$

Angle formed by three points:

$$(x_2, y_2)$$

$$(0, 0) \qquad \ell_1 \qquad (x_1, y_1)$$

$$\cos \theta = \frac{(x_1, y_1) \cdot (x_2, y_2)}{\ell_1 \ell_2}.$$

Line through two points  $(x_0, y_0)$ and  $(x_1, y_1)$ :

$$\begin{vmatrix} x & y & 1 \\ x_0 & y_0 & 1 \\ x_1 & y_1 & 1 \end{vmatrix} = 0.$$

Area of circle, volume of sphere:

$$A = \pi r^2, \qquad V = \frac{4}{3}\pi r^3.$$

If I have seen farther than others, it is because I have stood on the shoulders of giants.

- Issac Newton