UPLB Eliens ICPC Notebook (C++)

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Data Structures

1.1 Disjoint Set Union

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```
struct DSU {
  vector<int> parent, size;
  DSU(int n) {
    parent.resize(n);
    size.resize(n);
    for (int i = 0; i < n; i++) make_set(i);</pre>
 void make_set(int v) {
    parent[v] = v;
    size[v] = 1;
 bool is_same(int a, int b) { return find_set(a) ==
     find set(b); }
  int find_set(int v) { return v == parent[v] ? v :
     parent[v] = find_set(parent[v]); }
 void union sets(int a, int b) {
    a = find_set(a);
    b = find set(b);
    if (a != b) {
      if (size[a] < size[b]) swap(a, b);</pre>
      parent[b] = a;
      size[a] += size[b];
};
```

1.2 Minimum Queue

```
11 get_minimum(stack<pair<11, 11>> &s1, stack<pair<11,</pre>
   11>> &s2) {
  if (s1.empty() || s2.empty()) {
    return s1.empty() ? s2.top().second : s1.top().
       second:
  } else {
    return min(s1.top().second, s2.top().second);
void add_element(ll new_element, stack<pair<ll, 11>> &s1
  11 minimum = s1.empty() ? new_element : min(
     new_element, s1.top().second);
  s1.push({new element, minimum});
11 remove element(stack<pair<11, 11>> &s1, stack<pair<11</pre>
   , 11>> \&s2) {
```

1.3 Range Add Point Query

```
template<class T>
class SegTreeNode {
public:
  const T IDENTITY = 0;
  int i, j;
  T value:
  SegTreeNode<T>* lc, * rc;
  SegTreeNode(int i, int j) : i(i), j(j) {
    value = IDENTITY;
    if (i - i == 1) {
     lc = rc = nullptr;
      return;
    int k = (i + j) / 2;
    lc = new SegTreeNode(i, k);
    rc = new SegTreeNode(k, j);
  SegTreeNode(const vector<T>& a, int i, int j) : i(i),
     i(i) {
    value = j - i == 1 ? a[i] : IDENTITY;
    if (i - i == 1) {
     lc = rc = nullptr;
      return;
    int k = (i + j) / 2;
    lc = new SegTreeNode(a, i, k);
    rc = new SegTreeNode(a, k, j);
  void range_update(int 1, int r, T x) {
    if (r <= i || j <= 1) return;</pre>
    if (1 <= i && j <= r) {
      value = op(value, x);
      return;
```

```
lc->range_update(l, r, x);
    rc->range_update(l, r, x);
  T point_query(int k) {
    if (k < i | | j <= k) return IDENTITY;</pre>
    if (i - i == 1) {
      return value;
    return op(value, op(lc->point_query(k), rc->
       point query(k)));
  T \circ p(T \times, T y) \{ return \times + y; \}
};
template<class T>
class SegTree {
public:
  SegTreeNode<T> root;
  SegTree(int n) : root(0, n) {}
  SegTree(const vector<T>& a) : root(a, 0, a.size()) {}
  void range_update(int 1, int r, T x) { root.
     range_update(l, r, x); }
  T point_query(int k) { return root.point_query(k); }
};
```

1.4 Segment Tree

```
template < class T>
class SegTreeNode {
public:
  const T IDENTITY = 0, DEFAULT = 0;
  int i, j;
  T val:
  SegTreeNode<T>* lc, * rc;
  SegTreeNode(int i, int j) : i(i), j(j) {
    if (j - i == 1) {
      lc = rc = nullptr;
      val = DEFAULT;
      return;
    int k = (i + j) / 2;
    lc = new SegTreeNode(i, k);
    rc = new SegTreeNode(k, j);
    val = op(lc->val, rc->val);
  SegTreeNode(const vector<T>& a, int i, int j) : i(i),
     j(j) {
    if (i - i == 1) {
      lc = rc = nullptr;
      val = a[i];
      return;
```

```
int k = (i + j) / 2;
    lc = new SegTreeNode(a, i, k);
    rc = new SegTreeNode(a, k, j);
    val = op(lc->val, rc->val);
  void set (int k, T x) { // update a[k] := x
    if (k < i \mid | i \le k) return;
    if (j - i == 1) {
        val = x;
        return;
    lc->set(k, x);
    rc \rightarrow set(k, x);
    val = op(lc->val, rc->val);
  T range_query(int 1, int r) { // [1, r)
    if (1 <= i && j <= r) return val;
    if (j <= l || r <= i) return IDENTITY;</pre>
    return op(lc->range_query(l, r), rc->range_query(l,
        r));
  T \circ p(T \times, T y) \{\}
template<class T>
class SegTree {
public:
  SegTreeNode<T> root;
  SegTree(int n) : root(0, n) {}
  SegTree(const vector<T>& a) : root(a, 0, a.size()) {}
  void set(int k, T x) { root.set(k, x); }
  T range_query(int 1, int r) { return root.range_query(
     1, r); }
};
```

1.5 Sparse Table

```
11 log2_floor(ll i) {
    return i ? __builtin_clzll(1) - __builtin_clzll(i) :
        -1;
}
vector<vector<ll>> build_sum(ll N, ll K, vector<ll> &
        array) {
    vector<vector<ll>> st(K + 1, vector<ll>(N + 1));
    for (ll i = 0; i < N; i++) st[0][i] = array[i];
    for (ll i = 1; i <= K; i++)
        for (ll j = 0; j + (1 << i) <= N; j++)
        st[i][j] = st[i - 1][j] + st[i - 1][j + (1 << (i - 1))];
    return st;</pre>
```

```
11 sum_query(11 L, 11 R, 11 K, vector<vector<11>> &st) {
 11 \text{ sum} = 0;
  for (11 i = K; i >= 0; i--) {
    if ((1 << i) <= R - L + 1) {
      sum += st[i][L];
      L += 1 << i;
  return sum;
vector<vector<ll>> build_min(ll N, ll K, vector<ll> &
   array) {
  vector<vector<ll>> st(K + 1, vector<ll>(N + 1));
  for (ll i = 0; i < N; i++) st[0][i] = array[i];</pre>
  for (ll i = 1; i \le K; i++)
    for (11 \ j = 0; \ j + (1 << i) <= N; \ j++)
      st[i][j] = min(st[i - 1][j], st[i - 1][j + (1 << (
         i - 1))]);
  return st;
11 min_query(11 L, 11 R, vector<vector<11>> &st) {
 ll i = log2\_floor(R - L + 1);
  return min(st[i][L], st[i][R - (1 << i) + 1]);
```

2 Dynamic Programming

2.1 Divide And Conquer

```
11 m, n;
vector<ll> dp_before(n), dp_cur(n);
11 C(11 i, 11 j);
void compute(ll l, ll r, ll optl, ll optr) {
  if (1 > r) {
    return;
  11 \text{ mid} = (1 + r) >> 1;
  pair<ll, ll > best = \{LLONG\_MAX, -1\};
  for (ll k = optl; k <= min(mid, optr); k++) {</pre>
    best = min(best, \{(k ? dp_before[k - 1] : 0) + C(k,
       mid), k});
  dp_cur[mid] = best.first;
  11 opt = best.second;
  compute(1, mid - 1, optl, opt);
  compute(mid + 1, r, opt, optr);
11 solve() {
```

```
for (ll i = 0; i < n; i++) {
   dp_before[i] = C(0, i);
}
for (ll i = 1; i < m; i++) {
   compute(0, n - 1, 0, n - 1);
   dp_before = dp_cur;
}
return dp_before[n - 1];</pre>
```

2.2 Edit Distance

2.3 Knapsack

2.4 Knuth Optimization

```
11 solve() {
  11 N:
  // read N and input
  \ensuremath{\text{vector}}\ensuremath{\text{ll>}}\ensuremath{\text{(N)}}\ensuremath{\text{)}}, \ensuremath{\text{opt}}\ensuremath{\text{(N, vector)}}
      <11>(N);
  auto C = [\&](11 i, 11 j) {
    // Implement cost function C.
  };
  for (11 i = 0; i < N; i++) {
     opt[i][i] = i;
     ... // Initialize dp[i][i] according to the problem
  for (11 i = N - 2; i >= 0; i--) {
     for (11 \ j = i + 1; \ j < N; \ j++)  {
       ll mn = ll\_MAX, cost = C(i, j);
       for (ll k = opt[i][j-1]; k <= min(j-1, opt[i+1])
            1][\dot{1}]; k++) {
         if (mn >= dp[i][k] + dp[k + 1][j] + cost) {
            opt[i][j] = k;
            mn = dp[i][k] + dp[k + 1][j] + cost;
       dp[i][j] = mn;
  cout << dp[0][N - 1] << ' \n';
```

2.5 Longest Common Subsequence

```
ll LCS(string x, string y, ll n, ll m) {
  vector<vector<ll>> dp(n + 1, vector<ll>(m + 1));
  for (ll i = 0; i <= n; i++) {
    for (ll j = 0; j <= m; j++) {
      if (i == 0 || j == 0) {
         dp[i][j] = 0;
      } else if (x[i - 1] == y[j - 1]) {
         dp[i][j] = dp[i - 1][j - 1] + 1;
      } else {
         dp[i][j] = max(dp[i - 1][j], dp[i][j - 1]);
      }
    }
    ll index = dp[n][m];
    vector<char> lcs(index + 1);
    lcs[index] = '\0';
```

```
ll i = n, j = m;
while (i > 0 && j > 0) {
    if (x[i - 1] == y[j - 1]) {
        lcs[index - 1] = x[i - 1];
        i--;
        j--;
        index--;
    } else if (dp[i - 1][j] > dp[i][j - 1]) {
        i--;
    } else {
        j--;
    }
}
return dp[n][m];
}
```

2.6 Longest Increasing Subsequence

```
11 get ceil idx(vector<11> &a, vector<11> &T, 11 1, 11 r
   , 11 x) {
  while (r - 1 > 1) {
   11 m = 1 + (r - 1) / 2;
    if (a[T[m]] >= x) {
     r = m:
    } else {
      1 = m;
  return r;
11 LIS(ll n, vector<ll> &a) {
  11 len = 1;
  vector<11> T(n, 0), R(n, -1);
  T[0] = 0;
  for (ll i = 1; i < n; i++) {
    if (a[i] < a[T[0]]) {</pre>
      T[0] = i;
    } else if (a[i] > a[T[len - 1]]) {
      R[i] = T[len - 1];
      T[len++] = i;
    } else {
      ll pos = get_ceil_idx(a, T, -1, len - 1, a[i]);
      R[i] = T[pos - 1];
      T[pos] = i;
    }
  return len;
```

2.7 Subset Sum

```
bool subset_sum(ll n, vector<ll> &arr, ll sum) {
  vector<vector<ll>> dp(n + 1, vector<ll>(sum + 1, false
      ));
  dp[0][0] = true;
  for (ll i = 1; i <= n; i++) {
    for (ll j = 0; j <= sum; j++) {
      dp[i][j] = dp[i - 1][j];
      if (j >= arr[i]) {
        dp[i][j] |= dp[i - 1][j - arr[i]];
      }
    }
  }
  return dp[n][sum];
}
```

3 Geometry

3.1 Circle Line Intersection

```
double r, a, b, c; // given as input
double x0 = -a * c / (a * a + b * b);
double y0 = -b * c / (a * a + b * b);
if (c * c > r * r * (a * a + b * b) + EPS) {
  puts ("no points");
else if (abs (c *c - r * r * (a * a + b * b)) < EPS)
  puts ("1 point");
  cout << x0 << ' ' << y0 << '\n';
} else {
  double d = r * r - c * c / (a * a + b * b);
  double mult = sgrt (d / (a * a + b * b));
  double ax, ay, bx, by;
  ax = x0 + b * mult;
 bx = x0 - b * mult;
  ay = y0 - a * mult;
  by = y0 + a * mult;
  puts ("2 points");
  cout << ax << ' ' << ay << '\n' << bx << ' ' << by <<
     '\n';
```

3.2 Convex Hull

```
struct pt {
  double x, y;
```

```
};
11 orientation(pt a, pt b, pt c) {
  double v = a.x * (b.y - c.y) + b.x * (c.y - a.y) + c.x
      * (a.y - b.y);
  if (v < 0) {
    return -1;
  } else if (v > 0) {
    return +1;
  return 0;
bool cw(pt a, pt b, pt c, bool include_collinear) {
  11 o = orientation(a, b, c);
  return o < 0 || (include_collinear && o == 0);</pre>
bool collinear(pt a, pt b, pt c) {
  return orientation(a, b, c) == 0;
void convex_hull(vector<pt>& a, bool include_collinear =
    false) {
  pt p0 = *min_element(a.begin(), a.end(), [](pt a, pt b)
     ) {
    return make pair(a.y, a.x) < make pair(b.y, b.x);</pre>
  });
  sort(a.begin(), a.end(), [&p0](const pt& a, const pt&
     b) {
    11 o = orientation(p0, a, b);
    if (o == 0) {
      return (p0.x - a.x) * (p0.x - a.x) + (p0.y - a.y)
         * (p0.y - a.y)
           < (p0.x - b.x) * (p0.x - b.x) + (p0.y - b.y)
               * (p0.y - b.y);
    return o < 0;</pre>
  });
  if (include collinear) {
    11 i = (11) a.size()-1;
    while (i \ge 0 \&\& collinear(p0, a[i], a.back())) i--;
    reverse(a.begin()+i+1, a.end());
  vector<pt> st;
  for (ll i = 0; i < (ll) a.size(); i++) {</pre>
    while (st.size() > 1 && !cw(st[st.size() - 2], st.
       back(), a[i], include_collinear)) {
      st.pop_back();
    st.push_back(a[i]);
  a = st;
```

3.3 Line Sweep

```
const double EPS = 1E-9;
struct pt {
  double x, y;
struct seq {
  pt p, q;
  ll id;
  double get_y (double x) const {
    if (abs(p.x - q.x) < EPS) {
      return p.y;
    return p.y + (q.y - p.y) * (x - p.x) / (q.x - p.x);
} ;
bool intersect1d(double 11, double r1, double 12, double
    r2) {
  if (11 > r1) {
    swap(11, r1);
  if (12 > r2) {
    swap(12, r2);
  return max(11, 12) <= min(r1, r2) + EPS;
11 vec(const pt& a, const pt& b, const pt& c) {
  double s = (b.x - a.x) * (c.y - a.y) - (b.y - a.y) * (
     c.x - a.x);
  return abs(s) < EPS ? 0 : s > 0 ? +1 : -1;
bool intersect(const seg& a, const seg& b) {
  return intersect1d(a.p.x, a.q.x, b.p.x, b.q.x) &&
         intersect1d(a.p.y, a.q.y, b.p.y, b.q.y) &&
         vec(a.p, a.q, b.p) * vec(a.p, a.q, b.q) <= 0 &&</pre>
         vec(b.p, b.q, a.p) * vec(b.p, b.q, a.q) <= 0;
bool operator<(const seg& a, const seg& b) {
    double x = max(min(a.p.x, a.q.x), min(b.p.x, b.q.x))
    return a.get_y(x) < b.get_y(x) - EPS;</pre>
struct event {
  double x;
  ll tp, id;
  event() {}
  event (double x, ll tp, ll id) : x(x), tp(tp), id(id)
     { }
  bool operator<(const event& e) const {</pre>
```

```
if (abs(x - e.x) > EPS) {
      return x < e.x;
    return tp > e.tp;
 }
};
set<seg> s;
vector<set<seq>::iterator> where;
set<seg>::iterator prev(set<seg>::iterator it) {
  return it == s.begin() ? s.end() : --it;
set<seg>::iterator next(set<seg>::iterator it) {
  return ++it;
pair<11, 11> solve(const vector<seg>& a) {
 11 n = (11) a.size();
 vector<event> e;
  for (11 i = 0; i < n; ++i) {
    e.push_back(event(min(a[i].p.x, a[i].q.x), +1, i));
    e.push_back(event(max(a[i].p.x, a[i].q.x), -1, i));
  sort(e.begin(), e.end());
  s.clear();
  where.resize(a.size());
  for (size_t i = 0; i < e.size(); ++i) {</pre>
   ll id = e[i].id;
    if (e[i].tp == +1) {
      set<seg>::iterator nxt = s.lower_bound(a[id]), prv
          = prev(nxt);
      if (nxt != s.end() && intersect(*nxt, a[id])) {
        return make pair (nxt->id, id);
      if (prv != s.end() && intersect(*prv, a[id])) {
        return make pair (prv->id, id);
      where[id] = s.insert(nxt, a[id]);
    } else {
      set<seg>::iterator nxt = next(where[id]), prv =
         prev(where[id]);
      if (nxt != s.end() && prv != s.end() && intersect
         (*nxt, *prv)) {
        return make pair(prv->id, nxt->id);
      s.erase(where[id]);
  return make pair (-1, -1);
```

3.4 Nearest Points

```
struct pt {
  11 x, y, id;
};
struct cmp x {
  bool operator()(const pt & a, const pt & b) const {
    return a.x < b.x || (a.x == b.x && a.y < b.y);
};
struct cmp_y {
  bool operator()(const pt & a, const pt & b) const {
    return a.y < b.y;</pre>
};
11 n;
vector<pt> a;
double mindist;
pair<11, 11> best_pair;
void upd ans(const pt & a, const pt & b) {
  double dist = sqrt((a.x - b.x) * (a.x - b.x) + (a.y - b.x)
     b.y) * (a.y - b.y));
  if (dist < mindist) {</pre>
    mindist = dist;
    best_pair = {a.id, b.id};
vector<pt> t;
void rec(ll l, ll r) {
  if (r - 1 \le 3) {
    for (ll i = 1; i < r; ++i) {
      for (11 j = i + 1; j < r; ++j) {
        upd_ans(a[i], a[j]);
    sort(a.begin() + 1, a.begin() + r, cmp_y());
    return;
  11 m = (1 + r) >> 1, midx = a[m].x;
  rec(1, m);
  rec(m, r);
  merge(a.begin() + l, a.begin() + m, a.begin() + m, a.
     begin() + r, t.begin(), cmp_y();
  copy(t.begin(), t.begin() + r - l, a.begin() + l);
  11 tsz = 0;
  for (11 i = 1; i < r; ++i) {
    if (abs(a[i].x - midx) < mindist) {</pre>
      for (11 \dot{j} = tsz - 1; \dot{j} >= 0 && a[i].y - t[\dot{j}].y <
         mindist; --j) {
```

```
upd_ans(a[i], t[j]);
}
    t[tsz++] = a[i];
}
}
t.resize(n);
sort(a.begin(), a.end(), cmp_x());
mindist = 1E20;
rec(0, n);
```

4 Graph Theory

4.1 Articulation Point

```
void APUtil(vector<vector<11>> &adj, ll u, vector<bool>
   &visited.
vector<ll> &disc, vector<ll> &low, ll &time, ll parent,
   vector<bool> &isAP) {
  11 \text{ children} = 0;
  visited[u] = true;
  disc[u] = low[u] = ++time;
  for (auto v : adj[u]) {
    if (!visited[v]) {
      children++;
      APUtil(adj, v, visited, disc, low, time, u, isAP);
      low[u] = min(low[u], low[v]);
      if (parent !=-1 \&\& low[v] >= disc[u]) {
        isAP[u] = true;
    } else if (v != parent) {
      low[u] = min(low[u], disc[v]);
    }
  if (parent == -1 && children > 1) {
    isAP[u] = true;
void AP(vector<vector<ll>>> &adj, ll n) {
  vector<ll> disc(n), low(n);
  vector<bool> visited(n), isAP(n);
  11 time = 0, par = -1;
  for (11 u = 0; u < n; u++) {
    if (!visited[u]) {
      APUtil(adj, u, visited, disc, low, time, par, isAP
         );
  for (11 u = 0; u < n; u++) {
```

```
if (isAP[u]) {
    cout << u << " ";
    }
}</pre>
```

4.2 Bellman Ford

```
void bellman ford(vector<vector<ll>> &edges, ll n, ll m,
    ll src, vector<ll> &dis) {
  for (ll i = 0; i < n; i++) {
    dis[i] = INF;
  for (11 i = 0; i < n - 1; i++) {
    for (11 j = 0; j < m; j++) {
      ll\ u = edges[j][0],\ v = edges[j][1],\ w = edges[j]
         1[2];
      if (dis[u] < INF) {</pre>
        dis[v] = min(dis[v], dis[u] + w);
  for (11 i = 0; i < m; i++) {
    ll u = edges[i][0], v = edges[i][1], w = edges[i]
       1[2];
    if (dis[u] < INF && dis[u] + w < dis[v]) {</pre>
      cout << "The graph contains a negative cycle." <<</pre>
          '\n';
```

4.3 Bridge

```
void bridge_util(vector<vector<1l>> &adj, ll u, vector<
   bool> &visited, vector<1l> &disc, vector<1l> &low,
   vector<1l> &parent) {
   static ll time = 0;
   visited[u] = true;
   disc[u] = low[u] = ++time;
   list<1l>::iterator i;
   for (auto v : adj[u]) {
      if (!visited[v]) {
        parent[v] = u;
        bridge_util(adj, v, visited, disc, low, parent);
      low[u] = min(low[u], low[v]);
      if (low[v] > disc[u]) {
        cout << u << ' ' << v << '\n';
    }
}</pre>
```

```
}
    else if (v != parent[u]) {
        low[u] = min(low[u], disc[v]);
    }
}

void bridge(vector<vector<ll>> &adj, ll n) {
    vector<bool> visited(n, false);
    vector<ll> disc(n), low(n), parent(n, -1);
    for (ll i = 0; i < n; i++) {
        if (!visited[i]) {
            bridge_util(adj, i, visited, disc, low, parent);
        }
    }
}</pre>
```

4.4 Dijkstra

```
void dijkstra(ll n, vector<vector<pair<ll, ll>>> &adj,
   vector<ll> &dis) {
  priority_queue<pair<11, 11>, vector<pair<11, 11>>,
     greater<pair<11, 11>>> pq;
  for (int i = 0; i < n; i++) {
    dis[i] = INF;
  dis[0] = 0;
  pq.push({0, 0});
  while (!pq.empty()) {
    auto p = pq.top();
    pq.pop();
    11 u = p.second;
    if (dis[u] != p.first) {
      continue;
    for (auto x : adj[u]) {
     ll v = x.first, w = x.second;
      if (dis[v] > dis[u] + w) {
        dis[v] = dis[u] + w;
        pq.push({dis[v], v});
```

4.5 Find Cycle

```
bool dfs(ll v) {
  color[v] = 1;
```

```
for (ll u : adj[v]) {
    if (color[u] == 0) {
      parent[u] = v;
      if (dfs(u)) {
        return true;
    } else if (color[u] == 1) {
      cycle end = v;
      cycle_start = u;
      return true;
  color[v] = 2;
  return false;
void find cycle() {
  color.assign(n, 0);
  parent.assign(n, -1);
  cycle_start = -1;
  for (11 v = 0; v < n; v++) {
    if (color[v] == 0 && dfs(v)) {
      break;
    }
  if (cycle start == -1) {
    cout << "Acyclic" << endl;</pre>
  } else {
    vector<ll> cycle;
    cycle.push back(cycle start);
    for (ll v = cycle_end; v != cycle_start; v = parent[
       v]) {
      cycle.push_back(v);
    cycle.push back(cycle start);
    reverse(cycle.begin(), cycle.end());
    cout << "Cycle found: ";</pre>
    for (11 v : cycle) {
      cout << v << ' ';
    cout << '\n';
```

4.6 Floyd Warshall

```
void floyd_warshall(vector<vector<ll>>> &dis, ll n) {
  for (ll i = 0; i < n; i++) {
    for (ll j = 0; j < n; j++) {
      dis[i][j] = (i == j ? 0 : INF);
    }</pre>
```

4.7 Hierholzer

```
void print circuit(vector<vector<ll>> &adj) {
  map<11, 11> edge count;
  for (ll i = 0; i < adj.size(); i++) {</pre>
    edge count[i] = adj[i].size();
  if (!adj.size()) {
    return;
  stack<ll> curr path;
  vector<ll> circuit;
  curr path.push(0);
  11 \text{ curr } v = 0;
  while (!curr_path.empty()) {
    if (edge count[curr v]) {
      curr_path.push(curr_v);
      11 next_v = adj[curr_v].back();
      edge_count[curr_v]--;
      adj[curr_v].pop_back();
      curr_v = next_v;
    } else {
      circuit.push_back(curr_v);
      curr v = curr path.top();
      curr_path.pop();
```

```
for (ll i = circuit.size() - 1; i >= 0; i--) {
   cout << circuit[i] << ' ';
}</pre>
```

4.8 Is Bipartite

```
bool is_bipartite(vector<ll> &col, vector<vector<ll>> &
   adj, 11 n) {
  queue<pair<ll, ll>> q;
  for (ll i = 0; i < n; i++) {
    if (col[i] == -1) {
      q.push(\{i, 0\});
      col[i] = 0;
      while (!q.empty()) {
        pair<11, 11> p = q.front();
        q.pop();
        11 v = p.first, c = p.second;
        for (ll i : adi[v]) {
          if (col[j] == c) {
            return false;
          if (col[j] == -1) {
            col[j] = (c ? 0 : 1);
            q.push({j, col[j]});
  return true;
```

4.9 Is Cyclic

```
bool is_cyclic_util(int u, vector<vector<int>> &adj,
    vector<bool> &vis, vector<bool> &rec) {
    vis[u] = true;
    rec[u] = true;
    for(auto v : adj[u]) {
        if (!vis[v] && is_cyclic_util(v, adj, vis, rec)) {
            return true;
        } else if (rec[v]) {
            return true;
        }
    }
    rec[u] = false;
    return false;
```

```
bool is_cyclic(int n, vector<vector<int>> &adj) {
  vector<bool> vis(n, false), rec(n, false);
  for (int i = 0; i < n; i++) {
    if (!vis[i] && is_cyclic_util(i, adj, vis, rec)) {
      return true;
    }
  }
  return false;
}</pre>
```

4.10 Kahn

```
void kahn(vector<vector<ll>>> &adj) {
  ll n = adj.size();
  vector<ll> in degree(n, 0);
  for (11 u = 0; u < n; u++) {
    for (ll v: adi[u]) {
      in degree[v]++;
  queue<11> q;
  for (11 i = 0; i < n; i++) {</pre>
    if (in degree[i] == 0) {
      q.push(i);
  11 cnt = 0;
  vector<ll> top_order;
  while (!q.empty()) {
    ll u = q.front();
    q.pop();
    top_order.push_back(u);
    for (ll v : adj[u]) {
      if (--in_degree[v] == 0) {
        q.push(v);
      }
    }
    cnt++;
  if (cnt != n) {
    cout << -1 << '\n';
    return;
  for (11 i = 0; i < (11) top_order.size(); i++) {</pre>
    cout << top_order[i] << ' ';</pre>
  cout << '\n';
```

4.11 Kruskal Mst

```
struct Edge {
  ll u, v, weight;
 bool operator<(Edge const& other) {</pre>
    return weight < other.weight;</pre>
};
11 n;
vector<Edge> edges;
11 cost = 0;
vector<ll> tree id(n);
vector<Edge> result;
for (11 i = 0; i < n; i++) {
  tree_id[i] = i;
sort(edges.begin(), edges.end());
for (Edge e : edges) {
  if (tree id[e.u] != tree id[e.v]) {
    cost += e.weight;
    result.push_back(e);
    ll old_id = tree_id[e.u], new_id = tree_id[e.v];
    for (11 i = 0; i < n; i++) {
      if (tree_id[i] == old_id) {
        tree id[i] = new id;
```

4.12 Lowest Common Ancestor

```
struct LCA {
  vector<1l> height, euler, first, segtree;
  vector<bool> visited;
  ll n;
  LCA(vector<vector<1l>> &adj, ll root = 0) {
    n = adj.size();
    height.resize(n);
    first.resize(n);
    euler.reserve(n * 2);
    visited.assign(n, false);
    dfs(adj, root);
    ll m = euler.size();
    segtree.resize(m * 4);
    build(1, 0, m - 1);
}
  void dfs(vector<vector<1l>> &adj, ll node, ll h = 0) {
```

```
visited[node] = true;
    height[node] = h;
    first[node] = euler.size();
    euler.push_back(node);
    for (auto to : adj[node]) {
      if (!visited[to]) {
        dfs(adj, to, h + 1);
        euler.push back(node);
    }
  void build(ll node, ll b, ll e) {
    if (b == e) {
      segtree[node] = euler[b];
    } else {
      11 \text{ mid} = (b + e) / 2;
      build(node << 1, b, mid);</pre>
      build(node << 1 | 1, mid + 1, e);</pre>
      11 l = segtree[node << 1], r = segtree[node << 1 |</pre>
           11;
      segtree[node] = (height[l] < height[r]) ? l : r;</pre>
  11 query(11 node, 11 b, 11 e, 11 L, 11 R) {
    if (b > R | | e < L) {
      return -1;
    if (b >= L && e <= R) {
      return segtree[node];
    11 \text{ mid} = (b + e) >> 1;
    11 left = query(node << 1, b, mid, L, R);</pre>
    11 right = query(node << 1 | 1, mid + 1, e, L, R);</pre>
    if (left == -1) return right;
    if (right == -1) return left;
    return height[left] < height[right] ? left : right;</pre>
  ll lca(ll u, ll v) {
    11 left = first[u], right = first[v];
    if (left > right) {
      swap(left, right);
    return query(1, 0, euler.size() - 1, left, right);
} ;
```

4.13 Maximum Bipartite Matching

```
bool bpm(ll n, ll m, vector<vector<bool>> &bpGraph, ll u
, vector<bool> &seen, vector<ll>> &matchR) {
```

```
for (11 v = 0; v < m; v++) {
    if (bpGraph[u][v] && !seen[v]) {
      seen[v] = true;
      if (matchR[v] < 0 || bpm(n, m, bpGraph, matchR[v],</pre>
           seen, matchR)) {
        matchR[v] = u;
        return true;
  return false;
11 maxBPM(11 n, 11 m, vector<vector<bool>> &bpGraph) {
 vector<ll> matchR(m, -1);
 11 \text{ result} = 0;
  for (11 u = 0; u < n; u++) {
    vector<bool> seen(m, false);
    if (bpm(n, m, bpGraph, u, seen, matchR)) {
      result++;
    }
  return result;
```

4.14 Max Flow

```
bool bfs(ll n, vector<vector<ll>> &r_graph, ll s, ll t,
   vector<ll> &parent) {
  vector<bool> visited(n, false);
  queue<11> q;
  q.push(s);
  visited[s] = true;
  parent[s] = -1;
  while (!q.empty()) {
    11 u = q.front();
    q.pop();
    for (11 \ v = 0; \ v < n; \ v++) {
      if (!visited[v] && r_graph[u][v] > 0) {
        if (v == t) {
          parent[v] = u;
          return true;
        q.push(v);
        parent[v] = u;
        visited[v] = true;
  return false;
```

```
11 fordFulkerson(ll n, vector<vector<ll>> graph, ll s,
   11 t) {
  11 u, v;
  vector<vector<ll>> r_graph;
  for (u = 0; u < n; u++) {
    for (v = 0; v < n; v++) {
      r_{graph}[u][v] = graph[u][v];
  vector<ll> parent;
  11 \text{ max flow} = 0;
  while (bfs(n, r_graph, s, t, parent)) {
    11 path flow = INF;
    for (v = t; v != s; v = parent[v]) {
      u = parent[v];
      path_flow = min(path_flow, r_graph[u][v]);
    for (v = t; v != s; v = parent[v]) {
     u = parent[v];
      r_graph[u][v] -= path_flow;
      r_graph[v][u] += path_flow;
    max_flow += path_flow;
  return max flow;
```

4.15 Prim Mst.

```
vector<ll> prim_mst(ll n, vector<vector<pair<ll, ll>>> &
   adj) {
  priority_queue<pair<11, 11>, vector<pair<11, 11>>,
     greater<pair<11, 11>>> pq;
  11 \text{ src} = 0;
  vector<ll> key(n, INF), parent(n, -1);
  vector<bool> in_mst(n, false);
  pq.push(make_pair(0, src));
  kev[src] = 0;
  while (!pq.empty()) {
   11 u = pq.top().second;
    pq.pop();
    if (in_mst[u]) {
      continue;
    in_mst[u] = true;
    for (auto p : adj[u]) {
      ll v = p.first, w = p.second;
      if (in_mst[v] == false && w < key[v]) {</pre>
        key[v] = w;
        pq.push(make_pair(key[v], v));
```

```
parent[v] = u;
}
}
return parent;
}
```

4.16 Strongly Connected Component

```
void dfs(ll u, vector<vector<ll>> &adj, vector<bool> &
   visited) {
 visited[u] = true;
  cout << u + 1 << ' ';
  for (ll v : adi[u]) {
    if (!visited[v]) {
      dfs(v, adj, visited);
vector<vector<ll>> get transpose(ll n, vector<vector<ll</pre>
   >> &adi) {
  vector<vector<ll>> res(n);
  for (11 u = 0; u < n; u++) {
    for (ll v : adj[u]) {
      res[v].push_back(u);
  return res;
void fill_order(ll u, vector<vector<ll>> &adj, vector<</pre>
   bool> &visited, stack<ll> &stk) {
  visited[u] = true;
  for(auto v : adj[u]) {
    if(!visited[v]) {
      fill_order(v, adj, visited, stk);
  stk.push(u);
void get_scc(ll n, vector<vector<ll>> &adj) {
  stack<ll> stk;
  vector<bool> visited(n, false);
  for (ll i = 0; i < n; i++) {
    if (!visited[i]) {
      fill_order(i, adj, visited, stk);
    }
  vector<vector<ll>>> transpose = get_transpose(n, adj);
  for (11 i = 0; i < n; i++) {
    visited[i] = false;
```

```
}
while (!stk.empty()) {
    ll u = stk.top();
    stk.pop();
    if (!visited[u]) {
        dfs(u, transpose, visited);
        cout << '\n';
    }
}</pre>
```

4.17 Topological Sort

```
void dfs(ll v) {
    visited[v] = true;
    for (ll u : adj[v]) {
        if (!visited[u]) {
            dfs(u);
        }
    }
    ans.push_back(v);
}

void topological_sort() {
    visited.assign(n, false);
    ans.clear();
    for (ll i = 0; i < n; ++i) {
        if (!visited[i]) {
            dfs(i);
        }
    }
    reverse(ans.begin(), ans.end());
}</pre>
```

5 Miscellaneous

5.1 Gauss

```
if (abs (a[sel][col]) < EPS) {</pre>
    continue;
 for (ll i = col; i <= m; ++i) {
    swap(a[sel][i], a[row][i]);
 where [col] = row;
 for (11 i = 0; i < n; ++i) {
    if (i != row) {
      double c = a[i][col] / a[row][col];
      for (ll j = col; j <= m; ++j) {
        a[i][j] -= a[row][j] * c;
  ++row;
ans.assign(m, 0);
for (11 i = 0; i < m; ++i) {
 if (where[i] != -1) {
    ans[i] = a[where[i]][m] / a[where[i]][i];
for (11 i = 0; i < n; ++i) {
  double sum = 0;
  for (11 j = 0; j < m; ++j) {
    sum += ans[j] * a[i][j];
  if (abs (sum - a[i][m]) > EPS) {
    return 0;
for (ll i = 0; i < m; ++i) {
 if (where [i] == -1) {
    return INF;
return 1;
```

5.2 Ternary Search

```
double ternary_search(double 1, double r) {
  double eps = 1e-9;
  while (r - 1 > eps) {
    double m1 = 1 + (r - 1) / 3;
    double m2 = r - (r - 1) / 3;
    double f1 = f(m1);
```

```
double f2 = f(m2);
if (f1 < f2) {
    1 = m1;
} else {
    r = m2;
}
return f(1);
}</pre>
```

6 Number Theory

6.1 Extended Euclidean

```
1l gcd_extended(ll a, ll b, ll &x, ll &y) {
   if (b == 0) {
      x = 1;
      y = 0;
      return a;
   }
   ll x1, y1, g = gcd_extended(b, a % b, x1, y1);
   x = y1;
   y = x1 - (a / b) * y1;
   return g;
}
```

6.2 Find All Solutions

```
11 find_all_solutions(ll a, ll b, ll c, ll minx, ll maxx
   , ll miny, ll maxy) {
  11 x, y, q;
  if (!find_any_solution(a, b, c, x, y, g)) {
    return 0;
  a /= q;
  b /= q;
  11 \text{ sign}_a = a > 0 ? +1 : -1;
  shift_solution(x, y, a, b, (minx - x) / b);
  if (x < minx) {
    shift_solution(x, y, a, b, sign_b);
  if (x > maxx) {
    return 0;
  11 1x1 = x;
  shift_solution(x, y, a, b, (maxx - x) / b);
  if (x > maxx) {
    shift_solution(x, y, a, b, -sign_b);
  11 rx1 = x;
  shift_solution(x, y, a, b, -(miny - y) / a);
  if (y < miny) {
    shift_solution(x, y, a, b, -sign_a);
  if (y > maxy) {
    return 0;
  11 \ 1x2 = x;
  shift_solution(x, y, a, b, -(maxy - y) / a);
  if (y > maxy) {
    shift_solution(x, y, a, b, sign_a);
  11 \text{ rx2} = x;
  if (1x2 > rx2) {
    swap(1x2, rx2);
  11 1x = max(1x1, 1x2), rx = min(rx1, rx2);
  if (lx > rx) {
    return 0;
  return (rx - lx) / abs(b) + 1;
```

6.3 Linear Sieve

void linear_sieve(ll N, vector<ll> &lowest_prime, vector

```
<1l> &prime) {
for (ll i = 2; i <= N; i++) {
   if (lowest_prime[i] == 0) {
     lowest_prime[i] = i;
     prime.push_back(i);
   }
   for (ll j = 0; i * prime[j] <= N; j++) {
     lowest_prime[i * prime[j]] = prime[j];
     if (prime[j] == lowest_prime[i]) {
        break;
     }
   }
}</pre>
```

6.4 Miller Rabin

```
bool check_composite(u64 n, u64 a, u64 d, 11 s) {
  u64 x = binpower(a, d, n);
  if (x == 1 | | x == n - 1)  {
    return false:
  for (ll r = 1; r < s; r++) {
    x = (u128) x * x % n;
    if (x == n - 1) {
      return false;
  return true;
bool miller_rabin(u64 n) {
  if (n < 2) {
    return false;
  11 r = 0;
  u64 d = n - 1;
  while ((d \& 1) == 0) {
    d >>= 1;
    <u>r</u>++;
  for (11 a : {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31,
     37}) {
    if (n == a) {
      return true;
    if (check_composite(n, a, d, r)) {
      return false;
    }
  return true;
```

6.5 Modulo Inverse

```
11 mod_inv(ll a, ll m) {
   if (m == 1) {
      return 0;
   }
   ll m0 = m, x = 1, y = 0;
   while (a > 1) {
      ll q = a / m, t = m;
      m = a % m;
      a = t;
      t = y;
      y = x - q * y;
      x = t;
   }
   if (x < 0) {
      x += m0;
   }
   return x;
}</pre>
```

6.6 Pollard Rho Brent

```
11 mult(11 a, 11 b, 11 mod) {
  return ( int128 t) a * b % mod;
11 f(11 x, 11 c, 11 mod) {
  return (mult(x, x, mod) + c) % mod;
ll pollard rho brent(ll n, ll x0 = 2, ll c = 1) {
  11 x = x0, g = 1, q = 1, xs, y, m = 128, l = 1;
  while (q == 1) {
    v = x;
    for (11 i = 1; i < 1; i++) {
      x = f(x, c, n);
    11 k = 0;
    while (k < 1 \&\& q == 1) {
      xs = x;
      for (ll i = 0; i < m && i < l - k; i++) {
        x = f(x, c, n);
        q = mult(q, abs(y - x), n);
      g = \underline{gcd}(q, n);
      k += m;
```

```
1 *= 2;
}
if (g == n) {
    do {
        xs = f(xs, c, n);
        g = __gcd(abs(xs - y), n);
    } while (g == 1);
}
return g;
}
```

6.7 Range Sieve

```
vector<bool> range_sieve(ll l, ll r) {
  ll n = sqrt(r);
  vector<bool> is_prime(n + 1, true);
  vector<ll> prime;
  is_prime[0] = is_prime[1] = false;
  prime.push back(2);
  for (11 i = 4; i \le n; i += 2) {
    is_prime[i] = false;
  for (11 i = 3; i \le n; i += 2) {
    if (is_prime[i]) {
      prime.push back(i);
      for (11 \ j = i * i; \ j <= n; \ j += i)
        is prime[j] = false;
    }
  vector<bool> result(r - 1 + 1, true);
  for (ll i : prime) {
    for (11 j = max(i * i, (1 + i - 1) / i * i); j <= r;
        i += i) {
      result[j - l] = false;
  if (1 == 1) {
    result[0] = false;
  return result;
```

6.8 Segmented Sieve

```
vector<1l> segmented_sieve(ll n) {
  const ll S = 10000;
  ll nsqrt = sqrt(n);
```

```
vector<char> is_prime(nsqrt + 1, true);
vector<ll> prime;
is prime[0] = is prime[1] = false;
prime.push_back(2);
for (11 i = 4; i \le nsqrt; i += 2) {
  is prime[i] = false;
for (11 i = 3; i <= nsgrt; i += 2) {
  if (is_prime[i]) {
    prime.push back(i);
    for (11 j = i * i; j \le nsqrt; j += i) {
      is_prime[j] = false;
 }
}
vector<ll> result;
vector<char> block(S);
for (11 k = 0; k * S <= n; k++) {
  fill(block.begin(), block.end(), true);
  for (ll p : prime) {
    for (ll j = max((k * S + p - 1) / p, p) * p - k *
       S; j < S; j += p) {
      block[j] = false;
  if (k == 0) {
    block[0] = block[1] = false;
  for (ll i = 0; i < S && k * S + i <= n; i++) {
    if (block[i]) {
      result.push back(k * S + i);
return result;
```

6.9 Tonelli Shanks

```
1l legendre(ll a, ll p) {
    return bin_pow_mod(a, (p - 1) / 2, p);
}
1l tonelli_shanks(ll n, ll p) {
    if (legendre(n, p) == p - 1) {
        return -1;
    }
    if (p % 4 == 3) {
        return bin_pow_mod(n, (p + 1) / 4, p);
    }
    ll Q = p - 1, S = 0;
```

```
while (Q % 2 == 0) {
 Q /= 2;
 S++;
11 z = 2;
for (; z < p; z++) {
 if (legendre(z, p) == p - 1) {
   break;
Q, p), R = bin_pow_mod(n, (Q + 1) / 2, p);
while (t % p != 1) {
 if (t % p == 0) {
   return 0;
 11 i = 1, t2 = t * t % p;
 for (; i < M; i++) {
  if (t2 % p == 1) {
     break;
   t2 = t2 * t2 % p;
 ll b = bin_pow_mod(c, bin_pow_mod(2, M - i - 1, p),
    p);
 M = i;
 c = b * b % p;
 t = t * c % p;
 R = R * b % p;
return R;
```

7 Strings

7.1 Hashing

```
11 compute_hash(string const& s) {
  const 11 p = 31, m = 1e9 + 9;
  11 hash_value = 0, p_pow = 1;
  for (char c : s) {
    hash_value = (hash_value + (c - 'a' + 1) * p_pow) %
          m;
    p_pow = (p_pow * p) % m;
  }
  return hash_value;
}
```

7.2 Knuth Morris Pratt

7.3 Rabin Karp

```
vector<ll> rabin karp(string const& s, string const& t)
  const 11 p = 31, m = 1e9 + 9;
  11 S = s.size(), T = t.size();
  vector<ll> p pow(max(S, T));
  p pow[0] = 1;
  for (ll i = 1; i < (ll) p_pow.size(); i++) {</pre>
    p_pow[i] = (p_pow[i-1] * p) % m;
  vector<ll> h(T + 1, 0);
  for (11 i = 0; i < T; i++) {
    h[i + 1] = (h[i] + (t[i] - 'a' + 1) * p_pow[i]) % m;
  11 h_s = 0;
  for (ll i = 0; i < S; i++) {
    h_s = (h_s + (s[i] - 'a' + 1) * p_pow[i]) % m;
  vector<ll> occurences;
  for (11 i = 0; i + S - 1 < T; i++) {
   11 \text{ cur } h = (h[i + S] + m - h[i]) \% m;
    if (cur_h == h_s * p_pow[i] % m) {
      occurences.push_back(i);
    }
  return occurences;
```

7.4 Suffix Array

```
vector<ll> sort cyclic shifts(string const& s) {
  ll n = s.size();
  const 11 alphabet = 256;
  vector<ll> p(n), c(n), cnt(max(alphabet, n), 0);
  for (ll i = 0; i < n; i++) {
    cnt[s[i]]++;
  for (ll i = 1; i < alphabet; i++) {</pre>
    cnt[i] += cnt[i - 1];
  for (ll i = 0; i < n; i++) {
    p[--cnt[s[i]]] = i;
  }
  c[p[0]] = 0;
  ll classes = 1;
  for (ll i = 1; i < n; i++) {
    if (s[p[i]]] != s[p[i - 1]]) {
      classes++;
    c[p[i]] = classes - 1;
  vector<ll> pn(n), cn(n);
  for (11 h = 0; (1 << h) < n; ++h)
    for (ll i = 0; i < n; i++) {</pre>
      pn[i] = p[i] - (1 << h);
      if (pn[i] < 0) {
        pn[i] += n;
    fill(cnt.begin(), cnt.begin() + classes, 0);
    for (ll i = 0; i < n; i++) {
      cnt[c[pn[i]]]++;
    for (ll i = 1; i < classes; i++) {</pre>
      cnt[i] += cnt[i - 1];
    for (ll i = n-1; i >= 0; i--) {
      p[--cnt[c[pn[i]]]] = pn[i];
    cn[p[0]] = 0;
```

```
classes = 1;
    for (ll i = 1; i < n; i++) {</pre>
      pair < ll, ll > cur = {c[p[i]], c[(p[i] + (1 << h)) %}
          n]};
      pair < 11, 11 > prev = \{c[p[i-1]], c[(p[i-1] + (1
           << h)) % n]};
      if (cur != prev) {
        ++classes;
      cn[p[i]] = classes - 1;
    c.swap(cn);
  return p;
vector<ll> build_suff_arr(string s) {
  s += (char) 0;
  vector<ll> sorted shifts = sort cyclic shifts(s);
  sorted_shifts.erase(sorted_shifts.begin());
  return sorted shifts;
```

7.5 Z Function

```
vector<ll> z_function(string s) {
    ll n = (ll) s.length();
    vector<ll> z(n);
    for (ll i = 1, l = 0, r = 0; i < n; ++i) {
        if (i <= r) {
            z[i] = min (r - i + 1, z[i - 1]);
        }
        while (i + z[i] < n && s[z[i]] == s[i + z[i]]) {
            ++z[i];
        }
        if (i + z[i] - 1 > r) {
            l = i, r = i + z[i] - 1;
        }
    }
    return z;
}
```

<u> </u>		
f(n) = O(g(n))	iff \exists positive c, n_0 such that $0 \le f(n) \le cg(n) \ \forall n \ge n_0$.	$\sum_{i=1}^{n} i = \frac{n(n+1)}{2}, \sum_{i=1}^{n} i^2 = \frac{n(n+1)(2n+1)}{6}, \sum_{i=1}^{n} i^3 = \frac{n^2(n+1)^2}{4}.$
$f(n) = \Omega(g(n))$	iff \exists positive c, n_0 such that $f(n) \ge cg(n) \ge 0 \ \forall n \ge n_0$.	i=1 $i=1$ $i=1$ In general:
$f(n) = \Theta(g(n))$	iff $f(n) = O(g(n))$ and $f(n) = \Omega(g(n))$.	$\sum_{i=1}^{n} i^{m} = \frac{1}{m+1} \left[(n+1)^{m+1} - 1 - \sum_{i=1}^{n} \left((i+1)^{m+1} - i^{m+1} - (m+1)i^{m} \right) \right]$
f(n) = o(g(n))	iff $\lim_{n\to\infty} f(n)/g(n) = 0$.	$\sum_{i=1}^{n-1} i^m = \frac{1}{m+1} \sum_{k=0}^m \binom{m+1}{k} B_k n^{m+1-k}.$
$\lim_{n \to \infty} a_n = a$	iff $\forall \epsilon > 0$, $\exists n_0$ such that $ a_n - a < \epsilon$, $\forall n \ge n_0$.	Geometric series:
$\sup S$	least $b \in \mathbb{R}$ such that $b \ge s$, $\forall s \in S$.	$\sum_{i=0}^{n} c^{i} = \frac{c^{n+1} - 1}{c - 1}, c \neq 1, \sum_{i=0}^{\infty} c^{i} = \frac{1}{1 - c}, \sum_{i=1}^{\infty} c^{i} = \frac{c}{1 - c}, c < 1,$
$\inf S$	greatest $b \in \mathbb{R}$ such that $b \le s$, $\forall s \in S$.	$\sum_{i=0}^{n} ic^{i} = \frac{nc^{n+2} - (n+1)c^{n+1} + c}{(c-1)^{2}}, c \neq 1, \sum_{i=0}^{\infty} ic^{i} = \frac{c}{(1-c)^{2}}, c < 1.$
$ \liminf_{n \to \infty} a_n $	$\lim_{n\to\infty}\inf\{a_i\mid i\geq n, i\in\mathbb{N}\}.$	Harmonic series: $n = n + 1 =$
$\limsup_{n \to \infty} a_n$	$\lim_{n \to \infty} \sup \{ a_i \mid i \ge n, i \in \mathbb{N} \}.$	$H_n = \sum_{i=1}^n \frac{1}{i}, \qquad \sum_{i=1}^n iH_i = \frac{n(n+1)}{2}H_n - \frac{n(n-1)}{4}.$
$\binom{n}{k}$	Combinations: Size k subsets of a size n set.	$\sum_{i=1}^{n} H_i = (n+1)H_n - n, \sum_{i=1}^{n} {i \choose m} H_i = {n+1 \choose m+1} \left(H_{n+1} - \frac{1}{m+1} \right).$
$\begin{bmatrix} n \\ k \end{bmatrix}$	Stirling numbers (1st kind): Arrangements of an <i>n</i> element set into <i>k</i> cycles.	1. $\binom{n}{k} = \frac{n!}{(n-k)!k!}$, 2. $\sum_{k=0}^{n} \binom{n}{k} = 2^{n}$, 3. $\binom{n}{k} = \binom{n}{n-k}$,
${n \brace k}$	Stirling numbers (2nd kind): Partitions of an n element set into k non-empty sets.	$4. \binom{n}{k} = \frac{n}{k} \binom{n-1}{k-1}, \qquad \qquad 5. \binom{n}{k} = \binom{n-1}{k} + \binom{n-1}{k-1}, \\ 6. \binom{n}{m} \binom{m}{k} = \binom{n}{k} \binom{n-k}{m-k}, \qquad \qquad 7. \sum_{k=0}^{n} \binom{r+k}{k} = \binom{r+n+1}{n},$
$\langle {n \atop k} \rangle$	1st order Eulerian numbers: Permutations $\pi_1\pi_2\pi_n$ on $\{1, 2,, n\}$ with k ascents.	$8. \sum_{k=0}^{n} \binom{k}{m} = \binom{n+1}{m+1}, \qquad 9. \sum_{k=0}^{n} \binom{r}{k} \binom{s}{n-k} = \binom{r+s}{n},$
$\left\langle\!\left\langle {n\atop k}\right\rangle\!\right\rangle$	2nd order Eulerian numbers.	10. $\binom{n}{k} = (-1)^k \binom{k-n-1}{k}$, 11. $\binom{n}{1} = \binom{n}{n} = 1$,
C_n	Catalan Numbers: Binary trees with $n+1$ vertices.	12. $\binom{n}{2} = 2^{n-1} - 1$, 13. $\binom{n}{k} = k \binom{n-1}{k} + \binom{n-1}{k-1}$,
14. $\begin{bmatrix} n \\ 1 \end{bmatrix} = (n-1)^n$	15. $\begin{bmatrix} n \\ 2 \end{bmatrix} = (n - 1)^n$	$16. \begin{bmatrix} n \\ n \end{bmatrix} = 1, \qquad \qquad 17. \begin{bmatrix} n \\ k \end{bmatrix} \ge \begin{Bmatrix} n \\ k \end{Bmatrix},$
$18. \begin{bmatrix} n \\ k \end{bmatrix} = (n-1)$	$\binom{n-1}{k} + \binom{n-1}{k-1}, 19. \ \binom{n}{n-1}$	
$22. \ \left\langle \begin{matrix} n \\ 0 \end{matrix} \right\rangle = \left\langle \begin{matrix} n \\ n \end{matrix} \right\rangle$	$\binom{n}{-1} = 1,$ 23. $\binom{n}{k} = \binom{n}{k}$	$\binom{n}{n-1-k}$, $24. \left\langle \binom{n}{k} \right\rangle = (k+1) \left\langle \binom{n-1}{k} \right\rangle + (n-k) \left\langle \binom{n-1}{k-1} \right\rangle$,
25. $\left\langle {0\atop k}\right\rangle = \left\{ {1\atop 0}\right\}$	if $k = 0$, otherwise 26. $\begin{cases} r \\ 1 \end{cases}$	$\binom{n}{2} = 2^n - n - 1,$ 27. $\binom{n}{2} = 3^n - (n+1)2^n + \binom{n+1}{2},$
28. $x^n = \sum_{k=0}^n \binom{n}{k}$	$\left. \left\langle {x+k \atop n} \right\rangle, \qquad $ 29. $\left\langle {n \atop m} \right\rangle = \sum_{k=1}^m$	
		32. $\left\langle \left\langle \begin{array}{c} n \\ 0 \end{array} \right\rangle = 1,$ 33. $\left\langle \left\langle \begin{array}{c} n \\ n \end{array} \right\rangle = 0$ for $n \neq 0$,
$34. \; \left\langle \!\! \left\langle \!\! \begin{array}{c} n \\ k \end{array} \!\! \right\rangle = (k + 1)^n$	$+1$ $\binom{n-1}{k}$ $+(2n-1-k)$ $\binom{n-1}{k}$	
$36. \left\{ \begin{array}{c} x \\ x-n \end{array} \right\} = \sum_{k}^{n} \left\{ \begin{array}{c} x \\ x \end{array} \right\}$	$\sum_{k=0}^{n} \left\langle \!\! \left\langle n \atop k \right\rangle \!\! \right\rangle \!\! \binom{x+n-1-k}{2n},$	37. $\binom{n+1}{m+1} = \sum_{k} \binom{n}{k} \binom{k}{m} = \sum_{k=0}^{n} \binom{k}{m} (m+1)^{n-k},$

The Chinese remainder theorem: There exists a number C such that:

$$C \equiv r_1 \bmod m_1$$

: : :

$$C \equiv r_n \bmod m_n$$

if m_i and m_j are relatively prime for $i \neq j$. Euler's function: $\phi(x)$ is the number of positive integers less than x relatively prime to x. If $\prod_{i=1}^{n} p_i^{e_i}$ is the prime factorization of x then

$$\phi(x) = \prod_{i=1}^{n} p_i^{e_i - 1} (p_i - 1).$$

Euler's theorem: If a and b are relatively prime then

$$1 \equiv a^{\phi(b)} \bmod b$$
.

Fermat's theorem:

$$1 \equiv a^{p-1} \bmod p.$$

The Euclidean algorithm: if a > b are integers then

$$gcd(a, b) = gcd(a \mod b, b).$$

If $\prod_{i=1}^{n} p_i^{e_i}$ is the prime factorization of x

$$S(x) = \sum_{d|x} d = \prod_{i=1}^{n} \frac{p_i^{e_i+1} - 1}{p_i - 1}.$$

Perfect Numbers: x is an even perfect number iff $x = 2^{n-1}(2^n - 1)$ and $2^n - 1$ is prime. Wilson's theorem: n is a prime iff $(n-1)! \equiv -1 \mod n$.

$$\mu(i) = \begin{cases} (n-1)! = -1 \bmod n. \\ \text{M\"obius inversion:} \\ \mu(i) = \begin{cases} 1 & \text{if } i = 1. \\ 0 & \text{if } i \text{ is not square-free.} \\ (-1)^r & \text{if } i \text{ is the product of} \\ r & \text{distinct primes.} \end{cases}$$
 If

 If

$$G(a) = \sum_{d|a} F(d),$$

$$F(a) = \sum_{d|a} \mu(d) G\left(\frac{a}{d}\right).$$

Prime numbers:

$$p_n = n \ln n + n \ln \ln n - n + n \frac{\ln \ln n}{\ln n}$$

$$+O\left(\frac{n}{\ln n}\right),$$

$$\pi(n) = \frac{n}{\ln n} + \frac{n}{(\ln n)^2} + \frac{2!n}{(\ln n)^3} + O\left(\frac{n}{(\ln n)^4}\right).$$

-	`	0		
				ns

Loop An edge connecting a vertex to itself.

Directed Each edge has a direction. SimpleGraph with no loops or multi-edges.

WalkA sequence $v_0e_1v_1\ldots e_\ell v_\ell$. TrailA walk with distinct edges. Pathtrail with distinct

vertices.

ConnectedA graph where there exists a path between any two

vertices.

connected ComponentA maximal subgraph.

TreeA connected acyclic graph. Free tree A tree with no root. DAGDirected acyclic graph. EulerianGraph with a trail visiting each edge exactly once.

Hamiltonian Graph with a cycle visiting each vertex exactly once.

CutA set of edges whose removal increases the number of components.

Cut-setA minimal cut. Cut edge A size 1 cut.

k-Connected A graph connected with the removal of any k-1vertices.

k-Tough $\forall S \subseteq V, S \neq \emptyset$ we have $k \cdot c(G - S) \le |S|$.

A graph where all vertices k-Regular have degree k.

k-Factor Α k-regular spanning subgraph.

Matching A set of edges, no two of which are adjacent.

CliqueA set of vertices, all of which are adjacent.

Ind. set A set of vertices, none of which are adjacent.

Vertex cover A set of vertices which cover all edges.

Planar graph A graph which can be embeded in the plane.

Plane graph An embedding of a planar

$$\sum_{v \in V} \deg(v) = 2m.$$

If G is planar then n - m + f = 2, so

$$f \le 2n - 4, \quad m \le 3n - 6.$$

Any planar graph has a vertex with degree ≤ 5 .

Notation:

E(G)Edge set Vertex set V(G)

c(G)Number of components

G[S]Induced subgraph deg(v)Degree of v

Maximum degree $\Delta(G)$ $\delta(G)$ Minimum degree

 $\chi(G)$ Chromatic number $\chi_E(G)$ Edge chromatic number

 G^c Complement graph K_n Complete graph

 K_{n_1,n_2} Complete bipartite graph

Ramsev number

Geometry

Projective coordinates: (x, y, z), not all x, y and z zero.

$$(x, y, z) = (cx, cy, cz) \quad \forall c \neq 0.$$

Cartesian Projective (x, y)(x, y, 1)y = mx + b(m, -1, b)x = c(1,0,-c)

Distance formula, L_p and L_{∞}

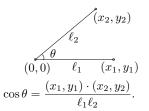
$$\sqrt{(x_1 - x_0)^2 + (y_1 - y_0)^2},$$
$$\left[|x_1 - x_0|^p + |y_1 - y_0|^p \right]^{1/p},$$

 $\lim \left[|x_1 - x_0|^p + |y_1 - y_0|^p \right]^{1/p}.$

Area of triangle $(x_0, y_0), (x_1, y_1)$ and (x_2, y_2) :

$$\frac{1}{2} \operatorname{abs} \begin{vmatrix} x_1 - x_0 & y_1 - y_0 \\ x_2 - x_0 & y_2 - y_0 \end{vmatrix}.$$

Angle formed by three points:



Line through two points (x_0, y_0) and (x_1, y_1) :

$$\begin{vmatrix} x & y & 1 \\ x_0 & y_0 & 1 \\ x_1 & y_1 & 1 \end{vmatrix} = 0.$$

Area of circle, volume of sphere:

$$A = \pi r^2, \qquad V = \frac{4}{3}\pi r^3.$$

If I have seen farther than others, it is because I have stood on the shoulders of giants.

- Issac Newton

Taylor's series:

$$f(x) = f(a) + (x - a)f'(a) + \frac{(x - a)^2}{2}f''(a) + \dots = \sum_{i=0}^{\infty} \frac{(x - a)^i}{i!}f^{(i)}(a).$$

Expansions:

Expansions:
$$\frac{1}{1-x} = 1 + x + x^2 + x^3 + x^4 + \cdots = \sum_{i=0}^{\infty} x^i,$$

$$\frac{1}{1-cx} = 1 + x + x^2 + x^3 + x^4 + \cdots = \sum_{i=0}^{\infty} c^i x^i,$$

$$\frac{1}{1-cx} = 1 + cx + c^2 x^2 + c^3 x^3 + \cdots = \sum_{i=0}^{\infty} c^i x^i,$$

$$\frac{1}{1-x^n} = 1 + x^n + x^{2n} + x^{3n} + \cdots = \sum_{i=0}^{\infty} x^{ni},$$

$$\frac{x}{(1-x)^2} = x + 2x^2 + 3x^3 + 4x^4 + \cdots = \sum_{i=0}^{\infty} i^n x^i,$$

$$x^k \frac{dx^n}{dx^n} \left(\frac{1}{1-x}\right) = x + 2^n x^2 + 3^n x^3 + 4^n x^4 + \cdots = \sum_{i=0}^{\infty} i^n x^i,$$

$$e^x = 1 + x + \frac{1}{2}x^2 + \frac{1}{6}x^3 + \cdots = \sum_{i=0}^{\infty} (-1)^{i+1} \frac{x^i}{i},$$

$$\ln(1+x) = x - \frac{1}{2}x^2 + \frac{1}{3}x^3 - \frac{1}{4}x^4 - \cdots = \sum_{i=0}^{\infty} (-1)^{i+1} \frac{x^i}{i},$$

$$\sin x = x - \frac{1}{3}x^3 + \frac{1}{13}x^5 - \frac{1}{71}x^7 + \cdots = \sum_{i=0}^{\infty} (-1)^i \frac{x^{2i+1}}{(2i+1)!},$$

$$\cos x = 1 - \frac{1}{2}x^2 + \frac{1}{4}x^4 - \frac{1}{6!}x^6 + \cdots = \sum_{i=0}^{\infty} (-1)^i \frac{x^{2i+1}}{(2i+1)!},$$

$$\tan^{-1} x = x - \frac{1}{3}x^3 + \frac{1}{5}x^5 - \frac{1}{7}x^7 + \cdots = \sum_{i=0}^{\infty} (-1)^i \frac{x^{2i+1}}{(2i+1)!},$$

$$(1+x)^n = 1 + nx + \frac{n(n-1)}{2}x^2 + \cdots = \sum_{i=0}^{\infty} (-1)^i \frac{x^{2i+1}}{(2i+1)!},$$

$$\frac{1}{(1-x)^{n+1}} = 1 + (n+1)x + \binom{n+2}{2}x^2 + \cdots = \sum_{i=0}^{\infty} \binom{i}{i}x^i,$$

$$\frac{x}{e^x - 1} = 1 - \frac{1}{2}x + \frac{1}{12}x^2 - \frac{1}{720}x^4 + \cdots = \sum_{i=0}^{\infty} \binom{i+n}{i}x^i,$$

$$\frac{1}{\sqrt{1-4x}} = 1 + x + 2x^2 + 5x^3 + \cdots = \sum_{i=0}^{\infty} \frac{1}{i+1} \binom{2i}{i}x^i,$$

$$\frac{1}{\sqrt{1-4x}} = 1 + x + 2x^2 + 6x^3 + \cdots = \sum_{i=0}^{\infty} \frac{1}{i+1} \binom{2i}{i}x^i,$$

$$\frac{1}{\sqrt{1-4x}} = 1 + (2+n)x + \binom{4+n}{2}x^2 + \cdots = \sum_{i=0}^{\infty} \frac{1}{i+1} \binom{2i}{i}x^i,$$

$$\frac{1}{1-x} \ln \frac{1}{1-x} = x + \frac{3}{2}x^2 + \frac{1}{16}x^3 + \frac{25}{12}x^4 + \cdots = \sum_{i=0}^{\infty} \frac{1}{i+1} \binom{2i}{i}x^i,$$

$$\frac{1}{1-x} \ln \frac{1}{1-x} = x + \frac{3}{2}x^2 + \frac{1}{16}x^3 + \frac{25}{12}x^4 + \cdots = \sum_{i=0}^{\infty} \frac{1}{i+1} \binom{2i}{i}x^i,$$

$$\frac{1}{1-x} \ln \frac{1}{1-x} = x + \frac{3}{2}x^2 + \frac{1}{16}x^3 + \frac{25}{12}x^4 + \cdots = \sum_{i=0}^{\infty} \frac{1}{i},$$

$$\frac{x}{1-x} = x^2 + x^2 + 2x^3 + 3x^4 + \cdots = \sum_{i=0}^{\infty} F_{i}x^i.$$

$$\frac{x}{1-x} = x^2 + x^2 + 2x^3 + 3x^4 + \cdots = \sum_{i=0}^{\infty} F_{i}x^i.$$

Ordinary power series:

$$A(x) = \sum_{i=0}^{\infty} a_i x^i.$$

Exponential power series:

$$A(x) = \sum_{i=0}^{\infty} a_i \frac{x^i}{i!}.$$

Dirichlet power series:

$$A(x) = \sum_{i=1}^{\infty} \frac{a_i}{i^x}.$$

Binomial theorem

$$(x+y)^n = \sum_{k=0}^n \binom{n}{k} x^{n-k} y^k.$$

$$x^{n} - y^{n} = (x - y) \sum_{k=0}^{n-1} x^{n-1-k} y^{k}.$$

For ordinary power se

$$\alpha A(x) + \beta B(x) = \sum_{i=0}^{\infty} (\alpha a_i + \beta b_i) x^i,$$

$$x^k A(x) = \sum_{i=k}^{\infty} a_{i-k} x^i,$$

$$\frac{A(x) - \sum_{i=0}^{k-1} a_i x^i}{x^k} = \sum_{i=0}^{\infty} a_{i+k} x^i,$$

$$A(cx) = \sum_{i=0}^{\infty} c^i a_i x^i,$$

$$A'(x) = \sum_{i=0}^{\infty} (i+1) a_{i+1} x^i,$$

$$xA'(x) = \sum_{i=1}^{\infty} ia_i x^i,$$

$$\int A(x) dx = \sum_{i=1}^{\infty} \frac{a_{i-1}}{i} x^{i},$$

$$\frac{A(x) + A(-x)}{2} = \sum_{i=0}^{\infty} a_{2i} x^{2i},$$

$$\frac{A(x) - A(-x)}{2} = \sum_{i=0}^{\infty} a_{2i+1} x^{2i+1}.$$

Summation: If $b_i = \sum_{i=0}^i a_i$ then

$$B(x) = \frac{1}{1-x}A(x).$$

Convolution:

$$A(x)B(x) = \sum_{i=0}^{\infty} \left(\sum_{j=0}^{i} a_j b_{i-j}\right) x^i.$$

God made the natural numbers; all the rest is the work of man. Leopold Kronecker