

### Understanding MySQL Group Replication

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## Program Agenda



#### Program Agenda

- Background
- <sup>2</sup> Use cases
- 3 Deployment modes
- 4 Features
- 5 Performance
- 6 Architecture
- 7 Conclusion



# Background

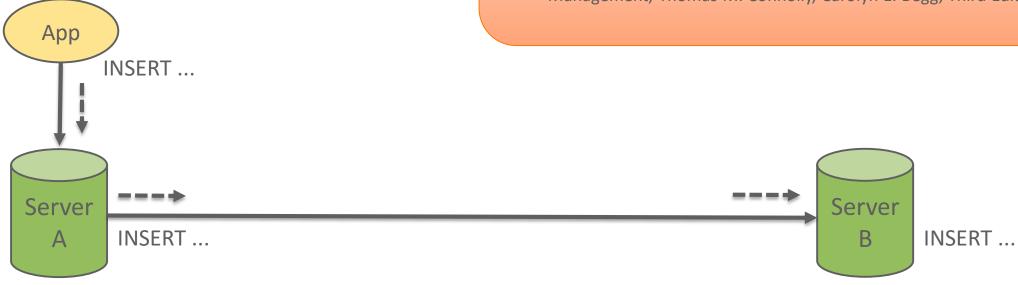


#### **Database Replication**

#### Replication

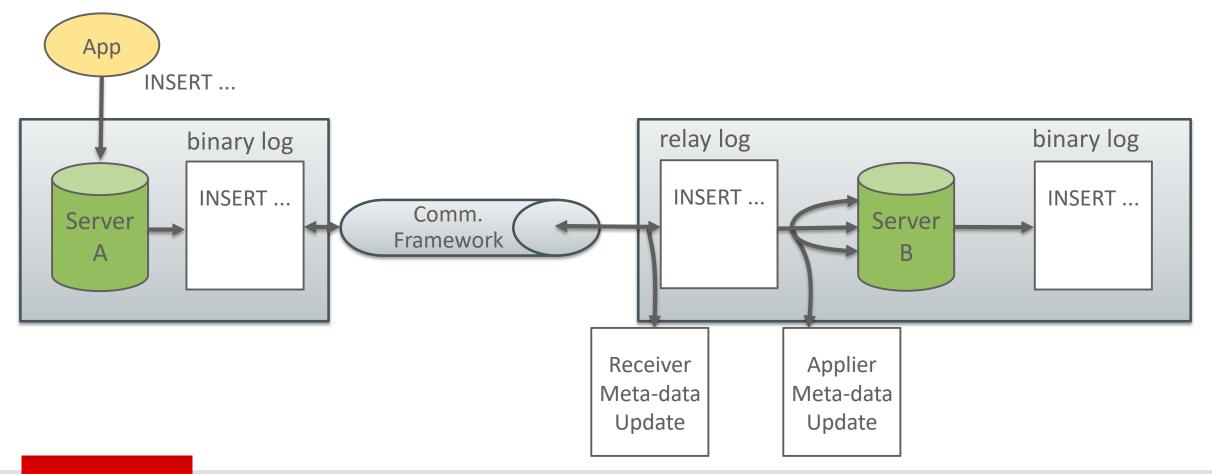
"The process of generating and reproducing multiple copies of data at one or more sites.",

Database Systems: A Practical Approach to Design, Implementation, and Management, Thomas M. Connolly, Carolyn E. Begg, Third Edition, 2002.





### MySQL Database Replication: Overview



#### MySQL Database Replication: Some Notes

#### **Coordination Between Servers**



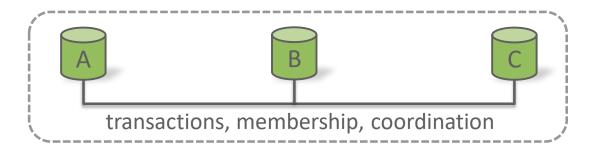
Since 3.23

asynchronous (native)



Since 5.5

semi-synchronous (plugin)



Since 5.7.17

And now in MySQL 8 as of 8.0.1

group replication (plugin)



#### MySQL Group Replication

#### • What is MySQL Group Replication?

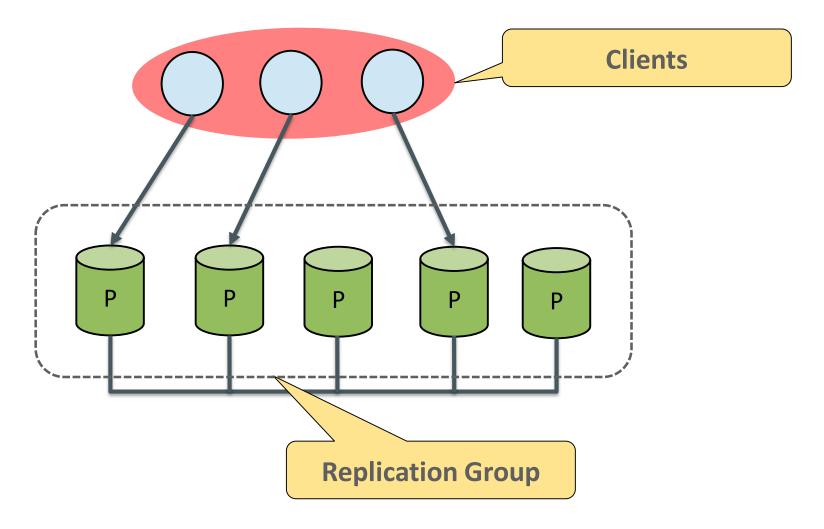
"Single/Multi-primary update everywhere replication plugin for MySQL with built-in automatic distributed recovery, conflict detection and group membership."

#### What does the MySQL Group Replication plugin do for the user?

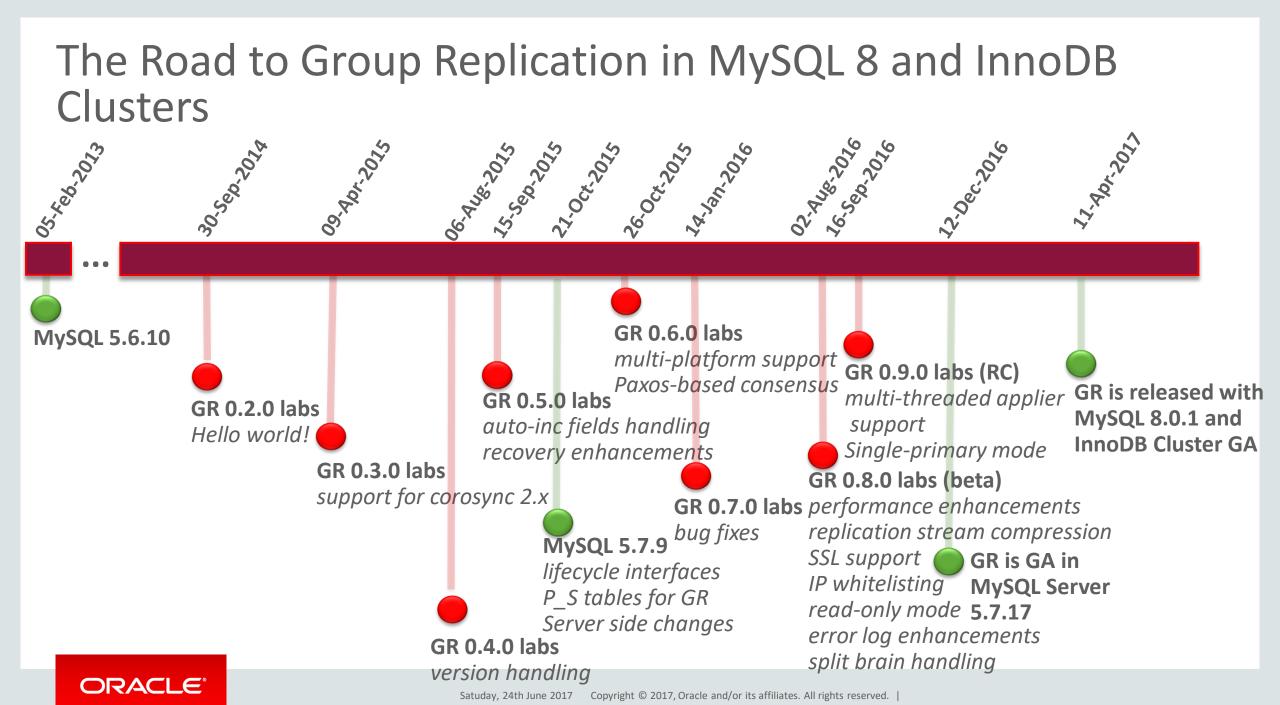
- Removes the need for handling server fail-over.
- Provides fault tolerance.
- Enables update everywhere setups.
- Automates group reconfiguration (handling of crashes, failures, re-connects).
- Provides a highly available replicated database.



### MySQL Group Replication







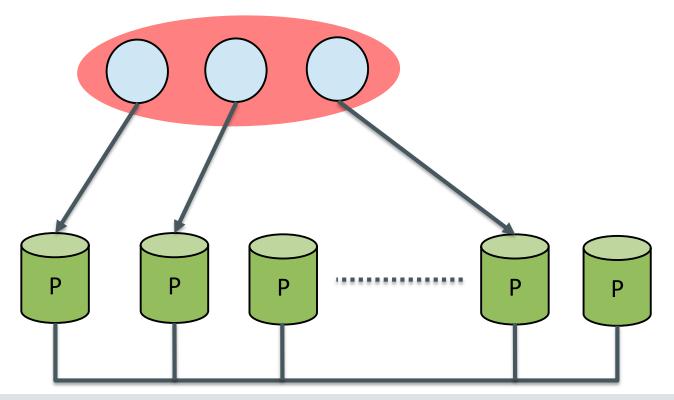
## Use cases



#### **Use Cases**

#### Elastic Replication

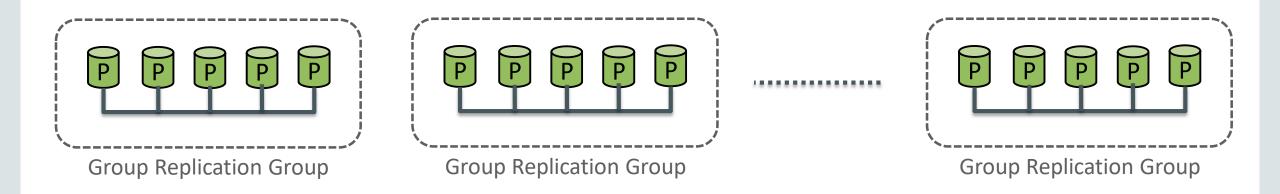
 Environments that require a very fluid replication infrastructure, where the number of servers has to grow or shrink dynamically and with as little pain as possible.



#### **Use Cases**

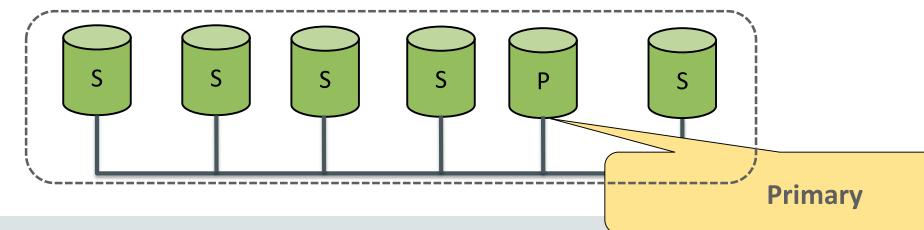
#### Highly Available Shards

Sharding is a popular approach to achieve write scale-out. Users can use MySQL
 Group Replication to implement highly available shards. Each shard can map into a Replication Group.



#### **Use Cases**

- Alternative to Master-Slave replication
- Single-primary mode provides further automation on such setups
  - Automatic PRIMARY/SECONDARY roles assignment
  - Automatic new PRIMARY election on PRIMARY failures
  - Automatic setup of read/write modes on PRIMARY and SECONDARIES
  - Global consistent view of which server is the PRIMARY



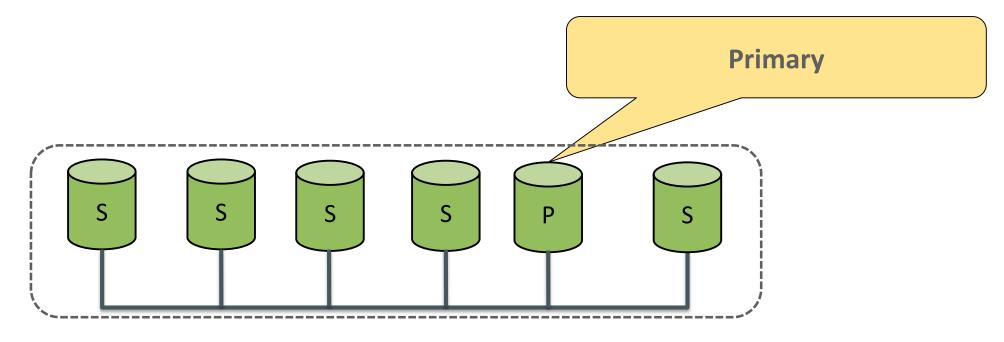
## Deployment modes

3.1 Single-primary

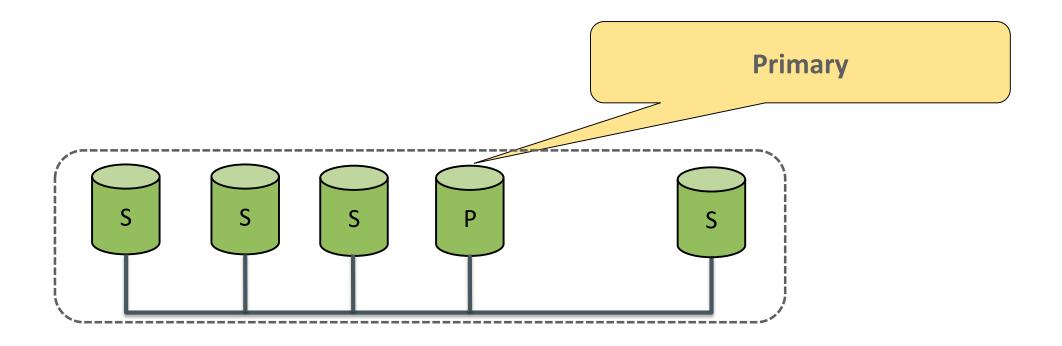


- Configuration mode that makes a single member act as a writeable master (PRIMARY) and the rest of the members act as hot-standbys (SECONDARIES).
  - The group itself coordinates automatically to figure out which is the member that will act as the PRIMARY, through a primary election mechanism.
- Single-primary mode is the default mode
  - Closer to classic asynchronous replication setups, simpler to reason about from the beginning.
  - Avoids some of the limitations of multi-primary mode by default.

- Automatic primary promotion election.
- Secondaries are automatically set to read-only.



• Automatic primary election mechanism.



 The current primary member UUID can be known by executing the following SQL statement.

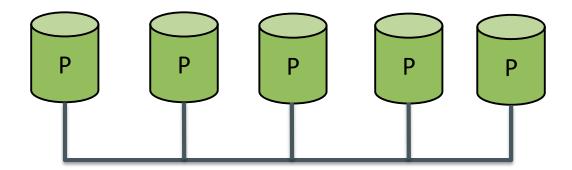
```
mysql> SELECT * FROM performance_schema.global_status WHERE
      VARIABLE_NAME='group_replication_primary_member';
VARIABLE NAME
                                  VARIABLE VALUE
group_replication_primary_member dcd3b36b-79c5-11e6-97b8-00212844d44e
```

## Deployment modes

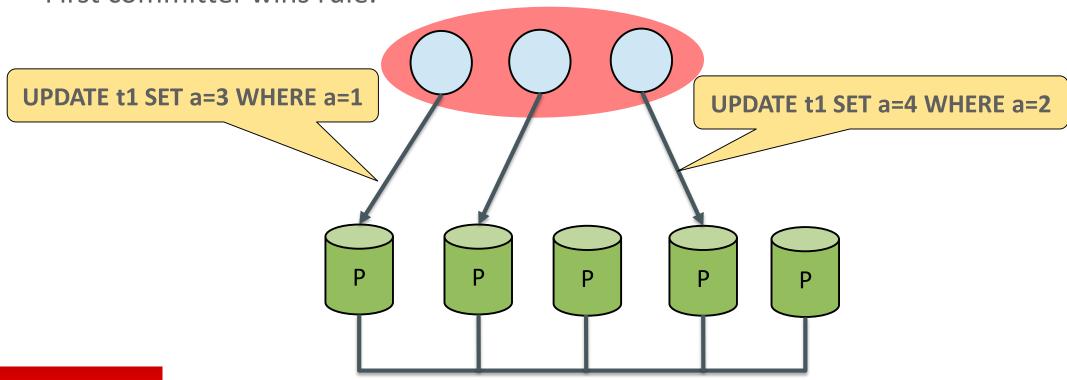
- 3.1 Single-primary
- 3.2 Multi-primary



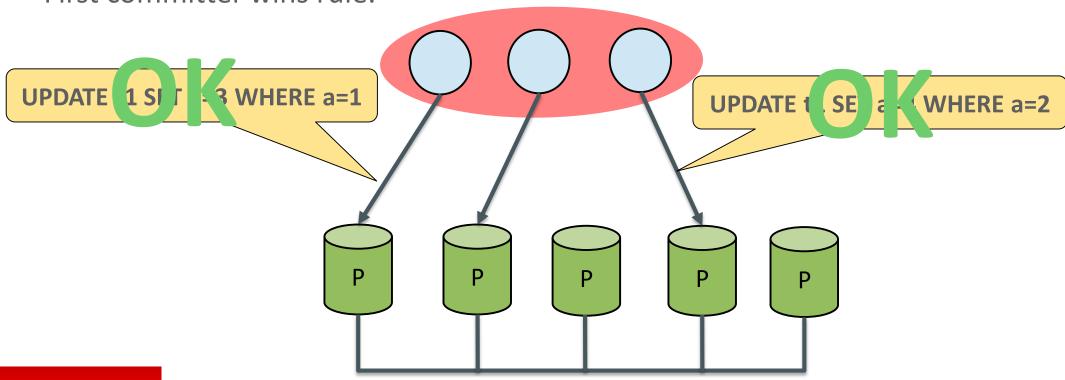
- Configuration mode that makes all members writable
  - Enabled by setting option --group\_replication\_single\_primary\_mode to OFF



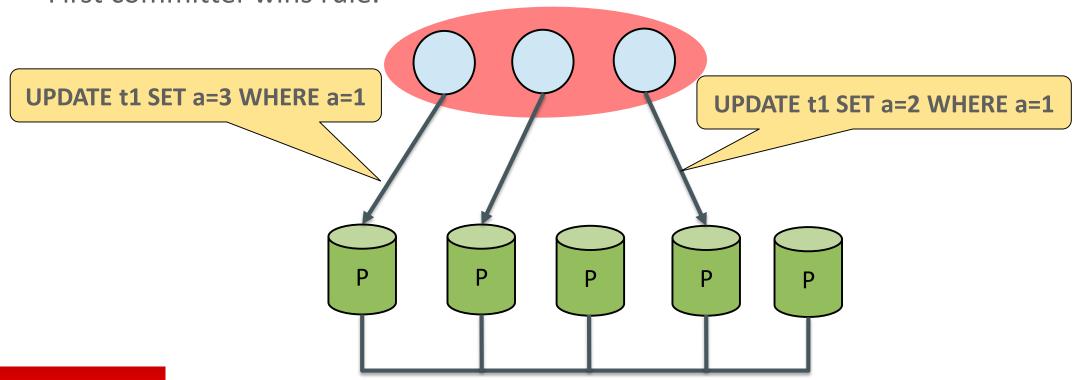
- Any two transactions on different servers can write to the same tuple.
- Conflicts will be detected and dealt with.
  - First committer wins rule.



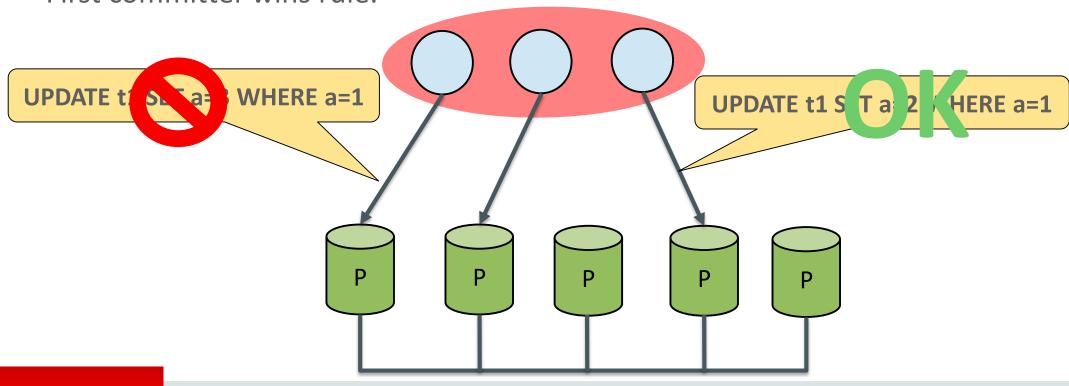
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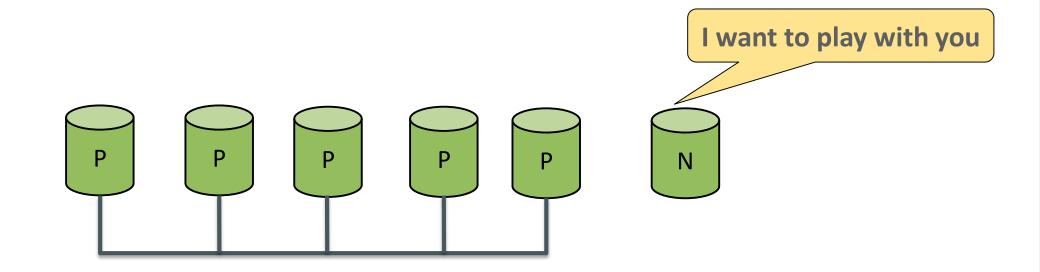


### Features

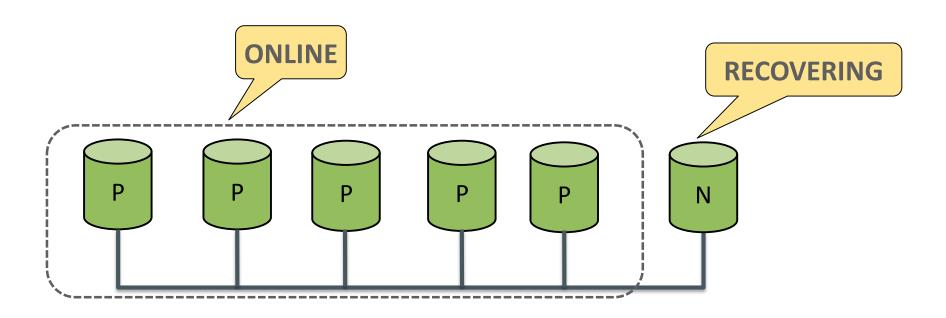
4.1 Automatic distributed server recovery



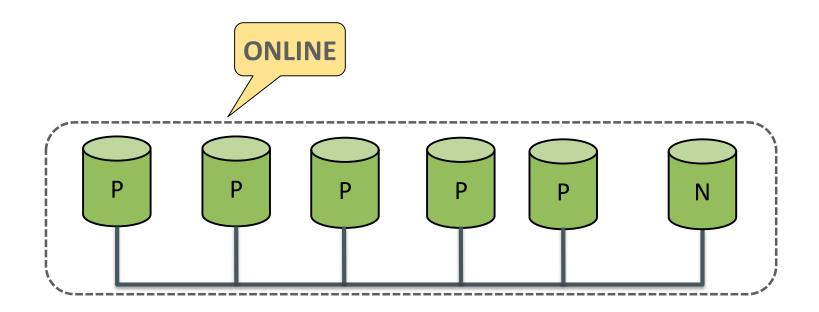
• Server that joins the group will automatically synchronize with the others.



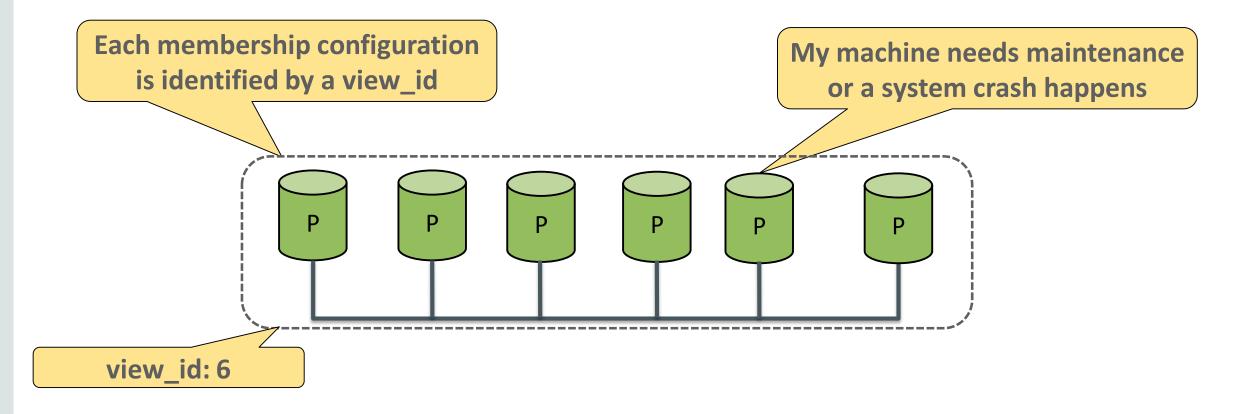
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• Server that joins the group will automatically synchronize with the others.

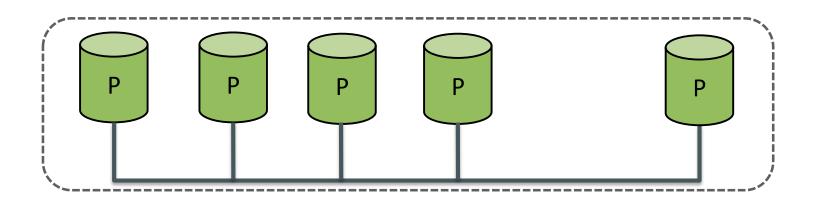


• If a server leaves the group, the others will automatically be informed.



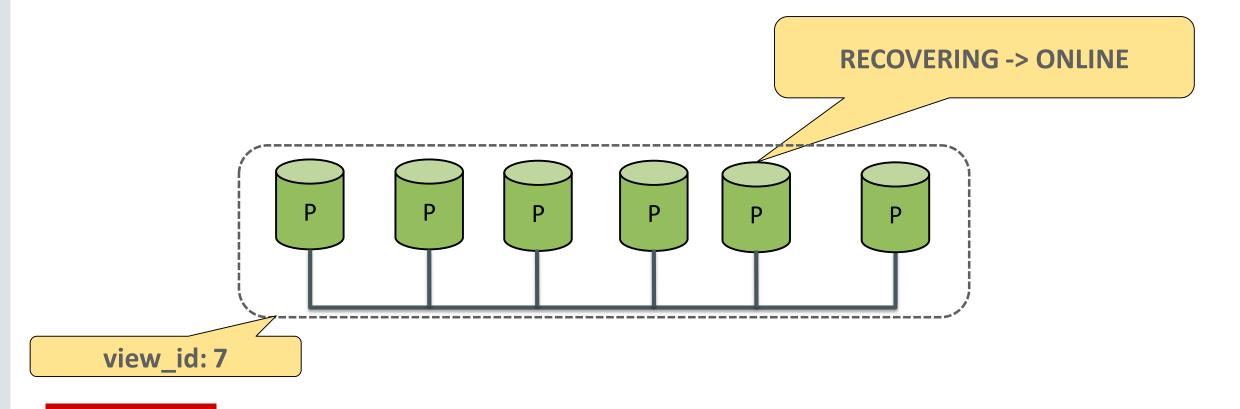


• If a server leaves the group, the others will automatically be informed.





• Server that (re)joins the group will automatically synchronize with the others.



### Features

- 4.1 Automatic distributed server recovery
- 4.2 MySQL Look & Feel



### MySQL Look & Feel!

- MySQL Plugin
  - Regular MySQL Plugin. Nothing new.
- MySQL InnoDB
  - Use InnoDB as normally you would. Nothing new.
  - Transparent optimizations in InnoDB to better support Group Replication.
- MySQL Performance Schema
  - Monitor Group Replication using regular Performance Schema tables. Nothing new.

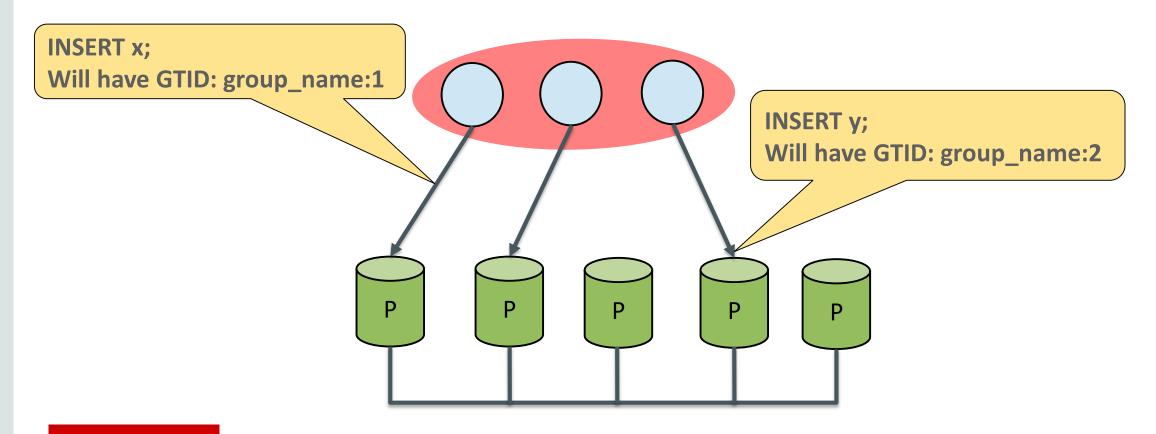
### MySQL Look & Feel!

- Outcome
  - Group Replication is no alien component.
  - Existing MySQL users feel right at home.
  - New MySQL users only have to learn MySQL tech, nothing else.

- 4.1 Automatic distributed server recovery
- 4.2 MySQL Look & Feel
- 4.3 Full GTID support

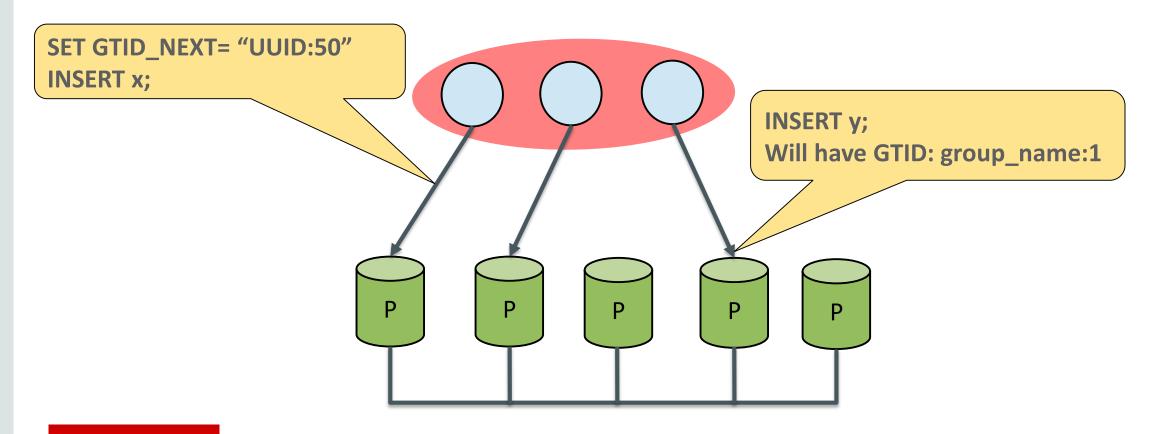


• All group members share the same UUID, the group name.





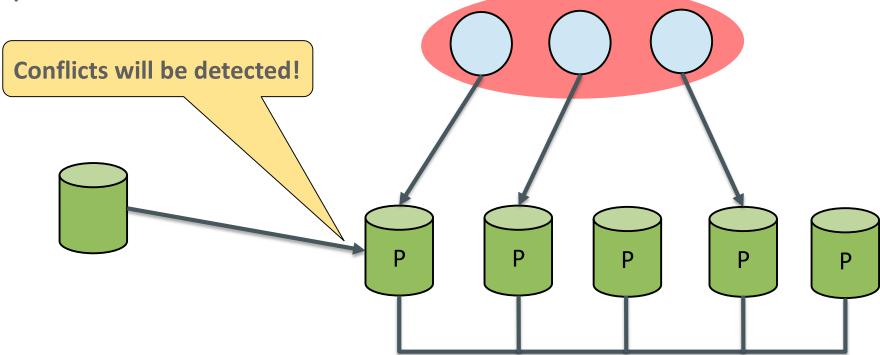
• Users can specify the identifier for the transaction.





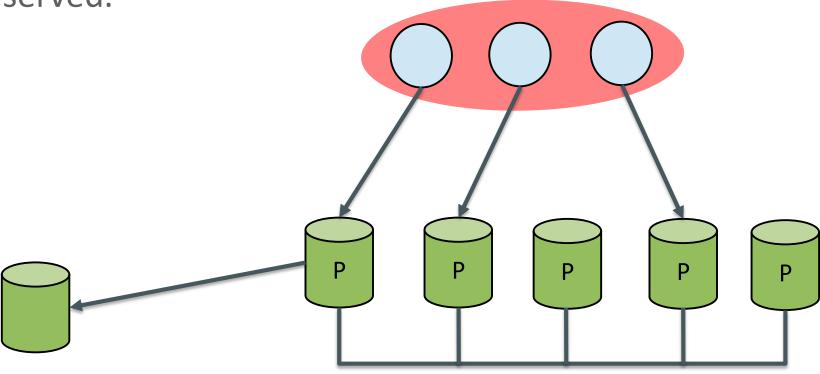
• You can even replicate from a outside server to a group, global identifiers

will be preserved.



• You can also replicate from a group to a outside server, global identifiers

will be preserved.

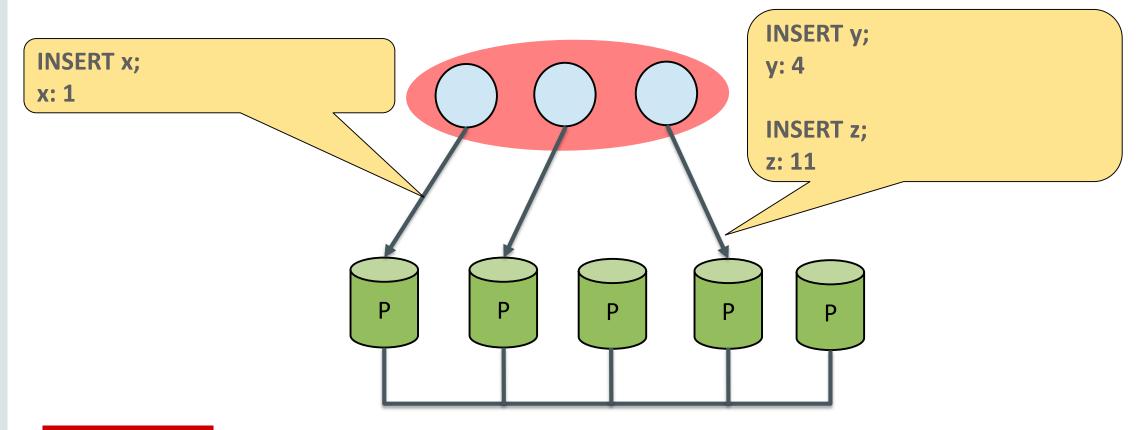


- 4.1 Automatic distributed server recovery
- 4.2 MySQL Look & Feel
- 4.3 Full GTID support
- 4.4 Auto-increment configuration/handling



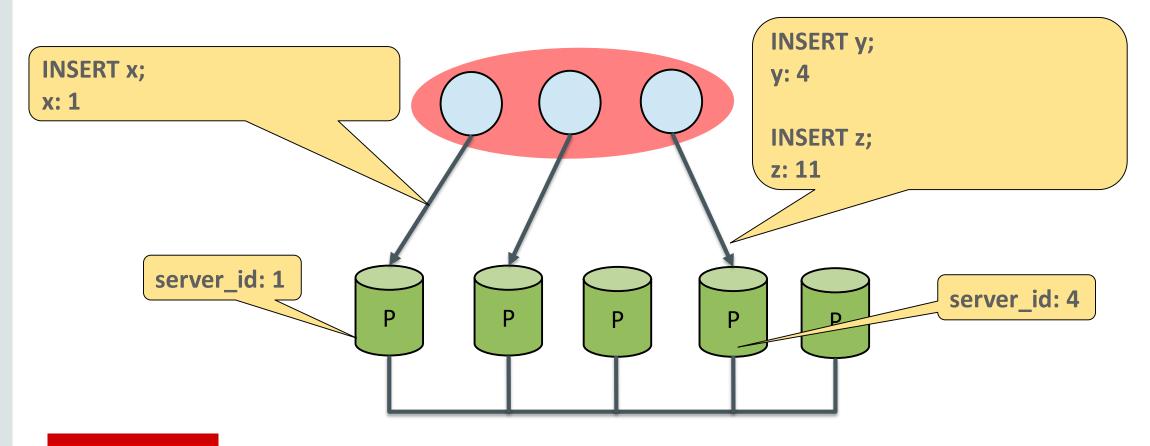
### Auto-increment configuration/handling

• Group is configured to not generate the same auto-increment value on all members.



### Auto-increment configuration/handling

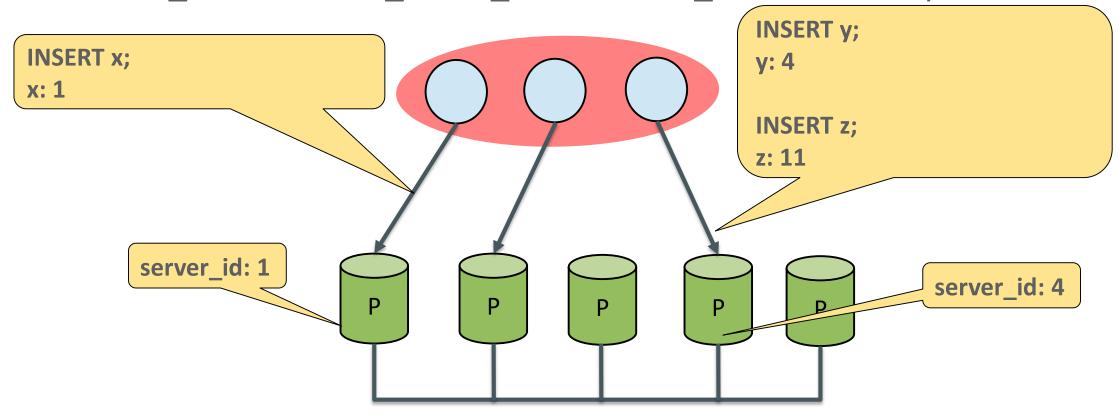
• By default, the offset is provided by server\_id and increment is 7.





#### Auto-increment configuration/handling

 Users can change the increment size to their needs using GROUP\_REPLICATION\_AUTO\_INCREMENT\_INCREMENT option.

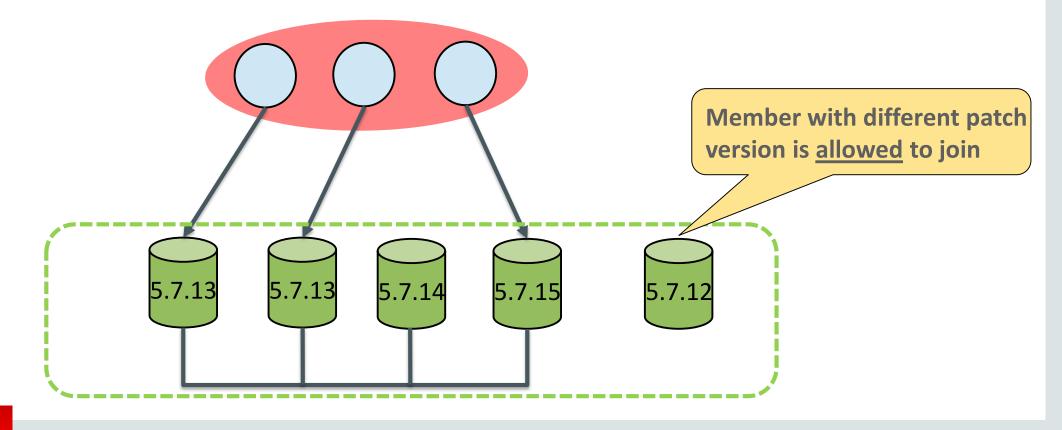


- 4.1 Automatic distributed server recovery
- 4.2 MySQL Look & Feel
- 4.3 Full GTID support
- 4.4 Auto-increment configuration/handling
- 4.5 Plugin version access control



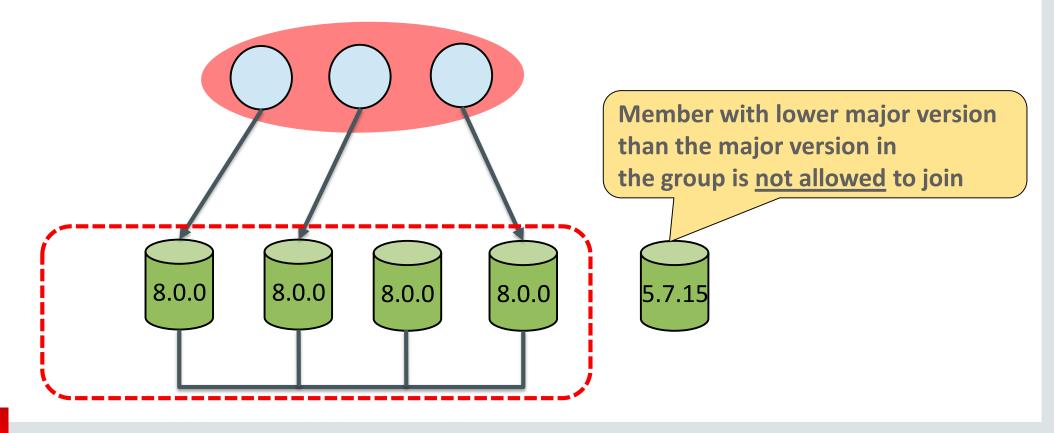
### Plugin Version Access Control

• When joining, versions are crucial when determining if a member is compatible with a group.



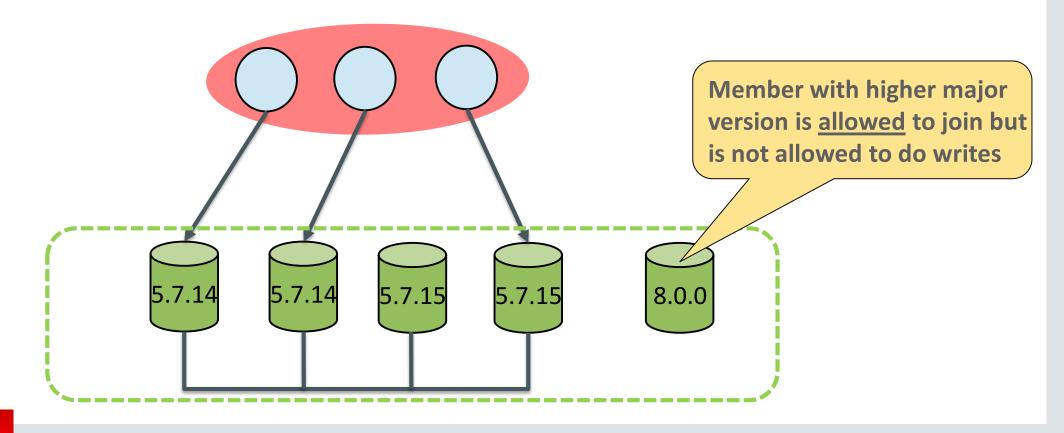
### Plugin Version Access Control

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- 4.1 Automatic distributed server recovery
- 4.2 MySQL Look & Feel
- 4.3 Full GTID support
- 4.4 Auto-increment configuration/handling
- 4.5 Plugin version access control
- 4.6 Built-in communication engine



### **Built-in Communication Engine**

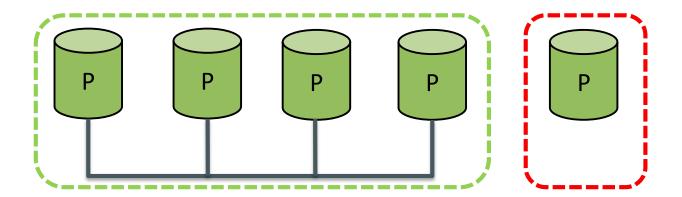
- Feature rich new replication plugin based on proven distributed systems algorithms (Paxos).
  - Compression, multi-platform, dynamic membership, distributed agreement, quorum based message passing, SSL, IP whitelisting.
- No third-party software required.
- No network multicast support required.
  - MySQL Group Replication can operate on cloud based installations where multicast is unsupported.

- 4.1 Automatic distributed server recovery
- 4.2 MySQL Look & Feel
- 4.3 Full GTID support
- 4.4 Auto-increment configuration/handling
- 4.5 Plugin version access control
- 4.6 Built-in communication engine
- 4.7 Read-only mode



# Read-only mode

- When a member joins the group, during distributed recovery, read-only mode is automatically set.
- On the unlikely event of a member failure, read-only mode is set automatically to prevent inconsistency with the group and member state changes to ERROR.





4.8 Full stack secure connections



#### Full stack secure connections

- Group Replication supports secure connections along the complete stack:
  - Distributed recovery connections
  - Connections between members
  - Client connections
- IP Whitelisting
  - Restrict which hosts are allowed to connect to the group
  - By default it is set to the value AUTOMATIC, which allows connections from private subnetworks active on the host

- 4.8 Full stack secure connections
- 4.9 Parallel applier support



# Parallel applier support

- Reduces applier lag and improves replication performance considerably.
- The same configuration options as asynchronous replication.

```
--slave_parallel_workers=NUMBER
```

--slave\_preserve\_commit\_order=ON



<sup>--</sup>slave\_parallel\_type=logical\_clock

### Parallel applier support

- Write set Based Transaction Dependencies
  - Already used on Group Replication from the beginning
  - Speedup distributed recovery time

```
master> SET @@GLOBAL.binlog_transaction_dependency_tracking=WRITESET;
Query OK, 0 rows affected (0,00 sec)

master> SET @@GLOBAL.binlog_transaction_dependency_tracking=WRITESET_SESSION;
Query OK, 0 rows affected (0,00 sec)

master> SET @@GLOBAL.binlog_transaction_dependency_tracking=COMMIT_ORDER; -- default
Query OK, 0 rows affected (0,00 sec)
```

- 4.8 Full stack secure connections
- 4.9 Parallel applier support
- 4.10 Transaction SAVEPOINT support



#### Transaction SAVEPOINT support

```
mysql> BEGIN;
Query OK, 0 rows affected (0,00 sec)
mysql> INSERT INTO t1 VALUES(1);
Query OK, 1 row affected (0,00 sec)
mysql> SAVEPOINT S1;
Query OK, 0 rows affected (0,00 sec)
mysql> INSERT INTO t1 VALUES(2);
Query OK, 1 row affected (0,00 sec)
mysql> ROLLBACK TO S1;
Query OK, 0 rows affected (0,00 sec)
mysql> COMMIT;
Query OK, 0 rows affected (0,00 sec)
```

- 4.8 Full stack secure connections
- 4.9 Parallel applier support
- 4.10 Transaction SAVEPOINT support
- 4.11 Requirements



# Requirements (by design)

- Requires InnoDB storage engine
- Primary key is required on every table
- Requires global transaction identifiers turned on
- Requires binary log turned on
- Requires binary log row format
- Optimistic execution: transactions may abort on COMMIT due to conflicts with concurrent transactions on other members
- Up to 9 servers in the group

#### Forbidden

- Serializable (on multi-primary)
- Cascading Foreign Keys (on multi-primary)
- Binary log events checksum

#### Warnings

- Concurrent DDL (on multi-primary)
- SELECT \*\*\* FOR UPDATE does not have group locking (on multi-primary)

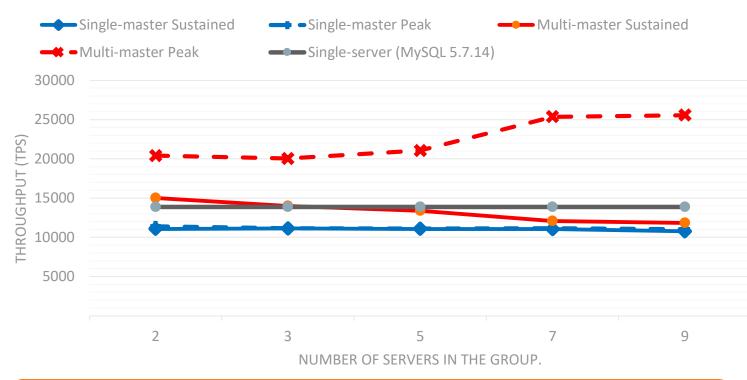


# Performance



#### Performance

# Group Replication Throughput (as perceived by the client application)



More on this subject on the series of replication performance blogs at: <a href="http://mysqlhighavailability.com/category/performance/">http://mysqlhighavailability.com/category/performance/</a>

#### Peak Throughput (i.e., no flow control)

The number of transactions that writers can propagate to the group (per second).

#### **Sustained Throughput (i.e., flow control)**

The number of transactions that can be propagated to the group without increasing the replication lag on any member (per second).

#### Servers

9 Dual Xeon E5-2660-v3 Enterprise SSD Storage 10Gbps Ethernet Network

#### Client

1 Dual Xeon E5-2699-v3 10Gbps Ethernet Network Sysbench 0.5 RW workload



#### Performance

- On a sustained throughput:
  - Multi-primary performance degrades gracefully while going from a group with 2 servers to a group with 9 servers.
  - Single-primary performance degrades marginally when growing the group size.
- On a peak throughput:
  - Multi-primary exhibits 1.8X speedup when compared to the single server.
    - Read load is balanced across the servers in the group.
    - Write load is lower since execution is balanced across the group, whereas in single-primary mode the primary becomes a bottleneck.
  - With a single-primary there is no lag on the other members.

# • Architecture

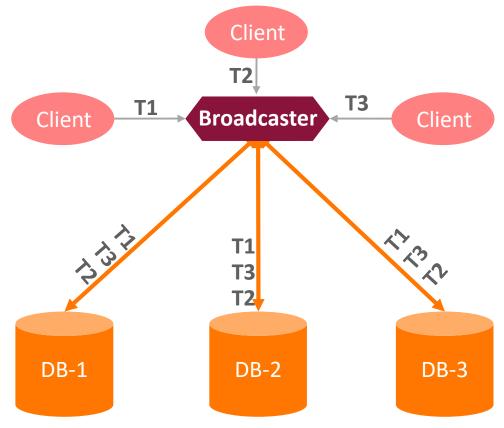
6.1 The Theories of Group Replication



#### The Theories of Group Replication

#### State Machine Replication

- All servers are initialized at same state.
- Same inputs in same order generate same output state.



**State Machine Replication Model** 



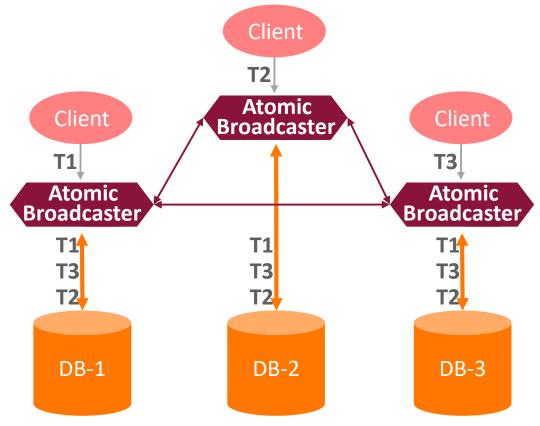
#### The Theories of Group Replication

#### State Machine Replication

- All servers are initialized at same state.
- Same inputs in same order generate same output state.

#### Atomic Broadcast System

- Messages are totally ordered
- All servers receive same messages in same order



**Distributed State Machine Replication Model** 



# Architecture

- 6.1 The Theories of Group Replication
- 6.2 Group Replication Architecture



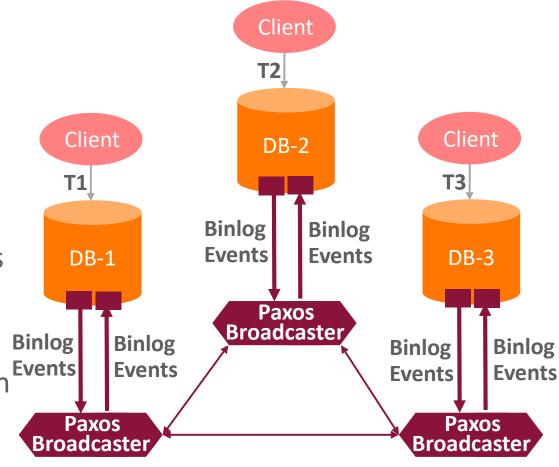
#### Group Replication Architecture

#### Broadcast Binary Log Events

- Row based data log
- Captures and broadcasts binary log events after transaction execution

#### Certification(Conflict Detection)

- Execution before atomic broadcasting causes conflict
- All conflicting transactions are broadcasted
- But only the first broadcasted transaction can commit, others are rolled back



**MySQL Group Replication Architecture** 

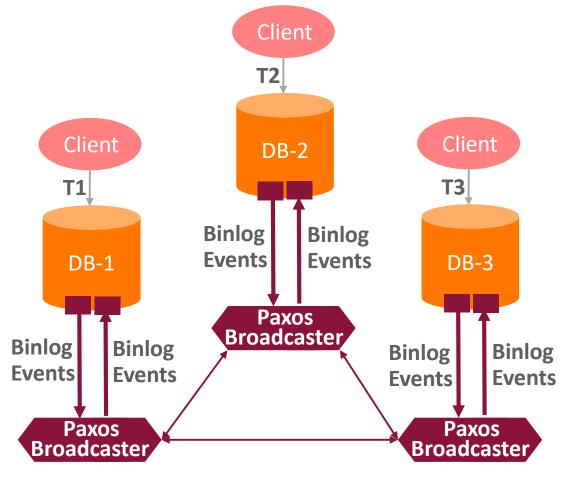
#### **Group Replication Architecture**

#### Binlog Event Applier

- Injects binlog events into relay log
- Controls binlog event automatically

#### Paxos Based Atomic Broadcaster

- Concensus
- Require majority servers available



**MySQL Group Replication Architecture** 



# Architecture

- 6.1 The Theories of Group Replication
- 6.2 Group Replication Architecture
- 6.3 Transaction Life Circle

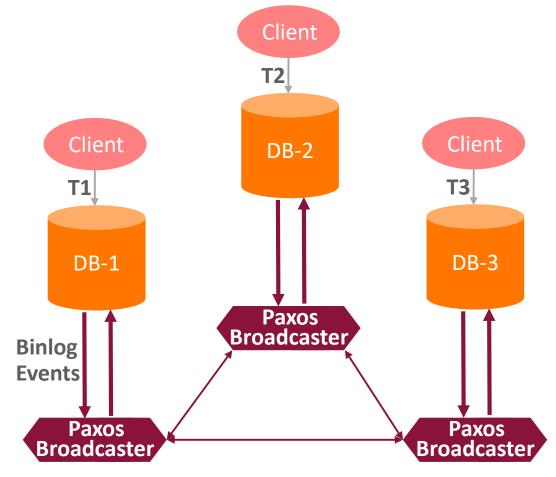


#### Local Transaction

- Executed by the client
- -T1 is local transaction on DB-1

#### Remote Transaction

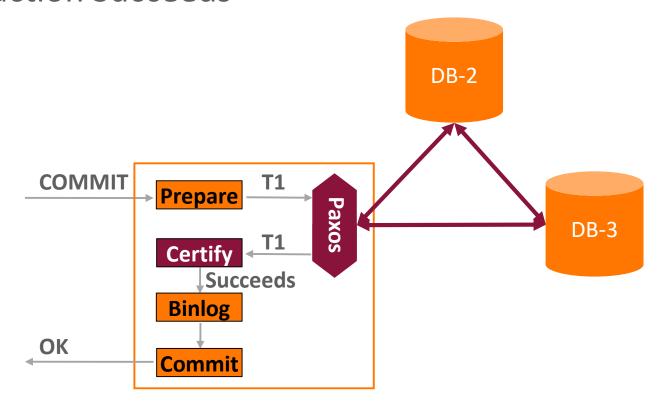
- Replicated from other servers
- T2,T3 are remote transaction on DB-1



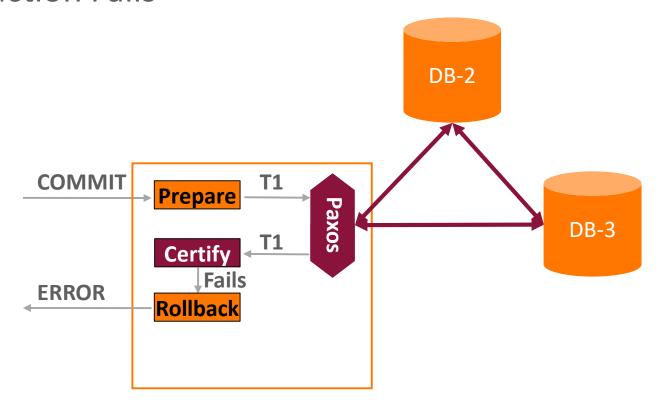
**MySQL Group Replication Architecture** 



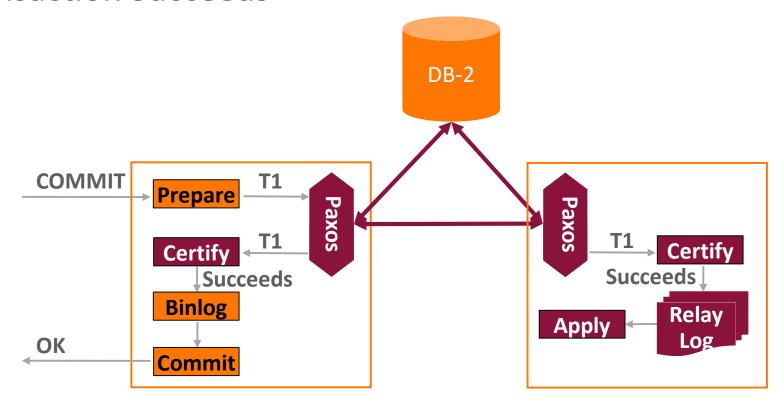
Local Transaction Succeeds



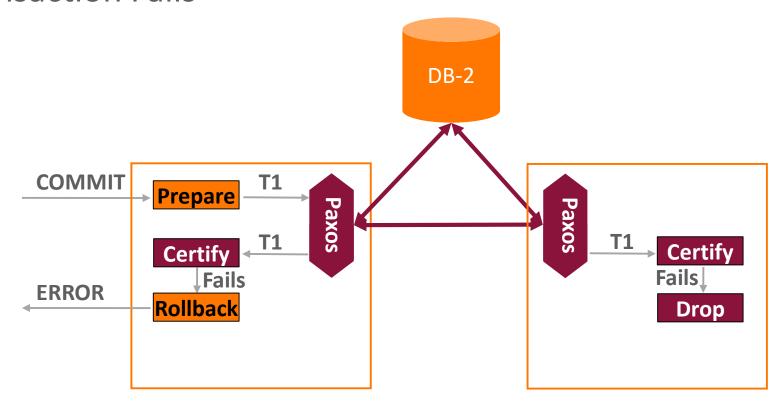
Local Transaction Fails



Remote Transaction Succeeds



Remote Transaction Fails

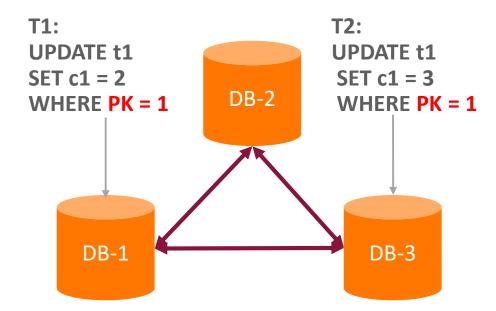


# 6 Architecture

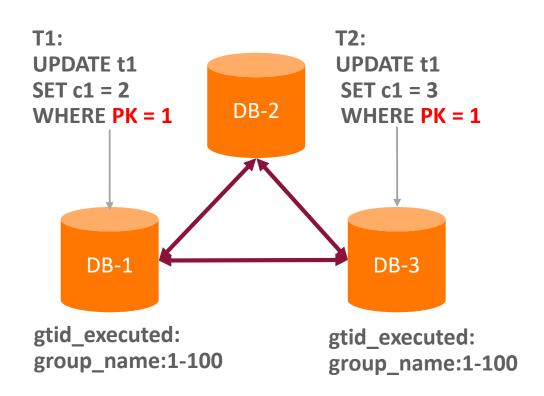
- 6.1 The Theories of Group Replication
- 6.2 Group Replication Architecture
- 6.3 Transaction Life Circle
- 6.4 Certification(Conflict Detection)



- Primary Key Based Conflict Detection
  - Possible conflict if two transactions updated same rows
  - Never conflict if two transactions updated different rows
  - Don't support DDL yet

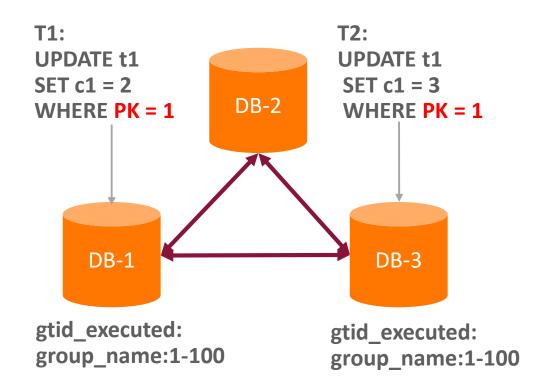


- Snapshot Version
  - Snapshot version is the value of gtid\_executed variable before broadcasting the transaction
  - Snapshot of certifying transaction must include the GTIDs of the certified transactions which updated the same rows



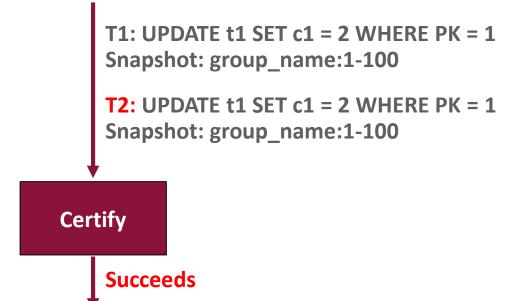
- Conflict Detection Database
  - A list of certified transactions' PK HASH and GTID set pairs.
  - Every member has one.

PK HASH	GTID SET
***	
db1.t1.pk=1	group_name:1-50
•••	



- Conflict Detection
  - T2's snapshot includes the GTID set in conflict detection database, so certification succeeds.

group\_name:1-100 > group\_name:1-50



PK HASH	GTID SET
db1.t1.pk = 1	group_name:1-50
•••	

- Conflict Detection
  - T2's snapshot includes the GTID set in conflict detection database, so certification succeeds.

group\_name:1-100 > group\_name:1-50

- Update Detection Database
  - Fill GTID SET with T2's snapshot and GTID
  - Suppose T2's GTID is group\_name:101

T1: UPDATE t1 SET c1 = 2 WHERE PK = 1
Snapshot: group\_name:1-100

T2: UPDATE t1 SET c1 = 2 WHERE PK = 1
Snapshot: group\_name:1-100

Certify

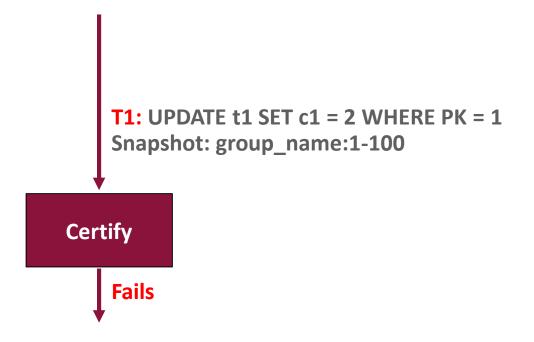
Succeeds

PK HASH	GTID SET
db1.t1.pk = 1	group_name:1-101
***	



- Conflict Detection
  - T1's snapshot do NOT include the GTID set in conflict detection database, so certification fails.

group\_name:1-100 < group\_name:1-101



PK HASH	GTID SET
db1.t1.pk = 1	group_name:1-101
***	



# Conclusion



## Summary

#### Cloud Friendly

 Great technology for deployments where elasticity is a requirement, such as cloud based infrastructures.

#### Integrated

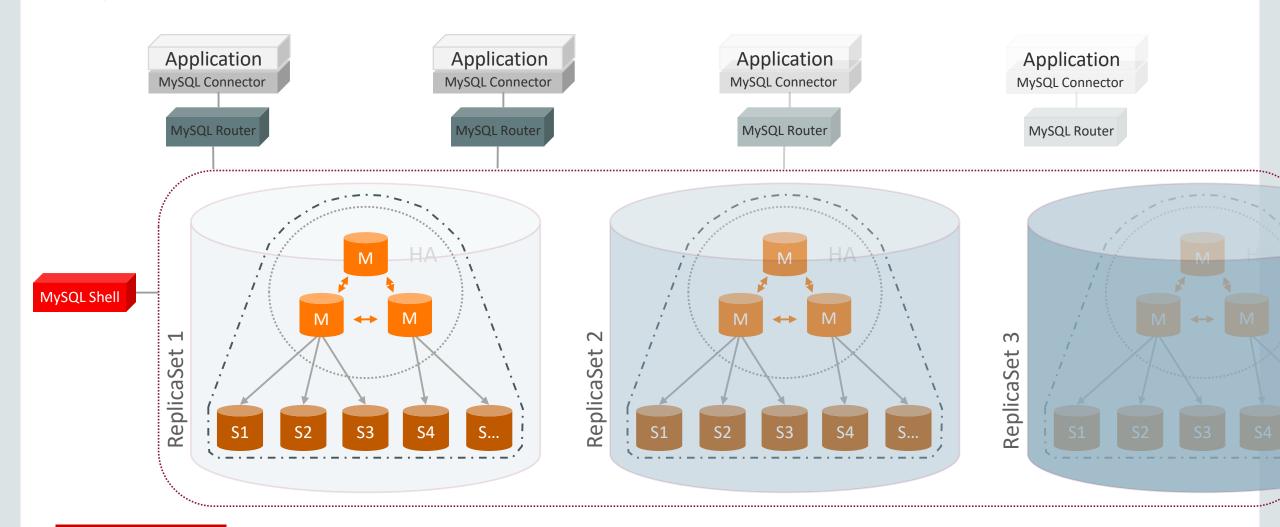
- With server core through a well defined API.
- With GTIDs, row based replication, performance schema tables.

#### Autonomic and Operations Friendly

- It is self-healing: no admin overhead for handling server fail-overs.
- Provides fault-tolerance, enables multi-primary update everywhere and a dependable MySQL service.
- Plugin **GA version** available with MySQL 5.7.17+, available on 8.0.1+



# MySQL InnoDB Cluster: The End Goal





# Where to go from here?

- Packages
  - http://www.mysql.com/downloads/
- Documentation
  - http://dev.mysql.com/doc/refman/5.7/en/group-replication.html
  - http://dev.mysql.com/doc/refman/8.0/en/group-replication.html
- Blogs from the Engineers (news, technical information, and much more)
  - http://mysqlhighavailability.com



# ORACLE®