Advanced Security Protocols - SSL/TLS

Design and Verification of Security Protocols and Security

Ceremonies

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- Major websites use TLS to secure all communications between their servers and web browsers;
- Although not common, TLS can be used for mutual authentication.
- TLS aims primarily to provide privacy, data integrity and peer authentication between two communicating computer applications.

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- It was first defined in 1999 and updated in RFC 5246 (August 2008) and RFC 6176 (March 2011);
- It builds on the earlier SSL specifications (1994, 1995, 1996) developed by Netscape Communications for adding the HTTPS protocol to their Navigator web browser.

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 - Loss or lack of authentication usually leads to major attacks on TLS;

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- TLS can also provide forward secrecy, ensuring that any future disclosure of encryption keys cannot be used to decrypt any TLS communications recorded in the past.
- Most properties of the protocol are verified by different methods, ensuring no major flaws exist at conceptual level;

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 - Beast (Browser Exploit Against SSL/TLS), POODLE (Padding Oracle On Downgraded Legacy Encryption), Sweet32 (attacks on half of the CBC block), Heartbleed (Implementation bug);
 - All of these have to do with implementation of the pre-conditions that fail in practice.

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- Record protocol is a simple symmetric encryption engine.

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- Each record has a content type field that designates the type of data encapsulated, a length field and a TLS version field;
- The data encapsulated may be control or procedural messages of the TLS itself, or simply the application data needed to be transferred by TLS.

TLS Record Frame

+	Byte +0	Byte +1	Byte +2	Byte +3
Byte 0	Content type			
Bytes 14	Version		Length	
	(Major)	(Minor)	(bits 158)	(bits 70)
Bytes 5(<i>m</i> -1)	Protocol message(s)			
Bytes m(p-1)	MAC (optional)			
Bytes p(q-1)	Padding (block ciphers only)			

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- The specifications (cipher suite, keys etc.) required to exchange application data by TLS, are agreed upon in the "TLS handshake":
- This happens between the client requesting the data and the server responding to requests;

Basic TLS Handshake

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- The server responds with a ServerHello message, containing the chosen protocol version, a random number, CipherSuite and compression method from the choices offered by the client;
- The server also sends its Certificate message, its ServerKeyExchange message and a ServerHelloDone message, indicating it is done with handshake negotiation;

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- This PreMasterSecret is encrypted using the public key of the server certificate;
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- All other key data for this connection is derived from this master secret;

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- Application phase starts using TLS record and parameters established at handshake time.

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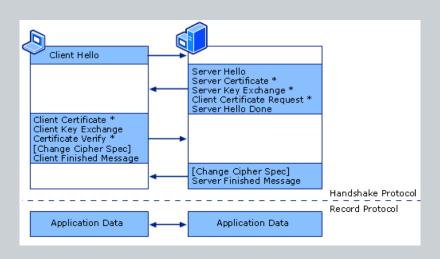
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- We can also have a Resumed handshake which avoids the use of asymmetric crypto;
- They are based on sessionId or sessions tickets established at ChangeCipherSpec time;
- This is very handy for busy server and it heavily depended on forward secrecy.



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Byte 0	22			
Bytes 14	Version		Length	
	(Major)	(Minor)	(bits 158)	(bits 70)
Bytes	Message type	Handshake message data length		
58		(bits 2316)	(bits 158)	(bits 70)
Bytes 9(<i>n</i> -1)	Handshake message data			
Bytes	Message type	Handshake message data length		
n(n+3)		(bits 2316)	(bits 158)	(bits 70)
Bytes (n+4)	Handshake message data			

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- Paulson proved correctness of TLS using Isabelle/HOL and the Inductive method.

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- Is up to-date one of the best abstractions and most precise abstract verification of the suite.

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- How come deeply verified and widely deployed protocols can still have problems?

Questions????



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