

Protocol Verification Techniques - Theorem Provers

Design and Verification of Security Protocols and Security Ceremonies

Programa de Pós-Graduação em Ciências da Computação
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Attention!

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This topic will be divided into two lectures. One will deal with automatic theorem provers using FOL and the second will deal with theorem provers using HOL

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- HOL is any predicate logic that has greater order than Second-Order Logic;

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- Higher-order logic admits quantification over sets that are nested arbitrarily deeply.

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- This is the definition of the set of Natural Numbers.

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- <http://www.geocities.ws/robrich18/Larry.html>



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- Last work published in 2015: "Verifying multicast-based security protocols using the inductive method. Martina, J.E., Paulson, L.C."

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- We stated the goals and theorems and lemmas;
- We prove the theorems inductively to demonstrate correctness.

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 - Use a theorem prover (Isabelle) to write the proofs.

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 - Generic treatment of inference rules;
- Powerful simplifier, classical reasoner and connected tools;
- Strong support for inductive definitions.

Inductive Method Support in Isabelle

Due to my lack of time we will jump to my Ph.D Thesis to get the explanation from there. I promise next time it will be everything on the slides..

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- Can this be extended?
- What this strategy can not do?

Questions????



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