Reference C++ Code A

```
if (fork() == 0) {
    // do stuff
} else if (fork() == 0) {
    // do more stuff..
} else if (fork() == 0) {
    // do more "stuff"..
} else if (fork() == 0) {
    // noooo
}
```

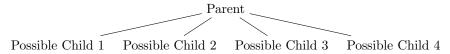
Reference C++ Code B

Reference C++ Code C

These code snippets were obtained from the Lab 6 assignment page.

1. What should be the resulting process *FAMILY TREES* from these two (Code A and Code B) code snippets? Illustrate.

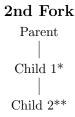
In reference to code A, all fork() calls are performed at the same "hierarchy level" in the if-else tree. The entry point starts at if (fork() == 0), which produces the first child process. However, should this fork() fail, the succeeding fork() calls in the following else if (fork() == 0) statement produces a child process accordingly. Once a fork() successfully produces a child, the program enters the respective instruction block and stops forking. The resulting process family tree from code A looks like:

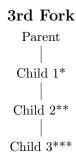


In this diagram, Possible Child n for $n \geq 2$, only gets instantiated if the first n-1 fork(s) fail.

In reference code B fork() calls are performed sequentially, with further forks being contained in the block associated with fork() == 0. The fork system call returns multiple values: (i) -1 if the fork was unsuccessful, (ii) 0 is returned to the child process, (iii) all positive nonzero values are returned to the parent. Each succeeding fork() is contained within an if fork() == 0 block—meaning that each succeeding fork() is only performed by the child process. We can illustrate the process family tree per stage up to the third fork:







In this diagram each * denotes the nth fork at which the child was instantiated, where n represents the amount of asterisks trailing a child label.

2. Will the last line in the sample code below (Code C) still be printed? How about when using execl()? Why or why not? Explain.

What Code C effectively does is to replace the current process with the gedit program from the user's programs (/usr/bin/... serves that purpose in UNIX-based environments) directory. Should the exec() call succeed then the current process is replaced with the program located within the specified path—if the call fails then the statement resolves to -1.

Assuming that gedit is properly installed within the machine, or that there is a program to be executed in whatever path is passed into execv(), then the program will never execute the instructions: cout << "will this line still be printed?" << endl, since the process containing those instructions would have already been replaced with what was passed into execv().

On the otherhand if the path passed into execv() does not exist, the current process does not get replaced and instead continues executing instructions stored within the current *.cpp file. In this case there will be two outputs to the terminal Error. Booooo! and Will this line still be printed?—since the process was not replaced by the call.

The use of execl() will not make any major changes to the logic of whether or not the last line will be printed to the terminal or not, the only thing that changes is how the path is passed into the exec() function. Whereas in execv() we passed the file path with an array of arguments, execl() requires that the file path is passed followed by each argument as separate parameters. The logic of whether or not the process gets replaced still is the same assuming that execl() is used properly, the only way it differs is in syntax.