```
Memory - Static + Dynamic
10-1/
           Statiz - an known position in program text to
                      call free (eig. when we're done with arg
                      or En)
                   - not state, the free position is a property of
                      evaluation.
          "Stack Memory" an implementation
                                               #Static memory
          ('s local wars
          "Heap memory" is dynamic memory
                                           dyn = {clo, E[XEV]}
          CEK = static = { ang, fn }
          Strategy for youin C:
          - I make mostates w/ off-by-1 (sometimes reference undefined memory)
          - Imput Letermines your allocation pattern
          - Allocation is Letermined by a computation
             Linervous about a hases (other pointers to the values)
          - It's a global decision
          - Freeing too early, may cause reuse, may cause inconsistent use => crash
 Soundness
          - Free too late => vie too much memory
 Completeness
            Soundness - never free too early
          f(10); ( what if I never retires?
         return x[2];
```

	Memory Management:
	- Necessary to be sound
	- As close to complete and little impact on the aspossible
	MM on space
alloc	peak memory use m
	time
	MMZ MMZ
	imagine MMm where pmu (MMm) < pmu (MMx) Y MMx
	MM on time minimize
	total time = your + mm
	your Winimize just mm
	Program & minimize % in mm
entrance and participation considerable	minimize the maximum mm slice
and the second s	maximize the mini yar slice
	Common C MM
	- Insert calls to free() when you think it's unight
The second secon	- Put in lots of copies to remove aliases Lityavine wrong,
or completely, to you also from the production and as a second	L-7 spend more space (i.e. less complete) Hen - sand
	- Pury Lea alloc - runs in Ign where n are the # of objects
heer 1 ook ediginaan oor - 12 eelhiib eelalah ee	free - lg n
	space - at worst 2x space
	"Garbage Collector"

10-2