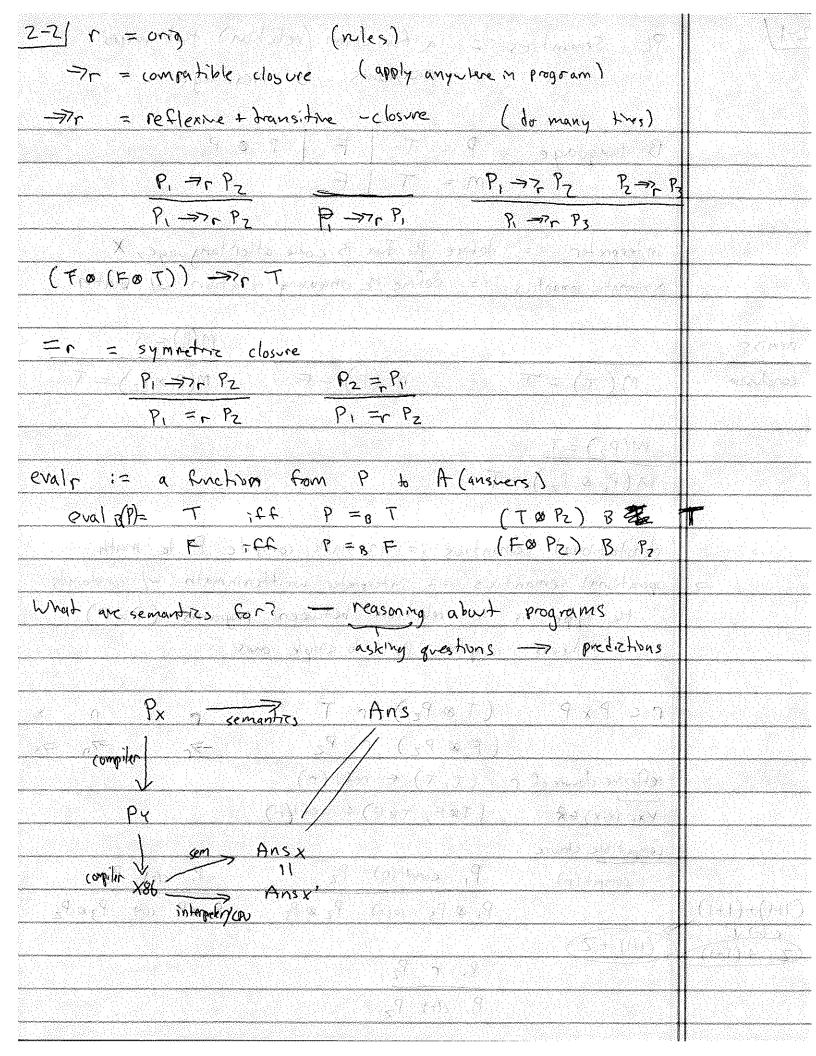
<i>IJ</i>	PL Semantics: = a function (relation) that wayses?			
	PL Semantics:= a function (relation) that maps? programs to meanings			
ant faire i de la companya de la co	Blangrage P=T F P & P			
	M = T + F			
	interpreter := define the fun is come other language, X			
	axiomatic semantics == define the maguing relation W/ Amouth			
premise	$M(P_i) = 7$			
Conclusion	$M(\tau) = \tau$ $M(F) = F$ $M(P_1 \otimes P_2) = \tau$			
ada a mara nga paga at a managana at a mara 1911 na 1,19 kao jiya na isan at a mana 1900.				
	$M(\Omega) = T$			
	$M(P_2) = T$ $M(P_1 \otimes P_2) = T$			
	denotational semantics:= in math, compile P to month			
	denotational semantics:= in math, compile P to month			
	denotational semantics := in math, compile P to month operational semantics := interpreter writteninmath w/ constaints			
	denotational semantics:= in math, compile P to month operational semantics:= interpreter writteninmath w/ constands the type is a relation between programs (PxP)			
	denotational semantics:= in math, compile P to month operational semantics:= interpreter writteninmath w/ constaints the type is a relation between programs (PXP) "reducess complex Ps to simple ones			
	denotational semantics:= in math, compile P to month operational semantics:= interpreter writteninmath w/ constaints the type is a relation between programs (PXP) "reducess complex Ps to simple ones TCPXP (T&Pz) TT T			
	denotational semantics:= in math, compile P to month operational semantics:= interpreter writteninmath w/ constants the type is a relation between programs (PxP) "reducess complex Ps to simple ones TCPXP (T&Pz) r T r r n x (F&Pz) r Pz -7n ->x			
	denotational semantics:= in math, compile P to month operational semantics:= interpreter written in math w/ constaints the type is a relation between programs (PxP) "reducess complex Ps to simple ones TCPXP (TOPZ) TT TOPZ TRESPENSE closure of T(T,T) & refl(T)			
	denotational semantics:= in math, compile P to month operational semantics:= interpreter writteninmath w/ constaints the type is a relation between programs (PxP) "reducess complex Ps to simple ones reflexive clusure of r (T, T) & refl(r) \(\text{Kx.} (x,x) \) & (ToF, ToF) \) reflexive			
	denotational semantics:= in math, compile P to month operational semantics:= interpreter writteninmath w/ constaints the type is a relation between programs (PxP) "reducess" complex Ps to simple ones rc PxP (T&Pz) r T reformation in x (F&Pz) r Pz -> -> -> -> -> -> -> -> -> -> -> -> ->			
	denotational semantics:= in math, compile P to mostly operational semantics:= interpreter writteninmath w/ constaints the type is a relation between programs (PXP) "reducess complex Ps to simple ones reflexive closure of r (T, T) & refl(r) TX. (X,X) & (T&F, T&F) & refl(r) compatible closure remet(r) P, compation Pz P, c(r) Pz			
(1+1)+(1+1)	denotational semantics:= in math, compile P to month operational semantics:= interpreter writteninmath w/ constaints the type is a relation between programs (PxP) "reducess" complex Ps to simple ones TCPXP (T&Pz) r T (F&Pz) r Pz reflexive closure of r (T, T) & refl(r) Vx. (x,x) & (T&F, T&F) & refl(r) compatible closure rompatible closure P1 & P3 &			
(I+1)+(I+1)	denotational semantics:= in math, compile P to moth operational semantics:= interpreter written in math w/ constaints the type is a relation between programs (PxP) "reducess" complex Ps to simple ones TCPXP (T&Pz) r T r n x (F&Pz) r Pz - 7n ->x reflexive closure of r (T,T) & refl(r) *** **Tx, Ix,x) & Tx, Tx, F, Compatible closure compatible closure **Type Pi compatible closure **Typ			
(1+1)+(1+1)	denotational semantics:= in math, compile P to moth operational semantics:= interpreter writteninmath w/ constaints the type is a relation between programs (PxP) "reducess complex Ps to simple ones TCPXP (T&Pz) TPz TOPXP (T&Pz) TPz reflexive closure of r (T,T) & refl(r) Wx. (x,x) & (T&F, T&F) & refl(r) compatible closure compatible closure Pl & P3 & P3 & P3 & P3 & P3 & P3 & P2			



1-3/	Semantics define program meaning which facilitates verification of programs
	which facilitates verification of Programs
	+ program fools (like outlies)
And grapher to Superior to Superior or Superior	
	Desirenthle: Programs are leterministic.
California de Alexandro de La California de	$\forall P. \forall A \forall A. , eval r(P) = A. Specialized for one P$ $A = val r(P) = A.$
) フ 1 emlr (P) = Az
	for all programs $\Rightarrow A_1 = A_2$ in language
mangan (Maria (M.)) (Mangangan) (Mangan) (Mangan) (Mangangan) (Mangangan) (Mangangan) (Mangangan) (Mangangan) (Mangangan) (Man	main: f(in+*x) { g(m+*x) { }
az 2 kippi, kemisti kasa Casapat ke endermining dan kemisi kempanya sa dalam dan dan dan dalam dan dan dalam d Barangan dan dan dan dan dan dan dan dan dan d	main: $f(m+*x) \in g(m+*x) \in lock \Rightarrow x = lock \Rightarrow x = *x + 5;$ $f(m+*x) \in lock \Rightarrow x = *x + 2;$ $f(m+*x) \in lock \Rightarrow x = *x + 2;$ $f(m+*x) \in lock \Rightarrow x = *x + 2;$ $f(m+*x) \in lock \Rightarrow x = *x + 2;$ $f(m+*x) \in lock \Rightarrow x = *x + 2;$
	thead-areate (f, 8x) 3
	Hneal-creak (4, dx)
	nail_theadslocking access to X
	nail_Heads locking access to X ret X => 30,25/15,20
	competible doesn't "where" work happens
	(1+71) + (3+4) Church-Rosser
مناسبته والمناسلة والمناسبة والمناسب	/ Ym,n.
	$\frac{1}{(1+2)+7} / If M=rN, \exists L,$
	3+(3+4) / M → 7 L
	and Narl.
	3+7
	Diamond Diamond
	then FL', m'>>>L'
	m N - and N' >>> L'

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