

1-1

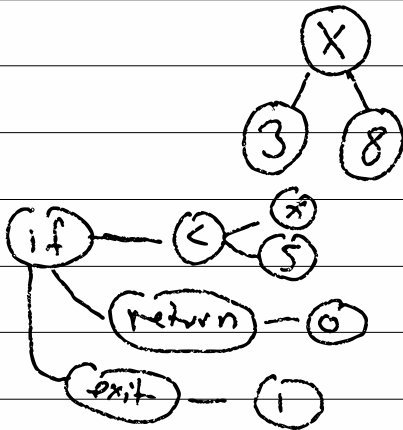
1 + 1

5

1 +

1 x 3

" 3 x 8 "

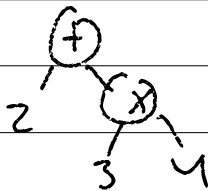


```
if (x < 5) {  
    return 0; }  
else {  
    exit(1); }
```

(+ 1 1)
bp → → → children



(+ 2 (x 3 4))



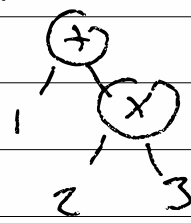
$$\frac{}{J_0 \Rightarrow e ::= v \mid (+ e e)}$$

$$v ::= \text{number} \mid (* e e)$$

$$(+ 1 (+ 2 3)) \in J_0$$

```

interface Joe { }
class JNumber implements Joe {
    int n; JNumber(m) { n = -m; }
}
class JPlus imp Joe {
    Joe left, right; JPlus(...) }
class JMult imp Joe {
    Joe l, r; JMult(...) { } }
    
```

$$(+ 1 (* 2 3)) \xrightarrow{\text{Sexpr}} =$$


$$=$$

$$\text{new JPlus}($$

$$\text{new JNum}(1),$$

$$\text{new JMult}($$

$$\text{new JNum}(2))$$

$$\text{new JNum}(3)))$$

$$= \text{JP}(\text{JN}(1), \text{JM}(\text{JN}(2), \text{JN}(3)))$$

```

class JPlus:
    def __init__(l, r):
    
```

$$\text{BST } n ::= m + \mid (\text{br num } n \ n)$$

```

        this.l = l;
        this.r = r;
    
```

1-3/0 pp = J0 \Rightarrow string

③ pp n = intos(n)

⑦ pp (+ eL eR) = "(" + pp(eL) + " + " + pp(eR) + ")"

④ pp (x eL eR) = "(" + pp(eL) + " * " + pp(eR) + ")"

① interface J0 { public String pp(); }

② class JNum { ...

public String pp() {

return intToStr(n); }

③ class JPlus {

public String pp() {

return this.left.pp() + " + " + this.right.pp(); }

1-4/ big-step interpreter

interp : $e \rightarrow v$

interp $n = n$

interp $(+ e_L e_R) = \text{interp } e_L + \text{interp } e_R$

① interp $(* e_L e_R) = \text{interp } e_L * \text{interp } e_R$

→ class JMult {
 public ^{int} interp() {
 return this.left.interp() * this.right.interp();
 }}

$(+ 1 2 3) = (+ 1 (+ 2 3))$
 └─ desugar →

se = empty | (cons ^{main} se se) | string
 $(a\ b\ c) = (\text{pair } "a" (\text{pair } "b" (\text{pair } "c" \text{mt})))$
 $(+ 1 2) = (\text{p } "+" (\text{p } "1" (\text{p } "2" \text{mt})))$
 $(+ 1 (+ 2 3)) = (\text{p } "+" (\text{p } "1" (\text{p } (+ 2 3) \text{mt})))$

1-5 / desugar for \mathcal{J}_0

$$("-" e) \Rightarrow (* -1 (\text{desugar } e))$$

$$("-" e_1 e_2) \Rightarrow (+ (d e_1) (d e ("-" e_2)))$$

$$("+") \Rightarrow 0$$

$$("+" e_1 \text{ more } \dots) \Rightarrow$$

$$(+ (d e_1) (d ("+" \text{ more } \dots))))$$

$$\begin{array}{ccc} 'x' & & = 1 \\ "x" & \downarrow & \downarrow \downarrow \\ (x & & "x" \end{array}$$

$$\text{se} \Rightarrow \mathcal{J}_0 \Rightarrow v$$

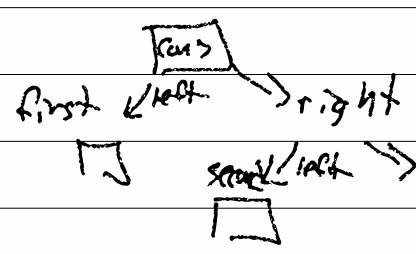
desugar interp

$$\begin{array}{c} \searrow \\ \text{compile} \end{array} \quad bc \Rightarrow_{vm} v$$

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desugarer $(- e_1) \Rightarrow (\hat{x} -1 e_1')$
 $(- e_1 e_2) \Rightarrow (\hat{x} e_1' (-e_2))$
 $(+) \Rightarrow 0$
 $(+ e_1 e_2 \dots) \Rightarrow$
 $(\hat{x} e_1' (+ e_2 \dots))$

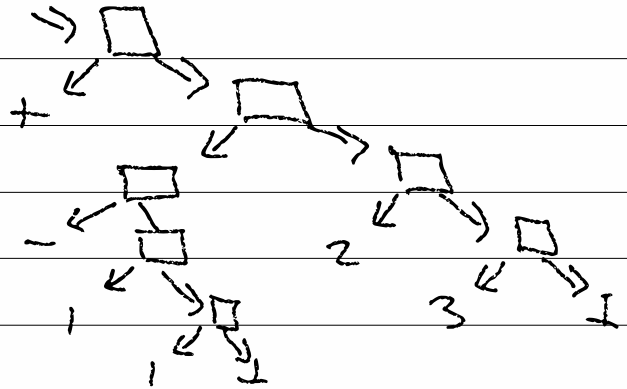
```
def
  desugar (se) :
    if isList(se) && length(se) = 2 &&
      first(se) == "-" then
      > def length(se) :
      >   if isNull(se) : ~return 0
      >   else if Cons(se) : return 1 + length(right(se))
      >   else false
      new JMult ( new JNum(-1), desugar(
        second(se)))
```



if isList(se) && len(se) = 3 && first(se) == "+"
 then : return new JAdd (desugar(second(se)),
 desugar(new Cons("-", ~~new~~
 new Cons(third(se), null))

22 | $\text{cons}(a, \text{cons}(b, \text{conc}(c, \text{null}))))$
 $\text{len} = 3$

$(+ \quad (- \quad 1 \quad 1)) \quad \text{len} = 4$
 $\begin{array}{ccccccc} & \nearrow & & \nearrow & & \nearrow & \\ 1 & & 2 & & 2 & \leftarrow & 3 \\ & & & & 3 & & \\ & & & & \nearrow & & \\ & & & & 4 & & \end{array}$



$\text{isList}(se):$

$\text{Null} = \text{True}$

$\text{cons} : \text{isList}(\text{right}(se))$

$\text{ow} : \text{False}$

$\text{Len}(se):$

$\text{Null} : 0$

$\text{cons} : 1 + \text{len}(\text{right}(se))$

2-3 / $J_0 \rightarrow J_1$

\swarrow fun \swarrow args
 $e ::= v \quad | \quad (e \ e \ \dots)$
 $\quad \quad \quad | \quad (\text{if } e \ e \ e)$
 $\quad \quad \quad \quad \quad \nearrow \quad \nearrow \quad \nearrow$
 $\quad \quad \quad \quad \quad c \quad + \quad f$

$v ::= b$

$b ::=$ some set of constants

// in J_0 , $b = \text{num} \mid + \mid *$

numbers \mid bools \mid prim

prim = $+, -, *, /, \leq, <, =, >, \geq, \dots$

interp $v = v$

interp $(\text{if } e_c \ e_t \ e_f) = \text{interp } e_k$

where $e_k = \text{if } \text{interp } e_c \text{ then}$
 $\quad \quad \quad e_t \quad \text{o.w. } e_f$

interp $(e_f \ e_a \ \dots) = \delta(p, v_a \ \dots)$

where $p = \text{interp } e_f$

$v_a = \text{interp } e_a \ \dots$

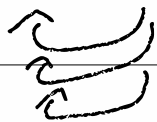
$\delta : b \ \dots \rightarrow b$

$\delta(+, 1, 2) = 3$ $\delta(/, 1, 0) = \perp$

$\delta(<, 1, 3) = \text{true}$

24/ "small step interp"

$e \rightarrow e$



until its the same

interp

"big step"

$e \rightarrow v$

Interp

Interp $e =$

let $e' = \text{interp}(e)$

if $e == e'$ then

ret e

O.V.

Interp (e')

$(+ (+ 1 1) 2) \leftarrow (+ (+ 1 1) (+ 1 1))$
 $\Rightarrow (+ 2 2)$
 $\Rightarrow 4$

\swarrow

$(+ 2 (+ 1 1))$

int $x = 1;$

f ($x--$, $x++$)

(1, 0)

(2, 1)

2-5 step : $e \rightarrow e$

$$\left. \begin{aligned} \text{step } (\text{if true } e_1 \ e_2) &= e_1 \\ \text{step } (\text{if false } e_1 \ e_2) &= e_2 \\ \text{step } (p \ v_1 \ \dots) &= \delta(p, v_1 \ \dots) \\ \text{step } v &= v \end{aligned} \right\}$$

$$\left[\begin{aligned} \text{step } (\text{if } e(\&v) \ e_1 \ e_2) &= \\ &(\text{if } (\text{step } e) \ e_1 \ e_2) \text{ on } (\text{if } e \ (\text{step } e_1)) \\ \text{step } (v_1 \ \dots \ e(\&v) \ e_1 \ \dots) &= \\ &(v_1 \ \dots \ (\text{step } e) \ e_1 \ \dots) \end{aligned} \right]$$

A context

$$\begin{aligned} C &::= \text{hole} & | & \text{if0 } C \ e \ e \\ & & | & \text{if1 } e \ C \ e \\ & & | & \text{if2 } e \ e \ C \\ & & | & (e \ \dots \ C \ e \ \dots) \end{aligned}$$

$$\text{plug } C \ e \quad (C[e])$$

$$\text{plug } \text{hole } x = x$$

$$\text{plug } (\text{if0 } C \ e_1 \ e_2) x = \text{if } x \ e_1 \ e_2$$

$$\text{plug } (\text{if1 } e_1 \ C \ e_2) x = \text{if } e_1 \ x \ e_2$$

$$\text{plug } (e_1 \ \dots \ C \ e_2 \ \dots) x = (e_1 \ \dots \ x \ e_2 \ \dots)$$

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$$\text{step } C[\text{if true } e_1 \text{ et } e_2] = C[e_1]$$

$$\text{step } C[\text{if false } e_1 \text{ et } e_2] = C[e_2]$$

$$\text{step } C[p \text{ va } \dots] = C[\mathcal{S}(p, \text{va} \dots)]$$

~~step~~ "parse" : $e \rightarrow C \times e$

step $\xrightarrow{\quad \quad \quad} e$

$$\text{interp } e = \text{if } e \in v \text{ then } e$$

$$C, e' = \text{parse } e$$

$$e'' = \text{step } e'$$

$$\text{plug } C \ e''$$

parse : $e \rightarrow C \times e$ ← redex

$$\text{parse } (\text{if } e_c \text{ et } e_e) =$$

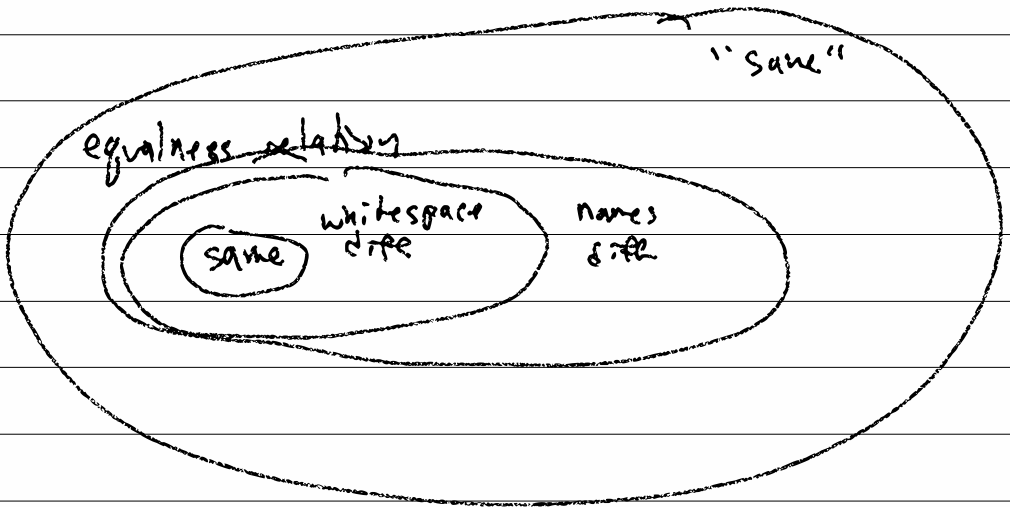
$$\text{if } e_c \in v \text{ then } (\text{hole}, e)$$

$$\text{o.w. let } C', e' = \text{parse } e_c$$

$$(\text{if } C' \text{ et } e_e, e')$$

2-7 Answer: Contexts

Question: How do I know when
two programs do the same thing?



$$x = y$$

$$\forall x. f x = g x$$

$$\forall c. c[x] = c[y]$$

$$C = \text{hole} \quad x = y$$

$$C = (+ \text{ hole } z) \quad x + z = y + z$$

$$C = (\text{map hole } (1 : + : z))$$

....

Observational Equivalence

2-8) $C ::= \text{hole} \mid \text{if } C \ e \ e$
 $\mid \text{if } e \ C \ e$
 $\mid \text{if } e \ e \ C$
 $\mid (e \dots C \ e \dots)$

$E ::= \text{hole} \mid \text{if } E \ e \ e$
 $\mid (v \dots E \ e \dots)$

"unique decomposition" \downarrow unique E
 $\forall e. e \in v \text{ or } e = E[e'] \text{ where } e' \in v$

$$g(1, 2) = 1$$

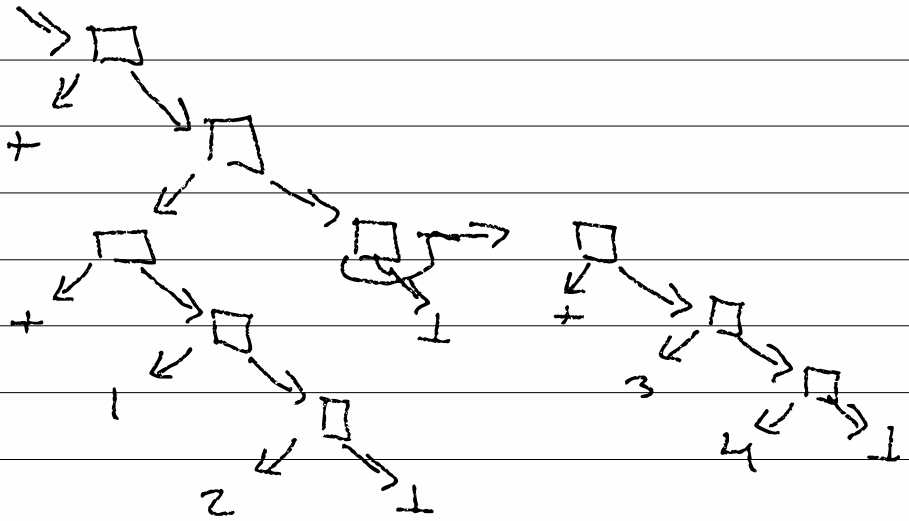
$$g(1, 1, 0) = 1$$

$(+ (+ 1 2) (+ 3 4))$ \leftarrow tree
 \leftarrow java tree

`new JPlus (new JPlus (new JNum(1),
new JNum(2))`

`new JPlus (new JNum(3),
new JNum(4)))`

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(1 (1 1 2)
(+ 3 4))