Control flow analysis

Control Flow Graph Basic Blocks

> · Segments of code without transfers of control.

2 -techniques

- · (ontinuation Passing Style (CPS) · Non-Standard Abstract Semantics
- · all functions get a new argument, K

All transfers of control are uniform

(2(V)

Calefine  $(f \times)$   $\rightarrow$   $(define Cf \times K)$   $(+ (g (h \times)) (g \times))) (h \times (2 (h \times))$   $(+ (g h \times) (g \times)))$ 

(define (f x K)

(h x (lambda (nx)

(g hx (lambda (ghx)

(+ ghx (gx)))))) (define (f x 5)

(hx (12 (hx)

(ghx()(ghx)

(gx, (2 cgx) (+ anx ax)))))

deline (h x K)

Stop Continuation: (2(V) V)

NSAS - Key behind abstraction is that the accuracy is traded for compile time computability.

Expr:== i/(+ e, e2)/(\*e1 e2)/(+e)

e,, ez E Expr i & Integer

 $\begin{array}{ll}
\Gamma(1) = i \\
\Gamma(+e_1e_2) J = \Gamma[e_1] + \Gamma[e_2] \\
\Gamma(*e_1e_2) J = \Gamma[e_1] * \Gamma[e_2] \\
\Gamma(-e) J = \Gamma[e_1] * \Gamma[e_2] \\
\Gamma(-e) J = \Gamma[e_1] * \Gamma[e_2] \\
\Gamma(+e_1e_2) J = \Gamma[e_1e_1] * \Gamma[e_1e_2] \\
\Gamma(+e_1e_2) J = \Gamma[e_1e_2] * \Gamma[e_1e_2] \\
\Gamma(+e_1e_2) J = \Gamma[e_1e_2] * \Gamma[e_1e_2] \\
\Gamma(+e_1e_2) J = \Gamma[e_1e_2$ 

Ssign [(+2 (&3 (-4)))]= -

Non-Standard Abstract Semantics Ssign  $Ssign[i] = Ssign[i] = Ssign[i] = Ssign[i] = Ssign[e_2] = Ssign[e_2]$ (0+) (r,0,13(0) (-10,13(-0,+3 Ssign [ (+2 (\*3 (-4)))]= (+ {+3 (x {+3 (- {+3))) (+ 8+3 (x & f3 (-3)) (+ 2+3 2-3)

2-0,+3

## (ontour s

Variable environment.

(define (f x K)