

jMeta 0.5

Design

Version: 0.2
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1 Introduction

This document specifies the technical design of `jMeta 0.5`.

Part I

Requirements

This chapter defines the most important functional and non-functional requirements for **jMeta**. Most relevant input is the document [?].

2 Scope

2.1 In Scope

`jMeta` is a Java library for reading and writing container and metadata formats. `jMeta` has the following goals:

- Define a robust, easy to use and generic API for reading and writing multimedia (i.e. audio, video and image) metadata
- In addition, define a way to also parse and change typical container formats that embed or surround multimedia metadata
- Be easily extensible with additional metadata or container formats, also by third party

Thus, clearly, `jMeta` targets applications and users in the multimedia editing area, e.g. software to manage an audio, video or image collection.

Special strenght of the library should be its versatility (in a sense of supported formats and generality) as well as its extensibility. At the same time, it targets to offer access to ALL features of each specific format, even the low-level ones, and thus does not want to hide anything from expert users who need fine-grained control.

Using `jMeta`, an application is allowed to access metadata quite generically and comfortably, or it can also explore deep specifics of each format down to the bit and byte level.

2.2 Out of Scope

Metadata not in the audio, video or image domain are not directly in scope of `jMeta` and thus the library might not offer the right abstractions for those. However, it is at least likely that other binary or textual formats can also be parsed quite the same way using `jMeta`.

Furthermore, `jMeta` clearly is no encoding or decoding library, it cannot understand codecs. It might however help to encode or decode by providing access to the low-level details of the container formats. It might be used as basis for extensible encoders or decoders, however, not exactly in high-performance areas where speed is of most importance.

Further things out of scope of specific versions are treated in “?? ??”.

2.3 Features of version 0.5

Here we give a brief overview of the features the library needs to offer in version 0.5. It is the first version of the library, so its goal is the bare minimum: Implement all core features and support some of the most important data formats. Image and video formats are not yet in scope of this version.

2.3.1 Supported Platforms

In general, `jMeta` in its current version 0.5 can be used with any platform that also supports Java SE in Java SE 9. However, it needs to be specified for which platforms tests were run explicitly and thus it is proven that specific features work on these platforms. We target to support the following platforms like this:

- Windows 10
- Ubuntu 16.04

2.3.2 Supported Metadata Formats

This version adds support for the following audio metadata formats:

- ID3v1 and ID3v1.1
- ID3v2.3
- APEv2
- Lyrics3v2

2.3.3 Supported Container Formate

This version adds support for the following audio container formats:

- MPEG-1 Audio Layer 3 (MP3)

2.3.4 Supported Input Media

This version adds support for the following data sources for reading data:

- Files
- Java `InputStreams`
- Java byte arrays

2.3.5 Supported Output Media

This version adds support for the following output media for writing data:

- Files
 - Java byte arrays
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2.4 Related Libraries

Here, related Java libraries are treated that are related to what **jMeta** offers. This includes existing productive Java libraries in the area of multimedia, which can be considered as open source competitors to **jMeta** on the one hand, on the other hand, they provide inspiration, possible reuse and also warning examples how to not do it.

2.4.1 Metadata Libraries

A selection of libraries offering access to multimedia metadata as of end of 2017, mostly audio metadata - One source for this e.g. is <http://id3.org/Implementations> and Google:

Library	URL	Supported Formats	Lines of Code	Comments
jaudiotagger (Java)	http://www.jthink.net/jaudiotagger/	MP3, MP4, Ogg Vorbis, Flac, WMA, partly WAV and Real Audio	> 86000	Obscure class and package structure, more than 500 classes, just a few testcases only (?), still actively developed
jLayer and Java-Zoom (Java)	http://www.javazoom.net/index.shtml	MP3, WAV, ID3v1, ID3v2	> 13600	Project targeting MP3 decoding and playing of MP3s
mp3agic (Java)	https://github.com/mpatric/mp3agic	ID3 Tags (1.0, 1.1, 2.3, 2.4), Read-only ID3 v2.2, MP3 low-level reading (incl. VBR)	> 5500	Still actively developed
BeagleBuddy (Java)	http://www.beaglebuddy.com/	Reads and writes ID3 Tags (1.0, 1.1, 2.3, 2.4), Lyrics3v2, Lyrics3v1, APEv1, APEv2, reads MP3 files CBR and VBR, Xing, LAME, and VBRI Header	> 40700	Focus on being “easy to use”; Last version in beginning of 2015
JID3 (Java)	https://blinkenlights.org/jid3/ , https://java.net/projects/jid3lib	ID3v1, ID3v2.2 and ID3v2.3; no ID3v2.4 support	> 30000	Seems to be dead (last version 0.46 in 2005)

Library	URL	Supported Formats	Lines of Code	Comments
Javamusictag (Java)	http://javamusictag.sourceforge.net/	MP3, ID3v1, ID3v1.1, Lyrics3v1, Lyrics3v2, ID3v2.2, ID3v2.3, and ID3v2.4	> 27000	Seems to be the same thing or a fork of JID3
Xuggler (Java)	http://www.xuggle.com/xuggler	MP3, Ogg (Vorbis, Speex), Flac, AAC, ID3v2 and others	-	Decoder and editing API for various audio and video formats
jFlac (Java)	http://jflac.sourceforge.net/	FLAC, VorbisComment	13000	Especially encoding and decoding of FLAC
MPEG-7 Audio Encoder (Java)	http://mpeg7audioenc.sourceforge.net/	MPEG-7	-	Creates MPEG-7 metadata
JAI Image I/O (Java)	https://jai-imageio.dev.java.net/binary-builds.html	EXIF	-	Can read and write EXIF tags from a variety of embedding formats
jmac (Java)	http://sourceforge.net/projects/jmac/	APE, MonkeyAudio	-	
MyID3 (Java)	https://github.com/jkauflin/jkMP3	Audio-Metadaten (ID3)	-	Private engagement, not really to be taken seriously

Library	URL	Supported Formats	Lines of Code	Comments
id3lib (C/C++)	http://id3lib.sourceforge.net/	Audio-Metadaten (ID3)	-	
Mutagen (Python)	http://pypi.python.org/pypi/mutagen/1.12	Audio-Metadaten und -Container-Daten (ID3)		

Table 2.1: Competitor libraries compared

The analysis shows that as of now, the biggest competitors would be jAudio-Tagger, mp3agic, BeagleBuddy and JID (which seems dead) from a metadata perspective and Xuggler as well as jLayer from decoding and editing perspective (if jMeta should ever try to go into this direction).

Drawbacks of Existing Metadata Libraries

Each of the mentioned libraries specializes itself to just a few formats. Saying this, applications are only likely to be satisfied with them if they just need to support those formats. If they do require more, they need to use multiple libraries or extend the existing ones or wait for new features.

Second, the architecture and extensibility of the existing libraries is not convincing. Most of the ID3 libraries e.g. open-heartedly show their very internals, which does not make it clear what really distinguishes “public API” and the private parts - with the usual consequences. Users might rely on private implementation aspects that are changed later, or if it is known that users do so for specific parts, the library evolution might get constrained by this, forcing duplication and maintenance of actually obsolete code.

Another smell of most of the libraries is their sheer volume: While mp3agic still seems humble in terms of “only” about 6000 lines of code (including comments and blank lines), it also does not offer so much features. The other libraries can really be considered big, topping to about over 86000 lines of code for jaudiotagger. These libraries also do not show an obvious structure or core concepts springing to one’s eyes. Maintainability and especially extensibility of such rather “hardcoded” libraries might be not too easy.

Users / software needing to support a broad variety of formats - when using these formats - likely need to implement a kind of plugin mechanism for easier extensibility themselves.

Advantages of jMeta

There are already so much library - why do we need jMeta? If it wants to do things better, a lot of effort is required. Is that worth it?

There are multiple aspects where jMeta can do better than existing libraries:

- jMeta should target container formats on a very abstract and generic level, yet still providing access to the bare details of each format; this can go as far as providing a generic framework for parsing and writing any multimedia container format, may it be audio, video, image or text content.
 - Applications whose core asset is extensibility and rich variety of supported formats do not have to use a dozen or so third party libraries having completely different programming, license and support models as well as performance characteristics - provided jMeta actually offers support for all formats, which is of course not necessarily the case
 - However, at least such applications do not need to reinvent the wheel by coding their own extensibility mechanism, but they just rely on the extensi-
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bility mechanism by using **jMeta**; they can code their own **jMeta** extension to support another format. The goal of **jMeta** is that doing so is easier for such applications than using another 3rd party library to enable support for such a format.

- Still, **jMeta** is not a huge Uber JAR containing all supported formats, but it is modular, where applications just need the core of the library plus an extension module per format they want to support.
- **jMeta** also targets to offer an easy to learn and use API for such applications.

All that said, the benchmark for how good **jMeta** really is, is as follows:

- It should be easier to use as BeaglyBuddy which claims to be easy to use
- It should be much less complex (means to say: much less classes and lines of code) than jaudiogetter or BeagleBuddy
- Supporting a new, averagely complex, previously unsupported data format does not take an experienced Java programmer (who never did write a **jMeta** extension) more than one man day!
- Likewise, we also want to compete with other libraries in terms of performance: **jMeta** must at least be comparatively fast as the most important competitors for Java, ideally of course it should be faster

2.4.2 Multimedia Libraries

Can **jMeta** cooperate, i.e. reuse or integrate with existing Java multimedia libraries? This question is not really answered in the following sections. The general answer is: No, **jMeta** will neither be some kind of plugin for the existing libraries, nor will it depend on them, nor will it consider the requirement of easy cooperation or integration with this libraries in any way.

For now, we just want to see what is existing in the wild in terms of Java multimedia libraries, mostly audio and video. This evaluation might point future directions for **jMeta** if it goes to probably superseding or extending the existing options.

All or most of the multimedia libraries listed here seem to offer some way of reading or writing multimedia metadata. But mostly, these libraries focus on streaming, playing, encoding and decoding on high performance. The metadata is just a side-product that seems to be offered most of the time via a rather inconvenient API.

Xuggler and Humble Video

Cite from the ebsite: “Xuggler is the easy way to uncompress, modify, and re-compress any media file (or stream) from Java. ... Xuggler allows Java programs to decode, encode, and experience (almost) any video format.”

It does not seem to be based on GStreamer or FFMPEG or JMF. It does not become clear which formats are supported by the library.

The official Github page of Xuggler states it is outdated and humble video should be used. The name suggests Humble Video is mainly focused on video, not audio such as MP3.

JLayer

Cite for JLayer: “JLayer is a JAVA library that decodes, converts and plays MP3 files in real-time. JLayer supports MPEG 1/2/2.5 Layer 1/2/3 audio format. JLayer doesn’t need JMF.”

It does also not seem to rely on GStreamer or FFMPEG but seems to be a “native” Java only implementation for encoding and decoding MP3 audio data.

Java FX and GStreamer

JavaFX is an Oracle library to support with platform-independent handling of user interfaces as well as multimedia content. So, it is not only “GUI” that actually replaces Swing and AWT, but also handling of images, video and audio content as well as even 3D graphics, thus somewhat replacing the obsolete Java 3D.

So it has a quite broad usage space. However, there is also some muttering that it might be dead or at least does not have a bright future.

For local GUIs it might still be a good option. What about multimedia playback, editing etc.? The media engine part of JavaFX is based on GStreamer. GStreamer is an extensive audio and video library written in C and nevertheless platform-independent by being ported to a lot of platforms and architectures. It also supports reading and writing tags, but in a very specific generic API which does not offer too much comfort. But how does this go together with JavaFX? First of all, GStreamer also offers a Java binding, i.e. can be called from Java applications where any installed GStreamer plugins for formats can also be used. This interface is most probably also used by JavaFX. JavaFX then offers a class named `MediaPlayer` with a method chain `getMedia().getMetadata()`, just returning a map (String to Object). There seems to be no documentation which objects might get returned. Furthermore it just seems to be read-only, there does not seem to be any way of writing tags with JavaFX.

FFMPEG

FFMPEG can be seen as direct competition to GStreamer, being a platform-independent framework and toolkit to record, convert and stream audio and video content. It is also written in C/C++ and supports a variety of formats.

JMF and FMJ

The Java Media Framework is an official Java library that is delivered with J2SE *desktop* technology. The most recent version is 2.1.1e and dating back to 2001. Which actually means: It is dead. However, we still consider it a bit here. JMF

can be used by Client as well as server applications. JMF was packaged with an MP3 decoder and encoder until 2002, but removed due to licensing issues. Since 2004 there is an MP3 playback only plug-in.¹

FMJ is an Open Source Alternative, that is API compatible: <http://www.fmj-sf.net/>. But FMJ also seems to be dead.

JMF comes in four JAR files:²

- JMStudio: Simple multimedia player application
- JMFRegistry: An application to manage different JMF settings and plug-ins
- JMFCustomizer: Allows creation of a simple JMF JAR file containing only those JMF classes needed by the client application
- JMFInit: Initializes a JMF application

JMF contains platform-specific *performance packs*, i.e. optimized packtes for Linux, Solaris or Windows.

Features: JMF deals with time-based media. The JMF features can be summed up as follows:³

- Capture: Read multimedia frame data of a given audio or video signal and encode it into a specific codec in realtime.
- Playback: Play multimedia data, i.e. display videos on screen or play music on audio output devices.
- Stream: Access multimedia streams
- Transcode: Convert media data of a given codec into another codec without first needing to decode

Criticism: [?] summarizes some negative feedback for JMF:

- A lot of codecs such as MPEG-4, MPEG-2, RealAudio and WindowsMedia are not supported, MP3 only as plug-in
- No maintenance or extension by Oracle, it is dead
- No editing functionality, i.e. modification of multimedia content
- Performance packs only for just a few platforms

¹see [?].

²Siehe [?].

³Siehe [?].

Basic concepts of the API Reading of multimedia data is abstracted using `DataSources`, while output goes to `DataSinks`. No specifics of supported formats are provided for direct API access, they can just be played, processed and transcoded, while the latter is not supported for all formats. A `Manager` class is the primary API for JMF users.⁴

The API documentation shows that JMF is quite complex and essentially time- and event-based.⁵ It offers possibilities to read raw binary data via a method `read` of `PullInputStream`. However, JMF controls processing starting at the source, i.e. from a file or stream.

JavaSound

JavaSound is Oracle's sound processing library. It has some things in common with JMF, but can be considered quite low-level, as it also offers modification functionality for audio data. It also supports MIDI devices.⁶ It can also be considered dead, unfortunately.

Basic concepts: JavaSound essentially offers the classes `Line` (representing an element in the audio processing pipeline), the derived classes `Clip` (for playing audio data) and `Mixer` (for editing audio data). The library can read from streams, files as well as in-memory bytes. It also offers "transcoding" functionality to convert between different formats.

⁴see [?].

⁵See [?].

⁶See [?].

3 Basic Terms

Here we define basic terms used throughout the whole design concept.

3.1 Metadata

Metadata in this document is short for digital metadata that are not necessary to parse the actual described (audio, video or image) data. Metadata semantically and structurally describes other data. The goal of **jMeta** is especially reading of metadata for audio and video data sets, e.g. title, artist etc. The structure of metadata is defined by a METADATA-FORMAT.

If it is specifically about technical metadata needed to parse a data structure, e.g. in the container header, we call it *Parsing Metadata*.

3.2 Data-Formats, Metadata-Formats and Container-Formats

A DATA-FORMAT defines the structure and interpretation of data: Which bytes or characters of which value and in which order have what kind of meaning? Usually, a data format describes how a consecutive block of bytes (i.e. a DATABLOCK) is built up by a number of so-called FIELDS or child DATABLOCKS.

METADATA-FORMATS are data formats that define the structure of digital metadata. Examples include: ID3v1, ID3v2.3, APEv1, MPEG-7, RDF/XML, Vorbis-Comment and others.

CONTAINER-FORMATS are a more general form of DATA-FORMAT optimized for storing, transporting, editing and seeking multimedia PAYLOAD data. Examples are: MP3, Ogg, TIFF, QuickTime, JPEG 2000, PDF and others. Metadata formats can also be considered as container formats.

An example for other DATA-FORMATS is HTML. You can argue that it is neither a METADATA-FORMAT nor a CONTAINER-FORMAT. XML is a DATA-FORMAT that itself can be used to define further XML DATA-FORMATS. Some XML DATA-FORMATS are METADATA-FORMATS, e.g. MPEG-7, MPEG-21 or P_Meta.

3.3 Binary vs. Textual Data Formats?

The terms *binary format* and *textual format* are sometimes rather misleading and not very clear to distinguish. A textual format is clearly a format where all information, might it be parsing metadata or content, must be interpreted as text of a specified encoding. So you basically need text processing for this. In contrast, binary formats might use chunks of bytes that need to be interpreted

as integers or other numbers according to a specific byte order. However, there are also formats such as Lyrics3v2 which fully consists of textual characters, no integers in the way. Nevertheless, it has a lot of similarities with e.g. ID3v2.

So a better distinction is:

- CONTAINER-FORMATS in **jMeta** are data formats that define payload lengths in separate parsing metadata fields; their most common representation is as binary format, but some formats such as Lyrics3v2 use text characters also for parsing metadata. Saying this, all common METADATA-FORMATS are seen as specific CONTAINER-FORMATS.
- In contrast to this, a *textual delimited data format* is a format where all bytes need to be interpreted as characters, and thus also parsing metadata is represented as strings. In addition, these formats do not have fields that specify lengths, but they use metadata delimiters to mark start end (optionally) end of a field. Examples are XML, HTML, JSON, LaTeX or YAML.

3.4 Transformationen

A DATA-FORMAT may define TRANSFORMATIONS. A TRANSFORMATION describes a way how read or to be written data needs to be transformed to fulfill specific needs. You can envision this as kind of encoding of the data. In contrast to the fixed data format specification which describes in detail how binary data is coded and needs to be interpreted, TRANSFORMATIONS are optional features that are dynamically applied to certain areas of the data. Partly, TRANSFORMATIONS can also be defined by users of the data. Examples are the TRANSFORMATIONS defined by ID3v2: Unsynchronization, Encryption and Compression.

3.5 Datablocks

A DATABLOCK is a sequence of bytes that together form a logical unit (i.e. an object or entity with a specific meaning) in terms of the underlying DATA-FORMAT. Each DATABLOCK belongs to exactly one DATA-FORMAT. It can be assigned a current length in bytes. There are several concrete types of DATABLOCKS that are described in the following sections.

The actual detailed meta model of DATABLOCKS is defined later in the detailed **jMeta** design sections of this document for **DataFormats**.

3.5.1 Container: Payload, Header, Footer

An important type of DATABLOCK is a CONTAINER: It consists of zero, one or several HEADERS, exactly one PAYLOAD and zero, one or several FOOTERS. All of these are *child* DATABLOCKS. CONTAINERS are a common concept for container formats: The HEADERS and FOOTERS describe the CONTAINER in terms of its length, size and other properties. The PAYLOAD contains the interesting data, e.g. the multimedia data to be extracted, played or viewed. FOOTER essentially

allow for backward or reverse reading. Most of the DATA-FORMATS specify a generic structure of a CONTAINER, with specific containers with specific purpose, but allowing user-defined new CONTAINERS at the same time, i.e. the format is extensible.

3.5.2 Tag

A TAG is a special CONTAINER whose purpose is to store metadata. It can either belong to a standalone METADATA-FORMAT ore to a more general CONTAINER-FORMAT. Especially audio, video and image metadata formats use this term when talking about such a DATABLOCK in a file or MEDIA STREAM, e.g. the ID3 or APE TAGS. It term originates from “tagging” something with additional meta-information, as you’d attach a label to describe the song, video or image.

The following figure shows the basic structure of a TAG, showing other basic terms:

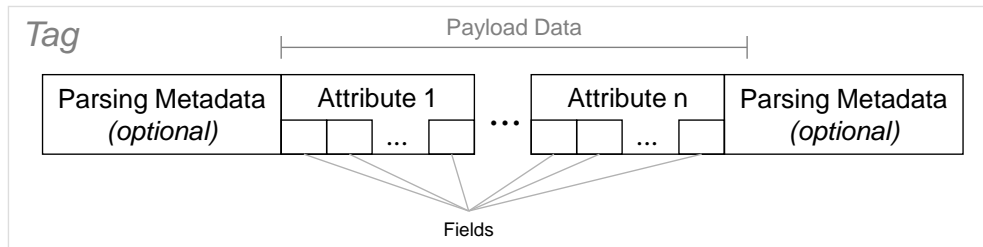


Figure 3.1: Structure of a TAG

The most important parts of a TAG are the ATTRIBUTES.

3.5.3 Attribute

An ATTRIBUT is a part of a TAG that contains the valuable metadata information in a key-value manner. Common examples are artist, title, album, composer etc. of a piece of audio. Often, an ATTRIBUT is also a CONTAINER in a sense that it has a HEADER and PAYLOAD. The HEADER may help to define the type (artist, title, album) of the ATTRIBUT as well as the size of the PAYLOAD. The PAYLOAD contains the actual information in an encoded way, e.g. the name of the artist or title of the piece of audio.

Most of the ATTRIBUTES have a simple main value that can be given. However, there are also more complex ATTRIBUTES that consist of many values in form of child CONTAINERS or FIELDS.

In each metadata format, an ATTRIBUT has a format-specific name, e.g.:

- ID3v1, Lyrics3: Field
- ID3v2: Frame

- APE: Item
- Matroska: SimpleTag
- VorbisComment: User Comment

3.5.4 Fields

A FIELD is a sequence of bits that together have a specific meaning in a given DATA-FORMAT. They form the leafs of the data block hierarchy by ultimately containing the raw data. A DATA-FORMAT describes how a specific DATABLOCK is built up by a specific sequence of FIELDS. A FIELD has a range of possible values and interpretations of these values. Often, one part of the value range is defined as “reserved” to ensure a bit of flexibility in extending the data format. Fields can be fixed-size or of variable length. If they have a variable length, there is often either another field specifying its length, or it is terminated by a specific sequence of bytes or characters.

3.6 Medium

A MEDIUM incorporates both the physical location of the data containing DATA-BLOCKS to read and write as well as the way of accessing this data. It may be a file, a stream or even plain memory.

4 Requirements and Exclusions

Here, all explicit high-level requirements for **jMeta** in version 0.5 are listed, including some explicit exclusions. Exclusions have the meaning that a given feature is explicitly not supported in the current version, and it might be like that also in future versions. No design is made for supporting them. However, it might well be that for future versions of the library, these decisions are changed and previously excluded features are then supported. The question is: Why listing exclusions at all? How far do they go, so will there be an excluding stating that **jMeta** is not an operating system, e.g.? Not that far. Exclusions are only meant to demonstrate lack of features that popped up during the design phase and that could be desirable for **jMeta** users.

In contrast to exclusions, there are potential features such as new metadata or container formats to support. These are neither explicitly excluded nor already planned for future releases.

4.1 REQ 001: Reading and writing of human-readable meta data

jMeta can read and write meta data in a *human-readable format*. That means librar users do not have to read (and interpret) or write raw binary data.

Rationale: This is the general core functionality of the library and must be supported.

A list of supported data formats in the current version is given in “?? ??”.

4.2 REQ 002: Read container formats

jMeta must be at least able to read popular and wide-spread container formats. Writing support is optional.

Rationale: Multimedia meta data is often embedded in container formats or an integral part of their specification. **jMeta** must be able to identify and parse container segments in order to skip them to find the actual start of a meta data tag, or to read their meta data content. Writing is optional because **jMeta** cannot offer encoding functionality, so just writing header information and raw binary data that has been encoded by another library could be offered.

4.3 REQ 003: Fulfill specification of supported data formats

As far as there is an official specification for a supported data format, this specification must be fully supported by **jMeta**, that means all encoding types, optional headers, data transformations (e.g. encryption), padding and other features that are included in the specification, as well as any generic attributes or containers must be readable and writable. It is not necessary that **jMeta** explicitly supports unspecified attribute types that might be popular but not defined by the specification. Such attributes should however be accessible using the generic attribute mechanism offered by most data formats.

Rationale: This ensures that all meta data or container format blocks that adhere to their specification can be read by **jMeta**, and furthermore, that **jMeta** itself generates data that is compliant to the specification. Second, the user can use **jMeta** to use specific data format features in a convenient way without needing to implement this himself.

4.4 REQ 004: Access to raw data via jMeta

In addition to the access to human-readable meta data (“?? ??”), it must be possible to explicitly read and write raw data on byte level. **jMeta** must offer fine-grained access to fields of binary data.

Rationale: Thus, user could implement parsing themselves without needing to use high-level functions. In some rare cases this might be necessary to circumvent some issues with the data written by some other tool. This channel of access is offered to users such that they are not forced to completely skip **jMeta** and implement raw access to supported media themselves, leading to also sync and consistency issues when doing this in parallel to the access by **jMeta**.

4.5 REQ 005: Performance as good as other Java meta data libraries

The performance of **jMeta** on all supported platforms must be as least as good as the performance of competitor Java meta data libraries such as **jaudiotagger**, **mp3agic** or **BeagleBuddy**. A corresponding reasonable benchmark needs to be set up with published code to demonstrate this. It is out of scope to make **jMeta** as fast as high-performance encoding or decoding libraries on C++ basis such as **gstreamer** or **ffmpeg**.

Rationale: **jMeta** targets to be a good replacement for any existing Java multimedia tagging library, and thus it is clear it should at least perform at the same level or even better.

4.6 REQ 006: Fault recognition, fault tolerance, fault correction

In addition to “?? ??”, the library also has to be *fault tolerant* as much as possible. That means that specification violations and incorrect parsing metadata must be detected and - as far as possible - this should not lead to aborting parsing with an error. All violations can be protocolled and ideally automatically corrected (optional, if the user wants it).

Rationale: Old applications or other libraries write data formats sometimes not 100% compliant to the specification. In addition, sometimes specifications are ambiguous or not accurate enough such that alternatives are possible. If there is a chance to still read the data, the user should be able to do it even in the event of faulty data.

4.7 REQ 007: Extensibility for new Metadata and Container Formats

jMeta must be comfortably extensible with new container or metadata formats. As the minimum level, easy extensibility by the library developers must be possible. As the maximum level of extensibility, also end users with programming experience must be able to easily write extensions without too much configuration or boilerplate code.

Rationale: A wealth of already existing meta data and container formats is yet out there. Just the formats with major importance are supported as of now by jMeta. Furthermore, we might expect new meta data or container formats in future. The extensibility mechanism first of all guarantees an easy extensibility by a 3rd party, e.g. lib name users or also other vendors. The extensibility also ensures a longer life time of the library. In the maximum level “Extensibility by end users”, this is a clear differentiation criterion to other libraries that do not offer this level of extensibility. Last but not least, this allows easier maintenance (i.e. extension) by the library developers itself if they decide to bundle additional format support extensions with the core library in a new release.

4.8 REQ 008: Read and write large data blocks

jMeta must be able to read and write large amounts of data efficiently. With “large” we mean that jMeta must support data blocks with maximum size of $2^{63} - 1$ bytes. The length of payloads must be interpreted correctly. Because of “?? ??”, also reading and writing of raw data must be supported without necessarily causing scarcity of available memory, e.g. by chunk-wise reading possibilities. In general, jMeta must use mechanisms to avoid `OutOfMemoryErrors`, unless the user forces this by reading the whole binary data of a block into memory.

Rationale: Especially in the domain of video container formats, payloads or at least files with multiple gigabytes are very common. Yet, still $2^{63} - 1$ bytes as upper limit should suffice for some years or decades to come. Supporting this is for

sure also a plus compared to other competitor libraries that might not explicitly support large files.

4.9 REQ 009: Selective Format Choice

An application using `jMeta` must be able to selectively choose those formats that it wants to support. This is not only necessary for runtime, but also for the library extension packages it wants to use.

Rationale: Audio applications do not need extensions for video or image formats. Applications can minimize the runtime and memory (both HDD and RAM) overhead by choosing as few extensions as really needed.

4.10 EXCL 001: Reading media streams

`jMeta` in its current version 0.5 does not explicitly support reading meta data or containers from media streams consumed e.g. from the internet. It might offer optional starting points for this by e.g. in principle supporting access to `InputStreams`. But it does not offer an explicit API or examples, nor do we create specific test cases for media streams. Furthermore, media stream specific tag formats such as `IcyTag` are not in scope of the current library version.

Rationale: Combined applications such as recorders or players whose main task is playing make more sense for this, as they are not only interested in meta data. Supporting streaming media in addition might overly complicate the design and API of the current library version, so this might be left for future versions to come.

4.11 EXCL 002: Support for XML meta data

There are also XML meta data formats out in the wild. They are not used very often for multimedia meta data, but cases exist. `jMeta` does not support reading and writing of XML meta data format in the current version, but only specializes to binary formats. **Rationale:** Binary meta data formats still seem to rule the scene due to their compactness. It would be absurd to design a generic library that can parse both XML and binary formats with the same code, as XML is usually efficiently read and written using streaming parsers and validated using schemas, where excellent APIs already exist. It is left for the future to offer built-in support for XML metadata in `jMeta` via an easy API.

4.12 EXCL 003: No user extensions for jMeta media

Extending `jMeta` with new media to support - on top of the out-of-the box supported media - is not supported.

Rationale: The mechanisms available should cover 80 to 90 percent of the use cases. Extensibility by new media might increase the complexity of `jMeta` itself

and the extension mechanism in specific, without adding real-world use cases really needing it. It is currently not clear which media beyond streams, byte arrays and files might be candidates for extending `jMeta`. Should new media make sense in future, a new core release can be created to also support media extensions.

4.13 EXCL 004: `jMeta` is not a (high-performance) encoder or decoder

Even though it would be tempting, `jMeta` currently is not a encoder or decoder for multimedia content, especially not a high-performance one. It is not the target to extend this library to be that, but you never know. However, it can be the basis for a Java-based encoder or decoder pack, as it offers bit-wise parsing and other raw data access functionality for container formats. But there is no specific design into that direction, especially no means for specific multi-threading built-in support, means users are left to leverage multiple cores themselves as necessary.

Rationale: Audio, video and image formats as well as corresponding encoders are extensively complex. A huge variety of great libraries and tools is already existing. For high-performance needs, especially GStreamer and FFmpeg seem to rule the scene. `jMeta` can never compete with that and it would mean to *just stop* the project when trying to do just that. Instead, we can offer the primitives for Java developers who might want to develop their own (probably not very high-performance) codecs. Maybe in the very far future, `jMeta` could offer integration functionality to plug-in existing codec implementations somehow conveniently, such that it can be used on its own as a basis for Java transcoding and player/viewer applications.

4.14 EXCL 005: Scanning for data formats not supported

`jMeta` in its current version does not yet support scanning a chunk of bytes for the occurrence of a data format. This functionality might be added in later versions.

Rationale: This simplifies implementation of data format discovery for now. Furthermore, support for streaming media (where scanning is most likely to be necessary) is out-of-scope, too, see “?? ??”.

5 Reference Examples

To proof conformance with the **jMeta** requirements, a set of (mostly) real-life examples for all supported formats is used. The examples are used to illustrate design decisions and also to verify them. In this chapter, these examples are presented for short. The concrete detailed structures of each data format is described in [?].

Note that the sizes of the data blocks in the following illustrating example figures do not have any specific meaning.

5.1 Example 1: MP3 File with ID3v2.3, ID3v1.1 and Lyrics3

The following figure shows the first example, an MP3 file with three TAGs, ID3v2.3, Lyrics3v2 and ID3v1.1. All are located at the end of the file:

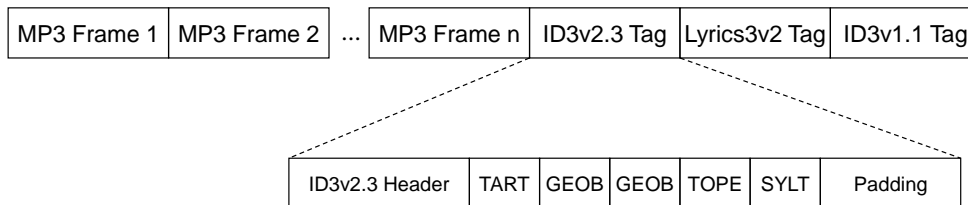


Figure 5.1: Example 1: MP3 file with two tags

The ID3v2.3 tag has several frames, including two **GEOB** frames. Furthermore, it has some padding within. Each of the MP3 frames corresponds to the MPEG-1 elementary stream audio format.

5.2 Example 2: MP3 File with two ID3v2.4 Tags

The following figure shows the second example, an MP3 file with two ID3v2.4 TAGs, one at the beginning, the other one at the end of the file:

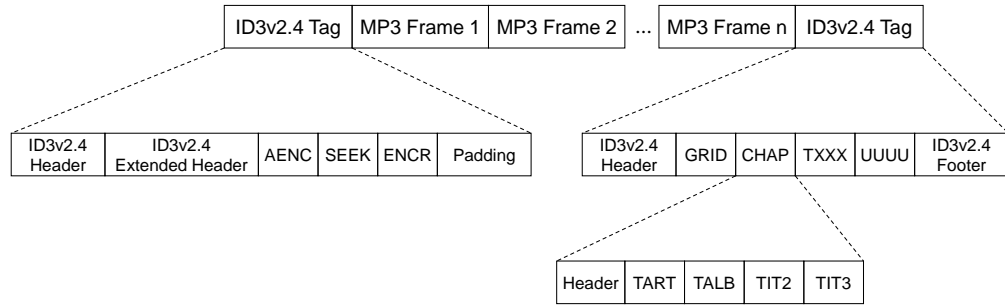


Figure 5.2: Example 2: MP3 file with two ID3v2.4 tags

The two ID3v2.4 tags are virtually connected by a **SEEK** frame. Both have several specialties described in [?].

5.3 Example 3: Ogg Bitstream with Theora and VorbisComment

The following figure shows the fourth example, an Ogg bitstream that contains Theora payload data with a corresponding vorbis comment:

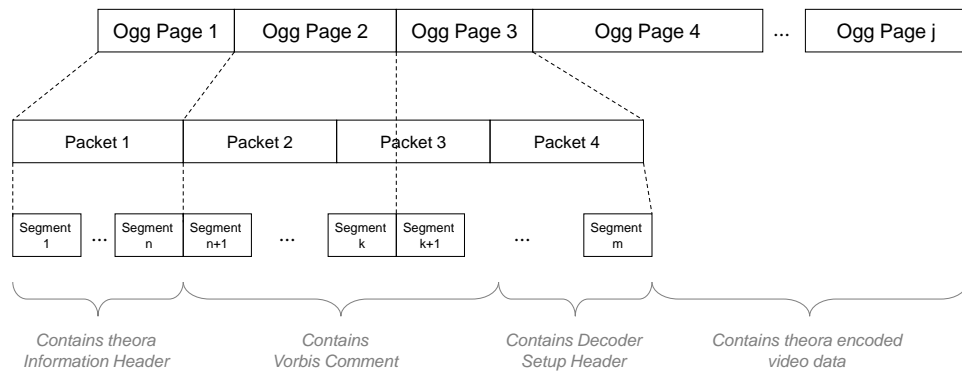


Figure 5.3: Example 4: Ogg Bitstream with Theora and VorbisComment

The example is complex on first sight. In an Ogg bitstream, physical and logical structure are not necessarily the same. The physical structure is built by pages, packets and segments, while the logical structure is the structure of the wrapped data. We took theora video data as an example, but it's basically arbitrary, as the codec does not really matter. What matters is where the Vorbis Comment, one of the supported data formats, is stored. This unfortunately depends on the embedded codec. In this example, the vorbis comment starts in the second page and spans over two packets. The second of these packets spans over two Ogg pages.

Part II

Architecture

In this part, the high-level structure of `jMeta` is defined. This is done in two refinement levels:

- The technical architecture shows the technical environment and parts of `jMeta`
- The functional architecture shows and describes the functional components of `jMeta`

6 General Design Decisions

This chapter presents all general design decisions that have a big impact on the overall architecture. It is usually very hard to revert any of those design decisions. As they generally affect the whole library, they have no direct relation to specific requirements, but build the framework to be able to fulfill those requirements.

6.1 Using Java SE 9

DES 001: jMeta is based on Java SE 9

jMeta is developed based on the “latest update” of Java SE 9. Migrating jMeta towards newer Java versions is considered on a regular basis.

Rationale: Java as programming language is well established and in wide use by a big number of world-wide developers. Furthermore, Java is platform-independent. In principle, porting jMeta to other platforms, to Java ME or to Java for e.g. smart phones (e.g. Android) should be less effort than porting an equivalent C/C++ library. The author has long experience with Java programming, so this should facilitate productivity. The competitor libraries in the Java environment are good, but jMeta can be even better. The current (state 28th of November 2017) version Java SE 9 is used, because we want to create a modern library using the newest language features, libraries and programming models in the Java world. Furthermore, it is guaranteed that we get longer support for this Java version than for the previous ones.

Disadvantages: Applications being based on Java 7, Java 8 or older are not supported by jMeta. This would especially effect Java EE applications where it seems not yet possible to use Java SE 9 for some application servers.

6.2 Use of 3rd Party Libraries

DES 002: jMeta uses as few 3rd party libraries as possible

jMeta uses mostly only the libraries offered by Java SE. In the productive code of the library, no dependencies to other 3rd party libraries are used except for logging.

Rationale: Additional dependencies could create overhead or question marks when it comes to building, shipping, testing or licensing. Furthermore, you might at one point in time reuse a well-tested and documented 3rd party library, and in the next year, it might get abandoned or bugs show up that do not get fixed. We ensure that bugs are just bugs that we can fix ourselves and we do not depend on other library developers.

Disadvantages: Sometimes this might mean “do it yourself” instead of reuse.

6.3 Component-based Library

DES 003: jMeta is component-based

jMeta is consisting of so-called components as defined in the follow-up sub-section. Each component has a well defined responsibility, a public API that can be used by all other components as well as a private implementation part that must not be used by other components.

Rationale: Structuring a big library into modular parts helps to manage the complexity of the library during planning, design, implementation and maintenance. By using a reasonable component structure, clear responsibilities, dependencies and decoupling of unrelated parts is achievable. Changes in implementation parts of a component do have less probability to affect other components, as they do only depend on the API of the component.

Disadvantages: Probably a little bit more overhead due to sometimes following “conventions” where an easier structure would be possible

6.3.1 Definition of the component term

A *component* in jMeta is a *self-contained software module with clearly defined task*. It offers *services* that can be used via a *well-defined interface (API)*. These services are used both by the users of jMeta as well as by other components of the library.

The following figure shows a component in jMeta schematically:

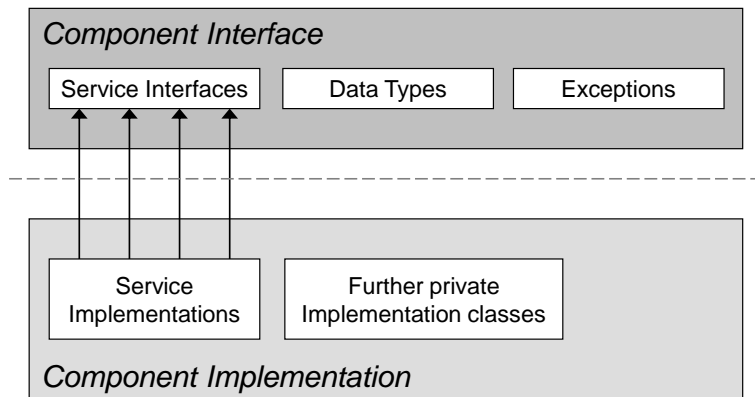


Figure 6.1: Structure of a component

A component interface consists of one or more Java interfaces, the so-called service interfaces, exceptions as well as data types.

- *Service interfaces* offer the functional API of the component, each method corresponds to a service of the component. A single Java interface is use for a clear defined subtask of the component. Some service interface have factory methods that provide access to further service interfaces of the component.
- *Data types* are (mostly) constant Java classes that are usually used to hold data for service input and output parameters. Usually they are value types, but sometimes they are also representing a specific entity such as a medium.
- *Exceptions* are functional errors. These are mostly unchecked exceptions signaling a specific error condition, in rare cases we use unchecked exceptions if the caller can handle the situation.

In addition, a component might have data ownership for a specific kind of relevant data. This means that this data can only be read or manipulated via using the public API of this component, and by no other means.

6.3.2 Design Decision for using components

DES 004: Components decoupled by a light-weight service locator pattern

To ensure low coupling between components, they must know each other only via their interfaces. A component can obtain a reference to a service implementation without directly depending on the implementation class itself by using a service locator pattern, specifically the Java SE **ServiceLoader**.

Rationale: We want to decouple components. To achieve this, it must not be allowed for components to instantiate service implementations of other components directly, because service implementations are not part of the public interface of a component. A service locator exactly offers to obtain a service implementation reference without depending on it. For a library such as **jMeta**, neither JavaEE nor Spring or CDI dependency injection are valid options. Google Guice is not chosen because of ?? . Java SE **ServiceLoader** is a very easy to use and fully sufficient facility.

Disadvantages: No disadvantages known

DES 005: Singleton Components

Each component offers a single *main API service interface* that itself offers access to any other service interface the component provides.

From a runtime perspective, each main API service interface thus has a single (ideally stateless) implementation, which you could refer to as “singleton”.

Rationale: A main API service interfaces is purely procedural. There is just one implementation of component allowed at runtime. We just need one instance for each component at runtime, as it is stateless. This safes memory and initialization overhead.

Disadvantages: No disadvantages known

DES 006: API and Implementation Layers

Each component offers its services, exceptions and data types on an *API layer*. This layer is a horizontal layer cutting through the whole library and representing the public API of the library. Also per component, it is the public API of a component. Compile-time dependencies from user code to `jMeta` as well as between components are only allowed via this API layer. The implementation layer depends on the API layer and contains any private classes. An outside component and user code is not allowed to have a compile-time dependency to implementation layer classes

Rationale: Clear separation of concerns and information hiding. What implementation layer classes do is their private business, nobody from the outside world is allowed to directly work with and depend on them.

Disadvantages: No disadvantages known

DES 007: No further subdivision of implementation layer

It is not necessary to further subdivide the implementation layer into sublayers. Instead, each component can have an arbitrary package structure in its implementation layer code.

Rationale: We do not see any reasonable sublayers for `jMeta` that could be relevant for all components of the library.

Disadvantages: No disadvantages known

6.4 Development Environment

DES 008: Development Environment

As development environment, we use Maven, Eclipse and SubVersion for `jMeta`.

Rationale: Amazing toolset that leaves not much else to wish for.

Disadvantages: No disadvantages known

6.5 Multithreading

DES 009: jMeta is not thread-safe

jMeta is not a thread-safe library and (usually) does not use any Java APIs that are thread-safe.

Rationale: In some cases, thread-safety means performance degradation due to synchronisation points. In addition, it almost always means higher complexity. In a non-functional programming language such as Java with a lot of mutable state, it is hard to get thread-safety right. If users need to work with jMeta in multiple threads, they have to ensure thread-safety themselves.

Disadvantages: No disadvantages known

6.6 Architecture

DES 010: Architecture of jMeta

jMeta's architecture adheres to the technical and functional architecture as described in chapters “?? ??” and “?? ??”.

Rationale: See architecture discussion in detail in the next chapters

Disadvantages: See architecture discussion in detail in the next chapters

7 Technical Architecture

The technical architecture allows functional components (i.e. components offering the actual functionality of `jMeta`) to operate in their technical infrastructure. In addition, the technical architecture consists of the layers of each component. We have a component-based layered architecture as defined in “?? ??”. The layering has already been defined in the design decisions ?? and ??. Thus, we just concentrate on a brief overview of the technical infrastructure and the technical base components here.

7.1 Technical Infrastructure

The technical infrastructure describes the environment required to work with `jMeta` as shown in the following figure:

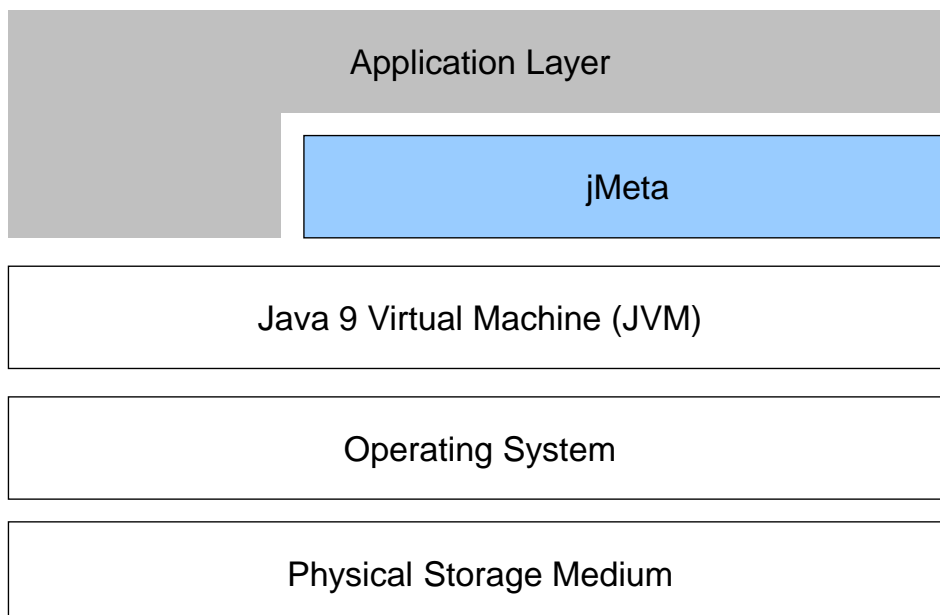


Figure 7.1: Technical infrastructure of `jMeta`

The technical layers can be interpreted as dependency and communication structure. The application layer is based on `jMeta`, while `jMeta` as well as the

application layer uses Java SE 9. Thanks to this, the library and its using applications are platform-independent and can be used on any platform currently supported by Java. Java SE 9 accesses the operating system, which offers services to access the physical medium.

However, `jMeta` is only developed for Java SE 9 and cannot be used by applications requiring earlier versions.

Each layer only accesses the neighbouring layer.

7.2 Technical Base Components

Technical base components are just helpers providing a framework and basis for the functional components of `jMeta`. With “functional” we are referring to the tasks of the library, i.e. reading and writing metadata and reading container data. The necessary base components are:

- Logging
- Service-Locator to achieve a component oriented architecture
- Utility
- Maintaining extensions

Details of these components are given in “?? ??”.

8 Functional Architecture

The functional Architecture defines the structuring of the library into functional units, so called functional components according to “?? ??”. For completeness, also the dependencies of the functional components with the technical base components is shown in functional architecture diagrams. In addition, the functional architecture also contains basic functional design decisions influencing the direction of the whole library.

8.1 Fundamental design decisions of jMeta functionality

The following design decisions are fundamental for the whole library and mostly have a cross-cutting effect. They form the basis for the functional design of all functional library components.

DES 011: Generic reading and writing according to a format specification

Reading and writing different CONTAINER-FORMATS is done generically by a central component. A format specification describes which features a CONTAINER-FORMAT has, especially how its DATABLOCKS are structured and how they need to be interpreted. Thus it can be seen as a kind of generic parsing (and writing and validating) instruction of the data format.

Rationale: According to the document [?], CONTAINER-FORMATS and their more specific children METADATA-FORMATS have a lot of things in common, in specific the way how they are structured and parsed. These common properties can be described using a generic format specification. Instead of rewriting specific parsing, writing and validation code for each new supported CONTAINER-FORMAT, all of them can be implemented by the same generic parsing, writing and validation code that is based on the generic format specification.

This way, in the end, there should be less code, less duplication, and also a higher quality of code as it can be tested easier.

This design decision also essentially enables to meet the requirement “?? ??”, as this way it is much easier to extend the library with new CONTAINER-FORMATS.

Disadvantages: The generic code can probably not be written to support any and every possible future or existing CONTAINER-FORMAT out of the box. Thus there might still be changes required in future. We have to carefully avoid to add format specific code to this generic core code, such that it not again becomes a complex monster, where in the end it would have been better to rewrite code for each extension. The latter risk exists if the formats do not have as much in common as assumed by this design decision. This can make testing or even understanding the generic code extensively hard.

In addition, there might be some small performance drawbacks, as generic code can of course not as heavily and specifically tuned as format-specific code can.

To mitigate the possible disadvantages of ??, we define the following design decision:

DES 012: It is possible to override generic reading and writing in extensions

jMeta CONTAINER-FORMAT extensions are allowed to override specific parts of reading, writing and validation code to adapt them to specifics of their data format.

Rationale: This minimizes the disadvantages of ???. In special cases, this can lead to easier or faster implementations. Furthermore, using such hooks, it might not be necessary to blow up the core code by specifics that are actually only relevant for one or just a few concrete formats. Those specifics can then be implemented in the corresponding format's code.

Disadvantages: It might not be easy to clearly define and communicate which parts can be overridden and which not.

To also mitigate the risk of too big and complex generic core code, we define:

DES 013: Format specifics are only defined in extensions

Any special features of a CONTAINER-FORMAT or METADATA-FORMAT that are no common properties/features of a lot of container formats are only defined in the corresponding extensions. That said, the generic core code never directly depends on specifics of data formats, neither syntactically (e.g. compile-time dependences to format extension; by using if-then-else to do things different depending on the current format) or semantically (e.g. implementing a special feature for just one format, but claim it is “generic”). Only if a special feature of a format can be brought down to a common denominator for all supported formats, it is implemented generically.

Rationale: There is a strict separation between core implementation and format specifics, which allows better handling of the overall complexity.

Disadvantages: No disadvantages known

DES 014: High-level and low-level API

The generic, format-independent core parts of `jMeta` form a kind of *low-level API* for parsing, writing and validating data of a `CONTAINER-FORMAT` quite generically down to the byte level. The end user of `jMeta` is allowed to use this part of the library directly.

However, on top of that there is a more convenient and format-specific *high-level API* the user usually uses to access data.

Rationale: According to ??, `jMeta` has a generic format-independent core part. This could be made private, not to be used by end users of `jMeta`. However, this would not make it possible to fulfill “?? ??”. At the same time, format-specific extensions need to reuse the core part. We do not want to throw the core part together with these higher-level extensions.

Disadvantages: No disadvantages known

8.2 Extensions

An *extension* summarizes format-specific content for a given data format, or multiple variations of the data format. Per data format defined by the extension, these are (compare ??, ?? and ??):

- A data format specification which defines `DATABLOCKS` and their structure and relations.
- An end-user high-level API to comfortably work with the data format, e.g. for creation of `DATABLOCKS` etc.; this API especially holds the data format identifier which the user can use to work with other components.
- Implementation extensions that might override generic parsing, writing and validation.

As a ground rule, every single extension can but should not define extension code for more than one data format. It might expose more data formats only if these two data formats heavily overlap and are based on each other, e.g. `ID3v1` and `ID3v1.1`. This is to fulfill the requirement “?? ??”.

Moreover, distinct extensions should never depend on each other, but only on the `jMeta` core to avoid any confusion for the user.

8.3 Functional Architecture Components

The following component diagram shows the functional and technical base components of `jMeta` and their dependencies. An arrow indicates a necessary compile-time dependency. For each format extension, there is exactly one component.

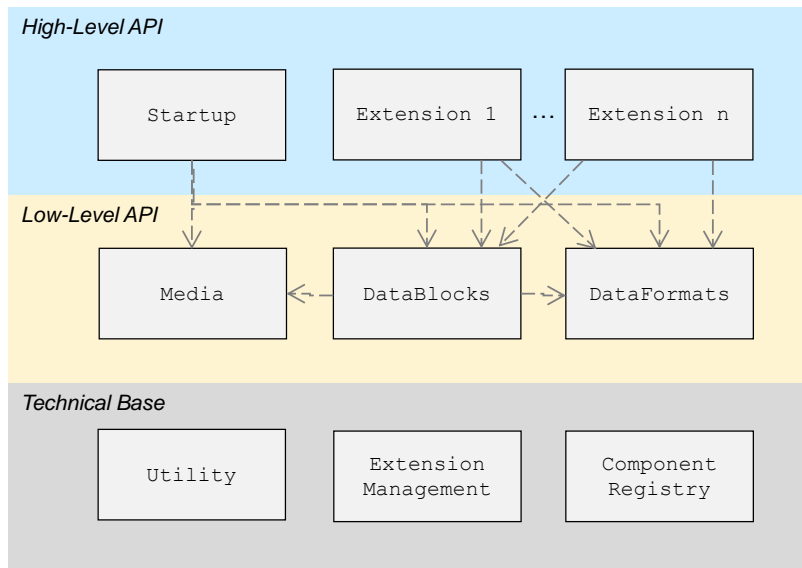


Figure 8.1: Component diagram of jMeta

You can identify a kind of “component layering” (not to be confused with the architecture layering, it is just a grouping of components). At the lowest level we have the technical base components, then followed up by the low-level API components and the high-level API (according to ??) on top.

Dependencies to the technical base components are not shown in this diagram, as anybody can (and usually does) depend on them. Note that, due to this, we only list dependencies to technical base components in the component characteristics sections to follow if these dependencies are important and need to be specifically mentioned.

All the components displayed are just briefly described in the next sections.

8.4 Component Characteristics

All components are briefly described by a common structure that looks like this:

Component Name: The name of the component.

Task: The main responsibility of the component.

Controlled Data: Any data the component has ownership for.

Depends on <Component Name>: Occurring multiple times, once for each component it depends on with a reason why this dependency is necessary.

8.5 Component EasyTag

Component Name: EasyTag

Task: The entry point for every user to `jMeta`. It loads all extensions and provides access to all other components of the library.

Controlled Data: Keine.

Depends on ExtensionManagement: `EasyTag` loads all extensions by using this component.

Depends on ComponentRegistry: `EasyTag` uses this component to instantiate all component singletons the user asks for.

8.6 Any extension component

Component Name: Arbitrary, usually the name of the data format the extension implements

Task: Offers the high-level API to work with the data format it implements.

Controlled Data: None.

Depends on ExtensionManagement: It needs this dependency to the extension framework because it is an extension.

Depends on DataFormats: It implements a specific data format, does needs to depend on data format base classes defined in `DataFormats`.

Depends on DataBlocks: Extensions are allowed to override reading, writing and validation of data, thus they can depend on `DataBlocks`.

8.7 Component DataBlocks

Component Name: `DataBlocks`

Task: Implements generic reading and writing of `CONTAINER-FORMATS` and `METADATA-FORMATS` according to their format specification.

Controlled Data: Accessing raw `CONTAINER` or `TAG` data on byte or bit level is only allowed by using this component.

Depends on Media: `DataBlocks` must read and write data portions from an external medium, for this it uses `Media`.

Depends on DataFormats: Needs access to the data format specification by depending on `DataFormats`. This is necessary to implement the generic reading and writing.

8.8 Component DataFormats

Component Name: `DataFormats`

Task: Manages the generic data format definitions for `CONTAINER-FORMATS`.

Controlled Data: Has control over the data format specifications.

Depends on: No other components except technical base components.

8.9 Component Media

Component Name: `Media`

Task: Offers primitive to read and write all supported physical MEDIUM. Hides medium access differences for other components.

Controlled Data: Accessing data of external media is only allowed by using Media.

Depends on: No other components except technical base components.

8.10 Component ExtensionManagement

Component Name: ExtensionManagement

Task: Manages all extensions of jMeta generically, i.e. loading and verification.

Controlled Data: Descriptive data for extensions can only be obtained using ExtensionManagement.

Depends on: No other components except utility.

8.10.1 Component ComponentRegistry

Component Name: ComponentRegistry

Task: Implements the service locator functionality to enable components to depend only on the API of another component.

Controlled Data: Only via ComponentRegistry, access to metadata about components of jMeta is possible.

Depends on: No other components except utility.

8.10.2 Component Utility

Component Name: Utility.

Task: Offers diverse (technical) cross-functional stuff that needs to be used by most if not all other components, e.g. helper functionality to implement design-by-contract.

Controlled Data: None.

Depends on: No other components.

Part III

jMeta Design

This part defines the design of the components per subsystem, based on the components and subsystems defined in the previous chapters.

9 Cross-functional Aspects

This chapter covers cross-functional aspects of the **jMeta** design which are not implemented in a single component only but cover the whole library, i.e. these are the so called *cross-cutting concerns*.

9.1 General Error Handling

Here, the general approach to error handling is presented. The basics of error handling are only slightly touched here. This section deals with errors that occur during the runtime of **jMeta**. This section does not describe concrete errors of concrete **jMeta** components. Instead, it defines guidelines used throughout **jMeta**.

9.1.1 Abnormal Events vs. Operation Errors

According to [?], errors can be categorized as follows:

- *Abnormal events*: Events that should only rarely occur and require a special handling.
 - Connection to a server is lost
 - A file operation fails
 - A configuration file, table or any other data item that is expected to exist is not present
 - An external memory or the main memory lacks enough space to handle a request
- *Operation Errors*: An operation can be completed with success or error. Errors can be distinguished from abnormal events in that they are more likely or even expected to happen and that they can be handled directly by the caller of the operation.
 - Invalid user input, e.g. an input value is out of range
 - The state to call the operation is not yet established internally, the method must be called later or multiple times. This is e.g. very common in embedded programming.
 - The previous case can be seen as a precondition that must be established to call the method. However, a precondition is rather something *the user can establish* from the outside by e.g. calling another operation before.

9.1.2 Error Handling Approaches

There are several different approaches to deal with the error categories, that are sometimes even used intermixed:

- *Error codes:* In procedural system programming APIs, e.g. the Linux or Windows API, operations usually return error codes. One code is reserved with the semantics "no error". Other codes are defined with special semantics that describe the kind of error.
- *Error handlers:* Some APIs allow to specify error handlers that are called back by the system in case of an error. The error handler, as the name suggests, needs to do something with the error condition. For this to work, it usually needs context information such as the kind of error that occurred.
- *Exceptions:* Mostly in object-oriented programming, exceptions are objects that usually represent errors and can contain context information about the error. They are thrown by an operation which changes the order the operations code is executed. They need to be caught somewhere in the call hierarchy of the operation, either within the operation itself or in one of its active parent callers. If they are not caught, they terminate the process the operation was running in. So throwing is the raising of the error, thereby assuming that further code in the operation is not allowed to be executed as the current state prohibits that. Catching equals handling the error. Some object-oriented programming languages like Java and C++ distinguish between checked and unchecked exceptions.

Each approach has pros and cons. In object-oriented programs, usually the exception mechanism is used.

9.1.3 Error Handling in jMeta

jMeta as an object-oriented application uses the exception error handling approach. There are no error handlers in jMeta. The return values of jMeta methods will never have the semantics of an error.

General Exception Usage Guidelines

For each special category of error, a single exception class is defined. The exception class has a meaningful name that describes the error category. The name must end with **Exception**. The class stores as much context information as necessary to be able to handle the exception. This information is passed to the constructor of the exception when throwing it. It can be requested using special getters in the exception class.

jMeta does not forward checked exceptions that are coming from internally used 3rd party libraries, such as the Java standard API. Instead, such errors are always wrapped by an own, meaningful jMeta exception.

jMeta cannot handle all unchecked exceptions that might occur when calling a 3rd party library. So it might happen that such an exception propagates up to the

`jMeta` library user. There is no means of a "catch all" implemented in `jMeta`, as there is not even a central common ground this could be implemented in.

If a `jMeta` exception was caused by another exception, the causing exception must be specified in the `jMeta` exception.

Treating Abnormal Events

In case that `jMeta` detects an abnormal event, it throws a `jMeta` unchecked runtime exception. `jMeta` unchecked runtime exceptions must provide information about causing exceptions and a message. The message text may further describe the exact properties of the error. By now, message texts must be given in british English.

As mentioned before, the use of `jMeta` might also cause 3rd party unchecked runtime exceptions that are not under control of `jMeta`.

How should `jMeta` users treat abnormal events? As the name and description of abnormal events suggest, there are usually only a few possible actions, but there might be also specialized handlings. A well-written application should not simply terminate in case of an abnormal event. It should at least have a catch all section that handles unchecked runtime errors. This section may show the error to their users in a dialog box or at least log them. In some cases, the users of that application might be able to retry the action, or choose another element on which to apply the action (e.g. in case of a locked or read-only file). Depending on the kind of runtime error, a more sophisticated automated error handling can be set up by such an application. E.g. if a server connection had been lost, it might try to reconnect in a parallel thread, if the event happens during a database transaction, the database is of course rolled back and so on.

Treating Operation Errors

Treating operation errors is not only done with the approach that every operation error is signaled by a checked exception. Most cases where checked exceptions could be used can well be covered with the design-by-contract paradigm. Some rare cases require checked exceptions, though. This is treated in the following sections.

Design-by-Contract Design-by-contract, as originally created by Bertrand Meyer, states that each interface a client uses¹ must define a contract. The contract includes the constraints and possibilities for the user (i.e. what he must provide or which state he has to establish) as well as the duty of the interface to behave as specified. This has multiple advantages:

- The interface clearly defines how it needs to be used. It is therefore easier to understand and also easier to test. The contract is most of the time the basis for test cases of the interface.

¹Might be a OO class or interface as well as the interface of a module etc.

- If the interface implementation directly rejects wrong input according to its precondition, it strongly supports a defensive programming style. E.g. there are no `NullPointerException`s that are thrown from just anywhere in the implementation. Instead, the user input is checked immediately and pinpointed as the source of error.

Violating the contract on the side of the user is clearly a programming error.² If the implementation violates the contract, it is simply not implementing its interface correctly.

Design-by-contract distinguishes:

- *Preconditions* that must be established by the user before calling a method
- *Postconditions* that the implementation must establish after a method call
- *Invariants*, i.e. state conditions for the interface that must be fulfilled before and after a method is called, but not during a method call.

jMeta and Design-by-Contract The use of design-by-contract is basically described in the design decision “?? ??”.

In the Java implementation of **jMeta**, an invalid parameter range leads to an `IllegalArgumentException`, while a precondition that can only be checked by calling a method before-hand causes a `PreconditionUnfulfilledException` if not established. Every violation of a precondition contains a british English message that describes the error using detailed argument values.

Design-by-Contract vs. Multithreading In case of multithreaded use of **jMeta**, the approach of actively checking the precondition at the beginning of each method can be outperformed. If a first thread finds that the precondition is established, but a second thread e.g. removes an item after that, leading to the precondition not being established anymore, the first thread will run into an error. The design-by-contract front-line is therefore not water-proof. However, this is not a problem of the approach itself, but a general multithreading issue. **jMeta** itself refuses to lock the code of every method for multiple threads at once to avoid decreased performance. Instead, multithreaded applications must ensure that multiple threads properly work together.

Checked Exceptions Checked exceptions are only rarely used in **jMeta**, in exactly those cases where neither runtime exceptions nor preconditions are appropriate or applicable.

Whenever the caller can handle an error that is quite likely to happen, this is reflected as a checked exception. The caller catches the exception and does an appropriate handling. An example for this is an end of file condition, which may always occur and simply indicates that the client should stop reading. Another example is a class loading some data in its constructor. However, the program

²In theory, with a formal approach, such programming errors could even be checked at compile-time.

functions even if this data contains errors and cannot be loaded. In that case, the constructor could throw a checked exception.

9.2 Logging in jMeta

Logging is also used in jMeta as suggested by the following design decision:

DES 015: jMeta maintains a log file logging the most important events and errors

Logging is used in jMeta at least in EasyTag to protocol the startup loading of the library. In other components, logging is just used in exceptional cases, e.g. for error handling. The logging can be set on class basis by configuration done by the user, and it can also be fully disabled.

Rationale: In sufficiently complex system it is simply necessary to log status information. The user can disable and fully configure it to his needs.

Disadvantages: No disadvantages known

The question now is: What are the logging levels used?

DES 016: Informative output is given on INFO level, details on DEBUG and errors on ERROR level

jMeta logs the following output:

- **INFO:** Every informative output related to system environment, versioning etc.
- **DEBUG:** Detailed output for progress of specific startup activities or complex operations are logged during every call.
- **ERROR:** If necessary, additional diagnostic information is logged in case of errors on ERROR level.

Rationale: There is no arbitrary but unnecessary logging. Complex background operations might require detailed log output for debugging. When errors at runtime occur within the library itself, debugging can be made easier by including addition information in logfiles.

Disadvantages: No disadvantages known

In addition, we have to define how the logging framework looks like:

DES 017: There is no self-made logging component, but jMeta uses *slf4j* wherever necessary

Instead of a dedicated, hand-crafted and tested technical logging component which is known by every other component, and which encapsulates another logging library, jMeta code directly uses slf4j with a user-provided provider anywhere necessary. I.e. a direct dependency to slf4j is allowed from everywhere in the library.

Rationale: An own component wastes time and effort, in addition, every other component must somehow get an instance of it. Furthermore, innovations of the logging framework encapsulated usually remain encapsulated until you also provide it at your outside interface. slf4j is already a wrapper shielding you from details of other logging frameworks. You do not need to rewrite a logging framework today yourself, there is more precious time to spent if you do not. jMeta reuses the great logging providers already out there directly.

Disadvantages: No disadvantages known

The question is how to deal with technical runtime exceptions? Should they be logged, if so, where? This is clarified in the following design decision:

DES 018: jMeta code logs an error message before throwing a runtime exception

Before jMeta code itself throws a runtime exception, it logs a corresponding error message, preferably the exception itself, such that its message and stack trace get protocolled in the logs. This is done on ERROR level. For jMeta's own exceptions - including precondition checks - the logging is done in the exception constructor itself. A generic base class JMetaRuntimeException is used for this. jMeta code does not log anything if it throws a checked exception.

Rationale: There is no "catch all" in jMeta possible or reasonable, so if runtime exceptions are thrown, they are not captured in the jMeta logs. This is usually fine, unless jMeta itself throws such exceptions. We do not want to rely on the arbitrary logs of users of jMeta, but we want to be able to analyze what was going on in the jMeta code. The most important aspect of this is seeing where which exceptions originated from the library itself.

Not logging anything for checked exceptions because they are in some way expected. If they become unexpected, a JMeta runtime exception can do the logging again.

Disadvantages: Even during usual unit tests, there is logging output, i.e. probably file I/O

The last topic to handle here is whether we use tracing or not:

DES 019: jMeta does not use tracing anywhere

With tracing, we mean very fine grained log statements usually at the beginning and end of every method. This can be e.g. service facade methods only, or in the most extreme case any public, protected, default or private method called. Some applications log the input and output parameters in addition.

Rationale: Even though in rare cases tracing can give a lot of very important hints to debug a problem based on logs, it has also several drawbacks. Either you have to add it all manually. If you do, you might forget it, or you might write it in a way that is “repeating itself”, making general changes harder. And, in addition, the code is cluttered with statements, probably statements such as “if logging enabled, then log”. This makes the code barely readable. An alternative is to use aspect-oriented libraries. However, here it might be not always entirely clear how to only log *those specific methods* you want to trace. And there is a third-party dependency.

Disadvantages: In some cases, it might be hard to track down runtime errors if no tracing logs can be created and analyzed

9.3 Configuration

With the term *configuration* in jMeta, we understand the change of specific parameters, which customize the behaviour of jMeta at runtime. This is quite vague and the difference to usual specialized setters is unclear.

Thus we want to give some more criteria for configuration:

- Configuration is part of the public API of jMeta, i.e. it can be changed by the user.
- It is generic and thus extensible for further releases, i.e. adding new configuration parameters means to add a new constant to the API and update the documentation; the API for setting and retrieving the configuration parameter - as it is generic - does not change by adding new, changing existing or removing old parameters
- Configuration parameters usually have global effect and usually are just used once; but, of course in theory we could also make them dynamically changeable by users at runtime, and their scope could also be limited, e.g. to individual media

For now it turned out that there is no need in jMeta for such dynamic configuration parameters a.k.a properties, neither globally, nor for individual components. We thus summarize as follows:

DES 020: jMeta offers no generic configuration mechanism

jMeta offers no generic mechanism for (global or local) configuration, especially it currently defines no properties that influence the behaviour of the library. All necessary “configurations” for customizing the behaviour of a component are passed by the user via specialized (i.e. non-generic) setters or constructor parameters.

Rationale: There is no need for the bigger flexibility yet.

Disadvantages: No disadvantages known.

9.4 Naming Conventions and Project Structure

9.4.1 Java Naming Conventions

We use the standard Java naming and formatting conventions as defined by Sun. Only small addition: If there is a service interface with only one implementation, this implementation is having the prefix **Standard**. E.g. the default implementation for the **MediumStore** interface is named **StandardMediumStore**.

9.4.2 Package Naming Conventions

Here, we quickly define the package naming to be used for the library and extensions of the library.

Core Library

All packages for the core library start with `com.github.jmeta`, according to the Java naming conventions. The reason is that the library is hosted on Github. Below this core name, the following sub-packages are used:

- `.library.<component>` for all core components of the library, where “<component>” is replaced by the actual component name
- `.utility.<component>` for all technical utility code needed for the library, where “<component>” is replaced by the actual category name
- `.defaultextensions.<extension name>` for all extensions delivered by default bundled with jMeta, where “<extension name>” is the name of the extension, usually a single data format name; see also the next section for third-party extensions

Below these package names, there is another hierarchy:

- `.api` containing the public interface of the component, utility or extension; this package itself should not contain classes directly, but is further subdivided into:
 - `.services` containing the function interfaces the user needs to use to work with the component, utility or extension

- **.types** containing any transport objects, data types or the like for the component, utility or extension
- **.exceptions** containing any public exceptions of the component, utility or extension
- **.impl** containing the private implementation of the component, utility or extension; sub-divisions into further packages below it are possible, but component-specific and not standardized

Extensions

As mentioned already, the default extensions delivered together with **jMeta** of course also use a package naming convention quite similar to the core library package names. For any third party extensions, there is of course no definition of how the packages should be named, but it is up to the creators of the extension.

However, extension providers should follow the standard Java package naming conventions, and they should ensure that their package names do not clash with package names of other **jMeta** extensions.

9.4.3 Project Naming and Structure

Every IDE project containing code for **jMeta** starts has the form **jMeta<LibraryPart>[Extension]**, where **<LibraryPart>** is:

- **Library**: The core library with the main user API, without any extension specific parts
- **Utility**: Any utility the core library and extensions need, but no classes the library user must depend on, i.e. just internal library utility
- **DefaultExtension**: For all projects building an extension that is already bundled by default with the core library.
- **Tools**: For any supporting, development or testing tools, i.e. mostly Java applications that can either be used by library users or by library developers
- **Docs**: For any documentation of the library, no source code here (except for docs, e.g. tex files)

[Extension] is only used for the **DefaultExtension** library part. Each extension needs its own project and build unit. The reason is requirement “?? ??”. This allows us to depoly individual JARs for each single extension. In addition, there is an integration testing projects for testing media that contain multiple different formats.

Below, you can see the actual projects currently existing and their allowed dependencies: TODO

9.4.4 Build Module Structure and Dependencies

The Maven build module structure is of course aligned with the project structure: There is exactly one Maven module per project, having the same dependencies as depicted above. One exception might be the Docs project as it contains no sources. In addition, there is a single parent POM for general dependencies and aggregate builds.

10 Subsystem Technical Base

10.1 Utility Design

In this section, the design of the component **Utility** is described. Basic task of the component is to offer generic cross-functions that are independent of metadata reading and writing, i.e. independent of the functional target of **jMeta**. As there is currently no need for detailed design for this component, nothing is described here for now.

10.2 ComponentRegistry Design

In this section, the design of the component **ComponentRegistry** is described. Basic task of this component is the implementation of design decisions ?? and ??.

We just give the basic design decisions here. Before the current solution, we used a relatively complex custom-made component mechanism with caching, XML configuration of interface and implementation as well as registration of the components at startup. This code was organized in an eclipse project named “ComponentRegistry”, but was later a bit simplified later as “SimpleComponentRegistry”. However, after realizing that the built-in Java SE **ServiceLoader** mechanism is fully suitable and even better than the “SimpleComponentRegistry”, this mechanism is used now.

According to ??, a component is a “Singleton”. However, its life cycle must also be clearly:

DES 021: ComponentRegistry uses Java's ServiceLoader mechanism

`ComponentRegistry` uses Java's `ServiceLoader` class to load the implementation for a given component interface. All component interfaces and their implementations need to be configured in META-INF configuration files as indicated by the Javadocs of `ServiceLoader`.

Rationale: The `ServiceLoader` mechanism is quite easy to setup and requires nearly no specific coding. No code needs to be written to read corresponding configuration. A specific description / documentation mechanism of each component at runtime is not necessary, thus it was omitted (in contrast to previous implementation).

Disadvantages: Of course, `ServiceLoader` has its limitations: You might not be easily able to add components at runtime without writing additional code or not at all. And it does not have any idea of the term of a component, it just knows interfaces and implementations. But this is all not a requirement for `jMeta`.

Here is just a short summary of small amount the self-written code in `ComponentRegistry`:

DES 022: `ComponentRegistry` consists of a single non-thread-safe class for looking up implementations and cache handling

`ComponentRegistry` solely consist of a single class. This class caches `ServiceLoaders`. It offers the following methods:

- **lookupService:** Loads a service implementation (if called the first time for an interface) and returns it, adds its `ServiceLoader` to the cache afterwards
- **clearServiceCache:** Clears the internal cache of `ServiceLoaders` such that the next call re-reads and re-instantiates new implementations again.

The class offers this as static methods, but it is not thread-safe.

Rationale: Why do we cache `ServiceLoaders`? Because the `ServiceLoader.load` method always creates a new instance of the `ServiceLoader` class, which has the following effects:

- Components are singletons, thus, if two different other components need an implementation of the component at different points in time, they would get different instance of the service implementation, as they would always create a new `ServiceLoader` instance.
- `ServiceLoader.load` does file I/O, which is not what we want every time we request a new implementation.

The method `clearServiceCache` is necessary for test cases which might need to reset the `ComponentRegistry` state after each test case. As it is a static class, this method is necessary.

Why static methods? Because we do not want to care about instantiating it, passing the same instance everywhere needed.

Why is it not thread-safe? Because the whole library is not intending to be thread-safe.

Disadvantages: No disadvantages known

10.3 ExtensionManagement Design

This section describes the most important design decisions for the technical helper component `ExtensionManagement`. The main task of this component is discovery of extensions and to make their interfaces available to the whole library, preferably in a very generic way.

In the first version of this components, extensions were loaded by a relatively complex `URLClassLoader` approach and used a main and an extension-specific XML configuration file. However, this was more complex than necessary, thus we define here a more lightweight approach:

DES 023: Extensions are discovered and loaded at library startup directly from the class path, no dynamic loading of extensions is required

When the library is first used in a Java application, it scans the class path for any available extensions. The extensions that are available at that time are available for further use. There is no possibility for the user itself to add extensions. Furthermore, it is not possible to add, exchange or update extensions at runtime in any way.

Rationale: It is quite convenient for users to simply put extension onto the class path, e.g. using Maven or Gradle, and it is auto-detected. This fulfills the requirement “?? ??”. Dynamic loading or updating of extensions automatically or triggered by the `jMeta` user is not necessary, as in most use cases it should be clear at compile-time which extensions are required. As extensions should not be introducing a big overhead, applications can easily provide all in total possible extensions at build-time already.

Last but not least, for the `jMeta` implementation, complexity is largely reduced if we omit the requirement to load, update or exchange extensions at runtime.

Disadvantages: No disadvantages known

How does the library identify extensions on the class path?

DES 024: An extension is identified by implementing `IExtension` and being configured by the `ServiceLoader` facility

An extensions is identified by each implementation of the `IExtension` interface that is configured in a JAR file on the class path as a service provider. It is found on the class path by the library core by using the Java `ServiceLoader` class. It does not matter if the extensions is located in a separate JAR file or even already in the core, and how much extensions are in a single JAR file.

Rationale: `ServiceLoader` is a quite convenient and very easy to use mechanism.

Disadvantages: No disadvantages known

Regarding the configuration of extensions:

DES 025: There is no central configuration file listing available extensions; extensions themselves use code to provide a description of the extension and no config files

Extensions are just discovered via classpath. There is no central configuration file listing all available extensions for the `jMeta` core. Furthermore, extensions themselves might provide a description, but this description is also not contained in a configuration file, but rather returned by the `IExtension` implementation directly in Java code.

Rationale: A central configuration of all available extensions for the `jMeta` core would be very inflexible. It would need to be extended anytime a new extension is available. And then, by whom? The end user, the developers of the extensions or even the `jMeta` developers? This would not make much sense.

For the extension specific details, a configuration file could be used, but it again makes it harder to create an extension and we would need additional generic code to parse the configurations. Configurability for such things is not necessary at all. If a change in the description is necessary, it would be related to changes in the extension anyway and the JAR for the extension can be updated to have the updated description. An XML file embedded in the JAR file would not bring any differences or even advantages.

Disadvantages: No disadvantages known

Which functional interfaces do extensions provide in which way?

DES 026: Extensions can be queried dynamically by the library core to provide 0 to N implementations for an arbitrary Java interface

`ExtensionManagement` provides a generic method in `IExtension` where `jMeta` core components can get all implementations of any Java interface they request which are provided by the extension. The extension of course might not have such implementations. But it might also have more than one implementation.

Rationale: Although in section “?? ??”, specific cases for the extensibility of `jMeta` were mentioned, we do not want to hard-wire this into a generic component such as `ExtensionManagement`. If we would do this, there would be most probably a cyclic dependency from `ExtensionManagement` to these components of the `jMeta` core that are extensible and thus need to load extensions.

Furthermore, we do not need to change the interface of `ExtensionManagement` whenever extensions might need to return implementations for any other interfaces that can be extended in future.

Disadvantages: No disadvantages known

What is the lifecycle of an extension, then, and is there anything else they should do?

DES 027: An extension is instantiated just once and then used throughout the whole application, offering a method to retrieve implementations

According to ??, an extension is loaded at `jMeta` startup. It can be considered as singleton, as just a single instance of it is loaded using `ServiceLoader`. Extensions should have the sole purpose of returning new implementation instances, for which they implement the method

`IExtension.getExtensionProviders`. Anytime this method is called, they provide a list of new implementations they provide for the given interface. If they do not provide an implementation, they return an empty list. Besides that, there is a method `getExtensionDescription` where extensions can provide additional detail information for the extension.

Extensions should not do anything else. Especially, they should not have any internal mutable state or perform any background processing.

Rationale: There is no reason why extension providers should implement any special behaviour in addition to just returning implementations to use by the core, except, maybe harmful hacking code. Of course, `jMeta` cannot really prevent that, but at least it should be made clear.

Disadvantages: No disadvantages known

We should now go into more detail how the library core interacts with extensions. So far, we defined that extensions will be loaded once at startup using the `ServiceLoader` facility. But what is the public interface of the `ExtensionManagement` component?

DES 028: ExtensionManagement provides the interface

`IExtensionsManagement` which provides access to all extensions found. `ExtensionManagement` provides the interface `IExtensionsManagement` for the library core with the method `getAllExtensions`. This method returns all available `IExtension` implementations found.

Rationale: An easy to use `ExtensionManagement` without any surprises.

Disadvantages: No disadvantages known

11 Subsystem Container API

11.1 Media Design

In this section, the design of the component **Media** is described. Basic task of the component is to provide access to memory areas which contain multimedia data. Primarily these are files.

The term MEDIUM needs to be sharpened here: In “?? ??” we had defined: “A MEDIUM defines the storage medium of DATABLOCKS. It can be a file or a MEDIA STREAM, or the main memory itself.”

In detail, the term summarizes the aspects “physical storage” and “access mechanism” (e.g. file-based random-access, or byte stream). Thus there might perfectly be two different media which access the same physical storage, but using different access mechanism. The term MEDIUM is an abstraction and potentially allows even more special possibilities, like media streams, databases etc.

11.1.1 Basic Design Decisions Media

Here, the fundamental design decisions of the component **Media** beschrieben.

Supported Media

This section lists decisions about supported MEDIA. To start with, it is clear that **jMeta** must support files as basic medium.

DES 029: Support for random-access file access

jMeta supports the use of files as input and output medium via **Media** with access mechanism “Random Access”.

Rationale: Files are *the* fundamental and most common digital media containers, even in 2016. Of course MP3 files, AVI files etc. with multimedia content are wide-spread. A library such as **jMeta** must support files as core element. To more efficiently process files, random-access is inevitable. Especially reading at arbitrary offsets - e.g. tags at end of file - as well as skipping of unimportant content is efficient to implement with random-access.

Disadvantages: No disadvantages known

But also reading streams shall be supported to increase the flexibility of the library:

DES 030: Support for sequential, reading byte streams

jMeta supports the use of reading byte streams, i.e. `InputStream`s for input in mode “sequential access”.

Rationale: `InputStream` represents the most general alternative of a `MEDIUM` from Java perspective, which ensures a potentially higher flexibility for using jMeta. E.g. multimedia files can be read from ZIP or JAR archives using streams, and support for media streams might be easier to implement in later releases - However: To state clearly: media streams do have nothing to do with this design decision. They might be implemented completely different in upcoming releases.

Disadvantages: An `InputStream` supports by definition only sequential access and no random-access (e.g. via `FileInputStream`). Thus there might be higher complexity for implementation, as well as significant performance drawbacks because of lacking random-access.

Last but not least, the library offers access to RAM contained data, due to flexibility:

DES 031: Support for random-access to byte arrays

jMeta allows for random-access to byte arrays as input medium and output medium.

Rationale: Already loaded memory content can be parsed with jMeta without need for artistic climbs, increasing flexibility of the library.

Disadvantages: No disadvantages known

What about `OutputStream`s? That is discussed in the following:

DES 032: No support for writing byte streams

jMeta does not support writing byte streams, i.e. `OutputStream`s.

Rationale: `OutputStream`s are write-only, but still not random-access. Thus we would need - provided we want to access random-access media in a random-access style - a second implementation next to writing random-access. A combined usage of `InputStream`s and `OutputStream`s for Read-/Write access on the same medium is not designed into the Java API and leads to diverse problems. As jMeta already implements writing to output files and byte arrays, for reasons of effort, `OutputStream`s are not supported as output media. The user might implement `OutputStream`s easily by him- or herself, e.g. by first writing into byte arrays, then into an `OutputStream`.

Disadvantages: No disadvantages known

Consistency of Medium Accesses

Parallel access to the same medium from different processes or threads, reading by one and writing by the other, might lead to unpredictable difficulties - even without using any caching. If you e.g. have some parsing metadata like the length of a block in bytes at hand, but a parallel process shortens the block, your read access trying to fetch the whole block will run into unexpected end of file or read inconsistent data.

To avoid such problems, there are special locking mechanisms for exclusive access to the bottleneck resource, at least for files. We define:

DES 033: Locking of files during jMeta access

Files are *always* locked during access by jMeta explicitly. File content is protected by exclusive locks from corruption by other processes and threads. See [?], where we show that a file in Java must be explicitly opened for writing to be able to lock it. “During access” means: After opening it and until closing it. The lock thus might be long-term. jMeta opens a file for writing (and locking) even if the user explicitly requested read access only.

Rationale: Other processes and threads of the same JVM cannot access the files and corrupt any data, which avoids consistency problems.

Disadvantages: It is not possible to access the same file in parallel threads when using jMeta. It seems rather unlikely that such parallel access to the same file (e.g. reading at different places) can speedup an application. But for future media this might indeed be a drawback.

The locking of byte streams or memory regions does not make sense, as discussed in the following design decisions:

DES 034: No locking of byte streams

Byte streams are not locked

Rationale: The interface `InputStream` does not offer any locking mechanisms. jMeta will not try to guess the kind of stream and lock it (e.g. by checking if it is a `FileInputStream`).

Disadvantages: No disadvantages known

For different processes, the os usually protects access of memory regions. The question is whether jMeta should protect access to byte arrays:

DES 035: No locking of byte arrays

Byte arrays are not locked

Rationale: This makes not much sense as the user anyways gets a reference to the byte array by the API, and thus can access and manipulate the raw bytes arbitrarily in a multi- or single-threaded way. Protecting it by thread locking mechanisms increases complexity and does not seem to generate any benefits whatsoever.

Disadvantages: No disadvantages known

Unified API for Media Access

In the wiki article [?], we have shown clearly the differences between byte streams and random-access files. With so many difference the question arises: Can this be unified at all and does the effort make sense here? The least common demoninator for random-file-access and `InputStreams` is the linear reading of all bytes in the medium. This is clearly too less. It denies all advantages of random-access. The intersection of features for a unification is therefore not making sense.

Moreover, we want a unifying combination of both approaches:

DES 036: Unified access to all supported media types in one API

Media offers a common abstraction for accessing files via random-access, **InputStreams** as well as byte arrays. This API provides the advantages of both access mechanisms via a common interface. The implementation throws exceptions of kind “Operation not supported” in some cases, if a feature is not supported by the medium. In other cases, a meaningful alternative behaviour is implemented. The using code must perform branch decisions at some places depending on the medium type.

While byte arrays are no problem for the abstraction, even random-access files and **InputStreams** have more in common as you might think at first glance:

- The operations Open, (sequential) Read, Close.
- **InputStreams** can also (at least technically) be assigned a beginning, offsets and an end.
- Files can be read-only, too, which **InputStreams** are always by definition.

Writing access to a read-only medium is acquitted with a runtime exception, especially for an **InputStream**.

The main difference between files and **InputStreams** is of course: Random access is possible for files, while **InputStreams** can only be read sequentially. This difference can be potentially decreased using mechanisms such as buffering.

Rationale: The API of the component **Media** gets easier for outside users, its usage feels more comfortable. Using components of **Media** can offer their users in turn an easier interface. At the same time, the advantages of both approaches (random-access and better performance for files, generality and flexibility for streams) are still available.

Disadvantages: A few operations of the API cannot be implemented for both media types, which makes case decisions in the client code necessary in some cases.

Two-Stage Write Protocol

When Writing, it is all about bundeling accesses and buffering. We want optimum performance und thus want to implement these mechanisms. Therefore we commit to following design decisionfor implementing writing in **jMeta** in general:

DES 037: Media uses a two-stage write protocol controlled by the user

The first stage is the mere registration of changes, that need to be written to the external medium. In this first stage, there is no access to the external medium yet. The second stage is the operation *flush*, the final writing and committing of all changes to the external medium. The underlying implementation bundles the write actions according to its needs into one or several packets and executes the write only in the second stage.

Rationale: An efficient write implementation is possible. Internally, write actions can be bundled as needed to perform better. And this can be done without forcing the user to do it himself. The user can perform write (registration) actions whenever his code architecture needs it. Saying this, the user code is not burdened with too much restrictions or rules. Furthermore, the potential possibility of an “undo” of already registered actions comes into view.

Disadvantages: Errors that occur when actually flushing changes to the external medium are recognized potentially quite late. Thus the registration of changes is quite fast while the flush itself can be a long taking process. Bugs might be introduced by user code forgetting to implement the second step, the flush.

Even if we implement this, it must be clearly stated that this is not in any way a transaction protocol as implemented by some O/R mappers (e.g. hibernate) or application servers. The mentioned protocol is much simpler and not in the least capable to provide ACID! Thus the following exclusion:

DES 038: Writing in Media does not guarantee ACID, in case of errors during *flush*, there is no rollback

ACID (atomicity, consistency, isolation and durability) is not ensured neither by the implementation of *Media* nor in general by the Java File I/O. If e.g. an error occurs during Writing in the *flush* stage, some data has been written already, while upcoming data will not get written anymore. There is no undo of already written data. The operation *undo* must not be mixed up with a rollback and it is no action that is done automatically. While isolation and durability can be more or less provided, the user is responsible for consistency and atomicity himself.

Rationale: A transaction manager that guarantees ACID, and this for files, is really hard to implement (correctly). This requirement is somehow out of scope, no other competing library is doing something similar. *jMeta* will not be a database!

Disadvantages: No disadvantages known

Requirements for the Two-Stage Write Protocol

Which writing operations must be offered? One method `write()` - at the end it is the only really writing primitive of the Java File I/O - is not sufficient. How do you remove with this method? `write()` equals *overwriting*, which means you have to do a lot of manual work to implement insertion, removal and replacement with this operation, it is not convenient at all. **Media** must offer a better API, taken some of the burdens of I/O from the user. Here, we only specify the necessary operations, without going into details with their implementation - this will be done later.

To develop a good design, however, you must first list down the user's requirements to **Media**. This especially includes the requirements for a two-stage write protocol. Main users of the component is definitely the component **DataBlocks**. It uses **Media** to extract and write metadata from and into tags. Without going into the design details of **DataBlocks**, here we nevertheless list detailed requirements that **DataBlocks** has for **Media** regarding two-stage writing, see table ??.

ID	Requirement	Motivation
AMed01	It must be possible to insert bytes	Formats such as ID3v2 can be dynamically extended and have a payload of flexible length. Before an already present data block, it must be possible to insert another one. There is especially a need for an insertion operation in the case when metadata with dynamic length need to be written at the beginning of a file.
AMed02	It must be possible to remove bytes	With the same motivation as for insertion. It must be especially possible to remove entire metadata tags.
AMed03	It must be possible to replace bytes and not only overwrite, but also grow or shrink an existing byte area with replacement bytes	In metadata formats, there are both static fields with fixed length as well as dynamic fields such as null-terminated strings. If bytes are already present, it must be possible to overwrite them to save costly remove and insert operations. The growing and shrinking is especially useful and represents a higher level of abstraction. If this would not be possible, replacing a previous small string value by a new longer or shorter one would need to be implemented with two operations (overwrite and insert or remove, respectively).
AMed04	Inserted data (requirement AMed01) must be changeable before a flush e.g. by extending, overwriting or removing of child fields inside the inserted data block	Based on the two-stage write protocol, an arbitrary number of writing changes can be made before a flush , and these might correct each other. E.g. a new ID3v2 tag footer is inserted, that stores the length of the tag. Assume that after this, a new frame is inserted into the tag, before the flush. This requires the size field in the firstly inserted footer to be changed afterwards again, before the flush.

ID	Requirement	Motivation
AMed05	Replaced data (requirement AMed03) must be changeable before a flush e.g. by extending, overwriting or removing of child fields inside the replaced data block	A prominent example is insertion of and step-by-step extension of a frame into an ID3v2 tag: For the first creation as well as each extension, the size field of the tag must be changed, which induces a replace operation each time. E.g. it is allowed that users first only create and insert the new frame, and then insert new child fields afterwards, step by step.
AMed06	The padding feature of several data formats should be used by jMeta	Formats such as ID3v2 allow padding, i.e. using an overwrite buffer are to avoid newly writing the whole file. jMeta must use this feature when writing data, such that e.g. an insert only affects the file content until the padding area, effectively decreasing the padding, while a remove increases the padding, but the overall tag size remains the same. It is rather an indirect requirement which needs not necessarily be implemented by Media only.
AMed07	The operations replace, remove and insert must be undoable before a flush	This allows to avoid unnecessary accesses to the medium and to undo mistakes by end users.
AMed08	It must be possible to insert multiple consecutive blocks of data	E.g. you want to add an APEv2 tag first, then an ID3v2.3 tag and then an ID3v1 tag.

ID	Requirement	Motivation
AMed09	It must be possible to insert a block of data before an existing block, then replace or remove this existing block with something else; the same must be possible for first remove or replace, then insert (i.e., the inverse order)	E.g. you want to add an APEv2 tag first, then replace the follow-up, already existing ID3v1 tag with some different content.
AMed10	It must be possible to modify (remove, insert, replace) a block of data, then remove an enclosing block of data	E.g. the code flow first decides it is necessary to add a new frame to an existing ID3v2.3 tag, modify some of its frames, remove another one; but then due to some decision in the code, it becomes necessary to remove the entire tag.
AMed11	It must be possible to modify (remove, insert, replace) a block of data, then replace an enclosing block of data	E.g. the code flow first decides it is necessary to add a new frame to an existing ID3v2.3 tag, modify some of its frames, remove another one; but then due to some decision in the code, it becomes necessary to replace the entire tag with different content.
AMed12	Except AMed10, Media does not need to support overlapping removes or replaces	Data blocks are never overlapping, except AMed10, there is no such case that you remove a block, then remove an overlapping portion of it.
AMed13	Media does not need to support modifications (remove, insert, replace) within an already removed or replaced block of data	It is not clear how to treat changes in an already removed or replaced parent object consistently, so we just reject them.

Table 11.1: Requirements for the two-stage write protocol by DataBlocks

Based on these requirements, we can first define the following basic design decisions for writing:

DES 039: Media offers the writing operations *insert*, *remove* and *replace*

The user can:

- *insert* N bytes at a given offset (addresses requirement AMed01)
- *remove* N bytes at a given offset (addresses requirement AMed02)
- *replace* N bytes at a given offset by M new bytes (addresses requirement AMed03). Thus, *replace* has three flavors:
 1. replace N bytes by $M > N$ new bytes actually behaves like an insertion, replacing the first N bytes with new ones, and inserting additional $M - N$ bytes behind. We call it an *inserting replace*.
 2. replace N bytes by $M < N$ new bytes actually behaves like a removal, replacing the first M bytes with new ones, and removing additional $N - M$ existing bytes behind. We call it a *removing replace*.
 3. replace N bytes by $M = N$ new bytes actually behaves like a usual write. We call it an *overwriting replace*.

Rationale: See requirements AMed01, AMed02 and AMed03. The burden to implemen these convenient operations using the Java File I/O, which essentially only offers `write()` and `truncate()`, is taken over by **Media**, such that the user need not care.

Disadvantages: No disadvantages known

How is the padding requirement addressed? Like this:

DES 040: Requirement AMed06 must be addressed by components using Media, it is not built into Media itself

Media does not implement AMed06 itself, it does not know something like “Padding”, but just the three primitive change operations. Using them, a third party component can do the following using Media:

- If a new field is added to the ID3v2 tag (with an *insert*) and padding that is longer than the newly inserted content is present, the 3rd party component has to do a **remove** of a corresponding number of padding bytes, or a **replace** of old padding bytes by new (shorter) padding bytes.
- If a field is removed from the ID3v2 tag (with a *remove*), the 3rd party component can do an **insert** of a corresponding number of padding bytes at the end of the tag, or a **replace** of old padding bytes (if present) by new (longer) padding bytes.

Rationale: Media should stay clean of any “functional” terms defined by some formats, but it should just stick to the primitives it offers to modify media content. As shown above, the requirement can quite easily be implemented by tag-specific logic on top of Media.

Disadvantages: No disadvantages known

As we have a two-stage write protocol, the *undo* of not yet flushed changes is possible, and according to the requirements also necessary.

DES 041: Writing operations on a medium can be undone with *undo* before a *flush*, addressing requirement AMed07

Writing operations lead to pending changes according to ???. These can be undone according to the requirements defined above.

Rationale: The application logic can require *undo* in some cases, e.g. for corrections of mistakes done by an end user. Instead of requiring to call the inverse operation (if any at all), the user is much more convenient with undoing the operation itself directly. This also ensures that the using code does not need to trace changes to be able to find which is the inverse operation.

Disadvantages: No disadvantages known

This leaves open AMed04 and AMed05 as well as AMed08 to AMed13 which we will address later when discussing the details of the two-stage write protocol.

Caching

The component Media takes over I/O tasks with potentially slow input and output media. Thus, it is here where the basic performance problems of the whole library need to be solved. We will approach these topics with some motivation and

deductions.

In [?], basic stuff regarding performance with file access is discussed. The ground rule for performant I/O is minimizing accesses to the external, potentially slow medium. For writing, we already introduced ??. A similar important question is: How can you make reading perform better?

At first it is quite clear that for reading, you should help yourself with buffering to improve performance:

DES 042: Reading access can be done using a buffering mechanism, controlled by the component's user

For each reading access, the calling code can specify the number of bytes to read, which corresponds to a buffering. The code controls the size of the buffer by itself. It potentially can also read only one byte. It lies in the responsibility of the calling code to read an amount of bytes that makes sense and minimizes read accesses.

Rationale: A hardcoded fixed length buffering would rather lead to performance disadvantages, as there might be read too much bytes, more than necessary in average. Furthermore, depending on the fixed size, it would be necessary to read two or even several times when a bigger chunk of data is needed. According to [?], there is no “one size fits all” for buffer sizes in file I/O. The logical consequence is to let the user decide.

Disadvantages: No disadvantages known.

A further important aspect is caching: With *Caching*, we refer to the possibly long-term storage of MEDIUM contents in RAM to support faster access to the data. Buffering differs from caching in a sense that buffering is only a short-lived temporary storage without the necessity of synchronisation.

To start with, we can see - besides the already mentioned buffering - different kinds of “Caching” in an application that is based on jMeta:

- During file access there are caches on hardware level, in the OS and file system.
- Java supports temporary buffering via the `BufferedInputStream`, and caching explicitly via `MappedByteBuffer`.
- Applications that mostly are interested in human-readable metadata read it, convert it via jMeta into a clear-text representation and show this representation in their GUI. In this case the GUI model represents a kind of caching that also allows changes to this metadata in the GUI, it is not necessary to re-read again from the medium.

If we look at all these alternatives, the question arises why at all an additional built-in caching in jMeta would be needed? For answering this question, we should look at some use case scenarios for the library: One scenario is the already mentioned reading of metadata from a file to display it in a GUI. For such a case,

caching would usually not be very useful. You read once, and maybe twice, if the user wishes to update the screen. It would be an acceptable performance without caching. Another use case is the arbitrary jumping between parts of a container format file using a low-level API to process specific contents. This is true “random access”. The question is: Do you want to read the same place twice? Sometimes possibly yes. Instead, would you want a direct medium access again? Possibly yes or no.

The last question brings up another general problem with caching: The problem of synchronicity with the external medium. If the medium has been changed in between, the cache content is probably aged and invalid. Code that accesses the cache can usually not recognize this.

We first sum up the identified advantages and disadvantages:

Advantages	Disadvantages
+ Performance improvement when reading the same data multiple times, as cache access is much faster than the access to the external medium	– When only accessing once there is of course no performance improvement
+ In a cache you are - in principle - more flexible to reorganize data than on an external medium, making it easier to correct, undo or bundle changes.	– External 3rd party changes on the MEDIUM cannot be recognized and lead to invalid cache content that might lead to erroneous behaviour or data corruption in follow-up write actions.
	– There is additional code necessary for caching, e.g. questions such as “when is the allocated memory freed?” must be answered. For consistency topics, even more complex code is necessary.
	– More heap space required

Table 11.2: Advantages and disadvantages of caching in `jMeta`

The disadvantages outweigh the advantages. Why should you then use caching in `jMeta` at all? Because some of the previous design decisions combine well with a caching approach:

- ?? can be achieved using a cache, as we see just a little later
- ?? can be implemented with a cache, i.e. anything that has been buffered should directly go into the cache for subsequent read actions
- ?? may or may not be easier to implement using a cache. In this case the

cache would be used to store the registered changes before a flush. However, if it would only be this, a cache would be greatly too complex. Easier solutions are possible for holding the not-yet-flushed data.

Thus we decide:

DES 043: Media keeps medium data read in a cache

Media uses a RAM-based, potentially long-lived fast storage (cache) to store already read content of the **MEDIUM**. Subsequent read accesses request the cache content (if present) only.

Rationale:

- We provide faster repeated read access to already read data to the end-user
- This can be used for direct implementation of ?? in a sense of buffering when reading. i.e. the cache works as the buffer for ??
- Implementation of ?? can be done using a cache

Disadvantages: Were given in table ??. The alternative is a direct medium access. To summarize the disadvantages against a direct medium access:

- Higher code complexity
- More heap required, the cache is durable
- The medium might change by external processes, such that the cache content is not in synch anymore.

Note that this last mentioned disadvantage is mostly mitigated by ??.

How to use caching to better achieve ???

DES 044: Caching is used to better mitigate the differences between **InputStreams and files according to ??.**

The data that has been read from an **InputStream** is always put into a cache. Reading actions are therefore allowed to “go back” to already read data, by not issuing another direct access (which is anyway not possible using an **InputStream**), but by taking the data from the cache. “Read ahead” for areas that have not yet been reached on the **InputStream** lead to the behaviour that all data up to the given higher offset is read and cached.

Rationale: This implements ?? nearly entirely, “transparent” to the user.

Disadvantages: Even more heap space is necessary for **InputStreams**, as in extreme cases the whole medium might end up in the cache, which might lead to **OutOfMemoryErrors**.

Of course, the disadvantages mentioned in ?? are heavy-weight. If you wouldn't do anything about it to mitigate these disadvantages, then ?? would be nonsense, as the advantages of this approach would be dramatically overshadowed by its disadvantages.

DES 045: The user can set the maximum size of the cache per medium, the cache keeps only the newest added data

The user is allowed to limit the maximum size of the cache per medium by configuring it before creating the cache. The cache implementation ensures that at any time, the cache does not contain more bytes than which correspond to the maximum size. This is done using a FIFO (first in, first out) mechanism. I.e. the bytes added first are first discarded whenever trying to add new bytes to a cache.

Rationale: This way, the user can influence the maximum heap size required for caching of a single medium. At the same time, he needs not control the cache size himself, but it is internally managed by **Media**

Disadvantages: No disadvantages known

DES 046: Performance drawbacks of `InputStreams` are explicitly documented

Reducing differences between `InputStreams` and files by caching in accordance to ?? means: You must deal with the fact that you cannot store virtually unlimited `InputStreams` in a cache. For file access, the user can also choose between `FileInputStream` and more direct access via `RandomAccessFiles`. The performance drawbacks induced by using `FileInputStream` compared to random-access - which are introduced by a unified API according to ?? - are explicitly described in the **jMeta** documentation. The mitigation mechanisms (setting maximum cache size, disabling caching) are explicitly described with their corresponding consequences.

Rationale: There are no wrong expectations by providing the unified API. The contract is described clearly enough to the user. He must choose the medium best suited for his purpose.

Disadvantages: No disadvantages known

DES 047: The user cannot free up cache data in a fine-grained way

Media will not provide any mechanisms for the user to free up cached data himself in a fine-grained way (e.g. range-based).

Rationale: This functionality needs to be implemented and tested (additional effort). It is quite unlikely that it is actually needed when having implemented `??`. Furthermore, when should the user free the data? A monitoring by the using code is necessary which also makes the using code more complex.

Disadvantages: No disadvantages known

After these results, the disadvantages previously listed in table `??` and in `??` need a closing look:

- **Code Complexity:** The higher code complexity cannot be disregarded. You have to accept it when implementing a permanent caching. It implies you need very good unit and integration tests to ensure it works as expected.
- **More Heap Memory:** It is usual to achieve a better runtime performance by increasing the memory footprint. So it is here. To nevertheless avoid `OutOfMemoryErrors`, we have defined `??` and thus give enough room for the user to avoid these situations.
- **Data Corruption due to Out-Of-Synch Medium:** The cache might receive updates that are not yet persisted on the external medium, as explicitly allowed by `??`. A problem might arise due to changes by other processes or threads. These cannot be handled in a general way by `jMeta`. Thus we have introduced the locking of media in `??`. Even this cannot give a full protection for some OSs. The user is in any case informed about irresolvable inconsistencies by a runtime exception.

Finally, we want to list a design decision rejected during a proof of concept, which was to provide a mechanism of skipping bytes for `InputStreams` instead of reading them into the cache, here is why:

DES 048: The Media API does not provide the option to skip bytes of an `InputStream` instead of reading them into the cache

It sounds like a good idea to offer the configurable possibility to skip bytes from an `InputStream` instead of unnecessarily reading them (and put them into a cache). This mainly applies for the case when you are currently at offset x , but now you want to read bytes from the stream starting at offset $x + 100$. You might never want to access the bytes between x and $x + 100$. So it would be worthwhile to allow the possibility to simply skip them instead of reading them into memory and into the cache.

However, `InputStream.skip` is not as straightforward as one could wish. It does not guarantee to skip the number of bytes given, but it might skip fewer bytes or even return a negative number. It might also throw an `IOException`, which might or might not include the case of reaching the end of medium. Although not specified, there might be the same behavior as for `read`, that `InputStream.skip` might block. Furthermore the javadocs specify that it is only supported for streams that support seeking. All in all, this gives the impression that the method is highly platform and implementation-dependent. How to bring this implementation into a reliable form that guarantees to always skip the given number of bytes or block (possibly with timeout)?

Rationale: It is too complex to implement the method reliably, thus we omit it and only offer the possibility to read ahead until a given offset. All bytes read are possibly read into the cache. However, the cache is “self-cleaning” according to ??, if a maximum size is configured. This mechanism should be sufficient for achieving a moderate heap usage.

Disadvantages: Probably more heap usage and bytes might be unnecessary read and kept in memory.

Implementing the discussed caching mechanisms is a big challenge. It will be detailed more in the implementation part of this component. Here, we can only exclude one way of implementing it:

DES 049: `MappedByteBuffer` will not be used to implement caching

You could come with the idea to use the Java NIO class `MappedByteBuffer` for implementing the caching of ??. However, we do not use it and implement another solution “by hand”.

Rationale: It is not guaranteed, that each OS supporting Java also supports a `MappedByteBuffer`. It is also not guaranteed that the data “cached” is really present in RAM and thus accessible faster. Furthermore, it is indicated that for each consecutive region of a medium a new `MappedByteBuffer` instance including new OS call would need to be created. Thus this approach is unpredictable and might lack the desired benefits.

Disadvantages: The “by hand” caching is harder to implement.

At the end we shortly list a special case of caching for byte array media:

DES 050: For byte array media, caching is always disabled

For byte array media, caching is always disabled

Rationale: byte arrays already are in RAM, caching would just be unnecessary overhead

Disadvantages: No disadvantages known

Reading Access to the Medium

The two-stage write protocol introduced in “?? ??” brings up some questions regarding reading the data. The most important among these: What does the user need to consider after calling a writing operation (stage 1) and before *flushing* these changes (stage 2)? Especially: What do reading calls return after already having made changes, that are however not yet *flushed*? The possibilities we have:

1. Either the last persisted state on the medium after opening it or the last successful *flush*, respectively,
2. Or already a state the includes any “pending” changes introduced by writing calls before the *flush*?

You could base your answer on the following: For sure alternative (2), as it is this way that e.g. transactions mostly work for code accessing databases. What you have already written during the transaction, you re-read later, too, even if the transaction is not yet persisted. However, the view of `jMeta` is different:

DES 051: The user can only read what is currently persisted on the medium

Even if there are pending changes not yet *flushed* (e.g. inserts, removes), with `Media` the user can only see the latest flushed state.

Rationale: The changes the user has registered are coming from the user, and he thus could potentially keep bookmarks of them. Therefore their sole management by `Media` is - at this point in time - not strictly necessary. Furthermore it must still be possible to read the data of a datablock that is threatened by a pending remove.

Another good reason for this behaviour is that the reading operations are much less complex, as they do not need to consider any pending changes. The code that reads data can be sure to always only work on a persistent (or at least a cached) state. Thus it cannot occur that logic is basing on data that is not yet persisted.

Disadvantages: The expectation that “what I have written before - even if pending - I can re-read afterwards” is not fulfilled. The user must manage this for changed data by himself, at least temporarily until the next flush.

In “?? ??”, caching has been discussed in detail. For buffering, we first need an operation to do buffering without actually returning the buffered data:

DES 052: Explicit operation for buffering of media data

There is an operation *cache* which buffers n data bytes starting at a given offset, without returning this data. Additionally, there is an operation to query the number of bytes buffered concecutively starting at a given offset. Of course, this might lead to an “end of medium” situation indicated by the method to the caller.

Rationale: Necessary for implementing ??. Of course one could ask: Why isn’t it sufficient to just provide a single `getData` operation that returns data, either from the cache or from the medium. If the data was not yet in cache, it also adds it to the cache. The reason is: The code might not want the bytes yet! It just wants to give a statement like “At this point in time, I know how much bytes I might need to read later, so please already buffer them. But do not give the bytes to me, because at this point in the code I cannot really use them and it would complicate my code strucutre.” For instance, if only providing one method returning the data, the code must pass the read data as parameter to other methods to read it. Instead of doing this, the code can just buffer the needed number of bytes, then call other code to get just the data it needs without needing to fiddle keeping track of the last read byte.

Disadvantages: No disadvantages known

Additionally, you must be able to get your hands at the buffered data:

DES 053: Explicit operation to get medium bytes

There is an operation *getData* which returns n data bytes starting at a given offset.

Rationale: Without it there would not be any possibility to read data from a medium, as *cache* only buffers it internally without returning it.

Disadvantages: No disadvantages known

Now the question arises, how the operation *getData* interacts with the cache, which is answered here:

DES 054: *getData* combines data from the cache with data from the medium and updates the cache thereby, if necessary

getData reads data from the cache, if the given range is entirely contained in it. If it is not entirely contained in the cache, *getData* reads the bytes present in the cache, and reads the non-present ones from the medium, adding it to the cache afterwards. Thus data is combined from the two sources.

Rationale: To ensure efficient reading, *getData* can be used as such to fetch data from the cache, if anyhow possible. Only if there is at least one byte not in the cache, the medium must be accessed directly. Updating the cache by read data is useful, to ensure later calls to query the same data can fully leverage the cache and do not need another access to the external medium.

Disadvantages: No disadvantages known

Should we provide a way to force medium access when getting data?

DES 055: *getData* provides no specific mode or configuration to ignore the cache and always read from the medium

getData does not provide a forced read mechanism, but it always only reads data from the external medium that it cannot find in the cache.

Rationale: It is not clear when using code should use the forced or the unforced mode. If the using code somehow can magically see that the underlying medium was changed by other processes, it could do a forced direct read. But in this case, everything is lost already, because you can never know what the other process did. A better way to ensure integrity and consistency would be for the using code to close the medium and reopen another access instance.

Disadvantages: No disadvantages known

We want to define now how the read media data is represented:

DES 056: Read media data is represented as read-only `ByteBuffer`

Read data is not returned in the form of `byte` arrays, but as `ByteBuffer` instances that are read-only.

Rationale: Firstly, users can directly gain profit from conversion functions offered by `ByteBuffer`, on the other hand the implementation is more flexible when it comes to the content of the `ByteBuffer`, as only the bytes between `position()` and `limit()` can be read. Using this e.g. an internally managed, much bigger `ByteBuffer` object can be returned as a read-only view instead of copying it.

Disadvantages: No disadvantages known

For efficiency reasons, we define the following:

DES 057: Handle reading from a single already cached region or a size smaller than the maximum read-write block size as special case

If a read offset and size for *getData* is exactly hitting a single cache region, and this single region fully contains the range to read, the **ByteBuffer** corresponding to this region with adapted position and limit is returned instead of allocating a new **ByteBuffer** and copying bytes, which will happen in any other case.

Similarly, if a range is not cached and needs to be read, and if the range's size is below the maximum read-write block size, again no copying happens, but the read **ByteBuffer** is directly returned.

Rationale: The case of an already cached region fully covering the range to read can be considered as at least the use case covering 80% of the typical uses. The reason is: Usually, the users should use cached media. And usually, they should call *cache* before *getData*, covering e.g. a whole header. Thus, reading individual fields later with *getData* will always hit the cache. Thus, for this 80% case, we could optimize the library performance a lot by simply not copying.

Disadvantages: Slightly more complexity and testing effort involved

Let's discuss the topic of timeouts. In Java, each reading and writing I/O call might block. How to deal with this? The following design decision clearly states it.

DES 058: jMeta does not support any timeouts, neither reading nor writing

Media does not offer the possibility to configure timeouts for any reading or writing actions, and likewise, it does not implement any measures to prevent these actions from blocking arbitrarily long

Rationale: Both for file or stream access, read and write operations, including determining where we are or truncation might or might not block. This is - unfortunately - OS and implementation dependent and cannot be predicted. It is possible and was tried to implement a parallel thread to execute the action against the medium and the main thread to monitor the time taken by the other thread and retrieve the result after a given timeout. While this is easily possible using Java's **Futures**, it is not so easy to really terminate the blocking action and its thread. In reality, on some OSs the action is really interruptible, on others it is not. This might also depend on the implementation. Thus, **jMeta** won't implement anything that would try to convince the user it can actually do it, leading also to increased complexity. Instead, the user must use his own mechanisms, i.e. an own custom **InputStream** implementation, using multithreading and monitoring of calls from the outside, or using implementation-specific timeouts such as for **SocketInputStream**. Furthermore, it is also problematic what state the medium is in after the timeout was detected. What shall the library or the user do with this implementation? Retry, fail, close and retry? There is not always an easy answer to it.

Disadvantages: Users might need to write more custom code themselves, if they strictly require this for mediums they use.

11.1.2 API Design

On the basis of the design decisions made in the previous section, we can now develop an API design for the component **Media**. The API is the public interface of the component, i.e. all classes that can be used by other components to access the **Media** functionality.

Representation of a Medium

The medium has appeared a lot of times already as a term, thus a representation as a class makes sense.

DES 059: Media are represented as interface and implementation class with following properties

A medium is represented as Java interface **Medium** and allows users of **jMeta** to specify a concrete physical medium (i.e. the implementations of the interface **Medium**). As implementations we support a **FileMedium** according to ??, according to ?? a **InputStreamMedium** and according to ?? a **InMemoryMedium**. A medium has the following properties:

- Is random-access: Yes/No - **Motivation:** This property has strong impact on the read and write process, yet it is an intrinsic property of the **MEDIUM** itself and not of the access mechanism. Thus it is directly available for at a **MEDIUM**.
- Read-only: Yes/No - **Motivation:** This property disables writing in practice if set to “Yes”. Some **MEDIA** can never be written (e.g. **InputStreams**), for others it is possible. This flag shall be used to also give the **jMeta** user a possibility to signal he wants to only access read-only.
- Requires caching: Yes/No - **Motivation:** This property states if the medium requires to be backed by a cache or not. Some **MEDIA** such as in-memory media do not require to be backed by a cache and must set this property to No.
- Current length in bytes (only relevant for random-access) - **Motivation:** Java offers queries for each kind of **Medium** except **InputStream**. Thus this should be implemented directly in the **Medium** implementation. For **InputStream** and non-random-access media in general, terms like length to not make much sense. Thus here there is no value, but a constant indicating an unknown length. In spirit pf design decision ??, it is a currently persisted lenght and not a length including any not-yet persisted changes.
- A clear text name of the **MEDIUM** - **Motivation:** This is helpful for identification purposes of the **Medium** e.g. in log output. It can be derived from e.g. a file name, depending on the medium type.
- The “wrapped” object representing the raw medium or its access mechanism, e.g. the file, the **InputStream** or the byte array.

Rationale: It can be controlled in detail which medium types are supported. The user can specify the medium to use in a comfortable way. Further API parts get more easier, as their interfaces must not distinguish between different media types, but rather only use the abstaction that **Medium** offers. Motivation for each of the properties see the listing above.

Disadvantages: No disadvantages known

Due to consistency reasons there are some restrictions regarding the manipula-

tion of media properties:

DES 060: If a `Medium` implementation is writable, it must also be random-access

Every in principle writable `Medium` implementation must be random-access, too.

Rationale: The `jMeta` APIs for writing content can thus concentrate on random-access output media. No separate API design and implementation for output media that are not random-access is necessary. The API gets easier for end-users. Lack of non-random-access output media such as `OutputStreams` can be mitigated via the examples in ??.

Disadvantages: No disadvantages known

Here is a very importante note for byte array media:

DES 061: For byte array media, a writing method for resetting the whole byte array is necessary

The user can reset the bytes of the medium via a public method `setBytes` of class `InMemoryMedium`.

Rationale: It is mostly not harmful to offer the method as public, it is even an advantage for the users, as he can set the bytes himself. Only between registering write operations and a flush, this call leads to unexpected behaviour. This methode is very important for the implementation: Via writing actions, the byte array must be extended or shrunk in some situations. This basically means recreating and copying the array. The method must thus be public, as the corresponding implementation functionality will be for sure placed in another package.

Disadvantages: No disadvantages known

Positions on and Lengths of a Medium

In each case where dare is read from or written to a `MEDIUM`, the question “where?” arises. Usually libraries use integer or long variables to represent offsets. It must be said however: Offsets do not only make sense for random-access media. You could also interpret them as offset since start of reading from an `InputStream`, which is actually the way it is done in `jMeta`. We decide:

DES 062: Byte offsets are used for any kind of MEDIA

Byte offsets that refer to a position on a MEDIUM are used for all media types: random-access and non-random-access. For byte streams they refer to the position of the current byte since start of reading the first byte after opening the stream, which has offset 0. The offset-based reading is simulated as specified in ?? and ??, because when directly reading from an `InputStream` via Java API, offsets are not needed, it is always read from the current position of the stream.

Rationale: We must therefore distinguish between random-access and non-random-access only at a few places in the implementation. Users can use the API uniformly and irrespective of the actual medium type (with restrictions: see ??).

Disadvantages: No disadvantages known

It makes not so much sense to represent offsets only via a primitive data type. Instead, the representation as a user-defined data type offers some advantages:

DES 063: jMeta uses the interface `MediumOffset` to represent offsets on a MEDIUM.

The interface binds both the MEDIUM and the offset on this MEDIUM together, and thus is a kind of “global” address of a byte. Next to reading medium and of offset, it offers some helper methods:

- **behindOrEqual:** Returns true if another `MediumOffset` is located on the same medium behind of at the same position as this instance.
- **before:** Returns true if another `MediumOffset` is located on the same medium before the position of this instance.
- **advance:** Creates a new `MediumOffset` that is located the given number of bytes before (negative argument) or behind (positive argument) this instance. Also zero is possible as argument.

Rationale: We clearly state how the library deals with offsets. We can thus implement some helper functions into the datatype (e.g. validation, offset comparison, advance etc.) which ensure reuse and ease working with offsets in general.

Disadvantages: No disadvantages known

Now we come to a central decision when it comes to dealing with lengths and offsets:

DES 064: jMeta uses long for length and offset specifications, byte is always its unit

In **jMeta**, lengths and offsets are always specified using the Java datatype `long`. The length is in any case the number of bytes, offsets are zero-based, linearly increasing byte offsets.

Rationale: This guarantees uniformity. However, we also want to meet the requirement “?? ??”. Integer with a maximum of 4.3 GB is already too limited, which leaves only `long` as a viable option. The datatype `long` allows for positive numbers up to $2^{63} - 1 = 9223372036854775807$, i.e. approximately $9 \cdot 10^{18}$ bytes, which is 9 exabytes or 9 billion gigabytes. From current point of view, such lengths and offsets for input media, even for streams, should be sufficient for some decades to come. Furthermore, big data chunks are almost in any case subdivided in small units that can be easier handled, and these small units will not have big lengths. Even the Java file I/O uses `long` as offset and length datatype in most cases.

Disadvantages: More memory for saving offsets and lengths is necessary. If we look at the development of storage media, storage needs and processing speed it might be that in 100 years the maximum data volume of `long` will be reached. If **jMeta** is still used in these future scenarios, a change request would be worth it!

A rather seldom special case is dealt with in the following design decision:

DES 065: No special handling of long overflows

For unusual `long` uninterrupted reading from `InputStream` you could think that even when using `long`, it could come to an overflow in some time. This is however very unlikely, thus this case is not treated. The implementation always assumes that the current offset is positive and can be incremented without reaching the max `long` number.

Rationale: Even here it holds true: The datatype `long` allows positive numbers up to $2^{63} - 1 = 9223372036854775807$. Let us assume that an implementation could make it to process 10 GB per second, then it would still need 9 billion seconds, i.e. nearly 30 years, to reach the offset limit and create an overflow.

Disadvantages: No disadvantages known

Now the problem arises that the medium changes due to writing access. How do the offsets change in this case? Is it necessary to update already created **MediumOffset** instances according to the changes on the mediums, or not? If yes, when this needs to happen? In principle, we could see following alternatives:

1. Never update already created **MediumOffset** instances
2. Update already created **MediumOffset** instances directly for each pending change registered (see ??)

3. Update already created `MediumOffset` instances only when an explicit *flush* according to ?? occurs

Assume that `MediumOffset` instances are not updated when writing. That means the user code remembers a position of an element in form of a `MediumOffset` instance, and uses it to read or write data. If e.g. an insert operation takes place on the medium before the offset of the `MediumOffset` instance, then the instance refers to another data byte than before, and thus not anymore to the object it was referring to initially. We should not only think of raw bytes but - as necessary for data formats - *objects*, i.e. parts of the binary data that form a specific unit which has a specific meaning, representing something. Then failure to update the offset is fatal. Code using `Media` locates an object at the wrong place if the medium changed before that offset meanwhile. The second point why we need it: If caching data, this data also refers to `MediumOffsets`. If we do not update them automatically, an additional functionality must be added to perform the updates based on the changes, increasing the complexity of the caching implementation.

To formulate the following design decision a bit easier, the vague term of “object” used above is now defined a bit sharper: An object is a consecutive byte unit starting at a specific offset x and it has a length of n bytes. **remove**, **insert** and **replace** in the offset interval $[x, x + n]$ change these objects, which cannot be in any case specifically treated by `Media`.

DES 066: Media needs to automatically update MediumOffset instances after medium changes: remove and insert

All **MediumOffset** instances ever created for a medium must be updated automatically whenever this medium changes. The kind of update needed is more complex than you would think on first glance.

Let x be the start offset of a collection of n bytes (an “object”), let y be the insert or remove offset and k the number of bytes to insert or remove. Let \bar{x} be the offset of the **MediumOffset** instance after updating. Then the following detailed rules apply:

- **insert before the object start offset x :** Is $y \leq x$, then $\bar{x} := x + k$, including the case $y = x$.
- **insert behind the object start offset x :** Is $y > x$, then $\bar{x} := x$, i.e. such insertions of course do not change the offset of byte x . Insertions might happen before $x + n$, essentially splitting the object.
- **remove before the object start offset x without overlap:** Is $y + k \leq x$, then $\bar{x} := x - k$. Thus k bytes are removed before the object, however the removed region does not overlap with the object.
- **remove before the object start offset x with overlap:** Is $y \leq x$, but $y + k > x$, then the removed region overlaps the object. It is thus a *truncation* of the object starting at front, and it might even reduce the object to length 0. Thus the start offset of the object shifts $x - y$ to be equal to y , i.e. $\bar{x} := y$.
- **remove behind the object start offset x :** Is $y > x$, then $\bar{x} := x$, i.e. the object start offset remains unchanged, of course also in the case that $y < x + n$. In the latter case, however, the object is truncated at its end.

Rationale: **MediumOffsets** are both handed out to the user as well as used to manage cache objects. If changes to the medium happens, all **MediumOffsets** that are already in user hands or are internally used to manage data are rendered invalid. Thus, we would need to write code to detect such **MediumOffsets** as invalid whenever used, or to invalidate the current cache content. Instead of doing so, we simply update all **MediumOffsets** previously created according to the kind of change, thus practically still letting them point to the same “object” as before (despite removing the medium bytes as special case). The **MediumOffsets** can thus be used further on.

Disadvantages: A central management of **MediumOffset** instances must be implemented (see ??).

replaces are discussed in the next design decision:

DES 067: Media needs to automatically update MediumOffset instances after medium changes: replace

For **replace**, the update of already created **MediumOffset**s is similar, but any offsets within the replaced region remain unchanged. In detail: Let x be the start offset of a collection of n bytes (an “object”):

- **replace behind the object start offset x :** The same as insert and remove behind the start offset, i.e. does not change it but might truncate or split an existing object starting earlier.
- **replace before the object start offset x without overlap:** For an inserting replace, The same as for **insert** before x . For a removing replace, the same as for **remove** before x without overlap. For an overwriting replace, no changes are applied to x .
- **replace before the object start offset x with overlap:** The replaced region overlaps with an existing object. Here, x remains unchanged by the replacement.

Rationale: There is no surprise for replaces behind the object start offset and those before the start offset without overlap. The only question is: If there is an overlap, i.e. x lies within the replaced region of bytes, why do we not update it? The reason is: Because of the replacement by other bytes, the object is anyway changed, the correct semantics seems to be to let the **MediumOffset**s point to x still point to x .

Disadvantages: A central management of **MediumOffset** instances must be implemented (see ??).

As already indicated by ?? the automatic updating of already created **MediumOffset** instances requires that only **Media** may create **MediumOffset** instances. These must be managed in a kind of factory to be able to automatically update them in case of writing operations.

DES 068: MediumOffset instances are centrally managed by Media and cannot be directly created by users of the component

The lifecycle of **MediumOffset** instances is controlled by **Media**. They are created and returned to the user via a factory method. The user has no direct access to the implementation class of **MediumOffset**.

Rationale: It is strictly required to implement ?. Instances that have been created by the user cannot be updated automatically, thus we have to ensure the manual creation by the user does not happen.

Disadvantages: More complex instantiation of **MediumOffset** instances.

The question *when* to update **MediumOffset** instances has still not been an-

swered yet. The following design decision clearly defines this:

DES 069: MediumOffset instances are only updated after a *flush*

According to ?? MediumOffset instances are automatically updated in case of medium changes. This automatic update only happens at *flush* time.

Rationale: Assumed that MediumOffset instances would already be updated whenever a pending change is registered using *insert*, *replace* or *remove*. In this case the following would be necessary:

- When reading data, this data is not necessarily in a cache. This indicates that reading from external medium is necessary. If all MediumOffset instances would reflect a state including any pending changes, they would no longer correspond to the state of the external medium. If you would now want to read or write to the external medium, the real offset on the external medium would need to be “reconstructed” based on the changes made so far, everytime you want to know where current data resides on the medium. This implies a complex coding overhead that would not ease debugging errors or understanding the current state of instances.
- The operation *undo* according to ?? requires that offsets would need to be “re-adapted” if a pending changes is undone again. Again this is additional complexity.

If MediumOffset instances in contrast are first updated after a *flush*, then no reconstruction of original offsets based on already made changes is necessary.

Disadvantages: No disadvantages known

This directly implies the following design decision:

DES 070: In Media, offset specifications always are offsets on the external medium as it was looking like after the last *flush* or after opening it initially

For operations of Media that take an offset as argument, this offset refers to a location on the external MEDIUM after the last *flush* or the initial opening - in case no *flush* has occurred yet. These offsets must especially be located within the interval $[0, \text{length}]$, where “length” is the current length of the medium in bytes.

Rationale: Naturally follows from ?. For users, the offset situation remains stable and logical, he does not need to maintain a history of *insert*, *replace* and *remove* operations. Likewise, the offsets stay stable for the implementation, too: Checking offsets and organisation of internal data structures can be based on this invariant.

Disadvantages: No disadvantages known

Semantics of Writing Operations

In ??, we have defined the primitive writing operations *insert*, *replace* and *remove* that are essential for **Media**. In the same section, we have listed basic requirements for writing that lead to these operations. The API of **Media** includes their guaranteed behavior. It is very important to define interrelations between these operations, especially how they behave for overlapping offset areas of the medium. These things are defined by the following design decisions. The behaviors are part of the API contract and must be documented as such for the API users.

We start with how multiple *insert* operations behave:

DES 071: *insert* concatenates insertions in call order

1. Step 1: *insert* at offset x , Step 2: *insert* at different offset: Both calls are allowed, they do not influence each other.
2. Step 1: *insert*, Step 2: *insert* at the same offset: *insert* concatenates, each call with the same offset determines a new, consecutive insertion, i.e. insertions at the same medium location are done with increasing offsets. Let “Call 1” be the earlier call with insertion length n_1 , “Call 2” the later call at x with insertion length n_2 . The end result on the medium after *flush* is:
 - At offset x , the insertion data of “Call 1” is located
 - At offset $x + n_2$, the insertion data of “Call 2” is located.

Rationale:

1. See requirement **AMed01**
2. See requirement **AMed08**

You could alternatively interpret “insert” as such that the second call inserts data *before* the previous earlier one. However, the design decision says that “insert” inserts before the currently persisted medium byte at the insertion offset. Concatenating allows user code to linearly insert stuff with increasing offsets, which is in most cases the convenient and expected behavior.

Disadvantages: The second possible interpretation sketched above might lead to surprises in a few use scenarios.

Before looking at the interactions between *inserts*, *removes* and *replaces* in detail, we have to deeper think about **AMed09**. The requirement is to allow *inserts* and one of these other operations at the same offset in arbitrary order. What is the meaning behind? This is clarified by the following design decision:

DES 072: *insert* at the same offset as *remove* and *replace* means to insert before remove and replace, irrespective of schedule order

No matter if you first call *insert* at offset x , then *remove* or *replace* at offset x , or the other way round: The meaning is that you want to insert new bytes before the bytes to remove or replace.

Rationale: If you want to insert something behind the replacement or removed region, the insertion offset must point to the offset of the first byte behind the replacement or removed region. Thus, calling *insert* after *remove* or *replace* cannot have the meaning “first execute the replace or remove, then insert the bytes behind the removed or replaced region”. Thus, we have to ignore the sequence and ensure *inserts* at the same offsets as *remove* or *replace* are always executed first.

Disadvantages: No disadvantages known

We now look at how *insert* and *remove* interact with each other:

DES 073: *insert* is revoked by overlapping *removes*, *inserts* within already removed regions are not allowed

The following cases can be identified:

1. Step 1: *insert* at offset x , Step 2: *remove* n bytes at offset y , where x is not located within the removed range $[y, y + n)$: Both calls are allowed, they do not influence each other.
2. Step 1: *remove* n bytes at offset y , Step 2: *insert* at offset x , where x is not located within the removed range $[y, y + n)$: Both calls are allowed, they do not influence each other.
3. Step 1: *insert* at offset x , Step 2: *remove* at same offset: Both calls are allowed, with the meaning to first insert new bytes at offset x , then removing existing bytes starting at offset x (see ??).
4. Step 1: *remove* at offset x , Step 2: *insert* at same offset: Both calls are allowed, with the meaning to first insert new bytes at offset x , then removing existing bytes starting at offset x (see ??).
5. Step 1: *insert* at offset x , Step 2: *remove* n bytes at offset y , where x is contained in the removed range $[y, y + n)$: The insert is revoked by the remove, i.e. will not be executed during a *flush*.
6. Step 1: *remove* n bytes at offset y , Step 2: *insert* at offset x , where x is contained in the removed range $[y, y + n)$: The second call is rejected with an exception, as it is pointing into an already removed region.

Rationale:

1. See requirement AMed01 and AMed02
2. See requirement AMed01 and AMed02
3. See requirement AMed09
4. See requirement AMed09
5. See requirement AMed10
6. See requirement AMed13

Disadvantages: Then, how to implement write requirement AMed04? See ??.

Now we are left with clarifying how *insert* and *replace* interact with each other. This is essentially the same as between *insert* and *remove*:

DES 074: *insert* is revoked by overlapping *replaces*, *inserts* within already replaced regions are not allowed

The following cases can be identified:

1. Step 1: *insert* at offset x , Step 2: *replace* n bytes by m new bytes at offset y , where x is not located within the replaced range $[y, y + n)$: Both calls are allowed, they do not influence each other.
2. Step 1: *replace* n bytes by m new bytes at offset y , Step 2: *insert* at offset x , where x is not located within the replaced range $[y, y + n)$: Both calls are allowed, they do not influence each other.
3. Step 1: *insert* at offset x , Step 2: *replace* at same offset: Both calls are allowed, with the meaning to first insert new bytes at offset x , then replacing existing bytes starting at offset x (see ??).
4. Step 1: *replace* at offset x , Step 2: *insert* at same offset: Both calls are allowed, with the meaning to first insert new bytes at offset x , then replacing existing bytes starting at offset x (see ??).
5. Step 1: *insert* at offset x , Step 2: *replace* n bytes by m new bytes at offset y , where x is contained in the replaced range $[y, y + n)$: The insert is revoked by the replace, i.e. will not be executed during a *flush*.
6. Step 1: *replace* n bytes by m new bytes at offset y , Step 2: *insert* at offset x , where x is contained in the replaced range $[y, y + n)$: The second call is rejected with an exception, as it is pointing into an already replaced region.

Rationale:

1. See requirement AMed01 and AMed03
2. See requirement AMed01 and AMed03
3. See requirement AMed09
4. See requirement AMed09
5. See requirement AMed11
6. See requirement AMed13

Disadvantages: Then, how to implement write requirement AMed05? See ??.

Saying all this, it is not possible to implement AMed04 (“it is possible to change or remove parts of an already done insertion before a flush”) and AMed05 (“it is possible to change or remove parts of an already done replacement before a flush”) by using subsequent inserts, removes and replaces to modify an already inserted or replaced block of data before a flush. This is of cause already clear when looking at ??: Offets of these operations always only refer to currently persisted data, so

you e.g. cannot change a prior insert not yet flushed with a remove. Then how are these requirements met? Here is how it could be done:

DES 075: Meeting requirements AMed04 and AMed05 is done outside of Media, yet using its *undo* feature

AMed04 and AMed05 can only be met outside Media. Using code must detect that a parent of the changed field was already scheduled for insertion or replacement. If that is the case, they have to “re-compute” the overall insertion or replacement bytes, undo the previous insertion or replacement using Media and finally schedule a new insertion or replacement containing the changed bytes.

Alternatively, the implementation using Media might freeze a data block for change once its scheduled for insertion or replacement and throws an exception if someone tries to modify it afterwards. The details of this are defined in the design of DataBlocks.

Rationale: Media could somehow support here by offering modification functions for already done *inserts* or *replaces*. However, this would also at least require the using code to detect the change in a parent, and also recompute. It would make the API more complex. Using *undo*, the using code has just slightly more work to do.

Disadvantages: No disadvantages known

We already clarified how insert interacts with insert, remove and replace, in any order. We are now only left to clarify how multiple removes and replaces interact with each other. Multiple *removes* interact with each other as follows:

DES 076: *remove* allows no overlaps by other *removes*, only later calls with bigger region make earlier calls obsolete

The following cases can be identified:

1. Step 1: *remove* n bytes at offset x , Step 2: *remove* m bytes at offset y without any overlaps to $[x, x + n)$: Both calls are allowed, they do not influence each other.
2. Step 1: *remove* n bytes at offset x , Step 2: *remove* $m \geq n$ bytes at offset $y \leq x$ with $[x, x + n) \subseteq [y, y + m)$: The second call to *remove* fully encloses the removed range of the previous *remove* call, so the previous call is revoked, i.e. it will not be executed during a *flush*.
3. Step 1: *remove* n bytes at offset x , Step 2: *remove* m bytes at offset y with $[y, y + m) \subseteq [x, x + n)$: The second region to be removed is fully contained in the region already removed by the previous call. The second call is rejected with an exception.
4. Step 1: *remove* n bytes at offset x , Step 2: *remove* m bytes at offset $y \in (x, x + n)$ with $y + m > x + n$: The second call is an overlapping call, that removes additional bytes behind the first *remove* call, but still includes some bytes of the region removed with the first call. The second call is rejected with an exception.
5. Step 1: *remove* n bytes at offset x , Step 2: *remove* m bytes at offset $y < x$ with $y > x, y + m < x + n$: The second call is an overlapping call, that removes additional bytes before the first *remove* call, but still includes some bytes of the region removed with the first call. The second call is rejected with an exception.

Rationale:

1. See requirement AMed02
2. See requirement AMed10
3. See requirement AMed13
4. See requirement AMed12
5. See requirement AMed12

Disadvantages: No disadvantages known

Interactions between *remove* and *replace* are identical to this, as defined in the following design decision:

DES 077: *remove* allows no overlaps by *replaces*, only later calls with bigger region make earlier calls obsolete

The following cases can be identified:

1. Step 1: *remove* n bytes at offset x , Step 2: *replace* m bytes by r new bytes at offset y without any overlaps to $[x, x + n)$: Both calls are allowed, they do not influence each other.
2. Previous case, exchange order of *remove* and *replace*: Same result
3. Step 1: *remove* n bytes at offset x , Step 2: *replace* $m \geq n$ bytes by r new bytes at offset $y \leq x$ with $[x, x + n) \subseteq [y, y + m)$: The replaced region of *replace* fully encloses the removed range of the previous *remove* call, so the previous call is revoked, i.e. it will not be executed during a *flush*.
4. Previous case, exchange order of *remove* and *replace*: Same result
5. Step 1: *remove* n bytes at offset x , Step 2: *replace* m bytes by r new bytes at offset y with $[y, y + m) \subseteq [x, x + n)$: The second region to be replaced is fully contained in the region already removed by the previous call. The second call is rejected with an exception.
6. Previous case, exchange order of *remove* and *replace*: Same result
7. Step 1: *remove* n bytes at offset x , Step 2: *replace* m bytes by r new bytes at offset $y \in (x, x + n)$ with $y + m > x + n$: The second call is an overlapping call, that replaces additional bytes behind the first *remove* call, but still includes some bytes of the region removed with the first call. The second call is rejected with an exception.
8. Previous case, exchange order of *remove* and *replace*: Same result
9. Step 1: *remove* n bytes at offset x , Step 2: *replace* m bytes by r new bytes at offset $y < x$ with $y > x, y + m < x + n$: The second call is an overlapping call, that replaces additional bytes before the first *remove* call, but still includes some bytes of the region removed with the first call. The second call is rejected with an exception.
10. Previous case, exchange order of *remove* and *replace*: Same result

Rationale: For (1.) and (2.): See requirement AMed02 and AMed03. For (3.): See requirement AMed11; For (4.): See requirement AMed10. For (5.) and (6.): See requirement AMed13. For (7.) to (10.): See requirement AMed12.

Disadvantages: No disadvantages known

Now we are just left to clarify how multiple *replace* operations interact with each other:

DES 078: *replace* allows no overlaps by other *replaces*, only later calls with bigger region make earlier calls obsolete

The following cases can be identified:

1. Step 1: *replace* n bytes by r_1 new bytes at offset x , Step 2: *replace* m bytes by r_2 new bytes at offset y without any overlaps to $[x, x + n)$: Both calls are allowed, they do not influence each other.
2. Step 1: *replace* n bytes by r_1 new bytes at offset x , Step 2: *replace* $m \geq n$ bytes by r_2 new bytes at offset $y \leq x$ with $[x, x + n) \subseteq [y, y + m)$: The second call to *replace* fully encloses the replaced range of the previous *replace* call, so the previous call is revoked, i.e. it will not be executed during a *flush*.
3. Step 1: *replace* n bytes by r_1 new bytes at offset x , Step 2: *replace* m bytes by r_2 new bytes at offset y with $[y, y + m) \subseteq [x, x + m)$: The second region to be replaced is fully contained in the region already replaced by the previous call. The second call is rejected with an exception.
4. Step 1: *replace* n bytes by r_1 new bytes at offset x , Step 2: *replace* m bytes by r_2 new bytes at offset $y \in (x, x + n)$ with $y + m > x + n$: The second call is an overlapping call, that replaces additional bytes behind the first replace call, but still includes some bytes of the region replaced with the first call. The second call is rejected with an exception.
5. Step 1: *replace* n bytes by r_1 new bytes at offset x , Step 2: *replace* m bytes by r_2 new bytes at offset $y < x$ with $y > x, y + m < x + n$: The second call is an overlapping call, that replaces additional bytes before the first replace call, but still includes some bytes of the region replaced with the first call. The second call is rejected with an exception.

Rationale:

1. See requirement AMed02
2. See requirement AMed11
3. See requirement AMed13
4. See requirement AMed12
5. See requirement AMed12

Disadvantages: No disadvantages known

Another problem: How to map the actual data to a data block before a *flush*? The offset alone is not unique anymore in the face of multiple *inserts* at the same offset before their *flush*. If you now want to relate an action (*insert*, *remove*, *replace*) data to its (current or future) offset, the offset alone is not sufficient to

state clearly which action happened first at the offset.

An answer is given in the following design decision that explains how action, data and offset are brought into close relation:

DES 079: Each of the writing operations returns an instance of a class `MediumAction` to describe the action in more detail

This class contains following data:

- The kind of action (*insert*, *replace*, *remove*); **Motivation:** The user must be able to know the kind of change
- `MediumOffset` of the action; **Motivation:** It must be clear where the change happened or must happen
- Data of the action (length or bytes to write); **Motivation:** It must be clear what needs to be changed.
- Number of actually affected medium bytes (0 for *insert*, number of bytes to remove for *remove*, number of bytes to replace for *replace*; **Motivation:** For **replace** only one length is not enough as the number of bytes to replace may be different from the length of the replacement bytes.
- Validity: Handle is already persisted by a *flush* or still pending; **Motivation:** Thus it can be clearly state from outside whether the data must still be persisted and thus must be taken from the instance of the user requires to read them, or the data has already been written to the external medium.

Rationale: According to ??, the user can only read data that is currently persisted.

Nevertheless it is necessary that application code can also re-read data previously registered for writing, but not yet persisted by a *flush*. That now becomes possible with the `MediumAction` class that is returned by *insert*, *replace* and *remove*. Is the `MediumAction` still pending, the application code can return the pending bytes from the `MediumAction`, otherwise by directly accessing the `Media` read functionality to fetch the currently persisted bytes.

Instances of `MediumAction` can ideally be used for internal data management by `Media`.

Disadvantages: No disadvantages known

End medium access

We still need an operation to end the medium access, thus we define:

DES 080: Operation *close* ends the medium access and empties the cache, consecutive operations are not possible on the medium

A user can manually end medium access by calling *close*. It is then no longer possible to access the medium via the closed access way. The user must explicitly reopen the medium to access it again.

Rationale: The implementation works with OS resources such as files that must be closed. Furthermore cache content and other memory is not freed anytime if you could not close the medium. Closing cannot be implemented automatically but must be explicitly called by the user.

Disadvantages: No disadvantages known

The public API of medium access

Based on the previous design decisions we now design the public API of the component **Media**. Until now, we only introduced the classes **MediumOffset**, **Medium** and **MediumAction** as well as some abstract operations to deal with media. How do we offer these operations to users? This is explained by the following design decision.

DES 081: Access to a medium is done using the **MediumStore**

The reading and writing access to a medium (both random-access as well as byte stream according to ??) is offered via interface **MediumStore** with the operations listed in table ??.

Rationale: A further subdivision of functionality into more than one interface is neither necessary nor helpful. It would just be an unnecessary complex API, and despite the single interface, the implementation can still be modularized as needed. The individual operations are motivated in the table itself.

Disadvantages: No disadvantages known

Operation	Description	Features and motivation random-access	Features and motivation InputStream
<code>cache(x, n)</code>	Buffers n bytes according to ?? and ?? permanently in the internal cache, starting at the given offset x . If caching is disabled (specifically for byte array media, see ??), this call is ignored. The method might reach the end of medium which it signals as an exception.	User code can prefetch data to work on it later, he himself gets the possibility to efficiently read and process data.	If x is bigger than the last read end position, the bytes up to the new offset are read and possibly cached. If x is smaller instead, a <code>InvalidMediumReferenceException</code> is thrown. The last read end position is advanced correspondingly.
<code>getCachedByteCountAt(x)</code>	According to ?? and ??, it provides the number of bytes that are consecutively cached at offset x .	See <code>cache</code> . User (and test) code must be able to additionally check whether enough data is already buffered or not.	See random-access.
<code>getMedium()</code>	Returns the <code>Medium</code> instance this <code>MediumStore</code> is associated with.		
<code>getData(x, n)</code>	Reads n bytes at offset x according to ??, ??, ??.	Calls <code>cache</code> to ensure all bytes are cached, then returns the data in a consolidated <code>ByteBuffer</code> .	Same as for random-access media. This indicates that this method also might throw an <code>InvalidMediumReferenceException</code> .

Operation	Description	Features and motivation random-access	Features and motivation InputStream
<code>isAtEndOfMedium(x)</code>	Checks if offset x is at the end of the MEDIUM.	Based on this knowledge, the user code can skip further reading and does not need to check for exceptions.	The provided offset is ignored, it is tried to read bytes from current offset. Is this resulting in return code -1, we are at the end of the stream, otherwise the byte is added to the cache.
<code>insertData(x, data)</code>	Implements the writing operation <i>insert</i> according to <code>??, ??, ??, ??, ??</code> : Adds data at the given offset x , consecutive bytes are shifted “to the back”, the changes first get written only with <code>flush()</code> .	The insertion of new metadata is a common case in <code>jMeta</code> and must be supported.	<code>ReadOnlyMediumException</code>
<code>removeData(x, n)</code>	Implements the writing operation <i>remove</i> according to <code>??, ??, ??, ??, ??</code> : Removes n bytes at offset x . Consecutive bytes are shifted “to front”, the changes first get written only with <code>flush()</code>	Removing existing metadata is a standard case in <code>jMeta</code> and must be supported.	<code>ReadOnlyMediumException</code>

Operation	Description	Features and motivation random-access	Features and motivation InputStream
<code>replaceData(x, n, data)</code>	Implements the writing operation <i>replace</i> according to <code>??, ??, ??, ??, ??</code> : Replaces <i>n</i> bytes at offset <i>x</i> with new bytes of length <i>m</i> . The changes first get written only with <code>flush()</code>	It often happens that – instead of entirely removing or newly inserting data – existing data must be overwritten. This is - especially at the beginning of a file - a much more efficient operation than insertion and deletion and must thus directly be supported.	<code>ReadOnlyMediumException</code>
<code>flush()</code>	Implements the writing operation <i>flush</i> according to <code>??</code> : All <i>changed</i> data currently in the temporary buffer are written in a suitable way to the external medium. It cannot be guaranteed that this operation is atomic.	This is a practical implementation of <code>??</code> : While <code>insertData()</code> , <code>removeData()</code> and <code>replaceData()</code> only write into a temporary buffer, this call directly writes to the external medium.	<code>ReadOnlyMediumException</code>
<code>createMediumOffset(x)</code>	Creates a new <code>MediumOffset</code> instance for the given offset <i>x</i> according to <code>??</code>	Is needed for random-access	Is needed for <code>InputStreams</code>
<code>undo(mediumAction)</code>	Undoes the changes of the given <code>MediumAction</code> according to <code>??</code> , as far as it is still pending.	Is needed for random-access	<code>InvalidMediumActionException</code>
<code>open()</code>	Opens access to the medium	Explicit open instead of constructor better for testing	Explicit open instead of constructor better for testing

Operation	Description	Features and motivation random-access	Features and motivation InputStream
<code>close()</code>	Closes all internal resources according to ??, clears the complete cache contents and other internal data structures	See ??	See ??
<code>isClosed()</code>	Allows users to check if the medium is already closed and thus cannot be accessed anymore (returns true), or it is still accessible (returns false).	See ??	See ??

Table 11.3: Operations of the Media API

The component interface

How are the public functions exposed to the outside world? There must be a functionality that creates a `MediumStore` instance for a given `Medium`. The API for this looks as follows:

DES 082: MediaAPI is the central entry point with creation functions for `MediumStores`

The interface `MediaAPI` offers the central entry point for the component `Media`. Using the method `createMediumStore()`, users can create an `MediumStore` instance.

Rationale: The necessity for a further interface in addition to `MediumStore` is clear enough: A `MediumStore` refers to just a single “medium access session” for a medium, and of course users want to be able to open multiple media at the same time using `Media`. Pushing creation functions into `MediumStore` is not considered good practice as it would decrease comprehensibility.

Disadvantages: No disadvantages known

Error Handling

In general violations of the interface contract according to ?? are acquitted with a runtime exception.

The following table summarizes all further error situations when working with `Media`:

Error Scenario	Description	Reaction jMeta	API method
Medium is already locked	The medium is already locked by another process. <code>jMeta</code> cannot work with the medium. It is the burden of the caller to ensure, that the medium is not used in parallel.	It is an abnormal situation, thus a <code>MediumAccessException</code> runtime exception is thrown.	<code>MediaAPI</code> <code>.createMediumStore()</code>
Unknown media type	A <code>Medium</code> implementation is specified by the caller which is unsupported.	This is an abnormal situation violating the interface contract, thus the same exception as for contract violations is thrown.	<code>MediaAPI</code> <code>.createMediumStore()</code>
Write to read-only medium	The user-provided <code>Medium</code> implementation is read-only, thus it only allows read access.	If the user nevertheless tries to write, this is an abnormal situation and is signalled by a <code>ReadOnlyMediumException</code> runtime exception.	<code>MediumStore</code> <code>.flush()</code> , <code>MediumStore</code> <code>.insertData()</code> , <code>MediumStore</code> <code>.removeData()</code> , <code>MediumStore</code> <code>.replaceData()</code>
Consecutive write calls overlap	Consecutive calls to <code>insertData</code> , <code>removeData</code> or <code>replaceData</code> before a <code>flush</code> overlap in an invalid way, see <code>??</code> , <code>??</code> , <code>??</code> , <code>??</code>	Is acquitted with an <code>InvalidOverlappingWriteException</code> runtime exception.	<code>MediumStore</code> <code>.removeData()</code> , <code>MediumStore</code> <code>.replaceData()</code>

Error Scenario	Description	Reaction jMeta	API method
Invalid cache offset	With <code>cache()</code> it is tried for an <code>InputStream</code> to cache an offset that is smaller than the last read offset.	This is a wrong usage of the API, thus it results in an <code>InvalidMediumReferenceException</code> runtime exception.	<code>MediumStore</code> <code>.cache()</code> , <code>MediumStore .getData()</code>
End of medium during reading	Of course each medium has an end sometimes. When reading, this end can be reached. When writing, this is not actually possible - there, we assume that all output media virtually are unlimited. If there is no more memory for writing, <code>Media</code> usually reacts with a <code>MediumAccessException</code> following an <code>IOException</code> . It cannot be requested from the calling code during reading, that it knows where the end of the input medium is. Reaching its end during reading can be an error, but it needs not - this depends on the current usage situation. The calling code thus must handle this depending on current context.	Because it is not necessarily an abnormal situation, a <code>EndOfMediumException</code> checked exception is thrown.	<code>MediumStore</code> <code>.cache()</code> , <code>MediumStore .getData()</code>

Error Scenario	Description	Reaction jMeta	API method
Stream data not available	Data requested with <code>getData()</code> for a given offset is not available (anymore) in the cache, and the underlying medium is an <code>InputStream</code> .	This is a wrong usage of the API, thus it results in an <code>InvalidMediumReferenceException</code> , a runtime exception.	<code>MediumStore.getData()</code>
Unknown or invalid <code>MediumAction</code>	The user passes an unknown or invalid <code>MediumAction</code> to any operation	This is an abnormal situation and is acquitted with the runtime exception <code>InvalidMediumActionException</code>	<code>MediumStore.undo()</code>
<code>IOException</code> in the implementation	The Java implementation use throws an <code>IOException</code> , at any place where none of the already presented error situations are involved.	It is an abnormal situation, thus a <code>MediumAccessException</code> runtime exception is thrown.	<code>MediumStore.cache()</code> , <code>MediumStore.getData()</code> , <code>MediumStore.isAtEndOfMedium()</code> , <code>MediumStore.flush()</code>
The <code>MediumStore</code> was already closed using <code>close()</code>	<code>MediumStoreClosedException</code> , a runtime exception	<code>MediumStoreClosedException</code> , a runtime exception	all

Table 11.4: Error handling in the component Media

To summarize:

DES 083: Reaching the end of medium is not necessarily an error, other problems with I/O are seen as abnormal events.

Media sees achieving at the end of a medium not as abnormal event, but it must be handled according to the current context by the user code. **jMeta** assumes output media to be virtually unlimited and thus does not implement any means of treating end of medium situations when writing (also in accordance to the Java API).

All other error situations in **Media** are abnormal situations according to table ??.

Rationale: See table

Disadvantages: No disadvantages known

11.1.3 Implementation Design

Management of `MediumOffset` instances

According to ??, `MediumOffset` instances must be maintained centrally by **MediumStore**. At the same time, we can have multiple `MediumOffset` instances referring to the same offset of the same medium, but actually refer to different data - in the special case of the method `insertData`. New data is inserted subsequently at the same offset with `flush()`, see ??. In the end, the `MediumOffset` instances must be automatically updated with `flush()`, as described in ?? and ??.

Thus, these things follow:

DES 084: No pooling of `MediumOffset` instances for the same offsets is possible

In contrast to Java strings, the reuse of `MediumOffset` instances in `createMediumOffset()` for same offsets is not possible, i.e. if an `MediumOffset` instance for offset x has already been created, the same instance cannot be returned for the same offset on the next call. A new `MediumOffset` instance must be created instead.

Rationale: Assume we use pooling. Furthermore, assume there are two inserts of lengths n_1 and n_2 at the same offset x scheduled via `insertData()`. The internal implementation would only have a single instance of `MediumOffset` for offset x . `flush()` must keep the offset of the first insertion unchanged, as the data of this first insertion remains there. However, the offset of the second insertion must be changed, as these inserted bytes are actually located at offset $x + n_1$ after `flush()`.

Disadvantages: For many created `MediumOffset` objects, there could be a bigger memory footprint.

Furthermore, the `MediumOffset` objects must be maintained in a dedicated data

structure:

DES 085: All instances ever created with `createMediumOffset()` are maintained in a dedicated data structure that allows duplicates
All instances are held in a data structure that allows duplicates. It must be dedicated, i.e. only be used for storing all ever created `MediumOffset` objects.

Rationale: Duplicates must be possible due to `??`. We cannot mix that data structure with the caching or pending change data structures, as on the one hand side there will be always more `MediumOffset` instances, than cache entries or pending changes, and on the other hand cache and pending change list could be cleared, while the `MediumOffset` instances must be kept until the explicit close of the mediums.

Disadvantages: No disadvantages known

Are these entries of the dedicated data structure indicated by `??` in any way sorted?

DES 086: Unsorted ArrayList for keeping the MediumOffset instances

We use an `ArrayList` as data structure for maintaining all `MediumOffset` instances. Insertion of `MediumOffset` instances is done in creation order. The list is unsorted.

Rationale: A list allows for duplicates. Sorting of the list after each addition would lead to $O(n \log n)$ runtime complexity on average when creating a new `MediumOffset` instance, if n is the current size of the list. Ascending sort order by offset would be a bit more efficient for `flush()`, but does hardly justify the effort during adding it, because the creation of a `MediumOffset` instance is usually much more commonly done than a `flush()`. Following approaches have been rejected especially:

- A `Map<Long, List<MediumOffset>>` with offsets as key and all `MediumOffset` instances for the offset as value won't work, as the offsets would need to be shifted on each insertion, i.e. the keys of the map would either need to be `MediumOffset` instances again, or would need to be updated expensively.
- A `TreeSet<MediumOffset>` can be excluded as it does not allow duplicates.
- A combination of a `TreeSet` with a special `Comparator`, which assumes `MediumOffset` instances only to be equal, if they are the same objects, and as bigger, if the offset is the same, but the object is different, would accept such "duplicates" with the same offsets and ensure correct sorting for every insert. However, the `Set` then is incompatible with `equals`, what is not a best practice according to the javadocs. Secondly: More time required for insertion as mentioned.

Disadvantages: Finding all `MediumOffset` instances within the `flush()` implementation, that are bigger than a given offset has $O(n)$ complexity, which can however be tolerated.

To handle the complexity of `MediumStore` we decide:

DES 087: The class `MediumOffsetFactory` is used for maintaining `MediumOffset` instances

`MediumOffsetFactory` implements the design decisions ??, ?? as well as ??. It offers following methods:

- `createMediumOffset()` to create `MediumOffset` instances
- `updateOffsets()` for implementation of ??, where a `MediumAction` instance is passed
- `getAllOffsets()` returns all maintained instances
- `getAllOffsetsInRegion()` returns all maintained instances with given offset range
- `getAllOffsetsBehindOrEqual()` returns all maintained instances with offsets bigger than the given offset
- `clear()` removes all maintained instances

Rationale: Reduction of total complexity of the `MediumStore`.

Disadvantages: No disadvantages known

How to deal with the method `advance()` mentioned in ??. It creates `MediumOffset` instances. Do they have to be maintained in `MediumOffsetFactory`, too?

DES 088: The `MediumOffset` instances created with `MediumOffset.advance()` need to be maintained with `MediumOffsetFactory`, too.

Initially, the new `MediumOffset` instance must get a reference to its creator, i.e. `MediumOffsetFactory`, during construction. To still enable a simple creation of instances of `MediumOffset` implementation classes (e.g. in unit tests), the constructor is public.

Rationale: The client code can arbitrarily use `MediumOffset.advance()`, and the returned references can be used as usual, e.g. for newly created data blocks. Thus it is clear that even these references are auto-corrected with `MediumOffsetFactory.updateOffsets()`.

Disadvantages: Close coupling between `MediumOffset` and `MediumOffsetFactory`, as both are knowing each other.

Now we still have the issue with the memory footprint as mentioned in ??, as we cannot pool `MediumOffsets` referring to the same offset. Furthermore, tests so far showed that without any special measures, e.g. for only one insert (and nothing else) to process during `flush`, there are `MediumOffsets` added to the `MediumOffsetFactory`. For the biggest test case having 34 different changes with

a flush, it resulted in a `MediumOffsetFactory` with 4398 reference in case of using a cache. This is still somehow tolerable, but might not scale that good. Imagine you are editing a large file with several gigabytes, adding and removing data in each frame. This would result in probably millions of `MediumOffsets` being stored in the `MediumOffsetFactory`. Thus, there is some motivation to reduce the number of `MediumOffsets` maintained in a `MediumOffsetFactory`. What can we do? We already saw the pooling is no option. The second thing that is not possible is the following:

DES 089: The end `MediumOffset` of a `MediumRegion` must not be stored, but be always recalculated

The end `MediumOffset` of a `MediumRegion` must not be stored within `MediumRegion`, but must always be recalculated whenever `MediumRegion.calculateEndOffset()` is called.

Rationale: The end `MediumOffset` must be managed in `MediumOffsetFactory`, because `MediumRegion` is a public class and you never know how the client uses the created end offset in his code. This would not be a reason against calculating it right at creation time of the `MediumRegion`.

Of course, first of all it is good style to avoid redundancy, as we anyway store start offset and size in the `MediumRegion`.

But the major reason for not storing it is the following case: Imagine an `insert` happens at an offset that is in the middle of an existing cached `MediumRegion`. When storing the end `MediumOffset` of the region, it would be shifted by the number of inserted bytes after a `flush`, while the start offset would not be changed. This would essentially punish the use of redundancy: The end offset would not fit anymore to the size also stored in `MediumRegion`, thus leading to subtle and interesting bugs in connection with flushing changes, e.g. for the `MediumRegions` stored in the cache.

Disadvantages: Each call to `MediumRegion.calculateEndOffset()` creates a new entry to be maintained in `MediumOffsetFactory`

A second thing we unfortunately cannot do is the following:

DES 090: It is not safe to create unmanaged `MediumOffsets` in internal code

Whenever the internal implementation of scheduling or flushing requires `MediumOffsets` advanced, it might be an idea that it only uses a variant of the `MediumOffset.advance()` named `MediumOffset.advanceUnmanaged()` that does not associate the advanced `MediumOffset` instance with the `MediumOffsetFactory`.

However, this will always include the risk that some unmanaged `MediumOffset` “leaps” out to the client or somehow comes into the cache.

Rationale: It might be true that internally created `MediumOffset` for flushing and scheduling never go out to the client user of `Media`. However, still some of the `MediumOffsets` created internally using `advance` might come into the cache. We cannot afford to risk stability and reliability for things that might only occur in “corner cases”. This is extremely hard to debug and puts the overall image of the library at risk.

Disadvantages: No disadvantages known

After having eliminated three possibilities for improvement already, what are the remaining options to keep the number of `MediumOffsets` maintained in `MediumOffsetFactory` to a minimum? Basically, the calls to `MediumOffset.advance()`, especially via `MediumRegion.calculateEndOffset()`, must be minimized in the implementation without sacrificing readability of the code. We define:

DES 091: `MediumRegion` is optimized to create as few new `MediumOffset` as possible

`MediumRegion` only rarely or not at all calls `MediumOffset.advance()`, thus avoiding creation of new managed `MediumOffset` instances. Specifically, a new method `calculateEndOffsetAsLong` is introduced that can be used as replacement for `calculateEndOffset` and does not create new `MediumOffsets`.

Rationale: It showed that `MediumRegion` was responsible for the greater part of `MediumOffsets` created, because it is so heavily used during flushing, offset updates and cache management. In tests we could show that for cached media, it only created six times less `MediumOffsets` using this approach.

Disadvantages: No disadvantages known

Now there is but one topic left: A special case for `MediumOffsetFactory.updateOffsets()`, as indicated in the following design decision:

DES 092: `updateOffsets()` must not increase the `MediumOffsets` auf the `insert MediumAction` causing the shift

According to ??, `insert MediumActions` shift managed `MediumOffsets` with an offset *equal to* or bigger than the insert offset by the number of inserted bytes towards higher offsets. This is necessary for *any other insert* at the same offset, but not for the causing `MediumActions` insert.

Rationale: First of all, only `MediumActions` of type `insert` are affected, as only for these also `MediumOffsets` at the same offset need to be updated.

If we would also shift the causing `MediumActions MediumOffset`, it would not point to the same bytes anymore. E.g. when adding the insertion bytes to the cache after updateing all offsets, we would need to recalculate the original offset again by subtraction.

Disadvantages: No disadvantages known

Medium Access

Access to the MEDIUM is implemented in own classes, as the following design decision states:

DES 093: Interface `MediumAccessor` with implementations for each `MEDIUM` type access

Access to a `MEDIUM` is possible via an interface `MediumAccessor`, doing a great deal of ensuring `??`. It offers the following primitives:

- **`getMedium`**: Returns the medium this instance is working on.
- **`open`**: Opens the medium for access. Opening creates an exclusive lock on the medium, as far as the medium supports this (see `??`, `??`, `??`).
- **`isOpen`**: Check if the medium is opened.
- **`close`**: Close the medium for access. A closed medium cannot be used anymore.
- **`setCurrentPosition`**: Sets the current position for the next `read`, `write`, `truncate` or `isAtEndOfMedium` call. Is ignored for non-random-access media.
- **`getCurrentPosition`**: gets the current position on the medium where the next `read`, `write`, `truncate` or `isAtEndOfMedium` will be executed.
- **`read` at the current position**: Read *n* bytes from external medium by explicit access from the current position, return the bytes in a `ByteBuffer`. Advances the current position by the number of read bytes.
- **`write` at the current position**: Writes *n* bytes to the external medium at the current position passed in a `ByteBuffer`. Throws an exception for read-only media, especially for `InputStream`. Advances the current position by the number of read bytes.
- **`truncate` to new end offset at the current position**: Truncates the medium to end at the current position. Throws an exception for read-only media, especially for `InputStream`.
- **`isAtEndOfMedium`**: Check if the current position is at end of the `MEDIUM`.

There is one implementation of this interface for each distinct `MEDIUM` type. `MediumStore` exclusively accesses the `MEDIUM` only via this interface.

Rationale: `MediumStore` itself can deal with caching and the complex implementation of writing functionality independently from the concrete medium, while the actual medium access can be abstracted away by concrete implementations generically. Classical separation of concerns to increase maintainability and comprehensibility of the solution.

Regarding the offset handling for read: This is necessary to ensure readable and non-confusing use for non-random-access media. In that way, the user does not assume the medium is actually read-only.

Disadvantages: No disadvantages known

After having discussed the implementation of `MediumOffsetFactory` in the last subsection, it should be stated what kind of current positions are stored in an `MediumAccessor` implementation? This is made clear by the following design decision:

DES 094: The current position of an `MediumAccessor` instance is not maintained in `MediumOffsetFactory`

The current position of an `MediumAccessor` instance is not maintained in `MediumOffsetFactory`, such that it won't be changed by a `flush` of changes that would otherwise modify it.

Rationale: It might work with using a managed `MediumOffset` for storing the current position of the `MediumAccessor` implementation. This is because for stream-based media, you cannot flush, and for all random-access media the current position can be arbitrary.

However, making the current position managed would first of all confuse outside users who might expect that the current position is only modified by read and write operations of the `MediumAccessor`, and not in addition by flushing.

Second, we want to keep the memory footprint of `MediumOffset` in `MediumOffsetFactory` to a minimum. Last but not least, we want to avoid a direct dependency from `MediumAccessor` to `MediumOffsetFactory`.

Disadvantages: No disadvantages known

Internal Data Structures for Caching

The cache on first sight maintains data bytes per offset, always assuming that the data in cache is exactly identical to the data on the external medium, according to ???. It was already defined that there is an explicit `MediumStore.cache()` method for reading data into the cache and a method `MediumStore.getCachedByteCount()` for querying connected bytes from the cache (see table ??, as well as ?? and ??). Querying data from the cache is done using `MediumStore.getData()`, which can on the one hand skip the cache, on the other hand it can automatically update the cache with missing data (?? and ??). Finally, the maximum cache size can be set (see ??).

In most cases we want to query which parts of data to read, write or remove is within the cache. The cache might be fragmented arbitrarily due to multiple fill operations. To make the code easier to understand and maintain, we create a specific class for representing the so called regions, i.e. added cache fragments:

DES 095: Class `MediumRegion` for consecutive regions of a medium, especially also for cache regions

The class `MediumRegion` represents a consecutive byte range with start offset, size and contained bytes of a medium, that may or may not be cached. It offers the following methods:

- `getStartOffset()`: Query start `MediumOffset` of the region
- `getSize()`: Query length of the region
- `isCached()`: Returns true if cached, false otherwise
- `getBytes()`: Returns null if the region is not cached, otherwise the `ByteBuffer` with the cached region data
- `isContained()`: Returns true if the given `MediumOffset` is contained within the region, false otherwise
- `overlapsOtherRegionAtBack()`: Returns true if the this region overlaps the other region at its back, i.e. shares bytes with it at its end, otherwise returns false.
- `overlapsOtherRegionAtFront()`: Returns true if the this region overlaps the other region at its front, i.e. shares bytes with it at its front, otherwise returns false.
- `getOverlappingByteCount()`: Returns the number of bytes that this region shares with another region or zero if it does not share any bytes with it.
- `split()`: Splits this region into two regions at the given offset
- `calculateEndOffset()`: Convenience method for calculating the end `MediumOffset` of the region

Rationale: The cache regions and the non-cached regions of the medium can be treated in the same way. Implementing caching is easier when not based on primitive types (e.g. `byte[]`) only, but using this helper class.

Disadvantages: No disadvantages known

Maintenance of cache content is done by a helper class:

DES 096: Cache maintenance by MediumCache, ensuring maximum cache size and no overlapping regions

Cache maintenance is done by the class `MediumCache`. This class has the following invariants at any point in time before and after public method calls:

- The `MediumRegions` stored in the cache do never overlap
- The maximum cache size is never exceeded
- The maximum size of a cache region is never exceeded by any region

It offers the following methods:

- `getAllCachedRegions()`: Returns a list of all cached `MediumRegion` instances currently contained in this cache, ordered by offset ascending
- `getCachedByteCountAt()`: Returns the number of bytes cached consecutively starting from offset x
- `getRegionsInRange()`: Returns a list of `MediumRegion` instances overlapping the offset range $[x, x + n]$, which represent the cached and non-cached ranges within the offset range. I.e. whenever there is a gap in the cache, this gap is also represented by a single `MediumRegion` instance without data
- `addRegion()`: Adds a `MediumRegion` to the cache. Previously cached data is overridden.
- `clear()`: Frees all data in the cache
- `getMaxRegionSizeInBytes()`: Returns the maximum size of a region in bytes. Default is `Integer.MAX_VALUE`, which is represented by a constant named `UNLIMITED_CACHE_REGION_SIZE`. The cached regions will have at most the given size.
- `getMaxCacheSizeInBytes()`: Returns the maximum size of the cache in bytes. Default is `Long.MAX_VALUE`, which is represented by a constant named `UNLIMITED_CACHE_SIZE`. The cache size will at no time exceed this number of bytes.
- `getCurrentCacheSizeInBytes()`: Returns the current size of the cache in bytes.

Rationale: Reducing complexity of `MediumStore`

Disadvantages: No disadvantages known

For the cache sizes, we saw already that we can query the maximum and current cache size as well as the maximum region size. However, why can we not change the maximum sizes dynamically?

DES 097: The maximum cache and maximum region size of a cache instance can never be changed throughout its life time.

The maximum cache size and the maximum region size are passed to the constructor of the `MediumCache` class and they must not be changed later.

Rationale: Otherwise complex methods for reorganizing the cache are necessary, splitting and discarding existing regions. It is very unlikely that the user needs to change the cache size during accessing the medium. It is fully sufficient that he can configure the maximum size of the cache and its regions before the first access to the medium.

Disadvantages: No disadvantages known

How are the cache regions internally managed?

DES 098: A `TreeMap` is used for managing the cache contents, an additional `LinkedList` for freeing up according to FIFO

A `TreeMap<MediumOffset, MediumRegion>` is used for managing the cache contents. An additional `LinkedList` is used for freeing up cache data according to FIFO.

Rationale: Content must be read based on offsets. Thus a data structure sorted by offset is necessary. It allows the efficient retrieval of all `MediumRegions` bigger or smaller than a given offset. This operation should most probably need a runtime complexity of only $O(\log(n))$ instead of $O(n)$.

The `LinkedList` is necessary to efficiently remove the first added cache regions when freeing up is necessary due to max cache size reached, see ??.

Disadvantages: No disadvantages known

We should also think about cache fragmentation. Let us assume that, by subsequent calls to `addRegion()`, we would add a single byte to the cache 20 times for a consecutive offset range. Will this result in 20 different `MediumRegions` with a length of one byte each? The same question arises for calls to `MediumStore.getData()`, which spans over an offset range with gaps in the cache coverage. Assume we have a cache that contains 20 bytes starting at offset x , further 50 bytes at offset $y := x + 30$, i.e. it has a gap of 10 bytes between the two regions. If there is now a call to `MediumStore.getData()` for range $x - 10$ with length of 100, we would result in five regions:

- Region 1: $[x - 10, x)$ not in the cache
 - Region 2: $[x, x + 20)$ in the cache
 - Region 3: $[x + 20, y)$ not in the cache
 - Region 4: $[y, y + 50)$ in the cache
-

- Region 5: $[y + 50, y + 60)$ not in the cache

`MediumStore.getData()` will then add regions 1, 3 and 5 to the cache, according to ???. So, do we have 5 `MediumRegions` in the cache after the call? Another extreme case is that the user reads 20 single bytes each with an offset gap of just one byte to the next one. This would result in 20 cache regions each with a size of one byte.

To ease the implementation complexity, we nevertheless decide:

DES 099: Fragmentation of cache regions is not actively avoided

`Media` does not try to mitigate or avoid the following situations:

- Direct consecutive cache regions are not joined, even if their total size is smaller than the maximum allowed region size
- “Nearby” but unconnected (i.e. non-consecutive) cache regions of small sizes are not handled in any special way.

Rationale: It would lead to even higher complexity of the implementation of `addRegion()`.

In contrast to that high complexity, we can ask the caller to ensure that he only adds data for reasonably connected offset ranges.

Disadvantages: No disadvantages known

Internal Data Structures for Managing Pending Changes

For the two-stage write protocol, changes scheduled via `insertData()`, `removeData()` and `replaceData()` must be managed in a reasonable way, such that they later can be processed during `flush()`. Any change is represented as a `MediumAction`.

We first state the following:

DES 100: `MediumAction` has a schedule sequence number for knowing the schedule order of actions

`MediumAction` defines a schedule sequence number (starting at 0) for distinguishing which action has been scheduled earlier. It is incremented by any scheduled action (no matter which type) by 1. The schedule sequence number is “global” for the same medium in a sense that it is simply incremented for each change, no matter what change it is and if there was a flush already or not.

Rationale: Two distinct `MediumAction` instances of type `insert` referring to the same offset cannot be sorted without this mechanism (see ??).

Disadvantages: No disadvantages known

We now have to define comparisons between two `MediumActions` in a reasonable way, as sorted data structures for efficient retrieval and iteration in order can be used correspondingly:

DES 101: Comparisons of MediumActions is done based on MediumOffset and - for insert only - the sequence number

A MediumAction `a` is smaller than another MediumAction `b` according to `Java's compareTo`, if and only if one of the following criteria are met:

- The `MediumOffset` of `a` is smaller than the `MediumOffset` of `b`,
- Or only in the case that the `MediumOffset` belongs to a MediumAction of type `insert`: if `a` is equal to the `MediumOffset` of `b`, and the sequence number of `a` is smaller than the sequence number of `b`

A MediumAction `a` is equal to a MediumAction `b` according to `Java's equals` and `compareTo`, if all attributes of `a` are equal to all attributes of `b` (in terms of `equals`).

A MediumAction `a` is bigger than a MediumAction `b` according to `Java's compareTo`, if `a` is neither smaller than `b` nor equals `b`. This is especially true if the `MediumOffset` is bigger, or for equal `MediumOffsets` of type `insert`, if the sequence number is bigger. Furthermore it is defined: If `MediumOffsets` and sequence numbers of `inserts` are identical, then `a` is still bigger than `b` in the case that any of the other attribute of the MediumActions differs.

Rationale: Sorting of MediumActions should be done according to their order on the medium (i.e. their `MediumOffsets`) and only for `inserts` in creation order (i.e. their sequence number), because behaviour of consecutive operations is based on call order, according to `??`, `??` and `??`. In addition, the MediumActions must be processed in a defined order during `flush`. `equals` must return true exactly in the case that `compareTo` returns 0.

Disadvantages: No disadvantages known

Now we can define in which way the MediumActions are stored:

DES 102: MediumActions are stored in a sorted data structure without duplicates

The MediumActions created before a `flush()` are held in a datastructure sorted by offset and sequence number (according to `??`), which does not allow duplicates (e.g. `TreeSet`).

Rationale: The sort order is necessary for in-order-processing via `flush()`, two MediumActions with same offset, same sequence number and same type should not show up twice.

Disadvantages: No disadvantages known

Now regarding management of MediumActions:

DES 103: For managing MediumActions, the class MediumChangeManager is responsible

For managing MediumActions, the class MediumChangeManager is defined, which internally implements ??, with following methods:

- **scheduleInsert()** for scheduling an *insert*, implements ??, gets a MediumRegion as parameter and returns a MediumAction of the given type and with correct sequence number, i.e. should there be other actions at the same offset, the new action is guaranteed to have a sequence number higher than the action with the biggest sequence number already present for the same offset.
- **scheduleRemove()** for scheduling an *remove*, implements ??, gets a MediumRegion as parameter and returns a MediumAction of the given type and with correct sequence number, i.e. should there be other actions at the same offset, the new action is guaranteed to have a sequence number higher than the action with the biggest sequence number already present for the same offset.
- **scheduleReplace()** for scheduling an *replace*, implements ??, gets a MediumRegion and the length of the range to replace as parameter, returns a MediumAction of the given type and with correct sequence number, i.e. should there be other actions at the same offset, the new action is guaranteed to have a sequence number higher than the action with the biggest sequence number already present for the same offset.
- **undo()** for undoing actions, implements ??
- **iterator()** returns an `Iterator<MediumAction>` for reading traversal of the changes in correct order, `Iterator.remove()` is not implemented
- **clearAll()** removes all changes

Here, the three **schedule** methods create MediumActions according to ?? and ??.

Rationale: Reduction of overall complexity of MediumStore.

The iterator allows reading of changes in order, but is not needed for processing, as we see later. **remove** on this iterator is not necessary, as **undo()** can undo an action.

Disadvantages: No disadvantages known

At the end, we highlight some commonalities between MediumActions and MediumRegions, and define:

DES 104: MediumAction aggregates a MediumRegion instance

MediumAction aggregate a **MediumRegion** instance which holds that offset and length of the action, BUT NOT the bytes related to the action itself, which are held in a separate attribute of the **MediumAction**.

Rationale: **MediumAction** needs a start **MediumOffset**, a length of the change as well as possibly the bytes to change, if any. However, we only implement start and length in form of a **MediumRegion** instance. You can interpret **MediumAction** as a class that refers to a **MediumRegion**. It is a classical “has a” instead of an “is a” relationship, which requires aggregation instead of inheritance. The reason to not keep the bytes in the aggregated **MediumRegion** but in the **MediumAction** itself lies in the special form of the **replace** operation. For detecting non-allowed overlaps, only the number of bytes to replace is important (see ?? and ??) and not the replacement bytes themselves. In that way, we can get a common implementation of overlap detection.

Disadvantages: No disadvantages known

Implementation of flush

The most complex functionality of **MediumStore** is **flush()**, because:

- **flush()** must iterate all pending changes,
- align them with the current cache content,
- cut reasonable blocks of the data bytes to write,
- update all **MediumOffset** and **MediumAction** instances,
- update the cache

We can add that the writing operations have some additional oddities:

- **insertData()** and **removeData()** require that data behind the insertion or removal offset is read and written again
 - **replaceData()** should also not be underestimated, as depending on the number of existing bytes to replaces with a different number of new bytes, it can either lead to no shifts at all (number of bytes to replace equal to number of replacement bytes), to an **insert** (number of bytes to replace smaller than number of replacement bytes) or a **remove** (number of bytes to replace bigger than number of replacement bytes)
 - If these operations are done at the beginning of files, reading data up to the end of file is possibly not possible in one large chunk, as it could lead to **OutOfMemoryError** for large files, thus block-wise reading is necessary
 - For this block-wise reading, the writing operations differ from each other:
-

- For an `insertData()` of n bytes at offset x , data must be read and written block-wise starting from the end of medium down to offset x . I.e. first the last k bytes of the medium at offset r are read and then written to offset $r + k$, then k bytes at offset $r - k$ are read, to be written at r and so on until offset x . Another way is not working if you do not want to overwrite and thus lose existing bytes.
- For `removeData()` of n bytes at offset x , we must read and write data starting at offset $x + n$ block-wise until the end of the medium. First, k bytes at offset $x + n$ are read and then written at x , then k bytes at offset $x + n + k$ are read to be written to offset $x + k$ and so on, until the last byte of the medium.
- For `replaceData()` we have to distinguish corresponding cases, it could either behave like `insert` or `remove`, or a simple overwrite (if number of bytes to replace equals number of replacement bytes).

We can first confirm that the configuration of a maximum read write block size for writing is necessary:

DES 105: A maximum read-write block size must be configurable for the user

It must be between 1 and N (any integer) bytes

Rationale: Necessary due to the block-wise reading and writing operations behind affected regions caused by inserts and removes. By making it configurable, the user can himself decide how much bytes should be processed in a single pass.

Disadvantages: No disadvantages known

Finding out which operations must be executed during a flush is a complex task. We need to perform this complex task within a method:

DES 106: `createFlushPlan()` in `MediumChangeManager` creates a read-write-plan for a flush in form of a `List<MediumAction>`
`createFlushPlan()` creates a read-write-plan for a flush and returns a `List<MediumAction>`. The read-write-plan contains the actions to be executed in the given order. The list of possible actions is extended by `READ`, `WRITE` and `TRUNCATE`. Which are defined as:

- `READ`: primitive reading of n bytes starting at offset x
- `WRITE`: primitive writing (i.e. overwriting) of n bytes starting at offset x
- `TRUNCATE`: explicit shortening of a file, which is especially necessary for removing data

Each returned `READ` action must be followed by a `WRITE` action, otherwise the plan is invalid. `INSERT` and `REPLACE` operations lead to `WRITE` actions. For all `READ` and `WRITE` actions: the number of bytes is between 0 and the (configured) maximum read-write block size. `createFlushPlan()` also returns the original `REMOVE`, `REPLACE` and `INSERT` actions explicitly in the plan, although they are implemented implicitly by `READ` and `WRITE` actions.

The read-write-plan thus contains `MediumActions` in addition to those caused by scheduling actions by a user. These additional actions are, however, not added to the internal data structures of the `MediumChangeManager`.

The created plan is the basis for processing in `flush` afterwards.

Rationale: Determining the necessary operations is a complex process, which should be done separately. A direct execution of the plan would be an alternative, but testability of the code would heavily suffer.

`MediumChangeManager` is the correct place for this operation, as here all `MediumActions` are managed anyways. The additional `MediumActions` in the plan are not added to any internal data structures to ensure the operation is stateless and ideally repeatable.

Adding another `WRITE` primitive seems unnecessary, as there is already an operation named `REPLACE`. However, `WRITE` differs insofar as the bytes to write might not yet be known when the action is created, in contrast to `REPLACE`.

Disadvantages: No disadvantages known

Next, we should clarify how the cache is involved in `flush()`. In case caching is enabled for the medium, we have the following points to look at:

- If bytes are added to the medium with `INSERT` or `REPLACE`, these bytes must also be added to the cache
- If bytes are removed from the medium with `REMOVE` or `REPLACE`, these bytes must also be removed from the cache
- If bytes behind a change must be read (and written), they should also be directly added to the cache to save future medium accesses in these ranges

- Furthermore, if bytes behind a change must be read, it must also be checked if they are already contained in the cache, such that an explicit **READ** on the external medium is only necessary if these bytes are not cached
- It is not yet clear if a specific handling of **TRUNCATE** is necessary

For the last point, we define the following:

DES 107: READ operations in flush prefer cached data and update the cache if not already contained - We implement this using `getData`

Whenever **flush** processes a **READ** action coming from the flush plan, it first checks if the data to read is fully cached. If it is, it takes the data from the cache. If it is not, it reads the data directly from the external medium and adds it to the cache.

This is exactly what `getData` already does - so we simply call this method. If an **EndOfMediumException** occurs during reading, this is unexpected and will lead to an **IllegalStateException** thrown.

Rationale: We want to use cached data as much as possible and avoid unnecessary medium accesses.

Disadvantages: No disadvantages known

Correspondingly, this is how we treat **INSERT** and inserting **REPLACE** in **flush** regarding caching:

DES 108: INSERT and inserting REPLACE operations in flush add their data to the cache

Whenever **flush** processes an **INSERT** or **REPLACE** action coming from the flush plan, it adds all bytes to insert or replacement bytes from these actions to the cache.

Rationale: We want to use cached data as much as possible and avoid unnecessary medium accesses.

Disadvantages: No disadvantages known

Very similar for **REMOVE** and **REPLACE**:

DES 109: REMOVE and REPLACE operations in flush remove their removed or replaced ranges from the cache

Whenever `flush` processes a `REMOVE` or `REPLACE` action coming from the flush plan, it removes all bytes to remove or replace from the cache.

Rationale: Avoid cache corruption if old data is still present. In best case, the code directly fails when trying to add new data, in worst case the stale removed data remains in the cache and is later returned by read operations despite the medium already changed.

Disadvantages: No disadvantages known

Do we have to do anything with the cache regarding a `TRUNCATE` action? This is answered in the following design decision:

DES 110: For a TRUNCATE, nothing is to be done related to caching

The cache does not require any updates when processing a `TRUNCATE` action in flush.

Rationale: Although it first might seem that, because `TRUNCATE` shortens the medium, no cache updates are necessary for any cached bytes at the end of the medium. This is because the cache is implicitly cleaned by handling the `REMOVE` that leads to the truncate. This consists of removing cached bytes for the removed region as well as updating `MediumOffsets` of any data cached behind the removed region.

Disadvantages: No disadvantages known

We still have the topic to update any already created `MediumOffsets` when processing `INSERTs`, `REMOVEs` and `REPLACEs` according to design decision ???. When do we need to perform these changes? The first part of the answer is given in the following design decision:

DES 111: flush processes the flush plan in two phases: Medium access phase and cache update phase

First of all, flush iterates all `MediumActions` in the flush plan in order and executes any `READ`, `WRITE` and `TRUNCATE` operations. So this can be called the *medium access phase*. In the second phase, all `MediumActions` are iterated again, while the cache is updated, the `MediumActions` are set to done and the `MediumOffsets` are updated in correspondence with the changes.

Rationale: Why not performing everything within the same loop, i.e. just once iterating the `MediumActions`? This would mean to e.g. first execute an `INSERT` action, and then update all references. The problem: The flush plan might still contain `READ` and `WRITE` operations which are referring the offsets on the current medium. Calling `MediumOffsetFactory.updateOffsets()` already now would also update these references, and they would be pointing to the wrong bytes! Thus, we must first process all `READ` (including adding of the read bytes to the cache), `WRITE` and `TRUNCATE` operations with their original offsets. For the caching updates in terms of `INSERT`, `REPLACE` And `REMOVE`, there are additional constraints that make it necessary to do it in the second phase. These are detailed in the next design decision.

Disadvantages: No disadvantages known

The second part of the answer when to perform `MediumOffsetFactory.updateOffsets()` comes here. Some things need to be kept in mind when updating the cache, which depends on the action type:

DES 112: For INSERTs and the replacement bytes of a REPLACE, MediumOffsets are updated *BEFORE* the cache is updated, for REMOVEs and the replaced bytes of a REPLACE they are updated *AFTERWARDS*. For INSERTs and REPLACEs, any MediumOffsets already created and maintained in the MediumOffsetFactory are updated first, which essentially increases the offsets of any MediumOffsets behind the insertion or replacement offset by the number of bytes inserted or the number of replacement bytes. Only after this, the data newly inserted or the replacement bytes can be added to the cache. In contrast to this, when REMOVEs and the replaced region of a REPLACE are processed, first of all, the bytes to remove or replace, if present in the cache, need to be removed from it. Only then, all MediumOffsets behind the remove or replace offset need to be decreased by the number of removed or replaced bytes (or fall back to remove or replace offset, if within the removed or replaced range).

Rationale: For INSERTs: If we add the inserted (or replacement) data to the cache before updating reference, the following would happen: Assume there are bytes already cached directly starting at the insertion offset. They would be superseded by the inserted data, i.e. thrown out of the cache. Then only all remaining data with higher offsets in the cache would be updated. This way, we would essentially lose cached data unnecessarily and increase the complexity of cache management needed at runtime a bit.

For REMOVEs: If we would first update the MediumOffsets here, the following could happen: Assume we have two consecutive cached regions within the region to remove. For both regions, its offset would be updated to point to the remove offset. So we would have two regions cached with the same start offset, which leaves the cache inconsistent. Instead, we first remove all cached regions within the removed range, then we update any remaining regions (those behind the removed range).

Disadvantages: No disadvantages known

The question for the complex `flush` operation now is: Do we need to do anything special for error handling? First of all we should think of: What could happen? In detail, there are two reasonable scenarios what could go wrong:

- There is a programming error in the library, and this leads to a runtime exception
- There is a problem during accessing the medium, may it be a HD failure or that another process accesses the medium in parallel

Both cases can be considered fatal. Would a retry make sense? Probably yes, but most often: Probably no. Should we support in performing a retry? Probably also no:

DES 113: No specific error handling is done in flush itself

There is no try-catch block within `flush` trying to handle any runtime errors. These errors are directly thrown up to the using code. The `flush` implementation does not perform any specific measures to ensure or allow retry or “recovery” mechanisms.

Rationale: First of all, flush does not and cannot really support ACID, see ?? . Second, the error situations one could think of are in no way expected, but refer to where unexpected, abnormal situation. In these situations, anyways the state of the overall system might be critical. We cannot foresee any possible error situations and we also cannot really pretend we can handle them correctly. The only advise for the user when an error in flush happens: If it might be a programming error, issue a bug. Otherwise, try closing and reopening the medium and redo the changes. In both cases, of course, due to ?? , it could be that some changes have already been performed, which might leave the external medium in a corrupted state.

Disadvantages: No disadvantages known, as we hope there are really no cases where a retry would absolutely make sense

Now we have all credentials to implement writing in `flush()`:

1. Create the flush plan according to ??
 2. **Phase 1 - Medium access phase:** Iterate all actions of the flush plan in order, but only handle the action types mentioned below in a specific way:
 - If action = `READ`: Call `getData` at the indicated offset to read n bytes. These bytes are written in the subsequent `WRITE` action. If an `EndOfMediumException` occurs, this is considered as impossible and an `IllegalStateException` is thrown. Store the read bytes to be used by the upcoming `WRITE` action. Set the action to done afterwards.
 - If action = `WRITE`: If the `WRITE` action contains bytes already, it as associated to an `INSERT` or `REPLACE`. In this case, directly write the contained bytes with direct access to the medium via `MediumAccessor`. If the action does not contain any bytes, the previous action must have been a `READ` reading exactly the same amount of bytes to write. If this is not the case, an `IllegalStateException` is thrown. Otherwise these previously read bytes are written with direct access to the medium via `MediumAccessor`. Set the action to done afterwards.
 - If action = `TRUNCATE`: Execute a truncation of the medium with direct access to the medium via `MediumAccessor`. Set the action to done afterwards.
 3. **Phase 2 - Cache update phase:** Iterate all actions of the flush plan in order, but only handle the action types mentioned below in a specific way:
-

- If action = INSERT: First call `MediumChangeManager.undo` on the action. If there is an existing cache region containing the insertion offset, we have to split this region now at the insertion offset. The reason is: Otherwise the already cached portion of the cached region behind the insertion offset would not be shifted by the number of insertion bytes, and the cache would get corrupted. Then call `MediumOffsetFactory.updateOffsets()` for the action. Finally add the insertion bytes to the cache. Here it must be noted that, as the cache add is done after `updateOffsets()`, the insertion offset was increased already before, and we must decrease it again to add the cached bytes at the original insertion offset!
- If action = REMOVE: First call `MediumChangeManager.undo` on the action. Then remove the bytes in the removal region from the cache. Finally call `MediumOffsetFactory.updateOffsets()` for the action.
- If action = REPLACE: First call `MediumChangeManager.undo` on the action. Then remove the bytes in the replaced region from the cache. Finally call `MediumOffsetFactory.updateOffsets()` for the action, then add the replacement bytes to the cache.

DES 114: `flush()` is implemented according to the process defined above

`flush()` is implemented according to the process defined above

Rationale: The flush plan must contain all operations explicitly, i.e. including those triggered by the user - REPLACE, INSERT and REMOVE, even if READ and WRITE would be sufficient for their implementation. The reason for that is that the actions of the user must explicitly be removed from the `MediumChangeManager` and their influence on `MediumOffset` instances behind must explicitly be executed. For this, you need the concrete types, WRITE is not sufficient.

Throwing an `IllegalStateException` in case of end of medium encountered is necessary, as the creation of the flush plan must have considered the medium size. Thus, if it occurs, this is either a programming error in `createFlushPlan` or another process as changed the medium meanwhile.

`undo()` is only executed for user operations, as only those are maintained in the internal data structures of `MediumChangeManager`, according to ???. Also, `undo()` must be executed first, otherwise `updateOffsets()` changes offsets which would lead to a serious problem in a `TreeSet` as used in `MediumChangeManager` to find the correct reference with `contains`.

The cache update is divided into two parts as mentioned in previous design decisions.

The cache management is completely done within `MediumCache.addRegion()`, such that maximum size and consecutive regions can be optimized there.

Disadvantages: No disadvantages known

DES 115: Flushing progress is logged on DEBUG level

Flushing and probably even scheduling of changes is logged on DEBUG level

Rationale: The process of flushing is the most complex process in **Media**. If there are any subtle bugs for end-users, they might be incredibly hard to track down without proper logging.

Disadvantages: No disadvantages known

Implementation of createFlushPlan

The implementation of `createFlushPlan` is anything but trivial. The general algorithm is formally described in a separate paper (including examples), see [?].

Here, we only list some testcases that should be implemented to demonstrate the intended behaviour, see table ?? . If we say

contains read/write-blocks from $\langle startOfs \rangle$ to $\langle endOfs \rangle$ $\langle backwards/forwards \rangle$
in the table, we are referring to the following:

- The generated flush plan contains a sequence of **READ** and **WRITE** pairs starting at **start offset**, each **WRITE** following a **READ** with the same size.
- Each **READ** and each **WRITE** have a maximum size $s \leq m$, where m is the configured maximum read-write block size according to ??
- With $N := endOfs - startOfs$, there are exactly

$$\frac{N}{m} + N \bmod m$$

such pairs present

- They are either read “backward”, i.e. chunk-wise starting at highest offset down to the lowest offset of the range, or “forward”, i.e. chunk-wise starting at lowest offset up to the highest offset of the range

Also note that these tests do not contain any negative cases for scheduling actions that are overlapping in any invalid way. This must already be tested for the schedule operations of `MediumChangeManager`.

ID	Testcase	Variations	Expectation
CF0	No operation	-	The plan created is empty
CF1	Single insert n bytes at offset x	<ul style="list-style-type: none"> a. x is start, intermediate or end offset of the medium b. Bytes behind: None, or whole-numbered multiple of the maximum read-write block size m, or no whole-numbered multiple of m, or fewer bytes than m c. Insertion bytes: The same cases as for the bytes behind 	<ul style="list-style-type: none"> • No read/write operations before insert offset • Bytes behind remain unchanged and get shifted by n towards higher offsets, such that medium length grows by n • Contains read/write-blocks from x to the end of medium backwards • The insert action follows after these actions in the plan

ID	Testcase	Variations	Expectation
CF2	Single remove n bytes at offset x	<p>a. x is start, intermediate offset or $x + n$ is the end offset of the medium</p> <p>b. Bytes behind: None, or whole-numbered multiple of the maximum read-write block size m, or no whole-numbered multiple of m, or fewer bytes than m</p> <p>c. Extreme case: All bytes of the medium are removed</p>	<ul style="list-style-type: none"> • No read/write operations before remove offset • Bytes behind remain unchanged and get shifted by n towards smaller offsets, such that medium length shrinks by n • Contains read/write-blocks from x to the end of medium forwards • The remove action follows after these actions in the plan • At the end there is a truncate operation

ID	Testcase	Variations	Expectation
CF3	Single replace of n bytes by r new bytes at offset x	<ol style="list-style-type: none"> x is start, intermediate offset or $x+n$ is the end offset of the medium Bytes behind: None, or whole-numbered multiple of the maximum read-write block size m, or no whole-numbered multiple of m, or fewer bytes than m Replacement bytes: Same cases as for bytes behind r smaller, equal to or bigger than n Extreme case: All bytes of the medium are replaced 	<ul style="list-style-type: none"> No read/write operations before replace offset If $n > r$: Bytes behind remain unchanged and get shifted by $n-r$ towards smaller offsets, such that medium length shrinks by $n-r$, If $r \geq n$: Bytes behind remain unchanged and get shifted by $r-n$ towards higher offsets, such that medium length grows by $r-n$, Contains read/write-blocks from $x+n$ to the end of medium forwards (if $n > r$) or backwards (if $r > n$) The replace action follows after these actions in the plan If $n > r$: At the end there is a truncate operation

ID	Testcase	Variations	Expectation
CF4	Multiple inserts of in total n bytes	<ul style="list-style-type: none"> a. At same or different offsets b. Bytes in-between/behind: None, whole-numbered multiple of the maximum read-write block size m, or no whole-numbered multiple of m, or fewer bytes than m c. Bytes to insert: Same cases as for bytes behind 	<ul style="list-style-type: none"> • No read/write operations before the first insert offset • Bytes behind last insertion remain unchanged and get shifted by n towards higher offsets, such that medium length grows by n • Bytes in between insertions: remain unchanged and are shifted to the back by the number of up-to-then inserted bytes • Each range between or behind inserts contains read/write-blocks from insertion offset x_i to the start of the next insertion (or end of medium) backwards • Each insert action follows after these block actions in the plan

ID	Testcase	Variations	Expectation
CF5	Multiple removes of in total n bytes	<ol style="list-style-type: none"> All removes refer to consecutive regions or have gaps between Bytes in-between/behind: None, whole-numbered multiple of the maximum read-write block size m, or no whole-numbered multiple of m, or fewer bytes than m Extreme case: All bytes of the medium are removed 	<ul style="list-style-type: none"> No read/write operations before the first remove offset Bytes behind last remove remain unchanged and get shifted by n towards smaller offsets, such that medium length shrinks by n Bytes in between removes: remain unchanged and are shifted to the front by the number of up-to-then removed bytes Each range between or behind removes contains read/write-blocks from offset $x_r + n_r$ up to the start of the next remove (or end of medium) forwards Each remove action follows after these block actions in the plan At the end there is a truncate operation

ID	Testcase	Variations	Expectation
CF6	Multiple replaces of in total n bytes by in total r new bytes	<ul style="list-style-type: none"> a. All replaces refer to consecutive regions or have gaps between b. Bytes in-between/behind: None, whole-numbered multiple of the maximum read-write block size m, or no whole-numbered multiple of m, or fewer bytes than m c. Replacement bytes: Same cases as for bytes behind d. r smaller, equal to or bigger than n e. Extreme case: All bytes of the medium are replaced 	<ul style="list-style-type: none"> • No read/write operations before the first replace offset • If $n > r$: Bytes behind remain unchanged and get shifted by $n - r$ towards smaller offsets, such that medium length shrinks by $n - r$, • If $r \geq n$: Bytes behind remain unchanged and get shifted by $r - n$ towards higher offsets, such that medium length grows by $r - n$, • Bytes in between replaces: remain unchanged and are shifted to the front or back by the number of up-to-then replaced minus inserted bytes • Each range between or behind replaces contains read/write-blocks from offset $x_r + n_r$ up to the start of the next replace (or end of medium) forwards or backwards • Each replace action follows after these block actions in the plan • If $n > r$: At the end there is a truncate operation

ID	Testcase	Variations	Expectation
CF7	Multiple removes and inserts	<ul style="list-style-type: none">a. At same offsets, remove first then consecutive insert or at different offsetsb. Mutually compensating or notc. First remove, then insert, or first insert, then remove - Note that remove and insert in any order at the same offset can be considered consecutived. Extreme case: All bytes of the medium are removed, then new bytes inserted	<ul style="list-style-type: none">• For read/write blocks and bytes between and behind: Same behaviour as for other test cases• For mutually compensating actions: Bytes behind the last action must not be read or written• The order of actions does not matter• For the extreme case: All bytes are actually “replaced” by the new inserted bytes

ID	Testcase	Variations	Expectation
CF8	Multiple replaces and inserts	<ul style="list-style-type: none"> a. At same offsets, replace first then consecutive insert or at different offsets b. Mutually compensating or not c. First replace, then insert, or first insert, then replace - Note that remove and insert in any order at the same offset can be considered consecutive d. r smaller, equal to or bigger than n e. Extreme case: All bytes of the medium are replaced, then new bytes inserted 	<ul style="list-style-type: none"> • For read/write blocks and bytes between and behind: Same behaviour as for other test cases • For mutually compensating actions: Bytes behind the last action must not be read or written • The order of actions does not matter • For the extreme case: All bytes are actually “replaced” by the replacement bytes as well as the new inserted bytes

ID	Testcase	Variations	Expectation
CF9	Multiple removes , replaces and inserts	a. At same offsets b. Mutually compensating or not c. In different orders d. Growing or shrinking the medium	<ul style="list-style-type: none"> • For read/write blocks and bytes between and behind: Same behaviour as for other test cases • For mutually compensating actions: Bytes behind the last action must not be read or written • The order of actions does not matter

Table 11.5: Test cases for checking `createFlushPlan`

Configuring Medium Access

So far, we have identified several mechanisms the user can use to influence the behaviour of medium access: He can set the maximum cache size and he can also change the maximum read-write block size. But how? In ??, we explicitly said that for now, `jMeta` will not offer the possibility to change such parameters dynamically. Still we have to offer the user the corresponding API without exposing him to the internal details of the implementation.

We thus first define:

DES 116: The configuration of the medium access parameters is done per `Medium` instance

The configuration of `Media` is done per `Medium` instance. By doing this, all configuration parameters correlate to an `Medium`, have the same lifetime and scope. The setting of the parameters is done by passing them directly to the constructor of the media created by the user, with constructors setting the parameters to sensitive defaults.

Rationale: The whole internal implementation of the most important classes, i.e. `MediumStore`, `MediumAccessor`, `MediumCache` and so on works on exactly one medium. The user can create `Medium` instances directly himself and configure them as needed, independent of other `Medium` instances. As these parameters must not change after creation of the medium, it is correct to pass them to the constructor and to not offer setters.

Disadvantages: No disadvantages known

Now the question arises: Which configuration parameters are needed in detail? We summarize them again in detail in table ??.

First of all, we exclude one potential parameter:

DES 117: The maximum cache region size is not configurable by the user, but is automatically set to the maximum read-write block size

The user can only configure the maximum read-write block size, but not the maximum cache region size which is just set to the configured maximum read-write block size.

Rationale: For a usual user, it is not clear what the maximum cache region size really is and which size it should have; there is not much “tuning” potential in setting it to another value than the maximum read-write block size. By doing so, we have usually a one to one match between a read block of bytes and the medium region added to the cache. This is good performance-wise, as it does not require to e.g. read the content of a not-yet-cached medium region below maximum cache region size block-wise, in case the maximum read-write block size is smaller than the maximum cache region size. Instead, we just automatically set them to the same value, the user can only influence the maximum read-write block size. It is clearly documented to the user what it does. As additional plus, the configuration interface of a medium gets easier.

Disadvantages: No disadvantages known

Medium	Parameter Name	Type	Default Value	Description
File, InputStream	maxCacheSize	long > 65535	1 MB	Sets the maximum cache size according to <code>??</code> . Changes to this parameter do not have any effect after the <code>MediumStore</code> was already created, so callers need to ensure to set this before creating the <code>MediumStore</code> .
File, InputStream, byte-Array	maxReadWriteBlockSize	int > 0	8192	The maximum size of read-write-actions in bytes, that is triggered by <code>INSERTs</code> or <code>REMOVEs</code> during a <code>flush()</code> , see <code>??</code> .

Table 11.6: Configuration parameters of medium access

11.2 DataFormats Design

In this section, the design of the component **DataFormats** is described. Basic task of the component is to describe the structure of a container or metadata format in a way that is useful for parsing them generically.

11.2.1 Format Comparison

Before any real design decisions can be found, we have to look what the formats are that we try to generically describe with the **DataFormats** component. The main source of this whole chapter is [?].

We fully focus on the following data formats here, ignoring any other possible formats as out of scope:

- Multimedia container formats: Ogg, QuickTime, MPEG Layer I Elementary Stream (for MP3), Matroska, RIFF
- Metadata container formats: ID3v1 and ID3v1.1, ID3v2.2, ID3v2.3, ID3v2.4, VorbisComment, Lyrics3v1, Lyrics3v1, APEv1 and APEv2

As first introduction, we have to compare those formats according to their most important structural aspects. These aspects are as follows:

- Basic encoding they use: **Byte orders** and **character encodings** supported
- Building blocks: What are the **containers** in this format? Are there other sub-structures? Can they be embedded in each other? Is there a hierarchical structure with arbitrary nesting? For metadata container formats, we call the elements actually holding the metadata *attributes*, which is nothing else than a special case of a container.
- Headers and Footers: Despite the payload, are there **headers** and/or **footers** in the containers or attributes? Are they fixed or variable size?
- Lengths of containers: How is the **length of containers** or of their payload defined?
- Lengths of attributes: How is the **length of attributes** for metadata container formats defined?
- Padding: Does the format have some special way of supporting **padding**?

The following table provides all these properties for multimedia container formats, skipping attribute topics:

TODO Table for multimedia formats tab:DFcompareMult

The following table compares these properties for the metadata container formats considered:

TODO Table for metadata formats tab:DFcompareMeta

Note that ID3v2x is a very special format in multiple dimensions: It is incredibly overloaded with a lot of features you might never dare dream to use, say: over-designed. It offers e.g. support for encryption, grouping, compression of tags, embedded lyrics or audio guide codec data which basically contain spoken versions of the metadata stored and other things. On top of that, it has a strange conversion scheme called the *unsynchronization scheme*.

11.2.2 The Container Metamodel

In order to give a feasible model of the structure of a container format (which includes metadata container formats), a metamodel of the typical data structure of a container format is developed here.

The metamodel is ultimately shown in the following figure, the corresponding design decisions will follow:

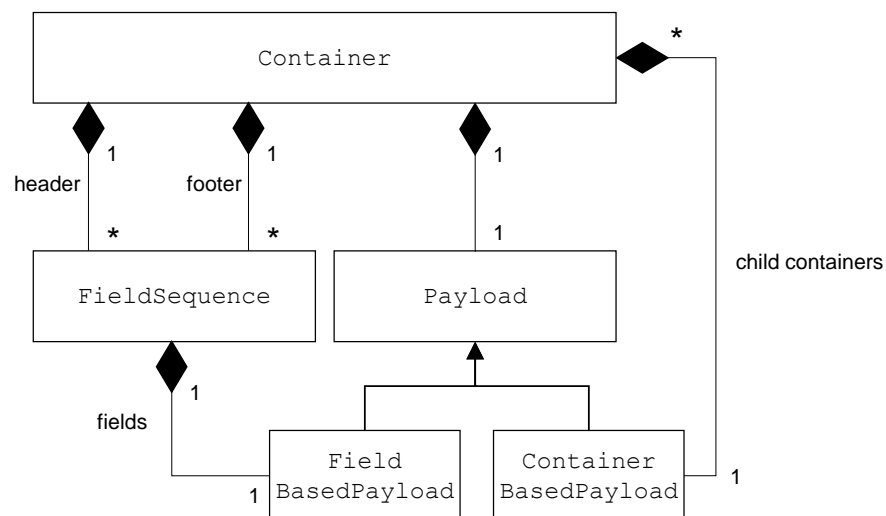


Figure 11.1: The container metamodel

DES 118: Containers as basic top-level unit, containing nested containers

All data that can be processed by **jMeta** consists of containers. A container is a sequence of bytes with a specific structure belonging to exactly one container data format. However, consecutive containers might belong to different data formats.

In addition, a container - as the name suggests - contains other data. This data might be containers of the same format again. The nesting might be arbitrarily deep. A nested container is able to contain data of other formats.

Rationale: Every format considered here has such basic building blocks, called differently in each format specification (atom, frame, page, element, item, tag etc.). The same file or media stream might consist of containers belonging to different formats. One example is an MP3 file with an ID3v2 tag. See also [?], part IV for more details.

Most of the supported data formats have some notion of nesting child containers (or, for metadata formats: attributes) within parent containers. Those child containers have either exactly the same structure as the top-level containers, or they differ slightly. Some formats such as RIFF explicitly define specific containers that embed data belonging to other metadata formats.

Disadvantages: No disadvantages known

How are these containers structured in detail?

DES 119: A container consists of an optional header, payload and an optional footer

Every container - no matter which data format - follows the same basic structure: One or more optional headers start the container's byte sequence. Following this, there is a sequence of bytes called the payload. At the end of the container, there are one or more optional footers. A container must either have at least a single header or at least a single footer.

Rationale: Considering table ?? and ??, every container in every format has at least one header or one footer, including nested containers. Only ID3v23 as sole exception has a second, so-called extended header. Thus we have to support multiple headers and footers, even though this might be a singular case. But future formats might come to bring multiple headers, too. Modelling the middle part of a container as payload explicitly instead of just as list of sub-containers or fields makes sense because it allows to lazily read substructures, or entirely skip the payload.

Disadvantages: No disadvantages known

We saw that containers have a structure, and containers can live within other containers. But on the lowest level, there are leaf nodes, the fields:

DES 120: Fields as the atomic unit of all byte sequences

Fields are the leaf nodes of the container metamodel: A field is a sequence of 1 to N consecutive bytes. A field has a **binary value** as well as an **interpreted value**, i.e. a human-understandable value with a specific meaning. We say that a field has a specific type, which basically describes the mapping between binary and interpreted value as well as the allowed values or format of a field.

Rationale: The term “field” is quite commonly used in binary data formats. It does not make sense to go down to individual byte or bit level. The smallest level of semantic in every binary or textual format is a field.

Disadvantages: No disadvantages known

Now towards headers and footers: Are they inherently different or basically the same? Fixed or variable size?

DES 121: Headers and footers as sequences of fields are structurally identical

Headers and footers are sequences of fields, i.e. they must not nest a container. They are either variable or fixed-size. For both header and footer, the same model class called `FieldSequence` is used.

Rationale: No data format specifies a header or footer nesting a container-like data structure, they all consist of plain fields. Despite the location (before or behind the payload), there is nothing really distinguishing headers or footers from a structural perspective. Thus the same model class can be used for both. However, there is no english term for that class that seems to clearly be a grouping term for headers and footers.

Disadvantages: No disadvantages known

So far, we just introduced containers, fields, headers, footers and payload. We learned that containers might embed other containers, while headers and footers must only consist of fields. But what about the structure of the payload of a container? Of course, if the headers and footers cannot embed child containers, they must hide within the payload. Is this the only possibility? Of course not:

DES 122: Payload either solely consists of containers or fields

The payload of a container either fully consists of a sequence of child containers or of a sequence of fields, but never both.

Rationale: Ultimately, we must boil payload down of fields, or put otherwise: We cannot recurse indefinitely into sub-containers. So at last, there must be a final sequence of child containers having payloads solely consisting of fields. So it is clear that we have both cases. But what about the mixed case: Should we have containers in the model that have both child fields neighbouring child containers? In rare cases such as - again - ID3v2, it would be possible to think of having the frames as child containers of the tag, followed by the padding child field of the tag behind it. However, we avoid such complexities, but payload is either `ContainerBasedPayload` or `FieldBasedPayload`. See next design decision of how we handle the ID3v2 special case.

Disadvantages: No disadvantages known

How to handle padding? If we look at the data format comparisons in ?? and ??, we see that padding is just a special case in ID3v2 and MP3. All other formats solve it “more naturally” by embedding it in special child container. ID3v2 just adds some nullbytes at the end of the tag payload behind the frames. How to model this? There are basically three possibilities:

- The padding is a part of the payload, basically a field. This seems most intuitive, but is already excluded by the previous design decision, because then we would need to support mixed payloads, consisting both of child containers and fields.
- The padding is part of the payload, but considered as a very special child container.
- The padding is an optional part of the tag footer.

Based on these alternatives, we can really only decide the following:

DES 123: Padding in ID3v2 is modeled as very special child container of the tag

Padding in ID3v2 is modeled as very special child container of the tag, where the header is the first null byte, and all follow-up bytes form the payload. Note that for this to work, **jMeta** must support containers with no (i.e. empty) payload for the very special case that the padding only consists of one byte: The header magic key identifying the padding.

Rationale: Modeling it as child field of the payload is not possible as it would contradict **??**. Modeling as part of the footer would be odd, as not every version of ID3v2 defines a footer at all, and we would have a variable sized footer optionally starting with an arbitrary number of nullbytes, which would not be very intuitive from a library users perspective, too.

Disadvantages: The only disadvantage possibly to be identified that there could be padding just consisting of a single byte. This would mean we'd have a container with just a header, but an empty payload, which actually is quite degenerate.

The overall term for containers, headers, footers, fields and payload is a **data-block**. This term is used further on throughout this document.

There is another ID3v2 specific that we directly want to exclude here:

DES 124: ID3v2 transformations such as unsynchronization, compression, encryption are not handled generically

ID3v2 is clearly overdesigned. It includes an outdated mechanism called *unsynchronization* to ensure players do not interpret ID3v2 tags as MP3 frames and try to play them. ID3v2 also defines frame-based compression and encryption possibilities.

These facilities are not handled generically in **jMeta**, they are just specifically implemented in the ID3v2 extensions.

Rationale: Unsynchronization is not needed anymore, as current players all know about ID3v2 tags. Even if it would be needed: There is no other data format known to date that defines any such transformation schemes, not even for compression or encryption.

Thus, declaring this all as generic concept is useless as only one format defines such things, and these things are even only very rarely used.

Disadvantages: No disadvantages known

11.2.3 Representing a Data Format

Data Formats need some kind of representation:

DES 125: Class for data formats and especially for container formats

There is a base class called `AbstractDataFormat` representing a general data format, which has a subclass called `ContainerDataFormat` that is used for all data formats listed in “?? ??”. Each individual format is represented as an instance of `ContainerDataFormat`. These two classes are neither enums nor “dynamic enums”, but just plain old java classes.

Rationale: We need some way to identify which data format a container belongs to. A class allows to store additional format-wide properties for logging or even with a functional notion behind. Why not using an enum to identify each supported data format? Because enums are static, and we do not want to change the enum whenever a new format comes around. And it is not clear which formats are out there by 3rd party. A dynamic enum pattern is unnecessarily complexity here, so just a plain class. Why a just an instance of `ContainerDataFormat` rather than a new subclass for each format supported? First of all, data formats must be used as tokens by end-users of the library to refer to which tags or containers they want to read. Creating a new sub-class instance everytime they want to have a token is strange. Instead, they can use a public constant instance as token. The next question: Why at all a `ContainerDataFormat` subclass? The reason is that `jMeta` might want to support even more different formats in future, e.g. XML. This can be a sibling subclass of `AbstractDataFormat`. From now on, everything described in this section is only covering container data formats.

Disadvantages: No disadvantages known

But that’s not all of it. If we structure data belonging to a format into containers, headers, footers and so on according to the generic metamodel, *how exactly* is a correct datablock of a given format structured? Of course, each format individually specifies the exact structure of its containers (how many headers, footers etc., what the payload can contain), headers (which fields make up a header), fields (which type a field has, how long it is and which values it can have) and so on. Thus, we need a place where we describe the structure of all datablocks belonging to a data format in more detail, otherwise we could not parse this data generically.

We define:

DES 126: `DataFormatSpecification` provides `DataBlockDescriptions`, `DataFormatRepository` maintains all `DataFormatSpecifications`

The exact properties of a datablock is described by an instance of the class `DataBlockDescription`. Which properties are needed in detail is determined later in “?? ??”. The interface `DataFormatSpecification` provides access to all `DataBlockDescriptions` offered by one specific data format. Last but not least there is a `DataFormatRepository` interface which basically only maintains a list of all supported data formats and their corresponding `DataFormatSpecifications`.

A concrete `DataFormatSpecification` is defined in an extension and loaded when starting `jMeta`.

Rationale: For generically parsing data blocks according to a format definition, the parser code needs information about the exact structure of a data block. Thus, we need a place to maintain all data formats, their specifications and a description of each possible individual container, field etc.

Disadvantages: No disadvantages known

11.2.4 `DataBlockDescription` Common Properties

First of all, we need a clear way to identify each specific datablock that might occur in the data format. Every container type, header, footer, field and payload needs an identifier:

DES 127: Each datablock of a format has a specification id unique across all supported formats

Each distinctly defined container, header, footer, field or payload block of a data format is assigned a unique specification id. Using this specification id, a parser implementation can lookup the `DataBlockDescription` of the block from its associated `DataFormatSpecification`. Specification ids are human-readable rather than generated GUIDs. The uniqueness is guaranteed by using a top-level prefix identifier that is directly checked to be unique upon all loaded data formats when loading the extension defining the data format. The ids reflect the hierarchical nature of a data format according to the container meta model, following these rules:

- Each id consists of one or more *segments* separated by a dot.
- A segment name does only contain ASCII alphanumeric characters, starting always with a lower-case letter and preferring lower-case only letters, i.e. no special characters, no blanks etc. So a datablock id has much in common with a Java package name.
- The first segment is the data format top-level prefix which equals the (single) top-level container id of the data format. So we have the implicit decision that each format starts off with a single container on top-level enclosing everything else. This is true for all data formats considered here.
- Each container consists of an arbitrary number of headers and footers as well as exactly one payload. The segment name of a container payload is always `payload`. If there is just a single header or footer, it is called header or footer respectively. Multiple headers or footers bear more specific, distinct names.
- Child fields of a header or footer have descriptive names and so on

Rationale: When parsing, the format structure of a datablock must be clearly visible and available. By an id, this format structure is clearly adressable and the corresponding `DataBlockDescription` can be obtained from a `DataFormatSpecification`. The uniqueness of the id guarantees that the parsing logic cannot unintentionally try to parse a block of bytes according to the incorrect format.

The reason for not using generated unique ids such as counters or GUIDs is that we want to have memorable and recognizable names, both for library developers (e.g. during debugging) or `jMeta` end-users that need to work with datablock ids directly.

Disadvantages: No disadvantages known

Besides their id, what else might be desirable on first glance when looking at a datablock taken from a data format's specification? Some descriptive stuff comes

to mind. Second, it seems like every datablock must have exactly one type that is useful to store. So we can state:

DES 128: Basic properties held in a DataBlockDescription

Each `DataBlockDescription` instance holds the datablock id of the datablock according to `??`. In addition, it has:

- A more human-readable name of the datablock, if available; only for display or documentation purposes
- A description of the datablock, probably taken from the format specification, preferably in English. This is meant to give more information about the purpose of the datablock, but has no other meaning.
- The `DataBlockType` of the datablock, which is: `HEADER`, `FOOTER`, `CONTAINER`, `FIELD`, as well as - according to `??` - `FIELD_BASED_PAYLOAD` or `CONTAINER_BASED_PAYLOAD`.

Rationale: These properties are clear from the beginning, the concrete is known statically, and it is needed to take decisions during generic parsing.

Disadvantages: No disadvantages known

We defined that datablock ids are organized hierarchically as their real blocks are organized hierarchically. However, still, it is not sufficient to somehow maintain the “hierarchy” indicated by the metamodel only in the datablock ids. Instead, we define:

DES 129: `DataBlockDescriptions` contain a list of all child ids they might contain, where order matters

A `DataBlockDescription` instance contains a list of all child datablock ids that form the structure of the datablock described. While headers, footers, containers and payload usually has child ids, `DataBlockDescriptions` of fields must not define child datablock ids. The list of child ids is ordered with the following meaning:

- For containers, the child id list exactly matches the correct order of headers, payload and footers that build the container
- For headers, footers and field-based payload, the child id list exactly matches the order of child fields within the header, footer or field-based
- For container-based payloads, the child id list order does not have a meaning, i.e. if multiple different container types are allowed, they might appear in any order

Rationale: The parsing code somehow needs information about what child blocks of which format to expect within a given datablock, otherwise parsing it correctly is simply impossible. Each data format clearly defines the order of fields, headers and footers. Usually, most container formats define just exactly one child container that can appear in a container-based payload. However, there are cases such as QuickTime (atom vs. QT atom container) that define multiple different container types that might appear in any possible order.

Disadvantages: No disadvantages known

We first have to define some limitations for these datablock ids that should be made clear:

DES 130: Datablock ids and their segments MUST not be used for any logic

Especially, no assumptions about the hierarchical relations of distinct blocks must be done based on their id. E.g. a block with specification id `id3v23.payload.talb.encoding` not necessarily is a child of a block with id `id3v23.payload.talb`, as there could be multiple `talb` containers within an ID3v23 tag. So datablock ids are only specification ids, not instance ids! Even within the same specification, a case might occur where the same chunk of data with the exact structure appears in several places, e.g. as child of the top-level container, and also as deeper nested child. Moreover, some container formats encourage to embed datablocks belonging to even completely different data formats. How do we handle such cases? In no special way. Even if a container might appear on several levels, we just define its structure once in the data format with just one id. And if a datablock of a data format is embedded in another container of a different format, it still keeps its well-defined datablock id according to its data format. So this again shows that based on the id, we can never determine the real hierarchy level or instance relations of two datablocks during parsing!

Rationale: Quite clear, otherwise we would implement a lot of bullshit code

Disadvantages: No disadvantages known

11.2.5 Generic and Concrete Containers

Data formats define different kinds of *container types*. A container type for us is defined as a data structure having a fixed number of headers and footers, where each header and footer has a strictly defined sequence of fields with strictly defined syntax and semantics, yet leaving undefined how the payload looks like. They just state that there is a payload consisting of bytes, but not how it is structured. Most often, there is just one header or just one footer. Here are some examples of container types:

- Most container formats such as RIFF or Matroska define exactly one container type. In these formats, those containers usually can be found on top-level as well as arbitrarily nested within parent containers
- Some formats that are no container formats in the narrower sense¹ such as ID3v23 use a slightly different container type for the top-level (the tag) container as for the child (the frame) containers
- Some container formats such as QuickTime even define multiple container types that can appear on the same level: First of all, there are the top-level containers called atoms; one extension of atoms are QT atoms, which just use a slightly extended header compared to the usual atoms. Both

¹In terms of their main purpose is to contain *arbitrary* multimedia data, not only data of a specific kind.

might appear on top-level. Furthermore, as child container types, not only atoms or QT atoms might appear, but also QT atom containers (a slightly confusing term here, but this is how QuickTime names them)

Container types are the pillars of the most basic property of most well-designed container formats: They define the structures of their basic buildings blocks (i.e. their container types) just once, and specialize this basic structure whenever needed. This at first allows to write generic parsers that may skip containers whose semantics and purpose they might not be aware of, but yet they still know their structure and thus are able to determine where the next container actually starts. This allows to extend data formats by new subtypes whereas existing parsers still work.

But this is not yet enough. Most container formats go much farther than only define a generic container type that can be parsed and defines a structure. They specialize in terms of defining specific containers having a specific meaning (data content) and whose payload follows a deeper defined format. These specialized containers can be said to have the same structure as a given container type, but they fill this container type with meaning and the payload with additional structure.

We define:

DES 131: DataFormats distinguishes generic and concrete containers

DataFormats distinguishes generic and concrete containers:

- A container is called a **generic container** if it is just known to follow a container type's structure, but a more specialized container type is not (yet) known; more technically speaking: A generic container is a container whose detailed payload structure is not yet known in advance, but can only be determined at runtime, usually by reading some specific `id` field of the container's header or footer. Examples for generic containers are: Generic ID3v23 frame, generic (QT) atom, generic Matroska segment, generic APEv2 item and others. We can say that generic containers form the archetypes for the even more useful concrete containers. DataFormats must maintain `DataBlockDescriptions` for generic containers.
- A **concrete container** is a container with a well-defined payload structure. It might be based on a generic container which means it entirely inherits its header and footer structure, while it might "override" specific field values in headers and footers. Examples for concrete containers are: ID3v23 tag (because its payload structure is already statically clearly defined), ID3v23 TALB frame, ID3v23 GEOB frame, moov QT atom, APEv2 title item and others. DataFormats must maintain `DataBlockDescriptions` for concrete containers.

Of course, there are much more concrete containers than there are generic containers.

Rationale: Generic containers are necessary such that `jMeta` can basically parse data format content, yet not necessarily understanding its payload. These "unknown" containers can safely be skipped, and `jMeta` can later be extended to also understand these concrete types.

Concrete containers are necessary such that `jMeta` can read and write data with specific semantics and inner structure. Otherwise it would not be possible to write album, track, artist etc. fields to metadata container formats and so on.

Disadvantages: No disadvantages known

How exactly are these generic and concrete container `DataBlockDescriptions` defined in a `DataFormatSpecification` and used at runtime?

DES 132: DataFormatSpecifications list generic container as children, parsers resolve them to concrete containers at runtime

Generic container types might occur on arbitrary levels of the data block hierarchy, either on top-level or as child datablocks of a `CONTAINER_BASED_PAYLOAD` datablock. The semantics is: The file/media stream or payload contains a sequence of containers having a specific basic structure (i.e. they have a specific container type). This type can be found to be even more specific at runtime. For this, container formats define a special id field either in the container header or footer. This id field identifies the concrete type of container. After the parser has read the id field, it can lookup if there is a more specific `DataBlockDescription` for this id defined in its `DataFormatSpecification`. If so, the generic container is changed to the concrete container, thus the payload can now be parsed differently.

Rationale: All container formats define their `CONTAINER_BASED_PAYLOAD` to contain a multitude of different concrete containers that all follow (usually) a single generic container type.

Disadvantages: No disadvantages known

The actual detailed parsing process is described within the design of `Data-Blocks`.

Now we already mentioned in ?? that there is “usually” only a single generic container defined within a `CONTAINER_BASED_PAYLOAD`. However, there are also already two exceptions we can list here:

- A (QT) atom generic container having a `CONTAINER_BASED_PAYLOAD` could have two different generic container types as child containers: Either (QT) atoms again, or QT atom containers that have a totally different header.
- We already “artificially” introduced a second case with ??: An ID3v23 tag `CONTAINER_BASED_PAYLOAD` either contains a generic ID3v23 frame or a padding container.

If a `CONTAINER_BASED_PAYLOAD` might contain different types of containers in arbitrary order, how is a parser able to tell which generic container type he is handling currently? This is all based on magic keys and described in “?? ??”.

11.2.6 Data Format Identification and Magic Keys

This section deals with the question how to identify that a given chunk of bytes belongs to a specific data format. This is not an easy thing because of the following aspects:

- There are different media types such as files, continuous streams and in-memory bytes
- Different strategies of reading data formats are (partly) supported in `jMeta`: Most data formats are designed for the default way of reading, i.e. **forward**

reading starting from the beginning of the medium from offset zero towards higher offsets. However, sometimes **backward reading** is much more performant, especially for multimedia tags. Last but not least, you can think of **scanning** to search a chunk of bytes for the first occurrence of a known format.

- Some data formats are extremely similar to each other, either due to the fact that they are different revisions of the same family of formats (e.g. ID3v23 and ID3v24) or due to decisions of standardization boards (e.g. the ISO base multimedia format with its children JPEG 2000, MPEG-4 and others which are nearly identical to QuickTime) - how to distinguish them.
- A few data formats such as MP3 allow coexistence of multiple unrelated data formats on top-level of a medium, usually a file. An MP3 file not only contains MP3 frames, but also ID3v1, ID3v3, LyricsV3 or APE tags, all of which are entirely different data formats specified on their own. This is a special case as most other multimedia data formats such as RIFF, QuickTime, Matroska and Ogg require to only have data blocks of RIFF, QuickTime, Matroska and Ogg on top-level in the same medium.

What are the different possibilities of identifying the data format of a medium?

- *Based on the file extension:* In 80 percent of the cases, you could think of identifying a data format based on its file extension. This requires a mapping of extensions to data formats. While this could be a feasible approach, its downsides are clear: If the file has no extension (e.g. on Unix systems this could be the case), you cannot identify it. Furthermore, this approach is not feasible for streaming media and in-memory bytes.
- *Using other external hints:* You can use other external hints, e.g. external metadata (provided in separate files, in an HTTP header for streams) or let the user select or hint at the data format. Obviously, this could never be a possibility covering all cases and requires a lot of very diverse strategies to implement.
- *Based on information in the data:* You read the bytes and try to identify the data format based on the bytes in the data itself. Mostly, this translates into using magic keys which we will cover immediately or scanning. While this might seem strenuous, it is actually the only proven way to ensure proper identification that even works for all different media.

So we decide:

DES 133: Default identification mechanism uses magic keys

jMeta uses magic keys to identify the data format of a given chunk of data bytes. Scanning is not used for identification for now.

Rationale: It's the only reliable way to identify a data format, and we anyway need to parse the data. It works for all media types. Regarding scanning: It is currently excluded, see “?? ??”.

Disadvantages: No disadvantages known

Now we want to deal with the very basic and central concept common to most if not all of those data formats: **Magic Keys**. A magic key is a sequence of bytes with a fixed, well-known value, which allows to identify a block of bytes as belonging to a data format, enabling generic parsers to understand how to parse a byte sequence. Thus, it is usually located at the start of the medium. For example, an Ogg page is identified by the magic key string “OggS”, encoded as ASCII. Of course, this implies that different data formats use different magic keys.

These are the main properties of magic keys we are most interested in:

- A magic key has a single fixed byte identification sequence that is usually an ASCII encoded string
- It is either starting at the beginning of a container, or at a well-defined fixed offset from the start of the container, thereby being a field in the container's header or footer
- That said, also footers have magic keys to be used for backward reading, e.g. for ID3v23, APEv2 or Lyrics3v2
- Data formats might have multiple magic keys identifying them - one example is RIFF where the magic key also identifies the byte order of the container

Let's list all magic keys of the data formats we care about in table ??.

Data Format	Header Keys	Footer Keys
Ogg	“OggS” at start of page header	-
Matroska	1A45DFA3h	-
RIFF	“RIFF” (big endian) or “RIFX” (little endian) at start of chunk header	-
AIFF	“FORM” at start of chunk header	-

Data Format	Header Keys	Footer Keys
QuickTime	“ftyp” or “moov” or “pnot” or “mdat” or “skip” as atom type (starting with 5th byte) in atom header	-
MP3	Eleven 1 bits at start of frame header	-
ID3v1 and ID3v1.1	“TAG” at start of tag header	-
ID3v22	“ID3” followed by version number 0200h at start of tag header	-
ID3v23	“ID3” followed by version number 0300h at start of tag header	-
ID3v24	“ID3” followed by version number 0400h at start of tag header	“3DI” followed by version number 0400h at start of tag footer
APEv1	“APETAGEX” followed by version number 1000 at start of tag header	“APETAGEX” followed by version number 1000 at start of tag footer
APEv2	“APETAGEX” followed by version number 2000 at start of tag header	“APETAGEX” followed by version number 2000 at start of tag footer
Lyrics3v1	“LYRICSBEGIN” at start of tag header	“LYRICSEND” starting with 7th byte of tag footer
Lyrics3v2	“LYRICSBEGIN” at start of tag header	“LYRICS200” starting with 7th byte of tag footer
VorbisComment	<i>No magic key</i>	-

Table 11.7: Magic keys of all relevant data formats

Let us look at the odd exceptions in this list:

- (1) We see that the only format not having a magic key is the VorbisComment
- (2) MP3 as the only format does not have a string magic key, but rather a bit sequence
- (3) RIFF as the only format defines multiple different magic keys in its header
- (4) Lyrics3v2 and ID3v24 are the only formats where header and footer magic keys differ
- (5) QuickTime and Lyrics3v2 are the only formats where the magic key does not start at the beginning of its header/footer, but at a fixed byte offset
- (6) QuickTime is very vague about its magic key, it actually defines all top-level atoms as optional, but recommends to usually start a QuickTime file with the ftyp atom
- (7) Matroska seems to have an integer magic key

Gradually, we will develop a strategy that includes all of those special cases. First of all, we summarize and define:

DES 134: A top-level container defines one or several magic keys in headers or footers

Every top-level container must define one or multiple magic keys identifying the container and data format. The magic key is either at the start (header) or at the end (footer) of a container. We introduce a class `MagicKey` for describing the properties of a magic key. It defines:

- The magic key bytes
- A bit length of the magic key
- Optional: A string representation of the magic key
- Optional: A fixed byte-length distance from start of header (positive) or end of footer (negative)

Rationale: Every top-level container must at least have one magic key, because otherwise we cannot identify it. The possibility of multiple different magic keys is implied by e.g. RIFF and QuickTime. ID3v24 defines a header magic key as well as a different footer magic key for backward reading. The bit length in contrast to a byte length is necessary for the MP3 magic key which is not a multiple of 8 bits, but 11 bits long. The fixed byte-length distance from start of header or footer is necessary for formats such as QuickTime.

Disadvantages: No disadvantages known

Having defined this, we already covered the exceptions number (2), (3), (4) and (5). Now let's tackle QuickTime:

DES 135: ftyp, pnot, mdat and moov as magic keys for QuickTime

ftyp, pnot, mdat and moov are defined as the magic keys for QuickTime. If none of these atoms is found, a byte chunk cannot be identified as QuickTime.

Rationale: The vagueness of the specification is a problem, but we can reasonably handle it like that and be quite sure that in 90 percent of the cases, we can read QuickTime files properly.

Disadvantages: Some very special QuickTime byte sequence might not be readable with jMeta.

Now we are left with the VorbisComment special case. Unfortunately, you need external context knowledge to just know that the chunk of bytes is a Vorbis Comment. As described in [?, MC17]:

- In Vorbis bitstreams, it always is embedded in the second page of the stream
- In FLAC, Speex and theora bitstreams, it is always starting with the second packet of the stream, it is actually not quite clear if this is the same thing as for Vorbis, but note that Vorbis says page and the other say packet which is conceptually different in Ogg!

The VorbisComment, to say it clearly, seems quite oddly designed. We will not invent a generic solution for this, thus we define:

DES 136: VorbisComment oddities are only handled in the Ogg and VorbisComment implementations

The generic parsing algorithm of jMeta does not provide any support for data formats that have no magic keys such as VorbisComment. Instead, such oddities need to be handled in the individual implementations. We assume that this can be done quite easily in the Ogg and VorbisComment parsing implementations.

Rationale: Handling such a singular special case in a generic way is quite hard and would greatly complicate the generic parser. Implementations should be able to handle it with their special knowledge, which does not need to leak into the generic parser.

Disadvantages: No disadvantages known

We have already covered top-level containers and data format identification. But what about nested containers? Do they also need to have magic keys? In most cases, we have generic nested containers whose exact structure is only revealed at runtime during parsing, e.g. ID3v2 frames, RIFF chunks, Matroska elements and so on. Having an id that is always different, they usually do not qualify to contain magic keys. If all data formats would define a generic nested container type, all would be fine. But again, we have candidates for problems:

- QuickTime defines two different types of nested containers that have a completely different structure: atoms and QT atom containers

- ID3v2 has an odd use of padding by defining it as special null bytes behind the payload and before the footer (if any), thus he have defined to treat padding as special container, see ??

Here, we can use magic keys to the rescue:

DES 137: Nested containers define magic keys if not generic

Usually, nested containers are generic. Non-generic nested containers are allowed to define one or multiple magic keys in their headers or footers.

Rationale: QT atom containers have 10 zero-bytes at start which can be used as magic key. Likewise, padding always must start with a zero byte which is then the magic key.

Parsers reading payload bytes of a top-level container can first check for the presence of any magic key bytes of defined child containers. Only if these are not found, the single generic child container is assumed to be present.

Disadvantages: No disadvantages known

Note that there are still two problems with the formats to solve when doing it like this:

- QuickTime (again) defines atoms and QT atoms as generic container types that can also be nested containers. QT atoms start with the same structure as atoms, but add some header fields.
- When parsing child content like defined in ??, there is a special case for Lyrics3v2: Unfortunately, the payload size of this tag is not known in advance, but its end is only identified by the magic key in its footer. This special case is treated in ?? in another section.

So let us check how to handle yet another QuickTime oddity:

DES 138: QT atoms can be identified at runtime

QuickTime as only format relevant for us defines two different generic container types: Atoms and QT atoms. Thus, we define these two types of generic containers and use the generic container identification schemes to handle the situation. At runtime, if a specific id is found, the atom type is known to be an atom or QT atom, and further parsing is going on correspondly. See ?? for the runtime concrete container identification.

Rationale: Luckily, QT atoms just extend the header of atoms and otherwise are structure the same way. Thus this mechanism will resolve the problem.

Disadvantages: No disadvantages known

Based on these definitions, we can be more detailed about the parsing strategies:

DES 139: Parsing top-level containers (forward reading)

For top-level containers, the data format identification is the first thing that happens. Each data format must define at least one top-level container which can either be generic or concrete. Each top-level container type must at least define one header magic key that indicates presence of the format. If the current bytes match the magic key, the data format is identified. Further on, if the top-level container is a generic container, the runtime identification of its concrete type is following as defined in ???. Data formats that require deviating behaviour (e.g. VorbisComment) can override some aspects of this behaviour.

Rationale: This is directly following from all previous design decisions

Disadvantages: No disadvantages known

For nested containers, forward reading basically works like this:

DES 140: Parsing nested containers (forward reading)

For nested containers, no data format identification is needed. The data format defines all possible containers that might occur in the container-based payload of the parent container. Each of these containers can either be generic or concrete. Any defined nested container might define a magic key. The parser first gets all child containers of the current payload who have a magic key defined and tries to match the current bytes with these magic key bytes. If they match, the child container is found. Further on, if the nested container is a generic container, the runtime identification of its concrete type is following as defined in ???. If none of the child containers with magic keys match, the parser gets the *default nested container* which can be obtained from the `DataFormatSpecification`, if set. If this is the case, the parser assumes that the next container is the default nested container, otherwise it cannot identify the next container. If a container is found and this container is generic, the runtime identification of its concrete type is following as defined in ???. This might e.g. reveal that an atom (default nested container for QuickTime) is actually a QT atom of a concrete type, see ??.

Rationale: This is directly following from all previous design decisions. We have to define a default nested container just because of QuickTime as sole exception defining two very similar generic nested container types.

Disadvantages: No disadvantages known

Now the topic of backward-reading containers is still open. Again, there is a special case here: ID3v1 has no footer, but still is actually backward-readable. The reasons: It is defined to always be at the end of the file. Even if it is not, it could still be identified, as it has a fixed length. This means we would just need to look for the header magic key at a fixed offset from the end of the tag. However, there is a rather hidden special logic that additionally needs to be applied: If you identify the presence of a container by its footer magic key, you yet have to still

read backward to obtain its payload. In case of ID3v1, you have arrived at the header already, such that you actually need to read forward to get its payload.

To handle this, again, we have two alternatives: Either bring this special case into the generic layer, e.g. by adding a field to the magic key, which points to the offset for backward reading. Or we handle it only in the ID3v1 implementation. A third esoteric possibility would be to have a fake “footer” in ID3v1, such that we could just say: Whenever backward reading is necessary, just rely on footers only. But this would be very artificial and would lead to other problems. As was the case previously, we decide:

***DES 141:* Backward-reading for ID3v1 is implemented in the ID3v1 parser**

The special case for ID3v1 is not considered in the generic parser code, but it is specifically overridden for this purpose in the ID3v1-specific parsing code.

Rationale: This decision is in line with the other design decisions. We want to keep the basic concepts and generic parser code and classes clear and clean of single-format special cases.

Disadvantages: No disadvantages known

Now, finally, let us define how backward reading is working overall. It is just a feature meant to be used for tags at the end of file. You cannot backward-read MP3 frames or Ogg pages, because they have no footer, such that you would need to “scan backwards” which we do not want to do. The second thought about backward-reading: It only affects the top-level, and not nested containers. To summarize:

DES 142: Parsing top-level containers (backward reading)

A prerequisite for backward-reading top-level containers is that the data format has a footer with a magic key. The user has to explicitly request backward reading. The parser checks if it finds a footer magic key for all backward-readable data formats, i.e. data format identification happens here. It starts reading bytes to compare at the end offset of the footer plus the magic key offset (i.e. it reuses the offset defined in the `MagicKey` class, which is negative in this case). If the read bytes match the magic key, the data format is identified. Further on, if the top-level container is a generic container, the runtime identification of its concrete type is following as defined in `??`. Data formats that require deviating behaviour (e.g. ID3v1) can override some aspects of this behaviour. Backward reading stops as soon as no further backward-readable data format could be identified. It is not necessary to backward-read nested child containers.

Rationale: This is directly following from all previous design decisions. The reason why nested containers do not require backward-reading: Once you know were a backward-read top-level container starts, you can again forward-read its nested children.

Disadvantages: No disadvantages known

Now we have to clarify how the parsing code can best access the information about magic keys. We decide:

DES 143: Obtain MagicKeys from DataBlockDescriptions of containers

For the parser code, `MagicKeys` can be obtained from a `DataBlockDescription` of a container data block. It offers a method for obtaining the backward-readable (i.e. footer) magic keys as well as a method for obtaining the forward-readable (i.e. header) magic keys.

Rationale: A magic key essentially belongs to a container. The top-level container `DataBlockDescriptions` and the nested container `DataBlockDescriptions` must anyway be available at the places where we need to parse them. By providing different methods for forward and backward reading, the intentions become very clear. Overall, it is most convenient for the parser code this way.

Disadvantages: No disadvantages known

Strongly connected to this is the question how to represent magic keys in a data format while avoiding redundancies. We already defined a specific class called `MagicKey` and methods for obtaining the magic keys for this, which is a good choice as it bundles all the relevant information (and nothing more) for the parser code. But in the end, a magic key is nothing but a header or footer field with a fixed value. Thus we define:

DES 144: MagicKey instances implicitly derived from field**DataBlockDescriptions**

The **MagicKey** instances that can be queried by the parsing code are not defined/added in parallel to the field properties anyway needed. In contrast to this, a **DataDataBlockDescription** of type **FIELD** additionally contains a flag saying if it is a magic key or not. If this flag is set to true, all other properties of the **MagicKey** can implicitly derived from the other field properties and a plausibility check can happen. If one of the following plausibility checks is failing, the **DataFormatSpecification** is considered invalid.

- A field marked as magic key must either have a fixed value, or it must be an enumerated field (i.e. with multiple alternative values).
- The header that contains the magic key field must be either be the first header or the headers before it must be fixed length headers. For footers: The footer that contains the magic key field must be either be last footer or the footers behind it must be fixed length footers. Otherwise the field does not qualify as magic key, because it might have a variable offset from the start or end of its container and the **DataFormatSpecification** cannot derive the necessary offset from start or end of container.
- For the same reason, we define for headers: The fields in front of the magic key field (if any) must all be fixed length fields. And for footers: The fields behind the magic key field (if any) must all be fixed length fields.
- There must be only at maximum one header magic key field and only at maximum one footer magic key field

Rationale: Redundancy is avoided, and it is more convenient to define a concrete **DataFormatSpecification**, as we do not need to repeat ourselves.

Disadvantages: A **DataFormatSpecification** implementation must be slightly more complex to derive and check the necessary properties of the **MagicKey** instances.

There is a special case with MP3 here: It has an odd number of bits. This makes defining the magic key as single field very hard. Instead, it needs to be defined as flag field and cannot be derived automatically:

DES 145: MP3 magic key is added manually to the specification

The MP3 magic key with 11 one bits is added manually to the MP3 `DataFormatSpecification`. Thus `DataBlockDescription` offers a method for adding a header or footer magic key manually. So auto-derival of magic keys is not used for MP3.

Rationale: The odd bit length does not allow to define the 11 bit field as field on its own, but it needs to use flags, and flags do not have individual default values. Thus it is not possible in the generic code to derive the magic key for MP3 automatically. Instead, the field is not tagged as magic key, but it is manually created in the MP3 extension and added using the `DataBlockDescription` method.

Disadvantages: No disadvantages known

Last but not least, we want to discuss the topic of different data format versions and their magic keys. Specifically, one could ask: Does the version number belong to the magic key of a data format or not? We define:

DES 146: Magic keys do not include a version number

Magic keys never include the version number of the data format, even if it directly follows or precedes the magic key and it might appear as solution to the problem of how to distinguish similar data formats.

Rationale: First of all, a version number might be numeric, while magic keys are usually strings. Does the string representation of the version number would be slightly odd. Second and more importantly, it is actually a different field, and it could well be that a data format does not define the version field to be directly adjacent to the magic key field.

Disadvantages: No disadvantages known

11.2.7 Inheriting and Overriding Datablock Structures

TODO

11.2.8 Occurrences and Sizes

In this section, we discuss the topics of how often does a data block occur, and how long is it?

The “how often” (occurrences) part is treated as follows:

DES 147: DataBlockDescription contains minimum and maximum occurrences

For every data block, its `DataBlockDescription` can specify the number of occurrences by giving its minimum and maximum occurrences. The rules:

- Either both are null, or both are set to “unlimited occurrences”, or minimum occurrences is set to a positive number (including zero) and maximum occurrences is set to “unlimited occurrences”, or both are set to a positive number (including zero)
- If both are set, minimum occurrences must be smaller or equal compared to maximum occurrences

Rationale: A few formats have fields whose occurrence count instead of their length is given: Vorbis comment (number of user comments), APEv2 (number of items, in addition to size), Ogg (number of segments, which is equal to size, as each segment has one byte as length).

Disadvantages: No disadvantages known

For the size of a data block, we have basically three possibilities:

- The data block has a constant, fixed size which is always the same. Most fields have this property, but also static format containers such as ID3v1.
- The data block is a variable-size field and it is terminated, see “?? ??” for more details about this case. This leads to variable-size headers, footers, payload and containers that have a variable size because one of their fields is a variable-size field.
- The data block size is dynamically determined by a previously read field value or a combination of such field values. This case is further detailed in section “?? ??”.

For each case, we need the following:

DES 148: DataBlockDescription contains minimum and maximum byte length

`DataBlockDescription` allows to set a minimum and maximum length of the data block in bytes. Either both are null (meaning: “unknown length”), or the minimum length is set to a positive number (bigger than zero!) and the maximum length is set to “unknown length”, or both are non-null and set to a positive number (bigger than zero), where the maximum length must be equal to or bigger than the minimum length.

The fixed size case is covered by setting both to the same value. Differing values have no specific influence for the parser, but can be later used for format validation of a chunk of bytes. If not set or the values differ, this is a signal that the data block is either terminated or its size needs to be determined dynamically using other field values.

Rationale: `jMeta` needs to be able to reflect both static sizes as well as dynamic sizes, and to be future-proof, it is also good to have differeng min and max sizes given for data format validation.

Disadvantages: No disadvantages known

We just cover one additional special case here: In ?? and the follow-up discussion, we already noted that `Lyrics3v2` is somewhat special in terms of how to determine its payload size: It cannot be done before parsing it. Instead, it is only known to end when the footer’s magic key is found. Thus it is kind of “terminated”. In a first design, this was handled by some kind of special magic keys, so-called “exclusion keys” that where defined to state “if you find this key, a container of the current type IS NOT present”. However, we decided to redesign this to be less odd:

DES 149: Lyrics3v2 payload size is determined in custom implementation

`Lyrics3v2` special case is treated in a custom implementation by overriding the process of determining the payload length in the `Lyrics3v2` parser implementation.

Rationale: This case is only known for `Lyrics3v2`, thus we do not want to handle it in the generic parsing code. The `Lyrics3v2` implementation is the right place to handle it.

Disadvantages: No disadvantages known

11.2.9 Field Properties

Now we look into properties that are specific to data blocks of type `FIELD`. Here, we cover other properties specific to fields. However, we do not cover the details of how field conversion (interpreted value to binary value and vice versa) is working

and how some special cases of field interpretation are handled. This is a topic for the `DataBlocks` component.

We already covered sizes of data blocks (including fields) and we already defined one property of a field: It might be a magic key field, see `??`. The fields are the leafs of the data block hierarchy, thus, ultimately, they represent concrete values. While we could decide to treat them only as raw bytes, this is counter-productive as it would make parsing less convenient. Thus we define:

DES 150: Fields have one of the types binary, string, numeric, enumerated or flags

Every field has a field type, and thus not only a **binary value**, but also an **interpreted value**. The field types are:

- **String:** The field's value is representing a string in a specific encoding.
- **Binary:** The field's value is not further interpreted, i.e. the interpreted value is identical to the binary value of the field. I.e. this could be used to represent raw encoded bytes of a codec.
- **Numeric:** The field's value is representing a numeric value with a specific byte order.
- **Enumerated:** The field's value must be one out of a fixed set of values; the interpreted value is allowed to have an arbitrary type, but the field's specification must provide a static map of all allowed byte values (keys) mapped to their interpreted values.
- **Flags:** The field's value is representing flags, i.e. a set of bit toggles, e.g. vital for parsing.

Rationale: Parsers can directly read the interpreted values of the fields they need for parsing (e.g. sizes, flags or others) without needing to implement conversion to an interpreted value themselves. Interpretation should be kept inside the format specific code - in this case in the data format specification. All of the field types of our reference data formats can be covered by these types.

Disadvantages: No disadvantages known

First, we will care about the string fields:

DES 151: String fields support termination characters and fixed character encoding as specific properties

A non-fixed-size string field can be terminated by a single termination character given in the field's character encoding known at parsing time. Second, a string field might have a fixed character encoding to be used. Both are optional properties.

Rationale: String fields can be terminated by termination characters, this is needed e.g. in ID3v1 and ID3v23 data formats, where the actual number of termination bytes depends on the current character encoding. Thus, it is best to define a termination character and then decide at parsing time, which bytes according to the current character encoding are expected as termination bytes. See also ??.

Some specifications define fixed character encodings for specific string fields that might deviate from the data format's default character encoding.

Disadvantages: No disadvantages known

What about binary fields? Basically, they do not need any additional properties, but we want to make one thing clear about termination:

DES 152: Binary fields do not support termination bytes

While string fields can define a fixed termination character, binary fields do not support termination bytes.

Rationale: The reason is simple: Within our reference formats, there is no single case where a terminated field exists that is not a string field.

Disadvantages: No disadvantages known

Numeric fields require a byte order such that their numeric value can be interpreted:

DES 153: Numeric fields support a fixed byte order

Numeric fields define a fixed byte order to use during conversion. It is an optional property.

Rationale: To be convertible into a numeric value, the parsers need to have information about a byte order to use.

Disadvantages: No disadvantages known

For enumerate fields, we define the following:

DES 154: Enumerated fields require a static mapping of allowed enumerated values

Every enumerated field must define a static map of all enumerated values that are allowed for the field. The mapping contains an interpreted value as key and maps it to its byte representation. Not only the keys must be (of course) unique, but also the values.

Rationale: This is the nature of enumerations.

Disadvantages: No disadvantages known

For flags, it would be convenient to describe them in particular, e.g. which flags are there, where are they and what is their meaning. Thus we introduce a new class:

DES 155: Flag fields require an instance of FlagSpecification that describes all flags in a flags field

The class `FlagSpecification` describes a flags field having following properties:

- The fixed byte length of the flags field
- The positions (byte and bit adress) and unique names of flags together with their specification descriptions
- A byte order

It also supports multi-bit-flags.

Rationale: It is more convenient to access the flags, in contrast to doing bitwise OR and AND yourself.

Disadvantages: No disadvantages known

Finally, some fields - no matter what type - have always the same static value. We join this requirement with the notion of a default value, i.e. a value to be preferred when it is not known which would be the right value of a field.

DES 156: Each field has an optional default (interpreted) value

Each field has an optional default (interpreted) value.

Rationale: The default values are needed for fields with fixed values, and might be used to fill field values in cases where they are unknown.

Disadvantages: No disadvantages known

11.2.10 Field Functions

11.2.11 Header and Footer Properties

11.2.12 Container Properties

11.2.13 Payload Properties

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