



LCI Advanced Workshop 2025: Slurm Scheduler Tuning: Balancing HTC & Large Jobs

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Learning Goals



- Understand the differences between HTC and large-job workloads and their impact on scheduling
- Learn how to identify and mitigate large-job starvation in mixed environments
- Explore scheduler parameter tuning for different cluster sizes
- Compare the roles of the main scheduler vs backfill scheduler
- Explore best practices for partition design and public/private node management
- Gain strategies for iterative tuning and monitoring performance impact



Resource Allocation Models in Slurm



Whole Node Allocation (SelectType=select/linear)

- Job gets entire node regardless of requested resources
- Simple, predictable, no fragmentation
- Wasteful for small jobs



Resource Allocation Models in Slurm



Consumable Resources (SelectType=select/cons_tres)

- Allocates CPUs, GPUs, memory individually
- SelectTypeParameters: CR_CPU, CR_Core, CR_Socket, CR_Memory
- Maximum resource utilization but increased complexity

Scheduling Considerations:

- Job sizing and resource requests
- Node fragmentation
- Memory enforcement (MemEnforceLimit)
- GPU/accelerator placement
- NUMA awareness



Scheduler Types in Slurm



SchedulerType=sched/buildin

- FIFO scheduler initiates jobs in priority order.
- If any job in the partition can not be scheduled, no lower priority job in that partition will be scheduled.
- exception is made for jobs that can not run due to partition constraints (e.g. the time limit) or down/drained nodes.

SchedulerType=sched/backfill

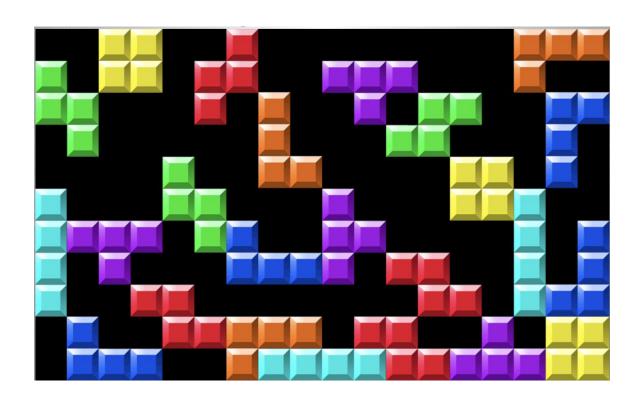
- Backfill scheduling module to augment the default FIFO scheduling
- Effectiveness dependent upon users specifying job time limits





Main vs Backfill Scheduler Roles

- Main scheduler: determines job start order based on priority
- Backfill scheduler: fills idle resources without delaying higher-priority jobs





Main Scheduler Loop



- Wake up every SchedulerTimeSlice seconds (default: 5)
- Load job queue sorted by priority
- For each job in priority order:
 - Check partition limits (MaxNodes, MaxCPUsPerUser, etc.)
 - Check association/QOS limits
 - Check job dependencies
 - Attempt resource allocation
 - If allocation succeeds: start job or make reservation
 - If allocation fails: continue to next job
- Stop when max_sched_time exceeded OR bf_max_job_test jobs evaluated
- Sleep until next cycle



Backfill Scheduler Loop



- Triggered every bf_interval seconds (default: 30)
- Build "reservation table" from main scheduler decisions
- Calculate earliest start times for all pending jobs
- For each job in queue (not priority order):
 - If job can start now without violating reservations: start it
 - Track jobs started (bf_max_job_start limit)
 - Track jobs tested (bf_max_job_test limit)
- Stop when limits hit or queue exhausted





The Value of Backfill Scheduling

Without Backfill

- Strict priority queue processing
- Resources idle while waiting for large jobs
- Poor cluster utilization
- Users frustrated by idle nodes they can't access

With Backfill

- Smaller jobs fill gaps between large jobs
- Better user experience (shorter wait times for small jobs)
- Maintains fairshare and priority goals



Balancing Partitions to Optimize Scheduling



Overlapping Resources

- Multiple partitions sharing same nodes
- Scheduler evaluates each partition separately
- Duplicate effort checking same resources
- Can cause scheduler performance issues

Minimize partition overlap where possible

- Use AllowAccounts/AllowGroups instead of separate partitions
- Partition by job characteristics (size, runtime, priority)
- Monitor sdiag for excessive partition evaluation times







```
# sdiag
```

Main schedule statistics (microseconds):

Last cycle: 2302

Max cycle: 62634

Total cycles: 26014

Mean cycle: 3323

Mean depth cycle: 69

Cycles per minute: 57

Last queue length: 68

Main scheduler exit:

End of job queue:26012

Hit default_queue_depth: 2

Hit sched max job start: 0

Blocked on licenses: 0

Hit max rpc cnt: 0

Timeout (max_sched_time): 0







```
# sdiag
Backfilling stats
                                         Backfill exit
     Total backfilled jobs: 243060
                                              End of job queue: 454
     Total cycles: 454
                                              Hit bf max job start: 0
     Last cycle: 609383
                                              Hit bf max job test: 0
     Max cycle: 853384
                                              System state changed: 0
     Mean cycle: 294402
                                              Hit table size limit: 0
     Last depth cycle: 130
                                              Timeout (bf max time): 0
     Last depth cycle (try sched): 130
     Depth Mean: 189
     Depth Mean (try depth): 188
```





Tuning Main Scheduler Cycle Timing

- max_sched_time Max time per main cycle (default: 2s), prevents
 scheduler from blocking too long
- sched_interval How often the main scheduling loop will run,
 considering all jobs while still honoring the partition_job_depth limit
- default_queue_depth Number of jobs to consider for scheduling on each event that may result in a job being scheduled. (Default 100)



Tuning Backfill Scheduler Cycle Timing on high performance clustered computing

bf_interval

- How often backfill runs (default: 30s)
- Independent of main scheduler cycle
- Reduce for fast-moving HTC workloads
- Increase to reduce overhead on large systems



Tuning Backfill Scheduler Cycle Timing on high performance clustered computing

bf_window

- Jobs needing resources beyond window ignored
- Increase for long jobs
- A value at least as long as the highest allowed time limit is generally advisable to prevent job starvation

bf_window_linear

- Setting this will create higher resolution of time precision in the backfill scheidler uses linear backoff instead of exponential back off
- Not recommended with more than a few hundred simultaneously executing jobs.



Tuning Backfill Scheduler Cycle Timing on high performance clustered computing

bf_max_job_test

- Max jobs evaluated per backfill cycle (default: 300)
- Increase for deeper backfill
- Watch cycle times can make cycles slow

bf_max_job_start

- Max jobs started per backfill cycle (default: 0 = unlimited)
- Limit to prevent backfill from being too aggressive
- Useful for protecting large job opportunities

bf_max_job_assoc

- How many jobs per user associona backfill can check
- Large impact when users submit many jobs







defer - Do not attempt to schedule jobs individually at submit time. Can be useful for high-throughput computing.

defer_batch - Like defer, but only will defer scheduling for batch jobs.

pack_serial_at_end - may reduce resource fragmentation by put serial
jobs at the end of the available nodes

SelectTypeParameters=CR_LLN - Schedule resources to jobs on the least loaded nodes



Workload Mix & Scheduling Pressure



HTC Workloads Want:

- Fast turnaround (seconds to minutes)
- High throughput (jobs/second)
- Frequent backfill opportunities
- Minimal wait times

Large Jobs Want:

- Resource guarantees
- Protected start times
- Priority enforcement
- Minimal fragmentation



Tuning for All HTC



```
# Fast, frequent cycles
SchedulerTimeSlice=2
SchedulerParameters=max sched time=1
# Aggressive backfill
SchedulerParameters=bf interval=15
SchedulerParameters=bf window=120 # 2 hours
SchedulerParameters=bf max job test=1000
SchedulerParameters=bf max job start=0 # unlimited
SchedulerParameters=bf continue
PriorityFavorSmall=Yes
```







```
# Standard cycles, deeper evaluation
SchedulerTimeSlice=5
SchedulerParameters=max sched time=4
# Conservative backfill
SchedulerParameters=bf interval=60
SchedulerParameters=bf window=10080 # 7 days
SchedulerParameters=bf max job test=300
SchedulerParameters=bf_max_job_start=5
SchedulerParameters=bf reserve nodes=1
```



Workload Mix & Scheduling Pressure



The Conflict:

- HTC jobs fill all available resources
- Large jobs can't find contiguous allocations
- Priority system gets subverted by constant backfill
- Users of both types dissatisfied

Solution Requires careful balance of main and backfill tuning, partition design, and policy enforcement





Tuning for Hybrid Workloads

```
# Moderate cycle timing
SchedulerTimeSlice=5
SchedulerParameters=max sched time=2
# Tuned backfill
SchedulerParameters=bf interval=30
SchedulerParameters=bf window=1440 # 24 hours
SchedulerParameters=bf max job test=500
SchedulerParameters=bf max job assoc=20 # CRITICAL
```





Scheduler Tuning at ASU

```
SchedulerParameters =
bf_window=43200
bf_resolution=300
bf_max_time=120
bf_max_job_assoc=25
bf_max_job_test=150000
bf_interval=30
max_rpc_cnt=1000
enable_user_top
bf_continue
nohold_on_prolog_fail
```



Monitoring & Iteration



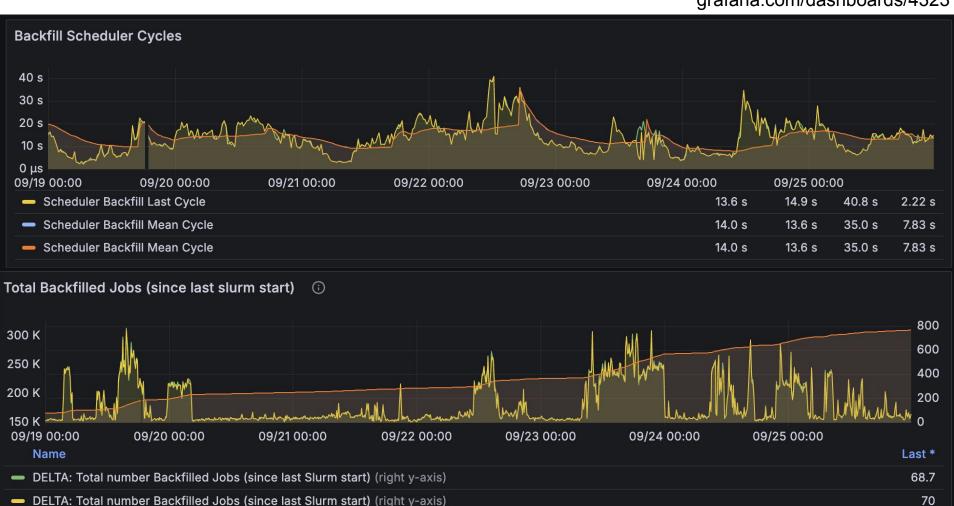
- Baseline before changes
- Metrics: wait times, utilization, fairshare
- Gradual adjustments
- User feedback loop
- Slurm Exporter:
 - github.com/vpenso/prometheus-slurm-exporter
 - github.com/ubccr/slurm-exporter (Uses Slurm API)
 - gitlab.com/SchedMD/slinky/slurm-exporter (Official)



Monitoring & Iteration



grafana.com/dashboards/4323





Runaway Jobs



sacctmgr show runawayjobs

NOTE: Runaway jobs are jobs that don't exist in the controller but have a start time and no end time in the database

[root@sol-slurm01:~]# sacctmgr show runawayjobs	
NOTE . Punaway jobs	and jobe that don't oxict in the	controll

NOTE: Runaway ID		jobs that do Partition	n't exist in Cluster	the controller State	but have a start TimeSubmit	t time and no end time in TimeStart	the database TimeEnd
28108882 35866256 35866258 35866283	/home/jeb+ ood-vscod+ ood-virtu+ interacti+	general	sol sol sol	PENDING 2025- PENDING 2025-	07-02T00:05:17 10-14T09:49:36 10-14T09:49:46 10-14T10:39:40	Unknown Unknown Unknown Unknown	Unknown Unknown Unknown Unknown

Would you like to fix these runaway jobs?

(This sets the end time for each job to the latest of the job's start, eligible, and submit times, and sets the state to completed.

Once corrected, this triggers the SlurmDBD to recalculate the usage from before the earliest submit time of all the r unaway jobs. Warning: This could take a long time and sreport may not return data until the recalculation is complete d.)

(You have 30 seconds to decide)



Key Takeaways



- Match tuning to workload type and site goals
- Use partitions and limits strategically
- Balance main scheduler and backfill tuning
- Iterate with data and feedback



Q & A



