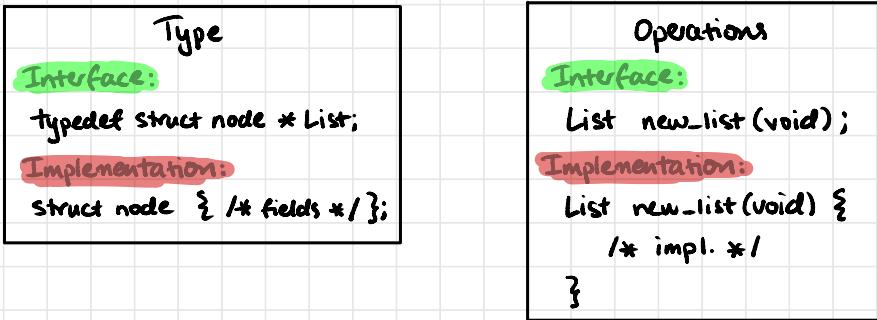


TUTORIAL 4: ADTs, BSTs

ABSTRACT DATA TYPES (ADTS)

Fundamentally, a data type in C will consist of 2 things:



But: to use a data type, we really just need to know its interface.

Also: an interface can have many possible implementations.

⇒ ADT: a data type which separates the interface from the implementation.

How? Put the interface in .h files, implementation in .c files.

Benefits:

- easy to swap out implementations without users needing to make changes to their code (the interface is still the same!)
- safe, because users can't accidentally alter internal ADT data and potentially cause problems inside it
- generalised nature of ADTs makes them reusable for other users

BINARY SEARCH TREES (BSTs)

Motivation: Searching for things in a sorted array is fast, because we can use **binary search**:

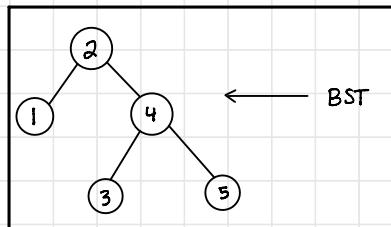
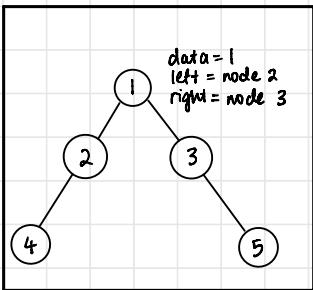


Problem: we now know sorting arrays is **expensive** — $\Omega(n \log n)$ at best. Inserting and removing values from a sorted array are $\Omega(n)$ as well.

Maintenance sounds like a pain. Can we do better?

Yes: think **hierarchically** instead of laterally!

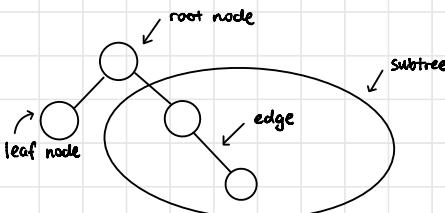
Binary tree: structure of nodes, where each has some associated value and up to 2 children:



Binary search tree: all nodes have the property **left data < data < right data**

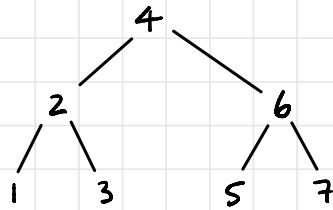
Selling point: If done properly, **search** and **insert** are both $O(\log n)$!

Terminology:



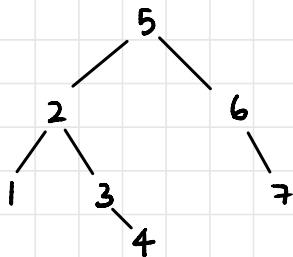
1. Insert values into the BST

(a) 4 6 5 2 1 3 7



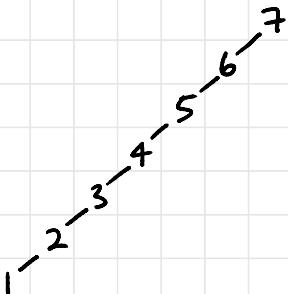
Note: This tree is balanced,
i.e. has minimal height for
the number of nodes
⇒ we get $O(\log n)$ search, inserts

(b) 5 2 3 4 6 1 7



This tree is less balanced,
but still not that bad
⇒ search, insert are somewhere
between $O(\log n)$ & $O(n)$

(c) 7 6 5 4 3 2 1



This is a really unbalanced
tree, since it is basically a
linked list (degenerate tree)
⇒ insert, search are $O(n)$:(